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ECC Cipher Suites for TLS  
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#### Abstract

This document describes new key exchange algorithms based on Elliptic Curve Cryptography (ECC) for the TLS (Transport Layer Security) protocol. In particular, it specifies the use of Elliptic Curve Diffie-Hellman (ECDH) key agreement in a TLS handshake and the use of

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Elliptic Curve Digital Signature Algorithm (ECDSA) as a new authentication mechanism.

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC 2119](#) [1].

Please send comments on this document to the TLS mailing list.

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## 1. Introduction

Elliptic Curve Cryptography (ECC) is emerging as an attractive public-key cryptosystem for mobile/wireless environments. Compared to currently prevalent cryptosystems such as RSA, ECC offers equivalent security with smaller key sizes. This is illustrated in the following table, based on [\[12\]](#), which gives approximate comparable key sizes for symmetric- and asymmetric-key cryptosystems based on the best-known algorithms for attacking them.

Symmetric		ECC		DH/DSA/RSA
-----	+	-----	+	-----
80		163		1024
112		233		2048
128		283		3072
192		409		7680
256		571		15360

Table 1: Comparable key sizes (in bits)

Smaller key sizes result in power, bandwidth and computational savings that make ECC especially attractive for constrained environments.

This document describes additions to TLS to support ECC. In particular, it defines

- o the use of the Elliptic Curve Diffie-Hellman (ECDH) key agreement scheme with long-term or ephemeral keys to establish the TLS premaster secret, and
- o the use of fixed-ECDH certificates and ECDSA for authentication of TLS peers.

The remainder of this document is organized as follows. [Section 2](#) provides an overview of ECC-based key exchange algorithms for TLS. [Section 3](#) describes the use of ECC certificates for client authentication. TLS extensions that allow a client to negotiate the use of specific curves and point formats are presented in [Section 4](#). [Section 5](#) specifies various data structures needed for an ECC-based handshake, their encoding in TLS messages and the processing of those messages. [Section 6](#) defines new ECC-based cipher suites and identifies a small subset of these as recommended for all implementations of this specification. [Section 7](#), [Section 8](#) and [Section 9](#) mention security considerations, intellectual property rights, and acknowledgments, respectively. This is followed by a list of references cited in this document and the authors' contact

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information.

Implementation of this specification requires familiarity with TLS [\[2\]](#), TLS extensions [\[3\]](#) and ECC [\[4\]](#)[\[5\]](#)[\[6\]](#)[\[8\]](#) .

## [2.](#) Key Exchange Algorithms

This document introduces five new ECC-based key exchange algorithms for TLS. All of them use ECDH to compute the TLS premaster secret and differ only in the lifetime of ECDH keys (long-term or ephemeral) and the mechanism (if any) used to authenticate them. The derivation of the TLS master secret from the premaster secret and the subsequent generation of bulk encryption/MAC keys and initialization vectors is independent of the key exchange algorithm and not impacted by the introduction of ECC.

The table below summarizes the new key exchange algorithms which mimic DH\_DSS, DH\_RSA, DHE\_DSS, DHE\_RSA and DH\_anon (see [\[2\]](#)), respectively.

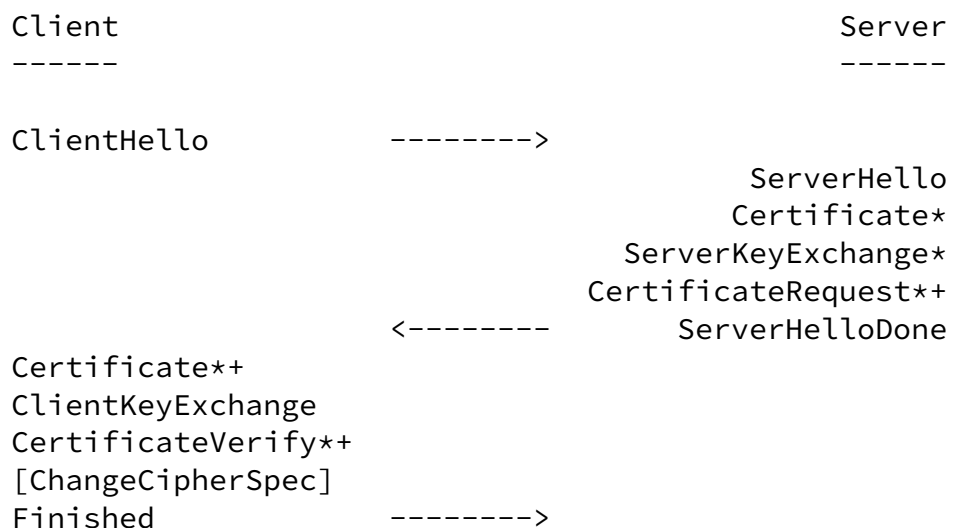
Key Exchange Algorithm -----	Description -----
ECDH_ECDSA	Fixed ECDH with ECDSA-signed certificates.

ECDHE_ECDSA	Ephemeral ECDH with ECDSA signatures.
ECDH_RSA	Fixed ECDH with RSA-signed certificates.
ECDHE_RSA	Ephemeral ECDH with RSA signatures.
ECDH_anon	Anonymous ECDH, no signatures.

Table 2: ECC key exchange algorithms

Note that the anonymous key exchange algorithm does not provide authentication of the server or the client. Like other anonymous TLS key exchanges, it is subject to man-in-the-middle attacks. Implementations of this algorithm SHOULD provide authentication by other means.

Note that there is no structural difference between ECDH and ECDSA keys. A certificate issuer may use X509.v3 keyUsage and extendedKeyUsage extensions to restrict the use of an ECC public key to certain computations. This document refers to an ECC key as ECDH-capable if its use in ECDH is permitted. ECDSA-capable is defined similarly.



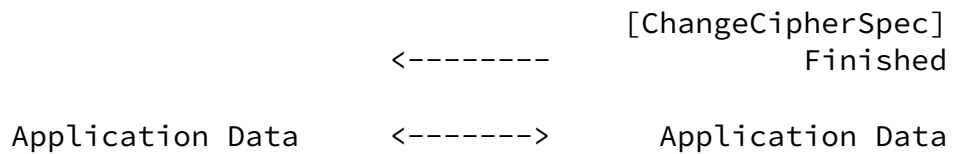


Figure 1: Message flow in a full TLS handshake  
 \* message is not sent under some conditions  
 + message is not sent unless the client is authenticated

Figure 1 shows all messages involved in the TLS key establishment protocol (aka full handshake). The addition of ECC has direct impact only on the ClientHello, the ServerHello, the server's Certificate message, the ServerKeyExchange, the ClientKeyExchange, the CertificateRequest, the client's Certificate message, and the CertificateVerify. Next, we describe each ECC key exchange algorithm in greater detail in terms of the content and processing of these messages. For ease of exposition, we defer discussion of client authentication and associated messages (identified with a + in Figure 1) until [Section 3](#) and of the optional ECC-specific extensions (which impact the Hello messages) until [Section 4](#).

## 2.1 ECDH\_ECDSA

In ECDH\_ECDSA, the server's certificate MUST contain an ECDH-capable public key and be signed with ECDSA.

A ServerKeyExchange MUST NOT be sent (the server's certificate contains all the necessary keying information required by the client to arrive at the premaster secret).

The client MUST generate an ECDH key pair on the same curve as the server's long-term public key and send its public key in the

ClientKeyExchange message (except when using client authentication algorithm ECDSA\_fixed\_ECDH or RSA\_fixed\_ECDH, in which case the modifications from section [Section 3.2](#) or [Section 3.3](#) apply).

Both client and server MUST perform an ECDH operation and use the resultant shared secret as the premaster secret. All ECDH calculations are performed as specified in [Section 5.10](#)

## [2.2](#) ECDHE\_ECDSA

In ECDHE\_ECDSA, the server's certificate MUST contain an ECDSA-capable public key and be signed with ECDSA.

The server MUST send its ephemeral ECDH public key and a specification of the corresponding curve in the ServerKeyExchange message. These parameters MUST be signed with ECDSA using the private key corresponding to the public key in the server's Certificate.

The client MUST generate an ECDH key pair on the same curve as the server's ephemeral ECDH key and send its public key in the ClientKeyExchange message.

Both client and server MUST perform an ECDH operation ([Section 5.10](#)) and use the resultant shared secret as the premaster secret.

## [2.3](#) ECDH\_RSA

This key exchange algorithm is the same as ECDH\_ECDSA except the server's certificate MUST be signed with RSA rather than ECDSA.

## [2.4](#) ECDHE\_RSA

This key exchange algorithm is the same as ECDHE\_ECDSA except the server's certificate MUST contain an RSA public key authorized for signing and the signature in the ServerKeyExchange message MUST be computed with the corresponding RSA private key. The server certificate MUST be signed with RSA.

## [2.5](#) ECDH\_anon

In ECDH\_anon, the server's Certificate, the CertificateRequest, the client's Certificate, and the CertificateVerify messages MUST NOT be sent.

The server MUST send an ephemeral ECDH public key and a specification of the corresponding curve in the ServerKeyExchange message. These parameters MUST NOT be signed.



The client MUST generate an ECDH key pair on the same curve as the server's ephemeral ECDH key and send its public key in the ClientKeyExchange message.

Both client and server MUST perform an ECDH operation and use the resultant shared secret as the premaster secret. All ECDH calculations are performed as specified in [Section 5.10](#)

### [3.](#) Client Authentication

This document defines three new client authentication mechanisms named after the type of client certificate involved: ECDSA\_sign, ECDSA\_fixed\_ECDH and RSA\_fixed\_ECDH. The ECDSA\_sign mechanism is usable with any of the non-anonymous ECC key exchange algorithms described in [Section 2](#) as well as other non-anonymous (non-ECC) key exchange algorithms defined in TLS [2]. The ECDSA\_fixed\_ECDH and RSA\_fixed\_ECDH mechanisms are usable with ECDH\_ECDSA and ECDH\_RSA. Their use with ECDHE\_ECDSA and ECDHE\_RSA is prohibited because the use of a long-term ECDH client key would jeopardize the forward secrecy property of these algorithms.

The server can request ECC-based client authentication by including one or more of these certificate types in its CertificateRequest message. The server MUST NOT include any certificate types that are prohibited for the negotiated key exchange algorithm. The client must check if it possesses a certificate appropriate for any of the methods suggested by the server and is willing to use it for authentication.

If these conditions are not met, the client should send a client Certificate message containing no certificates. In this case, the ClientKeyExchange should be sent as described in [Section 2](#) and the CertificateVerify should not be sent. If the server requires client authentication, it may respond with a fatal handshake failure alert.

If the client has an appropriate certificate and is willing to use it for authentication, it MUST send that certificate in the client's Certificate message (as per [Section 5.6](#)) and prove possession of the private key corresponding to the certified key. The process of determining an appropriate certificate and proving possession is different for each authentication mechanism and described below.

NOTE: It is permissible for a server to request (and the client to send) a client certificate of a different type than the server certificate.

#### [3.1](#) ECDSA\_sign

To use this authentication mechanism, the client MUST possess a certificate containing an ECDSA-capable public key and signed with ECDSA.

The client MUST prove possession of the private key corresponding to the certified key by including a signature in the CertificateVerify message as described in [Section 5.8](#).

### [3.2](#) ECDSA\_fixed\_ECDH

To use this authentication mechanism, the client MUST possess a certificate containing an ECDH-capable public key and that certificate MUST be signed with ECDSA. Furthermore, the client's ECDH key MUST be on the same elliptic curve as the server's long-term (certified) ECDH key.

When using this authentication mechanism, the client MUST send an empty ClientKeyExchange as described in [Section 5.7](#) and MUST NOT send the CertificateVerify message. The ClientKeyExchange is empty since the client's ECDH public key required by the server to compute the premaster secret is available inside the client's certificate. The client's ability to arrive at the same premaster secret as the server (demonstrated by a successful exchange of Finished messages) proves possession of the private key corresponding to the certified public key and the CertificateVerify message is unnecessary.

### [3.3](#) RSA\_fixed\_ECDH

This authentication mechanism is identical to ECDSA\_fixed\_ECDH except the client's certificate MUST be signed with RSA.

#### [4.](#) TLS Extensions for ECC

Two new TLS extensions --- (i) the Supported Elliptic Curves Extension, and (ii) the Supported Point Formats Extension --- allow a client to negotiate the use of specific curves and point formats (e.g. compressed v/s uncompressed), respectively. These extensions are especially relevant for constrained clients that may only support a limited number of curves or point formats. They follow the general approach outlined in [\[3\]](#). The client enumerates the curves and point formats it supports by including the appropriate extensions in its ClientHello message. By echoing that extension in its ServerHello, the server agrees to restrict its key selection or encoding to the choices specified by the client.

A TLS client that proposes ECC cipher suites in its ClientHello message SHOULD include these extensions. Servers implementing ECC cipher suites MUST support these extensions and negotiate the use of an ECC cipher suite only if they can complete the handshake while limiting themselves to the curves and compression techniques enumerated by the client. This eliminates the possibility that a negotiated ECC handshake will be subsequently aborted due to a client's inability to deal with the server's EC key.

These extensions MUST NOT be included if the client does not propose any ECC cipher suites. A client that proposes ECC cipher suites may choose not to include these extension. In this case, the server is free to choose any one of the elliptic curves or point formats listed in [Section 5](#). That section also describes the structure and processing of these extensions in greater detail.

## [5.](#) Data Structures and Computations

This section specifies the data structures and computations used by ECC-based key mechanisms specified in [Section 2](#), [Section 3](#) and [Section 4](#). The presentation language used here is the same as that used in TLS [[2](#)]. Since this specification extends TLS, these descriptions should be merged with those in the TLS specification and any others that extend TLS. This means that enum types may not specify all possible values and structures with multiple formats chosen with a `select()` clause may not indicate all possible cases.

### [5.1](#) Client Hello Extensions

When this message is sent:

The ECC extensions SHOULD be sent along with any ClientHello message that proposes ECC cipher suites.

Meaning of this message:

These extensions allow a constrained client to enumerate the elliptic curves and/or point formats it supports.

Structure of this message:

The general structure of TLS extensions is described in [3] and this specification adds two new types to ExtensionType.

```
enum { elliptic_curves(6), ec_point_formats(7) } ExtensionType;
```

`elliptic_curves`: Indicates the set of elliptic curves supported by the client. For this extension, the opaque `extension_data` field contains `EllipticCurveList`.

`ec_point_formats`: Indicates the set of point formats supported by the client. For this extension, the opaque `extension_data` field contains `ECPointFormatList`.

```
enum {
    sect163k1 (1), sect163r1 (2), sect163r2 (3),
    sect193r1 (4), sect193r2 (5), sect233k1 (6),
    sect233r1 (7), sect239k1 (8), sect283k1 (9),
    sect283r1 (10), sect409k1 (11), sect409r1 (12),
    sect571k1 (13), sect571r1 (14), secp160k1 (15),
    secp160r1 (16), secp160r2 (17), secp192k1 (18),
    secp192r1 (19), secp224k1 (20), secp224r1 (21),
    secp256k1 (22), secp256r1 (23), secp384r1 (24),
    secp521r1 (25), reserved (240..247),
    arbitrary_explicit_prime_curves(253),
    arbitrary_explicit_char2_curves(254),
    (255)
} NamedCurve;
```

sect163k1, etc: Indicates support of the corresponding named curve specified in SEC 2 [10]. Note that many of these curves are also recommended in ANSI X9.62 [6], and FIPS 186-2 [8]. Values 240 through 247 are reserved for private use. Values 253 and 254 indicate that the client supports arbitrary prime and charactersitic two curves, respectively (the curve parameters must be encoded explicitly in ECParameters).

```
struct {
    NamedCurve elliptic_curve_list<1..2^16-1>
} EllipticCurveList;
```

Items in `elliptic_curve_list` are ordered according to the client's preferences (favorite choice first).

As an example, a client that only supports `secp192r1` (aka NIST P-192) and `secp224r1` (aka NIST P-224) and prefers to use `secp192r1`, would include an `elliptic_curves` extension with the following octets:

```
00 06 00 02 13 15
```

A client that supports arbitrary explicit binary polynomial curves would include an extension with the following octets:

```
00 06 00 01 fe
```

```
enum { uncompressed (0), ansiX963_compressed (1), ansiX963_hybrid (2) }
ECPointFormat;
```

```
struct {
    ECPointFormat ec_point_format_list<1..2^16-1>
} ECPointFormatList;
```

Items in `ec_point_format_list` are ordered according to the client's

preferences (favorite choice first).

A client that only supports the uncompressed point format includes an extension with the following octets:

```
00 07 00 01 00
```

A client that prefers the use of the `ansiX963_compressed` format over uncompressed may indicate that preference by including an extension with the following octets:

```
00 07 00 02 01 00
```

Actions of the sender:

A client that proposes ECC cipher suites in its `ClientHello` appends these extensions (along with any others) enumerating the curves and point formats it supports.

Actions of the receiver:

A server that receives a `ClientHello` containing one or both of these extensions **MUST** use the client's enumerated capabilities to guide its selection of an appropriate cipher suite. One of the proposed ECC cipher suites must be negotiated only if the server can successfully complete the handshake while using the curves and point formats supported by the client.

NOTE: A server participating in an ECDHE-ECDSA key exchange may use different curves for (i) the ECDSA key in its certificate, and (ii) the ephemeral ECDH key in the `ServerKeyExchange` message. The server must consider the "elliptic\_curves" extension in selecting both of these curves.

If a server does not understand the "elliptic\_curves" extension or is unable to complete the ECC handshake while restricting itself to the enumerated curves, it **MUST NOT** negotiate the use of an ECC cipher suite. Depending on what other cipher suites are proposed by the client and supported by the server, this may result in a fatal handshake failure alert due to the lack of common cipher suites.



When this message is sent:

The ServerHello ECC extensions are sent in response to a Client Hello message containing ECC extensions when negotiating an ECC cipher suite.

Meaning of this message:

These extensions indicate the server's agreement to use only the elliptic curves and point formats supported by the client during the ECC-based key exchange.

Structure of this message:

The ECC extensions echoed by the server are the same as those in the ClientHello except the "extension\_data" field is empty.

For example, a server indicates its acceptance of the client's elliptic\_curves extension by sending an extension with the following octets:

```
00 06 00 00
```

Actions of the sender:

A server makes sure that it can complete a proposed ECC key exchange mechanism by restricting itself to the curves/point formats supported by the client before sending these extensions.

Actions of the receiver:

A client that receives a ServerHello with ECC extensions proceeds with an ECC key exchange assured that it will be able to handle the server's EC key(s).

### [5.3](#) Server Certificate

When this message is sent:

This message is sent in all non-anonymous ECC-based key exchange algorithms.

Meaning of this message:

This message is used to authentically convey the server's static public key to the client. The following table shows the server

certificate type appropriate for each key exchange algorithm. ECC public keys must be encoded in certificates as described in [Section 5.9](#).

NOTE: The server's Certificate message is capable of carrying a chain of certificates. The restrictions mentioned in Table 3 apply only to the server's certificate (first in the chain).

Key Exchange Algorithm	Server Certificate Type
-----	-----
ECDH_ECDSA	Certificate must contain an ECDH-capable public key. It must be signed with ECDSA.
ECDHE_ECDSA	Certificate must contain an ECDSA-capable public key. It must be signed with ECDSA.
ECDH_RSA	Certificate must contain an ECDH-capable public key. It must be signed with RSA.
ECDHE_RSA	Certificate must contain an RSA public key authorized for use in digital signatures. It must be signed with RSA.

Table 3: Server certificate types

Structure of this message:

Identical to the TLS Certificate format.

Actions of the sender:

The server constructs an appropriate certificate chain and conveys it to the client in the Certificate message.

Actions of the receiver:

The client validates the certificate chain, extracts the server's public key, and checks that the key type is appropriate for the negotiated key exchange algorithm.

#### [5.4](#) Server Key Exchange

When this message is sent:

This message is sent when using the ECDHE\_ECDSA, ECDHE\_RSA and ECDH\_anon key exchange algorithms.

Meaning of this message:

This message is used to convey the server's ephemeral ECDH public key (and the corresponding elliptic curve domain parameters) to the client.

Structure of this message:

```
enum { explicit_prime (1), explicit_char2 (2),
        named_curve (3), (255) } ECCurveType;
```

`explicit_prime`: Indicates the elliptic curve domain parameters are conveyed verbosely, and the underlying finite field is a prime field.

`explicit_char2`: Indicates the elliptic curve domain parameters are conveyed verbosely, and the underlying finite field is a characteristic 2 field.

`named_curve`: Indicates that a named curve is used. This option SHOULD be used when applicable.

```
struct {
    opaque a <1..2^8-1>;
    opaque b <1..2^8-1>;
} ECCurve;
```

a, b: These parameters specify the coefficients of the elliptic curve. Each value contains the byte string representation of a field element following the conversion routine in [Section 4.3.3](#) of ANSI X9.62 [6].

```
struct {
    opaque point <1..2^8-1>;
} ECPoint;
```

point: This is the byte string representation of an elliptic curve point following the conversion routine in [Section 4.3.6](#) of ANSI X9.62 [6]. Note that this byte string may represent an elliptic

curve point in compressed or uncompressed form. Implementations of this specification MUST support the uncompressed form and MAY support the compressed form.

```
enum { ec_basis_trinomial, ec_basis_pentanomial } ECBasisType;
```

ec\_basis\_trinomial: Indicates representation of a characteristic two field using a trinomial basis.

ec\_basis\_pentanomial: Indicates representation of a characteristic two field using a pentanomial basis.

```
struct {
    ECCurveType    curve_type;
    select (curve_type) {
        case explicit_prime:
            opaque    prime_p <1..2^8-1>;
            ECCurve    curve;
            ECPoint    base;
            opaque    order <1..2^8-1>;
            opaque    cofactor <1..2^8-1>;
        case explicit_char2:
            uint16    m;
            ECBasisType basis;
            select (basis) {
                case ec_trinomial:
                    opaque    k <1..2^8-1>;
                case ec_pentanomial:
                    opaque    k1 <1..2^8-1>;
                    opaque    k2 <1..2^8-1>;
            }
    }
}
```

```

        opaque k3 <1..2^8-1>;
    };
    ECCurve      curve;
    ECPoint      base;
    opaque       order <1..2^8-1>;
    opaque       cofactor <1..2^8-1>;
    case named_curve:
        NamedCurve namedcurve;
    };
} ECPParameters;

```

curve\_type: This identifies the type of the elliptic curve domain parameters.

prime\_p: This is the odd prime defining the field  $F_p$ .

curve: Specifies the coefficients  $a$  and  $b$  of the elliptic curve  $E$ .

base: Specifies the base point  $G$  on the elliptic curve.

order: Specifies the order  $n$  of the base point.

cofactor: Specifies the cofactor  $h = \#E(F_q)/n$ , where  $\#E(F_q)$  represents the number of points on the elliptic curve  $E$  defined over the field  $F_q$ .

$m$ : This is the degree of the characteristic-two field  $F_{2^m}$ .

$k$ : The exponent  $k$  for the trinomial basis representation  $x^m + x^k + 1$ .

$k_1, k_2, k_3$ : The exponents for the pentanomial representation  $x^m + x^{k_3} + x^{k_2} + x^{k_1} + 1$  (such that  $k_3 > k_2 > k_1$ ).

namedcurve: Specifies a recommended set of elliptic curve domain parameters. All enum values of NamedCurve are allowed except for arbitrary\_explicit\_prime\_curves(253) and arbitrary\_explicit\_char2\_curves(254). These two values are only allowed in the ClientHello extension.

```
struct {
    ECParameters    curve_params;
    ECPublicKey    public;
} ServerECDHParams;
```

`curve_params`: Specifies the elliptic curve domain parameters associated with the ECDH public key.

`public`: The ephemeral ECDH public key.

The `ServerKeyExchange` message is extended as follows.

```
enum { ec_diffie_hellman } KeyExchangeAlgorithm;
```

`ec_diffie_hellman`: Indicates the `ServerKeyExchange` message contains an ECDH public key.

```
select (KeyExchangeAlgorithm) {
    case ec_diffie_hellman:
        ServerECDHParams    params;
        Signature            signed_params;
} ServerKeyExchange;
```

`params`: Specifies the ECDH public key and associated domain parameters.

`signed_params`: A hash of the `params`, with the signature appropriate to that hash applied. The private key corresponding to the certified public key in the server's `Certificate` message is used for signing.

```
enum { ecdsa } SignatureAlgorithm;
```

```
select (SignatureAlgorithm) {
  case ecdsa:
    digitally-signed struct {
      opaque sha_hash[sha_size];
    };
} Signature;
```

NOTE: SignatureAlgorithm is 'rsa' for the ECDHE\_RSA key exchange algorithm and 'anonymous' for ECDH\_anon. These cases are defined in TLS [2]. SignatureAlgorithm is 'ecdsa' for ECDHE\_ECDSA. ECDSA signatures are generated and verified as described in [Section 5.10](#). As per ANSI X9.62, an ECDSA signature consists of a pair of integers r and s. These integers are both converted into byte strings of the same length as the curve order n using the conversion routine specified in Section 4.3.1 of [6]. The two byte strings are concatenated, and the result is placed in the signature field.

Actions of the sender:

The server selects elliptic curve domain parameters and an ephemeral ECDH public key corresponding to these parameters according to the ECKAS-DH1 scheme from IEEE 1363 [5]. It conveys this information to the client in the ServerKeyExchange message using the format defined above.

Actions of the recipient:

The client verifies the signature (when present) and retrieves the server's elliptic curve domain parameters and ephemeral ECDH public key from the ServerKeyExchange message.

## [5.5](#) Certificate Request

When this message is sent:

This message is sent when requesting client authentication.

Meaning of this message:

The server uses this message to suggest acceptable client authentication methods.

Structure of this message:

The TLS CertificateRequest message is extended as follows.

```
enum {  
    ecdsa_sign(?), rsa_fixed_ecdh(?),  
    ecdsa_fixed_ecdh(?), (255)  
} ClientCertificateType;
```

ecdsa\_sign, etc Indicates that the server would like to use the corresponding client authentication method specified in [Section 3](#).

EDITOR: The values used for ecdsa\_sign, rsa\_fixed\_ecdh, and ecdsa\_fixed\_ecdh have been left as ?. These values will be assigned when this draft progresses to RFC. Earlier versions of this draft used the values 5, 6, and 7 - however these values have been removed since they are used differently by SSL 3.0 [[13](#)] and their use by TLS is being deprecated.

Actions of the sender:

The server decides which client authentication methods it would like to use, and conveys this information to the client using the format defined above.

Actions of the receiver:

The client determines whether it has an appropriate certificate for use with any of the requested methods, and decides whether or not to proceed with client authentication.

## [5.6](#) Client Certificate

When this message is sent:

This message is sent in response to a CertificateRequest when a client has a suitable certificate.

Meaning of this message:

This message is used to authentically convey the client's static public key to the server. The following table summarizes what client



certificate types are appropriate for the ECC-based client authentication mechanisms described in [Section 3](#). ECC public keys must be encoded in certificates as described in [Section 5.9](#).

NOTE: The client's Certificate message is capable of carrying a chain of certificates. The restrictions mentioned in Table 4 apply only to the client's certificate (first in the chain).

Client Authentication Method -----	Client Certificate Type -----
ECDSA_sign	Certificate must contain an ECDSA-capable public key and be signed with ECDSA.
ECDSA_fixed_ECDH	Certificate must contain an ECDH-capable public key on the same elliptic curve as the server's long-term ECDH key. This certificate must be signed with ECDSA.
RSA_fixed_ECDH	Certificate must contain an ECDH-capable public key on the same elliptic curve as the server's long-term ECDH key. This certificate must be signed with RSA.

Table 4: Client certificate types

Structure of this message:

Identical to the TLS client Certificate format.

Actions of the sender:

The client constructs an appropriate certificate chain, and conveys it to the server in the Certificate message.

Actions of the receiver:

The TLS server validates the certificate chain, extracts the client's

public key, and checks that the key type is appropriate for the client authentication method.

### [5.7](#) Client Key Exchange

When this message is sent:

This message is sent in all key exchange algorithms. If client authentication with ECDSA\_fixed\_ECDH or RSA\_fixed\_ECDH is used, this message is empty. Otherwise, it contains the client's ephemeral ECDH public key.

Meaning of the message:

This message is used to convey ephemeral data relating to the key exchange belonging to the client (such as its ephemeral ECDH public key).

Structure of this message:

The TLS ClientKeyExchange message is extended as follows.

```
enum { yes, no } EphemeralPublicKey;
```

yes, no: Indicates whether or not the client is providing an ephemeral ECDH public key. (In ECC ciphersuites, this is "yes" except when the client uses the ECDSA\_fixed\_ECDH or RSA\_fixed\_ECDH client authentication mechanism.)

```
struct {
    select (EphemeralPublicKey) {
        case yes: ECPoint ecdh_Yc;
        case no: struct { };
    } ecdh_public;
} ClientECDiffieHellmanPublic;
```

ecdh\_Yc: Contains the client's ephemeral ECDH public key.

```
struct {
    select (KeyExchangeAlgorithm) {
        case ec_diffie_hellman: ClientECDiffieHellmanPublic;
    } exchange_keys;
} ClientKeyExchange;
```

Actions of the sender:

The client selects an ephemeral ECDH public key corresponding to the parameters it received from the server according to the ECKAS-DH1 scheme from IEEE 1363 [5]. It conveys this information to the client in the ClientKeyExchange message using the format defined above.

Actions of the recipient:

The server retrieves the client's ephemeral ECDH public key from the ClientKeyExchange message and checks that it is on the same elliptic curve as the server's ECDH key.

### [5.8](#) Certificate Verify

When this message is sent:

This message is sent when the client sends a client certificate containing a public key usable for digital signatures, e.g. when the client is authenticated using the ECDSA\_sign mechanism.

Meaning of the message:

This message contains a signature that proves possession of the private key corresponding to the public key in the client's Certificate message.

Structure of this message:

The TLS CertificateVerify message is extended as follows.

```
enum { ecdsa } SignatureAlgorithm;

select (SignatureAlgorithm) {
    case ecdsa:
        digitally-signed struct {
            opaque sha_hash[sha_size];
        };
} Signature;
```

For the ecdsa case, the signature field in the CertificateVerify message contains an ECDSA signature computed over handshake messages

exchanged so far. ECDSA signatures are computed as described in [Section 5.10](#). As per ANSI X9.62, an ECDSA signature consists of a pair of integers  $r$  and  $s$ . These integers are both converted into byte strings of the same length as the curve order  $n$  using the conversion routine specified in Section 4.3.1 of [6]. The two byte strings are concatenated, and the result is placed in the signature field.

#### Actions of the sender:

The client computes its signature over all handshake messages sent or received starting at client hello up to but not including this message. It uses the private key corresponding to its certified public key to compute the signature which is conveyed in the format defined above.

#### Actions of the receiver:

The server extracts the client's signature from the CertificateVerify message, and verifies the signature using the public key it received in the client's Certificate message.

### [5.9](#) Elliptic Curve Certificates

X509 certificates containing ECC public keys or signed using ECDSA MUST comply with [11] or another RFC that replaces or extends it. Clients SHOULD use the elliptic curve domain parameters recommended in ANSI X9.62 [6], FIPS 186-2 [8], and SEC 2 [10].

### [5.10](#) ECDH, ECDSA and RSA Computations

All ECDH calculations (including parameter and key generation as well as the shared secret calculation) MUST be performed according to [5] using

- o the ECKAS-DH1 scheme with the ECSVDP-DH secret value derivation primitive, the KDF1 key derivation function using SHA-1 [7], and null key derivation parameters "P" for elliptic curve parameters where field elements are represented as octet strings of length 24 or less (using the IEEE 1363 FE2OSP); in this case, the premaster secret is the output of the ECKAS-DH1 scheme, i.e. the 20-byte

SHA-1 output from the KDF.

- o the ECKAS-DH1 scheme with the identity map as key derivation function for elliptic curve parameters where field elements are represented as octet strings of length more than 24; in this case, the premaster secret is the x-coordinate of the ECDH shared secret elliptic curve point, i.e. the octet string Z in IEEE 1363 terminology.

Note that a new extension may be introduced in the future to allow the use of a different KDF during computation of the premaster secret. In this event, the new KDF would be used in place of the process detailed above. This may be desirable, for example, to support compatibility with the planned NIST key agreement standard.

All ECDSA computations MUST be performed according to ANSI X9.62 [6] or its successors. Data to be signed/verified is hashed and the result run directly through the ECDSA algorithm with no additional hashing. The default hash function is SHA-1 [7] and sha\_size (see [Section 5.4](#) and [Section 5.8](#)) is 20. However, an alternative hash function, such as one of the new SHA hash functions specified in FIPS 180-2 [7], may be used instead if the certificate containing the EC public key explicitly requires use of another hash function. (The mechanism for specifying the required hash function has not been standardized but this provision anticipates such standardization and obviates the need to update this document in response. Future PKIX RFCs may choose, for example, to specify the hash function to be used with a public key in the parameters field of subjectPublicKeyInfo.)

All RSA signatures must be generated and verified according to PKCS#1 [9].

## 6. Cipher Suites

The table below defines new ECC cipher suites that use the key exchange algorithms specified in [Section 2](#).

EDITOR: Most of the cipher suites below have been left as ??. The values 47-4C correspond to cipher suites which are known to have been implemented and are therefore proposed here. The final determination of cipher suite numbers will occur when this draft progresses to RFC. Implementers using the values 47-4C should therefore be wary that these values may change.

CipherSuite TLS_ECDH_ECDSA_WITH_NULL_SHA	= { 0x00, 0x47 }
CipherSuite TLS_ECDH_ECDSA_WITH_RC4_128_SHA	= { 0x00, 0x48 }
CipherSuite TLS_ECDH_ECDSA_WITH_DES_CBC_SHA	= { 0x00, 0x49 }
CipherSuite TLS_ECDH_ECDSA_WITH_3DES_EDE_CBC_SHA	= { 0x00, 0x4A }
CipherSuite TLS_ECDH_ECDSA_WITH_AES_128_CBC_SHA	= { 0x00, 0x4B }
CipherSuite TLS_ECDH_ECDSA_WITH_AES_256_CBC_SHA	= { 0x00, 0x4C }

CipherSuite TLS_ECDHE_ECDSA_WITH_NULL_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDHE_ECDSA_WITH_RC4_128_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDHE_ECDSA_WITH_3DES_EDE_CBC_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDHE_ECDSA_WITH_AES_128_CBC_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDHE_ECDSA_WITH_AES_256_CBC_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDH_RSA_WITH_NULL_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDH_RSA_WITH_RC4_128_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDH_RSA_WITH_3DES_EDE_CBC_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDH_RSA_WITH_AES_128_CBC_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDH_RSA_WITH_AES_256_CBC_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDHE_RSA_WITH_NULL_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDHE_RSA_WITH_RC4_128_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDHE_RSA_WITH_3DES_EDE_CBC_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDHE_RSA_WITH_AES_128_CBC_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDHE_RSA_WITH_AES_256_CBC_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDH_anon_NULL_WITH_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDH_anon_WITH_RC4_128_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDH_anon_WITH_3DES_EDE_CBC_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDH_anon_WITH_AES_128_CBC_SHA	= { 0x00, 0x?? }
CipherSuite TLS_ECDH_anon_WITH_AES_256_CBC_SHA	= { 0x00, 0x?? }

Table 5: TLS ECC cipher suites

The key exchange method, cipher, and hash algorithm for each of these cipher suites are easily determined by examining the name. Ciphers

other than AES ciphers, and hash algorithms are defined in [2]. AES ciphers are defined in [14].

Server implementations SHOULD support all of the following cipher suites, and client implementations SHOULD support at least one of them: TLS\_ECDH\_ECDSA\_WITH\_3DES\_EDE\_CBC\_SHA, TLS\_ECDH\_ECDSA\_WITH\_AES\_128\_CBC\_SHA, TLS\_ECDHE\_RSA\_WITH\_3DES\_EDE\_CBC\_SHA, and TLS\_ECDHE\_RSA\_WITH\_AES\_128\_CBC\_SHA.

## [7](#). Security Considerations

This document is based on [\[2\]](#), [\[5\]](#), [\[6\]](#) and [\[14\]](#). The appropriate security considerations of those documents apply.



For ECDH ([Section 5.10](#)), this document specifies two different ways to compute the premaster secret. The choice of the method is determined by the elliptic curve. Earlier versions of this specification used the KDF1 key derivation function with SHA-1 in all cases; the current version keeps this key derivation function only for curves where field elements are represented as octet strings of length 24 or less (i.e. up to 192 bits), but omits it for larger curves.

Rationale: Using KDF1 with SHA-1 limits the security to at most 160 bits, independently of the elliptic curve used for ECDH. For large curves, this would result in worse security than expected. Using a specific key derivation function for ECDH is not really necessary as TLS always uses its PRF to derive the master secret from the premaster secret. For large curves, the current specification handles ECDH like the basic TLS specification [\[14\]](#) handles standard DH. For smaller curves where the extra KDF1 step does not weaken security, the current specification keeps the KDF1 step to obtain compatibility with existing implementations of earlier versions of this specification. Note that the threshold for switching between the two ECDH calculation methods is necessarily somewhat arbitrary; 192-bit ECC corresponds to approximately 96 bits of security in the light of square root attacks, so the 160 bits provided by SHA-1 are comfortable at this limit.

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## [9.](#) Acknowledgments

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