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**Exported Authenticators in TLS**  
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Abstract

This document describes a mechanism in Transport Layer Security (TLS) to provide an exportable proof of ownership of a certificate that can be transmitted out of band and verified by the other party.

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## [1.](#) Introduction

This document provides a way to authenticate one party of a Transport Layer Security (TLS) communication to another using a certificate after the session has been established. This allows both the client and server to prove ownership of additional identities at any time after the handshake has completed. This proof of authentication can be exported and transmitted out of band from one party to be validated by the other party.

This mechanism provides two advantages over the authentication that TLS natively provides:

multiple identities - Endpoints that are authoritative for multiple identities - but do not have a single certificate that includes all of the identities - can authenticate with those identities over a single connection.

spontaneous authentication - Endpoints can authenticate after a connection is established, in response to events in a higher-layer protocol, as well as integrating more context.



This document intends to replace much of the functionality of renegotiation in previous versions of TLS. It has the advantages over renegotiation of not requiring additional on-the-wire changes during a connection. For simplicity, only TLS 1.2 and later are supported.

Post-handshake authentication is defined in TLS 1.3, but it has the disadvantage of requiring additional state to be stored in the TLS state machine and it composes poorly with multiplexed connection protocols like HTTP/2 [[RFC7540](#)]. It is also only available for client authentication. This mechanism is intended to be used as part of a replacement for post-handshake authentication in applications.

## **2. Conventions and Terminology**

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP 14](#) [[RFC2119](#)] [[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.

## **3. Authenticator Request**

The authenticator request is a structured message that can be exported from either party of a TLS connection. It can be transmitted to the other party of the TLS connection at the application layer. The application layer protocol used to send the authenticator SHOULD use TLS as its underlying transport to keep the request confidential.

An authenticator request message can be constructed by either the client or the server. This authenticator request uses the CertificateRequest message structure from Section 4.3.2 of [[TLS13](#)]. This message does not include the TLS record layer and is therefore not encrypted with a handshake key.

The CertificateRequest is used to define the parameters in a request for an authenticator. The definition for TLS 1.3 is:

```
struct {  
    opaque certificate_request_context<0..2^8-1>;  
    Extension extensions<2..2^16-1>;  
} CertificateRequest;
```

certificate\_request\_context: An opaque string which identifies the certificate request and which will be echoed in the authenticator message. A certificate\_request\_context value MUST NOT be repeated for multiple authenticator requests generated by the same party



within the scope of a connection (thus preventing replay of authenticators). The `certificate_request_context` SHOULD be chosen to be unpredictable to the peer (e.g., by randomly generating it) in order to prevent an attacker who has temporary access to the peer's private key from pre-computing valid authenticators. This value is unrelated to the `certificate_request_context` used in post-handshake authentication and collisions do not need to be avoided.

extensions: The extensions that are allowed in this structure include the extensions defined for `CertificateRequest` messages defined in Section 4.2. of [TLS13] and the `server_name` [RFC6066] extension, which is allowed for client-generated authenticator requests.

## **4. Authenticator**

The authenticator is a structured message that can be exported from either party of a TLS connection. It can be transmitted to the other party of the TLS connection at the application layer. The application layer protocol used to send the authenticator SHOULD use TLS as its underlying transport to keep the certificate confidential.

An authenticator message can be constructed by either the client or the server given an established TLS connection, a certificate, and a corresponding private key. Clients MUST NOT send an authenticator without a preceding authenticator request; for servers an authenticator request is optional.

The authenticator uses the message structures from Section 4.4 of [TLS13], but different parameters. These messages do not include the TLS record layer and are therefore not encrypted with a handshake key.

### **4.1. Authenticator Keys**

Each authenticator is computed using a Handshake Context and Finished MAC Key derived from the TLS session. These values are derived using an exporter as described in [RFC5705] (for TLS 1.2) or [TLS13] (for TLS 1.3). These values use different labels depending on the role of the sender:

- o The Handshake Context is an exporter value that is derived using the label "EXPORTER-client authenticator handshake context" or "EXPORTER-server authenticator handshake context" for authenticators sent by the client and server respectively.
- o The Finished MAC Key is an exporter value derived using the label "EXPORTER-client authenticator finished key" or "EXPORTER-server



authenticator finished key" for authenticators sent by the client and server respectively.

The context\_value used for the exporter is absent (length zero) for all four values. The length of the exported value is equal to the length of the output of the hash function selected in TLS for the pseudorandom function (PRF). Cipher suites that do not use the TLS PRF MUST define a hash function that can be used for this purpose or they cannot be used. For TLS 1.3 symmetric cipher suites, the hash algorithm used with HKDF is used.

If the connection is TLS 1.2, the master secret MUST have been computed with the extended master secret [RFC7627] to avoid key synchronization attacks.

## **4.2. Authenticator Construction**

An authenticator is formed from the concatenation of TLS 1.3 [TLS13] Certificate, CertificateVerify, and Finished messages.

If an authenticator request is present, the extensions used to guide the construction of these messages are taken from the authenticator request. If there is no authenticator request, the extensions are chosen from the TLS handshake. Only servers can provide an authenticator without a corresponding request. In such cases, ClientHello extensions are used to determine permissible extensions in the Certificate message.

### **4.2.1. Certificate**

The certificate to be used for authentication and any supporting certificates in the chain. This structure is defined in [TLS13], Section 4.4.2.

The certificate message contains an opaque string called certificate\_request\_context, which is extracted from the authenticator request if present. If no authenticator request is provided, the certificate\_request\_context can be chosen arbitrarily.

The certificates chosen in the Certificate message MUST conform to the requirements of a Certificate message in the negotiated version of TLS. In particular, the certificate chain MUST be valid for the signature algorithms indicated by the peer in the "signature\_algorithms" and "signature\_algorithms\_cert" extension, as described in Section 4.2.3 of [TLS13] for TLS 1.3 or the "signature\_algorithms" extension from Sections 7.4.2 and 7.4.6 of [RFC5246] for TLS 1.2.





In addition to "signature\_algorithms" and "signature\_algorithms\_cert", the "server\_name" [[RFC6066](#)], "certificate\_authorities" (Section 4.2.4. of [[TLS13](#)]), and "oid\_filters" (Section 4.2.5. of [[TLS13](#)]) extensions are used to guide certificate selection. The extensions, or others that might affect certificate selection, are taken from the authenticator request if present, or the TLS handshake if not.

Alternative certificate formats such as [[RFC7250](#)] Raw Public Keys are not supported in this version of the specification.

If an authenticator request was provided, the Certificate message MUST contain only extensions present in the authenticator request. Otherwise, the Certificate message MUST contain only extensions present in the TLS handshake.

#### **4.2.2. CertificateVerify**

This message is used to provide explicit proof that an endpoint possesses the private key corresponding to its certificate. The definition for TLS 1.3 is:

```
struct {  
    SignatureScheme algorithm;  
    opaque signature<0..2^16-1>;  
} CertificateVerify;
```

The algorithm field specifies the signature algorithm used (see Section 4.2.3 of [[TLS13](#)] for the definition of this field). The signature is a digital signature using that algorithm.

The signature scheme MUST be a valid signature scheme for TLS 1.3. This excludes all RSASSA-PKCS1-v1\_5 algorithms and combinations of ECDSA and hash algorithms that are not supported in TLS 1.3.

If an authenticator request is present, the signature algorithm MUST be chosen from one of the signature schemes in the authenticator request. Otherwise, the signature algorithm used should be chosen from the "signature\_algorithms" sent by the peer in the TLS handshake.

The signature is computed using the over the concatenation of:

- o A string that consists of octet 32 (0x20) repeated 64 times
- o The context string "Exported Authenticator" (which is not NULL-terminated)



- o A single 0 byte which serves as the separator
- o The hashed authenticator transcript

The authenticator transcript is the hash of the concatenated Handshake Context, authenticator request (if present), and Certificate message:

```
Hash(Handshake Context || authenticator request || Certificate)
```

Where Hash is the hash function negotiated by TLS. If the authenticator request is not present, it is omitted from this construction (that is, it is zero length).

If the party that generates the exported authenticator does so with a different connection than the party that is validating it, then the Handshake Context will not match, resulting in a CertificateVerify message that does not validate. This includes situations in which the application data is sent via TLS-terminating proxy. Given a failed CertificateVerify validation, it may be helpful for the application to confirm that both peers share the same connection using a value derived from the connection secrets before taking a user-visible action.

#### **4.2.3. Finished**

A HMAC [[HMAC](#)] over the hashed authenticator transcript, which is the concatenated Handshake Context, authenticator request (if present), Certificate, and CertificateVerify:

```
Hash(Handshake Context || authenticator request ||  
      Certificate || CertificateVerify)
```

The HMAC is computed using the same hash function using the Finished MAC Key as a key.

#### **4.2.4. Authenticator Creation**

An endpoint constructs an authenticator by serializing the Certificate, CertificateVerify, and Finished as TLS handshake messages and concatenating the octets:

```
Certificate || CertificateVerify || Finished
```

A given authenticator can be validated by checking the validity of the CertificateVerify message given the authenticator request (if used) and recomputing the Finished message to see if it matches.



## 5. Empty Authenticator

If, given an authenticator request, the endpoint does have an appropriate certificate or does not want to return one, it constructs an authenticated refusal called an empty authenticator. This is an HMAC over the hashed authenticator transcript with a Certificate message containing no CertificateEntries and the CertificateVerify message omitted:

```
"Hash(Handshake Context || authenticator request || Certificate) "
```

The HMAC is computed using the same hash function using the Finished MAC Key as a key.

## 6. API considerations

The creation and validation of both authenticator requests and authenticators SHOULD be implemented inside the TLS library even if it is possible to implement it at the application layer. TLS implementations supporting the use of exported authenticators MUST provide application programming interfaces by which clients and servers may request and verify exported authenticator messages.

Notwithstanding the success conditions described below, all APIs MUST fail if:

- o the connection uses a TLS version of 1.1 or earlier, or
- o the connection is TLS 1.2 and the extended master secret [[RFC7627](#)] was not used

The following sections describes APIs that are considered necessary to implement exported authenticators. These are informative only.

### 6.1. The "request" API

The "request" API takes as input:

- o certificate\_request\_context (from 0 to 255 bytes)
- o set of extensions to include (this MUST include signature\_algorithms)

It returns an authenticator request, which is a sequence of octets that includes a CertificateRequest message.



### **6.2. The "get context" API**

The "get context" API takes as input:

- o authenticator

It returns the `certificate_request_context`.

### **6.3. The "authenticate" API**

The "authenticate" takes as input:

- o a set of certificate chains and associated extensions (OCSP, SCT, etc.)
- o a signer (either the private key associated with the certificate, or interface to perform private key operation) for each chain
- o an optional authenticator request or `certificate_request_context` (from 0 to 255 bytes)

It returns either the exported authenticator or an empty authenticator as a sequence of octets. It is RECOMMENDED that the logic for selecting the certificates and extensions to include in the exporter is implemented in the TLS library. Implementing this in the TLS library lets the implementer take advantage of existing extension and certificate selection logic.

It is also possible to implement this API outside of the TLS library using TLS exporters. This may be preferable in cases where the application does not have access to a TLS library with these APIs or when TLS is handled independently of the application layer protocol.

### **6.4. The "validate" API**

The "validate" API takes as input:

- o an optional authenticator request
- o an authenticator

It returns the certificate chain and extensions and a status to indicate whether the authenticator is valid or not. If the authenticator was empty - that is, it did not contain a certificate - the certificate chain will contain no certificates.





## **7. IANA Considerations**

This document has no IANA actions.

## **8. Security Considerations**

The Certificate/Verify/Finished pattern intentionally looks like the TLS 1.3 pattern which now has been analyzed several times. In the case where the client presents an authenticator to a server, [[SIGMAC](#)] presents a relevant framework for analysis.

Authenticators are independent and unidirectional. There is no explicit state change inside TLS when an authenticator is either created or validated.

- o This property makes it difficult to formally prove that a server is jointly authoritative over multiple certificates, rather than individually authoritative over each.
- o There is no indication in the TLS layer about which point in time an authenticator was computed. Any feedback about the time of creation or validation of the authenticator should be tracked as part of the application layer semantics if required.

The signatures generated with this API cover the context string "Exported Authenticator" and therefore cannot be transplanted into other protocols.

## **9. Acknowledgements**

Comments on this proposal were provided by Martin Thomson. Suggestions for [Section 8](#) were provided by Karthikeyan Bhargavan.

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