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Record Size Limit Extension for Transport Layer Security (TLS)
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Abstract

An extension to Transport Layer Security (TLS) is defined that allows endpoints to negotiate the maximum size of protected records that each will send the other.

This replaces the maximum fragment length extension defined in [RFC 6066](#).

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[1.](#) Introduction

Implementing Transport Layer Security (TLS) [[I-D.ietf-tls-tls13](#)] for constrained devices can be challenging. However, recent improvements to the design and implementation of cryptographic algorithms have made TLS accessible to some highly limited devices (see for example [[RFC7925](#)]).

Receiving large protected records can be particularly difficult for a device with limited operating memory. TLS versions 1.2 and earlier [[RFC5246](#)] permit senders to generate records 16384 octets in size, plus any expansion from compression and protection up to 2048 octets (though typically this expansion is only 16 octets). TLS 1.3 reduces the allowance for expansion to 256 octets. Allocating up to 18K of memory for ciphertext is beyond the capacity of some implementations.

An Authentication Encryption with Additional Data (AEAD) ciphers (see [[RFC5116](#)]) API requires that an entire record be present to decrypt and authenticate it. Similarly, other ciphers cannot produce authenticated data until the entire record is present. Thus, incremental processing of records minimally exposes endpoints to the risk of forged data.

The "max_fragment_length" extension [[RFC6066](#)] was designed to enable

constrained clients to negotiate a lower record size. However, "max_fragment_length" suffers from several design problems (see [Section 3](#)).

This document defines a "record_size_limit" extension ([Section 4](#)). This extension replaces "max_fragment_length" [[RFC6066](#)], which this document deprecates. This extension is valid in all versions of TLS.

[2](#). Conventions and Definitions

The words "MUST", "MUST NOT", "SHOULD", and "MAY" are used in this document. It's not shouting; when they are capitalized, they have the special meaning defined in [[RFC2119](#)].

[3](#). Limitations of the "max_fragment_length" Extension

The "max_fragment_length" extension has several limitations that make it unsuitable for use.

A client that has no constraints preventing it from accepting a large record cannot use "max_fragment_length" without risking a reduction in the size of records. The maximum value that the extension permits is 2^{12} , much smaller than the maximum record size of 2^{14} that the protocol permits.

For large data transfers, small record sizes can materially affect performance. Every record incurs additional costs, both in the additional octets for record headers and for expansion due to encryption. Processing more records also adds computational overheads that can be amortized more effectively for larger record sizes. Consequently, clients that are capable of receiving large records could be unwilling to risk reducing performance by offering the extension, especially if the extension is rarely needed.

This would not be an issue if a codepoint were available or could be added for fragments of 2^{14} octets. However, [RFC 6066](#) requires that servers abort the handshake with an "illegal_parameter" alert if they receive the extension with a value they don't understand. This makes it impossible to add new values to the extension without risking connection attempts failing.

A server that negotiates "max_fragment_length" is required to echo the value selected by the client. The server cannot request a lower limit than the one the client offered. This is a significant problem if a server is more constrained than the clients it serves.

The "max_fragment_length" extension is also ill-suited to cases where the capabilities of client and server are asymmetric. Constraints on record size are often receiver constraints.

In comparison, an implementation might be able to send data incrementally. Encryption does not have the same atomicity

requirement. Some ciphers can be encrypted and sent progressively. Thus, an endpoint might be willing to send more than its receive limit.

If these disincentives are sufficient to discourage clients from deploying the "max_fragment_length" extension, then constrained servers are unable to limit record sizes.

[4.](#) The "record_size_limit" Extension

The ExtensionData of the "record_size_limit" extension is RecordSizeLimit:

```
uint16 RecordSizeLimit;
```

The value of RecordSizeLimit is the maximum size of record in octets that the endpoint is willing to receive. This value is used to limit the size of records that are created when encoding application data and handshake message into records.

When the "record_size_limit" extension is negotiated, an endpoint MUST NOT generate a protected record with plaintext that is larger than the RecordSizeLimit value it receives from its peer. Unprotected messages - handshake messages in particular - are not subject to this limit.

This value is the length of the plaintext of a protected record. The value includes the content type and padding added in TLS 1.3 (that is, the complete length of TLSInnerPlaintext). In TLS 1.2 and

earlier, the limit covers all input to compression and encryption, that is the data that ultimately produces TLSCiphertext.fragment. Padding added as part of encryption, such as that added by a block cipher, is not included in this count (see [Section 4.1](#)).

An endpoint that supports all record sizes can include any limit up to the protocol-defined limit for maximum record size. For TLS 1.3 and earlier, that limit is 2^{14} octets. Higher values are currently reserved for future versions of the protocol that may allow larger records; an endpoint MUST NOT send a value higher than the protocol-defined maximum record size unless explicitly allowed by such a future version or extension.

Even if a larger record size limit is provided by a peer, an endpoint MUST NOT send records larger than the protocol-defined limit, unless explicitly allowed by a future TLS version or extension.

The record size limit only applies to records sent toward the endpoint that advertises the limit. An endpoint can send records

that are larger than the limit it advertises as its own limit. An endpoint that receives a record larger than its advertised limit MUST generate a fatal "record_overflow" alert.

Clients SHOULD advertise the "record_size_limit" extension, even if they have no need to limit the size of records. This allows servers to apply a limit at their discretion. If this extension is not negotiated, endpoints can send records of any size permitted by the protocol or other negotiated extensions.

Endpoints MUST NOT send a "record_size_limit" extension with a value smaller than 64. An endpoint MUST treat receipt of a smaller value as a fatal error and generate an "illegal_parameter" alert.

In TLS 1.3, the server sends the "record_size_limit" extension in the EncryptedExtensions message.

During renegotiation, the record size limit is renegotiated. Records are subject to the limits that were set in the handshake that produces the keys that are used to protect those records. This admits the possibility that the extension might not be negotiated when a connection is renegotiated.

[4.1.](#) Record Expansion Limits

The size limit expressed in the "record_size_limit" extension doesn't account for expansion due to compression or record protection. It is expected that a constrained device will disable compression to avoid unpredictable increases in record size. Stream ciphers and existing AEAD ciphers don't permit variable amounts of expansion, but block ciphers do permit variable expansion.

In TLS 1.2, block ciphers allow between 1 and 256 octets of padding. When a limit lower than the protocol-defined limit is advertised, a second limit applies to the length of records that use block ciphers. An endpoint **MUST NOT** add padding to records that would cause the protected record to exceed the size of a protected record that contains the maximum amount of plaintext and the minimum permitted amount of padding.

For example, TLS_RSA_WITH_AES_128_CBC_SHA has 16 octet blocks and a 20 octet MAC. Given a record size limit of 256, a record of that length would require a minimum of 11 octets of padding (for [RFC5246](#) where the MAC is covered by encryption); or 15 octets if the "encrypt_then_mac" extension [RFC7366](#) is negotiated. With this limit, a record with 250 octets of plaintext could be padded to the same length by including at most 17 octets of padding; or 21 octets with "encrypt_then_mac".

An implementation that always adds the minimum amount of padding will always comply with this requirement.

[5.](#) Deprecating "max_fragment_length"

The "record_size_limit" extension replaces the "max_fragment_length" extension [RFC6066](#). A server that supports the "record_size_limit" extension **MUST** ignore and "max_fragment_length" that appears in a ClientHello if both extensions appear. A client **MUST** treat receipt of both "max_fragment_length" and "record_size_limit" as a fatal error, and **SHOULD** generate an "illegal_parameter" alert.

Clients that depend on having a small record size **MAY** continue to advertise the "max_fragment_length".

6. Security Considerations

Very small record sizes might generate additional work for senders and receivers, limiting throughput and increasing exposure to denial of service.

7. IANA Considerations

This document registers the "record_size_limit" extension in the TLS "ExtensionType Values" registry established in [RFC5246]. The "record_size_limit" extension has been assigned a code point of TBD; it is recommended and marked as "Encrypted" in TLS 1.3.

8. References

8.1. Normative References

[I-D.ietf-tls-tls13]

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8.2. Informative References

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Appendix A. Acknowledgments

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