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The Transport Layer Security (TLS) Protocol Version 1.3 **draft-ietf-tls-tls13-14**

Abstract

This document specifies version 1.3 of the Transport Layer Security (TLS) protocol. TLS allows client/server applications to communicate over the Internet in a way that is designed to prevent eavesdropping, tampering, and message forgery.

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[1](#). Introduction

DISCLAIMER: This is a WIP draft of TLS 1.3 and has not yet seen significant security analysis.

RFC EDITOR: PLEASE REMOVE THE FOLLOWING PARAGRAPH The source for this draft is maintained in GitHub. Suggested changes should be submitted as pull requests at <https://github.com/tlswg/tls13-spec>.

Instructions are on that page as well. Editorial changes can be managed in GitHub, but any substantive change should be discussed on the TLS mailing list.

The primary goal of TLS is to provide a secure channel between two communicating peers. Specifically, the channel should provide the following properties.

- Authentication: The server side of the channel is always authenticated; the client side is optionally authenticated. Authentication can happen via asymmetric cryptography (e.g., RSA [[RSA](#)], ECDSA [[ECDSA](#)]) or a pre-shared symmetric key.
- Confidentiality: Data sent over the channel is not visible to attackers.
- Integrity: Data sent over the channel cannot be modified by attackers.

These properties should be true even in the face of an attacker who has complete control of the network, as described in [[RFC3552](#)]. See [Appendix D](#) for a more complete statement of the relevant security properties.

TLS consists of two primary components:

- A handshake protocol ([Section 4](#)) which authenticates the communicating parties, negotiates cryptographic modes and

parameters, and establishes shared keying material. The handshake protocol is designed to resist tampering; an active attacker should not be able to force the peers to negotiate different parameters than they would if the connection were not under attack.

- A record protocol ([Section 5](#)) which uses the parameters established by the handshake protocol to protect traffic between the communicating peers. The record protocol divides traffic up into a series of records, each of which is independently protected using the traffic keys.

TLS is application protocol independent; higher-level protocols can layer on top of TLS transparently. The TLS standard, however, does not specify how protocols add security with TLS; the decisions on how to initiate TLS handshaking and how to interpret the authentication certificates exchanged are left to the judgment of the designers and implementors of protocols that run on top of TLS.

This document defines TLS version 1.3. While TLS 1.3 is not directly compatible with previous versions, all versions of TLS incorporate a versioning mechanism which allows clients and servers to interoperably negotiate a common version if one is supported.

1.1. Conventions and Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC 2119](#) [[RFC2119](#)].

The following terms are used:

client: The endpoint initiating the TLS connection.

connection: A transport-layer connection between two endpoints.

endpoint: Either the client or server of the connection.

handshake: An initial negotiation between client and server that establishes the parameters of their transactions.

peer: An endpoint. When discussing a particular endpoint, "peer" refers to the endpoint that is remote to the primary subject of discussion.

receiver: An endpoint that is receiving records.

sender: An endpoint that is transmitting records.

session: An association between a client and a server resulting from a handshake.

server: The endpoint which did not initiate the TLS connection.

1.2. Major Differences from TLS 1.2

[draft-14](#)

- Allow cookies to be longer (*).
- Remove the "context" from EarlyDataIndication as it was undefined and nobody used it (*).
- Remove 0-RTT EncryptedExtensions and replace the ticket_age extension with an obfuscated version. Also necessitates a change to NewSessionTicket (*).
- Move the downgrade sentinel to the end of ServerHello.Random to accomodate tlsv1_2 (*).
- Define ecdsa_sha1 (*).
- Allow resumption even after fatal alerts. This matches current practice.
- Remove non-closure warning alerts. Require treating unknown alerts as fatal.
- Make the rules for accepting 0-RTT less restrictive.
- Clarify 0-RTT backward-compatibility rules.
- Clarify how 0-RTT and PSK identities interact.
- Add a section describing the data limits for each cipher.
- Major editorial restructuring.
- Replace the Security Analysis section with a WIP draft.

(*) indicates changes to the wire protocol which may require implementations to update.

[draft-13](#)

- Allow server to send SupportedGroups.
- Remove 0-RTT client authentication
- Remove (EC)DHE 0-RTT.
- Flesh out 0-RTT PSK mode and shrink EarlyDataIndication
- Turn PSK-resumption response into an index to save room
- Move CertificateStatus to an extension
- Extra fields in NewSessionTicket.
- Restructure key schedule and add a resumption_context value.
- Require DH public keys and secrets to be zero-padded to the size of the group.
- Remove the redundant length fields in KeyShareEntry.
- Define a cookie field for HRR.

[draft-12](#)

- Provide a list of the PSK cipher suites.
- Remove the ability for the ServerHello to have no extensions (this aligns the syntax with the text).
- Clarify that the server can send application data after its first flight (0.5 RTT data)
- Revise signature algorithm negotiation to group hash, signature algorithm, and curve together. This is backwards compatible.
- Make ticket lifetime mandatory and limit it to a week.
- Make the purpose strings lower-case. This matches how people are implementing for interop.
- Define exporters.
- Editorial cleanup

[draft-11](#)

- Port the CFRG curves & signatures work from RFC4492bis.

- Remove sequence number and version from `additional_data`, which is now empty.
- Reorder values in `HkdfLabel`.
- Add support for version anti-downgrade mechanism.
- Update IANA considerations section and relax some of the policies.
- Unify authentication modes. Add post-handshake client authentication.
- Remove `early_handshake` content type. Terminate 0-RTT data with an alert.
- Reset sequence number upon key change (as proposed by Fournet et al.)

[draft-10](#)

- Remove `ClientCertificateTypes` field from `CertificateRequest` and add extensions.
- Merge client and server key shares into a single extension.

[draft-09](#)

- Change to RSA-PSS signatures for handshake messages.
- Remove support for DSA.
- Update key schedule per suggestions by Hugo, Hoeteck, and Bjoern Tackmann.
- Add support for per-record padding.
- Switch to encrypted record `ContentType`.
- Change HKDF labeling to include protocol version and value lengths.
- Shift the final decision to abort a handshake due to incompatible certificates to the client rather than having servers abort early.
- Deprecate SHA-1 with signatures.
- Add MTI algorithms.

[draft-08](#)

- Remove support for weak and lesser used named curves.
- Remove support for MD5 and SHA-224 hashes with signatures.
- Update lists of available AEAD cipher suites and error alerts.
- Reduce maximum permitted record expansion for AEAD from 2048 to 256 octets.
- Require digital signatures even when a previous configuration is used.
- Merge EarlyDataIndication and KnownConfiguration.
- Change code point for server_configuration to avoid collision with server_hello_done.
- Relax certificate_list ordering requirement to match current practice.

[draft-07](#)

- Integration of semi-ephemeral DH proposal.
- Add initial 0-RTT support.
- Remove resumption and replace with PSK + tickets.
- Move ClientKeyShare into an extension.
- Move to HKDF.

[draft-06](#)

- Prohibit RC4 negotiation for backwards compatibility.
- Freeze & deprecate record layer version field.
- Update format of signatures with context.
- Remove explicit IV.

[draft-05](#)

- Prohibit SSL negotiation for backwards compatibility.

- Fix which MS is used for exporters.

[draft-04](#)

- Modify key computations to include session hash.
- Remove ChangeCipherSpec.
- Renumber the new handshake messages to be somewhat more consistent with existing convention and to remove a duplicate registration.
- Remove renegotiation.
- Remove point format negotiation.

[draft-03](#)

- Remove GMT time.
- Merge in support for ECC from [RFC 4492](#) but without explicit curves.
- Remove the unnecessary length field from the AD input to AEAD ciphers.
- Rename {Client,Server}KeyExchange to {Client,Server}KeyShare.
- Add an explicit HelloRetryRequest to reject the client's.

[draft-02](#)

- Increment version number.
- Rework handshake to provide 1-RTT mode.
- Remove custom DHE groups.
- Remove support for compression.
- Remove support for static RSA and DH key exchange.
- Remove support for non-AEAD ciphers.

[2.](#) Protocol Overview

The cryptographic parameters of the session state are produced by the TLS handshake protocol. When a TLS client and server first start communicating, they agree on a protocol version, select cryptographic

algorithms, optionally authenticate each other, and establish shared secret keying material. Once the handshake is complete, the peers use the established keys to protect application layer traffic.

TLS supports three basic key exchange modes:

- Diffie-Hellman (of both the finite field and elliptic curve varieties).
- A pre-shared symmetric key (PSK)
- A combination of a symmetric key and Diffie-Hellman

Which mode is used depends on the negotiated cipher suite. Conceptually, the handshake establishes three secrets which are used to derive all the keys.

Figure 1 below shows the basic full TLS handshake:

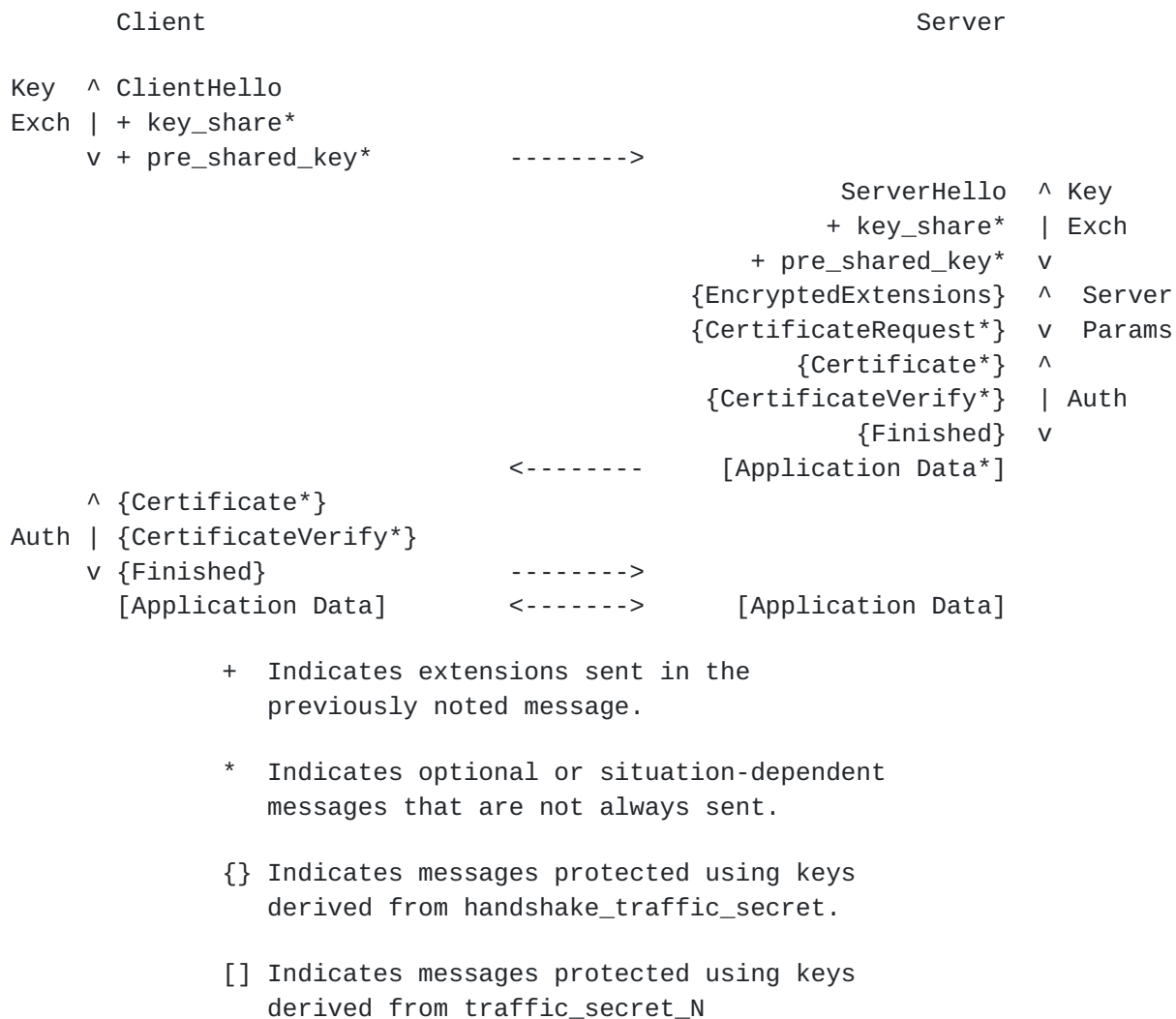


Figure 1: Message flow for full TLS Handshake

The handshake can be thought of as having three phases, indicated in the diagram above.

- Key Exchange: Establish shared keying material and select the cryptographic parameters. Everything after this phase is encrypted.
- Server Parameters: Establish other handshake parameters. (whether the client is authenticated, application layer protocol support, etc.)
- Authentication: Authenticate the server (and optionally the client) and provide key confirmation and handshake integrity.

In the Key Exchange phase, the client sends the ClientHello ([Section 4.1.1](#)) message, which contains a random nonce (ClientHello.random), its offered protocol version, cipher suite, and extensions, and in general either one or more Diffie-Hellman key shares (in the "key_share" extension [Section 4.2.4](#)), one or more pre-shared key labels (in the "pre_shared_key" extension [Section 4.2.5](#)), or both.

The server processes the ClientHello and determines the appropriate cryptographic parameters for the connection. It then responds with its own ServerHello which indicates the negotiated connection parameters. [[Section 4.1.2](#)]. The combination of the ClientHello and the ServerHello determines the shared keys. If either a pure (EC)DHE or (EC)DHE-PSK cipher suite is in use, then the ServerHello will contain a "key_share" extension with the server's ephemeral Diffie-Hellman share which MUST be in the same group as one of the client's shares. If a pure PSK or an (EC)DHE-PSK cipher suite is negotiated, then the ServerHello will contain a "pre_shared_key" extension indicating which of the client's offered PSKs was selected.

The server then sends two messages to establish the Server Parameters:

EncryptedExtensions. responses to any extensions which are not required in order to determine the cryptographic parameters. [[Section 4.2.8](#)]

CertificateRequest. if certificate-based client authentication is desired, the desired parameters for that certificate. This message will be omitted if client authentication is not desired.

Finally, the client and server exchange Authentication messages. TLS uses the same set of messages every time that authentication is needed. Specifically:

Certificate. the certificate of the endpoint. This message is omitted if the server is not authenticating with a certificate (i.e., with PSK or (EC)DHE-PSK cipher suites). Note that if raw public keys [[RFC7250](#)] or the cached information extension [[I-D.ietf-tls-cached-info](#)] are in use, then this message will not contain a certificate but rather some other value corresponding to the server's long-term key. [[Section 4.3.1](#)]

CertificateVerify. a signature over the entire handshake using the public key in the Certificate message. This message is omitted if the server is not authenticating via a certificate (i.e., with PSK or (EC)DHE-PSK cipher suites). [[Section 4.3.2](#)]

Finished. a MAC (Message Authentication Code) over the entire handshake. This message provides key confirmation, binds the endpoint's identity to the exchanged keys, and in PSK mode also authenticates the handshake. [[Section 4.3.3](#)]

Upon receiving the server's messages, the client responds with its Authentication messages, namely Certificate and CertificateVerify (if requested), and Finished.

At this point, the handshake is complete, and the client and server may exchange application layer data. Application data **MUST NOT** be sent prior to sending the Finished message. Note that while the server may send application data prior to receiving the client's Authentication messages, any data sent at that point is, of course, being sent to an unauthenticated peer.

[2.1.](#) Incorrect DHE Share

If the client has not provided a sufficient "key_share" extension (e.g. it includes only DHE or ECDHE groups unacceptable or unsupported by the server), the server corrects the mismatch with a HelloRetryRequest and the client will need to restart the handshake with an appropriate "key_share" extension, as shown in Figure 2. If no common cryptographic parameters can be negotiated, the server will send a "handshake_failure" or "insufficient_security" fatal alert (see [Section 6](#)).



Figure 2: Message flow for a full handshake with mismatched parameters

Note: The handshake transcript includes the initial ClientHello/HelloRetryRequest exchange; it is not reset with the new ClientHello.

TLS also allows several optimized variants of the basic handshake, as described in the following sections.

2.2. Resumption and Pre-Shared Key (PSK)

Although TLS PSKs can be established out of band, PSKs can also be established in a previous session and then reused ("session resumption"). Once a handshake has completed, the server can send the client a PSK identity which corresponds to a key derived from the initial handshake (See [Section 4.4.1](#)). The client can then use that PSK identity in future handshakes to negotiate use of the PSK. If the server accepts it, then the security context of the new connection is tied to the original connection. In TLS 1.2 and below, this functionality was provided by "session IDs" and "session tickets" [[RFC5077](#)]. Both mechanisms are obsoleted in TLS 1.3.

PSK cipher suites can either use PSK in combination with an (EC)DHE exchange in order to provide forward secrecy in combination with shared keys, or can use PSKs alone, at the cost of losing forward secrecy.

Figure 3 shows a pair of handshakes in which the first establishes a PSK and the second uses it:

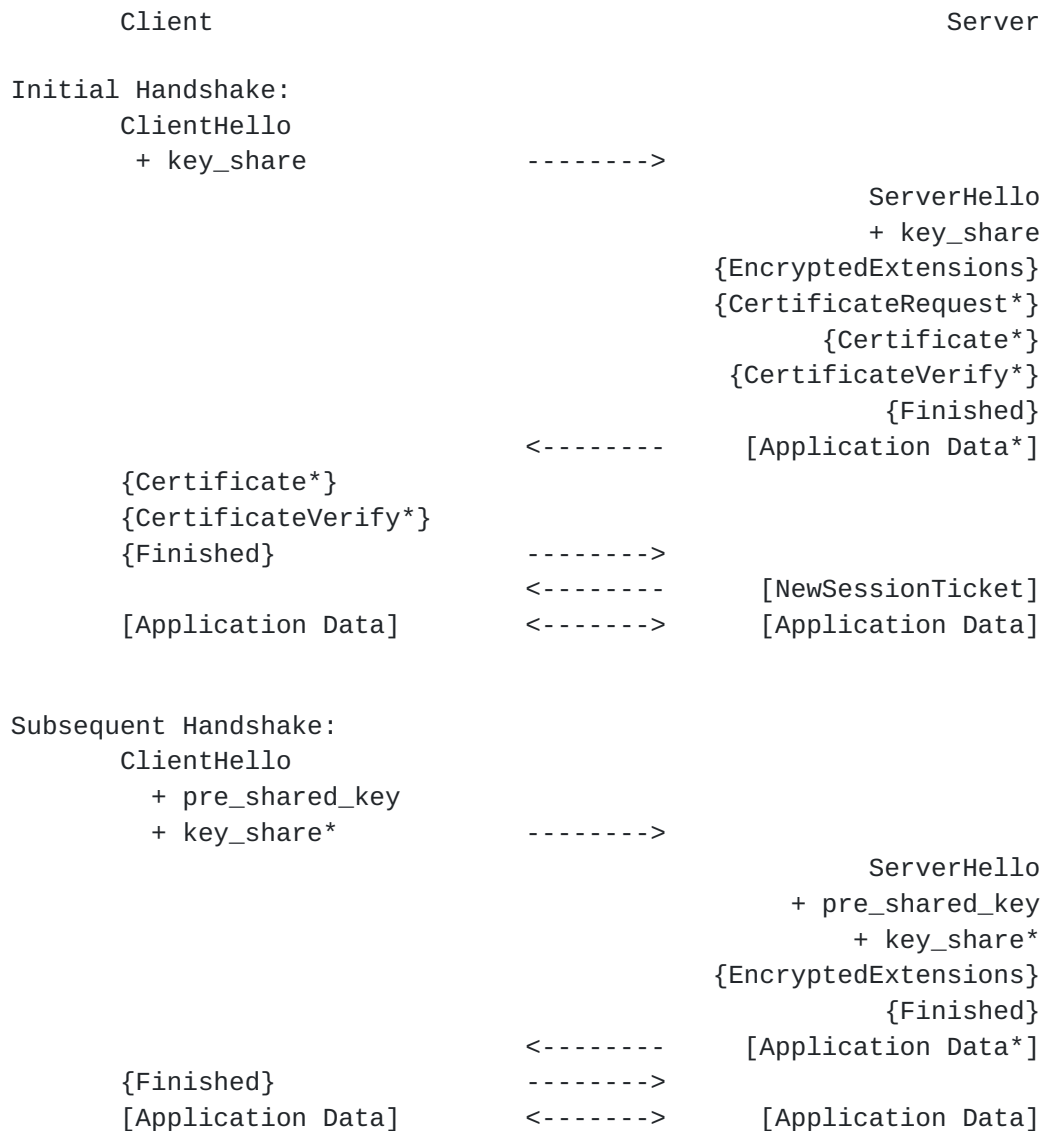


Figure 3: Message flow for resumption and PSK

As the server is authenticating via a PSK, it does not send a Certificate or a CertificateVerify. When a client offers resumption via PSK it SHOULD also supply a "key_share" extension to the server as well to allow the server to decline resumption and fall back to a full handshake, if needed. A "key_share" extension MUST also be sent if the client is attempting to negotiate an (EC)DHE-PSK cipher suite.

2.3. Zero-RTT Data

When resuming via a PSK with an appropriate ticket (i.e., one with the "allow_early_data" flag), clients can also send data on their first flight ("early data"). This data is encrypted solely under keys derived using the first offered PSK as the static secret. As shown in Figure 4, the Zero-RTT data is just added to the 1-RTT handshake in the first flight. The rest of the handshake uses the same messages.

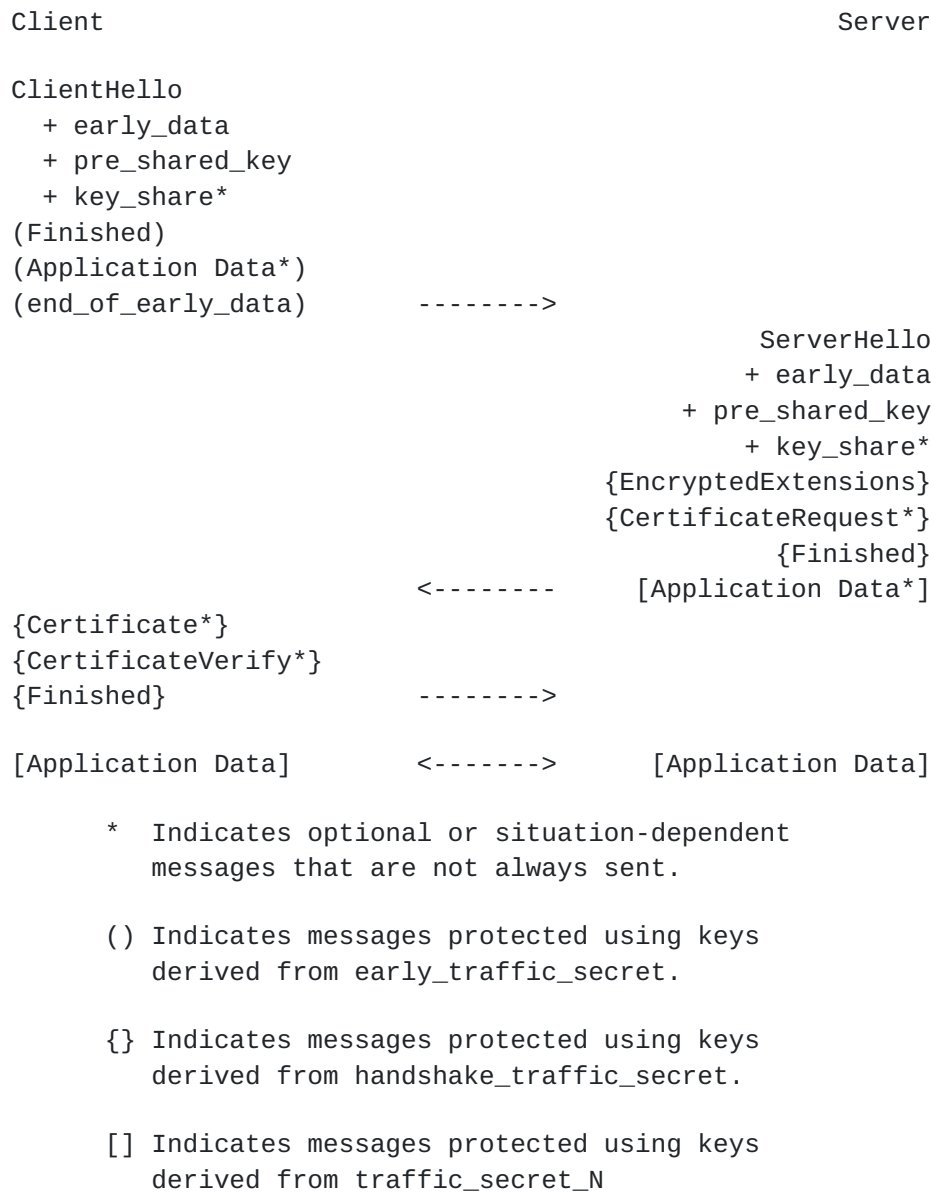


Figure 4: Message flow for a zero round trip handshake

[[OPEN ISSUE: Should it be possible to combine 0-RTT with the server authenticating via a signature <https://github.com/tlswg/tls13-spec/issues/443>]]

IMPORTANT NOTE: The security properties for 0-RTT data (regardless of the cipher suite) are weaker than those for other kinds of TLS data. Specifically:

1. This data is not forward secret, because it is encrypted solely with the PSK.
2. There are no guarantees of non-replay between connections. Unless the server takes special measures outside those provided by TLS, the server has no guarantee that the same 0-RTT data was not transmitted on multiple 0-RTT connections (See [Section 4.2.6.2](#) for more details). This is especially relevant if the data is authenticated either with TLS client authentication or inside the application layer protocol. However, 0-RTT data cannot be duplicated within a connection (i.e., the server will not process the same data twice for the same connection) and an attacker will not be able to make 0-RTT data appear to be 1-RTT data (because it is protected with different keys.)

The remainder of this document provides a detailed description of TLS.

3. Presentation Language

This document deals with the formatting of data in an external representation. The following very basic and somewhat casually defined presentation syntax will be used. The syntax draws from several sources in its structure. Although it resembles the programming language "C" in its syntax and XDR [[RFC4506](#)] in both its syntax and intent, it would be risky to draw too many parallels. The purpose of this presentation language is to document TLS only; it has no general application beyond that particular goal.

3.1. Basic Block Size

The representation of all data items is explicitly specified. The basic data block size is one byte (i.e., 8 bits). Multiple byte data items are concatenations of bytes, from left to right, from top to bottom. From the byte stream, a multi-byte item (a numeric in the example) is formed (using C notation) by:

```
value = (byte[0] << 8*(n-1)) | (byte[1] << 8*(n-2)) |  
        ... | byte[n-1];
```


This byte ordering for multi-byte values is the commonplace network byte order or big-endian format.

3.2. Miscellaneous

Comments begin with `/*` and end with `*/`.

Optional components are denoted by enclosing them in `"[]"` double brackets.

Single-byte entities containing uninterpreted data are of type opaque.

3.3. Vectors

A vector (single-dimensioned array) is a stream of homogeneous data elements. The size of the vector may be specified at documentation time or left unspecified until runtime. In either case, the length declares the number of bytes, not the number of elements, in the vector. The syntax for specifying a new type, `T'`, that is a fixed-length vector of type `T` is

```
T T'[n];
```

Here, `T'` occupies `n` bytes in the data stream, where `n` is a multiple of the size of `T`. The length of the vector is not included in the encoded stream.

In the following example, `Datum` is defined to be three consecutive bytes that the protocol does not interpret, while `Data` is three consecutive `Datum`, consuming a total of nine bytes.

```
opaque Datum[3];      /* three uninterpreted bytes */
Datum Data[9];         /* 3 consecutive 3 byte vectors */
```

Variable-length vectors are defined by specifying a subrange of legal lengths, inclusively, using the notation `<floor..ceiling>`. When these are encoded, the actual length precedes the vector's contents in the byte stream. The length will be in the form of a number consuming as many bytes as required to hold the vector's specified maximum (ceiling) length. A variable-length vector with an actual length field of zero is referred to as an empty vector.

```
T T'<floor..ceiling>;
```

In the following example, `mandatory` is a vector that must contain between 300 and 400 bytes of type opaque. It can never be empty. The actual length field consumes two bytes, a `uint16`, which is

sufficient to represent the value 400 (see [Section 3.4](#)). On the other hand, longer can represent up to 800 bytes of data, or 400 uint16 elements, and it may be empty. Its encoding will include a two-byte actual length field prepended to the vector. The length of an encoded vector must be an even multiple of the length of a single element (for example, a 17-byte vector of uint16 would be illegal).

```
opaque mandatory<300..400>;  
    /* length field is 2 bytes, cannot be empty */  
uint16 longer<0..800>;  
    /* zero to 400 16-bit unsigned integers */
```

[3.4.](#) Numbers

The basic numeric data type is an unsigned byte (uint8). All larger numeric data types are formed from fixed-length series of bytes concatenated as described in [Section 3.1](#) and are also unsigned. The following numeric types are predefined.

```
uint8 uint16[2];  
uint8 uint24[3];  
uint8 uint32[4];  
uint8 uint64[8];
```

All values, here and elsewhere in the specification, are stored in network byte (big-endian) order; the uint32 represented by the hex bytes 01 02 03 04 is equivalent to the decimal value 16909060.

Note that in some cases (e.g., DH parameters) it is necessary to represent integers as opaque vectors. In such cases, they are represented as unsigned integers (i.e., additional leading zero octets are not used even if the most significant bit is set).

[3.5.](#) Enumerateds

An additional sparse data type is available called enum. A field of type enum can only assume the values declared in the definition. Each definition is a different type. Only enumerateds of the same type may be assigned or compared. Every element of an enumerated must be assigned a value, as demonstrated in the following example. Since the elements of the enumerated are not ordered, they can be assigned any unique value, in any order.

```
enum { e1(v1), e2(v2), ... , en(vn) [[, (n)]] } Te;
```

An enumerated occupies as much space in the byte stream as would its maximal defined ordinal value. The following definition would cause one byte to be used to carry fields of type Color.


```
enum { red(3), blue(5), white(7) } Color;
```

One may optionally specify a value without its associated tag to force the width definition without defining a superfluous element.

In the following example, Taste will consume two bytes in the data stream but can only assume the values 1, 2, or 4.

```
enum { sweet(1), sour(2), bitter(4), (32000) } Taste;
```

The names of the elements of an enumeration are scoped within the defined type. In the first example, a fully qualified reference to the second element of the enumeration would be `Color.blue`. Such qualification is not required if the target of the assignment is well specified.

```
Color color = Color.blue;    /* overspecified, legal */  
Color color = blue;          /* correct, type implicit */
```

For enumerations that are never converted to external representation, the numerical information may be omitted.

```
enum { low, medium, high } Amount;
```

3.6. Constructed Types

Structure types may be constructed from primitive types for convenience. Each specification declares a new, unique type. The syntax for definition is much like that of C.

```
struct {  
    T1 f1;  
    T2 f2;  
    ...  
    Tn fn;  
} [[T]];
```

The fields within a structure may be qualified using the type's name, with a syntax much like that available for enumerations. For example, `T.f2` refers to the second field of the previous declaration. Structure definitions may be embedded.

3.6.1. Variants

Defined structures may have variants based on some knowledge that is available within the environment. The selector must be an enumerated type that defines the possible variants the structure defines. There must be a case arm for every element of the enumeration declared in

the select. Case arms have limited fall-through: if two case arms follow in immediate succession with no fields in between, then they both contain the same fields. Thus, in the example below, "orange" and "banana" both contain V2. Note that this is a new piece of syntax in TLS 1.2.

The body of the variant structure may be given a label for reference. The mechanism by which the variant is selected at runtime is not prescribed by the presentation language.

```
struct {  
    T1 f1;  
    T2 f2;  
    ....  
    Tn fn;  
    select (E) {  
        case e1: Te1;  
        case e2: Te2;  
        case e3: case e4: Te3;  
        ....  
        case en: Ten;  
    } [[fv]];  
} [[Tv]];
```

For example:

```
enum { apple, orange, banana } VariantTag;  
  
struct {  
    uint16 number;  
    opaque string<0..10>; /* variable length */  
} V1;  
  
struct {  
    uint32 number;  
    opaque string[10]; /* fixed length */  
} V2;  
  
struct {  
    select (VariantTag) { /* value of selector is implicit */  
        case apple:  
            V1; /* VariantBody, tag = apple */  
        case orange:  
        case banana:  
            V2; /* VariantBody, tag = orange or banana */  
    } variant_body; /* optional label on variant */  
} VariantRecord;
```


3.7. Constants

Typed constants can be defined for purposes of specification by declaring a symbol of the desired type and assigning values to it.

Under-specified types (opaque, variable-length vectors, and structures that contain opaque) cannot be assigned values. No fields of a multi-element structure or vector may be elided.

For example:

```
struct {  
    uint8 f1;  
    uint8 f2;  
} Example1;  
  
Example1 ex1 = {1, 4}; /* assigns f1 = 1, f2 = 4 */
```

4. Handshake Protocol

The handshake protocol is used to negotiate the secure attributes of a session. Handshake messages are supplied to the TLS record layer, where they are encapsulated within one or more TLSPlaintext or TLSCiphertext structures, which are processed and transmitted as specified by the current active session state.


```
enum {
    client_hello(1),
    server_hello(2),
    new_session_ticket(4),
    hello_retry_request(6),
    encrypted_extensions(8),
    certificate(11),
    certificate_request(13),
    certificate_verify(15),
    finished(20),
    key_update(24),
    (255)
} HandshakeType;

struct {
    HandshakeType msg_type;      /* handshake type */
    uint24 length;              /* bytes in message */
    select (HandshakeType) {
        case client_hello:      ClientHello;
        case server_hello:      ServerHello;
        case hello_retry_request: HelloRetryRequest;
        case encrypted_extensions: EncryptedExtensions;
        case certificate_request: CertificateRequest;
        case certificate:        Certificate;
        case certificate_verify: CertificateVerify;
        case finished:           Finished;
        case new_session_ticket: NewSessionTicket;
        case key_update:         KeyUpdate;
    } body;
} Handshake;
```

Protocol messages MUST be sent in the order defined below (and shown in the diagrams in [Section 2](#)). Sending handshake messages in an unexpected order results in an "unexpected_message" fatal error. Unneeded handshake messages are omitted, however.

New handshake message types are assigned by IANA as described in [Section 10](#).

[4.1](#). Key Exchange Messages

The key exchange messages are used to exchange security capabilities between the client and server and to establish the traffic keys used to protect the handshake and data.

[4.1.1.1.](#) Client Hello

When this message will be sent:

When a client first connects to a server, it is required to send the ClientHello as its first message. The client will also send a ClientHello when the server has responded to its ClientHello with a ServerHello that selects cryptographic parameters that don't match the client's "key_share" extension. In that case, the client MUST send the same ClientHello (without modification) except:

- Including a new KeyShareEntry as the lowest priority share (i.e., appended to the list of shares in the "key_share" extension).
- Removing the EarlyDataIndication [Section 4.2.6](#) extension if one was present. Early data is not permitted after HelloRetryRequest.

If a server receives a ClientHello at any other time, it MUST send a fatal "unexpected_message" alert and close the connection.

Structure of this message:

```
struct {
    uint8 major;
    uint8 minor;
} ProtocolVersion;

struct {
    opaque random_bytes[32];
} Random;

uint8 CipherSuite[2];    /* Cryptographic suite selector */

struct {
    ProtocolVersion client_version = { 3, 4 };    /* TLS v1.3 */
    Random random;
    opaque legacy_session_id<0..32>;
    CipherSuite cipher_suites<2..2^16-2>;
    opaque legacy_compression_methods<1..2^8-1>;
    Extension extensions<0..2^16-1>;
} ClientHello;
```

TLS allows extensions to follow the compression_methods field in an extensions block. The presence of extensions can be detected by determining whether there are bytes following the compression_methods at the end of the ClientHello. Note that this method of detecting optional data differs from the normal TLS method of having a

variable-length field, but it is used for compatibility with TLS before extensions were defined. As of TLS 1.3, all clients and servers will send at least one extension (at least "key_share" or "pre_shared_key").

client_version The latest (highest valued) version of the TLS protocol offered by the client. This SHOULD be the same as the latest version supported. For this version of the specification, the version will be { 3, 4 }. (See [Appendix C](#) for details about backward compatibility.)

random 32 bytes generated by a secure random number generator. See [Appendix B](#) for additional information.

legacy_session_id Versions of TLS before TLS 1.3 supported a session resumption feature which has been merged with Pre-Shared Keys in this version (see [Section 2.2](#)). This field MUST be ignored by a server negotiating TLS 1.3 and SHOULD be set as a zero length vector (i.e., a single zero byte length field) by clients which do not have a cached session ID set by a pre-TLS 1.3 server.

cipher_suites This is a list of the cryptographic options supported by the client, with the client's first preference first. Each cipher suite defines a key exchange algorithm, a record protection algorithm (including secret key length) and a hash to be used with HKDF. The server will select a cipher suite or, if no acceptable choices are presented, return a "handshake_failure" alert and close the connection. If the list contains cipher suites the server does not recognize, support, or wish to use, the server MUST ignore those cipher suites, and process the remaining ones as usual. Values are defined in [Appendix A.4](#).

legacy_compression_methods Versions of TLS before 1.3 supported compression with the list of supported compression methods being sent in this field. For every TLS 1.3 ClientHello, this vector MUST contain exactly one byte set to zero, which corresponds to the "null" compression method in prior versions of TLS. If a TLS 1.3 ClientHello is received with any other value in this field, the server MUST generate a fatal "illegal_parameter" alert. Note that TLS 1.3 servers might receive TLS 1.2 or prior ClientHellos which contain other compression methods and MUST follow the procedures for the appropriate prior version of TLS.

extensions Clients request extended functionality from servers by sending data in the extensions field. The actual "Extension" format is defined in [Section 4.2](#).

In the event that a client requests additional functionality using extensions, and this functionality is not supplied by the server, the client MAY abort the handshake. Note that TLS 1.3 ClientHello messages MUST always contain extensions, and a TLS 1.3 server MUST respond to any TLS 1.3 ClientHello without extensions with a fatal "decode_error" alert. TLS 1.3 servers may receive TLS 1.2 ClientHello messages without extensions. If negotiating TLS 1.2, a server MUST check that the amount of data in the message precisely matches one of these formats; if not, then it MUST send a fatal "decode_error" alert.

After sending the ClientHello message, the client waits for a ServerHello or HelloRetryRequest message.

4.1.2. Server Hello

When this message will be sent:

The server will send this message in response to a ClientHello message when it was able to find an acceptable set of algorithms and the client's "key_share" extension was acceptable. If the client proposed groups are not acceptable by the server, it will respond with a "handshake_failure" fatal alert.

Structure of this message:

```
struct {  
    ProtocolVersion server_version;  
    Random random;  
    CipherSuite cipher_suite;  
    Extension extensions<0..2^16-1>;  
} ServerHello;
```

server_version This field contains the version of TLS negotiated for this session. Servers MUST select the lower of the highest supported server version and the version offered by the client in the ClientHello. In particular, servers MUST accept ClientHello messages with versions higher than those supported and negotiate the highest mutually supported version. For this version of the specification, the version is { 3, 4 }. (See [Appendix C](#) for details about backward compatibility.)

random This structure is generated by the server and MUST be generated independently of the ClientHello.random.

cipher_suite The single cipher suite selected by the server from the list in ClientHello.cipher_suites. For resumed sessions, this

field is the value from the state of the session being resumed.
[[TODO: interaction with PSK.]]

extensions A list of extensions. Note that only extensions offered by the client can appear in the server's list. In TLS 1.3, as opposed to previous versions of TLS, the server's extensions are split between the ServerHello and the EncryptedExtensions [Section 4.2.8](#) message. The ServerHello MUST only include extensions which are required to establish the cryptographic context. Currently the only such extensions are "key_share", "pre_shared_key", and "early_data". Clients MUST check the ServerHello for the presence of any forbidden extensions and if any are found MUST terminate the handshake with a "illegal_parameter" alert. In prior versions of TLS, the extensions field could be omitted entirely if not needed, similar to ClientHello. As of TLS 1.3, all clients and servers will send at least one extension (at least "key_share" or "pre_shared_key").

TLS 1.3 has a downgrade protection mechanism embedded in the server's random value. TLS 1.3 server implementations which respond to a ClientHello with a client_version indicating TLS 1.2 or below MUST set the last eight bytes of their Random value to the bytes:

```
44 4F 57 4E 47 52 44 01
```

TLS 1.2 server implementations which respond to a ClientHello with a client_version indicating TLS 1.1 or below SHOULD set the last eight bytes of their Random value to the bytes:

```
44 4F 57 4E 47 52 44 00
```

TLS 1.3 clients receiving a TLS 1.2 or below ServerHello MUST check that the last eight octets are not equal to either of these values. TLS 1.2 clients SHOULD also perform this check if the ServerHello indicates TLS 1.1 or below. If a match is found, the client MUST abort the handshake with a fatal "illegal_parameter" alert. This mechanism provides limited protection against downgrade attacks over and above that provided by the Finished exchange: because the ServerKeyExchange includes a signature over both random values, it is not possible for an active attacker to modify the randoms without detection as long as ephemeral ciphers are used. It does not provide downgrade protection when static RSA is used.

Note: This is an update to TLS 1.2 so in practice many TLS 1.2 clients and servers will not behave as specified above.

4.1.3. Hello Retry Request

When this message will be sent:

Servers send this message in response to a ClientHello message if they were able to find an acceptable set of algorithms and groups that are mutually supported, but the client's KeyShare did not contain an acceptable offer. If it cannot find such a match, it will respond with a fatal "handshake_failure" alert.

Structure of this message:

```
struct {  
    ProtocolVersion server_version;  
    CipherSuite cipher_suite;  
    NamedGroup selected_group;  
    Extension extensions<0..2^16-1>;  
} HelloRetryRequest;
```

`selected_group` The mutually supported group the server intends to negotiate and is requesting a retried ClientHello/KeyShare for.

The `server_version`, `cipher_suite`, and `extensions` fields have the same meanings as their corresponding values in the ServerHello. [[NOTE: `cipher_suite` may disappear. <https://github.com/tlswg/tls13-spec/issues/528>]] The server SHOULD send only the extensions necessary for the client to generate a correct ClientHello pair (currently no such extensions exist). As with ServerHello, a HelloRetryRequest MUST NOT contain any extensions that were not first offered by the client in its ClientHello.

Upon receipt of a HelloRetryRequest, the client MUST first verify that the `selected_group` field corresponds to a group which was provided in the "supported_groups" extension in the original ClientHello. It MUST then verify that the `selected_group` field does not correspond to a group which was provided in the "key_share" extension in the original ClientHello. If either of these checks fails, then the client MUST abort the handshake with a fatal "handshake_failure" alert. Clients SHOULD also abort with "handshake_failure" in response to any second HelloRetryRequest which was sent in the same connection (i.e., where the ClientHello was itself in response to a HelloRetryRequest).

Otherwise, the client MUST send a ClientHello with an updated KeyShare extension to the server. The client MUST append a new KeyShareEntry for the group indicated in the `selected_group` field to the groups in its original KeyShare.

Upon re-sending the ClientHello and receiving the server's ServerHello/KeyShare, the client MUST verify that the selected CipherSuite and NamedGroup match that supplied in the HelloRetryRequest. If either of these values differ, the client MUST abort the connection with a fatal "handshake_failure" alert.

4.2. Hello Extensions

The extension format is:

```
struct {
    ExtensionType extension_type;
    opaque extension_data<0..2^16-1>;
} Extension;

enum {
    supported_groups(10),
    signature_algorithms(13),
    key_share(40),
    pre_shared_key(41),
    early_data(42),
    cookie(44),
    (65535)
} ExtensionType;
```

Here:

- "extension_type" identifies the particular extension type.
- "extension_data" contains information specific to the particular extension type.

The initial set of extensions is defined in [[RFC6066](#)]. The list of extension types is maintained by IANA as described in [Section 10](#).

An extension type MUST NOT appear in the ServerHello or HelloRetryRequest unless the same extension type appeared in the corresponding ClientHello. If a client receives an extension type in ServerHello or HelloRetryRequest that it did not request in the associated ClientHello, it MUST abort the handshake with an "unsupported_extension" fatal alert.

Nonetheless, "server-oriented" extensions may be provided within this framework. Such an extension (say, of type x) would require the client to first send an extension of type x in a ClientHello with empty extension_data to indicate that it supports the extension type. In this case, the client is offering the capability to understand the extension type, and the server is taking the client up on its offer.

When multiple extensions of different types are present in the ClientHello or ServerHello messages, the extensions MAY appear in any order. There MUST NOT be more than one extension of the same type.

Finally, note that extensions can be sent both when starting a new session and when in resumption-PSK mode. A client that requests session resumption does not in general know whether the server will accept this request, and therefore it SHOULD send the same extensions as it would send normally.

In general, the specification of each extension type needs to describe the effect of the extension both during full handshake and session resumption. Most current TLS extensions are relevant only when a session is initiated: when an older session is resumed, the server does not process these extensions in ClientHello, and does not include them in ServerHello. However, some extensions may specify different behavior during session resumption. [[TODO: update this and the previous paragraph to cover PSK-based resumption.]]

There are subtle (and not so subtle) interactions that may occur in this protocol between new features and existing features which may result in a significant reduction in overall security. The following considerations should be taken into account when designing new extensions:

- Some cases where a server does not agree to an extension are error conditions, and some are simply refusals to support particular features. In general, error alerts should be used for the former, and a field in the server extension response for the latter.
- Extensions should, as far as possible, be designed to prevent any attack that forces use (or non-use) of a particular feature by manipulation of handshake messages. This principle should be followed regardless of whether the feature is believed to cause a security problem. Often the fact that the extension fields are included in the inputs to the Finished message hashes will be sufficient, but extreme care is needed when the extension changes the meaning of messages sent in the handshake phase. Designers and implementors should be aware of the fact that until the handshake has been authenticated, active attackers can modify messages and insert, remove, or replace extensions.

4.2.1. Cookie

```
struct {  
    opaque cookie<0..2^16-1>;  
} Cookie;
```


Cookies serve two primary purposes:

- Allowing the server to force the client to demonstrate reachability at their apparent network address (thus providing a measure of DoS protection). This is primarily useful for non-connection-oriented transports (see [[RFC6347](#)] for an example of this).
- Allowing the server to offload state to the client, thus allowing it to send a HelloRetryRequest without storing any state. The server does this by pickling that post-ClientHello hash state into the cookie (protected with some suitable integrity algorithm).

When sending a HelloRetryRequest, the server MAY provide a "cookie" extension to the client (this is an exception to the usual rule that the only extensions that may be sent are those that appear in the ClientHello). When sending the new ClientHello, the client MUST echo the value of the extension. Clients MUST NOT use cookies in subsequent connections.

4.2.2. Signature Algorithms

The client uses the "signature_algorithms" extension to indicate to the server which signature algorithms may be used in digital signatures.

Clients which offer one or more cipher suites which use certificate authentication (i.e., any non-PSK cipher suite) MUST send the "signature_algorithms" extension. If this extension is not provided and no alternative cipher suite is available, the server MUST close the connection with a fatal "missing_extension" alert. (see [Section 8.2](#))

The "extension_data" field of this extension contains a "supported_signature_algorithms" value:


```
enum {  
    /* RSASSA-PKCS1-v1_5 algorithms */  
    rsa_pkcs1_sha1 (0x0201),  
    rsa_pkcs1_sha256 (0x0401),  
    rsa_pkcs1_sha384 (0x0501),  
    rsa_pkcs1_sha512 (0x0601),  
  
    /* ECDSA algorithms */  
    ecdsa_secp256r1_sha256 (0x0403),  
    ecdsa_secp384r1_sha384 (0x0503),  
    ecdsa_secp521r1_sha512 (0x0603),  
  
    /* RSASSA-PSS algorithms */  
    rsa_pss_sha256 (0x0700),  
    rsa_pss_sha384 (0x0701),  
    rsa_pss_sha512 (0x0702),  
  
    /* EdDSA algorithms */  
    ed25519 (0x0703),  
    ed448 (0x0704),  
  
    /* Reserved Code Points */  
    private_use (0xFE00..0xFFFF),  
    (0xFFFF)  
} SignatureScheme;  
  
SignatureScheme supported_signature_algorithms<2..2^16-2>;
```

Note: This enum is named "SignatureScheme" because there is already a "SignatureAlgorithm" type in TLS 1.2, which this replaces. We use the term "signature algorithm" throughout the text.

Each SignatureScheme value lists a single signature algorithm that the client is willing to verify. The values are indicated in descending order of preference. Note that a signature algorithm takes as input an arbitrary-length message, rather than a digest. Algorithms which traditionally act on a digest should be defined in TLS to first hash the input with a specified hash function and then proceed as usual. The code point groups listed above have the following meanings:

RSASSA-PKCS1-v1_5 algorithms Indicates a signature algorithm using RSASSA-PKCS1-v1_5 [[RFC3447](#)] with the corresponding hash algorithm as defined in [[SHS](#)]. These values refer solely to signatures which appear in certificates (see [Section 4.3.1.1](#)) and are not defined for use in signed TLS handshake messages.

ECDSA algorithms Indicates a signature algorithm using ECDSA [[ECDSA](#)], the corresponding curve as defined in ANSI X9.62 [[X962](#)] and FIPS 186-4 [[DSS](#)], and the corresponding hash algorithm as defined in [[SHS](#)]. The signature is represented as a DER-encoded [[X690](#)] ECDSA-Sig-Value structure.

RSASSA-PSS algorithms Indicates a signature algorithm using RSASSA-PSS [[RFC3447](#)] with MGF1. The digest used in the mask generation function and the digest being signed are both the corresponding hash algorithm as defined in [[SHS](#)]. When used in signed TLS handshake messages, the length of the salt MUST be equal to the length of the digest output.

EdDSA algorithms Indicates a signature algorithm using EdDSA as defined in [[I-D.irtf-cfrg-eddsa](#)] or its successors. Note that these correspond to the "PureEdDSA" algorithms and not the "prehash" variants.

The semantics of this extension are somewhat complicated because the cipher suite adds additional constraints on signature algorithms. [Section 4.3.1.1](#) describes the appropriate rules.

rsa_pkcs1_sha1, dsa_sha1, and ecdsa_sha1 SHOULD NOT be offered. Clients offering these values for backwards compatibility MUST list them as the lowest priority (listed after all other algorithms in the supported_signature_algorithms vector). TLS 1.3 servers MUST NOT offer a SHA-1 signed certificate unless no valid certificate chain can be produced without it (see [Section 4.3.1.1](#)).

The signatures on certificates that are self-signed or certificates that are trust anchors are not validated since they begin a certification path (see [[RFC5280](#)], [Section 3.2](#)). A certificate that begins a certification path MAY use a signature algorithm that is not advertised as being supported in the "signature_algorithms" extension.

Note that TLS 1.2 defines this extension differently. TLS 1.3 implementations willing to negotiate TLS 1.2 MUST behave in accordance with the requirements of [[RFC5246](#)] when negotiating that version. In particular:

- TLS 1.2 ClientHellos may omit this extension.
- In TLS 1.2, the extension contained hash/signature pairs. The pairs are encoded in two octets, so SignatureScheme values have been allocated to align with TLS 1.2's encoding. Some legacy pairs are left unallocated. These algorithms are deprecated as of TLS 1.3. They MUST NOT be offered or negotiated by any

implementation. In particular, MD5 [[SLOTH](#)] and SHA-224 MUST NOT be used.

- `ecdsa_secp256r1_sha256`, etc., align with TLS 1.2's ECDSA hash/signature pairs. However, the old semantics did not constrain the signing curve.

[4.2.3](#). Negotiated Groups

When sent by the client, the "supported_groups" extension indicates the named groups which the client supports for key exchange, ordered from most preferred to least preferred.

Note: In versions of TLS prior to TLS 1.3, this extension was named "elliptic_curves" and only contained elliptic curve groups. See [[RFC4492](#)] and [[I-D.ietf-tls-negotiated-ff-dhe](#)]. This extension was also used to negotiate ECDSA curves. Signature algorithms are now negotiated independently (see [Section 4.2.2](#)).

Clients which offer one or more (EC)DHE cipher suites MUST send at least one supported NamedGroup value and servers MUST NOT negotiate any of these cipher suites unless a supported value was provided. If this extension is not provided and no alternative cipher suite is available, the server MUST close the connection with a fatal "missing_extension" alert. (see [Section 8.2](#)) If the extension is provided, but no compatible group is offered, the server MUST NOT negotiate a cipher suite of the relevant type. For instance, if a client supplies only ECDHE groups, the server MUST NOT negotiate finite field Diffie-Hellman. If no acceptable group can be selected across all cipher suites, then the server MUST generate a fatal "handshake_failure" alert.

The "extension_data" field of this extension contains a "NamedGroupList" value:


```
enum {  
    /* Elliptic Curve Groups (ECDHE) */  
    secp256r1 (23), secp384r1 (24), secp521r1 (25),  
    x25519 (29), x448 (30),  
  
    /* Finite Field Groups (DHE) */  
    ffdhe2048 (256), ffdhe3072 (257), ffdhe4096 (258),  
    ffdhe6144 (259), ffdhe8192 (260),  
  
    /* Reserved Code Points */  
    ffdhe_private_use (0x01FC..0x01FF),  
    ecdhe_private_use (0xFE00..0xFEFF),  
    (0xFFFF)  
} NamedGroup;  
  
struct {  
    NamedGroup named_group_list<1..2^16-1>;  
} NamedGroupList;
```

Elliptic Curve Groups (ECDHE) Indicates support of the corresponding named curve. Note that some curves are also recommended in ANSI X9.62 [[X962](#)] and FIPS 186-4 [[DSS](#)]. Others are recommended in [[RFC7748](#)]. Values 0xFE00 through 0xFEFF are reserved for private use.

Finite Field Groups (DHE) Indicates support of the corresponding finite field group, defined in [[I-D.ietf-tls-negotiated-ff-dhe](#)]. Values 0x01FC through 0x01FF are reserved for private use.

Items in `named_group_list` are ordered according to the client's preferences (most preferred choice first).

As of TLS 1.3, servers are permitted to send the "supported_groups" extension to the client. If the server has a group it prefers to the ones in the "key_share" extension but is still willing to accept the ClientHello, it SHOULD send "supported_groups" to update the client's view of its preferences. Clients MUST NOT act upon any information found in "supported_groups" prior to successful completion of the handshake, but MAY use the information learned from a successfully completed handshake to change what groups they offer to a server in subsequent connections.

[4.2.4.](#) Key Share

The "key_share" extension contains the endpoint's cryptographic parameters for non-PSK key establishment methods (currently DHE or ECDHE).

Clients which offer one or more (EC)DHE cipher suites MUST send this extension and SHOULD send at least one supported KeyShareEntry value. Servers MUST NOT negotiate any of these cipher suites unless a supported value was provided. If this extension is not provided in a ServerHello or ClientHello, and the peer is offering (EC)DHE cipher suites, then the endpoint MUST close the connection with a fatal "missing_extension" alert. (see [Section 8.2](#)) Clients MAY send an empty client_shares vector in order to request group selection from the server at the cost of an additional round trip. (see [Section 4.1.3](#))

```
struct {  
    NamedGroup group;  
    opaque key_exchange<1..2^16-1>;  
} KeyShareEntry;
```

group The named group for the key being exchanged. Finite Field Diffie-Hellman [[DH](#)] parameters are described in [Section 4.2.4.1](#); Elliptic Curve Diffie-Hellman parameters are described in [Section 4.2.4.2](#).

key_exchange Key exchange information. The contents of this field are determined by the specified group and its corresponding definition. Endpoints MUST NOT send empty or otherwise invalid key_exchange values for any reason.

The "extension_data" field of this extension contains a "KeyShare" value:

```
struct {  
    select (role) {  
        case client:  
            KeyShareEntry client_shares<0..2^16-1>;  
  
        case server:  
            KeyShareEntry server_share;  
    }  
} KeyShare;
```

client_shares A list of offered KeyShareEntry values in descending order of client preference. This vector MAY be empty if the client is requesting a HelloRetryRequest. The ordering of values here SHOULD match that of the ordering of offered support in the "supported_groups" extension.

server_share A single KeyShareEntry value for the negotiated cipher suite.

Clients offer an arbitrary number of KeyShareEntry values, each representing a single set of key exchange parameters. For instance, a client might offer shares for several elliptic curves or multiple FFDHE groups. The key_exchange values for each KeyShareEntry MUST be generated independently. Clients MUST NOT offer multiple KeyShareEntry values for the same group. Clients MUST NOT offer any KeyShareEntry values for groups not listed in the client's "supported_groups" extension.

Servers offer exactly one KeyShareEntry value, which corresponds to the key exchange used for the negotiated cipher suite. Servers MUST NOT offer a KeyShareEntry value for a group not offered by the client in its corresponding KeyShare or "supported_groups" extension.

Implementations MAY check for violations of these rules and MAY abort the connection with a fatal "illegal_parameter" alert if one is violated.

If the server selects an (EC)DHE cipher suite and no mutually supported group is available between the two endpoints' KeyShare offers, yet there is a mutually supported group that can be found via the "supported_groups" extension, then the server MUST reply with a HelloRetryRequest. If there is no mutually supported group at all, the server MUST NOT negotiate an (EC)DHE cipher suite.

[[TODO: Recommendation about what the client offers. Presumably which integer DH groups and which curves.]]

4.2.4.1. Diffie-Hellman Parameters

Diffie-Hellman [DH] parameters for both clients and servers are encoded in the opaque key_exchange field of a KeyShareEntry in a KeyShare structure. The opaque value contains the Diffie-Hellman public value ($Y = g^X \bmod p$), encoded as a big-endian integer, padded with zeros to the size of p in bytes.

Note: For a given Diffie-Hellman group, the padding results in all public keys having the same length.

Peers SHOULD validate each other's public key Y by ensuring that $1 < Y < p-1$. This check ensures that the remote peer is properly behaved and isn't forcing the local system into a small subgroup.

4.2.4.2. ECDHE Parameters

ECDHE parameters for both clients and servers are encoded in the opaque key_exchange field of a KeyShareEntry in a KeyShare structure.

For secp256r1, secp384r1 and secp521r1, the contents are the byte string representation of an elliptic curve public value following the conversion routine in [Section 4.3.6](#) of ANSI X9.62 [[X962](#)].

Although X9.62 supports multiple point formats, any given curve MUST specify only a single point format. All curves currently specified in this document MUST only be used with the uncompressed point format (the format for all ECDH functions is considered uncompressed).

For x25519 and x448, the contents are the byte string inputs and outputs of the corresponding functions defined in [[RFC7748](#)], 32 bytes for x25519 and 56 bytes for x448.

Note: Versions of TLS prior to 1.3 permitted point negotiation; TLS 1.3 removes this feature in favor of a single point format for each curve.

[4.2.5](#). Pre-Shared Key Extension

The "pre_shared_key" extension is used to indicate the identity of the pre-shared key to be used with a given handshake in association with a PSK or (EC)DHE-PSK cipher suite (see [[RFC4279](#)] for background).

Clients which offer one or more PSK cipher suites MUST send at least one supported psk_identity value and servers MUST NOT negotiate any of these cipher suites unless a supported value was provided. If this extension is not provided and no alternative cipher suite is available, the server MUST close the connection with a fatal "missing_extension" alert. (see [Section 8.2](#))

The "extension_data" field of this extension contains a "PreSharedKeyExtension" value:

```
opaque psk_identity<0..2^16-1>;

struct {
    select (Role) {
        case client:
            psk_identity identities<2..2^16-1>;

        case server:
            uint16 selected_identity;
    }
} PreSharedKeyExtension;
```

identities A list of the identities (labels for keys) that the client is willing to negotiate with the server. If sent alongside

the "early_data" extension (see [Section 4.2.6](#)), the first identity is the one used for 0-RTT data.

`selected_identity` The server's chosen identity expressed as a (0-based) index into the identities in the client's list.

If no suitable identity is provided, the server MUST NOT negotiate a PSK cipher suite and MAY respond with an "unknown_psk_identity" alert message. Sending this alert is OPTIONAL; servers MAY instead choose to send a "decrypt_error" alert to merely indicate an invalid PSK identity or instead negotiate use of a non-PSK cipher suite, if available.

If the server selects a PSK cipher suite, it MUST send a "pre_shared_key" extension with the identity that it selected. The client MUST verify that the server's `selected_identity` is within the range supplied by the client. If the server supplies an "early_data" extension, the client MUST verify that the server selected the first offered identity. If any other value is returned, the client MUST generate a fatal "unknown_psk_identity" alert and close the connection.

Note that although 0-RTT data is encrypted with the first PSK identity, the server MAY fall back to 1-RTT and select a different PSK identity if multiple identities are offered.

[4.2.6](#). Early Data Indication

When PSK resumption is used, the client can send application data in its first flight of messages. If the client opts to do so, it MUST supply an "early_data" extension as well as the "pre_shared_key" extension.

The "extension_data" field of this extension contains an "EarlyDataIndication" value:

```
struct {
    select (Role) {
        case client:
            uint32 obfuscated_ticket_age;

        case server:
            struct {};
    }
} EarlyDataIndication;
```

`obfuscated_ticket_age` The time since the client learned about the server configuration that it is using, in milliseconds. This

value is added modulo 2^{32} to with the "ticket_age_add" value that was included with the ticket, see [Section 4.4.1](#). This addition prevents passive observers from correlating sessions unless tickets are reused. Note: because ticket lifetimes are restricted to a week, 32 bits is enough to represent any plausible age, even in milliseconds.

A server MUST validate that the ticket_age is within a small tolerance of the time since the ticket was issued (see [Section 4.2.6.2](#)).

The parameters for the 0-RTT data (symmetric cipher suite, ALPN, etc.) are the same as those which were negotiated in the connection which established the PSK. The PSK used to encrypt the early data MUST be the first PSK listed in the client's "pre_shared_key" extension.

0-RTT messages sent in the first flight have the same content types as their corresponding messages sent in other flights (handshake, application_data, and alert respectively) but are protected under different keys. After all the 0-RTT application data messages (if any) have been sent, an "end_of_early_data" alert of type "warning" is sent to indicate the end of the flight. 0-RTT MUST always be followed by an "end_of_early_data" alert.

A server which receives an "early_data" extension can behave in one of two ways:

- Ignore the extension and return no response. This indicates that the server has ignored any early data and an ordinary 1-RTT handshake is required.
- Return an empty extension, indicating that it intends to process the early data. It is not possible for the server to accept only a subset of the early data messages.

In order to accept early data, the server MUST have accepted a PSK cipher suite and selected the the first key offered in the client's "pre_shared_key" extension. In addition, it MUST verify that the following values are consistent with those negotiated in the connection during which the ticket was established.

- The TLS version number, symmetric ciphersuite, and the hash for HKDF.
- The selected ALPN [[RFC7443](#)] value, if any.
- The server_name [[RFC6066](#)] value provided by the client, if any.

Future extensions MUST define their interaction with 0-RTT.

If any of these checks fail, the server MUST NOT respond with the extension and must discard all the remaining first flight data (thus falling back to 1-RTT). If the client attempts a 0-RTT handshake but the server rejects it, it will generally not have the 0-RTT record protection keys and must instead trial decrypt each record with the 1-RTT handshake keys until it finds one that decrypts properly, and then pick up the handshake from that point.

If the server chooses to accept the "early_data" extension, then it MUST comply with the same error handling requirements specified for all records when processing early data records. Specifically, decryption failure of any 0-RTT record following an accepted "early_data" extension MUST produce a fatal "bad_record_mac" alert as per [Section 5.2](#).

If the server rejects the "early_data" extension, the client application MAY opt to retransmit the data once the handshake has been completed. TLS stacks SHOULD not do this automatically and client applications MUST take care that the negotiated parameters are consistent with those it expected. For example, if the ALPN value has changed, it is likely unsafe to retransmit the original application layer data.

[4.2.6.1](#). Processing Order

Clients are permitted to "stream" 0-RTT data until they receive the server's Finished, only then sending the "end_of_early_data" alert. In order to avoid deadlock, when accepting "early_data", servers MUST process the client's Finished and then immediately send the ServerHello, rather than waiting for the client's "end_of_early_data" alert.

[4.2.6.2](#). Replay Properties

As noted in [Section 2.3](#), TLS provides a limited mechanism for replay protection for data sent by the client in the first flight.

The "obfuscated_ticket_age" parameter in the client's "early_data" extension SHOULD be used by servers to limit the time over which the first flight might be replayed. A server can store the time at which it sends a session ticket to the client, or encode the time in the ticket. Then, each time it receives an "early_data" extension, it can subtract the base value and check to see if the value used by the client matches its expectations.

The ticket age (the value with "ticket_age_add" subtracted) provided by the client will be shorter than the actual time elapsed on the server by a single round trip time. This difference is comprised of the delay in sending the NewSessionTicket message to the client, plus the time taken to send the ClientHello to the server. For this reason, a server SHOULD measure the round trip time prior to sending the NewSessionTicket message and account for that in the value it saves.

To properly validate the ticket age, a server needs to save at least two items:

- The time that the server generated the session ticket and the estimated round trip time can be added together to form a baseline time.
- The "ticket_age_add" parameter from the NewSessionTicket is needed to recover the ticket age from the "obfuscated_ticket_age" parameter.

There are several potential sources of error that make an exact measurement of time difficult. Variations in client and server clocks are likely to be minimal, outside of gross time corrections. Network propagation delays are most likely causes of a mismatch in legitimate values for elapsed time. Both the NewSessionTicket and ClientHello messages might be retransmitted and therefore delayed, which might be hidden by TCP.

A small allowance for errors in clocks and variations in measurements is advisable. However, any allowance also increases the opportunity for replay. In this case, it is better to reject early data than to risk greater exposure to replay attacks.

4.2.7. OCSF Status Extensions

[RFC6066] and [[RFC6961](#)] provide extensions to negotiate the server sending OCSF responses to the client. In TLS 1.2 and below, the server sends an empty extension to indicate negotiation of this extension and the OCSF information is carried in a CertificateStatus message. In TLS 1.3, the server's OCSF information is carried in an extension in EncryptedExtensions. Specifically: The body of the "status_request" or "status_request_v2" extension from the server MUST be a CertificateStatus structure as defined in [[RFC6066](#)] and [[RFC6961](#)] respectively.

Note: This means that the certificate status appears prior to the certificates it applies to. This is slightly anomalous but matches

the existing behavior for SignedCertificateTimestamps [[RFC6962](#)], and is more easily extensible in the handshake state machine.

4.2.8. Encrypted Extensions

When this message will be sent:

In all handshakes, the server MUST send the EncryptedExtensions message immediately after the ServerHello message. This is the first message that is encrypted under keys derived from handshake_traffic_secret.

Meaning of this message:

The EncryptedExtensions message contains any extensions which should be protected, i.e., any which are not needed to establish the cryptographic context.

The same extension types MUST NOT appear in both the ServerHello and EncryptedExtensions. If the same extension appears in both locations, the client MUST rely only on the value in the EncryptedExtensions block. All server-sent extensions other than those explicitly listed in [Section 4.1.2](#) or designated in the IANA registry MUST only appear in EncryptedExtensions. Extensions which are designated to appear in ServerHello MUST NOT appear in EncryptedExtensions. Clients MUST check EncryptedExtensions for the presence of any forbidden extensions and if any are found MUST terminate the handshake with an "illegal_parameter" alert.

Structure of this message:

```
struct {  
    Extension extensions<0..2^16-1>;  
} EncryptedExtensions;
```

extensions A list of extensions.

4.2.9. Certificate Request

When this message will be sent:

A non-anonymous server can optionally request a certificate from the client, if appropriate for the selected cipher suite. This message, if sent, will follow EncryptedExtensions.

Structure of this message:


```
opaque DistinguishedName<1..2^16-1>;

struct {
    opaque certificate_extension_oid<1..2^8-1>;
    opaque certificate_extension_values<0..2^16-1>;
} CertificateExtension;

struct {
    opaque certificate_request_context<0..2^8-1>;
    SignatureScheme
        supported_signature_algorithms<2..2^16-2>;
    DistinguishedName certificate_authorities<0..2^16-1>;
    CertificateExtension certificate_extensions<0..2^16-1>;
} CertificateRequest;
```

certificate_request_context An opaque string which identifies the certificate request and which will be echoed in the client's Certificate message. The `certificate_request_context` MUST be unique within the scope of this connection (thus preventing replay of client CertificateVerify messages).

supported_signature_algorithms A list of the signature algorithms that the server is able to verify, listed in descending order of preference. Any certificates provided by the client MUST be signed using a signature algorithm found in `supported_signature_algorithms`.

certificate_authorities A list of the distinguished names [X501] of acceptable certificate authorities, represented in DER-encoded [X690] format. These distinguished names may specify a desired distinguished name for a root CA or for a subordinate CA; thus, this message can be used to describe known roots as well as a desired authorization space. If the `certificate_authorities` list is empty, then the client MAY send any certificate that meets the rest of the selection criteria in the CertificateRequest, unless there is some external arrangement to the contrary.

certificate_extensions A list of certificate extension OIDs [RFC5280] with their allowed values, represented in DER-encoded [X690] format. Some certificate extension OIDs allow multiple values (e.g. Extended Key Usage). If the server has included a non-empty `certificate_extensions` list, the client certificate MUST contain all of the specified extension OIDs that the client recognizes. For each extension OID recognized by the client, all of the specified values MUST be present in the client certificate (but the certificate MAY have other values as well). However, the client MUST ignore and skip any unrecognized certificate extension OIDs. If the client has ignored some of the required certificate

extension OIDs, and supplied a certificate that does not satisfy the request, the server MAY at its discretion either continue the session without client authentication, or terminate the session with a fatal `unsupported_certificate` alert. PKIX RFCs define a variety of certificate extension OIDs and their corresponding value types. Depending on the type, matching certificate extension values are not necessarily bitwise-equal. It is expected that TLS implementations will rely on their PKI libraries to perform certificate selection using certificate extension OIDs. This document defines matching rules for two standard certificate extensions defined in [\[RFC5280\]](#):

- o The Key Usage extension in a certificate matches the request when all key usage bits asserted in the request are also asserted in the Key Usage certificate extension.
- o The Extended Key Usage extension in a certificate matches the request when all key purpose OIDs present in the request are also found in the Extended Key Usage certificate extension. The special `anyExtendedKeyUsage` OID MUST NOT be used in the request.

Separate specifications may define matching rules for other certificate extensions.

Note: It is a fatal `"handshake_failure"` alert for an anonymous server to request client authentication.

[4.3.](#) Authentication Messages

As discussed in [Section 2](#), TLS uses a common set of messages for authentication, key confirmation, and handshake integrity: Certificate, CertificateVerify, and Finished. These messages are always sent as the last messages in their handshake flight. The Certificate and CertificateVerify messages are only sent under certain circumstances, as defined below. The Finished message is always sent as part of the Authentication block.

The computations for the Authentication messages all uniformly take the following inputs:

- The certificate and signing key to be used.
- A Handshake Context based on the hash of the handshake messages
- A base key to be used to compute a MAC key.

Based on these inputs, the messages then contain:

Certificate The certificate to be used for authentication and any supporting certificates in the chain. Note that certificate-based client authentication is not available in the 0-RTT case.

CertificateVerify A signature over the value Hash(Handshake Context + Certificate) + Hash(resumption_context) See [Section 4.4.1](#) for the definition of resumption_context.

Finished A MAC over the value Hash(Handshake Context + Certificate + CertificateVerify) + Hash(resumption_context) using a MAC key derived from the base key.

Because the CertificateVerify signs the Handshake Context + Certificate and the Finished MACs the Handshake Context + Certificate + CertificateVerify, this is mostly equivalent to keeping a running hash of the handshake messages (exactly so in the pure 1-RTT cases). Note, however, that subsequent post-handshake authentications do not include each other, just the messages through the end of the main handshake.

The following table defines the Handshake Context and MAC Base Key for each scenario:

Mode	Handshake Context	Base Key
0-RTT	ClientHello	early_traffic_secret
1-RTT (Server)	ClientHello ... later of EncryptedExtensions/CertificateRequest	handshake_traffic_secret
1-RTT (Client)	ClientHello ... ServerFinished	handshake_traffic_secret
Post-Handshake	ClientHello ... ClientFinished + CertificateRequest	traffic_secret_0

Note: The Handshake Context for the last three rows does not include any 0-RTT handshake messages, regardless of whether 0-RTT is used.

[4.3.1. Certificate](#)

When this message will be sent:

The server MUST send a Certificate message whenever the agreed-upon key exchange method uses certificates for authentication (this includes all key exchange methods defined in this document except PSK).

The client MUST send a Certificate message if and only if server has requested client authentication via a CertificateRequest message ([Section 4.2.9](#)). If the server requests client authentication but no suitable certificate is available, the client MUST send a Certificate message containing no certificates (i.e., with the "certificate_list" field having length 0).

Meaning of this message:

This message conveys the endpoint's certificate chain to the peer.

The certificate MUST be appropriate for the negotiated cipher suite's authentication algorithm and any negotiated extensions.

Structure of this message:

```
opaque ASN1Cert<1..224-1>;

struct {
    opaque certificate_request_context<0..28-1>;
    ASN1Cert certificate_list<0..224-1>;
} Certificate;
```

certificate_request_context If this message is in response to a CertificateRequest, the value of certificate_request_context in that message. Otherwise, in the case of server authentication this field SHALL be zero length.

certificate_list This is a sequence (chain) of certificates. The sender's certificate MUST come first in the list. Each following certificate SHOULD directly certify one preceding it. Because certificate validation requires that trust anchors be distributed independently, a certificate that specifies a trust anchor MAY be omitted from the chain, provided that supported peers are known to possess any omitted certificates.

Note: Prior to TLS 1.3, "certificate_list" ordering required each certificate to certify the one immediately preceding it, however some implementations allowed some flexibility. Servers sometimes send both a current and deprecated intermediate for transitional purposes, and others are simply configured incorrectly, but these cases can nonetheless be validated properly. For maximum compatibility, all implementations SHOULD be prepared to handle potentially extraneous

certificates and arbitrary orderings from any TLS version, with the exception of the end-entity certificate which **MUST** be first.

The server's certificate list **MUST** always be non-empty. A client will send an empty certificate list if it does not have an appropriate certificate to send in response to the server's authentication request.

4.3.1.1. Server Certificate Selection

The following rules apply to the certificates sent by the server:

- The certificate type **MUST** be X.509v3 [[RFC5280](#)], unless explicitly negotiated otherwise (e.g., [[RFC5081](#)]).
- The server's end-entity certificate's public key (and associated restrictions) **MUST** be compatible with the selected authentication algorithm (currently RSA or ECDSA).
- The certificate **MUST** allow the key to be used for signing (i.e., the `digitalSignature` bit **MUST** be set if the Key Usage extension is present) with a signature scheme indicated in the client's `"signature_algorithms"` extension.
- The `"server_name"` and `"trusted_ca_keys"` extensions [[RFC6066](#)] are used to guide certificate selection. As servers **MAY** require the presence of the `"server_name"` extension, clients **SHOULD** send this extension, when applicable.

All certificates provided by the server **MUST** be signed by a signature algorithm that appears in the `"signature_algorithms"` extension provided by the client, if they are able to provide such a chain (see [Section 4.2.2](#)). Certificates that are self-signed or certificates that are expected to be trust anchors are not validated as part of the chain and therefore **MAY** be signed with any algorithm.

If the server cannot produce a certificate chain that is signed only via the indicated supported algorithms, then it **SHOULD** continue the handshake by sending the client a certificate chain of its choice that may include algorithms that are not known to be supported by the client. This fallback chain **MAY** use the deprecated SHA-1 hash algorithm only if the `"signature_algorithms"` extension provided by the client permits it. If the client cannot construct an acceptable chain using the provided certificates and decides to abort the handshake, then it **MUST** send an `"unsupported_certificate"` alert message and close the connection.

If the server has multiple certificates, it chooses one of them based on the above-mentioned criteria (in addition to other criteria, such as transport layer endpoint, local configuration and preferences).

As cipher suites that specify new key exchange methods are specified for the TLS protocol, they will imply the certificate format and the required encoded keying information.

4.3.1.2. Client Certificate Selection

The following rules apply to certificates sent by the client:

In particular:

- The certificate type MUST be X.509v3 [[RFC5280](#)], unless explicitly negotiated otherwise (e.g., [[RFC5081](#)]).
- If the `certificate_authorities` list in the certificate request message was non-empty, one of the certificates in the certificate chain SHOULD be issued by one of the listed CAs.
- The certificates MUST be signed using an acceptable signature algorithm, as described in [Section 4.2.9](#). Note that this relaxes the constraints on certificate-signing algorithms found in prior versions of TLS.
- If the `certificate_extensions` list in the certificate request message was non-empty, the end-entity certificate MUST match the extension OIDs recognized by the client, as described in [Section 4.2.9](#).

Note that, as with the server certificate, there are certificates that use algorithm combinations that cannot be currently used with TLS.

4.3.1.3. Receiving a Certificate Message

In general, detailed certificate validation procedures are out of scope for TLS (see [[RFC5280](#)]). This section provides TLS-specific requirements.

If the server supplies an empty Certificate message, the client MUST terminate the handshake with a fatal "decode_error" alert.

If the client does not send any certificates, the server MAY at its discretion either continue the handshake without client authentication, or respond with a fatal "handshake_failure" alert. Also, if some aspect of the certificate chain was unacceptable (e.g.,

it was not signed by a known, trusted CA), the server MAY at its discretion either continue the handshake (considering the client unauthenticated) or send a fatal alert.

Any endpoint receiving any certificate signed using any signature algorithm using an MD5 hash MUST send a "bad_certificate" alert message and close the connection. SHA-1 is deprecated and therefore NOT RECOMMENDED. All endpoints are RECOMMENDED to transition to SHA-256 or better as soon as possible to maintain interoperability with implementations currently in the process of phasing out SHA-1 support.

Note that a certificate containing a key for one signature algorithm MAY be signed using a different signature algorithm (for instance, an RSA key signed with an ECDSA key).

Endpoints that reject certification paths due to use of a deprecated hash MUST send a fatal "bad_certificate" alert message before closing the connection.

[4.3.2](#). Certificate Verify

When this message will be sent:

This message is used to provide explicit proof that an endpoint possesses the private key corresponding to its certificate and also provides integrity for the handshake up to this point. Servers MUST send this message when using a cipher suite which is authenticated via a certificate. Clients MUST send this message whenever authenticating via a Certificate (i.e., when the Certificate message is non-empty). When sent, this message MUST appear immediately after the Certificate Message and immediately prior to the Finished message.

Structure of this message:

```
struct {  
    SignatureScheme algorithm;  
    opaque signature<0..2^16-1>;  
} CertificateVerify;
```

The algorithm field specifies the signature algorithm used (see [Section 4.2.2](#) for the definition of this field). The signature is a digital signature using that algorithm that covers the hash output described in [Section 4.3](#) namely:

Hash(Handshake Context + Certificate) + Hash(resumption_context)

the server MUST check any candidate cipher suites against the "signature_algorithms" extension before selecting them. This is somewhat inelegant but is a compromise designed to minimize changes to the original cipher suite design.

If sent by a client, the signature algorithm used in the signature MUST be one of those present in the supported_signature_algorithms field of the CertificateRequest message.

In addition, the signature algorithm MUST be compatible with the key in the sender's end-entity certificate. RSA signatures MUST use an RSASSA-PSS algorithm, regardless of whether RSASSA-PKCS1-v1_5 algorithms appear in "signature_algorithms". SHA-1 MUST NOT be used in any signatures in CertificateVerify. All SHA-1 signature algorithms in this specification are defined solely for use in legacy certificates, and are not valid for CertificateVerify signatures.

Note: When used with non-certificate-based handshakes (e.g., PSK), the client's signature does not cover the server's certificate directly, although it does cover the server's Finished message, which transitively includes the server's certificate when the PSK derives from a certificate-authenticated handshake. [\[PSK-FINISHED\]](#) describes a concrete attack on this mode if the Finished is omitted from the signature. It is unsafe to use certificate-based client authentication when the client might potentially share the same PSK/key-id pair with two different endpoints. In order to ensure this, implementations MUST NOT mix certificate-based client authentication with pure PSK modes (i.e., those where the PSK was not derived from a previous non-PSK handshake).

[4.3.3. Finished](#)

When this message will be sent:

The Finished message is the final message in the authentication block. It is essential for providing authentication of the handshake and of the computed keys.

Meaning of this message:

Recipients of Finished messages MUST verify that the contents are correct. Once a side has sent its Finished message and received and validated the Finished message from its peer, it may begin to send and receive application data over the connection.

The key used to compute the finished message is computed from the Base key defined in [Section 4.3](#) using HKDF (see [Section 7.1](#)). Specifically:


```
client_finished_key =  
    HKDF-Expand-Label(BaseKey, "client finished", "", Hash.Length)  
  
server_finished_key =  
    HKDF-Expand-Label(BaseKey, "server finished", "", Hash.Length)
```

Structure of this message:

```
struct {  
    opaque verify_data[Hash.length];  
} Finished;
```

The verify_data value is computed as follows:

```
verify_data =  
    HMAC(finished_key, Hash(  
        Handshake Context +  
        Certificate* +  
        CertificateVerify*  
    ) +  
    Hash(resumption_context)  
)
```

* Only included if present.

Where HMAC [[RFC2104](#)] uses the Hash algorithm for the handshake. As noted above, the HMAC input can generally be implemented by a running hash, i.e., just the handshake hash at this point.

In previous versions of TLS, the verify_data was always 12 octets long. In the current version of TLS, it is the size of the HMAC output for the Hash used for the handshake.

Note: Alerts and any other record types are not handshake messages and are not included in the hash computations.

4.4. Post-Handshake Messages

TLS also allows other messages to be sent after the main handshake. These messages use a handshake content type and are encrypted under the application traffic key.

4.4.1. New Session Ticket Message

At any time after the server has received the client Finished message, it MAY send a NewSessionTicket message. This message creates a pre-shared key (PSK) binding between the ticket value and the following two values derived from the resumption master secret:


```
resumption_psk = HKDF-Expand-Label(
    resumption_secret,
    "resumption psk", "", Hash.Length)

resumption_context = HKDF-Expand-Label(
    resumption_secret,
    "resumption context", "", Hash.Length)
```

The client MAY use this PSK for future handshakes by including the ticket value in the "pre_shared_key" extension in its ClientHello ([Section 4.2.5](#)) and supplying a suitable PSK cipher suite. Servers may send multiple tickets on a single connection, for instance after post-handshake authentication. For handshakes that do not use a resumption_psk, the resumption_context is a string of Hash.Length zeroes.

```
enum { (65535) } TicketExtensionType;

struct {
    TicketExtensionType extension_type;
    opaque extension_data<1..2^16-1>;
} TicketExtension;

enum {
    allow_early_data(1),
    allow_dhe_resumption(2),
    allow_psk_resumption(4)
} TicketFlags;

struct {
    uint32 ticket_lifetime;
    uint32 flags;
    uint32 ticket_age_add;
    TicketExtension extensions<2..2^16-2>;
    opaque ticket<0..2^16-1>;
} NewSessionTicket;
```

flags A 32-bit value indicating the ways in which this ticket may be used (as a bitwise OR of the flags values).

ticket_lifetime Indicates the lifetime in seconds as a 32-bit unsigned integer in network byte order from the time of ticket issuance. Servers MUST NOT use any value more than 604800 seconds (7 days). The value of zero indicates that the ticket should be discarded immediately. Clients MUST NOT cache session tickets for longer than 7 days, regardless of the ticket_lifetime. It MAY delete the ticket earlier based on local policy. A server MAY

treat a ticket as valid for a shorter period of time than what is stated in the `ticket_lifetime`.

`ticket_age_add` A randomly generated 32-bit value that is used to obscure the age of the ticket that the client includes in the "early_data" extension. The actual ticket age is added to this value modulo 2^{32} to obtain the value that is transmitted by the client.

`ticket_extensions` A placeholder for extensions in the ticket. Clients MUST ignore unrecognized extensions.

`ticket` The value of the ticket to be used as the PSK identifier. The ticket itself is an opaque label. It MAY either be a database lookup key or a self-encrypted and self-authenticated value. [Section 4 of \[RFC5077\]](#) describes a recommended ticket construction mechanism.

The meanings of the flags are as follows:

`allow_early_data` When resuming with this ticket, the client MAY send data in its first flight (early data) encrypted under a key derived from this PSK.

`allow_dhe_resumption` This ticket MAY be used with (EC)DHE-PSK cipher suite.

`allow_psk_resumption` This ticket MAY be used with a pure PSK cipher suite.

In all cases, the PSK or (EC)DHE-PSK cipher suites that the client offers/uses MUST have the same symmetric parameters (cipher/hash) as the cipher suite negotiated for this connection. If no flags are set that the client recognizes, it MUST ignore the ticket.

4.4.2. Post-Handshake Authentication

The server is permitted to request client authentication at any time after the handshake has completed by sending a `CertificateRequest` message. The client SHOULD respond with the appropriate Authentication messages. If the client chooses to authenticate, it MUST send `Certificate`, `CertificateVerify`, and `Finished`. If it declines, it MUST send a `Certificate` message containing no certificates followed by `Finished`.

Note: Because client authentication may require prompting the user, servers MUST be prepared for some delay, including receiving an arbitrary number of other messages between sending the

CertificateRequest and receiving a response. In addition, clients which receive multiple CertificateRequests in close succession MAY respond to them in a different order than they were received (the `certificate_request_context` value allows the server to disambiguate the responses).

4.4.3. Key and IV Update

```
struct {} KeyUpdate;
```

The KeyUpdate handshake message is used to indicate that the sender is updating its sending cryptographic keys. This message can be sent by the server after sending its first flight and the client after sending its second flight. Implementations that receive a KeyUpdate message prior to receiving a Finished message as part of the 1-RTT handshake MUST generate a fatal "unexpected_message" alert. After sending a KeyUpdate message, the sender SHALL send all its traffic using the next generation of keys, computed as described in [Section 7.2](#). Upon receiving a KeyUpdate, the receiver MUST update their receiving keys and if they have not already updated their sending state up to or past the then current receiving generation MUST send their own KeyUpdate prior to sending any other messages. This mechanism allows either side to force an update to the entire connection. Note that implementations may receive an arbitrary number of messages between sending a KeyUpdate and receiving the peer's KeyUpdate because those messages may already be in flight.

Note that if implementations independently send their own KeyUpdates and they cross in flight, this only results in an update of one generation; when each side receives the other side's update it just updates its receive keys and notes that the generations match and thus no send update is needed.

Note that the side which sends its KeyUpdate first needs to retain its receive traffic keys (though not the traffic secret) for the previous generation of keys until it receives the KeyUpdate from the other side.

Both sender and receiver MUST encrypt their KeyUpdate messages with the old keys. Additionally, both sides MUST enforce that a KeyUpdate with the old key is received before accepting any messages encrypted with the new key. Failure to do so may allow message truncation attacks.

5. Record Protocol

The TLS record protocol takes messages to be transmitted, fragments the data into manageable blocks, protects the records, and transmits the result. Received data is decrypted and verified, reassembled, and then delivered to higher-level clients.

TLS records are typed, which allows multiple higher level protocols to be multiplexed over the same record layer. This document specifies three content types: handshake, application data, and alert. Implementations **MUST NOT** send record types not defined in this document unless negotiated by some extension. If a TLS implementation receives an unexpected record type, it **MUST** send an "unexpected_message" alert. New record content type values are assigned by IANA in the TLS Content Type Registry as described in [Section 10](#).

Application data messages are carried by the record layer and are fragmented and encrypted as described below. The messages are treated as transparent data to the record layer.

5.1. Record Layer

The TLS record layer receives uninterpreted data from higher layers in non-empty blocks of arbitrary size.

The record layer fragments information blocks into TLSPlaintext records carrying data in chunks of 2^{14} bytes or less. Message boundaries are not preserved in the record layer (i.e., multiple messages of the same ContentType **MAY** be coalesced into a single TLSPlaintext record, or a single message **MAY** be fragmented across several records). Alert messages ([Section 6](#)) **MUST NOT** be fragmented across records.

```
enum {
    alert(21),
    handshake(22),
    application_data(23)
    (255)
} ContentType;

struct {
    ContentType type;
    ProtocolVersion record_version = { 3, 1 };    /* TLS v1.x */
    uint16 length;
    opaque fragment[TLSPlaintext.length];
} TLSPlaintext;
```


`type` The higher-level protocol used to process the enclosed fragment.

`record_version` The protocol version the current record is compatible with. This value MUST be set to { 3, 1 } for all records. This field is deprecated and MUST be ignored for all purposes.

`length` The length (in bytes) of the following `TLSPplaintext.fragment`. The length MUST NOT exceed 2^{14} .

`fragment` The data being transmitted. This value transparent and treated as an independent block to be dealt with by the higher-level protocol specified by the `type` field.

This document describes TLS Version 1.3, which uses the version { 3, 4 }. The version value 3.4 is historical, deriving from the use of { 3, 1 } for TLS 1.0 and { 3, 0 } for SSL 3.0. In order to maximize backwards compatibility, the record layer version identifies as simply TLS 1.0. Endpoints supporting other versions negotiate the version to use by following the procedure and requirements in [Appendix C](#).

Implementations MUST NOT send zero-length fragments of Handshake or Alert types, even if those fragments contain padding. Zero-length fragments of Application data MAY be sent as they are potentially useful as a traffic analysis countermeasure.

When record protection has not yet been engaged, `TLSPplaintext` structures are written directly onto the wire. Once record protection has started, `TLSPplaintext` records are protected and sent as described in the following section.

[5.2](#). Record Payload Protection

The record protection functions translate a `TLSPplaintext` structure into a `TLSCiphertext`. The deprotection functions reverse the process. In TLS 1.3 as opposed to previous versions of TLS, all ciphers are modeled as "Authenticated Encryption with Additional Data" (AEAD) [[RFC5116](#)]. AEAD functions provide a unified encryption and authentication operation which turns plaintext into authenticated ciphertext and back again. Each encrypted record consists of a plaintext header followed by an encrypted body, which itself contains a type and optional padding.


```
struct {
    opaque content[TLSPlaintext.length];
    ContentType type;
    uint8 zeros[length_of_padding];
} TLSInnerPlaintext;

struct {
    ContentType opaque_type = application_data(23); /* see fragment.type */
    ProtocolVersion record_version = { 3, 1 }; /* TLS v1.x */
    uint16 length;
    opaque encrypted_record[length];
} TLSCiphertext;
```

content The cleartext of TLSPlaintext.fragment.

type The content type of the record.

zeros An arbitrary-length run of zero-valued bytes may appear in the cleartext after the type field. This provides an opportunity for senders to pad any TLS record by a chosen amount as long as the total stays within record size limits. See [Section 5.4](#) for more details.

opaque_type The outer opaque_type field of a TLSCiphertext record is always set to the value 23 (application_data) for outward compatibility with middleboxes accustomed to parsing previous versions of TLS. The actual content type of the record is found in fragment.type after decryption.

record_version The record_version field is identical to TLSPlaintext.record_version and is always { 3, 1 }. Note that the handshake protocol including the ClientHello and ServerHello messages authenticates the protocol version, so this value is redundant.

length The length (in bytes) of the following TLSCiphertext.fragment, which is the sum of the lengths of the content and the padding, plus one for the inner content type. The length MUST NOT exceed $2^{14} + 256$. An endpoint that receives a record that exceeds this length MUST generate a fatal "record_overflow" alert.

encrypted_record The AEAD encrypted form of the serialized TLSInnerPlaintext structure.

AEAD ciphers take as input a single key, a nonce, a plaintext, and "additional data" to be included in the authentication check, as described in [Section 2.1 of \[RFC5116\]](#). The key is either the

client_write_key or the server_write_key, the nonce is derived from the sequence number (see [Section 5.3](#)) and the client_write_iv or server_write_iv, and the additional data input is empty (zero length). Derivation of traffic keys is defined in [Section 7.3](#).

The plaintext is the concatenation of TLSPlaintext.fragment, TLSPlaintext.type, and any padding bytes (zeros).

The AEAD output consists of the ciphertext output by the AEAD encryption operation. The length of the plaintext is greater than TLSPlaintext.length due to the inclusion of TLSPlaintext.type and however much padding is supplied by the sender. The length of the AEAD output will generally be larger than the plaintext, but by an amount that varies with the AEAD cipher. Since the ciphers might incorporate padding, the amount of overhead could vary with different lengths of plaintext. Symbolically,

```
AEADEncrypted =  
    AEAD-Encrypt(write_key, nonce, plaintext of fragment)
```

In order to decrypt and verify, the cipher takes as input the key, nonce, and the AEADEncrypted value. The output is either the plaintext or an error indicating that the decryption failed. There is no separate integrity check. That is:

```
plaintext of fragment =  
    AEAD-Decrypt(write_key, nonce, AEADEncrypted)
```

If the decryption fails, a fatal "bad_record_mac" alert MUST be generated.

An AEAD cipher MUST NOT produce an expansion of greater than 255 bytes. An endpoint that receives a record from its peer with TLSCipherText.length larger than $2^{14} + 256$ octets MUST generate a fatal "record_overflow" alert. This limit is derived from the maximum TLSPlaintext length of 2^{14} octets + 1 octet for ContentType + the maximum AEAD expansion of 255 octets.

[5.3](#). Per-Record Nonce

A 64-bit sequence number is maintained separately for reading and writing records. Each sequence number is set to zero at the beginning of a connection and whenever the key is changed.

The sequence number is incremented after reading or writing each record. The first record transmitted under a particular set of traffic keys record key MUST use sequence number 0.

Sequence numbers do not wrap. If a TLS implementation would need to wrap a sequence number, it MUST either rekey ([Section 4.4.3](#)) or terminate the connection.

The length of the per-record nonce (`iv_length`) is set to $\max(8 \text{ bytes}, N_{\text{MIN}})$ for the AEAD algorithm (see [\[RFC5116\] Section 4](#)). An AEAD algorithm where N_{MAX} is less than 8 bytes MUST NOT be used with TLS. The per-record nonce for the AEAD construction is formed as follows:

1. The 64-bit record sequence number is padded to the left with zeroes to `iv_length`.
2. The padded sequence number is XORed with the static `client_write_iv` or `server_write_iv`, depending on the role.

The resulting quantity (of length `iv_length`) is used as the per-record nonce.

Note: This is a different construction from that in TLS 1.2, which specified a partially explicit nonce.

5.4. Record Padding

All encrypted TLS records can be padded to inflate the size of the `TLSCipherText`. This allows the sender to hide the size of the traffic from an observer.

When generating a `TLSCiphertext` record, implementations MAY choose to pad. An unpadded record is just a record with a padding length of zero. Padding is a string of zero-valued bytes appended to the `ContentType` field before encryption. Implementations MUST set the padding octets to all zeros before encrypting.

Application Data records may contain a zero-length `fragment.content` if the sender desires. This permits generation of plausibly-sized cover traffic in contexts where the presence or absence of activity may be sensitive. Implementations MUST NOT send Handshake or Alert records that have a zero-length `fragment.content`.

The padding sent is automatically verified by the record protection mechanism: Upon successful decryption of a `TLSCiphertext.fragment`, the receiving implementation scans the field from the end toward the beginning until it finds a non-zero octet. This non-zero octet is the content type of the message. This padding scheme was selected because it allows padding of any encrypted TLS record by an arbitrary size (from zero up to TLS record size limits) without introducing new content types. The design also enforces all-zero padding octets, which allows for quick detection of padding errors.

Implementations MUST limit their scanning to the cleartext returned from the AEAD decryption. If a receiving implementation does not find a non-zero octet in the cleartext, it should treat the record as having an unexpected ContentType, sending an "unexpected_message" alert.

The presence of padding does not change the overall record size limitations - the full fragment plaintext may not exceed 2^{14} octets.

Selecting a padding policy that suggests when and how much to pad is a complex topic, and is beyond the scope of this specification. If the application layer protocol atop TLS has its own padding padding, it may be preferable to pad application_data TLS records within the application layer. Padding for encrypted handshake and alert TLS records must still be handled at the TLS layer, though. Later documents may define padding selection algorithms, or define a padding policy request mechanism through TLS extensions or some other means.

5.5. Limits on Key Usage

There are cryptographic limits on the amount of plaintext which can be safely encrypted under a given set of keys. [\[AEAD-LIMITS\]](#) provides an analysis of these limits under the assumption that the underlying primitive (AES or ChaCha20) has no weaknesses. Implementations SHOULD do a key update [Section 4.4.3](#) prior to reaching these limits.

For AES-GCM, up to $2^{24.5}$ full-size records may be encrypted on a given connection while keeping a safety margin of approximately 2^{57} for Authenticated Encryption (AE) security. For ChaCha20/Poly1305, the record sequence number will wrap before the safety limit is reached.

6. Alert Protocol

One of the content types supported by the TLS record layer is the alert type. Like other messages, alert messages are encrypted as specified by the current connection state.

Alert messages convey the severity of the message (warning or fatal) and a description of the alert. Warning-level messages are used to indicate orderly closure of the connection (see [Section 6.1](#)). Upon receiving a warning-level alert, the TLS implementation SHOULD indicate end-of-data to the application and, if appropriate for the alert type, send a closure alert in response.

Fatal-level messages are used to indicate abortive closure of the connection (See [Section 6.2](#)). Upon receiving a fatal-level alert, the TLS implementation SHOULD indicate an error to the application and MUST NOT allow any further data to be sent or received on the connection. Servers and clients MUST forget keys and secrets associated with a failed connection. Stateful implementations of session tickets (as in many clients) SHOULD discard tickets associated with failed connections.

All the alerts listed in [Section 6.2](#) MUST be sent as fatal and MUST be treated as fatal regardless of the AlertLevel in the message. Unknown alert types MUST be treated as fatal.


```
enum { warning(1), fatal(2), (255) } AlertLevel;
```

```
enum {  
    close_notify(0),  
    end_of_early_data(1),  
    unexpected_message(10),  
    bad_record_mac(20),  
    record_overflow(22),  
    handshake_failure(40),  
    bad_certificate(42),  
    unsupported_certificate(43),  
    certificate_revoked(44),  
    certificate_expired(45),  
    certificate_unknown(46),  
    illegal_parameter(47),  
    unknown_ca(48),  
    access_denied(49),  
    decode_error(50),  
    decrypt_error(51),  
    protocol_version(70),  
    insufficient_security(71),  
    internal_error(80),  
    inappropriate_fallback(86),  
    user_canceled(90),  
    missing_extension(109),  
    unsupported_extension(110),  
    certificate_unobtainable(111),  
    unrecognized_name(112),  
    bad_certificate_status_response(113),  
    bad_certificate_hash_value(114),  
    unknown_psk_identity(115),  
    (255)  
} AlertDescription;  
  
struct {  
    AlertLevel level;  
    AlertDescription description;  
} Alert;
```

6.1. Closure Alerts

The client and the server must share knowledge that the connection is ending in order to avoid a truncation attack. Failure to properly close a connection does not prohibit a session from being resumed.

close_notify This alert notifies the recipient that the sender will not send any more messages on this connection. Any data received after a closure **MUST** be ignored.

`end_of_early_data` This alert is sent by the client to indicate that all 0-RTT `application_data` messages have been transmitted (or none will be sent at all) and that this is the end of the flight. This alert **MUST** be at the warning level. Servers **MUST NOT** send this alert and clients receiving it **MUST** terminate the connection with an "unexpected_message" alert.

`user_canceled` This alert notifies the recipient that the sender is canceling the handshake for some reason unrelated to a protocol failure. If a user cancels an operation after the handshake is complete, just closing the connection by sending a "close_notify" is more appropriate. This alert **SHOULD** be followed by a "close_notify". This alert is generally a warning.

Either party **MAY** initiate a close by sending a "close_notify" alert. Any data received after a closure alert is ignored. If a transport-level close is received prior to a "close_notify", the receiver cannot know that all the data that was sent has been received.

Each party **MUST** send a "close_notify" alert before closing the write side of the connection, unless some other fatal alert has been transmitted. The other party **MUST** respond with a "close_notify" alert of its own and close down the connection immediately, discarding any pending writes. The initiator of the close need not wait for the responding "close_notify" alert before closing the read side of the connection.

If the application protocol using TLS provides that any data may be carried over the underlying transport after the TLS connection is closed, the TLS implementation must receive the responding "close_notify" alert before indicating to the application layer that the TLS connection has ended. If the application protocol will not transfer any additional data, but will only close the underlying transport connection, then the implementation **MAY** choose to close the transport without waiting for the responding "close_notify". No part of this standard should be taken to dictate the manner in which a usage profile for TLS manages its data transport, including when connections are opened or closed.

Note: It is assumed that closing a connection reliably delivers pending data before destroying the transport.

6.2. Error Alerts

Error handling in the TLS Handshake Protocol is very simple. When an error is detected, the detecting party sends a message to its peer. Upon transmission or receipt of a fatal alert message, both parties immediately close the connection. Whenever an implementation

encounters a condition which is defined as a fatal alert, it **MUST** send the appropriate alert prior to closing the connection. All alerts defined in this section below, as well as all unknown alerts are universally considered fatal as of TLS 1.3 (see [Section 6](#)).

The following error alerts are defined:

unexpected_message An inappropriate message was received. This alert should never be observed in communication between proper implementations.

bad_record_mac This alert is returned if a record is received which cannot be deprotected. Because AEAD algorithms combine decryption and verification, this alert is used for all deprotection failures. This alert should never be observed in communication between proper implementations, except when messages were corrupted in the network.

record_overflow A TLSCiphertext record was received that had a length more than $2^{14} + 256$ bytes, or a record decrypted to a TLSPlaintext record with more than 2^{14} bytes. This alert should never be observed in communication between proper implementations, except when messages were corrupted in the network.

handshake_failure Reception of a "handshake_failure" alert message indicates that the sender was unable to negotiate an acceptable set of security parameters given the options available.

bad_certificate A certificate was corrupt, contained signatures that did not verify correctly, etc.

unsupported_certificate A certificate was of an unsupported type.

certificate_revoked A certificate was revoked by its signer.

certificate_expired A certificate has expired or is not currently valid.

certificate_unknown Some other (unspecified) issue arose in processing the certificate, rendering it unacceptable.

illegal_parameter A field in the handshake was out of range or inconsistent with other fields.

unknown_ca A valid certificate chain or partial chain was received, but the certificate was not accepted because the CA certificate could not be located or couldn't be matched with a known, trusted CA.

access_denied A valid certificate or PSK was received, but when access control was applied, the sender decided not to proceed with negotiation.

decode_error A message could not be decoded because some field was out of the specified range or the length of the message was incorrect. This alert should never be observed in communication between proper implementations, except when messages were corrupted in the network.

decrypt_error A handshake cryptographic operation failed, including being unable to correctly verify a signature or validate a Finished message.

protocol_version The protocol version the peer has attempted to negotiate is recognized but not supported. (see [Appendix C](#))

insufficient_security Returned instead of "handshake_failure" when a negotiation has failed specifically because the server requires ciphers more secure than those supported by the client.

internal_error An internal error unrelated to the peer or the correctness of the protocol (such as a memory allocation failure) makes it impossible to continue.

inappropriate_fallback Sent by a server in response to an invalid connection retry attempt from a client. (see [\[RFC7507\]](#))

missing_extension Sent by endpoints that receive a hello message not containing an extension that is mandatory to send for the offered TLS version. [\[\[TODO: IANA Considerations.\]\]](#)

unsupported_extension Sent by endpoints receiving any hello message containing an extension known to be prohibited for inclusion in the given hello message, including any extensions in a ServerHello not first offered in the corresponding ClientHello.

certificate_unobtainable Sent by servers when unable to obtain a certificate from a URL provided by the client via the "client_certificate_url" extension [\[RFC6066\]](#).

unrecognized_name Sent by servers when no server exists identified by the name provided by the client via the "server_name" extension [\[RFC6066\]](#).

bad_certificate_status_response Sent by clients when an invalid or unacceptable OCSP response is provided by the server via the "status_request" extension [\[RFC6066\]](#). This alert is always fatal.

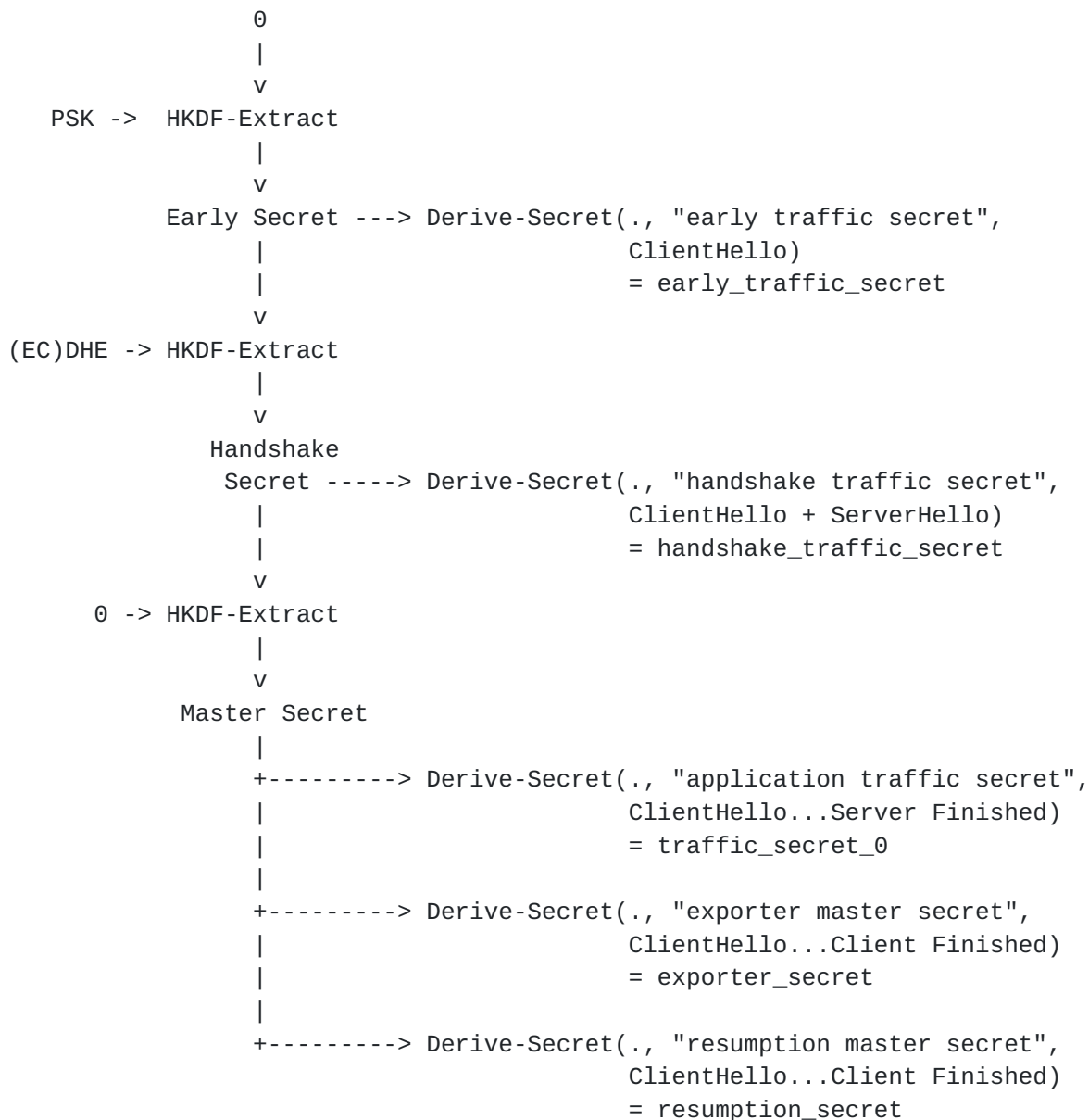
The Hash function and the HKDF hash are the cipher suite hash function. Hash.Length is its output length.

Given a set of n InputSecrets, the final "master secret" is computed by iteratively invoking HKDF-Extract with InputSecret_1, InputSecret_2, etc. The initial secret is simply a string of zeroes as long as the size of the Hash that is the basis for the HKDF. Concretely, for the present version of TLS 1.3, secrets are added in the following order:

- PSK
- (EC)DHE shared secret

This produces a full key derivation schedule shown in the diagram below. In this diagram, the following formatting conventions apply:

- HKDF-Extract is drawn as taking the Salt argument from the top and the IKM argument from the left.
- Derive-Secret's Secret argument is indicated by the arrow coming in from the left. For instance, the Early Secret is the Secret for generating the early_traffic-secret.



The general pattern here is that the secrets shown down the left side of the diagram are just raw entropy without context, whereas the secrets down the right side include handshake context and therefore can be used to derive working keys without additional context. Note that the different calls to Derive-Secret may take different Messages arguments, even with the same secret. In a 0-RTT exchange, Derive-Secret is called with four distinct transcripts; in a 1-RTT only exchange with three distinct transcripts.

If a given secret is not available, then the 0-value consisting of a string of Hash.length zeroes is used. Note that this does not mean skipping rounds, so if PSK is not in use Early Secret will still be HKDF-Extract(0, 0).

7.2. Updating Traffic Keys and IVs

Once the handshake is complete, it is possible for either side to update its sending traffic keys using the KeyUpdate handshake message defined in [Section 4.4.3](#). The next generation of traffic keys is computed by generating `traffic_secret_N+1` from `traffic_secret_N` as described in this section then re-deriving the traffic keys as described in [Section 7.3](#).

The next-generation `traffic_secret` is computed as:

```
traffic_secret_N+1 = HKDF-Expand-Label(  
    traffic_secret_N,  
    "application traffic secret", "", Hash.Length)
```

Once `traffic_secret_N+1` and its associated traffic keys have been computed, implementations SHOULD delete `traffic_secret_N`. Once the directional keys are no longer needed, they SHOULD be deleted as well.

7.3. Traffic Key Calculation

The traffic keying material is generated from the following input values:

- A secret value
- A phase value indicating the phase of the protocol the keys are being generated for
- A purpose value indicating the specific value being generated
- The length of the key

The keying material is computed using:

```
key = HKDF-Expand-Label(Secret,  
    phase + ", " + purpose,  
    "",  
    key_length)
```

The following table describes the inputs to the key calculation for each class of traffic keys:

Record Type	Secret	Phase
0-RTT Handshake	early_traffic_secret	"early handshake key expansion"
0-RTT Application	early_traffic_secret	"early application data key expansion"
Handshake	handshake_traffic_secret	"handshake key expansion"
Application Data	traffic_secret_N	"application data key expansion"

The following table indicates the purpose values for each type of key:

Key Type	Purpose
client_write_key	"client write key"
server_write_key	"server write key"
client_write_iv	"client write iv"
server_write_iv	"server write iv"

All the traffic keying material is recomputed whenever the underlying Secret changes (e.g., when changing from the handshake to application data keys or upon a key update).

7.3.1. Diffie-Hellman

A conventional Diffie-Hellman computation is performed. The negotiated key (Z) is converted to byte string by encoding in big-endian, padded with zeros up to the size of the prime. This byte string is used as the shared secret, and is used in the key schedule as specified above.

Note that this construction differs from previous versions of TLS which remove leading zeros.

7.3.2. Elliptic Curve Diffie-Hellman

For secp256r1, secp384r1 and secp521r1, ECDH calculations (including parameter and key generation as well as the shared secret calculation) are performed according to [\[IEEE1363\]](#) using the ECKAS-DH1 scheme with the identity map as key derivation function (KDF), so that the shared secret is the x-coordinate of the ECDH shared secret elliptic curve point represented as an octet string. Note that this octet string (Z in IEEE 1363 terminology) as output by FE2OSP, the Field Element to Octet String Conversion Primitive, has constant length for any given field; leading zeros found in this octet string MUST NOT be truncated.

(Note that this use of the identity KDF is a technicality. The complete picture is that ECDH is employed with a non-trivial KDF because TLS does not directly use this secret for anything other than for computing other secrets.)

ECDH functions are used as follows:

- The public key to put into the KeyShareEntry.key_exchange structure is the result of applying the ECDH function to the secret key of appropriate length (into scalar input) and the standard public basepoint (into u-coordinate point input).
- The ECDH shared secret is the result of applying ECDH function to the secret key (into scalar input) and the peer's public key (into u-coordinate point input). The output is used raw, with no processing.

For X25519 and X448, see [\[RFC7748\]](#).

7.3.3. Exporters

[RFC5705] defines keying material exporters for TLS in terms of the TLS PRF. This document replaces the PRF with HKDF, thus requiring a new construction. The exporter interface remains the same, however the value is computed as:

```
HKDF-Expand-Label(exporter_secret,  
                  label, context_value, key_length)
```

8. Compliance Requirements

8.1. MTI Cipher Suites

In the absence of an application profile standard specifying otherwise, a TLS-compliant application MUST implement the following cipher suites:

```
TLS_ECDHE_ECDSA_WITH_AES_128_GCM_SHA256
TLS_ECDHE_RSA_WITH_AES_128_GCM_SHA256
```

These cipher suites MUST support both digital signatures and key exchange with secp256r1 (NIST P-256) and SHOULD support key exchange with X25519 [[RFC7748](#)].

A TLS-compliant application SHOULD implement the following cipher suites:

```
TLS_ECDHE_ECDSA_WITH_AES_256_GCM_SHA384
TLS_ECDHE_ECDSA_WITH_CHACHA20_POLY1305_SHA256
TLS_ECDHE_RSA_WITH_AES_256_GCM_SHA384
TLS_ECDHE_RSA_WITH_CHACHA20_POLY1305_SHA256
```

8.2. MTI Extensions

In the absence of an application profile standard specifying otherwise, a TLS-compliant application MUST implement the following TLS extensions:

- Signature Algorithms ("signature_algorithms"; [Section 4.2.2](#))
- Negotiated Groups ("supported_groups"; [Section 4.2.3](#))
- Key Share ("key_share"; [Section 4.2.4](#))
- Pre-Shared Key ("pre_shared_key"; [Section 4.2.5](#))
- Server Name Indication ("server_name"; [Section 3 of \[RFC6066\]](#))
- Cookie ("cookie"; [Section 4.2.1](#))

All implementations MUST send and use these extensions when offering applicable cipher suites:

- "signature_algorithms" is REQUIRED for certificate authenticated cipher suites.
- "supported_groups" and "key_share" are REQUIRED for DHE or ECDHE cipher suites.

- "pre_shared_key" is REQUIRED for PSK cipher suites.
- "cookie" is REQUIRED for all cipher suites.

When negotiating use of applicable cipher suites, endpoints MUST abort the connection with a "missing_extension" alert if the required extension was not provided. Any endpoint that receives any invalid combination of cipher suites and extensions MAY abort the connection with a "missing_extension" alert, regardless of negotiated parameters.

Additionally, all implementations MUST support use of the "server_name" extension with applications capable of using it. Servers MAY require clients to send a valid "server_name" extension. Servers requiring this extension SHOULD respond to a ClientHello lacking a "server_name" extension with a fatal "missing_extension" alert.

Servers MUST NOT send the "signature_algorithms" extension; if a client receives this extension it MUST respond with a fatal "unsupported_extension" alert and close the connection.

9. Security Considerations

Security issues are discussed throughout this memo, especially in Appendices B, C, and D.

10. IANA Considerations

This document uses several registries that were originally created in [\[RFC4346\]](#). IANA has updated these to reference this document. The registries and their allocation policies are below:

- TLS Cipher Suite Registry: Values with the first byte in the range 0-254 (decimal) are assigned via Specification Required [\[RFC2434\]](#). Values with the first byte 255 (decimal) are reserved for Private Use [\[RFC2434\]](#). IANA [SHALL add/has added] a "Recommended" column to the cipher suite registry. All cipher suites listed in [Appendix A.4](#) are marked as "Yes". All other cipher suites are marked as "No". IANA [SHALL add/has added] add a note to this column reading:

Cipher suites marked as "Yes" are those allocated via Standards Track RFCs. Cipher suites marked as "No" are not; cipher suites marked "No" range from "good" to "bad" from a cryptographic standpoint.

- TLS ContentType Registry: Future values are allocated via Standards Action [[RFC2434](#)].
- TLS Alert Registry: Future values are allocated via Standards Action [[RFC2434](#)].
- TLS HandshakeType Registry: Future values are allocated via Standards Action [[RFC2434](#)]. IANA [SHALL update/has updated] this registry to rename item 4 from "NewSessionTicket" to "new_session_ticket".

This document also uses a registry originally created in [[RFC4366](#)]. IANA has updated it to reference this document. The registry and its allocation policy is listed below:

- TLS ExtensionType Registry: Values with the first byte in the range 0-254 (decimal) are assigned via Specification Required [[RFC2434](#)]. Values with the first byte 255 (decimal) are reserved for Private Use [[RFC2434](#)]. IANA [SHALL update/has updated] this registry to include the "key_share", "pre_shared_key", and "early_data" extensions as defined in this document.

IANA [shall update/has updated] this registry to include a "TLS 1.3" column with the following four values: "Client", indicating that the server shall not send them. "Clear", indicating that they shall be in the ServerHello. "Encrypted", indicating that they shall be in the EncryptedExtensions block, and "No" indicating that they are not used in TLS 1.3. This column [shall be/has been] initially populated with the values in this document. IANA [shall update/has updated] this registry to add a "Recommended" column. IANA [shall/has] initially populated this column with the values in the table below. This table has been generated by marking Standards Track RFCs as "Yes" and all others as "No".

Extension	Recommend ed	TLS 1.3
server_name [RFC6066]	Yes	Encrypted
max_fragment_length [RFC6066]	Yes	Encrypted
client_certificate_url [RFC6066]	Yes	Encrypted
trusted_ca_keys [RFC6066]	Yes	Encrypted

truncated_hmac [RFC6066]	Yes	No
status_request [RFC6066]	Yes	No
user_mapping [RFC4681]	Yes	Encrypted
client_authz [RFC5878]	No	Encrypted
server_authz [RFC5878]	No	Encrypted
cert_type [RFC6091]	Yes	Encrypted
supported_groups [RFC-ietf-tls-negotiated-ff-dhe]	Yes	Encrypted
ec_point_formats [RFC4492]	Yes	No
srp [RFC5054]	No	No
signature_algorithms [RFC5246]	Yes	Client
use_srtp [RFC5764]	Yes	Encrypted
heartbeat [RFC6520]	Yes	Encrypted
application_layer_protocol_negotiation [RFC7301]	Yes	Encrypted
status_request_v2 [RFC6961]	Yes	Encrypted
signed_certificate_timestamp [RFC6962]	No	Encrypted
client_certificate_type [RFC7250]	Yes	Encrypted
server_certificate_type [RFC7250]	Yes	Encrypted
padding [RFC7685]	Yes	Client
encrypt_then_mac [RFC7366]	Yes	No
extended_master_secret [RFC7627]	Yes	No
SessionTicket TLS [RFC4507]	Yes	No

renegotiation_info [RFC5746]	Yes	No
key_share [[this document]]	Yes	Clear
pre_shared_key [[this document]]	Yes	Clear
early_data [[this document]]	Yes	Encrypted
cookie [[this document]]	Yes	Encrypted/HelloRetryRequest

In addition, this document defines two new registries to be maintained by IANA

- TLS SignatureScheme Registry: Values with the first byte in the range 0-254 (decimal) are assigned via Specification Required [[RFC2434](#)]. Values with the first byte 255 (decimal) are reserved for Private Use [[RFC2434](#)]. This registry SHALL have a "Recommended" column. The registry [shall be/ has been] initially populated with the values described in [Section 4.2.2](#). The following values SHALL be marked as "Recommended":
ecdsa_secp256r1_sha256, ecdsa_secp384r1_sha384, rsa_pss_sha256, rsa_pss_sha384, rsa_pss_sha512, ed25519.

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11.3. URIs

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[Appendix A.](#) Protocol Data Structures and Constant Values

This section describes protocol types and constants. Values listed as `_RESERVED` were used in previous versions of TLS and are listed here for completeness. TLS 1.3 implementations **MUST NOT** send them but might receive them from older TLS implementations.

[A.1.](#) Record Layer

```
enum {
    invalid_RESERVED(0),
    change_cipher_spec_RESERVED(20),
    alert(21),
    handshake(22),
    application_data(23)
    (255)
} ContentType;

struct {
    ContentType type;
    ProtocolVersion record_version = { 3, 1 };    /* TLS v1.x */
    uint16 length;
    opaque fragment[TLSPplaintext.length];
} TLSPplaintext;

struct {
    opaque content[TLSPplaintext.length];
    ContentType type;
    uint8 zeros[length_of_padding];
} TLSInnerPlaintext;

struct {
    ContentType opaque_type = application_data(23); /* see fragment.type */
    ProtocolVersion record_version = { 3, 1 };    /* TLS v1.x */
    uint16 length;
    opaque encrypted_record[length];
} TLSCiphertext;
```

[A.2.](#) Alert Messages


```
enum { warning(1), fatal(2), (255) } AlertLevel;
```

```
enum {  
    close_notify(0),  
    end_of_early_data(1),  
    unexpected_message(10),  
    bad_record_mac(20),  
    decryption_failed_RESERVED(21),  
    record_overflow(22),  
    decompression_failure_RESERVED(30),  
    handshake_failure(40),  
    no_certificate_RESERVED(41),  
    bad_certificate(42),  
    unsupported_certificate(43),  
    certificate_revoked(44),  
    certificate_expired(45),  
    certificate_unknown(46),  
    illegal_parameter(47),  
    unknown_ca(48),  
    access_denied(49),  
    decode_error(50),  
    decrypt_error(51),  
    export_restriction_RESERVED(60),  
    protocol_version(70),  
    insufficient_security(71),  
    internal_error(80),  
    inappropriate_fallback(86),  
    user_canceled(90),  
    no_renegotiation_RESERVED(100),  
    missing_extension(109),  
    unsupported_extension(110),  
    certificate_unobtainable(111),  
    unrecognized_name(112),  
    bad_certificate_status_response(113),  
    bad_certificate_hash_value(114),  
    unknown_psk_identity(115),  
    (255)  
} AlertDescription;
```

```
struct {  
    AlertLevel level;  
    AlertDescription description;  
} Alert;
```


[A.3.](#) Handshake Protocol

```
enum {
    hello_request_RESERVED(0),
    client_hello(1),
    server_hello(2),
    new_session_ticket(4),
    hello_retry_request(6),
    encrypted_extensions(8),
    certificate(11),
    server_key_exchange_RESERVED(12),
    certificate_request(13),
    server_hello_done_RESERVED(14),
    certificate_verify(15),
    client_key_exchange_RESERVED(16),
    finished(20),
    key_update(24),
    (255)
} HandshakeType;

struct {
    HandshakeType msg_type;    /* handshake type */
    uint24 length;            /* bytes in message */
    select (HandshakeType) {
        case client_hello:      ClientHello;
        case server_hello:      ServerHello;
        case hello_retry_request: HelloRetryRequest;
        case encrypted_extensions: EncryptedExtensions;
        case certificate_request: CertificateRequest;
        case certificate:        Certificate;
        case certificate_verify: CertificateVerify;
        case finished:           Finished;
        case new_session_ticket: NewSessionTicket;
        case key_update:         KeyUpdate;
    } body;
} Handshake;
```

[A.3.1.](#) Key Exchange Messages

```
struct {
    uint8 major;
    uint8 minor;
} ProtocolVersion;

struct {
    opaque random_bytes[32];
} Random;
```



```
uint8 CipherSuite[2];    /* Cryptographic suite selector */

struct {
    ProtocolVersion client_version = { 3, 4 };    /* TLS v1.3 */
    Random random;
    opaque legacy_session_id<0..32>;
    CipherSuite cipher_suites<2..2^16-2>;
    opaque legacy_compression_methods<1..2^8-1>;
    Extension extensions<0..2^16-1>;
} ClientHello;

struct {
    ProtocolVersion server_version;
    Random random;
    CipherSuite cipher_suite;
    Extension extensions<0..2^16-1>;
} ServerHello;

struct {
    ProtocolVersion server_version;
    CipherSuite cipher_suite;
    NamedGroup selected_group;
    Extension extensions<0..2^16-1>;
} HelloRetryRequest;

struct {
    ExtensionType extension_type;
    opaque extension_data<0..2^16-1>;
} Extension;

enum {
    supported_groups(10),
    signature_algorithms(13),
    key_share(40),
    pre_shared_key(41),
    early_data(42),
    cookie(44),
    (65535)
} ExtensionType;

struct {
    NamedGroup group;
    opaque key_exchange<1..2^16-1>;
} KeyShareEntry;

struct {
    select (role) {
        case client:
```



```
        KeyShareEntry client_shares<0..2^16-1>;

        case server:
            KeyShareEntry server_share;
    }
} KeyShare;

opaque psk_identity<0..2^16-1>;

struct {
    select (Role) {
        case client:
            psk_identity identities<2..2^16-1>;

        case server:
            uint16 selected_identity;
    }
} PreSharedKeyExtension;

struct {
    select (Role) {
        case client:
            uint32 obfuscated_ticket_age;

        case server:
            struct {};
    }
} EarlyDataIndication;
```

[A.3.1.1.](#) Cookie Extension

```
struct {
    opaque cookie<0..2^16-1>;
} Cookie;
```

[A.3.1.2.](#) Signature Algorithm Extension


```
enum {
    /* RSASSA-PKCS1-v1_5 algorithms */
    rsa_pkcs1_sha1 (0x0201),
    rsa_pkcs1_sha256 (0x0401),
    rsa_pkcs1_sha384 (0x0501),
    rsa_pkcs1_sha512 (0x0601),

    /* ECDSA algorithms */
    ecdsa_secp256r1_sha256 (0x0403),
    ecdsa_secp384r1_sha384 (0x0503),
    ecdsa_secp521r1_sha512 (0x0603),

    /* RSASSA-PSS algorithms */
    rsa_pss_sha256 (0x0700),
    rsa_pss_sha384 (0x0701),
    rsa_pss_sha512 (0x0702),

    /* EdDSA algorithms */
    ed25519 (0x0703),
    ed448 (0x0704),

    /* Reserved Code Points */
    dsa_sha1_RESERVED (0x0202),
    dsa_sha256_RESERVED (0x0402),
    dsa_sha384_RESERVED (0x0502),
    dsa_sha512_RESERVED (0x0602),
    ecdsa_sha1_RESERVED (0x0203),
    obsolete_RESERVED (0x0000..0x0200),
    obsolete_RESERVED (0x0204..0x0400),
    obsolete_RESERVED (0x0404..0x0500),
    obsolete_RESERVED (0x0504..0x0600),
    obsolete_RESERVED (0x0604..0x06FF),
    private_use (0xFE00..0xFFFF),
    (0xFFFF)
} SignatureScheme;

SignatureScheme supported_signature_algorithms<2..2^16-2>;
```

[A.3.1.3.](#) Named Group Extension


```
enum {
    /* Elliptic Curve Groups (ECDHE) */
    obsolete_RESERVED (1..22),
    secp256r1 (23), secp384r1 (24), secp521r1 (25),
    obsolete_RESERVED (26..28),
    x25519 (29), x448 (30),

    /* Finite Field Groups (DHE) */
    ffdhe2048 (256), ffdhe3072 (257), ffdhe4096 (258),
    ffdhe6144 (259), ffdhe8192 (260),

    /* Reserved Code Points */
    ffdhe_private_use (0x01FC..0x01FF),
    ecdhe_private_use (0xFE00..0xFEFF),
    obsolete_RESERVED (0xFF01..0xFF02),
    (0xFFFF)
} NamedGroup;

struct {
    NamedGroup named_group_list<1..2^16-1>;
} NamedGroupList;
```

Values within "obsolete_RESERVED" ranges were used in previous versions of TLS and MUST NOT be offered or negotiated by TLS 1.3 implementations. The obsolete curves have various known/theoretical weaknesses or have had very little usage, in some cases only due to unintentional server configuration issues. They are no longer considered appropriate for general use and should be assumed to be potentially unsafe. The set of curves specified here is sufficient for interoperability with all currently deployed and properly configured TLS implementations.

A.3.1.4. Deprecated Extensions

The following extensions are no longer applicable to TLS 1.3, although TLS 1.3 clients MAY send them if they are willing to negotiate them with prior versions of TLS. TLS 1.3 servers MUST ignore these extensions if they are negotiating TLS 1.3: truncated_hmac [[RFC6066](#)], srp [[RFC5054](#)], encrypt_then_mac [[RFC7366](#)], extended_master_secret [[RFC7627](#)], SessionTicket [[RFC5077](#)], and renegotiation_info [[RFC5746](#)].

A.3.2. Server Parameters Messages


```
struct {
    Extension extensions<0..2^16-1>;
} EncryptedExtensions;

opaque DistinguishedName<1..2^16-1>;

struct {
    opaque certificate_extension_oid<1..2^8-1>;
    opaque certificate_extension_values<0..2^16-1>;
} CertificateExtension;

struct {
    opaque certificate_request_context<0..2^8-1>;
    SignatureScheme
        supported_signature_algorithms<2..2^16-2>;
    DistinguishedName certificate_authorities<0..2^16-1>;
    CertificateExtension certificate_extensions<0..2^16-1>;
} CertificateRequest;
```

[A.3.3.](#) Authentication Messages

```
opaque ASN1Cert<1..2^24-1>;

struct {
    opaque certificate_request_context<0..2^8-1>;
    ASN1Cert certificate_list<0..2^24-1>;
} Certificate;

struct {
    SignatureScheme algorithm;
    opaque signature<0..2^16-1>;
} CertificateVerify;

struct {
    opaque verify_data[Hash.length];
} Finished;
```

[A.3.4.](#) Ticket Establishment


```

enum { (65535) } TicketExtensionType;

struct {
    TicketExtensionType extension_type;
    opaque extension_data<1..2^16-1>;
} TicketExtension;

enum {
    allow_early_data(1),
    allow_dhe_resumption(2),
    allow_psk_resumption(4)
} TicketFlags;

struct {
    uint32 ticket_lifetime;
    uint32 flags;
    uint32 ticket_age_add;
    TicketExtension extensions<2..2^16-2>;
    opaque ticket<0..2^16-1>;
} NewSessionTicket;

```

[A.4.](#) Cipher Suites

A cipher suite defines a cipher specification supported in TLS and negotiated via hello messages in the TLS handshake. Cipher suite names follow a general naming convention composed of a series of component algorithm names separated by underscores:

```
CipherSuite TLS_KEA_AUTH_WITH_CIPHER_HASH = VALUE;
```

Component	Contents
TLS	The string "TLS"
KEA	The key exchange algorithm (e.g. ECDHE, DHE)
AUTH	The authentication algorithm (e.g. certificates, PSK)
WITH	The string "WITH"
CIPHER	The symmetric cipher used for record protection
HASH	The hash algorithm used with HKDF
VALUE	The two byte ID assigned for this cipher suite

The "CIPHER" component commonly has sub-components used to designate the cipher name, bits, and mode, if applicable. For example, "AES_256_GCM" represents 256-bit AES in the GCM mode of operation. Cipher suite names that lack a "HASH" value that are defined for use with TLS 1.2 or later use the SHA-256 hash algorithm by default.

The primary key exchange algorithm used in TLS is Ephemeral Diffie-Hellman [[DH](#)]. The finite field based version is denoted "DHE" and the elliptic curve based version is denoted "ECDHE". Prior versions of TLS supported non-ephemeral key exchanges, however these are not supported by TLS 1.3.

See the definitions of each cipher suite in its specification document for the full details of each combination of algorithms that is specified.

The following is a list of standards track server-authenticated (and optionally client-authenticated) cipher suites which are currently available in TLS 1.3:

Cipher Suite Name	Value	Specification
TLS_DHE_RSA_WITH_AES_128_GCM_SHA256	{0x00,0x9E}	[RFC5288]
TLS_DHE_RSA_WITH_AES_256_GCM_SHA384	{0x00,0x9F}	[RFC5288]
TLS_ECDHE_ECDSA_WITH_AES_128_GCM_SHA256	{0xC0,0x2B}	[RFC5289]
TLS_ECDHE_ECDSA_WITH_AES_256_GCM_SHA384	{0xC0,0x2C}	[RFC5289]
TLS_ECDHE_RSA_WITH_AES_128_GCM_SHA256	{0xC0,0x2F}	[RFC5289]
TLS_ECDHE_RSA_WITH_AES_256_GCM_SHA384	{0xC0,0x30}	[RFC5289]
TLS_DHE_RSA_WITH_AES_128_CCM	{0xC0,0x9E}	[RFC6655]
TLS_DHE_RSA_WITH_AES_256_CCM	{0xC0,0x9F}	[RFC6655]
TLS_DHE_RSA_WITH_AES_128_CCM_8	{0xC0,0xA2}	[RFC6655]
TLS_DHE_RSA_WITH_AES_256_CCM_8	{0xC0,0xA3}	[RFC6655]
TLS_ECDHE_RSA_WITH_CHACHA20_POLY1305_SHA256	{0xCC,0xA8}	[RFC7905]
TLS_ECDHE_ECDSA_WITH_CHACHA20_POLY1305_SHA256	{0xCC,0xA9}	[RFC7905]
TLS_DHE_RSA_WITH_CHACHA20_POLY1305_SHA256	{0xCC,0xAA}	[RFC7905]

Note: The values listed for ChaCha/Poly are preliminary but are being or will be used for interop testing and therefore are likely to be assigned.

Note: ECDHE AES GCM was not yet standards track prior to the publication of this specification. This document promotes the above-listed ciphers to standards track.

The following is a list of standards track ephemeral pre-shared key cipher suites which are currently available in TLS 1.3:

Cipher Suite Name	Value	Specification
TLS_DHE_PSK_WITH_AES_128_GCM_SHA256	{0x00, 0xAA}	[RFC5487]
TLS_DHE_PSK_WITH_AES_256_GCM_SHA384	{0x00, 0xAB}	[RFC5487]
TLS_DHE_PSK_WITH_AES_128_CCM	{0xC0, 0xA6}	[RFC6655]
TLS_DHE_PSK_WITH_AES_256_CCM	{0xC0, 0xA7}	[RFC6655]
TLS_PSK_DHE_WITH_AES_128_CCM_8	{0xC0, 0xAA}	[RFC6655]
TLS_PSK_DHE_WITH_AES_256_CCM_8	{0xC0, 0xAB}	[RFC6655]
TLS_ECDHE_PSK_WITH_AES_128_GCM_SHA256	{0xD0, 0x01}	[I-D.mattsson-tls-ecdhe-psk-aead]
TLS_ECDHE_PSK_WITH_AES_256_GCM_SHA384	{0xD0, 0x02}	[I-D.mattsson-tls-ecdhe-psk-aead]
TLS_ECDHE_PSK_WITH_AES_128_CCM_8_SHA256	{0xD0, 0x03}	[I-D.mattsson-tls-ecdhe-psk-aead]
TLS_ECDHE_PSK_WITH_AES_128_CCM_SHA256	{0xD0, 0x04}	[I-D.mattsson-tls-ecdhe-psk-aead]
TLS_ECDHE_PSK_WITH_AES_256_CCM_SHA384	{0xD0, 0x05}	[I-D.mattsson-tls-ecdhe-psk-aead]
TLS_ECDHE_PSK_WITH_CHACHA20_POLY1305_SHA256	{0xCC, 0xAC}	[RFC7905]
TLS_DHE_PSK_WITH_CHACHA20_POLY1305_SHA256	{0xCC, 0xAD}	[RFC7905]

Note: The values listed for ECDHE and ChaCha/Poly are preliminary but are being or will be used for interop testing and therefore are likely to be assigned.

Note: [\[RFC6655\]](#) is inconsistent with respect to the ordering of components within PSK AES CCM cipher suite names. The names above are as defined.

All cipher suites in this section are specified for use with both TLS 1.2 and TLS 1.3, as well as the corresponding versions of DTLS. (see [Appendix C](#))

New cipher suite values are assigned by IANA as described in [Section 10](#).

[A.4.1](#). Unauthenticated Operation

Previous versions of TLS offered explicitly unauthenticated cipher suites based on anonymous Diffie-Hellman. These cipher suites have been deprecated in TLS 1.3. However, it is still possible to negotiate cipher suites that do not provide verifiable server authentication by several methods, including:

- Raw public keys [\[RFC7250\]](#).
- Using a public key contained in a certificate but without validation of the certificate chain or any of its contents.

Either technique used alone is are vulnerable to man-in-the-middle attacks and therefore unsafe for general use. However, it is also possible to bind such connections to an external authentication mechanism via out-of-band validation of the server's public key, trust on first use, or channel bindings [\[RFC5929\]](#). [[NOTE: TLS 1.3 needs a new channel binding definition that has not yet been defined.]] If no such mechanism is used, then the connection has no protection against active man-in-the-middle attack; applications **MUST NOT** use TLS in such a way absent explicit configuration or a specific application profile.

[Appendix B](#). Implementation Notes

The TLS protocol cannot prevent many common security mistakes. This section provides several recommendations to assist implementors.

[B.1](#). Random Number Generation and Seeding

TLS requires a cryptographically secure pseudorandom number generator (PRNG). Care must be taken in designing and seeding PRNGs. PRNGs based on secure hash operations, most notably SHA-256, are acceptable, but cannot provide more security than the size of the random number generator state.

To estimate the amount of seed material being produced, add the number of bits of unpredictable information in each seed byte. For example, keystroke timing values taken from a PC compatible 18.2 Hz timer provide 1 or 2 secure bits each, even though the total size of the counter value is 16 bits or more. Seeding a 128-bit PRNG would thus require approximately 100 such timer values.

[RFC4086] provides guidance on the generation of random values.

[B.2.](#) Certificates and Authentication

Implementations are responsible for verifying the integrity of certificates and should generally support certificate revocation messages. Certificates should always be verified to ensure proper signing by a trusted Certificate Authority (CA). The selection and addition of trusted CAs should be done very carefully. Users should be able to view information about the certificate and root CA.

[B.3.](#) Cipher Suite Support

TLS supports a range of key sizes and security levels, including some that provide no or minimal security. A proper implementation will probably not support many cipher suites. Applications SHOULD also enforce minimum and maximum key sizes. For example, certification paths containing keys or signatures weaker than 2048-bit RSA or 224-bit ECDSA are not appropriate for secure applications. See also [Appendix C.4](#).

[B.4.](#) Implementation Pitfalls

Implementation experience has shown that certain parts of earlier TLS specifications are not easy to understand, and have been a source of interoperability and security problems. Many of these areas have been clarified in this document, but this appendix contains a short list of the most important things that require special attention from implementors.

TLS protocol issues:

- Do you correctly handle handshake messages that are fragmented to multiple TLS records (see [Section 5.1](#))? Including corner cases like a ClientHello that is split to several small fragments? Do you fragment handshake messages that exceed the maximum fragment size? In particular, the certificate and certificate request handshake messages can be large enough to require fragmentation.
- Do you ignore the TLS record layer version number in all TLS records? (see [Appendix C](#))

- Have you ensured that all support for SSL, RC4, EXPORT ciphers, and MD5 (via the "signature_algorithm" extension) is completely removed from all possible configurations that support TLS 1.3 or later, and that attempts to use these obsolete capabilities fail correctly? (see [Appendix C](#))
- Do you handle TLS extensions in ClientHello correctly, including unknown extensions or omitting the extensions field completely?
- When the server has requested a client certificate, but no suitable certificate is available, do you correctly send an empty Certificate message, instead of omitting the whole message (see [Section 4.3.1.2](#))?
- When processing the plaintext fragment produced by AEAD-Decrypt and scanning from the end for the ContentType, do you avoid scanning past the start of the cleartext in the event that the peer has sent a malformed plaintext of all-zeros?
- When processing a ClientHello containing a version of { 3, 5 } or higher, do you respond with the highest common version of TLS rather than requiring an exact match? Have you ensured this continues to be true with arbitrarily higher version numbers? (e.g. { 4, 0 }, { 9, 9 }, { 255, 255 })
- Do you properly ignore unrecognized cipher suites ([Section 4.1.1](#)), hello extensions ([Section 4.2](#)), named groups ([Section 4.2.3](#)), and signature algorithms ([Section 4.2.2](#))?

Cryptographic details:

- What countermeasures do you use to prevent timing attacks against RSA signing operations [[TIMING](#)]?
- When verifying RSA signatures, do you accept both NULL and missing parameters? Do you verify that the RSA padding doesn't have additional data after the hash value? [[FI06](#)]
- When using Diffie-Hellman key exchange, do you correctly preserve leading zero bytes in the negotiated key (see [Section 7.3.1](#))?
- Does your TLS client check that the Diffie-Hellman parameters sent by the server are acceptable, (see [Section 4.2.4.1](#))?
- Do you use a strong and, most importantly, properly seeded random number generator (see [Appendix B.1](#)) when generating Diffie-Hellman private values, the ECDSA "k" parameter, and other security-

critical values? It is RECOMMENDED that implementations implement "deterministic ECDSA" as specified in [[RFC6979](#)].

- Do you zero-pad Diffie-Hellman public key values to the group size (see [Section 4.2.4.1](#))?

[B.5.](#) Client Tracking Prevention

Clients SHOULD NOT reuse a session ticket for multiple connections. Reuse of a session ticket allows passive observers to correlate different connections. Servers that issue session tickets SHOULD offer at least as many session tickets as the number of connections that a client might use; for example, a web browser using HTTP/1.1 [[RFC7230](#)] might open six connections to a server. Servers SHOULD issue new session tickets with every connection. This ensures that clients are always able to use a new session ticket when creating a new connection.

[Appendix C.](#) Backward Compatibility

The TLS protocol provides a built-in mechanism for version negotiation between endpoints potentially supporting different versions of TLS.

TLS 1.x and SSL 3.0 use compatible ClientHello messages. Servers can also handle clients trying to use future versions of TLS as long as the ClientHello format remains compatible and the client supports the highest protocol version available in the server.

Prior versions of TLS used the record layer version number for various purposes. (TLSPlaintext.record_version & TLSCiphertext.record_version) As of TLS 1.3, this field is deprecated and its value MUST be ignored by all implementations. Version negotiation is performed using only the handshake versions. (ClientHello.client_version & ServerHello.server_version) In order to maximize interoperability with older endpoints, implementations that negotiate the use of TLS 1.0-1.2 SHOULD set the record layer version number to the negotiated version for the ServerHello and all records thereafter.

For maximum compatibility with previously non-standard behavior and misconfigured deployments, all implementations SHOULD support validation of certification paths based on the expectations in this document, even when handling prior TLS versions' handshakes. (see [Section 4.3.1.1](#))

TLS 1.2 and prior supported an "Extended Master Secret" [[RFC7627](#)] extension which digested large parts of the handshake transcript into

the master secret. Because TLS 1.3 always hashes in the transcript up to the server CertificateVerify, implementations which support both TLS 1.3 and earlier versions SHOULD indicate the use of the Extended Master Secret extension in their APIs whenever TLS 1.3 is used.

C.1. Negotiating with an older server

A TLS 1.3 client who wishes to negotiate with such older servers will send a normal TLS 1.3 ClientHello containing { 3, 4 } (TLS 1.3) in ClientHello.client_version. If the server does not support this version it will respond with a ServerHello containing an older version number. If the client agrees to use this version, the negotiation will proceed as appropriate for the negotiated protocol. A client resuming a session SHOULD initiate the connection using the version that was previously negotiated.

Note that 0-RTT data is not compatible with older servers. See [Appendix C.3](#).

If the version chosen by the server is not supported by the client (or not acceptable), the client MUST send a "protocol_version" alert message and close the connection.

If a TLS server receives a ClientHello containing a version number greater than the highest version supported by the server, it MUST reply according to the highest version supported by the server.

Some legacy server implementations are known to not implement the TLS specification properly and might abort connections upon encountering TLS extensions or versions which it is not aware of.

Interoperability with buggy servers is a complex topic beyond the scope of this document. Multiple connection attempts may be required in order to negotiate a backwards compatible connection, however this practice is vulnerable to downgrade attacks and is NOT RECOMMENDED.

C.2. Negotiating with an older client

A TLS server can also receive a ClientHello containing a version number smaller than the highest supported version. If the server wishes to negotiate with old clients, it will proceed as appropriate for the highest version supported by the server that is not greater than ClientHello.client_version. For example, if the server supports TLS 1.0, 1.1, and 1.2, and client_version is TLS 1.0, the server will proceed with a TLS 1.0 ServerHello. If the server only supports versions greater than client_version, it MUST send a "protocol_version" alert message and close the connection.

Note that earlier versions of TLS did not clearly specify the record layer version number value in all cases (TLSPlaintext.record_version). Servers will receive various TLS 1.x versions in this field, however its value MUST always be ignored.

C.3. Zero-RTT backwards compatibility

0-RTT data is not compatible with older servers. An older server will respond to the ClientHello with an older ServerHello, but it will not correctly skip the 0-RTT data and fail to complete the handshake. This can cause issues when a client attempts to use 0-RTT, particularly against multi-server deployments. For example, a deployment could deploy TLS 1.3 gradually with some servers implementing TLS 1.3 and some implementing TLS 1.2, or a TLS 1.3 deployment could be downgraded to TLS 1.2.

A client that attempts to send 0-RTT data MUST fail a connection if it receives a ServerHello with TLS 1.2 or older. A client that attempts to repair this error SHOULD NOT send a TLS 1.2 ClientHello, but instead send a TLS 1.3 ClientHello without 0-RTT data.

To avoid this error condition, multi-server deployments SHOULD ensure a uniform and stable deployment of TLS 1.3 without 0-RTT prior to enabling 0-RTT.

C.4. Backwards Compatibility Security Restrictions

If an implementation negotiates use of TLS 1.2, then negotiation of cipher suites also supported by TLS 1.3 SHOULD be preferred, if available.

The security of RC4 cipher suites is considered insufficient for the reasons cited in [RFC7465]. Implementations MUST NOT offer or negotiate RC4 cipher suites for any version of TLS for any reason.

Old versions of TLS permitted the use of very low strength ciphers. Ciphers with a strength less than 112 bits MUST NOT be offered or negotiated for any version of TLS for any reason.

The security of SSL 2.0 [SSL2] is considered insufficient for the reasons enumerated in [RFC6176], and MUST NOT be negotiated for any reason.

Implementations MUST NOT send an SSL version 2.0 compatible CLIENT-HELLO. Implementations MUST NOT negotiate TLS 1.3 or later using an SSL version 2.0 compatible CLIENT-HELLO. Implementations are NOT RECOMMENDED to accept an SSL version 2.0 compatible CLIENT-HELLO in order to negotiate older versions of TLS.

Implementations MUST NOT send or accept any records with a version less than { 3, 0 }.

The security of SSL 3.0 [[SSL3](#)] is considered insufficient for the reasons enumerated in [[RFC7568](#)], and MUST NOT be negotiated for any reason.

Implementations MUST NOT send a ClientHello.client_version or ServerHello.server_version set to { 3, 0 } or less. Any endpoint receiving a Hello message with ClientHello.client_version or ServerHello.server_version set to { 3, 0 } MUST respond with a "protocol_version" alert message and close the connection.

Implementations MUST NOT use the Truncated HMAC extension, defined in [Section 7 of \[RFC6066\]](#), as it is not applicable to AEAD ciphers and has been shown to be insecure in some scenarios.

[Appendix D](#). Overview of Security Properties

[[TODO: This section is still a WIP and needs a bunch more work.]]

A complete security analysis of TLS is outside the scope of this document. In this section, we provide an informal description the desired properties as well as references to more detailed work in the research literature which provides more formal definitions.

We cover properties of the handshake separately from those of the record layer.

[D.1](#). Handshake

The TLS handshake is an Authenticated Key Exchange (AKE) protocol which is intended to provide both one-way authenticated (server-only) and mutually authenticated (client and server) functionality. At the completion of the handshake, each side outputs its view on the following values:

- A "session key" (the master secret) from which can be derived a set of working keys.
- A set of cryptographic parameters (algorithms, etc.)
- The identities of the communicating parties.

We assume that the attacker has complete control of the network in between the parties [[RFC3552](#)]. Even under these conditions, the handshake should provide the properties listed below. Note that

these properties are not necessarily independent, but reflect the protocol consumers' needs.

Establishing the same session key. The handshake needs to output the same session key on both sides of the handshake, provided that it completes successfully on each endpoint (See [\[CK01\]](#); defn 1, part 1).

Secrecy of the session key. The shared session key should be known only to the communicating parties, not to the attacker (See [\[CK01\]](#); defn 1, part 2). Note that in a unilaterally authenticated connection, the attacker can establish its own session keys with the server, but those session keys are distinct from those established by the client.

Peer Authentication. The client's view of the peer identity should reflect the server's identity. If the client is authenticated, the server's view of the peer identity should match the client's identity.

Uniqueness of the session key: Any two distinct handshakes should produce distinct, unrelated session keys

Downgrade protection. The cryptographic parameters should be the same on both sides and should be the same as if the peers had been communicating in the absence of an attack (See [\[BBFKZG16\]](#); defns 8 and 9).

Forward secret If the long-term keying material (in this case the signature keys in certificate-based authentication modes or the PSK in PSK-(EC)DHE modes) are compromised after the handshake is complete, this does not compromise the security of the session key (See [\[DOW92\]](#)).

Protection of endpoint identities. The server's identity (certificate) should be protected against passive attackers. The client's identity should be protected against both passive and active attackers.

Informally, the signature-based modes of TLS 1.3 provide for the establishment of a unique, secret, shared, key established by an (EC)DHE key exchange and authenticated by the server's signature over the handshake transcript, as well as tied to the server's identity by a MAC. If the client is authenticated by a certificate, it also signs over the handshake transcript and provides a MAC tied to both identities. [\[SIGMA\]](#) describes the analysis of this type of key exchange protocol. If fresh (EC)DHE keys are used for each connection, then the output keys are forward secret.

The PSK and resumption-PSK modes bootstrap from a long-term shared secret into a unique per-connection short-term session key. This secret may have been established in a previous handshake. If PSK-(EC)DHE modes are used, this session key will also be forward secret. The resumption-PSK mode has been designed so that the resumption master secret computed by connection N and needed to form connection N+1 is separate from the traffic keys used by connection N, thus providing forward secrecy between the connections.

For all handshake modes, the Finished MAC (and where present, the signature), prevents downgrade attacks. In addition, the use of certain bytes in the random nonces as described in [Section 4.1.2](#) allows the detection of downgrade to previous TLS versions.

As soon as the client and the server have exchanged enough information to establish shared keys, the remainder of the handshake is encrypted, thus providing protection against passive attackers. Because the server authenticates before the client, the client can ensure that it only reveals its identity to an authenticated server. Note that implementations must use the provided record padding mechanism during the handshake to avoid leaking information about the identities due to length.

The 0-RTT mode of operation generally provides the same security properties as 1-RTT data, with the two exceptions that the 0-RTT encryption keys do not provide full forward secrecy and that the server is not able to guarantee full uniqueness of the handshake (non-replayability) without keeping potentially undue amounts of state. See [Section 4.2.6](#) for one mechanism to limit the exposure to replay.

The reader should refer to the following references for analysis of the TLS handshake [[CHSV16](#)] [[FGSW16](#)] [[LXZFH16](#)].

[D.2.](#) Record Layer

The record layer depends on the handshake producing a strong session key which can be used to derive bidirectional traffic keys and nonces. Assuming that is true, and the keys are used for no more data than indicated in [Section 5.5](#) then the record layer should provide the following guarantees:

Confidentiality. An attacker should not be able to determine the plaintext contents of a given record.

Integrity. An attacker should not be able to craft a new record which is different from an existing record which will be accepted by the receiver.

Order protection/non-replayability An attacker should not be able to cause the receiver to accept a record which it has already accepted or cause the receiver to accept record N+1 without having first processed record N. [[TODO: If we merge in DTLS to this document, we will need to update this guarantee.]]

Length concealment. Given a record with a given external length, the attacker should not be able to determine the amount of the record that is content versus padding.

Forward security after key change. If the traffic key update mechanism described in [Section 4.4.3](#) has been used and the previous generation key is deleted, an attacker who compromises the endpoint should not be able to decrypt traffic encrypted with the old key.

Informally, TLS 1.3 provides these properties by AEAD-protecting the plaintext with a strong key. AEAD encryption [[RFC5116](#)] provides confidentiality and integrity for the data. Non-replayability is provided by using a separate nonce for each record, with the nonce being derived from the record sequence number ([Section 5.3](#)), with the sequence number being maintained independently at both sides thus records which are delivered out of order result in AEAD deprotection failures.

The plaintext protected by the AEAD function consists of content plus variable-length padding. Because the padding is also encrypted, the attacker cannot directly determine the length of the padding, but may be able to measure it indirectly by the use of timing channels exposed during record processing (i.e., seeing how long it takes to process a record). In general, it is not known how to remove this type of channel because even a constant time padding removal function will then feed the content into data-dependent functions.

Generation N+1 keys are derived from generation N keys via a key derivation function [Section 7.2](#). As long as this function is truly one way, it is not possible to compute the previous keys after a key change (forward secrecy). However, TLS does not provide security for data which is sent after the traffic secret is compromised, even after a key update (backward secrecy); systems which want backward secrecy must do a fresh handshake and establish a new session key with an (EC)DHE exchange.

The reader should refer to the following references for analysis of the TLS record layer.

Appendix E. Working Group Information

The discussion list for the IETF TLS working group is located at the e-mail address tls@ietf.org [1]. Information on the group and information on how to subscribe to the list is at <https://www1.ietf.org/mailman/listinfo/tls>

Archives of the list can be found at: <https://www.ietf.org/mail-archive/web/tls/current/index.html>

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