Y. Li TRILL Working Group D. Eastlake INTERNET-DRAFT Intended Status: Standard Track W. Hao H. Chen Huawei Technologies S. Chatterjee

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TRILL: Data Label based Tree Selection for Multi-destination Data draft-ietf-trill-tree-selection-05

#### Abstract

TRILL uses distribution trees to deliver multi-destination frames. Multiple trees can be used by an ingress RBridge for flows regardless of the VLAN, Fine Grained Label (FGL), and/or multicast group of the flow. Different ingress RBridges may choose different distribution trees for TRILL Data packets in the same VLAN, FGL, and/or multicast group. To avoid unnecessary link utilization, distribution trees should be pruned based on VLAN and/or FGL and/or multicast destination address. If any VLAN, FGL, or multicast group can be sent on any tree, for typical fast path hardware, the amount of pruning information is multiplied by the number of trees, but there is a limited hardware capacity for such pruning information.

This document specifies an optional facility to restrict the TRILL Data packets sent on particular distribution trees by VLAN, FGL, and/or multicast group thus reducing the total amount of pruning information so that it can more easily be accommodated by fast path hardware.

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### 1. Introduction

### 1.1. Background Description

One or more distribution trees, identified by their root nickname, are used to distribute multi-destination data in a TRILL campus [RFC6325]. The RBridge having the highest tree root priority announces the total number of trees that should be computed for the campus. It may also specify the list of trees that RBridges need to compute using the Tree Identifiers (TREE-RT-IDs) sub-TLV [RFC7176]. Every RBridge can specify the trees it will use for multi-destination TRILL data packets it originates in the Trees Used Identifiers (TREE-USE-IDs) sub-TLV and the VLANs or fine grained labels (FGLs [RFC7172]) it is interested in are specified in Interested VLANs and/or Interested Labels sub-TLVs [RFC7176]. It is suggested that, by default, the ingress RBridge uses the distribution tree whose root is the closest [RFC6325]. Trees Used Identifiers sub-TLVs are used to build the RPF (Reverse Path Forwarding) Check table that is used for reverse path forwarding check, Interested VLANs and Interested Labels sub-TLVs are used for distribution tree pruning and the multidestination forwarding table with pruning info is built based on that RPF Check Table. To reduce unnecessary link loads, each distribution tree should be pruned per VLAN/FGL, eliminating branches that have no potential receivers downstream as specified in [RFC6325]. Further pruning based on Layer 2 or Layer 3 multicast address is also possible.

Defaults are provided but it depends on the implementation how many trees are calculated, where the tree roots are located, and which tree(s) are to be used by an ingress RBridge. With the increasing demand to use TRILL in data center networks, there are some features we can explore for multi-destination frames in the data center use case. In order to achieve non-blocking data forwarding, a fat tree structure is often used. Figure 1 shows a typical fat tree structure based data center network. RB1 and RB2 are aggregation switches and RB11 to RB14 are access switches. It is a common practice to configure the tree roots to be at the aggregation switches for efficient traffic transportation. Then all the ingress RBridges that are access switches have the same distance to all the tree roots.

## 1.2. Terminology Used in This Document

This document uses the terminology from [RFC6325] and [RFC7172], some of which is repeated below for convenience, along with some additional terms listed below:

Campus: Name for a TRILL network, like "bridged LAN" is a name for a bridged network. It does not have any academic implication.

Data Label: VLAN or FGL.

ECMP: Equal Cost Multi-Path [RFC6325].

FGL: Fine Grained Label [RFC7172].

IPTV: "Television" (video) over IP.

RBridge: An alternative name for a TRILL switch.

RPF: Reverse Path Forwarding.

TRILL: Transparent Interconnection of Lots of Links (or Tunneled Routing in the Link Layer).

TRILL switch: A device implementing the TRILL protocol. Sometimes called an RBridge.

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in <a href="https://recommended.org/recom

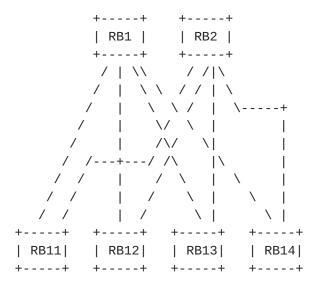


Figure 1. Fat Tree Structure based TRILL network

### 2. Motivations

In the structure of Figure 1, if we choose to put the tree roots at RB1 and RB2, the ingress RBridge (e.g. RB11) would find more than one equal cost closest tree root (i.e. RB1 & RB2). An ingress RBridge has two options to select the tree root for multi-destination frames: choose one and only one as distribution tree root or use ECMP-like algorithm to balance the traffic among the multiple trees whose roots

are at the same distance.

- For the former (one distribution tree root), a single tree used by each ingress RBridge, can have the problem of uneven or inefficient link usage. For example, if RB11 chooses the tree1 that is rooted at RB1 as the distribution tree, the link between RB11 and RB2 will not be used for multi-destination frames ingressed by RB11.
- For the latter (ECMP-Like algorithm), ECMP based tree selection results in a linear increase in multicast forwarding table size with the number of trees as explained in the next paragraph.

A multicast forwarding table at an RBridge is normally used to map the key of (distribution tree nickname + VLAN) to an index to a list of ports for multicast packet replication. The key used for mapping is simply the tree nickname when the RBridge does not prune the tree. The key could be the distribution tree nickname augmented by the Fine Grained Label (FGL) and/or Layer 2 or 3 multicast address when the RBridge supports FGL and/or Layer 2 or 3 pruning information.

For any RBridge RBn, for each VLAN x, if RBn is in a distribution tree t used by traffic in VLAN x, there will be an entry of (t, x, port list) in the multicast forwarding table on RBn. Typically each entry contains a distinct combination of (tree nickname, VLAN) as the lookup key. If there are n such trees and m such VLANs, the multicast forwarding table size on RBn is n\*m entries. If a fine-grained label is used [RFC7172] and/or finer pruning is used (for example, VLAN + multicast group address is used for pruning), the value of m increases. In the larger scale data center, more trees would be necessary for better load balancing purpose and this results in an increased value for n. In either case, the number of table entries n\*m will increase dramatically.

The left hand table in Figure 2 shows an example of the multicast forwarding table on RB11 in the Figure 1 topology with 2 distribution trees in a campus using typical fast path hardware. The number of entries is approximately 2 \* 4K in this case. If 4 distribution trees are used in a TRILL campus and RBn has 4K VLANs with downstream receivers, it consumes 16K table entries. Fast path TRILL multicast forwarding tables typically have a size limited by hardware. The table entries are a precious resource. In some implementations, the table is shared with Layer 3 IP multicast for a total of 16K or 8K table entries. Therefore we want to reduce the table size consumed for TRILL distribution trees as much as possible and at the same time maintain the load balancing among trees.

In cases where blocks of consecutive VLANs or FGLs can be assigned to a tree, the multicast forwarding table could be greatly compressed if

entries could have a Data Label value and mask with the fast path hardware doing the longest prefix matching. But few if any fast path implementations provide such logic.

A straightforward way to alleviate the limited table entries problem is not to prune the distribution tree. However this can only be used in restricted scenarios for the following reasons:

- Not pruning wastes bandwidth for multi-destination packets. There is normally broadcast traffic, like ARP and unknown unicast, that can be pruned on VLAN (or FGL) so it is not sent down branches of a distribution tree where it is not needed. In addition, if there is a lot of Layer 3 multicast traffic, no pruning may result in the worse consequence of that user data unnecessarily flooded all over the campus. The volume could be very large if certain applications like IPTV ("Television" (video) over IP) are supported. More precise pruning, such as pruning based on multicast group, may be desirable in this case.
- Not pruning is only useful at pure transit nodes. Edge nodes always need to maintain the multicast forwarding table with the key of (tree nickname + VLAN (or FGL)) since the edge node needs to decide whether and how to replicate the frame to local access ports. It is likely that edge nodes are relatively low end switches with a smaller shared table size, say 4K, available.
- Security concerns. VLAN (or FGL) based traffic isolation is a basic requirement in some scenarios. No pruning may increase the risk of leakage of the traffic. Misbehaved RBridges may take advantage of this leakage of traffic.

In addition to the multicast table size concern, some silicon does not currently support hashing-based tree nickname selection at the ingress RBridge but commonly uses VLAN based tree selection. If the control plane of the ingress RBridge maps the incoming VLAN x to a tree nickname t. Then the data plane will always use tree t for VLAN x multi-destination frames. Such an ingress RBridge may choose multiple trees to be used for load sharing, it can use one and only one tree for each VLAN. If we make sure all ingress RBridges campuswide send VLAN x multi-destination packets only using tree t, then there would be no need to store the multicast table entry with the key of (tree-other-than-t, x) on any RBridge.

This document describes the TRILL control plane support for distribution tree selection based on VLAN, FGL, and/or multicast address to reduce the multicast forwarding table size. It is compatible with the silicon implementations mentioned in the previous paragraph.

## 3. Data Label based Tree Selection

Data Label (VLAN or FGL) based tree selection can be used as a distribution tree selection mechanism, especially when the multicast forwarding table size is a concern. This section specifies that mechanism and how to extend it so that tree selection can be based on multicast group.

### 3.1. Overview of the Mechanism

The RBridge that has the highest priority to be a tree root announces the tree nicknames and the Data Labels allowed on each tree. Such tree to Data Label correspondence announcements can be based on static configuration or some predefined algorithm beyond the scope of this document. An ingress RBridge selects the tree-VLAN correspondence it wishes to use from the list announced by the highest priority tree root. It SHOULD NOT transmit VLAN x frame on tree y if the highest priority tree root does not say VLAN x is allowed on tree y.

If we make sure a particular VLAN is allowed on one and only one tree, we can keep the number of multicast forwarding table entries on any RBridge fixed at 4K maximum (or up to 16M in case of fine grained label). Take Figure 1 as example, two trees rooted at RB1 and RB2 respectively. The highest priority tree root appoints the tree1 to carry VLAN 1-2000 and tree2 to carry VLAN 2001-4094. With such announcement by the highest priority tree root, every RBridge which understands the announcement will not send VLAN 2001-4094 traffic on tree1 and not send VLAN 1-2000 traffic on tree2. Then no RBridge would need to store the entries for tree1/VLAN2001-4094 or tree2/VLAN1-2000. Figure 2 shows the multicast forwarding table on an RBridge before and after we use VLAN based tree selection. The number of entries is reduced by a factor f, f being the number of trees used in the campus. In this example, it is reduced from 2\*4094 to 4094. This affects both transit nodes and edge nodes. The data plane encoding does not change.

tree nickname	•	+  port list  +	
tree 1	1	+       +	tree 1   1
,	2		
tree 1		 	tree 1
tree 1		 	tree 1   1999
tree 1	•		tree 1   2000
tree 1	•		,
tree 1	4094		tree 2   2002
tree 2	•		tree 2
tree 2	,		
tree 2			
tree 2			+
tree 2			
tree 2		 	
tree 2	4093		
tree 2	4094		

Figure 2. Multicast forwarding table before (left) & after (right)

# 3.2. APPsub-TLVs Supporting Tree Selection

Six new APPsub-TLVs that can be carried in the TRILL GENINFO TLV [RFC7357] in E-L1FS FS-LSPs [rfc7780] are defined below. The first four can be considered analogous to finer granularity versions of the Tree Identifiers Sub-TLV and the Trees Used Identifiers Sub-TLV in [RFC7176]. Two APPsub-TLVs supporting VLAN based tree selection are specified in Sections 3.2.1 and 3.2.2. They are used by the highest priority tree root to announce the allowed VLANs on each tree in the campus and by an ingress RBridge to announce the tree-VLAN correspondence it selects from the list announced by the highest priority tree root. Two APPsub-TLVs supporting FGL based tree

selection are specified in <u>Section 3.2.3</u> and 3.2.4 for the same purpose. Sections 3.2.5 and 3.2.6 define two APPsub-TLVs to support finer granularity in selecting trees based on multicast group rather than Data Label.

New APPSubTLVs	Description
=======================================	=========
Tree and VLANS	announcement by the highest priority tree root of the VLANs allowed per tree
Tree and VLANS Used	tree-VLAN correspondence an ingress RBridge selects
Tree and FGLs	announcement by the highest priority tree root of the FGLs allowed per tree
Tree and FGLs Used	tree-FGL correspondence an ingress RBridge selects
Tree and GROUPs	announcement by the highest priority tree root of the multicast groups
Tree and GROUPs Used	allowed on each tree tree and multicast group correspondence
	an ingress RBridge selects

## 3.2.1. The Tree and VLANs APPsub-TLV

The RBridge that is the highest priority tree root announces the VLANs allowed on each tree with the Tree and VLANs (TREE-VLANS) APPsub-TLV. Multiple instances of this sub-TLV may be carried. The same tree nicknames may occur in multiple Tree-VLAN RECORDs within the same or across multiple sub-TLVs. The sub-TLV format is as follows:

```
1 1 1 1 1 1
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5
Type = tbd1
               (2 bytes)
| Length |
               (2 bytes)
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-
| Tree-VLAN RECORD (1)
                | (6 bytes)
Tree-VLAN RECORD (N) | (6 bytes)
```

where each Tree-VLAN RECORD is of the form:

+-+-+-+-+-+							
	Nickname						
+-+-+-+-+-+							
RESV	Start.VLAN		(2 bytes)				
+-+-+-+	+-+-+-+-+-+-	+-+-+-+					
RESV	End.VLAN		(2 bytes)				
+-+-+-+		+-+-+-+					

- o Type: TRILL GENINFO APPsub-TLV type, set to tbd1 (TREE-VLANS).
- o Length: 6\*n bytes, where there are n Tree-VLAN RECORDs. Thus the value of Length can be used to determine n. If Length is not a multiple of 6, the sub-TLV is corrupt and MUST be ignored.
- o Nickname: The nickname identifying the distribution tree by its root.
- o RESV: 4 bits that MUST be sent as zero and ignored on receipt.
- o Start.VLAN, End.VLAN: These fields are the VLAN IDs of the allowed VLAN range on the tree, inclusive. To specify a single VLAN, the VLAN's ID appears as both the start and end VLAN. If End.VLAN is less than Start.VLAN the Tree-VLAN RECORD MUST be ignored.

### 3.2.2. The Tree and VLANs Used APPsub-TLV

This APPsub-TLV has the same structure as the Tree and VLANs APPsub-TLV (TREE-VLANS) specified in <u>Section 3.2.1</u>. The differences are that its APPsub-TLV type is set to tbd2 (TREE-VLANS-USE) and the Tree-VLAN correspondences in the Tree-VLAN RECORDs listed are those the originating RBridge wants to use for multi-destination packets. This APPsub-TLV is used by an ingress RBridge to distribute the tree-VLAN correspondence it selects from the list announced by the highest priority tree root.

## 3.2.3. The Tree and FGLs APPsub-TLV

The RBridge that is the highest priority tree root can use the Tree and FGLs (TREE-FGLS) APPsub-TLV to announce the FGLs allowed on each tree. Multiple instances of this APPsub-TLV may be carried. The same tree nicknames may occur in the multiple Tree-FGL RECORDs within the same or across multiple APPsub-TLVs. Its format is as follows:

1 1 1 1	1 1	
0 1 2 3 4 5 6 7 8 9 0 1 2 3	3 4 5	
+-+-+-+-+-	+-+-+	
Type = tbd3		(2 bytes)
+-+-+-+-+-	+-+-+	
Length	I	(2 bytes)
+-+-+-+-+-	+-+-+-+	+-+
Tree-FGL RECORD (1)		(8 bytes)
+-+-+-+-+-+-	+-+-+-+	+ - +
+-+-+-+-+-	+-+-+-+	+ - +
Tree-FGL RECORD (N)		(8 bytes)
+-+-+-+-+-+-	+-+-+-+	+-+

where each Tree-VLAN RECORD is of the form:

+-+-+-+-+							
Nickname	(2	bytes)					
+-							
Start.FGL	(3	bytes)					
+-							
End.FGL	(3	bytes)					
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-++-+							

- o Type: TRILL GENINFO APPsub-TLV type, set to tbd3 (TREE-FGLS).
- o Length: 8\*n bytes, where there are n Tree-FGL RECORDs. Thus the value of Length can be used to determine n. If Length is not a multiple of 8, the sub-TLV is corrupt and MUST be ignored.
- o Nickname: The nickname identifying the distribution tree by its root.
- o RESV: 4 bits that MUST be sent as zero and ignored on receipt.
- o Start.FGL, End.FGL: These fields are the FGL IDs of the allowed FGL range on the tree, inclusive. To specify a single FGL, the FGL's ID appears as both the start and end FGL. If End.FGL is less than Start.FGL the Tree-FGL RECORD MUST be ignored.

## 3.2.4. The Tree and FGLs Used APPsub-TLV

This APPsub-TLV has the same structure as the Tree and FGLs APPsub-TLV (TREE-FGLS) specified in <u>Section 3.2.3</u>. The only difference is that its APPsub-TLV type is set to tbd4 (TREE-FGLS-USE), and the Tree-FGL RECORDs listed are those the originating RBridge wants to use for multi-destination packets. This APPsub-TLV is used by an ingress RBridge to distribute the tree-FGL correspondence it selects from the list announced by the highest priority tree root.

## 3.2.5. The Tree and Groups APPsub-TLV

Data Label based tree selection is easily extended to (Data Label + Layer 2 or 3 multicast group) based tree selection. We can appoint multicast group 1 in VLAN 10 to tree1 and appoint group 2 in VLAN 10 to tree2 for better load sharing.

The RBridge that is the highest priority tree root can announce the multicast groups allowed on each tree for each data label with the Tree and Groups (TREE-GROUPS) APPsub-TLV. Multiple instances of this sub-TLV may be carried. The sub-TLV format is as follows:

```
Type = tbd5
            | (2 byte)
| (2 byte)
Tree Nickname | (2 bytes)
Group Sub-Sub-TLVs
             (variable)
+-+-+-+-+-+-+-+....
```

- o Type: TRILL GENINFO APPsub-TLV type, set to tbd5 (TREE-GROUPS).
- o Length: 2 + the length of the Group Sub-Sub TLVs included
- o Nickname: The nickname identifying the distribution tree by its root.
- o Group Sub-Sub-TLVs: Zero or more of the TLV structures that are allowed as sub-TLVs of the GADDR TLV [RFC7176]. Each such TLV structure specifies a multicast group and either a VLAN or FGL. Although these TLV structure are considered sub-TLVs when they appear inside a GADDR TLV, they are technically sub-sub-TLVs when they appear inside a TREE-GROUPs APPsub-TLV which is in turn inside a TRILL GENINFO TLV [RFC7357].

## 3.2.6. The Tree and Groups Used APPsub-TLV

This APPsub-TLV has the same structure as the Tree and GROUPs APPsub-TLV (TREE-GROUPS) specified in <u>Section 3.2.5</u>. The only difference is that its APPsub-TLV type is set to tbd6 (TREE-GROUPS-USE), and the tree and multicast groups listed in this sub-TLV are those the originating RBridge wants to use for multi-destination packets. This APPsub-TLV is used by an ingress RBridge to distribute the tree-group correspondence it selects from the list announced by the highest priority tree root.

## 3.3. Detailed Processing

The highest priority tree root RBridge MUST include all the necessary tree related sub-TLVs defined in [RFC7176] as usual in its E-L1FS FS-LSP and MAY include the Tree and VLANs Sub-TLV (TREE-VLANs) and/or Tree and FGLs Sub-TLV (TREE-FGLs) in its E-L1FS FS-LSP [RFC7780]. In this way it MAY indicate that each VLAN and/or FGL is only allowed on one or some other number of trees less than the number of trees being calculated in the campus in order to save table space in the fast path forwarding hardware.

An ingress RBridge that understands the TREE-VLANs APPsub-TLV SHOULD select the tree-VLAN correspondences it wishes to use and put them in TREE-VLAN-USE APPsub-TLVs. If there are multiple tree nicknames announced in TREE-VLANs Sub-TLV for a VLAN x, ingress RBridge chooses one of them if it supports this feature. For example, the ingress RBridge may choose the closest (minimum cost) root among them. How to make such a choice is out of the scope of this document. It may be desirable to have some fixed algorithm to make sure all ingress RBs choose the same tree for VLAN x in this case. Any single Data Label that the ingress RBridge is interested in should be related to only one tree ID in TREE-VLAN-USE to minimize the multicast forwarding table size on other RBridges but as long as the Data Label is related to less than all the trees being calculated, it will reduce the burden on the forwarding table size.

When an ingress RBridge encapsulates a multi-destination frame for Data Label x, it SHOULD use a tree nickname that it selected previously in TREE-VLAN-USE or TREE-FGL-USE for Data Label x. However, that may not be possible because either (1) the RBridge may not have advertised such TREE-VLAN-USE or TREE-FGL-USE APPsub-TLVs, in which case it can use any tree that has been advertised as permitted for the Data Label by the highest priority tree root RBridge, or (2) the tree or trees it advertised might be unavailable due to failures.

If RBridge RBn does not perform pruning, it builds the multicast forwarding table as specified in [RFC6325].

If RBn prunes the distribution tree based on VLANs, RBn uses the information received in TREE-VLAN-USE APPsub-TLVs to mark the set of VLANs reachable downstream for each adjacency and for each related tree. If RBn prunes the distribution tree based on FGLs, RBn uses the information received in TRILL-FGL-USE APPsub-TLVs to mark the set of FLGs reachable downstream for each adjacency and for each related tree.

Logically, an ingress RBridge that does not support VLAN/FGL based

tree selection is equivalent to the one that supports it and announces all the combination pair of tree-id-used and interested-vlan/interested-fgl as TREE-VLAN-USE.

### 3.4. Failure Handling

This section discusses failure of a distribution tree root for the cases where that is not the highest priority root and the case where it is the highest priority root. It also discusses some other transient error conditions.

Failure of a tree root that is not the highest priority: It is the responsibility of the highest priority tree root to inform other RBridges of any change in the allowed tree-VLAN correspondence. When the highest priority tree root learns the root of tree t has failed, it should re-assign the VLANs allowed on tree t to other trees or to a tree replacing the failed one.

Failure of the highest priority tree root: It is suggested that the second highest priority tree root be pre-configured with the proper knowledge of the tree-VLAN correspondence allowed when the highest priority tree root fails. The information announced by the second priority tree root would be in the link state of all RBridges but would not take effect unless the RBridge noticed the failure of the highest priority tree root. When the highest priority tree root fails, the former second priority tree root will become the highest priority tree root of the campus. When an RBridge notices the failure of the original highest priority tree root, it can immediately use the stored information announced by the original second priority tree root. It is suggested that the tree-VLAN correspondence information be pre-configured on the second highest priority tree root to be the same as that on the highest priority tree root for the trees other than the highest priority tree itself. This can minimize the change to multicast forwarding tables in the case of highest priority tree root failure. For a large campus, it may make sense to pre-configure this information in a similar way on the third, fourth, or even lower priority tree root RBridges.

In some transient conditions or in case of misbehavior by the highest priority tree root, an ingress RBridge may encounter the following scenarios:

- No tree has been announced for which VLAN x frames are allowed.
- An ingress RBridge is supposed to transmit VLAN x frames on tree t, but root of tree t is no longer reachable.

For the second case, an ingress RBridge may choose another reachable

tree root which allows VLAN x according to the highest priority tree root announcement. If there is no such tree available, then it is the same as the first case above. Then the ingress RBridge should be 'downgraded' to a conventional RBridge with behavior as specified in [RFC6325]. A timer should be set to allow the temporary transient stage to complete before the change of responsive tree or 'downgrade' takes effect. The value of timer should be set to at least the LSP flooding time of the campus.

## 4. Backward Compatibility

RBridges MUST include the TREE-USE-IDs and INT-VLAN sub-TLVs in their LSPs when required by [RFC6325] whether or not they support the new TREE-VLAN-USE or TREE-FGL-USE sub-TLVs specified by this draft.

RBridges that understand the new TREE-VLAN-USE sub-TLV sent from another RBridge RBn should use it to build the multicast forwarding table and ignore the TREE-USE-IDs and INT-VLAN sub-TLVs sent from the same RBridge. TREE-USE-IDs and INT-VLAN sub-TLVs are still useful for some purposes other than building multicast forwarding table (E.g. RPF table building, spanning tree root notification, etc.) If the RBridge does not receive TREE-VLAN-USE sub-TLVs from RBn, it uses the conventional way described in [RFC6325] to build the multicast forwarding table.

For example, there are two distribution trees, tree1 and tree2, in the campus. RB1 and RB2 are RBridges that use the new APPsub-TLVs described in this document. RB3 is an old RBridge that is compatible with [RFC6325]. Assume RB2 is interested in VLANs 10 and 11 and RB3 is interested in VLANs 100 and 101. Hence RB1 receives ((tree1, VLAN10), (tree2, VLAN11)) as a TREE-VLAN-USE sub-TLV and (tree1, tree2) as a TREE-USE-IDs sub-TLV from RB2 on port x. And RB1 receives (tree1) as a TREE-USE-IDs sub-TLV and no TREE-VLAN-USE sub-TLV from RB3 on port y. RB2 and RB3 announce their interested VLANs in an INT-VLAN sub-TLV as usual. Then RB1 will build the entry of (tree1, VLAN10, port x) and (tree2, VLAN11, port x) based on RB2's LSP and the mechanism specified in this document. RB1 also builds entries of (tree1, VLAN100, port y), (tree1, VLAN101, port y), (tree2, VLAN100, port y), (tree2, VLAN101, port y) based on RB3's LSP in conventional way. The multicast forwarding table on RB1 with merged entry would be like the following.

+  tree nickname		
+	•	
tree 1	10	
	100	y
tree 1	101	
tree 2	11	
tree 2	100	
tree 2	101	

As expected, that table is not as small as the one where every RBridge supports the new TREE-VLAN-USE sub-TLVs. The worst case in a hybrid campus is the number of entries equal to the number in current practice which does not support VLAN based tree selection. Such an extreme case happens when the interested VLAN set from the new RBridges is a subset of the interested VLAN set from the old RBridges.

Data Label and multicast group based tree selection is compatible with the current practice. Its effectiveness increases with more RBridge supporting this feature in the TRILL campus.

# **5**. Security Considerations

This document does not change the general RBridge security considerations of the TRILL base protocol. The APPsub-TLVs specified can be secured using the IS-IS authentication feature [RFC5310]. See Section 6 of [RFC6325] for general TRILL security considerations.

## **6.** IANA Considerations

IANA is requested to assign six new TRILL APPsub-TLV type codes from the range less than 255 as specified in <u>Section 3</u> and update the "TRILL APPsub-TLV Types under IS-IS TLV 251 Application Identifier 1" Registry on the IANA TRILL Parameters web page as shown below.

Туре	Name of APPsub-TLV code	Reference
tbd1	Tree and VLANs	[this document 3.2.1]
tbd2	Tree and VLANs Used	[this document 3.2.2]
tbd3	Tree and FGLs	[this document 3.2.3]
tbd4	Tree and FGLs Used	[this document 3.2.4]
tbd5	Tree and Groups	[this document 3.2.5]
tbd6	Tree and Groups Used	[this document 3.2.6]

## 7. References

#### 7.1. Normative References

- [RFC6325] Perlman, R., Eastlake 3rd, D., Dutt, D., Gai, S., and A.
  Ghanwani, "Routing Bridges (RBridges): Base Protocol
  Specification", RFC 6325, July 2011, <a href="http://www.rfc-editor.org/info/rfc6325">http://www.rfc-editor.org/info/rfc6325</a>.
- [RFC7357] Zhai, H., Hu, F., Perlman, R., Eastlake 3rd, D., and O.
   Stokes, "Transparent Interconnection of Lots of Links
   (TRILL): End Station Address Distribution Information
   (ESADI) Protocol", RFC 7357, September 2014,
   <a href="http://www.rfc-editor.org/info/rfc7357">http://www.rfc-editor.org/info/rfc7357</a>>

#### 7.2. Informative References

[RFC5310] Bhatia, M., Manral, V., Li, T., Atkinson, R., White, R., and M. Fanto, "IS-IS Generic Cryptographic Authentication", <a href="https://www.rfc-editor.org/info/rfc5310">RFC 5310</a>, February 2009, <a href="http://www.rfc-editor.org/info/rfc5310">http://www.rfc-editor.org/info/rfc5310</a>.

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Authors' Addresses

Yizhou Li Huawei Technologies 101 Software Avenue, Nanjing 210012 China

Phone: +86-25-56624629 Email: liyizhou@huawei.com

Donald Eastlake Huawei R&D USA 155 Beaver Street Milford, MA 01757 USA

Phone: +1-508-333-2270 Email: d3e3e3@gmail.com

Weiguo Hao Huawei Technologies 101 Software Avenue, Nanjing 210012 China

Phone: +86-25-56623144 Email: haoweiguo@huawei.com

Hao Chen Huawei Technologies 101 Software Avenue, Nanjing 210012 China

Email: philips.chenhao@huawei.com

Somnath Chatterjee Cisco Systems, SEZ Unit, Cessna Business Park, Outer ring road, Bangalore - 560087 India

Email: somnath.chatterjee01@gmail.com