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## **Alternative Delta Time Encoding for CCNx Using Compact Floating-Point Arithmetic**

### **Abstract**

CCNx utilizes delta time for a number of functions. When using CCNx in environments with constrained nodes or bandwidth constrained networks, it is valuable to have a compressed representation of delta time. In order to do so, either accuracy or dynamic range has to be sacrificed. Since the current uses of delta time do not require both simultaneously, one can consider a logarithmic encoding such as that specified in [RFC5497] and [IEEE.754.2019]. This document updates *CCNx messages in TLV Format* [RFC8609] to specify this alternative encoding.

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## 1. Introduction

CCNx is well suited for Internet of Things (IoT) applications [RFC7927]. LowPAN adaptation layers (e.g., [RFC9139]) confirm the advantages of a space-efficient packet encoding for low-power and lossy networks. CCNx utilizes absolute and delta time values for a number of functions. When using CCNx in resource-constrained environments, it is valuable to have a compact representation of time values. However, any compact time representation has to sacrifice accuracy or dynamic range. For some time uses this is relatively straightforward to achieve, for other uses, it is not. As a result of experiments in bandwidth-constrained IEEE 802.15.4 deployments [ICNLOWPAN], this document discusses the various cases of time values, and proposes a compact encoding for delta times.

This document has received fruitful reviews by the members of the research group (see the Acknowledgments section). It is the consensus of ICNRG that this document should be published in the

IRTF Stream of the RFC series. This document does not constitute an IETF standard.

## 2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC 2119](#) [[RFC2119](#)].

This document uses the terminology of [[RFC8569](#)] and [[RFC8609](#)] for CCNx entities.

The following terms are used in the document and defined as follows:

**byte:** synonym for octet

**time value:** a time offset measured in seconds

**time code:** an 8-bit encoded time value

## 3. Usage of Time Values in CCNx

### 3.1. Relative Time in CCNx

CCNx, as currently specified in [[RFC8569](#)], utilizes delta time for only the lifetime of an Interest message (see sections 2.1, 2.2, 2.4.2, 10.3 of [[RFC8569](#)]). It is a hop-by-hop header value, and is currently encoded via the T\_INTLIFE TLV as a 64-bit integer ([[RFC8609](#)] section 3.4.1). While formally an optional TLV, in all but some corner cases every Interest message is expected to carry the Interest Lifetime TLV, and hence having compact encoding is particularly valuable for keeping Interest messages short.

Since the current uses of delta time do not require both accuracy and dynamic range simultaneously, one can consider a logarithmic encoding such as that specified in [[IEEE.754.2019](#)] and outlined in [Section 4](#). This document updates *CCNx messages in TLV Format* [[RFC8609](#)] to permit this alternative encoding for selected time values. See [Section 6](#) for the specific actions needed to register this alternative compact representation of Interest Lifetime.

### 3.2. Absolute Time in CCNx

CCNx, as currently specified in [[RFC8569](#)], utilizes absolute time for various important functions. Each of these absolute time usages poses a different challenge for a compact representation. These are discussed in the following subsections.

### 3.2.1. Signature Time and Expiry Time

*Signature Time* is the time the signature of a content object was generated (sections 8.2-8.4 [RFC8569]). *Expiry Time* indicates the expiry time of a content object (section 4 [RFC8569]). Both values are content message TLVs and represent absolute timestamps in milliseconds since the UTC epoch (i.e., an NTP timestamp). They are currently encoded via the T\_SIGTIME and T\_EXPIRY TLVs as 64-bit unsigned integers (see section 3.6.4.1.4.5 and 3.6.2.2.2 [RFC8609]).

Both time values could be in the past, or in the future, potentially by a large delta. They are also included in the security envelope of the message. Therefore, it seems there is no practical way to define an alternative compact encoding that preserves its semantics and security properties; hence we don't consider it further as a candidate.

### 3.2.2. Recommended Cache Time

*Recommended Cache Time* (RCT) for a content object (see section 4 [RFC8569]) is a hop-by-hop header stating the expiration time for a cached content object in milliseconds since the UTC epoch (i.e., an NTP timestamp). It is currently encoded via the T\_CACHETIME TLV as a 64-bit unsigned integer (see section 3.4.2 [RFC8609]).

A recommended cache time could be far in the future, but cannot be in the past and is likely to be a reasonably short offset from the current time. Therefore, this document allows the recommended cache time to be interpreted as a relative time value rather than an absolute time, since the semantics associated with an absolute time value do not seem to be critical to the utility of this value. This document therefore updates the recommended cache time with the following rule set:

- \*Use absolute time as per [RFC8609]

- \*Use relative time, if the compact time representation is used (see Section 4 and Section 5)

## 4. A Compact Time Representation with Logarithmic Range

This document uses the compact time representation of ICNLoWPAN (see section 7 of [RFC9139]) that is inspired by [RFC5497] and [IEEE.754.2019]. Its logarithmic encoding supports a representation ranging from milliseconds to years. Figure 1 depicts the logarithmic nature of this time representation.

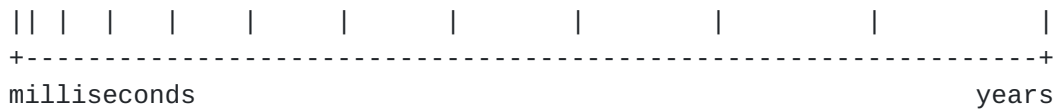


Figure 1: A logarithmic range representation allows for higher precision in the smaller time ranges and still supports large time deltas.

Time codes encode exponent and mantissa values in a single byte, but in contrast to the representation in [IEEE.754.2019], time codes only encode positive numbers and hence do not include an extra sign bit. Figure 2 shows the configuration of a time code: an exponent width of 5 bits, and a mantissa width of 3 bits.

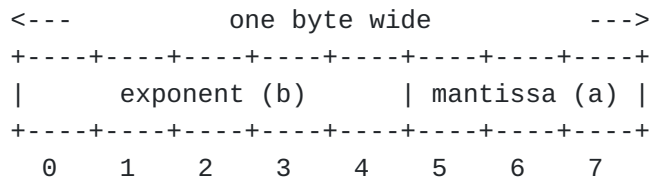


Figure 2: A time code with exponent and mantissa to encode a logarithmic range time representation.

The base unit for time values are seconds. A time value is calculated using the following formula (adopted from [RFC5497] and [RFC9139]), where (a) represents the mantissa, (b) the exponent, and (C) a constant factor set to  $C := 1/32$ .

**Subnormal (b == 0):**  $(0 + a/8) * 2 * C$

**Normalized (b > 0):**  $(1 + a/8) * 2^b * C$

The subnormal form provides a gradual underflow between zero and the smallest normalized number. Eight time values exist in the subnormal range [0s, ~0.054688s] with a step size of ~0.007812s between each time value. This configuration also encodes the following convenient numbers in seconds: [1, 2, 4, 8, 16, 32, 64, ...]. Appendix A further includes test vectors to illustrate the logarithmic range.

An example algorithm to encode a time value into the corresponding exponent and mantissa is given as pseudo code in Figure 3. Not all time values can be represented by a time code. For these instances, the closest time code is chosen that is smaller than the value to encode.

```

input: float v    // time value
output: int a, b // mantissa, exponent of time code

(a, b) encode (v):

    if (v == 0)
        return (0, 0)

    if (v < 2 * C) // subnormal
        a = floor (v * 4 / C) // round down
        return (a, 0)
    else // normalized
        if (v > (1 + 7/8) * 2^31 * C) // check bounds
            return (7, 31) // return maximum
        else
            b = floor (log2(v / C)) // round down
            a = floor ((v / (2^b * C) - 1) * 8) // round down
            return (a, b)

```

Figure 3: Algorithm in pseudo code.

As an example: No specific time code for 0.063 exists, but this algorithm maps to the closest valid time code that is smaller, i.e., exponent 1 and mantissa 0 (the same as for time value 0.0625).

## 5. Protocol Integration of the Compact Time Representation

A straightforward way to accommodate the compact time approach is to use a 1-byte length field to indicate this alternative encoding while retaining the existing TLV registry entries. This approach has backward compatibility problems, but is still considered for the following reasons:

- \*Both CCNx RFCs are experimental and not Standards Track, hence expectations for forward and backward compatibility are not as stringent. "Flag day" upgrades of deployed CCNx networks, while inconvenient, are still feasible.
- \*The major use case for these compressed encodings are smaller-scale IoT and/or sensor networks where the population of consumers, producers, and forwarders is reasonably small.
- \*Since the current TLVs have hop-by-hop semantics, they are not covered by any signed hash and hence may be freely re-encoded by any forwarder. That means a forwarder supporting the new encoding can translate freely between the two encodings.
- \*The alternative of assigning new TLV registry values does not substantially mitigate the interoperability problems anyway.



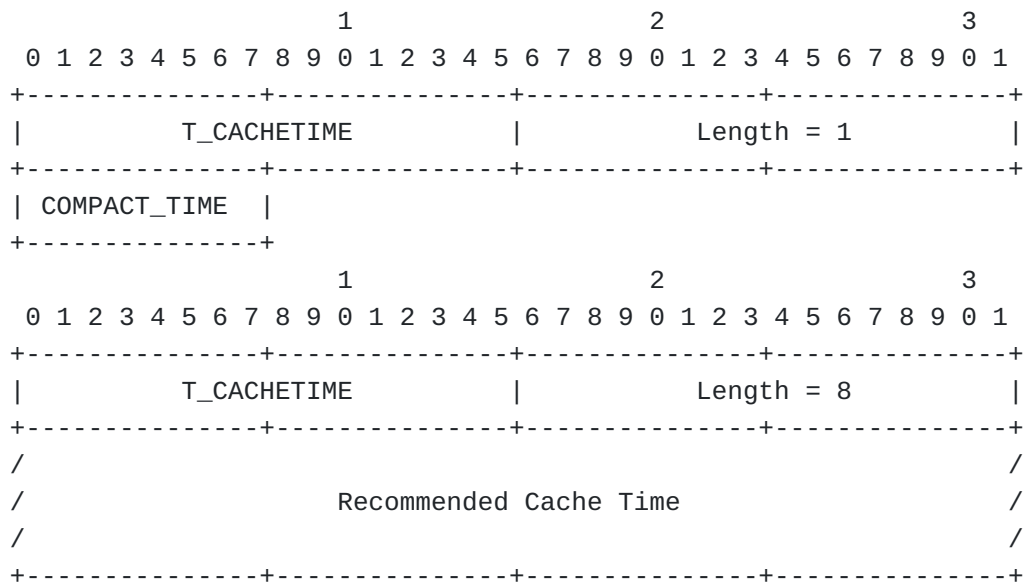


Figure 5: Changes to the definition of the Recommended Cache Time TLV.

The packet processing is adapted to calculate an absolute time from the relative time code based on the absolute reception time. On transmission, a new relative time code is calculated based on the current system time.

A note on legacy forwarders: A forwarder that does not support this compact time representation is expected to consider a Recommended Cache Time with length 1 as a structural or syntactical error and discard the packet. Otherwise, a forwarder interprets the compact 1-byte time value as an absolute time far in the past, which impacts cache utilization.

## 6. IANA Considerations

The protocol integration of the compact time encoding defined in [Section 5.1](#) and [Section 5.2](#) uses the following types registered in the CCNx Hop-by-Hop Types registry and assigned by IANA [[RFC8609](#)].

Type	Name
0x0001	T_INTLIFE
0x0002	T_CACHETIME

Table 1: Registered CCNx Hop-by-Hop Types

## 7. Security Considerations

This document makes no semantic changes to [[RFC8569](#)], nor to any of the security properties of the message encodings of [[RFC8609](#)], and



hence has the same security considerations as those two existing documents.

## 8. References

### 8.1. Normative References

- [RFC2119] Bradner, S. and RFC Publisher, "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, DOI 10.17487/RFC2119, March 1997, <<https://www.rfc-editor.org/info/rfc2119>>.
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- [RFC7927] Kutscher, D., Ed., Eum, S., Pentikousis, K., Psaras, I., Corujo, D., Saucez, D., Schmidt, T., Waehlich, M., and RFC Publisher, "Information-Centric Networking (ICN)

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[RFC9139] Gündoğan, C., Schmidt, T., Wählisch, M., Scherb, C., Marxer, C., Tschudin, C., and RFC Publisher, "Information-Centric Networking (ICN) Adaptation to Low-Power Wireless Personal Area Networks (LOWPANS)", RFC 9139, DOI 10.17487/RFC9139, November 2021, <<https://www.rfc-editor.org/info/rfc9139>>.

## Appendix A. Test Vectors

The test vectors in [Table 2](#) show sample time codes and their corresponding time values according to the algorithm outlined in [Section 4](#).

Time Code	Time Value (seconds)
0x00	0.000000
0x01	0.007812
0x04	0.031250
0x08	0.062500
0x15	0.203125
0x28	1.000000
0x30	2.000000
0xF8	67108864.000000
0xFF	125829120.000000

Table 2: Test Vectors

## Appendix B. Efficient Time Value Approximation

A forwarder frequently converts compact time into milliseconds to compare Interest lifetimes and the duration of cache entries. On many architectures, multiplication and division perform slower than addition, subtraction, and bit shifts. The following equations approximate the formulas in [Section 4](#), and scale from seconds into the milliseconds range by applying a factor of  $2^{10}$  instead of  $10^3$ . This results in an error of 2.4%.

$$\begin{aligned} \mathbf{b} == \mathbf{0}: & \quad 2^{10} * a * 2^{-3} * 2^1 * 2^{-5} \\ & = a \ll 3 \end{aligned}$$

$$\begin{aligned} \mathbf{b} > \mathbf{0}: & \quad (2^{10} + a * 2^{-3} * 2^{10}) * 2^b * 2^{-5} \\ & = (1 \ll 5 + a \ll 2) \ll b \end{aligned}$$

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