

v6ops
Internet-Draft
Intended status: Informational
Expires: 8 September 2022

G. Lencse
Szechenyi Istvan University
7 March 2022

Scalability of IPv6 Transition Technologies for IPv4aaS
draft-lencse-v6ops-transition-scalability-02

Abstract

Several IPv6 transition technologies have been developed to provide customers with IPv4-as-a-Service (IPv4aaS) for ISPs with an IPv6-only access and/or core network. All these technologies have their advantages and disadvantages, and depending on existing topology, skills, strategy and other preferences, one of these technologies may be the most appropriate solution for a network operator.

This document examines the scalability of the five most prominent IPv4aaS technologies (464XLAT, Dual Stack Lite, Lightweight 4over6, MAP-E, MAP-T) considering two aspects: (1) how their performance scales up with the number of CPU cores, (2) how their performance degrades, when the number of concurrent sessions is increased until hardware limit is reached.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of [BCP 78](#) and [BCP 79](#).

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at <https://datatracker.ietf.org/drafts/current/>.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on 8 September 2022.

Copyright Notice

Copyright (c) 2022 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to [BCP 78](#) and the IETF Trust's Legal Provisions Relating to IETF Documents (<https://trustee.ietf.org/license-info>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Revised BSD License text as described in Section 4.e of the [Trust Legal Provisions](#) and are provided without warranty as described in the Revised BSD License.

Table of Contents

- [1.](#) Introduction [2](#)
- [1.1.](#) Requirements Language [3](#)
- [2.](#) Scalability of iptables [3](#)
- [2.1.](#) Measurement Method [3](#)
- [2.2.](#) Performance scale up against the number of CPU cores . . [4](#)
- 2.3. Performance degradation caused by the number of sessions [6](#)
- [2.4.](#) Connection tear down rate [8](#)
- [3.](#) Scalability of Jool [9](#)
- [3.1.](#) Measurement Method [10](#)
- [3.2.](#) Performance scale up against the number of CPU cores . . [10](#)
- 3.3. Performance degradation caused by the number of sessions [11](#)
- [3.4.](#) Connection tear down rate [12](#)
- [4.](#) Acknowledgements [13](#)
- [5.](#) IANA Considerations [13](#)
- [6.](#) Security Considerations [13](#)
- [7.](#) References [13](#)
- [7.1.](#) Normative References [13](#)
- [7.2.](#) Informative References [14](#)
- [Appendix A.](#) Change Log [15](#)
- [A.1.](#) 00 [15](#)
- [A.2.](#) 01 [15](#)
- Author's Address [15](#)

[1.](#) Introduction

IETF has standardized several IPv6 transition technologies [[LEN2019](#)] and occupied a neutral position trusting the selection of the most appropriate ones to the market. [[I-D.ietf-v6ops-transition-comparison](#)] provides a comprehensive comparative analysis of the five most prominent IPv4aaS technologies

to assist operators with this problem. This document adds one more detail: measurement data regarding the scalability of the examined IPv4aaS technologies.

Currently, this document contains only the scalability measurements of the iptables stateful NAT44 implementation. It serves as a sample to test if the disclosed results are (1) useful and (2) sufficient for the network operators.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP14](#) [[RFC2119](#)] [[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.

2. Scalability of iptables

2.1. Measurement Method

[RFC8219] has defined a benchmarking methodology for IPv6 transition technologies. [[I-D.lencse-bmwg-benchmarking-stateful](#)] has amended it by addressing how to benchmark stateful NATxy gateways using pseudorandom port numbers recommended by [[RFC4814](#)]. It has defined a measurement procedure for maximum connection establishment rate and reused the classic measurement procedures like throughput, latency, frame loss rate, etc. from [[RFC8219](#)]. We used two of them: maximum connection establishment rate and throughput to characterize the performance of the examined system.

The scalability of iptables is examined in two aspects:

- * How its performance scales up with the number of CPU cores?
- * How its performance degrades, when the number of concurrent sessions is increased?

+-----+

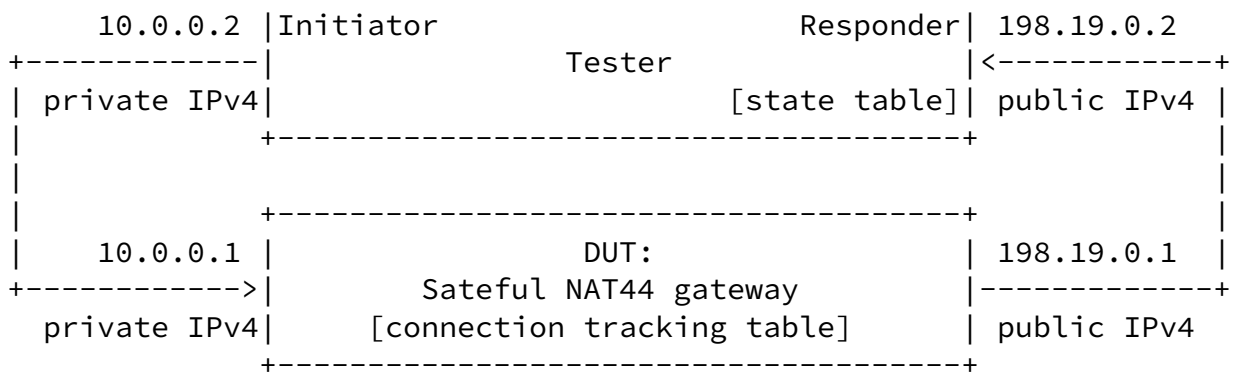


Figure 1: Test setup for benchmarking stateful NAT44 gateways

The test setup in Figure 1 was followed. The two devices, the Tester and the DUT (Device Under Test), were both Dell PowerEdge R430 servers having two 2.1GHz Intel Xeon E5-2683 v4 CPUs, 384GB 2400MHz DDR4 RAM and Intel 10G dual port X540 network adapters. The NICs of the servers were interconnected by direct cables, and the CPU clock frequency was set to fixed 2.1 GHz on both servers. They had Debian 9.13 Linux operating system with 4.9.0-16-amd64 kernel. The measurements were performed by siitperf [LEN2021] using the "stateful" branch (latest commit Aug. 16, 2021). The DPDK version was 16.11.11-1+deb9u2. The version of iptables was 1.6.0.

The ratio of number of connections in the connection tracking table and the value of the hashsize parameter of iptables significantly influences its performance. Although the default setting is `hashsize=nf_contrack_max/8`, we have usually set `hashsize=nf_contrack_max` to increase the performance of iptables, which was crucial, when high number of connections were used, because then the execution time of the tests was dominated by the preliminary phase, when several hundreds of millions connections had to be established. (In some cases, we had to use different settings due to memory limitations. The tables presenting the results always contain these parameters.)

The size of the port number pool is an important parameter of the benchmarking method for stateful NATxy gateways, thus it is also given for all tests.

[2.2.](#) Performance scale up against the number of CPU cores

To examine how the performance of iptables scales up with the number of CPU cores, the number of active CPU cores was set to 1, 2, 4, 8, 16 using the "maxcpus=" kernel parameter.

The number of connections was always 4,000,000 using 4,000 different source port numbers and 1,000 different destination port numbers. Both the connection tracking table size and the hash table size was set to 2^{23} .

The error of the binary search was chosen to be lower than 0.1% of the expected results. The experiments were executed 10 times.

Besides the connection establishment rate and the throughput of iptables, also the throughput of the IPv4 packet forwarding of the Linux kernel was measured to provide a basis for comparison.

The results are presented in Figure 2. The unit for the maximum connection establishment rate is 1,000 connections per second. The unit for throughput is 1,000 packets per second (measured with bidirectional traffic, and the number of all packets per second is displayed).

num. CPU cores	1	2	4	8	16
src ports	4,000	4,000	4,000	4,000	4,000
dst ports	1,000	1,000	1,000	1,000	1,000
num. conn.	4,000,000	4,000,000	4,000,000	4,000,000	4,000,000
contrack t. s.	2^{23}	2^{23}	2^{23}	2^{23}	2^{23}
hash table size	2^{23}	2^{23}	2^{23}	2^{23}	2^{23}
c.t.s/num.conn.	2.097	2.097	2.097	2.097	2.097
num. experiments	10	10	10	10	10
error	100	100	100	1,000	1,000
cps median	223.5	371.1	708.7	1,341	2,383
cps min	221.6	367.7	701.7	1,325	2,304
cps max	226.7	375.9	723.6	1,376	2,417
cps rel. scale up	1	0.830	0.793	0.750	0.666
throughput median	414.9	742.3	1,379	2,336	4,557
throughput min	413.9	740.6	1,373	2,311	4,436

throughput max	416.1	746.9	1,395	2,361	4,627
tp. rel. scale up	1	0.895	0.831	0.704	0.686
IPv4 packet forwarding (using the same port number ranges)					
error	200	500	1,000	1,000	1,000
throughput median	910.9	1,523	3,016	5,920	11,561
throughput min	874.8	1,485	2,951	5,811	10,998
throughput max	914.3	1,534	3,037	5,940	11,627
tp. rel. scale up	1	0.836	0.828	0.812	0.793
throughput ratio (%)	45.5	48.8	45.7	39.5	39.4

Figure 2: Scale up of iptables against the number of CPU cores
(Please refer to the next figure for the explanation of the abbreviations.)

abbreviation	explanation
-----	-----
num. CPU cores	number of CPU cores
src ports	size of the source port number range
dst ports	size of the destination port number range
num. conn.	number of connections = src ports * dst ports
conntrack t. s.	size of the connection tracking table of the DUT
hash table size	size of the hash table of the DUT
c.t.s/num.conn.	conntrack table size / number of connections
num. experiments	number of experiments
error	the difference between the upper and the lower bound of the binary search when it stops
cps (median/min/max)	maximum connection establishment rate (median, minimum, maximum)

cps rel. scale up	the relative scale up of the maximum connection establishment rate against the number of CPU cores
tp. rel. scale up	the relative scale up of the throughput
throughput ratio (%)	the ratio of the throughput of iptables and the throughput of IPv4 packet forwarding

Figure 3: Explanation of the abbreviations for the scale up of iptables against the number of CPU cores

Whereas the throughput of IPv4 packet forwarding scaled up from 0.91Mpps to 11.56Mpps showing a relative scale up of 0.793, the throughput of iptables scaled up from 414.9kpps to 4,557kpps showing a relative scale up of 0.686 (and the relative scale up of the maximum connection establishment rate is only 0.666). On the one hand, this is the price of the stateful operation. On the other hand, this result is quite good compared to the scale-up results of NSD (a high performance authoritative DNS server) presented in Table 9 of [LEN2020], which is only 0.52. (1,454,661/177,432=8.2-fold performance using 16 cores.) And DNS is not a stateful technology.

2.3. Performance degradation caused by the number of sessions

To examine how the performance of iptables degrades with the number connections in the connection tracking table, the number of connections was increased fourfold by doubling the size of both the source port number range and the destination port number range. Both the connection tracking table size and the hash table size was also increased four fold. However, we reached the limits of the hardware at 400,000,000 connections: we could not set the size of the hash table to 2^{29} but only to 2^{28} . The same value was used at 800,000,000 connections too, when the number of connections was only

doubled, because 1.6 billion connections would not fit into the memory.

The error of the binary search was chosen to be lower than 0.1% of the expected results. The experiments were executed 10 times (except for the very long lasting measurements with 800,000,000 connections).

The results are presented in Figure 4. The unit for the maximum

connection establishment rate is 1,000,000 connections per second. The unit for throughput is 1,000,000 packets per second (measured with bidirectional traffic, and the number of all packets per second is displayed).

num. conn.	1.56M	6.25M	25M	100M	400M	800M
src ports	2,500	5,000	10,000	20,000	40,000	40,000
dst ports	625	1,250	2,500	5,000	10,000	20,000
conntrack t. s.	2 ²¹	2 ²³	2 ²⁵	2 ²⁷	2 ²⁹	2 ³⁰
hash table size	2 ²¹	2 ²³	2 ²⁵	2 ²⁷	2 ²⁸	2 ²⁸
num. exp.	10	10	10	10	10	5
error	1,000	1,000	1,000	1,000	1,000	1,000
n.c./h.t.s.	0.745	0.745	0.745	0.745	1.490	2.980
cps median	2.406	2.279	2.278	2.237	2.013	1.405
cps min	2.358	2.226	2.226	2.124	1.983	1.390
cps max	2.505	2.315	2.317	2.290	2.050	1.440
throughput med.	5.326	4.369	4.510	4.516	4.244	3.689
throughput min	5.217	4.240	3.994	4.373	4.217	3.670
throughput max	5.533	4.408	4.572	4.537	4.342	3.709

Figure 4: Performance of iptables against the number of sessions

The performance of iptables shows degradation at 6.25M connections compared to 1.56M connections very likely due to the exhaustion of the L3 cache of the CPU of the DUT. Then the performance of iptables is fairly constant up to 100M connections. A small performance decrease can be observed at 400M connections due to the lower hash table size. A more significant performance decrease can be observed at 800M connections. It is caused by two factors:

- * on average, about 3 connections were hashed to the same place
- * non NUMA local memory was also used.

We note that the CPU has 2 NUMA nodes, cores 0, 2, ... 14 belong to NUMA node 0, and cores 1, 3, ... 15 belong to NUMA node 1. The maximum memory consumption with 400,000,000 connections was below 150GB, thus it could be stored in NUMA local memory.

Therefore, we have pointed out important limitations of the stateful

NAT44 technology:

- * there is a performance decrease, when approaching hardware limits
- * there is a hardware limit, beyond which the system cannot handle the connections at all (e.g. 1600M connections would not fit into the memory).

Therefore, we can conclude that, on the one hand, a well tailored hashing may guarantee an excellent scale-up of stateful NAT44 regarding the number of connections in a wide range, however, on the other hand, stateful operation has its limits resulting both in performance decrease, when approaching hardware limits and also in inability to handle more sessions, when reaching the memory limits.

[2.4.](#) Connection tear down rate

[I-D.lencse-bmwg-benchmarking-stateful] has defined connection tear down rate measurement as an aggregate measurement, that is, N number of connections are loaded into the connection tracking table of the DUT and then the entire content of the connection tracking table is deleted, and its deletion time is measured (T). Finally, the connection tear down rate is computed as: N/T .)

We have observed that the deletion of an empty connection tracking table of iptables may take a significant amount of time depending on its size. Therefore, we made our measurements more accurate by subtracting the deletion time of the empty connection tracking table from that of the filled one, thus we got the time spent with the deleting of the connections.

The same setup and parameters were used as in [Section 2.3](#) and the experiments were executed 10 times (except for the long lasting measurements with 800,000,000 connections).

The results are presented in Figure 5.

num. conn.	1.56M	6.35M	25M	100M	400M	800M
src ports	2,500	5,000	10,000	20,000	40,000	40,000
dst ports	625	1,250	2,500	5,000	10,000	20,000
contrack t. s.	2 ²¹	2 ²³	2 ²⁵	2 ²⁷	2 ²⁹	2 ³⁰
hash table size	2 ²¹	2 ²³	2 ²⁵	2 ²⁷	2 ²⁸	2 ²⁸
num. exp.	10	10	10	10	10	5
n.c./h.t.s.	0.745	0.745	0.745	0.745	1.490	2.980
full contr. del med	4.33	18.05	74.47	305.33	1,178.3	2,263.1
full contr. del min	4.25	17.93	72.04	299.06	1,164.0	2,259.6
full contr. del max	4.38	18.20	75.13	310.05	1,188.3	2,275.2
empty contr. del med	0.55	1.28	4.17	15.74	31.2	31.2
empty contr. del min	0.55	1.26	4.16	15.73	31.1	31.1
empty contr. del max	0.57	1.29	4.22	15.79	31.2	31.2
conn. deletion time	3.78	16.77	70.30	289.59	1,147.2	2,232.0
conn. tear d. rate	413,360	372,689	355,619	345,316	348,690	358,429

Figure 5: Connection tear down rate of iptables against the number of connections

The connection tear down performance of iptables shows significant degradation at 6.25M connections compared to 1.56M connections very likely due to the exhaustion of the L3 cache of the CPU of the DUT. Then it shows only a minor degradation up to 100M connections. A small performance increase can be observed at 400M connections due to the relatively lower hash table size. A more visible performance decrease can be observed at 800M connections. It is likely caused by keeping the hash table size constant and doubling the number of connections. The same thing that caused performance degradation of the maximum connection establishment rate and throughput, made now the deletion of the connections faster and thus caused an increase of the connection tear down rate.

We note that according to the recommended settings of iptables, 8 connections are hashed to each place of the hash table on average, but we wilfully used much smaller number (0.745 whenever it was possible) to increase the maximum connection establishment rate and thus to speed up experimenting. However, finally this choice significantly slowed down our experiments due to the very low connection tear down rate.

[3.](#) Scalability of Jool

[3.1.](#) Measurement Method

The same methodology was used as in [Section 2](#), but now the test setup in Figure 6 was followed. The same Tester and DUT devices were used as before, but the operating system of the DUT was updated to Debian 10.11 with 4.19.0-18-amd64 kernel to meet the requirement of the jool-tools package. The version of Jool was 4.1.6. (The most mature version of Jool at the date of starting the measurements, Release Date: 2021-12-10.)

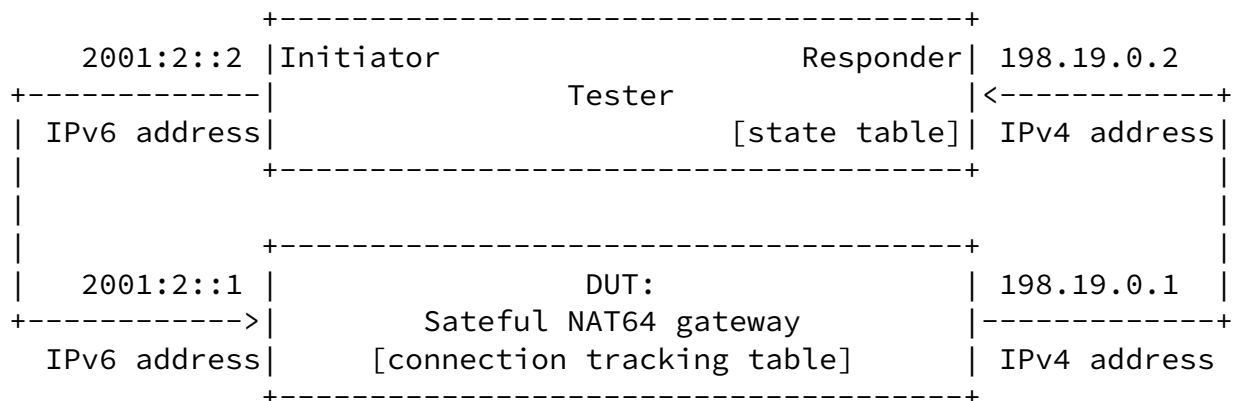


Figure 6: Test setup for benchmarking stateful NAT64 gateways

Unlike with iptables, we did not find any way to tune the hashsize or any other parameters of Jool.

[3.2.](#) Performance scale up against the number of CPU cores

The number of connections was always 1,000,000 using 2,000 different source port numbers and 500 different destination port numbers.

The error of the binary search was chosen to be lower than 0.1% of the expected results. The experiments were executed 10 times.

The results are presented in Figure 7. The unit for the maximum connection establishment rate is 1,000 connections per second. The

unit for throughput is 1,000 packets per second (measured with bidirectional traffic, and the number of all packets per second is displayed).

Internet-Draft Scalability of IPv4aaS Technologies March 2022

num. CPU cores	1	2	4	8	16
src ports	2,000	2,000	2,000	2,000	2,000
dst ports	500	500	500	500	500
num. conn.	1,000,000	1,000,000	1,000,000	1,000,000	1,000,000
num. experiments	10	10	10	10	10
error	100	100	100	100	100
cps median	228.6	358.5	537.4	569.9	602.6
cps min	226.5	352.5	530.7	562.0	593.7
cps max	230.5	362.4	543	578.3	609.7
cps rel. scale up	1	0.784	0.588	0.312	0.165
throughput median	251.8	405.7	582.4	604.1	612.3
throughput min	249.8	402.9	573.2	587.3	599.8
throughput max	253.3	409.6	585.7	607.2	616.6
tp. rel. scale up	1	0.806	0.578	0.300	0.152

Figure 7: Scale up of Jool against the number of CPU cores

Both the maximum connection establishment rate and the throughput scaled up poorly with the number of active CPU cores. The increase of the performance was very low above 4 CPU cores.

3.3. Performance degradation caused by the number of sessions

To examine how the performance of Jool degrades with the number connections, the number of connections was increased fourfold by doubling the size of both the source port number range and the destination port number range. We did not reach the limits of the hardware regarding the number of connections, because unlike iptables, Jool worked also with 1.6 billion connections.

The error of the binary search was chosen to be lower than 0.1% of

the expected results and the experiments were executed 10 times (except for the very long lasting measurements with 800 million and 1.6 billion connections to save execution time).

The results are presented in Figure 8. The unit for the maximum connection establishment rate is 1,000 connections per second. The unit for throughput is 1,000 packets per second (measured with bidirectional traffic, and the number of all packets per second is displayed).

num. conn.	1.56M	6.35M	25M	100M	400M	1600M
src ports	2,500	5,000	10,000	20,000	40,000	40,000
dst ports	625	1,250	2,500	5,000	10,000	40,000
num. exp.	10	10	10	10	5	5
error	100	100	100	100	1,000	1,000
cps median	480.2	394.8	328.6	273.0	243.0	232.0
cps min	468.6	392.7	324.9	269.4	243.0	230.5
cps max	484.9	397.4	331.3	280.6	244.5	233.6
throughput med.	511.5	423.9	350.0	286.5	257.8	198.4
throughput min	509.2	420.3	348.2	284.2	257.8	195.3
throughput max	513.1	428.3	352.5	290.8	260.9	201.6

Figure 8: Performance of Jool against the number of sessions

The performance of Jool shows degradation at the entire range of the number of connections. We did not analyze the root cause of the degradation yet. And we are not aware of the implementation of its connection tracking table. We also plan to check the memory consumption of Jool, what is definitely lower than that of iptables.

[3.4.](#) Connection tear down rate

Basically, the same measurement method was used as in [Section 2.4](#), however having no parameter of Jool to tune, only a single

measurement series was performed to determine the deletion time of the empty connection tracking table. The median, minimum and maximum values of the 10 measurements were 0.46s, 0.42s and 0.50s respectively.

The same setup and parameters were used as in [Section 2.3](#) and the experiments were executed 10 times (except for the long lasting measurements with 800,000,000 connections).

The results are presented in Figure 9. The unit for the connection tear down rate is 1,000,000 connections per second.

Num. conn.	1.56M	6.35M	25M	100M	400M	1600M
src ports	2,500	5,000	10,000	20,000	40,000	40,000
dst ports	625	1,250	2,500	5,000	10,000	40,000
num. exp.	10	10	10	10	10	5
full contr. del med	0.87	2.05	7.84	36.38	126.09	474.68
full contr. del min	0.80	2.02	7.80	36.27	125.84	473.20
full contr. del max	0.91	2.09	7.94	36.80	127.54	481.38
empty contr. del med	0.46	0.46	0.46	0.46	0.46	0.46
conn. deletion time	0.41	1.59	7.38	35.92	125.63	474.22
conn. t. d. r. (M)	3.811	3.931	3.388	2.784	3.184	3.374

Figure 9: Connection tear down rate of Jool against the number of connections

The connection tear down performance of Jool is excellent at any number of connections. It is about an order of magnitude higher than its connection establishment rate and than the connection tear down rate of iptables. (A slight degradation can be observed at 100M connections.)

4. Acknowledgements

The measurements were carried out by remotely using the resources of NICT StarBED, 2-12 Asahidai, Nomi-City, Ishikawa 923-1211, Japan. The author would like to thank Shuuhei Takimoto for the possibility to use StarBED, as well as to Satoru Gonno and Makoto Yoshida for their help and advice in StarBED usage related issues.

The author would like to thank Ole Troan for his comments on the v6ops mailing list, while the scalability measurements of iptables were intended to be a part of [[I-D.ietf-v6ops-transition-comparison](#)].

5. IANA Considerations

This document does not make any request to IANA.

6. Security Considerations

TBD.

7. References

7.1. Normative References

- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", [BCP 14](#), [RFC 2119](#), DOI 10.17487/RFC2119, March 1997, <<https://www.rfc-editor.org/info/rfc2119>>.
- [RFC4814] Newman, D. and T. Player, "Hash and Stuffing: Overlooked Factors in Network Device Benchmarking", [RFC 4814](#), DOI 10.17487/RFC4814, March 2007, <<https://www.rfc-editor.org/info/rfc4814>>.
- [RFC8174] Leiba, B., "Ambiguity of Uppercase vs Lowercase in [RFC 2119](#) Key Words", [BCP 14](#), [RFC 8174](#), DOI 10.17487/RFC8174, May 2017, <<https://www.rfc-editor.org/info/rfc8174>>.

Lencse

Expires 8 September 2022

[Page 13]

Internet-Draft

Scalability of IPv4aaS Technologies

March 2022

- [RFC8219] Georgescu, M., Pislaru, L., and G. Lencse, "Benchmarking Methodology for IPv6 Transition Technologies", [RFC 8219](#), DOI 10.17487/RFC8219, August 2017, <<https://www.rfc-editor.org/info/rfc8219>>.

7.2. Informative References

- [[I-D.ietf-v6ops-transition-comparison](#)]
Lencse, G., Martinez, J. P., Howard, L., Patterson, R., and I. Farrer, "Pros and Cons of IPv6 Transition Technologies for IPv4aaS", Work in Progress, Internet-

Draft, [draft-ietf-v6ops-transition-comparison-02](https://www.ietf.org/archive/id/draft-ietf-v6ops-transition-comparison-02), 3 March 2022, <<https://www.ietf.org/archive/id/draft-ietf-v6ops-transition-comparison-02.txt>>.

[I-D.lencse-bmwg-benchmarking-stateful]

Lencse, G. and K. Shima, "Benchmarking Methodology for Stateful NATxy Gateways using [RFC 4814](#) Pseudorandom Port Numbers", Work in Progress, Internet-Draft, [draft-lencse-bmwg-benchmarking-stateful-03](https://www.ietf.org/archive/id/draft-lencse-bmwg-benchmarking-stateful-03), 4 March 2022, <<https://www.ietf.org/archive/id/draft-lencse-bmwg-benchmarking-stateful-03.txt>>.

[LEN2019] Lencse, G. and Y. Kadobayashi, "Comprehensive Survey of IPv6 Transition Technologies: A Subjective Classification for Security Analysis", IEICE Transactions on Communications, vol. E102-B, no.10, pp. 2021-2035., DOI: 10.1587/transcom.2018EBR0002, 1 October 2019, <http://www.hit.bme.hu/~lencse/publications/e102-b_10_2021.pdf>.

[LEN2020] Lencse, G., "Benchmarking Authoritative DNS Servers", IEEE Access, vol. 8. pp. 130224-130238, DOI: 10.1109/ACCESS.2020.3009141, July 2020, <<https://ieeexplore.ieee.org/document/9139929>>.

[LEN2021] Lencse, G., "Design and Implementation of a Software Tester for Benchmarking Stateless NAT64 Gateways", IEICE Transactions on Communications, DOI: 10.1587/transcom.2019EBN0010, 1 February 2021, <<http://www.hit.bme.hu/~lencse/publications/IEICE-2020-siitperf-revised.pdf>>.

Initial version: scale up of iptables.

[A.2.](#) 01

Added the scale up of Jool.

Author's Address

Gabor Lencse
Szechenyi Istvan University
Gyor
Egyetem ter 1.
H-9026
Hungary
Email: lencse@sze.hu