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D. Migault
Ericsson
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LURK Extension version 1 for (D)TLS 1.3 Authentication
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Abstract

This document describes the LURK Extension 'tls13' which enables interactions between a LURK Client and a LURK Server in a context of authentication with (D)TLS 1.3.

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1. Introduction

This document defines a LURK extension for TLS 1.3 [[RFC8446](#)].

This document assumes the reader is familiar with TLS 1.3 the LURK architecture [[I-D.mglt-lurk-lurk](#)].

The motivations for the LURK Extension TLS 1.3 are similar to those for the LURK use cases [[I-D.mglt-lurk-tls-use-cases](#)].

Interactions with the Cryptographic Service (CS) can be performed by the TLS Client as well as by the TLS Server.

LURK defines an interface to a CS that stores the security credentials which include the PSK involved in a PSK or PSK-ECDHE authentication or the key used for signing in an ECDHE authentication. In the case of session resumption the PSK is derived from the resumption_master_secret during the key schedule [[RFC8446](#) [section 7.1](#)], this secret MAY require similar protection as well. On the other hand session resumption MAY also be delegated as in the LURK extension of TLS 1.2 [[I-D.mglt-lurk-tls12](#)].

The current document extends the scope of the LURK extension for TLS 1.2 in that it defines the CS on the TLS server as well as on the TLS client and the CS can operate in non delegating scenarios.

2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP 14](#) [[RFC2119](#)][[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.

This document uses the terms defined [[RFC8446](#)] and [[I-D.mgmt-lurk-tls12](#)].

3. LURK Header

LURK / TLS 1.3 is a LURK Extension that introduces a new designation "tls13". This document assumes that Extension is defined with designation set to "tls13" and version set to 1. The LURK Extension extends the LURKHeader structure defined in [[I-D.mgmt-lurk-lurk](#)] as follows:

```
enum {  
    tls13 (2), (255)  
} Designation;
```

```
enum {  
    capabilities(0),  
    ping(1),  
    s_init_early_secret(2),  
    s_init_cert_verify(3),  
    s_hand_and_app_secret(4),  
    s_new_ticket(5),  
    c_binder_key(6),  
    c_init_early_secret(7),  
    c_init_hand_secret(8),  
    c_hand_secret(9),  
    c_app_secret(10),  
    c_cert_verify(11),  
    c_register_ticket(12),  
    c_post_hand(13), (255)  
} TLS13Type;
```

```
enum {  
    // generic values reserved or aligned with the  
    // LURK Protocol  
    request (0), success (1), undefined_error (2),  
    invalid_payload_format (3),  
  
    invalid_psk  
    invalid_freshness_func  
  
    invalid_request  
    invalid_key_id_type  
    invalid_key_id  
    invalid_signature_scheme  
    invalid_certificate_type  
    invalid_certificate
```



```
    invalid_certificate_verify
    invalid_key_request
    invalid_handshake
    invalid_extension
    invalid_ephemeral
    invalid_cookie_h

}TLS13Status

struct {
    Designation designation = "tls13";
    int8 version = 1;
} Extension;

struct {
    Extension extension;
    select( Extension ){
        case ("tls13", 1):
            TLS13Type;
    } type;
    select( Extension ){
        case ("tls13", 1):
            TLS13Status;
    } status;
    uint64 id;
    uint32 length;
} LURKHeader;
```

4. Generic structures

The CS is not expected to perform any policies such as choosing the appropriated authentication method. Such choices are performed by the TLS client or TLS server that instruct the LURK client accordingly. Such enforced policies between the TLS client and the TLS server are those performed on a standard TLS exchange.

On the other hand, some CS MAY be optimized by implementing a subset of the specified possibilities described in this document. Typically some implementations MAY not implement the session resumption or the post handshake authentication to avoid keeping states of a given session once the handshake has been performed. These capabilities of the CS MAY also in return impact the policies of the TLS client or TLS server.

These limitations are mentioned throughout the document, and even represented in the state diagrams, the recommendation is that the CS SHOULD NOT impact the policies of the TLS client or TLS server. Instead they SHOULD be able to optimize the CS to their policies via

some configuration parameters presented in section [Section 10.1](#). Such parameters are implementation dependent and only provided here as informative.

This document defines the role to specify whether the CS runs on a TLS client or a TLS service. The CS MUST be associated a single role.

LURK exchanges falls into three categories: 1) request of keys or secrets, 2) request of signing operations, and 3) requests for ticket (NewSessionTicket) management purposes. In some cases, these operations are combined into a single LURK exchange. Table Figure 1 below summarizes the operations associated for each exchange.

Role	LURK exchange	secret	sign	ticket
server	s_init_early_secret	yes	-	-
server	s_init_cert_verify	yes	yes	-
server	s_hand_and_app_secret	yes	-	-
server	s_new_ticket	yes	-	yes
client	c_binder_key	yes	-	-
client	c_init_early_secret	yes	-	-
client	c_init_hand_secret	yes	-	-
client	c_hand_secret	yes	-	-
client	c_app_secret	yes	-	-
client	c_cert_verify	yes	yes	-
client	c_register_ticket	yes	-	yes
client	c_post_hand	-	yes	-

Figure 1: Operation associated to LURK exchange

This section describes structures that are widely re-used across the multiple LURK exchanges.

4.1. key_request

key_request is a 16 bit structure described in Table Figure 2 that indicates the requested key or secrets by the LURK client. The same structure is used across all LURK exchanges, but each LURK exchange only permit a subset of values described in Table Figure 3.

A LURK client MUST NOT set key_request to key or secrets that are not permitted. The CS MUST check the key_request has only permitted values and has all mandatory keys or secrets set. If these two criteria are not met the CS MUST NOT perform the LURK exchange and SHOULD return a invalid_key_request error. If the CS is not able to

compute an optional key or secret, the CS MUST proceed the LURK exchange and ignore the optional key or secret.

Bit	key or secret (designation)
0	binder_key (b)
1	client_early_traffic_secret (e_c)
2	early_exporter_master_secret (e_x)
3	client_handshake_traffic_secret (h_c)
4	server_handshake_traffic_secret (h_s)
5	client_application_traffic_secret_0 (a_c)
6	server_application_traffic_secret_0 (a_s)
7	exporter_master_secret (x)
8	resumption_master_secret (r)
9-15	reserved and set to zero

Figure 2: key_request structure

LURK exchange	Permitted key/secrets
s_init_early_secret	b, e_c*, e_x*
s_init_cert_verify	h_c, h_s, a_c*, a_s*, x*
s_hand_and_app_secret	h_c, h_s, a_c*, a_s*, x*
s_new_ticket	r*
c_binder_key	b
c_init_early_secret	e_c*, e_x*
c_init_hand_secret	h_c, h_s
c_hand_secret	h_c, h_s
c_app_secret	a_c*, a_s*, x*
c_cert_verify	a_c*, a_s*, x*
c_register_ticket	r*
c_post_hand	

(*) indicates an optional value, other values are mandatory

Figure 3: key_request permitted values per LURK exchange

4.2. secrets

The Secret structure carries a secret designated by its type and value.


```
enum {  
    binder_key (0),  
    client_early_traffic_secret(1),  
    early_exporter_master_secret(2),  
    client_handshake_traffic_secret(3),  
    server_handshake_traffic_secret(4),  
    client_application_traffic_secret_0(5),  
    server_application_traffic_secret_0(6),  
    exporter_master_secret(7),  
    esumption_master_secret(8),  
    (255)  
} SecretType;
```

```
struct {  
    SecretType secret_type;  
    opaque secret_data<0..2^8-1>;  
} Secret;
```

secret_type: The type of the secret or key

secret_data: The value of the secret.

4.3. handshake_context

Secrets derivation takes Handshake Context as input. It is the responsibility of the CS to maintain this variable in an internal variable. On the other hand, it is the responsibility of the LURK client to provide the necessary element so the Cryptographic Service got the necessary Handshake Context. The Handshake Context evolves during the key derivation schedule, the LURK client implements an incremental approach where only the missing part of the Handshake Context are provided. The main intention is to prevent the LURK client from providing multiple time the same information as well as to perform extensive compatibility checks between the duplicated information provided.

The handshake_context variable is based on the Handshake structure defined in [\[RFC8446\] section 4](#). The table below lists the values of the handshake_context associated to each LURK exchange.

+-----+	+-----+	+-----+
LURK exchange	handshake_context	
+-----+	+-----+	+-----+
s_init_early_secret	ClientHello	
s_init_cert_verify	ClientHello ... later of	
	server EncryptedExtensions /	
	CertificateRequest	
s_hand_and_app_secret	ServerHello ... later of	
	server EncryptedExtensions /	
	CertificateRequest	
s_new_ticket	earlier of client Certificate /	
	client CertificateVerify /	
	Finished ... Finished	
c_binder_key		
c_init_early_secret	ClientHello	
c_init_hand_secret	ClientHello ... ServerHello	
c_hand_secret	ServerHello	
c_app_secret	server EncryptedExtensions ...	
	server Finished	
c_cert_verify	server EncryptedExtensions ...	
	later of server Finished/	
	EndOfEarlyData	
c_register_ticket	earlier of client Certificate	
	client CertificateVerify ...	
	client Finished	
c_post_hand	CertificateRequest	
+-----+	+-----+	+-----+

Figure 4: handshake values per LURK exchange

4.4. Secret Sub Exchange

Secrets are derived from the key schedule of [\[RFC8446\] section 7](#).

The derivation of secrets requires an optional PSK that is provided in the psk_id extension described in section [Section 4.4.2](#) as well as a ECDHE value which is provided by the ecde extension described in section [Section 4.4.1](#).

The extensions considered in this document are defined as below:


```
enum { psk_id(1), ephemeral(2), freshness(3), session_id(4) ... (255)
} LURK13ExtensionType;

struct {
    LURK13ExtensionType extension_type;
    opaque extension_data<0..2^16-1>;
} LURK13Extension

struct {
    uint16 key_request;
    Handshake handshake_context<0..2^32> //RFC8446 section 4.
    LURK13Extension extension_list<0...2^16>
} SecretsRequest;

struct {
    Secret secret_list<0..2^16-1>;
    Extension extension_list<0...2^16>
} SecretsResponse;
```

key_request: designates the requested secrets (see section [Section 4.1](#)).

handshake_context: designates the necessary messages so the CS is aware of the appropriated Handshake Context to generate the secrets (see section [Section 4.3](#)).

extension_list: the list of extensions.

secret_list: the list of requested secrets (see section [Section 4.2](#)).

[4.4.1](#). Ephemeral Extension

The Ephemeral structure carries the necessary information to generate the ECDHE input used to derive the secrets. This document describes two ways the shared secret can be generated: `shared_secret_provided`: When Diffie Hellman or ECDH keys and shared secret are generated by the TLS server and the shared secret is provided to the CS

`secret_generated`: When the DH / ECDH keys and shared secret are generated by the CS.

[4.4.1.1](#). `shared_secret_provided`:

When ECDHE shared secret are not generated by the CS, the LURK client provides the shared secret value to the CS via the ephemeral extension. The shared secret is transmitted via the

EphemeralSharedSecret, constructed similarly to the key_exchange parameter of the KeyShareEntry described in [\[RFC8446\] section 4.2.8](#). The CS MUST NOT return any data.

```
struct {  
    NamedGroup group;  
    opaque shared_secret[coordinate_length];  
} EphemeralSharedSecret;
```

Where coordinate_length depends on the chosen group. For secp256r1, secp384r1, secp521r1, x25519, x448, the coordinate_length is respectively 32 bytes, 48 bytes, 66 bytes, 32 bytes and 56 bytes. Upon receiving the shared_secret, the CS MUST check group is proposed in the KeyShareClientHello and agreed in the KeyShareServerHello.

4.4.1.2. secret_generated:

When the ECDHE public/private keys are generated by the CS, the LURK client requests the CS the associated public value. Note that in such cases the CS would receive an incomplete Handshake Context from the LURK client with the public part of the ECDHE missing. Typically the ServerHello message would present a KeyShareServerHello that consists of a KeyShareEntry with an empty key_exchange field, but the field group is present.

The CS MUST check the group field in the KeyShareServerHello, and get the public value of the TLS client from the KeyShareClientHello. The CS performs the same checks as described in [\[RFC8446\] section 4.2.8](#). The CS generates the private and public Diffie Hellman or ECDH keys, computes the shared key and return the KeyShareEntry server_share structure defined in [\[RFC8446\] section 4.2.8](#).

Other methods may be defined in the future.

This extension MUST NOT be sent outside the LURK exchanges mentioned below. When received outside these exchanges, the CS SHOULD return an invalid_extension error. When the ephemeral is not supported, an invalid_ephemeral error SHOULD be returned. The ephemeral extension MUST NOT appear more than once in a LURK session. When the extensions appears in more than one LURK exchange an invalid_ephemeral error SHOULD be returned

+-----+-----+	
LURK exchange	Presence
+-----+-----+	
s_init_early_secret	-
s_init_cert_verify	M
s_hand_and_app_secret	*
s_new_ticket	-
c_binder_key	*
c_init_early_secret	-
c_init_hand_secret	*
c_hand_secret	*
c_app_secret	-
c_cert_verify	-
c_register_ticket	-
c_post_hand	-
+-----+-----+	

M indicates the extension is mandatory
 - indicates the extension MUST NOT be provided
 * indicates the extension MAY be provided

Figure 5: Ephemeral Extension presence per LURK exchange

The extension data is defined as follows:

```
enum { secret_provided(0), secret_generated(1) (255)} EphemeralMethod;
```

```
EphemeralDataRequest {
    EphemeralMethod method;
    select(method) {
        case secret_provided:
            EphemeralSharedSecret shared_secret<0..2^16>;
    }
}
```

```
EphemeralDataResponse {
    select(method) {
        case secret_generated:
            KeyShareEntry server_share
    }
}
```

4.4.2. PSK_id Extension

The psk_id indicates the identity of the PSK used in the key schedule.

The LURK client MUST provide this extension only when PSK or PSK-authentication is envisioned and when the PSK has not been provided

earlier. These exchanges are `s_init_early_secret` on the TLS server. On the TLS client side, these exchanges are `c_binder_key`, `c_init_early_secret` and `c_init_hand_secret`. The LURK client MUST NOT provide this extension outside these exchanges. When receiving the PSK extension outside these messages, the CS MUST NOT proceed to the exchange and SHOULD return a `invalid_format` error.

LURK exchange	Presence
<code>s_init_early_secret</code>	M
<code>s_init_cert_verify</code>	-
<code>s_hand_and_app_secret</code>	-
<code>s_new_ticket</code>	-
<code>c_binder_key</code>	M
<code>c_init_early_secret</code>	M
<code>c_init_hand_secret</code>	*
<code>c_hand_secret</code>	-
<code>c_app_secret</code>	-
<code>c_cert_verify</code>	-
<code>c_register_ticket</code>	-
<code>c_post_hand</code>	-

M indicates the extension is mandatory
 - indicates the extension MUST NOT be provided
 * indicates the extension MAY be provided

Figure 6: psk extension presence per LURK exchange

The extension data is defined as follows:

```
PsIdentity psk_id; //RFC8446 section 4.2.11
```

When the psk extension is provided in LURK exchange that is not permitted an `invalid_extension` error SHOULD be returned.

Upon receiving this extension in the permitted LURK exchange the CS checks the PSK is available. In case the PSK is not available, an `invalid_psk` error is returned. If the PSK is not provided, a default PSK is generated as described in [\[RFC8446\] section 7.1](#). If the default PSK is not allowed then an `invalid_psk` is returned.

4.4.3. Freshness Extension

The `freshness_function` provides perfect forward secrecy (PFS) and is used by the LURK client on the TLS client to generate the `ClientHello.random` or by the LURK client on the TLS server to

generate the `ServerHello.random`. When these randoms are provided to the CS, the `freshness_function` MUST be provided as well.

Table Figure 7 lists the LURK exchange that MUST include the freshness function extension as well as those where the extension may be provided.

LURK exchange	Presence
<code>s_init_early_secret</code>	-
<code>s_init_cert_verify</code>	M
<code>s_hand_and_app_secret</code>	M
<code>s_new_ticket</code>	-
<code>c_binder_key</code>	-
<code>c_init_early_secret</code>	M
<code>c_init_hand_secret</code>	M
<code>c_hand_secret</code>	-
<code>c_app_secret</code>	-
<code>c_cert_verify</code>	-
<code>c_register_ticket</code>	-
<code>c_post_hand</code>	-

* indicates the extension MAY be provided
M indicates the extension is mandatory
- indicates the extension MUST NOT be provided

Figure 7: `freshness_func` extension presence per LURK exchange

The extension data is defined as follows:

```
FreshnessFunc freshness_func; // {{I-D.mglt-lurk-tls12}} section 4.1
```

If the CS does not support the `freshness_func`, an `invalid_freshness_func` error is returned.

Perfect forward secrecy is implemented in a similar manner as with the TLS 1.2 extension described in [[I-D.mglt-lurk-tls12](#)] [section 4.1.1](#). As `ServerHello.random` in TLS 1.3 do not include time, it is not considered here. In addition, we use a specific context related to TLS 1.3.

As a result, the `ServerHello.random` is generated as follows on the TLS server.

```
ServerHello.random = freshness_func( server_random + "tls13 pfs srv" );
```


The ClientHello.random is generated as follows on the TLS client side:

```
ClientHello.random = freshness_func( server_random + "tls13 pfs clt" );
```

Perfect forward secrecy applies to the ServerHello.random on the TLS server and on the ClientHello.random on the TLS client. As a result, PFS is provided on the TLS server as long as the ServerHello is part of the Handshake Context. Similarly PFS is provided on the TLS client as long as ClientHello is part of the Handshake Context. On the TLS server, s_init_early_secret exchange do not have the ServerHello so this exchange is not protected by PFS later exchanges are. On the TLS client side, c_binder_key does not have any Handshake Context so this exchange is not protected by PFS. Later exchanges are.

4.4.4. Session ID Extension

A LURK client and the CS are likely to establish a LURK session. A session is mandatory to be set when a given TLS session requires multiple interactions between the Lurk client and the CS. Stateless interactions MAY be possible with ECDHE authentication without session resumption. Other configuration MAY also use other configurations to determine the session, such as a TCP session for example, in which case multiplexing LURK sessions over a common transport layer is not possible.

The LURK client indicates its willing to set a session as well as the session ID that should be used by the CS by inserting a Session ID Extension. The LURK client MAY only insert the Extension in a LURK message that initiates a session.

Upon receiving the Session ID Extension in an unexpected message, the CS MUST return an invalid_extension error. Upon receiving the Session ID Extension in an expected LURK message, The CS MAY ignored the received extension indicating thus to the LURK client that no session ID are needed. The LURK client MUST NOT insert a session ID in the following messages. If the CS agrees on setting a LURK session, the CS will return the Extension back with the expected Session ID use for the communication between the LURK client and the CS. When a session ID has been agreed all remaining exchange contain the session ID provided by the peer.

The CS MAY require a session ID being agreed between the LURK client and the CS. When the LURK client does not include the Session ID Extension, the CS MUST respond with an invalid_session_id error.

The policy to have session ID on LURK message is a policy that applies to the CS and LURK client cannot have different policies. When session ID are enabled, the CS expect every LURK message to have a session ID except for the initiating messages.

+-----+-----+		
LURK exchange	Presence	
+-----+-----+		
s_init_early_secret	*	
s_init_cert_verify	*	
s_hand_and_app_secret	-	
s_new_ticket	-	
c_binder_key	-	
c_init_early_secret	*	
c_init_hand_secret	*	
c_hand_secret		
c_app_secret	-	
c_cert_verify	-	
c_register_ticket	-	
c_post_hand	-	
+-----+-----+		
* indicates the extension MAY be provided		
- indicates the extension MUST NOT be provided		

Figure 8: Presence of the Session ID Extension in the various LURK exchanges

The extension data is defined as follows:

```
uint32 session_id
```

The session ID agreement leads to the definition of the following structure that will be embedded into any non initiating LURK exchange.

```
struct{
    select( session_id_agreed ){
        case True:
            uint32 session_id
        case False:
    }
} SessionID
```

session_id_agreed: indicates the LURK client and the CS have agreed on using session ID as well as the respective session ID value to use.

session_id can take the following values: 1. session_id_cs: the session ID provided by the CS to the LURK client in the Session ID extension. This is the value the LURK client MUST use in any subsequent exchange. That value will be used by the CS to associate its internal context to the session. 2. session_id_client: the session ID provided by the LURK client to the CS in the Session ID extension. This is the value the CS MUST use in any subsequent exchange. That value will be used by the LURK client to associate its internal context to the session.

4.5. Signing Sub-Exchange

The signature requires the signature scheme (sig_algo), the designated private key (key_id), as well as sufficient context to generate the necessary data to be signed. In our case the necessary context is provided by the LURKCertificate, assuming the CS will have the necessary Handshake Context. The latest may be provided in a combination of a secret request.

key_id is processed as described in [[I-D.mglt-lurk-tls12](#)] [section 4.1](#). If the CS does not support the KeyPairIdType an invalid_key_id_type is returned. If the CS does not recognize the key, an invalid_key_id error is returned.

sig_algo designates the signature algorithm scheme, and it is defined in [\[RFC8446\] section 4.2.3](#). When the CS does not support the signature scheme an invalid_signature_scheme error is returned.

The certificate is a public data that may repeat over multiple distinct TLS handshakes. To limit the load of unnecessary information being transmitted multiple times, the LURKCertificate enables to carry the index of the Certificate structure rather than the structure itself. When the lurk_certificate_type is set to sha256_32, the index of the Certificate structure is sent. The current specification generates the index using sha256_32, that is the first 32 bits of the hash of the Certificate structure using SHA256 as the hashing function. When lurk_certificate_type is set to X509 or RawPublicKey the full Certificate structure is expected. When the CS does not support the certificate_type, an invalid_certificate_type error is returned. When the Certificate structure does not match the private key, an invalid_certificate error is returned.

Signing operations are described in [[RFC8446](#)] [section 4.4.3](#). The context string is derived from the role and the type of the LURK exchange as described below. The Handshake Context is taken from the key schedule context.


```

+-----+-----+
| type           | context                               |
+-----+-----+
| s_init_cert_verify | "TLS 1.3, server CertificateVerify" |
| c_cert_verify     | "TLS 1.3, client CertificateVerify" |
+-----+-----+

```

The CS computes the signature as described in [\[RFC8446\] section 4.4.3](#). and returns signature in SigningResponse. When the CS does not have the necessary Handshake Context, context or is unable to proceed to the signing operation, an invalid_certificate_verify error is returned.

The structure is represented below:

```

enum { X509(0), RawPublicKey(1),
       sha256_32(128) (255)}; LURK13CertificateType

struct {
    LURK13CertificateType certificate_type;
    select (lurk_certificate_type) {
        case sha256_32:
            uint32 hash_cert;
        case X509, RawPublicKey:
            Certificate certificate; // RFC8446 section 4.4.2
    };
} LURK13Certificate;

struct {
    KeyPairId key_id; // draft-mglt-lurk-tls12 section 4.1
    SignatureScheme sig_algo; //RFC8446 section 4.2.3.
    LURKCertificate certificate;
} SigningRequest;

struct {
    opaque signature<0..2^16-1>; //RFC8446 section 4.4.3.
} SigningResponse;

```

5. LURK exchange on the TLS server

This section describes the LURK exchanges that are performed on the TLS server. The state diagram is provided in section [Section 10.2](#)

5.1. s_init_early_secret

A TLS server MAY receive a ClientHello that proposes PSK or PSK-ECDHE authentication via the `pre_shared_key` and `psk_key_exchange_modes` extensions. Depending on its policies, the TLS server MAY decide to proceed to such authentication. It chooses a PSK identity so the LURK client initiates a key schedule context (`ks_ctx`) that will manage the session with the CS. This session is initiated with a `s_init_early_secret` exchange.

The `binder_key` MUST be requested, since it is used to validate the PSK.

The TLS client MAY indicate support for early application data via the `early_data` extension. Depending on the TLS server policies, it MAY accept early data and request the `client_early_traffic_secret`.

The TLS server MAY have specific policies and request `early_exporter_master_secret`.

Upon receiving an `s_init_early_secret` request, the CS proceeds the `SecretRequest` as described in section [Section 4.4](#).

The CS MUST check `pre_shared_key` and `psk_key_exchange_modes` extensions are present in the ClientHello. If these extensions are not present, a `invalid_handshake` error SHOULD be returned.

The CS MUST ignore the `client_early_traffic_secret` if `early_data` extension is not found in the ClientHello. The Cryptographic Service MAY ignore the request for `client_early_traffic_secret`, in any case. The CS MAY ignored the request for `early_exporter_master_secret`.

```
struct{
    SecretRequest secret_request
} InitEarlySecretRequest
```

```
struct{
    SecretResponse secret_response
} InitEarlySecretResponse
```

`secret_request`: The structure associated to the secret request defined in section [Section 4.4](#)

`secret_response`: The structure associated to the secret request defined in section [Section 4.4](#).

5.2. s_init_cert_verify

A TLS server MAY receive a ClientHello that proposes ECDHE authentication with a key_share extension. Depending on its policies, the TLS server MAY decide to proceed to such authentication and indicate it to the LURK client so it initiates a key schedule context (ks_ctx) that will manage the session with the CS. This session is initiated with a s_init_cert_verify exchange.

The Cryptographic MUST ensure the ServerHello has selected the ECDHE authentication that is a key_share extension is present and no pre_shared_key extension is present. If these conditions are not met, a invalid_handshake error SHOULD be returned.

In order to provide generate the client_application_traffic_secret_0 and server_application_traffic_secret_0, the CS generates the server Finished. This value is computed to avoid multiple round trips. This value is not returned to the LURK client and needs to be computed again by the TLS server.

After the exchange is completed, the TLS server is able to build and return the ServeHello and complete the TLS handshake.

If the CS has been configured not to handle session resumption. The session is finished and ks_ctx SHOULD be deleted and some implementations MAY NOT create the ks_ctx.

```
struct{
    SecretRequest secret_request
    SigningRequest signing_request
}InitCertVerifyRequest

struct{
    SecretResponse secret_response
    SigningResponse signing_response
}InitCertVerifyResponse
```

secret_request and secret_response are defined in section [Section 5.1](#).

5.3. s_hand_and_app_secret

The s_hand_and_app_secret is necessary to complete the ServerHello and always follows an s_init_early_secret LURK exchange. Such sequence is guaranteed by the session_id and cookie mechanism. In case of unknown session_id or an unexpected cookie value, an invalid_request error SHOULD be returned.

The LURK client MUST ensure that PSK or PSK-ECDHE authentication has been selected via the presence of the `pre_shared_key` extension in the `ServerHello`. In addition, the selected identity MUST be the one provided in the `psk` extension of the previous `s_init_early_secret` exchange.

The LURK client MAY request the `exporter_master_secret` depending on its policies. The CS MAY ignore the request based on its policies.

Similarly to the `s_init_cert_verify`, if session resumption is not provided by the CS, the LURK session ends after this exchange and `ks_ctx` SHOULD be removed.

```
struct{
    SessionID session_id_cs
    SecretRequest secret_request
} HandAndAppRequest
```

```
struct{
    SessionID session_id_client
    SecretResponse secret_response
} HandAndAppResponse
```

5.4. s_new_tickets

`new_session_ticket` handles session resumption. It enables to retrieve `NewSessionTickets` that will be forwarded to the TLS client by the TLS server to be used later when session resumption is used. It also provides the ability to delegate the session resumption authentication from the CS to the TLS server. In fact, if the LURK client requests and receives the `resumption_master_secret` it is able to emit on its own `NewSessionTicket`. As a result `s_new_ticket` LURK exchanges are only initiated if the TLS server expects to perform session resumption and the CS responds only if if `session_resumption` is enabled. If session resumption is not enabled, the Cryptographic MAY have ended the LURK session and the `s_new_ticket` will be ignored or responded with a `invalid_request` error.

The CS MAY responds with a `resumption_master_secret` based on its policies.

The LURK client MAY perform multiple `s_new_ticket` exchanges before the session between the LURK client and the CS is in a finished state with `ks_ctx` deleted.


```

struct {
    SessionID session_id_cs
    uint8 ticket_nbr;
    uint16 key_request;
    Handshake handshake_context<0..2^32> //RFC8446 section 4.
} NewTicketRequest;

struct {
    SessionID session_id_client
    Secrets secrets
    NewSessionTicket ticket_list<0..2^16-1>; //RFC8446 section 4.6.1.
} NewTicketResponse;

crypto_service_session_id, crypto_service_cookie,
lurk_client_session_id, and lurk_client_cookie are defined in section
Section 5.3. key_request is defined in section Section 4.1.

```

ticket_nbr: designates the requested number of NewSessionTicket. In the case of delegation this number MAY be set to zero. The CS MAY responds with less tickets when the value is too high.

[6](#). LURK exchange on the TLS client

This section describes the LURK exchanges that are performed on the TLS server. The state diagram is provided in section [Section 10.1](#)

[6.1](#). c_binder_key

The c_binder_key LURK exchange is initiated when the TLS client is willing to propose a PSK for PSK or PSK-ECDHE authentication.

The handshake_context is empty as the ClientHello is under construction.

When a LURK client proposes multiple PSK, multiple binder_keys are requested.

The c_binder_key is equivalent to a secret request LURK exchange and there is no creation of a ks_ctx.

```

``` struct{ SecretRequest secret_request } BinderKeyRequest

struct{ SecretResponse secret_response } BinderKeyResponse ```

```



## **[6.2.](#) c\_init\_early\_secret**

c\_init\_early\_secret on the TLS client side works similarly as the s\_init\_early\_secret LURK exchange on the TLS server as described in section [Section 5.1](#). One key difference is that the c\_binder\_key is not requested during that LURK exchange, as a result, this LURK exchange MAY be omitted even when PSK or PSK-ECDHE authentication has been chosen by the TLS client. The c\_init\_early\_secret will only be performed in the case of 0-RTT handshake or when early exporters are required.

## **[6.3.](#) c\_hand\_secret**

The c\_hand\_secret is performed after an c\_init\_early\_secret LURK exchange. This exchange is performed in the case of an PSK or PSK-ECDHE authentication and coherence with the Handshake Context MUST be checked by the LURK client as well as by the CS as described in [Section 6.2](#) and section [Section 5.3](#).

The structures of the c\_hand\_secret follow those of the s\_hand\_and\_app\_secret described in section [Section 5.3](#).

## **[6.4.](#) c\_init\_hand\_secret**

Coherence between with the Handshake Context and the authentication (ECDHE versus PSK or PSK-ECDHE) is performed as described in section [Section 6.2](#) and section [Section 5.3](#). The LURK client and the CS MUST ensure such coherence. A Signing sub exchange MUST only be performed when ECDHE authentication has been selected which is determined by the presence of a key\_share extension as well as the absence of a pre\_shared\_key extension in the ServerHello.

Only the client\_handshake\_traffic\_secret\_0 and server\_handshake\_traffic\_secret\_0 secrets MAY be requested.



```
struct{
 SecretRequest secret_request
 select (handshake_context.ecdhe_selected){
 case :
 SigningRequest signing_request
 };
}InitHandshakeRequest
```

```
struct{
 SecretResponse secret_response
 select (handshake_context.ecdhe_selected){
 case :
 SigningResponse signing_response
 };
}InitHandshakeResponse
```

### **6.5. c\_app\_secret**

The c\_app\_secret LURK exchange is performed when no TLS client authentication has been requested, i.e. CertificateRequest message is not provided in the flight of the ServerHello. The LURK client and the CS MUST ensure no CertificateRequest is present in the Handshake Context.

Only the client\_application\_traffic\_secret\_0 and server\_application\_traffic\_secret\_0 secrets MAY be requested.

The structure follows the one of the c\_hand\_secret described in section [Section 6.3](#).

After the c\_app\_secret LURK exchange, unless the TLS client supports session resumption or post\_handshake, the LURK session is finished. The support for post\_handshake by the TLS client is indicated by the post\_handshake\_auth extension.

### **6.6. c\_cert\_verify**

The c\_cert\_verify LURK exchange is performed when TLS client authentication has been requested by the TLS server. When performed, the LURK client and the CS MUST check the presence of a CertificateRequest structure in the Handshake Context. When not present, a invalid\_handshake error SHOULD be returned.

After the c\_app\_secret LURK exchange, unless the TLS client supports session resumption or post\_handshake, the LURK session is finished.



The support for `post_handshake` by the TLS client is indicated by the `post_handshake_auth` extension.

The `CertVerifyRequest` and `CertVerifyResponse` structures are used for this LURK exchange.

```
struct{
 SessionID session_id_cs
 InitCertVerifyRequest cert_request
}CertVerifyRequest

struct{
 SessionID session_id_client
 InitCertVerifyRequest cert_response
}CertVerifyResponse
```

### **6.7. c\_register\_tickets**

The `c_register_ticket` is only used when the TLS client intend to perform session resumption. This LURK exchange has three functions. First, it is used to register the handshake in order to provide the full TLS handshake. Such information will be necessary to generate the PSK value during the future session resumptions. Second, the LURK client MAY provide one or multiple `NewSessionTickets`. These tickets will be helpful for the session resumption to bind the PSK value to some identities. Third, the LURK client MAY retrieve the `resumption_master_secret` when session resumption is being delegated by the CS to the TLS client.

The first `c_register_ticket` MUST carry the TLS handshake and future `c_register_ticket` LURK exchange MUST have a `handshake_context` of zero length. If these conditions are not met, the CS SHOULD return a `invalid_handshake` error.

The first `c_register_ticket` MAY request the `session_resumption_master`. Next `register_new_session` MUST not request that secret. If these conditions are not met, a `invalid_key_request` error is returned.

The `ticket_list` MAY have zero `NewSessionTickets` for the first `register_new_session_ticket`. Next LURK exchanges MUST have at least one `NewSessionTickets`.





```
struct {
 SessionID session_id_cs
 Handshake handshake_context<0..2^32>; //RFC8446 section 4.
 NewSessionTicket ticket_list<0..2^16-1>; //RFC8446 section 4.6.1.
 uint16 key_request;
} RegisterTicketRequest;

struct {
 SessionID session_id_client
} RegisterTicketResponse;
```

crypto\_service\_session\_id, crypto\_service\_cookie, lurk\_client\_session\_id, and lurk\_client\_cookie are defined in [section Section 5.3](#). handshake\_context is defined in [section Section 4.3](#). NewSessionTicket is defined in [\[RFC8466\] section 4.6.1](#). key\_request is defined in [Section 4.1](#).

## **6.8. c\_post\_hand**

The c\_post\_hand LURK exchange is performed in order to the client to authenticate after the TLS handshake has complete. The TLS client MUST NOT proceed to this exchange if post handshake support has not been announced in the ClientHello with the post\_handshake\_auth extension. When such extension is not found the CS MUST return a invalid\_handshake error.

```
struct {
 SessionID session_id_cs
 Handshake handshake_context<0..2^32>; //RFC8446 section 4.
 int16 app_n;
} PostHandRequest;

struct {
 SessionID session_id_client
} PostHandResponse;
```

handshake\_context is defined in [section Section 4.3](#)

app\_n: describes the number of iteration of the session keys.

## **7. Security Considerations**

Security credentials as per say are the private key used to sign the CertificateVerify when ECDHE authentication is performed as well as the PSK when PSK or PSK-ECDHE authentication is used.



The protection of these credentials means that someone gaining access to the CS MUST NOT be able to use that access from anything else than the authentication of an TLS being established. In other way, it MUST NOT leverage this for: \* any operations outside the scope of TLS session establishment. \* any operations on past established TLS sessions \* any operations on future TLS sessions \* any operations on establishing TLS sessions by another LURK client.

The CS outputs are limited to secrets as well as NewSessionTickets. The design of TLS 1.3 make these output of limited use outside the scope of TLS 1.3. Signature are signing data specific to TLS 1.3 that makes the signature facility of limited interest outside the scope of TLS 1.3. NewSessionTicket are only useful in a context of TLS 1.3 authentication.

ECDHE and PSK-ECDHE provides perfect forward secrecy which prevents past session to be decrypted as long as the secret keys that generated teh ECDHE share secret are deleted after every TLS handshake. PSK authentication does not provide perfect forward secrecy and authentication relies on the PSK remaining sercet. The Cryptographic Service does not reveal the PSK and instead limits its disclosure to secrets that are generated from the PSK and hard to be reversed.

Future session may be impacted if an attacker is able to authenticate a future session based on what it learns from a current session. ECDHE authentication relies on cryptographic signature and an ongoing TLS handshake. The robustness of the signature depends on the signature scheme and the unpredictability of the TLS Handshake. PSK authentication relies on not revealing the PSK. The CS does not reveal the PSK. TLS 1.3 has been designed so secrets generated do not disclose the PSK as a result, secrets provided by the Cryptographic do not reveal the PSK. NewSessionTicket reveals the identity (ticket) of a PSK. NewSessionTickets.ticket are expected to be public data. It value is bound to the knowledge of the PSK. The Cryptographic does not output any material that could help generate a PSK - the PSK itself or the resumption\_master\_secret. In addition, the Cryptographic only generates NewSessionTickets for the LURK client that initiates the key schedule with CS with a specific way to generate ctx\_id. This prevents the leak of NewSessionTickets to an attacker gaining access to a given CS.

If an the attacker get the NewSessionTicket, as well as access to the CS of the TLS client it will be possible to proceed to the establishment of a TLS session based on the PSK. In this case, the CS cannot make the distinction between the legitimate TLS client and teh attacker. This corresponds to the case where the TLS client is corrupted.



Note that when access to the CS on the TLS server side, a similar attack may be performed. However the limitation to a single re-use of the NewSessionTicket prevents the TLS server to proceed to the authentication.

Attacks related to other TLS sessions are hard by design of TLS 1.3 that ensure a close binding between the TLS Handshake and the generated secrets. In addition communications between the LURK client and the CS cannot be derived from an observed TLS handshake (freshness function). This makes attacks on other TLS sessions unlikely.

## **8. IANA Considerations**

## **9. Acknowledgments**

## **10. Annex**

### **10.1. LURK state diagrams on TLS client**

The state diagram sums up the LURK exchanges. The notations used are defined below:

LURK exchange indicates a LURK exchange is stated by the LURK client or is received by the CS ---> (resp. <---) indicates a TLS message is received (resp. received). These indication are informative to illustrates the TLS state machine.

CAPITAL LETTER indicates potential configuration parameters or policy applied by the LURK client or the CS. The following have been considered:

- o PSK, PSK-ECDHE, ECDHE that designates the authentication method. This choice is made by the LURK client. The choice is expressed by a specific LURK exchange as well as from the TLS Handshake Context.
- o SESSION\_RESUMPTION indicates the session resumption has been enabled on the LURK client or the CS. As a consequence the TLS client is considered performing session resumption and the TLS server MUST make session resumption possible.
- o POST\_HANDSHAKE\_AUTH indicates that post handshake authentication proposed by the TLS client in a post\_handshake\_auth extension is not ignored by the LURK client or on the CS.



Note that `SESSION_RESUMPTION`, `POST_HANDSAHKE_AUTH` are mostly informative and the current specification does not mandate to have such configuration parameters. By default, these `SHOULD` be enabled.

Other potential configuration could be proposed for configuring LURK client or CS policies. These have not been represented in the state diagram and the specification does not mandate to have these parameters implemented.

- o `CLIENT_EARLY_TRAFFIC` indicates that client early traffic `MAY` be sent by the TLS client and the notification by the TLS client in the `ClientHello` via the `early_data` extension `MUST` be considered.
- o `EARLY_EXPORTER_MASTER_SECRET` indicates whether or not `early_exporter_master_secret` `MUST` be requested by the LURK client and responded by the CS.
- o `MASTER_EXPORTER` indicates whether or not `exporter_master_secret` `MUST` be requested by the LURK client and responded by the CS.
- o `SESSION_RESUMPTION_DELEGATION` indicates whether or not `session_resumption_master` is requested by the LURK client and responded by the CS.
- o `MAX_SESSION_TICKET_NBR` indicates the maximum number of tickets that can be requested or provided by the LURK client and provided by the CS. It is strongly `RECOMMENDED` to have such limitations being configurable.

The analysis of the TLS Handshake Context enables to set some variables that can be used by the LURK client to determine which LURK exchange to proceed as well as by the CS to determine which secret `MAY` be responded. The following variables used are:

`psk_proposed`: The TLS Client is proposing PSK authentication by including a `pre_shared_key` and a `psk_key_exchange_mode` extensions in the `ClientHello`.

`dhe_proposed`: The received or to be formed `ClientHello` contains a `key_share` extensions.

`psk_accepted`: The chosen authentication method is pSK or PSK-ECDHE which is indicated via the `pre_shared_key` extension in the `ServerHello`.

`0rtt_proposed`: Indicates the TLS client supports early data which is indicated by the `early_data` extension in the `ClientHello`.





post\_handshake\_proposed: indicates the TLS client supports post handshake authentication which is indicated by the presence of a post\_handshake\_auth extension in the ClientHello.

finished: indicates that the LURK client or the CS has determined the session should be closed and ks\_ctx are deleted.

The CS contains three databases:

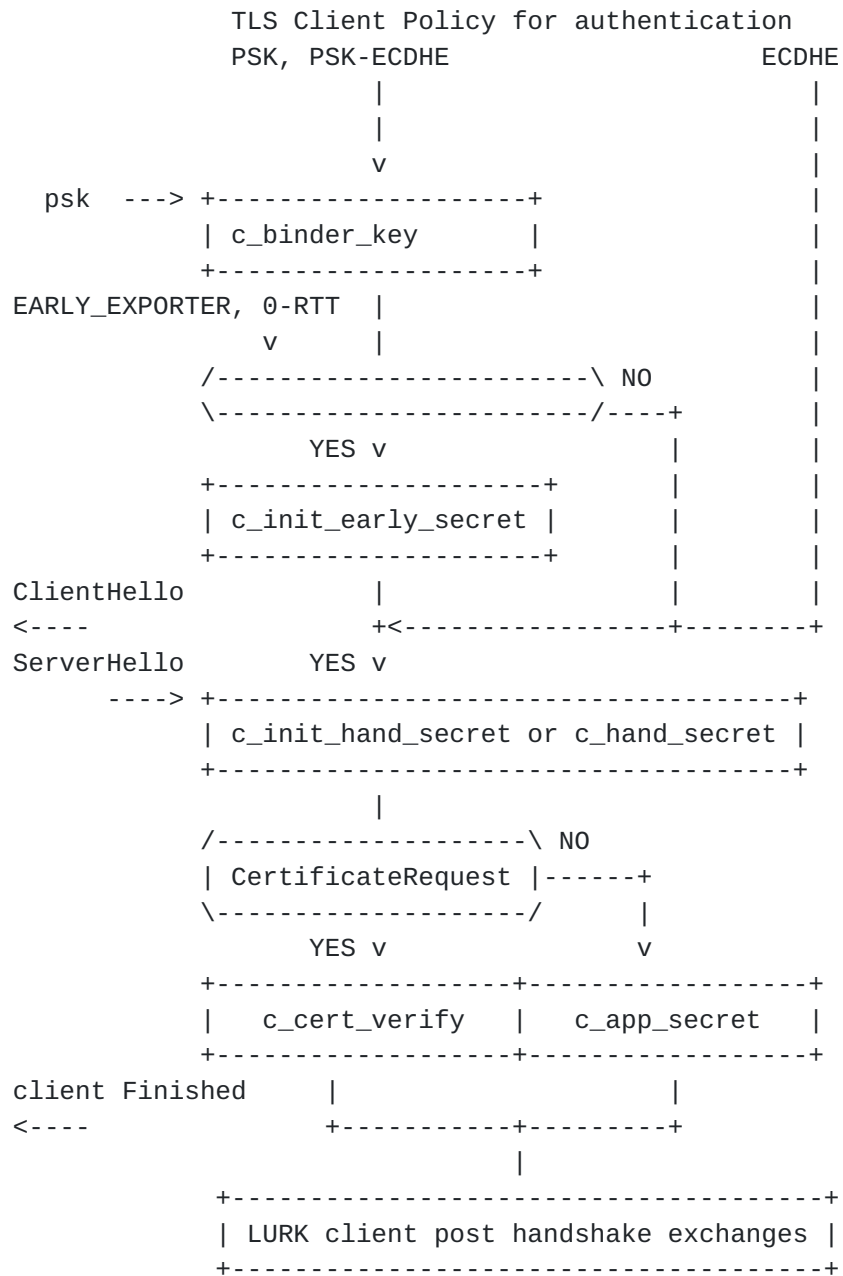
CTX\_ID\_DB: database that contains the valid ctx\_id of type opaque.

PSK\_DB: contains the list of PSKs, with associated parameters such as Hash function. This database includes the session resumption tickets.

Key\_DB: contains the asymmetric signing keys with supported signing algorithms.

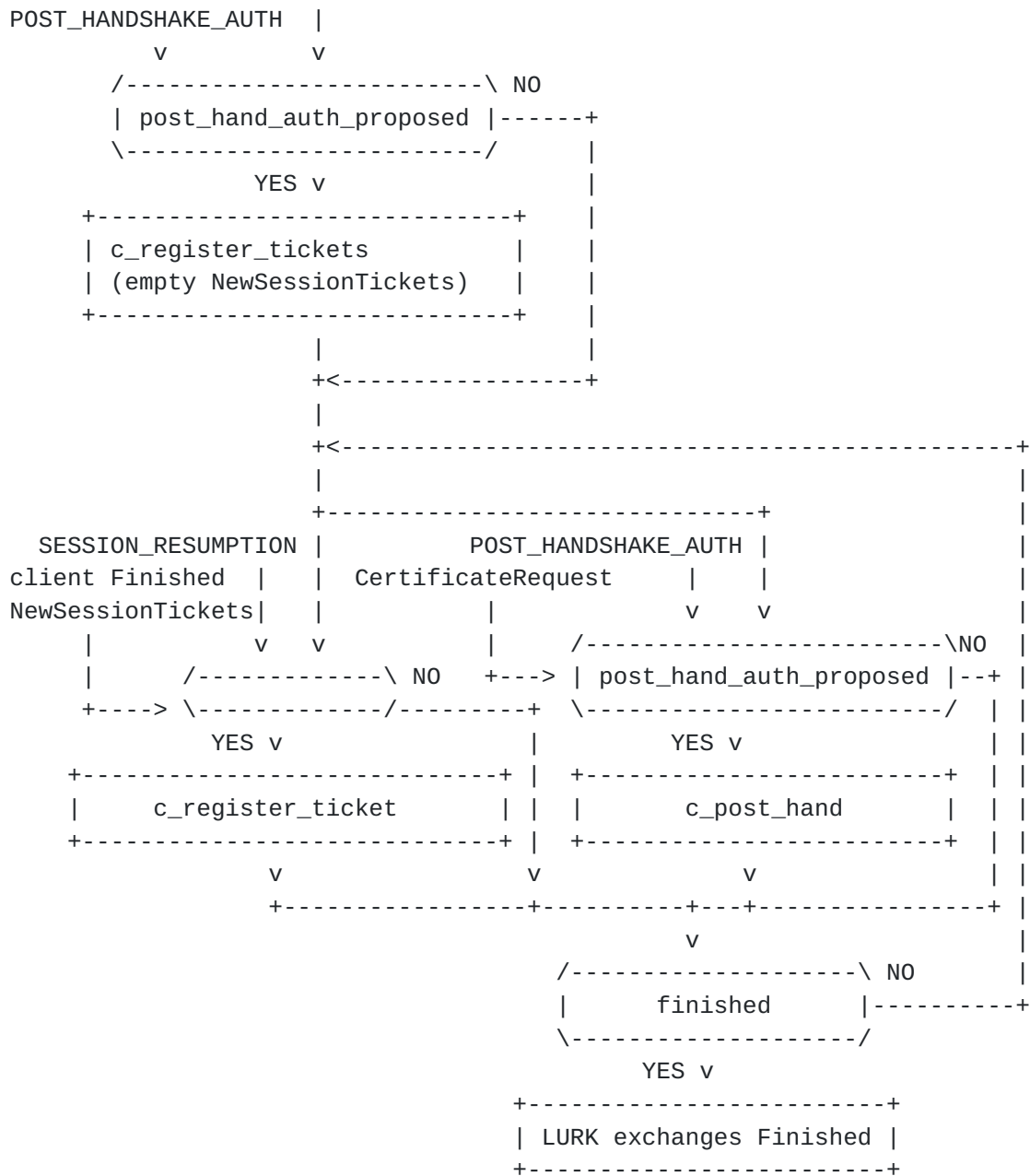
#### [10.1.1.](#) LURK client





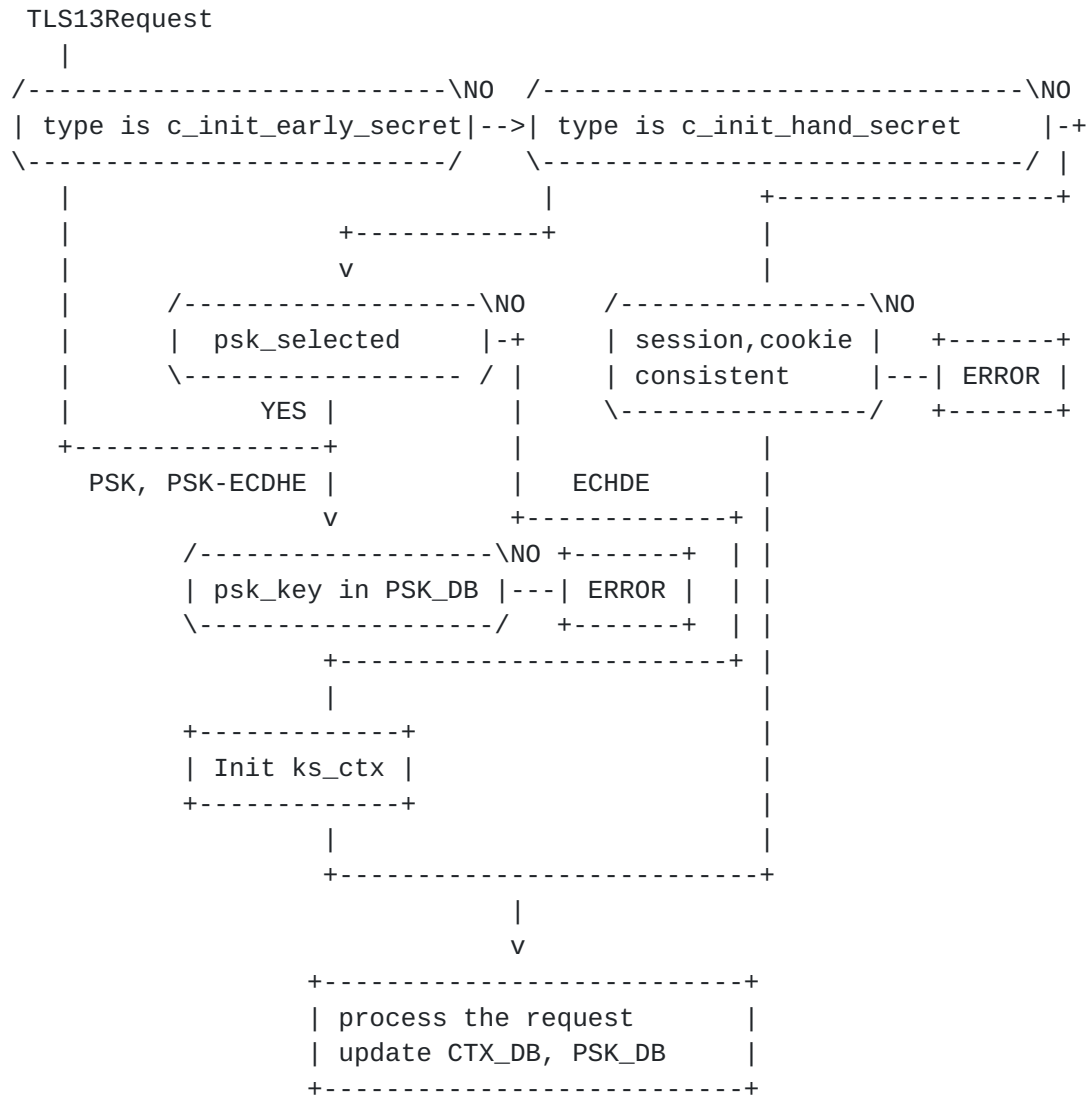
The LURK client post handshake diagram is represented below:





### 10.1.2. Cryptographic Service



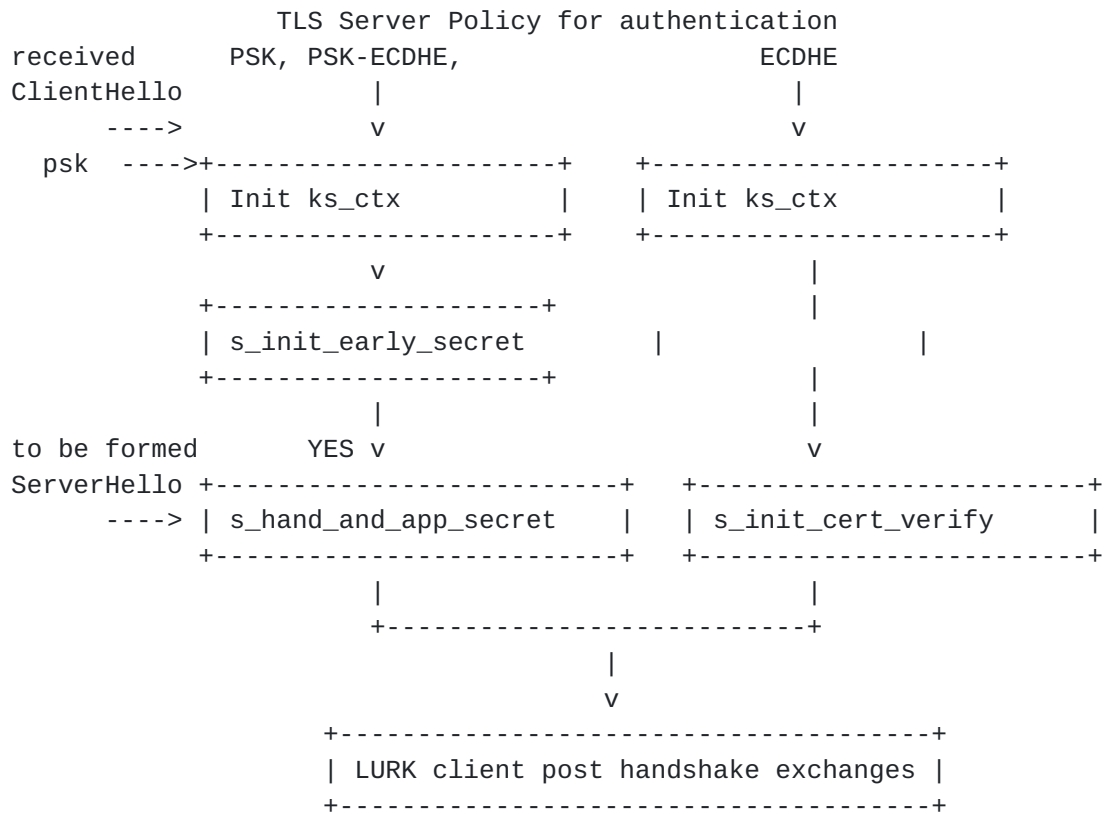


## [10.2.](#) LURK state diagrams on TLS server

### [10.2.1.](#) LURK client







### [10.2.2.](#) Cryptographic Service





### 10.3. TLS handshakes with Cryptographic Service

This section is non normative. It illustrates the use of LURK in various configurations.

The TLS client may propose multiple ways to authenticate the server (ECDHE, PSK or PSK-ECDHE). The TLS server may chose one of those, and this choice is reflected by the LURK client on the TLS server. In other words, this decision is out of scope of the CS.

The derivation of the secrets is detailed in [RFC8446](#) section 7.1. Secrets are derived using Transcript-Hash and HKDF, PSK and ECDHE secrets as well as some Handshake Context.

The Hash function: When PSK or PSK-ECDHE authentication is selected, the Hash function is a parameter associated to the PSK. When ECDHE, the hash function is defined by the cipher suite algorithm negotiated. Such algorithm is defined in the cipher\_suite extension provided in the ServerHello which is provided by the LURK client in the first request when ECDHE authentication is selected.

PSK secret: When PSK or PSK-ECDHE authentication is selected, the PSK is the PSK value identified by the identity. When ECDHE



authentication is selected, the PSK takes a default value of string of Hash.length bytes set to zeros.

ECDHE secret: When PSK or PSK-ECDHE authentication is selected, the ECDHE secret takes the default value of a string of Hash.length bytes set to zeros. The Hash is always known as a parameter associated to the selected PSK. When ECDHE authentication is selected, the ECDHE secret is generated from the secret key (ephemeral\_secret) provided by the LURK client and the counter part public key in the key\_share extension. When the LURK client is on the TLS client, the public key is provided in the ServerHello. When the LURK client is on the TLS Server, the public key is provided in the ClientHello. When ECDHE secret is needed, ClientHello...ServerHello is always provided to the CS.

Handshake Context: is a subset of Handshake messages that are necessary to generate the requested secrets. The various Handshake Contexts are summarized below:

Key Schedule secret or key	Handshake Context
binder_key	None
client_early_traffic_secret	ClientHello
early_exporter_master_secret	ClientHello
client_handshake_traffic_secret	ClientHello...ServerHello
server_handshake_traffic_secret	ClientHello...ServerHello
client_application_traffic_secret_0	ClientHello...server Finished
server_application_traffic_secret_0	ClientHello...server Finished
exporter_master_secret	ClientHello...server Finished
resumption_master_secret	ClientHello...client Finished

The CS has always the Hash function, the PSK and ECDHE secrets and the only remaining parameter is the Handshake Context. The remaining sections will only focus on checking the Handshake Context available to the CS is sufficient to perform the key schedule.

When ECDHE authentication is selected both for the TLS server or the TLS client, a CertificateVerify structure is generated as described in [\[RFC8446\] section 4.4.3](#). CertificateVerify consists in a signature over a context that includes the output of Transcript-Hash(Handshake Context, Certificate) as well as a context string. Both Handshake Context and context string depends on the Mode which is set to server in this case via the configuration of the LURK server. Similarly to the key schedule, the Hash function is defined by the PSK or the ServerHello. The values for the Handshake Context are represented below:



Mode	Handshake Context	Base Key
Server	ClientHello ... later of EncryptedExtensions/ CertificateRequest	server_handshake_traffic_ secret
Client	ClientHello ... later of server Finished/EndOfEarlyData	client_handshake_traffic_ secret
Post- Handshake	ClientHello ... client Finished + CertificateRequest	client_application_traffic_ secret_N

When ECDHE authentication is selected, the CS generates a Finished message, which is a MAC over the value Transcript-Hash(Handshake Context, Certificate, CertificateVerify) using a MAC key derived from the Base Key. As a result, the same Base Key and Handshake Context are required for its computation describe in [\[RFC8466\] section 4.4.4..](#)

#### [10.4.](#) TLS 1.3 ECDHE Full Handshake

This example illustrates the case of a TLS handshake where the TLS server is authenticated using ECDHE only, that is not PSK or PSK-ECDHE authentication is provided and so session resumption is provided either.

##### [10.4.1.](#) TLS Client: ClientHello

The TLS client does not provides any PSK and omits the `pre_shared_key` as well as the `psk_key_exchange_mode` extensions. Note that omitting the `psk_key_exchange_mode` extension prevents the TLS client to perform further session resumption.

The TLS client does not need any interaction with the Cryptographic Service to generate and send the ClientHello message to the TLS server.

TLS Client

TLS Server

```

Key ^ ClientHello
Exch | + key_share
 v + signature_algorithms ----->

```





#### 10.4.2. TLS Server: ServerHello

Upon receiving the ClientHello, the TLS server determines the TLS client requests an ECDHE authentication. The TLS server initiates a LURK session to provide ECDHE authentication as represented below:

TLS Client	TLS Server
	ServerHello ^ Key
	+ key_share   Exch
	{EncryptedExtensions} ^ Server
	{CertificateRequest*} v Params
	{Certificate} ^
	{CertificateVerify}   Auth
	{Finished} v
	<----- [Application Data*]

The LURK Client on the TLS server initiates a `s_init_cert_verify` to retrieve the necessary secrets to finish the exchange and request the generation of the signature (`certificate_verify`) carried by the `CertificateVerify` TLS structure.

The `s_init_cert_verify` request uses a `InitCertVerifyRequest` structure which is composed of two substructures: A `SecretRequest` structure (`secret_request`) is in charge of requesting the necessary secrets to decrypt and encrypt the TLS handshake as well as the applications carried over the TLS session. Finally a `SigningRequest` substructure (`signing_request`) is used to request the `certificate_verify` payload.

The `secret_request` carries the requested secrets as well as the necessary parameters to generate the secrets. In our case, the requested secrets are the handshake secrets (`h_c`, `h_s`) as well as the application secrets (`a_c`, `a_s`). This corresponds to the most expected use cases, though other use case may require different secrets to be requested. These requests are indicated in the `key_request`. The necessary Handshake Context is provided through `handshake_context` which is set to `ClientHello ...`

`EncryptedExtensions`. The ECDHE shared secret is provided in this example via the ephemeral extension. In our case, the secret key is provided directly though other means may be used. In particular providing the secret key implies the `dhe` parameters have been generated outside the CS. The freshness function is provided through the freshness extension.

The `signing_request` provides the `key_id` that identifies the private key used to generate the signature, the algorithm use to generate the signature (`sig_algo`) as well as the certificate. The certificate carries information to generate the `Certificate` structure of the



ServerHello, and may not be the complete certificate chain but only an index.

TLS Server

Lurk Client

CS

```

InitCertVerifyRequest
 secret_request
 key_request = h_c, h_s, a_c, a_s
 handshake_context = ClientHello ... EncryptedExtensions
 ext
 ephemeral = dhe_secret
 freshness
 session_id
 signing_request
 key_id,
 sig_algo
 certificate
----->
InitCertVerifyResponse
 secret_response
 keys
 ext
 session_id
 signing_response
 certificate_verify
<-----

```

Upon receiving the InitCertificateRequest, the CS initiates a context associated to the newly created LURK session.

The secrets are generated from the TLS 1.3 key schedule describe in [\[RFC8446\]](#) and requires as input PSK, ECDHE as well as some context handshake.

The CS determine that ECDHE without specific PSK is used from the ClientHello and associated extensions. As a result, the default PSK value is used. The ECDHE share secret is derived, in our case from the dhe\_secret of the TLS server and the public dhe value provided by the ClientHello shared\_key extension.

The CS reads the freshness extension and generates the handshake\_context that will be used further.

The necessary Handshake Context to generate the handshake secrets is ClientHello...ServerHello which is provided by the handshake\_context. The CS uses the freshness function provided in the freshness extension to derive the appropriated server.random.



The generation of the CertificateVerify is described in [\[RFC8446\]](#) [section 4.4.3](#). and consists in a signature over a context that includes the output of Transcript-Hash(Handshake Context, Certificate) as well as a context string. Both Handshake Context and context string depends on the Mode which is set to server in this case via the configuration of the LURK server.

The necessary Handshake Context to generate the CertificateVerify is ClientHello ... later of EncryptedExtensions / CertificateRequest. In our case, this is exactly handshake\_context, that is ClientHello ... EncryptedExtensions. The Certificate payload is generated from the information provided in the certificate extension.

Once the certificate\_verify value has been defined, the LURK server generates the server Finished message in order to have the necessary Handshake Context ClientHello...server Finished to generate the application secrets.

The LURK server returns the requested keys, the certificate\_verify in a InitCertVerifyResponse structure. This structure is composed of the two substructures: SecretResponse that contains the secrets and SigningResponse that contains the certificate\_verify.

The TLS server can complete the ServerHello response, that is proceed to the encryption and generates the Finished message.

As session resumption is not provided, the LURK server goes into a finished state and delete the ks\_ctx. The special case described in this session does not use LURK session and as such may be stateless.

#### **[10.4.3](#). TLS client: client Finished**

Upon receiving the ServerHello message, the TLS client retrieve the handshake and application secrets to decrypt the messages received from server as well as to encrypt its own messages and application data as represented below:

TLS Client		TLS Server
{Finished}	----->	
[Application Data]	<----->	[Application Data]

To retrieves these secrets, the TLS client proceeds successively to an c\_init\_hand\_secret LURK exchange followed by a c\_app\_secret LURK exchange.

The c\_init\_hand\_secret exchange is composed of one substructure: (secret\_request) to request the secrets. Optionally, a



SigningRequest (signing\_request) when the TLS server requests the TLS client to authenticate itself. The indication of a request for TLS client authentication is performed by the TLS server by providing a CertificateRequest message associated to the ServerHello. We consider that such request has not been provided here so the SigningRequest structure is not present.

The secret\_request specifies the secrets requested via the key\_request. In our case only the handshake secrets are requested (h\_c, h\_s). In this example the ECDHE share secret is provided via the ephemeral extension. In this case the ECDHE secrets have been generated by the TLS client, and the TLS client chooses to provide the ephemeral secret (dhe\_secret) to the CS via the ephemeral extension. The TLS client also provides the freshness function via the freshness extension so the handshake\_context can be appropriately be interpreted. The handshake context is provided via the handshake\_context and is set to ClientHello ... ServerHello.

Note that if the TLS client would have like the CS to generate the ECDHE public and private keys, the generation of the keys would have been made before the ClientHello is sent, that is in our case during a c\_init\_early\_secret LURK exchange. If that had been the case a c\_hand\_secret LURK exchange would have followed and not a c\_init\_hand\_secret exchange.





```

TLS Client
Lurk Client
 InitHandshakeSecretRequest
 secret_request
 key_request = h_c, h_s
 handshake_context = ClientHello ... ServerHello
 ext
 ephemeral = dhe_secret
 freshness
 session_id
 ----->

 InitHandshakeSecretResponse
 secret_response
 ext
 session_id
 <----- keys

TLS Client
Lurk Client
 AppSecretRequest
 session_id
 cookie
 secret_Request
 key_request
 handshake_context
 ----->

 AppSecretResponse
 session_id
 cookie
 secret_response
 <----- keys

```

Upon receiving the `InitHandshakeSecretRequest`, the servers initiates a LURK session context (`ks_ctx`) and initiates a key schedule. The key schedule requires PSK, ECDHE as well as Handshake Context to be complete. As no `pre_shared_key` and `psk_key exchange_modes` are found in the `ClientHello` the CS determines that ECDHE is used for the authentication. The PSK is set to its default value. The ECHDE shared secret is generated from the ephemeral extension as well as the public value provided in the `ClientHello`. The CS takes the freshness function and generates the appropriated handshake context. The necessary Handshake Context to generate handshake secrets is `ClientHello...ServerHello` which is provided by the `handshake_context`.

The handshake secrets are returned in the `secret_response` to the TLS client. The TLS client decrypt the encrypted extensions and messages of the `ServerHello` exchange.



As no CertificateRequest appears, the LURK client initiates an app\_secret LURK exchange decrypt and encrypt application data while finishing the TLS handshake.

The AppSecretRequest structure uses session\_id and cookies as agreed in the previous c\_init\_hand\_secret exchange. The AppSecretRequest embeds a SecretRequest sub structure. The application secrets requested are indicated by the key\_request (a\_s, a\_s). The Handshake Context (handshake\_context) is set to server EncryptedExtensions ... server Finished.

Upon receiving the AppSecretRequest, the CS checks the session\_id. The CS has now the ClientHello ... server Finished which enables it to compute the application secrets.

As no session resumption is provided, the CS and the LURK client goes into a finished state and delete their ks\_ctx.

### **10.5. TLS 1.3 Handshake with session resumption**

This scenario considers that the TLS server is authenticated using ECDHE only in the first time and that further TLS handshake use the session resumption mechanism. The first TLS Handshake is very similar as the previous one. The only difference is that psk\_key\_exchange\_mode extension is added to the ClientHello. However, as no PSK identity is provided, the Full exchange is performed as described in section [Section 10.4](#).

The only change is that session resumption is activated, and thus LURK client and LURK servers do not go in a finished state and close the LURK session after the exchanges are completed. Instead further exchanges are expected. Typically, on the TLS server side new\_Session\_ticket exchanges are expected while registered\_session\_ticket are expected on the client side.

When session resumption is performed, a new LURK session is initiated.

#### **10.5.1. Full Handshake**

The Full TLS Handshake use ECDHE authentication. It is very similar to the logic described in section [Section 10.4](#). The TLS handshake is specified below for convenience.



```

TLS Client TLS Server

Key ^ ClientHello
Exch | + key_share
 | + psk_key_exchange_mode
 v + signature_algorithms ----->
 ServerHello ^ Key
 + key_share | Exch
 {EncryptedExtensions} Server Param
 {Certificate} ^
 {CertificateVerify} | Auth
 {Finished} v
 <----- [Application Data*]
{Finished} ----->
[Application Data] <-----> [Application Data]

```

#### [10.5.2.](#) TLS server: NewSessionTicket

As session resumption has been activated by the `psk_key_exchange_mode`, the TLS Server is expected to provide the TLS client NewSessionTickets as mentioned below:

```

TLS Client TLS Server
 <----- [NewSessionTicket]

```

The LURK client and LURK server on the TLS server does not go into a finished state. Instead, the LURK client continues the LURK session with a NewTicketRequest to enable the CS to generate the `resumption_master_secret` necessary to generate the PSK and generate a NewTicketSession. `ticket_nbr` indicates the number of NewSessionTickets and `handshake_context` is set to earlier of client Certificate client CertificateVerify ... client Finished. As we do not consider TLS client authentication, the `handshake_context` is set to client Finished as represented below.

```

TLS Server Cryptographic Service
Lurk Client
 NewTicketRequest
 session_id
 cookie
 ticket_nbr
 handshake_context=client Finished ----->
 NewTicketResponse
 session_id
 cookie
 <----- tickets

```



The necessary Handshake Context to generate the `resumption_master_secret` is `ClientHello...client Finished`. From the `InitCertificateVerify` the `context_handshake` was set to `ClientHello...server Finished`. The additional `handshake_context` enables the CS to generate the `NewSessionTickets`.

Note that the LURK client on the TLS server may send multiple `NewTicketRequest`. Future request have an empty `handshake_context`.

Upon receiving the `NewTicketRequest`, the LURK server checks the `session_id` and `cookie`. It then generates the `resumption_master_secret`, `NewSessionTickets`. `NewSessionTickets` are stored into the `PSK_DB` under `NewSessionTicket.ticket`. Note that `PSK` is associated with the authentication mode as well as the Hash function negotiated for the cipher suite. The CS responds with `NewSessionTickets` that are then transmitted back to the TLS client. The TLS server is ready for session resumption.

### [10.5.3](#). TLS client: `NewSessionTicket`

Similarly, the LURK client on the TLS client will have to provide sufficient information to the CS the necessary `PSK` can be generated in case of session resumption. This includes the remaining Handshake Context to generate the `resumption_master_secret` as well as `NewSessionTickets` provided by the TLS server. The LURK client uses the `c_register_ticket` exchange.

Note that the LURK client may provide the `handshake_context` with an empty list of `NewSessionTickets`, and later provide the `NewSessionTickets` as they are provided by the TLS server. The Handshake Context only needs to be provided for the first `RegisterTicketRequest`.

TLS Client	
Lurk Client	Cryptographic Service
<code>NewTicketRequest</code>	
<code>session_id</code>	
<code>cookie</code>	
<code>handshake_context=client Finished</code>	
<code>ticket_list</code>	----->
	<code>NewTicketResponse</code>
	<code>session_id</code>
	<code>cookie</code>
	<----- tickets

Both TLS client and TLS Servers are ready for further session resumption. On both side the CS stores the `PSK` in a database designated as `PSK_DB`. Each `PSK` is associated to a Hash function as





well as authentication modes. Each PSK is designated by an identity. The identity may be a label, but in our case the identity is derived from the `NewSessionTicket.ticket`.

#### **10.5.4. Session Resumption**

Session resumption is initiated by the TLS client. Session resumption is based on PSK authentication and different PSK may be proposed by the TLS client. The TLS handshake is presented below.

TLS Client		TLS Server
ClientHello		
+ key_share		
+ psk_key_exchange_mode		
+ pre_shared_key	----->	
		ServerHello
		+ pre_shared_key
		+ key_share
		{EncryptedExtensions}
		{Finished}
	<-----	[Application Data*]

The TLS client may propose to the TLS Server multiple PSKs. Each of these PSKs is associated a `PskBindersEntry` defined in [\[RFC8446\]](#) [section 4.2.11.2](#). `PskBindersEntry` is computed similarly to the Finished message using the `binder_key` and the partial ClientHello.

The TLS server is expected to pick a single PSK and validate the binder. In case the binder does not validate the TLS Handshake is aborted. As a result, only one `binder_key` is expected to be requested by the TLS server as opposed to the TLS client.

In this example we assume the `psk_key_exchange_mode` indicated by the TLS client supports PSK-ECDHE as well as PSK authentication. The presence of a `pre_shared_key` and a `key_share` extension in the ServerHello indicates that PSK-ECDHE has been selected.

##### **10.5.4.1. TLS client: ClientHello**

To compute binders, the TLS Client needs to request the `binder_key` associated to each proposed PSK. These `binder_keys` are retrieved to the CS using the `BinderKeyRequest`. The `key_request` is set to `binder_key`, and the `PSK_id` extension indicates the PSK's identity (`PSKIdentity.identity` or `NewSessionTicket.ticket`). No Handshake Context is needed and `handshake_context` is empty.



```

TLS Client
Lurk Client
 BinderKeyRequest
 key_request=binder_key
 handshake_context=""
 ext
 PSK_id
 BinderKeyResponse
 <----- key

```

Upon receiving the BinderKeyRequest, the CS checks the psk is in the PSK\_DB and returns the binder\_key.

With the binder keys, the TLS Client is able to send it ClientHello message.

We assume in this example that the ECDHE secrets is generated by the TLS client and not the Cryptographic service. As a result, the TLS client does not need an extra exchange to request the necessary parameters to derive the key\_shared extension.

#### **10.5.4.2. TLS server: ServerHello**

The TLS server is expected to select a PSK, check the associated binder and proceed further. If the binder fails, it is not expected to proceed to another PSK, as a result, the TLS server is expected to initiates a single LURK session.

The binder\_key is requested by the TLS server via and s\_init\_early\_secret LURK exchange. The InitEarlySecretRequest structure is composed of a SecretRequest structure (secret\_request).

In our case, only the binder\_key is requested so key\_request is set to binder\_key only. Similarly, to the TLS client, the handshake\_context is not needed to generate the binder\_key. However, the EarlySecret exchange requires the ClientHello to be provided so early secrets may be computed in the same round during 0-RTT handshake. The chosen PSK is indicated in the PSK\_id extension and the freshness function is indicated in the freshness extension.



```

TLS Server
Lurk Client
 InitEarlySecretRequest
 secret_Request
 key_request=binder_key
 handshake_context=ClientHello
 ext
 freshenss
 PSK_id
 session_id

 Cryptographic Service
 InitEarlySecretResponse
 secret_response
 <----- key
 ext
 session_id

```

To complete to the ServerHello exchange, the TLS server needs the handshake and application secrets. These secrets are requested via an s\_hand\_and\_app\_secret LURK exchange. The HandshakeAndAppSecretRequest is composed of SecretRequest structure. The key\_request is set to handshake (h\_c, h\_s) and application secrets (a\_s, a\_c). The Handshake Context (handshake\_context) is set to ServerHello ... EncryptedExtensions as their is no authentication of the TLS client. Finally, the ephemeral ECDHE is provided or requested via the ephemeral extension. In our case, we assume the ephemeral secrets is generated by the tLS client is provided to the CS.

The necessary Handshake Context to generate the handshake secrets is ClientHello ... ServerHello, so the CS can generate the handshake secrets. The necessary Handshake Context to generate the application secrets is ClientHello ... server Finished. So the CS needs to generate the Finished message before as in the case of the InitCerificateVerify exchange detailed in [Section 10.5.1](#).



```

TLS Server
Lurk Client
 HandshakeAndAppRequest
 session_id
 secret_request
 key_request = h_c, h_s, a_c, a_s
 handshake_context = ServerHello ... EncryptedExtensions
 ext
 ephemeral = dhe_secret
 ----->
 HandshakeAndAppResponse
 session_id
 secret_response
 keys
 <-----

```

The CS returns the necessary secret to the TLS server to complete the ServerHello response.

The remaining of the TLS handshake is proceeded similarly as described in the Full Handshake in section [Section 10.5](#).

#### **[10.6](#). TLS 1.3 0-RTT handshake**

The 0-RTT Handshake is a PSK or PSK-ECDHE authentication that enables the TLS client to provide application data during the first round trip. The main differences to the PSK PSK-ECDHE authentication described in the case of session resumption is that:

- o Application Data is encrypted in the ClientHello based on the client\_early\_secret
- o Generation of the client\_early\_secret requires the Cryptographic Service to be provisioned with the ClientHello which does not need to be re-provisioned later to generate the handshake secrets
- o An additional message EndOfEarlyData needs to be considered to compute the client Finished message.





## TLS Client

## TLS Server

```

ClientHello
+ early_data
+ key_share*
+ psk_key_exchange_modes
+ pre_shared_key
(Application Data*) ----->

 ServerHello
 + pre_shared_key
 + key_share*
 {EncryptedExtensions}
 + early_data
 {Finished}
 [Application Data*]
 <-----
(EndOfEarlyData)
{Finished}
[Application Data] ----->
[Application Data] <-----> [Application Data]

```

**10.6.1. TLS client: ClientHello**

With 0-RTT handshake, the TLS client builds binders as in session resumption described in section [Section 10.5.4](#). The binder\_key is retrieved for each proposed PSK with a BinderKeyRequest. When early application data is sent it is encrypted using the client\_early\_traffic\_secret. This secret is retrieved using the c\_init\_early\_secret LURK exchange.

The InitEarlySecretRequest is composed of a SecretRequest (secret\_request) substructure. The TLS Client sets the key\_request to client\_early\_traffic\_secret (e\_s). The handshake is set to ClientHello. The PSK is indicated via the the PSK\_id extension, the freshness function is indicated via the freshness extension. If the TLS client is willing to have the ECDHE keys generated by the CS an ephemeral extension MAY be added also.

When multiple PSK are proposed by the TLS client, the first proposed PSK is used to encrypt the application data.





Upon receiving the `InitEarlySecretRequest`, the CS generates the `client_early_traffic_secret`.

The TLS client is able to send its `ClientHello` with associated binders and application data.

#### **10.6.2. TLS server: `ServerHello`**

If the TLS server accepts the early data. It proceeds as described in session resumption described in [Section 10.5.4](#). In addition to the `binder_key`, the TLS server also request the `client_early_traffic_secret` to decrypt the early data as well as to proceed to the `ServerHello` exchange.

#### **10.6.3. TLS client: `Finished`**

The TLS client proceeds as described in handshake based on ECDHE, PSK or PSK-ECDHE authentications described in [Section 10.4](#) and [Section 10.5](#). The main difference is that upon requesting handshake and application secrets, using an `HandAndAppRequest` the TLS client will not provide the `ClientHello` as part as the `handshake_context`. The Client as already been provided during the `EarlySecret` exchange.

#### **10.7. TLS client authentication**

TLS client authentication can be performed during the Full TLS handshake or after the TLS handshake as a post handshake authentication. In both cases, the TLS client authentication is initiated by the TLS server sending a `CertificateRequest`. The authentication is performed via a `CertificateVerify` message generated by the TLS client but such verification does not involve the CS on the TLS server.



### **10.8. TLS Client:Finished (CertificateRequest)**

The ServerHello MAY carry a CertificateRequest encrypted with the handshake secrets.

Upon receiving the ServerHello response, the TLS client decrypts the ServerHello response. If a CertificateRequest message is found, the TLS Client requests the Cryptographic to compute the CertificateVerify in addition to the application secrets via a certificate\_verify LURK exchange. The CertVerifyRequest is composed of a Secret Request structure and a SigningRequest structure.

The key\_request is set to the application secrets (a\_c, a\_s) and the handshake\_context is set to server EncryptedExtensions ... later of server Finished/EndOfEarlyData. As the request follows a (BinderKey, EarlySecret, HandshakeSecret) or HandshakeSecret the Handshake Context on the CS now becomes: ClientHello ... later of server Finished/EndOfEarlyData which is the Handshake Context required to generate the CertificateVerify on the TLS client side and includes the Handshake Context required to generate the application secrets (ClientHello...server Finished).

TLS Client

Lurk Client

Cryptographic Service

CertVerifyRequest

session\_id

secret\_request

key\_request

handshake\_context = EncryptedExtensions ...

later of server Finished/EndOfEarlyData

signing\_request

CertVerifyResponse

session\_id

secret\_response

keys

signing\_response

<----- certificate\_verify

Upon receiving the CertificateRequest, the CS checks the session\_id and cookie.

### **10.9. TLS Client Authentication (PostHandshake)**

When post-handshake is enabled by the TLS client, the TLS client may receive at any time after the handshake a CertificateRequest message. When post handshake is enabled by the TLS client, as soon as the client Finished message has been sent, the TLS client sends a RegisteredNewSessionTicketRequest with an empty NewSessionTicket to



register the remaining Handshake Context to the CS. ctx\_id is set to opaque, handshake\_context is set to earlier of client Certificate client CertificateVerify ... client Finished.

Upon receiving the RegisteredNewSessionTicketsRequest the Cryptographic is aware of the full Handshake Context. It updates ks\_ctx.next\_request to c\_post\_hand or c\_register\_ticket.

TLS Client

```

Lurk Client Cryptographic Service
 RegisteredNewSessionTicketRequest
 session_id
 handshake_context
 ticket_list (empty)
 <----- RegisteredNewSessionTicketResponse
 session_id
 cookie

```

When the TLS client receives a CertificateRequest message from the TLS server, the TLS client sends a PostHandshakeRequest to the Cryptographic Service to generate certificate\_verify. The handshake\_context is set to CertificateRequest. The index N of the client\_application\_traffic\_N key is provided as well as the Cryptographic so it can generate the appropriated key.

TLS Client

```

Lurk Client Cryptographic Service
 PostHandshakeRequest
 session_id
 handshake_context=CertificateRequest
 app_n=N
 PostHandshakeResponse
 session_id
 <----- certificate_verify

```

Upon receiving the PostHandshakeRequest the CS checks session\_id and cookie. The necessary Handshake Context to generate the certificate\_verify is ClientHello ... client Finished + CertificateRequest. Once the PostHandshakeResponse. Next requests expected are c\_post\_hand or c\_register\_ticket.

## [11.](#) References

### [11.1.](#) Normative References





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## **Author's Address**

Daniel Migault  
Ericsson  
8275 Trans Canada Route  
Saint Laurent, QC 4S 0B6  
Canada

EMail: [daniel.migault@ericsson.com](mailto:daniel.migault@ericsson.com)

