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## **End-to-End Object Encryption for the Extensible Messaging and Presence Protocol (XMPP)**

### **draft-miller-3923bis-02**

#### **Abstract**

This document defines a method of end-to-end object encryption for the Extensible Messaging and Presence Protocol (XMPP). The protocol defined herein is a simplified version of the protocol defined in RFC 3923.

#### **Status of this Memo**

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## 1. Introduction

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End-to-end encryption of traffic sent over the Extensible Messaging and Presence Protocol [\[XMPP-CORE\] \(Saint-Andre, P., "Extensible Messaging and Presence Protocol \(XMPP\): Core," May 2010.\)](#) is a desirable goal. Requirements and a threat analysis for XMPP encryption are provided in [\[E2E-REQ\] \(Saint-Andre, P., "Requirements for End-to-End Encryption in the Extensible Messaging and Presence Protocol \(XMPP\)," March 2010.\)](#). Many possible approaches to meet those (or similar) requirements have been proposed over the years, including methods based on PGP, S/MIME, SIGMA, and TLS.

The S/MIME approach defined in [\[RFC3923\] \(Saint-Andre, P., "End-to-End Signing and Object Encryption for the Extensible Messaging and Presence Protocol \(XMPP\)," October 2004.\)](#) has never been implemented in XMPP clients to the best of our knowledge, but has some attractive features, especially the ability to store-and-forward an encrypted message at a user's server if the user is not online when the message is received (in the XMPP community this is called "offline storage" and the message is referred to as an "offline message"). The authors surmise that RFC 3923 has not been implemented mainly because it adds several new dependencies to XMPP clients, especially MIME (along with the CPIM and MSGFMT media types). Therefore this document explores the possibility of an approach that is similar to but simpler than RFC 3923, while retaining the same basic object encryption model.

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## 2. Terminology

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This document inherits terminology defined in [\[XMPP-CORE\] \(Saint-Andre, P., "Extensible Messaging and Presence Protocol \(XMPP\): Core," May 2010.\)](#).

Security-related terms are to be understood in the sense defined in [\[SECTERMS\] \(Shirey, R., "Internet Security Glossary, Version 2," August 2007.\)](#).

The capitalized key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [\[KEYWORDS\] \(Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels," March 1997.\)](#).

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## 3. Securing XMPP Stanzas

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The process that a sending agent follows for securing stanzas is the same regardless of the form of stanza (i.e., <iq/>, <message/>, or <presence/>).

1. Constructs a cleartext version of the stanza, S.
2. Generates a session key R appropriate for the intended block cipher (e.g. AES-SHA-256).
3. Notes the current UTC date and time N when this stanza is constructed, formatted as described under [Section 6 \(Inclusion and Checking of Timestamps\)](#).
4. Converts the stanza to a UTF-8 encoded string, optionally removing line breaks and other insignificant whitespace between elements and attributes, i.e., S' = UTF8-encode(S). We call S' a "stanza-string" because for purposes of encryption and decryption it is treated not as XML but as an opaque string (this avoids the need for complex canonicalization of the XML input).

5. Constructs a plaintext envelope (E) <plain/> as follows:

\*The attribute 'timestamp' set to the UTC date and time value  
N

\*The XML character data set to the base64-encoded form of S'  
(where the encoding adheres to the definition in Section 4  
of [\[BASE64\] \(Josefsson, S., "The Base16, Base32, and Base64  
Data Encodings," October 2006.\)](#) and where the padding bits  
are set to zero). This encoding is necessary to preserve a  
canonicalized form of S'.

6. Converts the envelope (E) to a UTF-8 encoded string, optionally  
removing line breaks and other insignificant whitespace between  
elements and attributes, i.e., E' = UTF8-encode(E).

7. Encrypts the UTF8-encoded enveloped (E') using the intended  
block cipher, i.e. T = block-encrypt(R, E').

8. Generates a message authentication code (MAC) with a  
cryptographic hashing algorithm (e.g. HMACSHA256) using the  
encrypted data T as the salt and the session block cipher key R  
as the message, i.e., M = mac-hash(T, R)

9. Encrypts the session key (R) using the recipient's public key  
to produce encrypted data K. (Known issue: This step is under-  
specified and will be expanded in a later version of this  
document.)

10. Constructs an <e2e/> element qualified by the  
"urn:ietf:params:xml:ns:xmpp-objenc:0" namespace as follows:

\*The child element <key/> (implicitly qualified by the  
"urn:ietf:params:xml:ns:xmpp-objenc:0" namespace) as  
follows:

-The attribute 'cipher-algo' set to the asynchronous  
encryption scheme used in step 9;

-The XML character data set to the base64-encoded form of  
K.

\*The child element `<data/>` qualified by the  
"urn:ietf:params:xml:ns:xmpp-objenc:0" namespace as follows:

- The attribute 'mac-algo' set to the cryptographic hashing algorithm used to generate M in step 8;
- The attribute 'mac-hash' set to the base64-encoded result of the MAC, M;
- The attribute 'cipher-algo' set to the block encryption scheme used to generate the encrypted data T in step 7;
- The XML character data as the base64-encoded form of T.

11. Sends the `<e2e/>` element as the payload of a stanza that SHOULD match the stanza from step 1 in kind (e.g., `<message/>`), type (e.g., "chat"), and addressing (e.g. to="romeo@montague.net" from="juliet@capulet.net/balcony"). If the original stanza (S) has a value for the "id" attribute, this stanza MUST NOT use the same value for its "id" attribute.

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### 3.1. Example of Securing Messages

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The sender begins with the cleartext version of the `<message/>` stanza "S":

```
<message      xmlns='jabber:client'
              from='juliet@capulet.net/balcony'
              id='183ef129'
              to='romeo@montague.net'
              type='chat'>
  <thread>8996aef0-061d-012d-347a-549a200771aa</thread>
  <body>Wherefore art thou, Romeo?</body>
</message>
```

The sender then performs the steps 1 through 5 from above to generate:

```
<plain xmlns="urn:ietf:params:xml:ns:xmpp-objenc:0"
  timestamp="2010-06-29T02:15:21.012Z">
  PG1lc3NhZ2UgeG1sbnM9ImphYmJlcjpbGllbnQiIGZyb209Imp1bGll
  dEBjYXB1bGV0Lm5ldC9iYWxjb255IiB0bz0icm9tZW9AbW9udGVndWUubmV0IiB0eXB1PS
  JjaGF0Ij48dGhyZWFKPmM2MzczODI0LWEzMDctNDBkZC04ZmUwLWJhZDZlNzI5O
  WFKMDwvdGhyZWFKPjxib2R5PlldoZXJlZm9yZSBhcnQgdGhvdSwgUm9tZW8/PC9i
  b2R5PjwvbWVzc2FnZT4=
</plain>
```

Then performs steps 6 through 10, and sends the following:

```
<message xmlns='jabber:client'
  from='juliet@capulet.net/balcony'
  id='6410ed123'
  to='romeo@montague.net'
  type='chat'>
  <e2e xmlns="urn:ietf:params:xml:ns:xmpp-objenc:0">
    <key cipher-algo="RSAES-OAEP-SHA-256-MFG1">
      OPfr4zudqiEeLc0QazZJIB6B9gx3zrVbyHKTU8a/aDb0wiZevztxxCi8hto0+Qw
      Foyhcupj547WbFZJN1B2dsAPh1JzeH9SuGLJShjhbk0yKjmqZLLCZr30QtJjcTU
      sAVj7IZZs00PDmwsb4Dxv5sz+icsDpi5l+5APfthDaoHbcrvz2pA1CJ5IFQoob4
      a0i0WevcAFyB+vWXsRqQCxjn5sHdb6G4vjQ/m1lzTWahzKvi56pNUM71l18oI8L
      mPi1VWUEqH3aayGLVlJ9fhBDSSpw4jTQ/ts1nzPJwVlKdTqdgNBusFEhrRMhJD5
      1JdL0hxx+0v2Xbs22++XQ1tS8/A==
    </key>
    <data cipher-algo="AES-256-CBC-PKCS5-WITH-IV"
      mac-algo="HMACSHA256"
      mac-data="HSGmwUFd4sESB+012S32xsXVvMn04gjRPaQITirjWbs=">
      a8zpjjgRc01VHZ9CoqU19/jB7nn58Gzu5/sQm8YQe4F9zz+YKUfqTS9LaHcqdaAwa
      z8BG1a24Z72VYb5Ptjh7nQ19f5QQdA/P4lZ3oqeTJTsA4DkhvJaSUhrjYib/N0k
      3lkMoatR/OSbfvhpDqXQ/dutLuRFjkilXGVvNwkgLm3iSnKUiYSdUzWvj88RgR3
      ldVHFeyrdgufU9qu/Fy06MZxfEtD800+3ZBbESqllzmYFXnfkzBrhfi14iCba6
      /b5Io5zhFUyWaq5e6qq2z72a+1bjeWkG8F9XBimkyaxkB64wAS0o6aDpWdir50i
      +Rnms4LV/wxL4Is/oe8Fo9xR3UmrdlAiaehdGBh+EnJGqprKa9ecc0KqSu7/lJQ
      0bAdJGE0eAVs8JEkQkxw+qR8edkEDuv6ZXN7JCWQx9LNaiiwsfAzApJJbqfrtDx
      koQ3JaBbxQ+8FE3TM0E4Tbr9V8NDZC8abgBramlpUBfgknJvLYMTzx1lnsiCUxo
      6ezC0xqV
    </data>
  </e2e>
</message>
```

---

### 3.2. Example of Securing IQs

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The sender begins with the cleartext version of the <iq/> stanza "S":

```

<iq      xmlns="jabber:client"
      from="juliet@capulet.net/crypt"
      id="a543bc3ee"
      to="romeo@montague.net/crypt"
      type="result">
  <mood xmlns="http://jabber.org/protocol/mood">
    <dejected />
    <text>
      Romeo, what's here? Poison? Drunk all, and
      left no friendly drop to help me after?
    </text>
  </mood>
</iq>

```

The sender then performs the steps 1 through 5 from above to generate:

```

<plain  xmlns="urn:ietf:params:xml:ns:xmpp-objenc:0"
      timestamp="2010-06-29T02:15:21.012Z">
  PGlxIHhtbG5zPSJqYWJiZXI6aXElIGZyb209Imp1bGlldEBjYXB1bGV0Lm5ldC9
  jcnlwdCIgaWQ9ImE1NDNiYzNlZSIgdG89InJvbWVvQG1vbnRlZ3VlLm5ldC9jcn
  lwdCIgdHlwZT0icmVzdWx0Ij48bW9vZCB4bWxucz0iaHR0cDovL2phYmJlci5vc
  mcvCHJvdG9jb2wvbW9vZCI+PGRlamVjdGVkIC8+PHRleHQ+Um9tZW8sIHdoYXQn
  cyBoZXJlPyBQb2lzb24/IERYdW5rIGFsbCwgYW5kIGx1ZnQgbm8gZnJlbnRseSB
  kcm9wIHRvIGhlbHAgbWUgYWZ0ZXI/PC90ZXh0PjwvbW9vZD48L2lxPg==
</plain>

```

Then performs steps 6 through 10, and sends the following:

```

<iq      xmlns="jabber:client"
      type="result"
      to="romeo@montegue.net/crypt"
      id="42ca3de0345"
      from="juliet@capulet.net/crypt">
<e2e xmlns="urn:ietf:params:xml:ns:xmpp-objenc:0">
  <key      cipher-algo="RSAES-OAEP-SHA-256-MFG1">
    hOU+BRkEcCY0+eKTX9hzCbP30Ij0q5zZ9buFgk0Wu4LsVki920iH65SvYL/XCB6
    12sb9fhjkiAIeR0AySGiid+AeS7KZDzpcZ+ORg8j9CkEX/LeTYszBfZFiHzDFkh
    qtwu3s7QMAR0Bzxj9NVE7W8fSdleusvy00P5c0scrpRkXDMV02Z3/rTjC0xInx3
    XQUP+RlqFE7g1HCr01BjoPjI4p3N+fONVv0U9mwt1I5tJ4EXgTofUM0GMNGX1i
    NoNNjPDb9XsihpLvDIjMblXVHvYAIyPwCs2ZdDv7L5kmZ6U+35b7Qx8TdWUN2I4
    5fBbxczvKFN6+cx2h5uap0TxBkw==
  </key>
  <data      cipher-algo="AES-256-CBC-PKCS5-WITH-IV"
    mac-algo="HMACSHA256"
    mac-data="iKuTGRZNHe3PbZNdfxkFzwC1XLMQlhx0Y8BuYawbaho=">
    ksCAkoJeoymtf3ygzBJkrJYQV+g04CkAs5oSmej60GU89mRN3rKSX5FVfWo558W
    Bcn8mVUxFxWhSdNBrsW5GQS1EyygDT+yfJe60qzLTCqZn4iqaCyIPWM7XB/PolA
    fvELw7y3hf8JrEAM4JXIfxrc0YDqewr7zmamwuuos4B6qzgiNN9ZW2AfTyKL3+l
    twcmFvF/nWF1YN8CquGmBm83WFn7Ik9R+Nqq54+QNCABjSFPT25ZYquEhXk/RIS
    CDAIXF0aFBozGjC20M3UDqnuwLsUF+P4ucjybysxhHQlqL0ffX0Vhb1YesWaZac
    pvsj80vfpv+ESrWGptXr+8GMK1og69GHRrd2k2TonPFp1KwS5MkbEpP2tS7R+nT
    b9oGFojr6waNKhhhVmP/9FWRMi7C2KfLCHggAatLWDjBG8k7yd5DwdSqY7LwkWB
    hT6+iErRfhdvk1EVxn2TVqj fhsFh33XDqkRT4BhPJUjJPkwLZkQ03PVgHKluMsE
    JoUBS00xD7gE5q808hy3qA+r5PDowy6nQ9zbUaCu4JbvKV2moql7fgHUy8MZLIe
    DFVJ5A5z8Te6K4pFaQGAzxEOouS2A+BmvPAFcZFeL+QGy58RSNMiXJ9ZMpb+N2C
    1iDzPD80L
  </data>
</e2e>
</iq>

```

---

#### 4. Interaction with Stanza Semantics

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The following limitations and caveats apply:

- \*Undirected <presence/> stanzas MUST NOT be encrypted. Such stanzas are delivered to anyone the sender has authorized, and therefore it is highly unlikely that the sender can find an appropriate certificate.

- \*Stanzas directed to multiplexing services (e.g. multi-user chat) SHOULD NOT be encrypted, unless the sender has established an acceptable trust relationship with the multiplexing service.



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## 5. Handling of Inbound Stanzas

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Several scenarios are possible when an entity receives an encrypted stanza:

- \*The receiving application does not understand the protocol.

- \*The receiving application understands the protocol and is able to decrypt the payload.

- \*The receiving application understands the protocol and is able to decrypt the payload, but the timestamps fail the checks specified under [Checking of Timestamps \(Inclusion and Checking of Timestamps\)](#).

- \*The receiving application understands the protocol but is unable to decrypt the payload.

In Case #1, the receiving application MUST do one and only one of the following: (1) ignore the <e2e/> extension, (2) ignore the entire stanza, or (3) return a <service-unavailable/> error to the sender, as described in [\[XMPP-CORE\] \(Saint-Andre, P., "Extensible Messaging and Presence Protocol \(XMPP\): Core," May 2010.\)](#).

In Case #2, the receiving application MUST NOT return a stanza error to the sender, since this is the success case.

In Case #3, the receiving application MAY return a <not-acceptable/> error to the sender (as described in [\[XMPP-CORE\] \(Saint-Andre, P., "Extensible Messaging and Presence Protocol \(XMPP\): Core," May 2010.\)](#)), optionally supplemented by an application-specific error condition element of <bad-timestamp/> (previously defined in [\[RFC3923\] \(Saint-Andre, P., "End-to-End Signing and Object Encryption for the Extensible Messaging and Presence Protocol \(XMPP\)," October 2004.\)](#)):

```

<message from='romeo@example.net/orchard'
  id='6410ed123'
  to='juliet@capulet.net/balcony'
  type='error'>
  <e2e xmlns='urn:ietf:params:xml:ns:xmpp-objenc:0'>
    XML-character-data-here
  </e2e>
  <error type='modify'>
    <not-acceptable xmlns='urn:ietf:params:xml:ns:xmpp-stanzas' />
    <bad-timestamp xmlns='urn:ietf:params:xml:xmpp-e2e' />
  </error>
</message>

```

In Case #4, the receiving application SHOULD return a <bad-request/> error to the sender (as described in [\[XMPP-CORE\] \(Saint-Andre, P., "Extensible Messaging and Presence Protocol \(XMPP\): Core," May 2010.\)](#)), optionally supplemented by an application-specific error condition element of <decryption-failed/> (previously defined in [\[RFC3923\] \(Saint-Andre, P., "End-to-End Signing and Object Encryption for the Extensible Messaging and Presence Protocol \(XMPP\)," October 2004.\)](#)):

```

<message from='romeo@example.net/orchard'
  id='6410ed123'
  to='juliet@capulet.net/balcony'
  type='error'>
  <e2e xmlns='urn:ietf:params:xml:ns:xmpp-objenc:0'>
    XML-character-data-here
  </e2e>
  <error type='modify'>
    <bad-request xmlns='urn:ietf:params:xml:ns:xmpp-stanzas' />
    <decryption-failed xmlns='urn:ietf:params:xml:xmpp-e2e' />
  </error>
</message>

```

In addition to returning an error in Case #4, the receiving application SHOULD NOT present the stanza to the intended recipient (human or application) and SHOULD provide some explicit alternate processing of the stanza (which MAY be to display a message informing the recipient that it has received a stanza that cannot be decrypted).

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## 6. Inclusion and Checking of Timestamps

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Timestamps are included to help prevent replay attacks. All timestamps MUST conform to [\[DATETIME\] \(Klyne, G. and C. Newman, "Date and Time on the Internet: Timestamps," July 2002.\)](#) and be presented as UTC with no offset, and SHOULD include the seconds and fractions of a second to

three digits. Absent a local adjustment to the sending agent's perceived time or the underlying clock time, the sending agent MUST ensure that the timestamps it sends to the receiver increase monotonically (if necessary by incrementing the seconds fraction in the timestamp if the clock returns the same time for multiple requests). The following rules apply to the receiving application:

- \*It MUST verify that the timestamp received is within five minutes of the current time, except as described below for offline messages.

- \*It SHOULD verify that the timestamp received is greater than any timestamp received in the last 10 minutes which passed the previous check.

- \*If any of the foregoing checks fails, the timestamp SHOULD be presented to the receiving entity (human or application) marked as "old timestamp", "future timestamp", or "decreasing timestamp", and the receiving entity MAY return a stanza error to the sender.

The foregoing timestamp checks assume that the recipient is online when the message is received. However, if the recipient is offline then the server will probably store the message for delivery when the recipient is next online (offline storage does not apply to <iq/> or <presence/> stanzas, only <message/> stanzas). As described in [\[OFFLINE\] \(Saint-Andre, P., "Best Practices for Handling Offline Messages," January 2006.\)](#), when sending an offline message to the recipient, the server SHOULD include delayed delivery data as specified in [\[DELAY\] \(Saint-Andre, P., "Delayed Delivery," September 2009.\)](#) so that the recipient knows that this is an offline message and also knows the original time of receipt at the server. In this case, the recipient SHOULD verify that the timestamp received in the encrypted message is within five minutes of the time stamped by the recipient's server in the <delay/> element.

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## 7. Mandatory-to-Implement Cryptographic Algorithms

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All implementations MUST support the following algorithms. Implementations MAY support other algorithms as well.

- \*The RSA (PKCS #1 v2.1) key transport, as specified in [\[X509-ALGO\] \(Jonsson, J. and B. Kaliski, "Public-Key Cryptography Standards \(PKCS\) #1: RSA Cryptography Specifications Version 2.1," February 2003.\)](#).

\*The AES-128 encryption algorithm in CBC mode, as specified in [\[CMS-AES\] \(Schaad, J., "Use of the Advanced Encryption Standard \(AES\) Encryption Algorithm in Cryptographic Message Syntax \(CMS\)," July 2003.\)](#).

\*The HMACSHA256 hashing algorithm, as specified in [\[HMAC\] \(Eastlake, D. and T. Hansen, "US Secure Hash Algorithms \(SHA and HMAC-SHA\)," July 2006.\)](#).

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## 8. Certificates

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To participate in end-to-end encryption using the methods defined in this document, a client needs to possess an X.509 certificate [\[PKIX\] \(Cooper, D., Santesson, S., Farrell, S., Boeyen, S., Housley, R., and W. Polk, "Internet X.509 Public Key Infrastructure Certificate and Certificate Revocation List \(CRL\) Profile," May 2008.\)](#). It is expected that many clients will generate their own (self-signed) certificates rather than obtain a certificate issued by a certification authority (CA). In any case the certificate MUST include an XMPP address that is represented using the ASN.1 Object Identifier "id-on-xmppAddr" as specified in Section 5.1.1 of [\[XMPP-CORE\] \(Saint-Andre, P., "Extensible Messaging and Presence Protocol \(XMPP\): Core," May 2010.\)](#).

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## 9. Security Considerations

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The recipient's server might store any <message/> stanzas received until the recipient is next available; this duration could be anywhere from a few minutes to several months.

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## 10. IANA Considerations

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## 10.1. XML Namespace Name for e2e Data in XMPP

A URN sub-namespace of encrypted content for the Extensible Messaging and Presence Protocol (XMPP) is defined as follows.

**URI:** urn:ietf:params:xml:ns:xmpp-objenc:0

**Specification:** RFC XXXX

**Description:** This is an XML namespace name of signed and encrypted content for the Extensible Messaging and Presence Protocol as defined by RFC XXXX.

**Registrant Contact:** IESG, <iesg@ietf.org>

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## 11. References

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### 11.1. Normative References

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[BASE64]	Josefsson, S., " <a href="#">The Base16, Base32, and Base64 Data Encodings</a> ," RFC 4648, October 2006 ( <a href="#">TXT</a> ).
[CMS-AES]	Schaad, J., " <a href="#">Use of the Advanced Encryption Standard (AES) Encryption Algorithm in Cryptographic Message Syntax (CMS)</a> ," RFC 3565, July 2003 ( <a href="#">TXT</a> ).
[DATETIME]	Klyne, G. and C. Newman, " <a href="#">Date and Time on the Internet: Timestamps</a> ," RFC 3339, July 2002 ( <a href="#">TXT</a> ).
[DELAY]	<a href="#">Saint-Andre, P.</a> , " <a href="#">Delayed Delivery</a> ," XSF XEP 0203, September 2009.
[E2E-REQ]	Saint-Andre, P., " <a href="#">Requirements for End-to-End Encryption in the Extensible Messaging and Presence Protocol (XMPP)</a> ," draft-saintandre-xmpp-e2e-requirements-01 (work in progress), March 2010 ( <a href="#">TXT</a> ).
[KEYWORDS]	<a href="#">Bradner, S.</a> , " <a href="#">Key words for use in RFCs to Indicate Requirement Levels</a> ," BCP 14, RFC 2119, March 1997.
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[PKIX]	Cooper, D., Santesson, S., Farrell, S., Boeyen, S., Housley, R., and W. Polk, " <a href="#">Internet X.509 Public Key Infrastructure Certificate and Certificate Revocation List (CRL) Profile</a> ," RFC 5280, May 2008 ( <a href="#">TXT</a> ).
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[XMPP-CORE]	Saint-Andre, P., " <a href="#">Extensible Messaging and Presence Protocol (XMPP): Core</a> ," draft-ietf-xmpp-3920bis-08 (work in progress), May 2010 ( <a href="#">TXT</a> ).

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## 11.2. Informative References

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[OFFLINE]	<a href="#">Saint-Andre, P.</a> , " <a href="#">Best Practices for Handling Offline Messages</a> ," XSF XEP 0160, January 2006.
[RFC3923]	Saint-Andre, P., " <a href="#">End-to-End Signing and Object Encryption for the Extensible Messaging and Presence Protocol (XMPP)</a> ," RFC 3923, October 2004 ( <a href="#">TXT</a> ).

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## Appendix A. Schema for urn:ietf:params:xml:ns:xmpp-objenc:0

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The following XML schema is descriptive, not normative.

```

<?xml version='1.0' encoding='UTF-8'?>

<xs:schema
  xmlns:xs='http://www.w3.org/2001/XMLSchema'
  targetNamespace='urn:ietf:params:xml:ns:xmpp-objenc:0'
  xmlns='urn:ietf:params:xml:ns:xmpp-objenc:0'
  elementFormDefault='qualified'>

  <xs:element name='e2e'>
    <xs:complexType>
      <xs:sequence>
        <xs:element ref='key' minOccurs='1' maxOccurs='1' />
        <xs:element ref='data' minOccurs='1' maxOccurs='1' />
      </xs:sequence>
    </xs:complexType>
  </xs:element>

  <xs:element name='key'>
    <xs:complexType>
      <xs:simpleType>
        <xs:extension base='xs:string'>
          <xs:attribute name='cipher-algo'
            type='xs:string' />
        </xs:extension>
      </xs:simpleType>
    </xs:complexType>
  </xs:element>

  <xs:element name='data'>
    <xs:complexType>
      <xs:simpleType>
        <xs:extension base='xs:string'>
          <xs:attribute name='cipher-algo'
            type='xs:string' />
          <xs:attribute name='mac-algo'
            type='xs:string' />
          <xs:attribute name='mac-hash'
            type='xs:string' />
        </xs:extension>
      </xs:simpleType>
    </xs:complexType>
  </xs:element>

  <xs:element name='plain'>
    <xs:complexType>
      <xs:simpleType>
        <xs:extension base='xs:string'>
          <xs:attribute name='timestamp'

```

```

                                type='xs:string' />
        </xs:extension>
    </xs:simpleType>
</xs:complexType>
</xs:element>

<xs:element name='decryption-failed' type='empty' />
<xs:element name='bad-timestamp' type='empty' />

<xs:simpleType name='empty'>
    <xs:restriction base='xs:string'>
        <xs:enumeration value='' />
    </xs:restriction>
</xs:simpleType>

</xs:schema>

```

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## Authors' Addresses

[TOC](#)

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