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Xiaoyi Liu  
Cheryl Madson  
David McGrew  
Andrew Nourse  
Cisco Systems

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Cisco Systems' Simple Certificate Enrollment Protocol(SCEP):

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## Abstract

This document specifies the Cisco Simple Certificate Enrollment Protocol, a PKI communication protocol which leverages existing technology by using PKCS#7 and PKCS#10. SCEP is the evolution of the enrollment protocol developed by Verisign, Inc. for Cisco Systems, Inc. It now enjoys wide support in both client and CA implementations.

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## [Section 1.](#) Introduction

Public key technology is becoming more widely deployed and is becoming the basis for standards based security, such as the Internet Engineering Task Force's IPSEC and IKE protocols. With the use of public key certificates in network security protocols comes the need for a certificate management protocol that Public Key Infrastructure (PKI) clients and Certificate Authority servers can use to support certificate life cycle operations such as certificate enrollment and revocation, and certificate and CRL access.

In the following, [Section 2](#) gives an overview of the PKI operations, and [Section 2.4](#) describes the security goals of the protocol and the mechanisms used to achieve them. The transport protocol and the security protocol PKCS#7 are described at [Section 3](#) and [Section 4](#), respectively. The last section, [Section 5](#), specifies each PKI operation in terms of the message formats and the data structures of each operation.

The appendices provide detailed specifications and examples. Requestor subject names are specified in [Appendix A](#), attribute OIDs are specified in [Appendix C](#), and the SCEP state transitions are described in [Appendix E](#). **An example of a certificate enrollment request is provided in Appendix B**, and an example LDAP query URL encoding is provided in [Appendix D](#).

The authors would like to thank Peter William of ValiCert, Inc. (formerly of Verisign, Inc) and Alex Deacon of Verisign, Inc. and Christopher Welles of IRE, Inc. for their contributions to this protocol and to this document.



## **2.0 The Goal of SCEP**

The goal of SCEP is to support the secure issuance of certificates to network devices in a scalable manner, using existing technology whenever possible. The protocol supports the following operations:

- CA and RA public key distribution
- Certificate enrollment
- Certificate revocation
- Certificate query
- CRL query

Certificate and CRL access can be achieved by using the LDAP protocol (as specified in [Appendix D](#)), or by using the query messages defined in SCEP. The use of HTTP certificate and CRL access, and the support of CDP as specified in [RFC2459](#), will be specified in a future version of this document. In [Section 2.1](#), we first define PKI entity types as well as the properties of each entity type. In [Section 2.2](#), the PKI operations are described at functional level. [Section 2.3](#) describes the transaction behavior of each PKI operations. The complete PKI messages are covered in [Section 5](#).

## **2.1 SCEP Entity types**

The entity types defined in SCEP are the "requestor" type (i.e., IPSEC clients), the Certificate Authority (CA) entity type, and the Registration Authority entity type (RA). A requestor is sometimes called a "SCEP client" in the following.

### **2.1.1 Requestors**

A requestor is an entity whose name is defined in a certificate subject name field and optionally, in SubjectAltName, a X.509 certificate V3 extension. As a requestor, a SCEP client is identified by a subject name consisting of the following naming attributes:

- Fully qualified domain name, for example, router.cisco.com
- IP address, Serial number, and/or x.500 distinguished name

The fully qualified domain name is required for a requestor that intends to use the certificate for ISAKMP. The IP address, serial number, and x.500 distinguished name are optional name attributes. In the certificate enrollment request, the PKCS#10 subject field contains the required and optional name attributes. Based on the PKCS#10 subject name information, the certificate issued to the requestor must have the same name attributes set both in the subjectName field and in the SubjectAltName extension.

It is important to note that a client named as Alice.cisco.com is different than a client named as Alice.cisco.com plus the IP address name attribute 117.96.1.219. From CA point of view, the Distinguished

names assigned in these two cases are distinct names.

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Entity names which are specified as in the IPSEC profile (i.e., FQDN, IP address and User FQDN) must be presented in certificate's SubjectAltName extension. Multiple IPSEC entity names, (if any) are encoded as multiple values of a single SubjectAltName extension. The CA has the authority to assign a distinguished name to a requestor, whether or not one was included in the request. The assigned DN should contain the SCEP client names as the relative DN.

The attribute identifiers and an example of SCEP client subject name are specified in [Appendix A](#). [Appendix B](#) has an example from Cisco VPN Client enrollment request.

#### **2.1.1.1 Local Key/Certificate/CRL Storage and Certificate-name uniqueness**

A requestor is required to generate asymmetric key pairs and to provide storage to store its private keys. If the requestor does not have enough permanent memory to save its certificate, the it should be able to query its own certificate from the CA or an LDAP server, once the certificate has been issued. The public key pairs can be generated with a specific key usage. The key usage is conveyed to the CA through the certificate enrollment request. All current SCEP client implementations expect that there will be only one pair of keys for a given subject name and key usage combination and CA, at any time. This property is called the certificate-name uniqueness property, and it implies that a CA that implements SCEP will enforce the unique mapping between a SCEP client subject name and its key pairs with a given key usage. At any time, if the subject name is changed, or if the key is updated, the existing certificate would have to be revoked before a new one could be issued.

It is desirable that the CA enforce certificate-name uniqueness, but it is not mandatory. However a CA that does not enforce uniqueness must provide some other mechanism to prevent the re-transmission of an enrollment request by a SCEP client from creating a second certificate or certificate request, nor can the second request merely be rejected. If a client times out from polling for a pending request it can resynchronize by reissuing the original request with the original subject name and transaction ID. This must return the status of the original transaction, including the certificate if it was granted. It must not create a new transaction unless the original cert has been revoked, or the transaction arrives more than halfway through the validity time of the original certificate.

An enrollment request that occurs more than halfway through the validity time of an existing certificate for the same subject name and key usage MAY be interpreted as a renewal request and accepted regardless of the duplication of subject name. Certificate renewal can be done this way.

#### **2.1.1.2 Requestor authentication**

As with every protocol that uses public-key cryptography, the association between the public keys used in the protocol and the identities with which they are associated must be authenticated in a



cryptographically secure manner. This requirement is needed to prevent a "man in the middle" attack, in which an adversary that can manipulate the data as it travels between the protocol participants can subvert the security of the protocol. To satisfy this requirement, SCEP provides two authentication methods: manual authentication, and authentication based on pre-shared secret. In the manual mode, the requestor is required to wait until its identity can be verified by the CA operator using any reliable out-of-band method. To prevent a "man-in-the-middle" attack, an MD5 'fingerprint' generated on the PKCS#10 (before PKCS #7 enveloping and signing) must be compared out-of-band between the server and the requestor. SCEP clients and CAs (or RAs, if appropriate) must display this fingerprint to the operator to enable this verification if manual mode is used. Failing to provide this information leaves the protocol vulnerable to attack by sophisticated adversaries. When utilizing a pre-shared secret scheme, the server should distribute a shared secret to the requestor which can uniquely associate the enrollment request with the given end entity. The distribution of the secret must be private: only the end entity should know this secret. The actual binding mechanism between the requestor and the secret is subject to the server policy and implementation. When creating the enrollment request, the requestor is asked to provide a challenge password. When using the pre-shared secret scheme, the requestor must enter the re-distributed secret as the password. In the manual authentication case, the challenge password only used to authenticate a request for the certificate's revocation. This challenge password is included as a PKCS#10 attribute, and is sent to the server as encrypted data. The PKCS#7 envelope protects the privacy of the challenge password with DES encryption.

#### **2.1.1.3 Requestor Uses Existing CA-Issued or Self-Signed Certificates**

In this protocol, the communication between the requestor and the certificate authority is secured by using PKCS#7 as the messaging protocol. PKCS#7, however, is a protocol which assumes the communicating entities already possess the peer's certificates and requires both parties use the issuer names and issuer assigned certificate serial numbers to identify the certificate in order to verify the signature and decrypt the message. If the requesting system already has a certificate issued by the CA, that certificate may be presented as credentials for the renewal of that certificate if the policy of the CA permits this. If the requestor has no certificate issued by the CA, it must generate a self-signed certificate for its public key. In this self-signed certificate, the issuer name is the requestor subject name (the same name later used in the PKCS#10). During the certificate enrollment, the requestor will first post itself as the signing authority by attaching the self-signed certificate to the signed certificate request. When the Certificate Authority makes the envelope on the issued certificate using the public key included in the

self-signed certificate, it should use the same issuer name and serial number as conveyed in the self-signed certificate to inform the end entity on which private key should be used to open the envelope.

Note that when a client enrolls for separate encryption and signature certificates, it may use the signature certificate to sign both requests, and then expect its signature key to be used to encrypt both responses. In any case, the recipientinfo on the envelope should reflect the key used to encrypt the request.

### **[2.1.2](#) Certificate Authority**

A Certificate Authority(CA) is an entity whose name is defined in the certificate issuer name field. Before any PKI operations can begin, the CA generates its own public key pair and creates a self-signed CA certificate, or causes another CA to issue a certificate to it. Associated with the CA certificate is a fingerprint which will be used by the requestor to authenticate the received CA certificate if it is self-signed. The fingerprint is created by calculating a MD5 hash on the whole CA certificate. Before any requestor can start its enrollment, this root certificate has to be configured at the entity side securely. For IPSEC clients, the client certificates must have SubjectAltName extension. To utilize LDAP as a CRL query protocol, the certificates must have CRL Distribution Point. Key usage is optional. Without key usage, the public key is assumed as a general purpose public key and it can be used for all the purposes.

A Certificate Authority may enforce certain name policy. When using [X.500](#) directory name as the subject name, all the name attributes specified in the PKCS#10 request should be included as Relative DN. All the name attributes as defined in [RFC2459](#) should be specified in the SubjectAltName. An example is provided in [Appendix A](#).

If there is no LDAP query protocol support, the Certificate Authority should answer certificate and CRL queries, and to this end it should be online all the time.

The updating of the CA's public key is not addressed within the SCEP protocol. An SCEP client can remove its copy of a CA's public key and re-enroll under the CA's new public key.

### **[2.1.3](#) Registration Authorities**

In the environment where a RA is present, an requestor performs enrollment through the RA. In order to setup a secure channel with RA using PKCS#7, the RA certificate(s) have to be obtained by the client in addition to the CA certificate(s).

In the following, the CA and RA are specified as one entity in the context of PKI operation definitions.

### **[2.1.4](#) Trusted Root Store**

To support interoperability between IPSEC peers whose certificates are

issued by different CA, SCEP allows the users to configure multiple trusted roots. A root is a trusted root when its certificate has been

configured as such in the client. An SCEP client that supports multiple roots must associate with each root the information needed to query a CRL from each root.

Once a trusted root is configured in the client, the client can verify the signatures of the certificates issued by the given root.

## **2.2 SCEP Operations Overview**

In this section, we give a high level overview of the PKI operations as defined in SCEP.

### **2.2.1 Requestor Initialization**

The requestor initialization includes the key pair generation and the configuring of the required information to communicate with the certificate authority.

#### **2.2.1.1 Key Pair Generation**

Before an requestor can start PKI transaction, it first generates asymmetric key pairs, using the selected algorithm (the RSA algorithm is required in SCEP, and is the only algorithm in current implementations).

A requestor can create one or more asymmetric key pairs. The key pairs can be created for encryption only, signing only, or for all purposes.

The key pairs are saved by the client in NVRAM or other non-volatile media. The identification of a key pair is based on the FQDN assigned to the client and the selected key usage. Every time a new key pair is generated to replace the old key pair, the existing certificates have to be revoked from the CA and a new enrollment has to be completed.

#### **2.2.1.2 Required Information**

A requestor is required to have the following information configured before starting any PKI operations:

- 1. the certificate authority IP address or fully-qualified domain name,**
- 2. the certificate authority HTTP CGI script path, and  
the HTTP proxy information in case there is no direct Internet  
connection to the server,**
- 3. If the CRL is to be obtained by from a directory server by means  
of LDAP, and the CDP contains only the X.500 directory name, then  
the client will need to know the LDAP server fully-qualified domain  
name or IP address.**

### **2.2.2 CA/RA Certificate Distribution**

Before any PKI operation can be started, the requestor needs to get the CA/RA certificates. At this time, since no public key has been

exchanged between the requestor and the CA/RA, the message to get the CA/RA certificate can not be secured using PKCS#7 protocol. Instead, the CA/RA certificate distribution is implemented as a clear HTTP Get operation. After the requestor gets the CA certificate, it has to authenticate the CA certificate by comparing the finger print with the CA/RA operator. Since the RA certificates are signed by the CA, there is no need to authenticate the RA certificates.

This operation is defined as a transaction consisting of one HTTP Get message and one HTTP Response message:

REQUESTOR	CA SERVER
Get CA/RA Cert: HTTP Get message	
----->	
	CA/RA Cert download: HTTP Response message
	<-----
Compute finger print and call CA operator.	
	Receive call and check finger print

If an RA is in use, a degenerated PKCS#7 with a certificate chain consisting of both RA and CA certificates is sent back to the end entity. Otherwise the CA certificate is directly sent back as the HTTP response payload.

### **2.2.3 Certificate Enrollment**

A requestor starts an enrollment transaction by creating a certificate request using PKCS#10 and send it to the CA/RA enveloped using the PKCS#7. After the CA/RA receives the request, it will either automatically approve the request and send the certificate back, or it will require the requestor to wait until the operator can manually authenticate the identity of the requesting requestor. Two attributes are included in the PKCS#10 certificate request - a Challenge Password attribute and an optional ExtensionReq attribute which will be a sequence of extensions the requestor would like to be included in its V3 certificate extensions. The Challenge Password is used for revocation and may be used (at the option of the CA/RA) additionally as a one-time password for automatic enrollment.

In the automatic mode, the transaction consists of one PKCSReq PKI Message, and one CertRep PKI message. In the manual mode, the requestor enters into polling mode by periodically sending GetCertInitial PKI message to the server, until the server operator completes the manual authentication, after which the CA will respond to GetCertInitial by returning the issued certificate.

The transaction in automatic mode:

REQUESTOR

CA SERVER

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PKCSReq: PKI cert. enrollment msg

```
-----> CertRep: pkiStatus = SUCCESS
          certificate
```

attached

```
<-----
```

Receive issued certificate.

The transaction in manual mode:

REQUESTOR

CA SERVER

PKCSReq: PKI cert. enrollment msg

```
-----> CertRep: pkiStatus = PENDING
```

```
<-----
```

GetCertInitial: polling msg

```
-----> CertRep: pkiStatus = PENDING
```

```
<-----
```

```
..... <manual identity authentication.....
```

GetCertInitial: polling msg

```
-----> CertRep: pkiStatus = SUCCESS
```

```
certificate
```

attached

```
<-----
```

Receive issued certificate.

#### **2.2.4 Requestor Certificate Revocation**

A requestor should be able to revoke its own certificate. Currently the revocation is implemented as a manual process. In order to revoke a certificate, the requestor makes a phone call to the CA server operator. The operator will come back asking the ChallengePassword (which has been send to the server as an attribute of the PKCS#10 certificate request). If the ChallengePassword matches, the certificate is revoked. The reason of the revocation is documented by CA/RA.

#### **2.2.5 Certificate Access**

There are two methods to query certificates. The first method is to use LDAP as a query protocol. Using LDAP to query assumes the client understand the LDAP scheme supported by the CA. The SCEP client assumes that the subject DN name in the certificate is used as URL to query the certificate. The standard attributes (userCertificate and caCertificate) are used as filter.

For the environment where LDAP is not available, a certificate query message is defined to retrieve the certificates from CA.

To query a certificate from the certificate authority, a requestor sends a request consisting of the certificate's issuer name and the

serial number. This assumes that the requestor has saved the issuer

name and the serial number of the issued certificate from the previous enrollment transaction. The transaction to query a certificate consists of one GetCert PKI message and one CertRep PKI message:

```

      REQUESTOR                                CA SERVER
GetCert: PKI cert query msg
-----> CertRep:  pkiStatus = SUCCESS
              certificate
attached
                                <-----
      Receive the certificate.
```

### **2.2.6 CRL Distribution**

The CA/RA will not "push" the CRL to the end entities. The query of the CRL can only be initialized by the requestor.

There are three methods to query CRL.

The CRL may be retrieved by a simple HTTP GET. If the CA supports this method, it should encode the URL into a CRL Distribution Point extension in the certificates it issues. Support for this method should be incorporated in new and updated clients, but may not be in older versions.

The second method is to query CRL using LDAP. This assumes the CA server supports CRL LDAP publishing and issues the CRL Distribution Point in the certificate. The CRL Distribution Point is encoded as a DN. Please refer to [Appendix D](#) for the examples of CRL Distribution Point.

The third method is implemented for the CA which does not support LDAP CRL publishing or does not implement the CRL Distribution Point. In this case, a CRL query is composed by creating a message consists of the CA issuer name and the CA's certificate serial number. This method is deprecated because it does not scale well and requires the CA to be a high-availability service.

The message is send to the CA in the same way as the other SCEP requests: The transaction to query CRL consists of one GetCRL PKI message and one CertRep PKI message which have no certificates but CRL.

```

      REQUESTOR                                CA SERVER
GetCRL: PKI CRL query msg
-----> CertRep:  CRL attached
                                <-----
```

## **2.3 PKI Operation Transactional Behavior**

As described before, a PKI operation is a transaction consisting of the messages exchanged between a requestor and the CA/RA. This section

will specify the transaction behavior on both the requestor and the

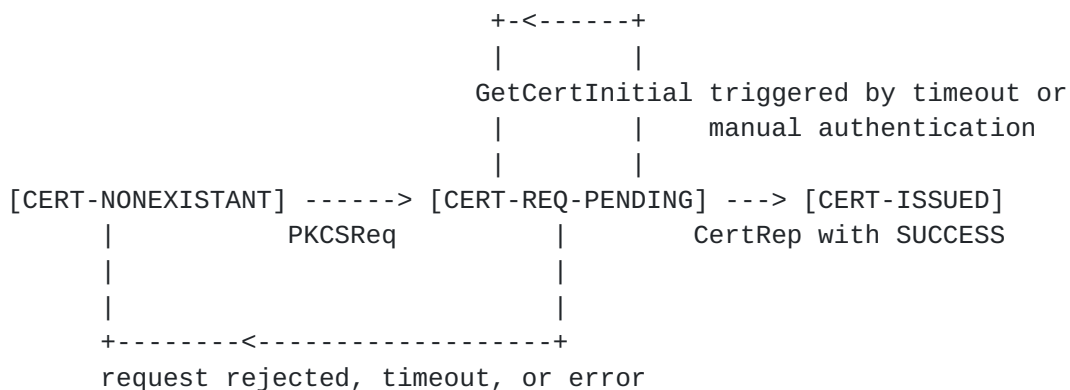
certificate authority server. Because the protocol is basically a two way communication protocol without a confirmation message from the initiating side, state and state resynchronization rules have to be defined, in case any error happens at either side. Before the state transition can be defined, the notion of transaction identifier has to be defined first.

### [2.3.1](#) Transaction Identifier

A transaction identifier is a string generated by the entity when starting a transaction. Since all the PKI operations defined in this protocol are initiated by the requestor, it is the responsibility of the requestor to generate a unique string as the transaction identifier. All the PKI messages exchanged for a given PKI operations must carry the same transaction identifier. The transaction identifier is generated as a MD5 hash on the public key value for which the enrollment request is made. This allows the SCEP client to reuse the same transaction identifier if it is reissuing a request for the same certificate (i.e. a certificate with the same subject, issuer, and key). The SCEP protocol requires that transaction identifiers be unique, so that queries can be matched up with transactions. For this reason, in those cases in which separate signing and encryption certificates are issued to the same requestor, the keys must be different.

### [2.3.2](#) State Transitions in Certificate Enrollment

The requestor state transitions during enrollment operation is indicated in the diagram below:



As described in the [section 2.2.3](#), certificate enrollment starts at the state CERT-NONEXISTANT. Sending PKCSReq changes the state to CERT-REQ-PENDING. Receiving CertRep with SUCCESS status changes the state to CERT-ISSUED. In the case the server sending back the response with pending status, the requestor will keep polling certificate response by sending GetCertInitial to the server, until either a CertRep with SUCCESS status is received, or the maximum polling number has been exceeded.

If an error or timeout occurs in the CERT-REQ-PENDING state, the end entity will transition to the CERT-NONEXISTANT state.

The client administrator will, eventually, start up another enrollment request. It is important to note that, as long as the requestor does not change its subject name or keys, the same transaction id will be used in the "new" transaction. This is important because based on this transaction id, the certificate authority server can recognize this as an existing transaction instead of a new one.

### **[2.3.3](#) Transaction Behavior of Certificate/CRL Access**

There is no state maintained during certificate access and CRL access transaction. When using the certificate query and CRL query messages defined in this protocol, the transaction identifier is still required so that the requestor can match the response message with the upstanding request message. When using LDAP to query the certificate and the CRL, the behavior is specified by the LDAP protocol.

## **[2.4](#) Security**

The security goals of SCEP are that no adversary can:

- o subvert the public key/identity binding from that intended,
- o discover the identity information in the enrollment requests and issued certificates,
- o cause the revocation of certificates with any non-negligible probability.

Here an adversary is any entity other than the requestor and the CA (and optionally the RA) participating in the protocol that is computationally limited, but that can manipulate data during transmission (that is, a man-in-the-middle). The precise meaning of 'computationally limited' depends on the implementer's choice of cryptographic hash functions and ciphers. The required algorithms are RSA, DES, and MD5.

The first and second goals are met through the use of PKCS#7 and PKCS#10 encryption and digital signatures using authenticated public keys. The CA's public key is authenticated via the checking of the CA fingerprint, as specified in [Section 2.1.2](#), and the SCEP client's public key is authenticated through the manual authentication or pre-shared secret authentication, as specified in [Section 2.1.1.2](#). The third goal is met through the use of a Challenge Password for revocation, that is chosen by the SCEP client and communicated to the CA protected by the PKCS#7 encryption, as specified in [Section 2.2.4](#).

The motivation of the first security goal is straightforward. The motivation for the second security goal is to protect the identity information in the enrollment requests and certificates. For example, two IPSEC hosts behind a firewall may need to exchange certificates, and

may need to enroll certificates with a CA that is outside of a firewall.  
Most networks with firewalls seek to prevent IP addresses and DNS



information from the trusted network leaving that network. The second goal enables the hosts in this example to enroll with a CA outside the firewall without revealing this information. The motivation for the third security goal is to protect the SCEP clients from denial of service attacks.

### [Section 3](#) Transport Protocol

In the SCEP protocol, HTTP is used as the transport protocol for the PKI messages.

#### [3.1](#) HTTP "GET" Message Format

In the PKI protocol, CA/RA certificates are sent to the requestor in clear, whereas the requestor certificates are sent out using the PKCS#7 secure protocol. This results in two types of GET operations. The type of GET operation is specified by augmenting the GET message with OPERATION and MESSAGE parameters in the Request-URL. OPERATION identifies the type of GET operation, and MESSAGE is actually the PKI message encoded as a text string.

The following is the syntax definition of a HTTP GET message sent from a requestor to a certificate authority server:

Request = "GET " CGI-PATH CGI-PROG "?operation=" OPERATION "&message=" MESSAGE  
where:

CGI-PATH defines the actual CGI path to invoke the CGI program which parses the request.

CGI-PROG is set to be the string "pkiclient.exe". This is intended to be the program that the CA will use to handle the SCEP transactions, though the CA may ignore CGI-PROG and use only the CGI-PATH.

OPERATION is set to be the string "PKIOperation" when the GET message carries a PKI message to request certificates or CRL; OPERATION is set to be the string "GetCACert" or "GetCACertChain" when the GET operation is used to get CA/RA certificate or the CA Cert chain (respectively).

When OPERATION is "PKIOperation", MESSAGE is a base64-encoded PKI message

when OPERATION is "GetCACert" or "GetCACertChain", MESSAGE is a string which represents the certificate authority issuer identifier.

For example, a requestor may submit a message via HTTP to the server as follows:

```
GET /cgi-bin/pkiclient.exe?operation=PKIOperation&message=MIAGCSqGSib3D
QEHA6CAMIACAQxgDCBzAIBADB2MGIXETAPBgNVBActCE .....AAAAAA==
```

#### [3.2](#) Response Message Format

For each GET operation, the CA/RA server will return a MIME object via

HTTP. For a GET operation with PKIOperation as its type, the response is tagged as having a Content Type of application/x-pki-message. The body of this message is a BER encoded binary PKI message. The following is an example of the response:

"Content-Type:application/x-pki-message\n\n"<BER-encoded PKI msg>

In the case of GET operation with a type of GetCACert, the MIME content type returned will depend on whether or not an RA is in use. If there is no RA, only the CA certificate is send back in the response, and the response has the content type tagged as application/x-x509-ca-cert. the body of the response is a DER encoded binary X.509 certificate. For example:

"Content-Type:application/x-x509-ca-cert\n\n"<BER-encoded X509>

If there is an RA, the RA certificates are send back together with the CA certificates, a certificate-only PKCS#7 SignedData is send back in the response where the SignerInfo is empty. [Section 5](#) has the detailed definition of the message format in this case. The content type is application/x-x509-ca-ra-cert.

#### [Section 4](#) Secure Transportation: PKCS#7

PKCS#7 is a general enveloping mechanism that enables both signed and encrypted transmission of arbitrary data. It is widely implemented and included in the RSA tool kit.

In this section, the general PKCS#7 enveloped PKI message format is specified. The complete PKCS#7 message format for each PKI transaction will be covered in [Section 5](#).

##### [4.1](#) SCEP Message Format

As a transaction message, a SCEP message has a set of transaction specific attributes and an information portion. Employing PKCS#7 protocol, the transaction specific attributes are encoded as a set of authenticated attributes of the SignedData. The information portion will first be encrypted to become Enveloped Data, and then the digest of the enveloped information portion is included as one of the message digest attributes and being signed together with the other transaction specific attributes.

By applying both enveloping and signing transformations, a SCEP message is protected both for the integrity of its end-end-transition information and the confidentiality of its information portion. The advantage of this technique over the conventional transaction message format is that, the signed transaction type information and the status of the transaction can be determined prior to invoke security handling procedures specific to the information portion being processed.

The following is an example of a SCEP message with its enveloped and signed data portion represented by pkcsPKISigned and pkcsPKIEnveloped. The out-most of any PKI message is a blob of ContentInfo, with its content type set to SignedData and the actual

signed data as the content.

```

pkiMessage ContentInfo ::= {
    contentType {pkcs-7 signedData(2)}
    content pkcsPKISigned
}
pkcsPKISigned SignedData ::= {
    version 1
    digestAlgorithm { iso(1) member-body(2) US(840) rsadsi(113549)
                    digestAlgorithm(2) 5}
    contentInfo {
        contentType {pkcs-7 1} -- data content identifier
        content pkcsPKIEnvelope -- enveloped information portion
    }
    certificates -- signer certificate chain
    signerInfo -- including signed transaction info and the digest
               -- of the enveloped information portion as the
               -- authenticated attributes
}
pkcsPKIEnvelope EnvelopedData ::= {
    version 0
    recipientInfos -- information required to open the envelop
    encryptedContentInfo {
        contentType {pkcs-7 1} -- data content identifier
        contentEncryptionAlgorithm
        encryptedContent -- encrypted information portion
    }
}

```

## 4.2 Signed Transaction Attributes

The following transaction attributes are encoded as authenticated attributes. Please refer to [Appendix B](#) for the OID definitions.

transactionID	PrintableString	-- Decimal value as a string
messageType	PrintableString	-- Decimal value as a string
pkiStatus	PrintableString	-- Decimal value as a string
failinfo	PrintableString	-- Decimal value as a string
senderNonce	Octet String	
recipientNonce	Octet String	

where:

The transactionID is an attribute which uniquely identify a transaction. This attribute is required in all PKI messages.

The messageType attribute specify the type of operation performed by the transaction. This attribute is required in all PKI messages. Currently, the following message types are defined:

PKCSReq (19) -- Permits use of PKCS#10 certificate request

CertRep (3) -- Response to certificate or CRL request  
GetCertInitial (20) -- Certificate polling in manual enrollment  
GetCert (21) -- Retrieve a certificate

GetCRL (22) -- Retrieve a CRL

All response message will include transaction status information which is defined as pkiStatus attribute:

SUCCESS (0) -- request granted  
FAILURE (2) -- request rejected  
PENDING (3) -- request pending for manual approval.

If the status in the response is FAILURE, the failinfo attribute will contain one of the following failure reasons:

badAlg (0) -- Unrecognized or unsupported algorithm ident  
badMessageCheck (1) -- integrity check failed  
badRequest (2) -- transaction not permitted or supported  
badTime (3) -- Message time field was not sufficiently close  
to the system time  
badCertId (4) -- No certificate could be identified matching  
the provided criteria

The attributes of senderNonce and recipientNonce are the 16 byte random numbers generated for each transaction to prevent the replay attack.

When a requestor sends a PKI message to the server, a senderNonce is included in the message. After the server processes the request, it will send back the requestor senderNonce as the recipientNonce and generates another nonce as the senderNonce in the response message. Because the proposed pki protocol is a two-way communication protocol, it is clear that the nonce can only be used by the requestor to prevent the replay. The server has to employ extra state related information to prevent a replay attack.

## [Section 5](#). SCEP Transaction Specification

In this section each SCEP transaction is specified in terms of the complete messages exchanged during the transaction.

### [5.1](#) Certificate Enrollment

The certificate enrollment transaction consists of one PKCSReq message send to the certificate authority from a requestor, and one CertRep message send back from the server. The pkiStatus returned in the response message is either SUCCESS, or FAILURE, or PENDING. The information portion of a PKCSReq message is a PKCS#10 certificate request, which contains the subject Distinguished Name, the subject public key, and two attributes, a ChallengePassword attribute to be used for revocation, and an optional ExtensionReq attribute which will be a sequence of extensions the requestor expects to be included in its V3 certificate extensions. One of the extension attribute specifies the key

usage. If the request is granted, the pkiStatus is set to SUCCESS, and the certificate is returned in CertRep; if the request is rejected, the



pkiStatus is set to FAILURE; if the server requires manual approval of the request, the pkiStatus is set to PENDING. The messages exchanged in the manual authentication mode is further specified in [Section 5.2](#).

**Precondition:**

Both the requestor and the certificate authority have completed their initialization process. The requestor has already been configured with the CA/RA certificate.

**Postcondition:**

Either the certificate is received by the requestor, or the end entity is notified to do the manual authentication, or the request is rejected.

### **[5.1.1](#) PKCSReq Message Format**

A PKCSReq message is created by following the steps defined below:

- [1.](#) Create a PKCS#10 certificate request which is signed by the end entity's private key, corresponding to the public key included in the PKCS#10 certificate request. This constitutes the information portion of PKCSReq.**
- [2.](#) Encrypt the PKCS#10 certificate request using a randomly generated content-encryption key. This content-encryption key is then encrypted by the CA's\* public key and included in the recipientInfo. This step completes the "envelope" for the PKCS#10 certificate request.**
- [3.](#) Generate a unique string as the transaction id.**
- [4.](#) Generate a 16 byte random number as senderNonce.**
- [5.](#) Generate message digest on the enveloped PKCS#10 certificate request using the selected digest algorithm.**
- [6.](#) Create SignedData by adding the requestor's self- or CA-certificate as the signer's public key certificate. Include the message type, transaction id, the senderNonce and the message digest as the authenticated attributes and sign the attributes using the end entity's private key. This completes the SignedData.**
- [7.](#) The SignedData is prepended with the ContentInfo blob which indicates a SignedData object. This final step completes the create of a complete PKCSReq PKI message.**

In the following, the PKCSReq message is defined following the ASN.1 notation.

For readability, the values of a field is either represented by a quoted

string which specifies the intended value, or a constant when the value is known.

```
-- PKCSReq information portion
pkcsCertReq CertificationRequest ::= { -- PKCS#10
    version 0
    subject "the requestor's subject name"
    subjectPublicKeyInfo {
        algorithm {pkcs-1 1} -- rsa encryption
        subjectPublicKey "DER encoding of the requestor's public key"
    }
    attributes {
        challengePassword {{pkcs-9 7} "password string" }
        extensions
    }
    signatureAlgorithm {pkcs-1 4} -- MD5WithRSAEncryption
    signature "bit string which is created by signing inner content
              of the defined pkcsCertReq using requestor's private
              key, corresponding to the public key included in
              subjectPublicKeyInfo."
}

-- Enveloped information portion
pkcsCertReqEnvelope EnvelopeData ::= { -- PKCS#7
    version 0
    recipientInfo {
        version 0
        issuerAndSerialNumber {
            issuer "the CA issuer name"
            serialNumber "the CA certificate serial number"
        }
        keyEncryptionAlgorithm {pkcs-1 1} -- rsa encryption
        encryptedKey "content-encryption key
                     encrypted by CA public key"
    }
    encryptedContentInfo {
        contentType {pkcs-7 1} -- data content
        contentEncryptionAlgorithm "object identifier
                                   for DES encryption"
        encryptedContent "encrypted pkcsCertReq using the content-
                          encryption key"
    }
}

-- Signed PKCSReq
pkcsCertReqSigned SignedData ::= { -- PKCS#7
    version 1
    digestAlgorithm {iso(1) member-body(2) US(840) rsadsi(113549)
                    digestAlgorithm(2) 5}
    contentInfo {
        contentType {pkcs-7 1} -- data content identifier
        content pkcsCertReqEnvelope
    }
    certificate { -- requestor self-signed or CA-issued certificate
```

```
version 3
serialNumber "the transaction id associated with enrollment"
signature {pkcs-1 4} -- md5WithRSAEncryption
```

```

    issuer " the requestor's subject name"
    validity {
        notBefore "a UTC time"
        notAfter  "a UTC time"
    }
    subject "the requestor's subject name"
    subjectPublicKeyInfo {
        algorithm {pkcs-1 1}
        subjectPublicKey "DER encoding of requestor's public key"
    }
    signatureAlgorithm {pkcs-1 4}
    signature "the signature generated by using the requestor's
        private key corresponding to the public key in
        this certificate."
}
signerInfo {
    version 1
    issuerAndSerialNumber {
        issuer "the requestor's subject name"
        serialNumber "the transaction id associated
            with the enrollment"
    }
    digestAlgorithm {iso(0) member-body(2) US(840) rsadsi(113549)
        digestAlgorithm(2) 5}
    authenticateAttributes {
        contentType {{pkcs-9 3} {pkcs-7 1}}
        messageDigest {{pkcs-9 4} "an octet string"}
        transaction-id {{id-attributes transId(7)} "printable
            string"}
            -- this transaction id will be used
            -- together with the subject name as
            -- the identifier of the requestor's key
            -- pair during enrollment
        messageType {{id-attributes messageType(2)} "PKCSReq"}
        senderNonce {{id-attributes senderNonce(5)}
            "a random number encoded as a string"}
    }
    digestEncryptionAlgorithm {pkcs-1 1} -- rsa encryption
    encryptedDigest "encrypted digest of the authenticated
        attributes using requestor's private key"
}
}
pkcsReq PKIMessage ::= {
    contentType {pkcs-7 2}
    content pkcsCertRepSigned
}

```



### [5.1.2](#) CertRep Message Format

The response to an SCEP enrollment request is a CertRep message.

#### [5.1.2.1](#) PENDING Response

When the CA is configured to manually authenticate the requestor, the CertRep is returned with the attribute `pkiStatus` set to `PENDING`. The data portion for this message is null. Only the transaction required attributes are sent back.

```

CertRepSigned SignedData ::= { -- PKCS#7
    version 1
    digestAlgorithm {iso(1) member-body(2) US(840) rsadsi(113549)
        digestAlgorithm(2) 5}
    contentInfo {contentType {pkcs-7 1} -- empty content
    }
    signerInfo {
        version 1
        issuerAndSerialNumber {
            issuer "name of CA that issued the CA [RA] cert"
            serialNumber "the serial number of the CA [RA] cert"
        }
        digestAlgorithm {iso(1) member-body(2) US(840) rsadsi(113549)
            digestAlgorithm(2) 5}
        authenticateAttributes {
            contentType {{pkcs-9 3} {pkcs-7 1}}
            messageDigest {{pkcs-9 4} NULL}
            messageType {{id-attribute messageType(0)} "CertRep"}
            transaction-id {{id-attributes transid(7)} "printablestring"}
            --- same transaction id used in PKCSReq
            pkiStatus {{id-attributes pkiStatus(3)} "PENDING"}
            recipientNonce {{id-attributes recipientNonce(6)} <16 bytes>}
            senderNonce {{id-attributes senderNonce(5)} <16 bytes>}
        }
        digestEncryptionAlgorithm {pkcs-1 1}
        encryptedDigest "encrypted message digest of the authenticated
            attributes using the CA's [RA's] private key"
    }
}

CertRep PKIMessage ::= {
    contentType {pkcs-7 2}
    content CertRepSigned
}

```

#### [5.1.2.2](#) Failure Response

In this case, the CertRep sent back to the requestor is same as in the `PENDING` case, except that the `pkiStatus` attribute is set to `FAILURE`,

and the failInfo attribute should be included:

```
pkistatus {{id-attributes pkiStatus(3)}} "FAILURE"  
failInfo {{id-attributes failInfo(4)}} "the reason to reject"}
```



**5.1.2.3 SUCCESS response**

In this case, the information portion of CertRep will be a degenerated PKCS#7 which contains the requestor's certificate. It is then enveloped and signed as below:

```
pkcsCertRep SignedData ::= { -- PKCS#7
  version 1
  digestAlgorithm {iso(1) member-body(2) US(840) rsadsi(113549)
    digestAlgorithm(2) 5}
  contentInfo { -- empty content since this is degenerated PKCS#7
    contentType {pkcs-7 1}
  }
  certificates {
    certificate { -- issued requestor's certificate // must be first
      version 3
      serialNumber "issued requestor's certificate serial number"
      signature {pkcs-1 4} -- md5WithRSAEncryption
      issuer "the certificate authority issuer name"
      validity {
        notBefore "UTC time"
        notAfter "UTC time"
      }
      subject "the requestor subject name as given in PKCS#10"
      subjectPublicKeyInfo {
        algorithm {pkcs-1 1}
        subjectPublicKey "a DER encoding of requestor public
          key as given in PKCS#10"
      }
      extensions " the extensions as given in PKCS#10"
      signatureAlgorithm {pkcs-1 4}
      signature " the certificate authority signature"
    }
    certificate "the certificate authority certificate" (optional)
    certificate "the registration authority certificate(s)" (optional)
  }
}

pkcsCertRepEnvelope EnvelopedData ::= { -- PKCS#7
  version 0
  recipientInfo {
    version 0
    issuerAndSerialNumber { -- use issuer name and serial number as
      -- conveyed in requestor's self-signed
      -- certificate, included in the PKCSReq
    }
    issuer "the requestor's subject name"
    serialNumber "the serial number defined by the requestor in
      its self-signed certificate"
  }
  keyEncryptionAlgorithm {pkcs-1 1}
```

```
        encryptedKey "content-encrypt key encrypted by the requestor's  
                        public key which is same key as authenticated in  
                        the requestor's certificate"  
    }
```

```

    encryptedContentInfo {
        contentType {pkcs-7 1} -- data content identifier
        contentEncryptionAlgorithm "OID for DES encryption"
        encryptedContent "encrypted pkcsCertRep using content encryption
                        key"
    }
}
pkcsCertRepSigned SignedData ::= { -- PKCS#7
    version 1
    digestAlgorithm {iso(1) member-body(2) US(840) rsadsi(113549)
                    digestAlgorithm(2) 5}
    contentInfo {
        contentType {pkcs-7 1}
        content pkcsCertRepEnvelope
    }
    signerInfo {
        version 1
        issuerAndSerialNumber {
            issuer "the certificate authority issuer name"
            serialNumber "the CA certificate's serial number"
        }
        digestAlgorithm {iso(1), member-body(2) US(840) rsadsi(113549)
                        digestAlgorithm(2) 5}
        authenticateAttributes {
            contentType {{pkcs-9 3} {pkcs-7 1}}
            messageDigest {{pkcs-9 4} "a octet string"}
            messageType {{id-attribute messageType(2)} "CertRep"}
            transaction-id {{id-attributes transId(7)} "printable
                                string"}
                                -- same transaction id as given in PKCSReq
            pkiStatus {{id-attributes pkiStatus(3)} "SUCCESS"}
            recipientNonce {{id-attribute recipientNonce(6)}<16 bytes>}
            senderNonce {{ id-attributes senderNonce(5) <16 bytes>}}
        }
        digestEncryptionAlgorithm {pkcs-1 1}
        encryptedDigest "encrypted digest of authenticate attributes
                        using CA's private key "
    }
}
CertRep PKIMessage ::= {
    contentType {pkcs-7 2}
    content pkcsCertRepSigned
}

```

## 5.2 Poll for Requestor Initial Certificate

Either triggered by the PENDING status received from the CertRep, or by the non-response timeout for the previous PKCSReq, a requestor will enter the polling state by periodically sending GetCertInitial to the

server, until either the request is granted and the certificate is sent back, or the request is rejected, or the the configured time limit for polling is exceeded.

Since GetCertInitial is part of the enrollment, the messages exchanged during the polling period should carry the same transaction identifier as the previous PKCSReq.

#### PreCondition

Either the requestor has received a CertRep with pkiStatus set to be PENDING, or the previous PKCSReq has timed out.

#### PostCondition

The requestor has either received the certificate, or been rejected of its request, or the polling period ended as a failure.

### 5.2.1 GetCertInitial Message Format

Since at this time the certificate has not been issued, the requestor can only use the requestor's subject name, combined with the transaction identifier, to identify the polled certificate request.

The certificate authority server must be able to uniquely identify the polled certificate request. A subject name can have more than one outstanding certificate request (with different key usage attributes).

-- Information portion

```
pkcsGetCertInitial issuerAndSubject ::= {
    issuer "the certificate authority issuer name"
    subject "the requestor subject name as given in PKCS#10"
}
pkcsGetCertInitialEnvelope EnvelopedData ::= {
    version 0
    recipientInfo {
        version 0
        issuerAndSerialNumber {
            issuer "the CA issuer name"
            serialNumber "the CA certificate serial number"
        }
        keyEncryptionAlgorithm {pkcs-1 1}
        encryptedKey "content-encrypt key encrypted by CA's public key"
    }
    encryptedContentInfo {
        contentType {pkcs-7 1} -- data content
        contentEncryptionAlgorithm "OID for DES encryption"
        encryptedContent "encrypted getCertInitial"
    }
}
pkcsGetCertInitialSigned SignedData ::= { -- PKCS#7
    version 1
    digestAlgorithm {iso(1) member-body(2) US(840) rsadsi(113549)}
```

```
                digestAlgorithm(2) 5}  
contentInfo {  
    contentType {pkcs-7 1}
```

```

    content pkcsGetCertInitialEnvelope
}
certificate { -- the requestor's self-signed certificate
    version 3
    serialNumber "the transaction id associated with enrollment"
    signature {pkcs-1 4} -- md5WithRSAEncryption
    issuer "the requestor's subject name"
    validity {
        notBefore "a UTC time"
        notAfter "a UTC time"
    }
    subject "the requestor's subject name"
    subjectPublicKeyInfo {
        algorithm {pkcs-1 1}
        subjectPublicKey "DER encoding of requestor's public key"
    }
    signatureAlgorithm {pkcs-1 4}
    signature "the signature generated by using the requestor's
        private key corresponding to the public key in
        this certificate."
}
signerInfo {
    version 1
    issuerAndSerialNumber {
        issuer "requestor's subject name"
        serialNumber "the transaction id used in previous PKCSReq"
    }
    digestAlgorithm {iso(1), member-body(2) US(840) rsadsi(113549)
        digestAlgorithm(2) 5}
    authenticateAttributes {
        contentType {{pkcs-9 3} {pkcs-7 1}}
        messageDigest {{pkcs-9 4} "an octet string"}
        -- digest of getCertInitial
        messageType {{id-attribute messageType(2)} "GetCertInitial"}
        transaction-id {{id-attributes transId(7)} "printable
            string"}
        -- same transaction id used in previous PKCSReq
        senderNonce {{id-attribute senderNonce(3)} 0x<16 bytes>}
    }
    digestEncryptionAlgorithm {pkcs-1 1}
    encryptedDigest "encrypted digest of authenticateAttributes"
}
}
GetCertInitial PKIMessage ::= {
    contentType {pkcs-7 2}
    content pkcsGetCertInitialSigned
}

```

### **5.2.2 GetCertInitial Response Message Format**

The response messages for GetCertInitial are the same as for PKCSReq.



### 5.3 Certificate Access

The certificate query message defined in this section is an option when the LDAP server is not available to provide the certificate query. An requestor should be able to query an issued certificate from the certificate authority, as long as the issuer name and the issuer assigned certificate serial number is known to the requesting end entity. This transaction is not intended to provide the service as a certificate directory service. A more complicated query mechanism would have to be defined in order to allow a requestor to query a certificate using various different fields.

This transaction consists of one GetCert message send to the server by a requestor, and one CertRep message send back from the server.

#### PreCondition

The queried certificate have been issued by the certificate authority and the issuer assigned serial number is known.

#### PostContition

Either the certificate is send back or the request is rejected.

#### 5.3.1 GetCert Message Format

The queried certificate is identified by its issuer name and the issuer assigned serial number. If this is a query for an arbitrary requestor's certificate, the requesting requestor should includes its own CA issued certificate in the signed envelope. If this is a query for its own certificate (assume the requestor lost the issued certificate, or does not have enough non-volatile memory to save the certificate), then the self-signed certificate has to be included in the signed envelope.

```
pkcsGetCert issuerAndSerialNumber ::= {
    issuer "the certificate issuer name"
    serialNumber "the certificate serial number"
}
pkcsGetCertEnvelope EnvelopedData ::= {
    version 0
    recipientInfo {
        version 0
        issuerAndSerialNumber {
            issuer "the CA [RA] issuer name"
            serialNumber "the CA [RA] certificate serial number"
        }
        keyEncryptionAlgorithm {pkcs-1 1}
        encryptedKey "content-encrypt key encrypted
                     by CA [RA] public key"
    }
}
```



```
    encryptedContentInfo {
        contentType {pkcs-7 1} -- data content
        contentEncryptionAlgorithm "OID for DES encryption"
        encryptedContent "encrypted pkcsGetCert using the content
                           encryption key"
    }
}
pkcsGetCertSigned SignedData ::= {
    version 1
    digestAlgorithm {iso(1) member-body(2) US(840) rsadsi(113549)
                    digestAlgorithm(2) 5}
    contentInfo {
        contentType {pkcs-7 1}
        content pkcsGetCertEnvelope
    }
    certificates {
        certificate "CA issued certificate"
                   or "self-signed certificate"
    }
    signerInfo {
        version 1
        issuerAndSerialNumber {
            issuer "the requestor's subject name"
            serialNumber "requestor's certificate serial number"
        }
        digestAlgorithm {iso(1), member-body(2) US(840) rsadsi(113549)
                        digestAlgorithm(2) 5}
        authenticateAttributes {
            contentType {{pkcs-9 3} {pkcs-7 1}}
            messageDigest {{pkcs-9 4} "an octet string"}
                           -- digest of pkcsGetCertEnvelope
            messageType {{id-attribute messageType(2)} "GetCert"}
            transaction-id {{id-attributes transId(7)} "printable
                           string"}
            senderNonce {{id-attribute senderNonce(3)} <16 bytes>}
        }
        digestEncryptionAlgorithm {pkcs-1 1}
        encryptedDigest "encrypted digest of authenticateAttributes"
    }
}
GetCert PKIMessage ::= {
    contentType {pkcs-7 2}
    content pkcsGetCertSigned
}
```



### **5.3.2 CertRep Message Format**

In this case, the CertRep from the server is same as the CertRep for the PKCSReq, except that the server will only either grant the request or reject the request. Also, the recipientInfo should use the CA issuer name and CA assigned serial number to identify the requestor's key pair since at this time, the requestor has received its own certificate.

## **5.4 CRL Access**

The CRL query message defined in this section is an option when the LDAP server is not available to provide the CRL query. In the PKI protocol proposed here, only the requestor can initiate the transaction to download CRL. A requestor send GetCRL request to the server and the server send back CertRep whose information portion is a degenerated PKCS#7 which contains only the most recent CRL. The size of CRL included in the CertRep should be determined by the implementation.

### **PreCondition**

The certificate authority certificate has been downloaded to the end entity.

### **PostCondition**

CRL send back to the requestor.

### **5.4.1 GetCRL Message format**

The CRL is identified by using both CA's issuer name and the CA certificate's serial number:

```
pkcsGetCRL issuerAndSerialNumber {  
    issuer "the certificate authority issuer name"  
    serialNumber "certificate authority certificate's serial number"  
}
```

When the CRLDistributionPoint is supported, the pkcsGetCRL is defined as the following:

```
pkcsGetCRL SEQUENCE {  
    crlIssuer issuerAndSerialNumber  
    distributionPoint CE-CRLDistPoints  
}
```

where CE-CRLDisPoints is defined in X.509, but must contain only one CRL distribution point.



```
pkcsGetCRLEnvelope EnvelopedData ::= {
  version 0
  recipientInfo {
    version 0
    issuerAndSerialNumber {
      issuer "the certificate authority (or RA) issuer name"
      serialNumber "the CA (RA) certificate's serial number"
    }
    keyEncryptionAlgorithm {pkcs-1 1}
    encryptedKey "content-encrypt key encrypted by CA (RA) public key"
  }
  encryptedContentInfo {
    contentType {pkcs-7 1} -- data content
    contentEncryptionAlgorithm "OID for DES encryption"
    encryptedContent "encrypted pkcsGetCRL"
  }
}

pkcsGetCRLSigned SignedData ::= {
  version 1
  digestAlgorithm {iso(1) member-body(2) US(840) rsadsi(113549)
    digestAlgorithm(2) 5}
  contentInfo {
    contentType {pkcs-7 1}
    content pkcsGetCRLEnvelope
  }
  certificates {
    certificate "CA-issued or self-signed requestor's certificate"
  }
  signerInfo {
    version 1
    issuerAndSerialNumber {
      issuer "the requestor's issuer name"
      serialNumber "the requestor's certificate serial number"
    }
    digestAlgorithm {iso(1), member-body(2) US(840) rsadsi(113549)
      digestAlgorithm(2) 5}
    authenticateAttributes {
      contentType {{pkcs-9 3} {pkcs-7 1}}
      messageDigest {{pkcs-9 4} 0x<16/20 bytes>}
        -- digest of pkcsGetCRLEnvelope
      messageType {{id-attribute messageType(2)} "CertCRL"}
      transaction-id {{id-attributes transId(7)} "printable
        string"}
      senderNonce {{id-attribute senderNonce(3)} <16 bytes>}
    }
    digestEncryptionAlgorithm {pkcs-1 1}
    encryptedDigest "encrypted digest of authenticateAttributes"
  }
}
```

```
GetCRL PKIMessage ::= {  
    contentType {pkcs-7 2}  
    content pkcsGetCRLSigned  
}
```



#### 5.4.2 CertRep Message Format

The CRL is send back to the requestor through CertRep message. The information portion of this message is a degenerated PKCS#7 SignedData which contains only a CRL.

```
pkcsCertRep SignedData ::= {
  version 1
  digestAlgorithm {iso(1) member-body(2) US(840) rsadsi(113549)
    digestAlgorithm(2) 5}
  contentInfo {
    contentType {pkcs-7 1}
  }
  crl {
    signature {pkcs-1 4}
    issuer "the certificate authority issuer name"
    lastUpdate "UTC time"
    nextUpdate "UTC time"
    revokedCertificate {
      -- the first entry
      userCertificate "certificate serial number"
      revocationData "UTC time"
      ....
      -- last entry
      userCertificate "certificate serial number"
      revocationData "UTC time"
    }
  }
}

pkcsCertRepEnvelope EnvelopedData ::= {
  version 0
  recipientInfo {
    version 0
    issuerAndSerialNumber {
      issuer "the requestor's issuer name"
      serialNumber "the requestor certificate serial number"
    }
    keyEncryptionAlgorithm {pkcs-1 1}
    encryptedKey "content-encrypt key encrypted by requestor's
      public key "
  }
  encryptedContentInfo {
    contentType {pkcs-7 1} -- data content
    contentEncryptionAlgorithm "OID for DES encryption"
    encryptedContent "encrypted pkcsCertRep using requestor's
      public key"
  }
}
```



```

pkcsCertRepSigned SignedData ::= { -- PKCS#7
    version 1
    digestAlgorithm {iso(1) member-body(2) US(840) rsadsi(113549)
        digestAlgorithm(2) 5}
    contentInfo {
        contentType {pkcs-7 1}
        content pkcsCertRepEnvelope
    }
    signerInfo {
        version 1
        issuerAndSerialNumber {
            issuer "the certificate authority issuer name"
            serialNumber "the CA certificate's serial number"
        }
        digestAlgorithm {iso(1), member-body(2) US(840) rsadsi(113549)
            digestAlgorithm(2) 5}
        authenticateAttributes {
            contentType {{pkcs-9 3} {pkcs-7 1}}
            messageDigest {{pkcs-9 4} "an octet string"}
                -- digest of pkcsCertRepEnvelope
            messageType {{id-attribute messageType(2)} "CertRep"}
            transaction-id {{id-attributes transId(7)} "printable
                string"}
                -- same transaction id as given in PKCSReq
            pkiStatus {{id-attributes pkiStatus(3)} "SUCCESS"}
            recipientNonce{{id-attribute recipientNonce(6)}<16 bytes>}
            senderNonce {{id-attribute senderNonce (5) 0x<16 bytes>}}
        }
        digestEncryptionAlgorithm {pkcs-1 1}
        encryptedDigest "encrypted digest of authenticatedAttributes
            using CA private key"
    }
}

```

NOTE: The PKCS#7 EncryptedContent is specified as an octet string, but SCEP entities must also accept a sequence of octet strings as a valid alternate encoding.

This alternate encoding must be accepted wherever PKCS #7 Enveloped Data is specified in this document.



## **5.5 Get Certificate Authority Certificate**

Before any transaction begins, end entities have to get the CA (and possibly RA) certificate(s) first. Since the requestor may have no CA certificates or CA public keys at all, this message can not be encrypted and the response must be authenticated by out-of-band means. These certs are obtained by means of an HTTP GET message. To get the CA certificate, the requestor does a "HTTP GET" with a URL that identifies a CGI script on the server and an optional CA issuer identifier as the parameter to the CGI script. The response is either a single X.509 CA certificate ("CA mode"), or a PKCS7 message containing the CA certificate and RA certificates ("RA mode"). The client can determine which mode the CA operates in by which response it gets. Once the CA certificate is received by the requestor, a fingerprint is generated using MD5 hash algorithm on the whole CA certificate. If the requestor does not have a certificate path to a trusted root certificate, this fingerprint may be used to verify the certificate, by some positive out-of-band means, such as a phone call.

### **5.5.1 GetCACert HTTP Message Format**

"GET" CGI-PATH CGI-PROG "?operation=GetCACert" "&message=" CA-IDENT where:

CGI-PATH defines the actual CGI path to invoke the CGI program which parses the request.

CGI-PROG is set to be the string "pkiclient.exe" and this is expected to be the program that the CA will use to handle the SCEP transactions.

CA-IDENT is any string which is understood by the CA.

For example, it could be a domain name like ietf.org.

If a certificate authority has multiple root certificates this field can be used to distinguish which is required.

Otherwise it may be ignored.

### **5.5.2 Response**

The response for GetCACert is different between the case where the CA directly communicated with the requestor during the enrollment, and the case where a RA exists and the requestor communicates with the RA during the enrollment.

#### **5.5.2.1 CA Certificate Only Response**

A binary X.509 CA certificate is send back as a MIME object with a Content-Type of application/x-x509-ca-cert.

#### **5.5.2.2 CA and RA Certificates Response**

When an RA exists, both CA and RA certificates must be sent back in the response to the GetCACert request. The RA certificate(s) must be

signed by the CA. A certificates-only PKCS#7 SignedData is used to carry the certificates to the requestor, with a Content-Type of application/x-x509-ca-ra-cert.

## 5.6 Get Certificate Authority Certificate Chain

In order to support Certificate Authority hierarchies, it is desirable to have a way to get the entire certificate chain. The following message has been added to SCEP for this purpose.

### [5.6.1](#) GetCACertChain HTTP Message Format

```
"GET" CGI-SCRIPT "?" "operation=GetCACertChain" "&" "message" CA-IDENT
where CGI-SCRIPT and CA-IDENT are as described for GetCACert.
```

### [5.6.2](#) Response

The response for GetCACertChain is a certificates-only PKCS#7 SignedData to carry the certificates to the requestor, with a Content-Type of application/x-x509-ca-ra-cert-chain.

### [5.6.3](#) Backwards Compatability

Versions of SCEP prior to revision 3 do not support GetCACertChain. Certificate Authorities written to these prior versions will not be able to process the message and may return an HTML error.

To avoid this, clients should send the GetCACert message first. If the returned certificate is self-signed or is signed by a Certificate Authority that is trusted by the client, then it is not necessary to send the GetCACertChain message and it should not be sent.

If a Certificate Authority is configured with a certificate that is not either self-signed or has a self-signed issuer, then it should support this message. In other words, it should be supported if the CA hierarchy is more than two-deep.

An old CA in a two-deep hierarchy might still get this message from a client if the client did not trust either that CA or its issuer. In that event, the certificate cannot be trusted anyway. In any case the CA must not crash or hang upon the receipt of the message and the client must be able to handle whatever error is returned by the CA, including an HTML error or an ungraceful disconnect.





The following is the ASN.1 definition of Cert-Only PKCS#7:

```
certOnly SignedData ::= {
    version 1
    digestAlgorithm {iso(1) member-body(2) US(840) rsadsi(113549)
        digestAlgorithm(2) 5}

contentInfo {
    contentType {pkcs-7 1} -- data content identifier
    content -- NULL
}
certificates -- the RA and CA certificates.
}

CARACerts PKIMessage ::= { -- special pki message sent in the clear
    contentType {pkcs-7 2}
    content certOnly
}
```

## **[6.0 Security Considerations](#)**

This entire document is about security. Common security considerations such as keeping private keys truly private and using adequate lengths for symmetric and asymmetric keys must be followed in order to maintain the security of this protocol.

## **[7.0 Intellectual Property](#)**

This protocol includes the optional use of Certificate Revocation List Distribution Point (CRLDP) technology, which is a patented technology of Entrust Technologies, Inc. (Method for Efficient Management of Certificate Revocation Lists and Update Information (U.S. Patent 5,699,431)). Please contact Entrust Technologies, Inc. ([www.entrust.com](http://www.entrust.com)) for more information on licensing CRLDP technology.

## **[8.0 References](#)**

[PKCS7] Kaliski, B., "PKCS #7: Cryptographic Message Syntax Version 1.5", [RFC 2315](#), March 1998.

[PKCS10] Kaliski, B., "PKCS #10: Certification Request Syntax Version 1.5", [RFC 2314](#), March 1998.

[RFC2459] Housley, R., et al., "Internet X.509 Public Key Infrastructure Certificate and CRL Profile", [RFC 2459](#), January 1999.



## Appendix A: Cisco Requestor Subject Name Definition

The ip address and the FQDN of a SCEP client should be included in the V3 extension subjectAltName. When the subjectAltName extension attribute is present, both the subjectAltName fields and the subjectName field could have the IP address and the FQDN information.

When the X.500 directory is used by the CA to define the name space, the subject name defined above become a RDN which is part of DN binded to the requestor's public key in the certificate.

A sample of DN assigned by Entrust CA is given below (assume the same ciscoRouterAlice is used as the requestor defined subject name):

OU = InteropTesting, O = Entrust Technologies, C = CA  
RDN = {"alice.cisco.com", "172.21.114.67", "22334455"}

## Appendix B: IPSEC Client Enrollment Certificate Request

The following is the certificate enrollment request (PKCS#10) as created by Cisco VPN Client:

```

-----END NEW CERTIFICATE REQUEST-----
 0 30 439: SEQUENCE {
 4 30 288:   SEQUENCE {
 8 02  1:     INTEGER 0
11 30  57:   SEQUENCE {
13 31  55:     SET {
15 30  53:       SEQUENCE {
17 06  3:         OBJECT IDENTIFIER commonName (2 5 4 3)
22 13 46:         PrintableString
                :         'For Xiaoyi, IPSEC attrs in alternate name
                :         extn'
                :       }
                :     }
                :   }
70 30 158: SEQUENCE {
73 30  13:   SEQUENCE {
75 06  9:     OBJECT IDENTIFIER rsaEncryption (1 2 840 113549 1
                :         1 1)
86 05  0:     NULL
                :   }
88 03 140:   BIT STRING 0 unused bits
                :     30 81 88 02 81 80 73 DB 1D D5 65 AA EF C7 D4 8E
                :     AA 6E EB 46 AC 91 2A 0F 50 51 17 AD 50 A2 2A F2
                :     CE BE F1 E4 22 8C D7 61 A1 6C 87 61 62 92 CB A6
                :     80 EA B4 0F 09 9D 18 5F 39 A3 02 0E DB 38 4C E4
                :     8A 63 2E 72 8B DC BE 9E ED 6C 1A 47 DE 13 1B 0F
                :     83 29 4D 3E 08 86 FF 08 2B 43 09 EF 67 A7 6B EA
                :     77 62 30 35 4D A9 0F 0F DF CC 44 F5 4D 2C 2E 19
                :     E8 63 94 AC 84 A4 D0 01 E1 E3 97 16 CD 86 64 18
                :     [ Another 11 bytes skipped ]
                :   }
231 A0 63: [0] {
233 30 61:   SEQUENCE {
235 06  9:     OBJECT IDENTIFIER extensionReq (1 2 840 113549 1 9
                :         14)
246 31 48:     SET {
248 30 46:       SEQUENCE {
250 30 44:         SEQUENCE {
252 06  3:           OBJECT IDENTIFIER subjectAltName (2 5 29 17)
257 04 37:           OCTET STRING
                :             30 23 87 04 01 02 03 04 81 0D 65 6D 61 69

```

```

        6C 40 69 72 65 2E 63 6F 6D 82 0C 66 71 64
        6E 2E 69 72 65 2E 63 6F 6D
:
:   }
:   }
:   }
:   }
:   }
:   }
:   }

296 30   13:   SEQUENCE {
298 06    9:     OBJECT IDENTIFIER md5withRSAEncryption (1 2 840 113549
                                     1 1 4)

309 05    0:     NULL
:         }
311 03  129:   BIT STRING 0 unused bits
:         19 60 55 45 7F 72 FD 4E E5 3F D2 66 B0 77 13 9A
:         87 86 75 6A E1 36 C6 B6 21 71 68 BD 96 F0 B4 60
:         95 8F 12 F1 65 33 16 FD 46 8A 63 19 90 40 B4 B7
:         2C B5 AC 63 17 50 28 F0 CD A4 F0 00 4E D2 DE 6D
:         C3 4F F5 CB 03 4D C8 D8 31 5A 7C 01 47 D2 2B 91
:         B5 48 55 C8 A7 0B DD 45 D3 4A 8D 94 04 3A 6C B0
:         A7 1D 64 74 AB 8A F7 FF 82 C7 22 0A 2A 95 FB 24
:         88 AA B6 27 83 C1 EC 5E A0 BA 0C BA 2E 6D 50 C7
:         }

```

## Appendix C: Private OID Definitions

The OIDs used in defining pkiStatus are VeriSign self-maintained OIDs. Please note, work is in progress to replace the VeriSign owned object identifiers with the standard object identifiers. Once the standardization is completed, this documentation will be updated.

```

id-VeriSign   OBJECT_IDENTIFIER ::= {2 16 US(840) 1 VeriSign(113733)}
id-pki        OBJECT_IDENTIFIER ::= {id-VeriSign pki(1)}
id-attributes OBJECT_IDENTIFIER ::= {id-pki attributes(9)}
id-messageType OBJECT_IDENTIFIER ::= {id-attributes messageType(2)}
id-pkiStatus  OBJECT_IDENTIFIER ::= {id-attributes pkiStatus(3)}
id-failInfo   OBJECT_IDENTIFIER ::= {id-attributes failInfo(4)}
id-senderNonce OBJECT_IDENTIFIER ::= {id-attributes senderNonce(5)}
id-recipientNonce OBJECT_IDENTIFIER ::= {id-attributes recipientNonce(6)}
id-transId    OBJECT_IDENTIFIER ::= {id-attributes transId(7)}
id-extensionReq OBJECT_IDENTIFIER ::= {id-attributes extensionReq(8)}

```

#### [Appendix D:](#) CRL Query by means of LDAP

In order to retrieve the CRL by means of LDAP, the client needs to know where in the directory it is stored. The certificate must contain a CRL Distribution Point extension encoded as a DN or as an LDAP URI.

For example, the certificate issued by Entrust VPN contains the following DN as the CRL distribution point:

CN = CRL1, O = cisco, C = US.

The asn.1 encoding of this distribution point is:

```
30 2C 31 0B 30 09 06 03 55 04 06 13 02 55 53 31 0E 30 0C 06
03 55 04 0A 13 05 63 69 73 63 6F 31 0D 30 0B 06 03 55 04 03
13 04 43 52 4C 31
```

The ldap form would be:

ldap://servername/CN=CRL1,O=cisco,C=US

#### Appendix E: SCEP State Transitions

SCEP state transitions are based on transaction identifier. The design goal is to ensure the synchronization between the CA and the requestor under various error situations.

An identity is defined by the combination of FQDN, the IP address and the client serial number. FQDN is the required name attribute. It is important to notice that, a client named as Alice.cisco.com is different from the client named as Alice.cisco.com plus IPAddress 117.96.1.219.

Each enrollment transaction is uniquely associated with a transaction identifier. Because the enrollment transaction could be interrupted by various errors, including network connection errors or client reboot, the SCEP client generates a transaction identifier by calculating MD5 hash on the public key value for which the enrollment is requested. This retains the same transaction identifier throughout the enrollment transaction, even if the client has rebooted or timed out, and issues a new enrollment request for the same key pair. It also provides the way for the CA to uniquely identify a transaction in its database. At the requestor side, it generates a transaction identifier which is included in PKCSReq. If the CA returns a response of PENDING, the requestor will poll by periodically sending out GetCertInitial with the same

transaction identifier until either a response other than PENDING is obtained, or the configured maximum time has elapsed.

If the client times out or the client reboots, the client administrator will start another enrollment transaction with the same key pair. The second enrollment will have the transaction identifier. At the server side, instead of accepting the PKCSReq as a new enrollment request, it should respond as if another GetCertInitial message had been sent with that transaction ID. In another word, the second PKCSReq should be taken as a resynchronization message to allow the enrollment resume as the same transaction.

It is important to keep the transaction id unique since CEP requires the same policy and same identity be applied to the same subject name and

key pair binding. In the current implementation, an SCEP client can only assume one identity. At any time, only one key pair, with a given key usage, can be associated with the same identity.

The following gives several example of client to CA transactions.

Client actions are indicated in the left column, CA actions are indicated in the right column. A blank action signifies that no message was received. Note that these examples assume that the CA enforces the certificate-name uniqueness property defined in [Section 2.1.1.1](#).

The first transaction, for example, would read like this:

"Client Sends PKCSReq message with transaction ID 1 to the CA. The CA signs the certificate and constructs a CertRep Message containing the signed certificate with a transaction ID 1. The client receives the message and installs the cert locally."

Successful Enrollment Case: no manual authentication

```
PKCSReq (1)           -----> CA Signs Cert
Client Installs Cert  <----- CertRep (1) SIGNED CERT
```

Successful Enrollment Case: manual authentication required

```
PKCSReq (10)          -----> Cert Request goes into Queue
Client Polls          <----- CertRep (10) PENDING
GetCertInitial (10)   -----> Still pending
Client Polls          <----- CertRep (10) PENDING
GetCertInitial (10)   -----> Still pending
Client Polls          <----- CertRep (10) PENDING
GetCertInitial (10)   -----> Still pending
Client Polls          <----- CertRep (10) PENDING
GetCertInitial (10)   -----> Cert has been signed
Client Installs Cert  <----- CertRep (10) SIGNED CERT
```

Resync Case - CA Receive and Signs PKCSReq, Client Did not receive CertRep:

```
PKCSReq (3)           -----> Cert Request goes into queue
                       <----- CertRep (3) PENDING
GetCertInitial (3)     ----->
                       <----- CertRep (3) PENDING
GetCertInitial (3)     ----->
                       <----- CA signed Cert and send back
                               CertRep(3)
(Time Out)
PKCSReq (3)           -----> Cert already signed, send back to
                               client
```



Client Installs Cert <----- CertRep (3) SIGNED CERT

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Case when NVRAM is lost and client has to generate a new key pair, there is no change of name information:

```
PKCSReq (4)          -----> CA Signs Cert
Client Installs Cert <----- CertRep (4) SIGNED CERT
(Client loses Cert)
PKCSReq (5)          -----> There is already a valid cert with
                        this DN.
Client Admin Revokes <----- CertRep (5) OVERLAPPING CERT ERROR
PKCSReq (5)          -----> CA Signs Cert
Client Installs Cert <----- CertRep (5) SIGNED CERT
```

Case when client admin resync the enrollment using a different PKCS#10:

```
PKCSReq (6)          -----> CA Signs Cert
                        <----- CertRep (6) SIGNED CERT
(Client timeout and admin starts another enrollment with a different
 PKCS#10, but the same transaction id)
PKCSReq (6) with different PKCS#10
                        -----> There is already a valid cert with
                        this entity (by checking FQDN).
                        <----- CertRep (6) INVALID PKCS#10 CERT
                        ERROR
```

Client admin either revokes the existing cert  
or corrects the error by enrolling with  
the same PKCS#10 as the first PKCSReq(6)

```
PKCSReq (6)          -----> CA find the existing Cert
Client Installs Cert <----- CertRep (6) SIGNED CERT
```

Resync case when server is slow in response:

```
PKCSReq (13)         -----> Cert Request goes into Queue
                        <----- CertRep (13) PENDING
GetCertInitial       -----> Still pending
                        <----- CertRep (13) PENDING
GetCertInitial       -----> Still pending
                        <----- CertRep (13) PENDING
GetCertInitial       -----> Still pending
                        <----- CertRep (13) PENDING
GetCertInitial       -----> Still pending
                        <----- CertRep (13) PENDING
GetCertInitial       -----> Still pending
(TimeOut)            <----- CertRep (13) PENDING
* Case 1
PKCSReq (13)         -----> Still pending
Client polls         <----- CertRep (13) PENDING
CertCertInitial      -----> Cert has been signed
Client Installs Cert <----- CertRep (13) SIGNED CERT
* Case 2
PKCSReq (13)         -----> Cert has been signed
```

Client Installs Cert <----- CertRep (13) SIGNED CERT

**Appendix F. Author Contact Information**

Xiaoyi Liu  
Cisco  
**510 McCarthy Drive**  
Milpitas, CA  
xliu@cisco.com

Cheryl Madson  
Cisco  
510 McCarthy Drive  
Milpitas, CA.  
cmadson@cisco.com

David McGrew  
Cisco  
**170 West Tasman Drive**  
San Jose, CA 94134  
mcgrew@cisco.com

Andrew Nourse  
Cisco  
510 McCarthy Drive  
Milpitas, CA.  
nourse@cisco.com

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