Reliable and Available Wireless Architecture

Abstract

Due to uncontrolled interferences, including the self-induced multipath fading, deterministic networking can only be approached on wireless links. The radio conditions may change -way- faster than a centralized routing can adapt and reprogram, in particular when the controller is distant and connectivity is slow and limited. RAW separates the routing time scale at which a complex path is recomputed from the forwarding time scale at which the forwarding decision is taken for an individual packet. RAW operates at the forwarding time scale. The RAW problem is to decide, within the redundant solutions that are proposed by the routing, which will be used for each individual packet to provide a DetNet service while minimizing the waste of resources.

Status of This Memo

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1. Introduction

Bringing determinism in a packet network means eliminating the statistical effects of multiplexing that result in probabilistic jitter and loss. This can be approached with a tight control of the physical resources to maintain the amount of traffic within a budgetted volume of data per unit of time that fits the physical capabilities of the underlying technology, and the use of time-shared resources (bandwidth and buffers) per circuit, and/or by shaping and/or scheduling the packets at every hop.

Wireless networks operate on a shared medium where uncontrolled interference, including the self-induced multipath fading, adds another dimension to the statistical effects that affect the delivery. Scheduling transmissions can alleviate those effects by leveraging diversity in the spatial, time, code, and frequency domains, and provide a Reliable and Available service while preserving energy and optimizing the use of the shared spectrum.

Deterministic Networking is an attempt to mostly eliminate packet loss for a committed bandwidth with a guaranteed worst-case end-to-end latency, even when co-existing with best-effort traffic in a shared network. This innovation is enabled by recent developments in technologies including IEEE 802.1 TSN (for Ethernet LANs) and IETF DetNet (for wired IP networks). It is getting traction in various industries including manufacturing, online gaming, professional A/V, cellular radio and others, making possible many cost and performance optimizations.

The ["Deterministic Networking Architecture"](#) [[RFC8655](#)] is composed of three planes: the Application (User) Plane, the Controller Plane, and the Network Plane. Reliable and Available Wireless (RAW) extends DetNet to focus on issues that are mostly a co"ern on wireless links, and inherits the architecture and the planes. A RAW Network Plane is thus a Network Plane inherited by RAW from DetNet, composed of one or multiple hops of homogeneous or heterogeneous technologies, e.g. a Wi-Fi6 Mesh or one-hop CBRS access links federated by a 5G backhaul.

RAW networking aims at providing highly available and reliable end-to-end performances in a network with scheduled wireless segments. Uncontrolled interference and transmission obstacles may impede the

transmission, and techniques such as beamforming with Multi-User MIMO can only alleviate some of those issues, so the term "deterministic" is usually not associated with short range radios, in particular in the ISM band. This uncertainty places limits to the amount of traffic that can be transmitted on a link while conforming to a RAW Service Level Agreement (SLA) that may vary rapidly.

The wireless and wired media are fundamentally different at the physical level, and while the generic ["Deterministic Networking Problem Statement"](#) [[RFC8557](#)] applies to both the wired and the wireless media, the methods to achieve RAW must extend those used to support time-sensitive networking over wires, as a RAW solution has to address less consistent transmissions, energy conservation and shared spectrum efficiency.

The development of RAW technologies has been lagging behind deterministic efforts for wired systems both at the IEEE and the IETF. But recent efforts at the IEEE and 3GPP indicate that wireless is finally catching up at the lower layer and that it is now possible for the IETF to extend DetNet for wireless segments that are capable of scheduled wireless transmissions.

The intent for RAW is to provide DetNet elements that are specialized for short range radios. From this inheritance, RAW stays agnostic to the radio layer underneath though the capability to schedule transmissions is assumed. How the PHY is programmed to do so, and whether the radio is single-hop or meshed, are unknown at the IP layer and not part of the RAW abstraction.

Still, in order to focus on real-worlds issues and assert the feasibility of the proposed capabilities, RAW will focus on selected technologies that can be scheduled at the lower layers: IEEE Std. 802.15.4 timeslotted channel hopping (TSCH), 3GPP 5G ultra-reliable low latency communications (URLLC), IEEE 802.11ax/be where 802.11be is extreme high throughput (EHT), and L-band Digital Aeronautical Communications System (LDACS). See [[RAW-TECHNOS](#)] for more.

The establishment of a path is not in-scope for RAW. It may be the product of a centralized Controller Plane as described for DetNet. As opposed to wired networks, the action of installing a path over a set of wireless links may be very slow relative to the speed at which the radio conditions vary, and it makes sense in the wireless case to provide redundant forwarding solutions along a complex path and to leave it to the Network Plane to select which of those forwarding solutions are to be used for a given packet based on the current conditions.

RAW distinguishes the longer time scale at which routes are computed from the the shorter forwarding time scale where per-packet

decisions are made. RAW operates at the forwarding time scale on one
DetNet flow over one path that is preestablished and installed by
means outside of the scope of RAW. The scope of the RAW WG comprises
Network plane protocol elements such as OAM and in-band control to
improve the RAW operation at the Service and at the forwarding sub-
layers, e.g., controlling whether to use packet replication, Hybrid
ARQ and coding, with a constraint to limit the use of redundancy
when it is really needed, e.g., when a spike of loss is observed.
This is discussed in more details in [Section 4.3](#) and the next
sections.

2. Terminology

RAW defines the following terms:

PAREO: Packet (hybrid) ARQ, Replication, Elimination and Ordering.
   PAREO is a superset Of DetNet's PREOF that includes radio-
   specific techniques such as short range broadcast, MUMIMO,
   constructive interference and overhearing, which can be leveraged
   separately or combined to increase the reliability.

Flapping: In the context of RAW, a link flaps when the wireless
   connectivity is interrupted for short transient times, typically
   of a subsecond duration.

RAW reuses terminology defined for DetNet in the ["Deterministic
Networking Architecture"](#) [[RFC8655](#)], e.g., PREOF for Packet
Replication, Elimination and Ordering Functions.

RAW also reuses terminology defined for 6TiSCH in [[6TiSCH-ARCH](#)] such
as the term Track. 6TiSCH defined a Track as a complex path with
associated PAREO operations.

This document reuses terms that are well-defined in the context of
automation to networking and packet delivery, in particular for
reliability and availability. In the context of the RAW work, they
are defined as follows:

Reliability: Reliability is a measure of the probability that an
   item will perform its intended function for a specified interval
   under stated conditions. For RAW, the service that is expected is
   delivery within a bounded latency and a failure is when the
   packet is either lost or delivered too late. RAW expresses
   reliability in terms of Mean Time Between Failure (MTBF) and
   Maximum Consecutive Failures (MCF).

Availability: Availability is a measure of the relative amount of
   time where a path operates in stated condition, in other words
   (uptime)/(uptime+downtime). Because a serial wireless path may
   not be good enough to provide the required availability, and even

2 parallel paths may not be over a longer period of time, the RAW
availability implies a path that is a lot more complex than what
DetNet typically envisages (a Track).

3. Use Cases and Requirements Served

[RFC8578] presents a number of wireless use cases including Wireless
for Industrial Applications, Pro-Audio and SmartGrid. [RAW-USE-
CASES] adds a number of use cases that demonstrate the need for RAW
capabilities in new use cases including Pro-Gaming and drones.

The use cases can be abstracted in two families, radio access
protection and Track protection in a wireless mesh.

3.1. Radio Access Protection

In order to maintain the committed reliability at all times, a
wireless station may use more than one Radio Access Network (RAN) in
parallel.

The RANs may be heterogeneous, e.g., 5G [I-D.farkas-5g-raw] and Wi-
Fi [RAW-TECHNOS] for high-speed communication, in which case a
Layer-3 abstraction becomes useful to select which of the RANs are
used at a particular point of time, and the amount of traffic that
is distributed over each RAN.

The idea is that the rest of the path to the destination(s) is
protected separately (e.g., uses non-congruent paths) and/or is a
lot more reliable, e.g., wired. In that case, RAW observes
reliability of the path through each of the RANs but only acts on
the first hop.

```
                                        ***    *
                    RAN 1   -----   ***      **   **
                 /                *     **         ****
       +----+  /                 *              **      ***
       |    |-                   *                         **
       |    |--zzz- RAN 2 -- *         Internet        ***
       |    |-                   *                    ***
       +----+   \                 *                ***
              \                ***    **        **
                  RAN n   --------   ***    ***


      zzz = flapping now


                  Figure 1: Radio Access Protection
```

3.2. Track Protection in a Wireless Mesh

   A Track

```
                      A-------B-------C-----D
                     /  \   /         /         \
                    I ----M-------N--zzzzz--- E
                     \       \   /               /
                      P--zzz--Q------------R


                zzz = flapping now

                   Figure 2: Track Protection
```
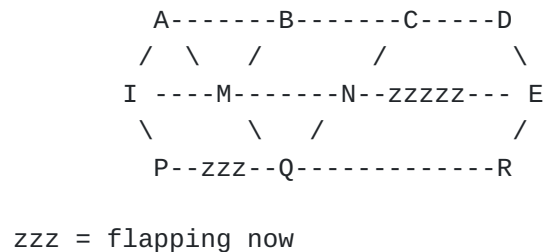
4. RAW Architecture Considerations

4.1. Related Work at The IETF

   RAW intersects with protocols or practices in development at the
   IETF as follows:

   *The Dynamic Link Exchange Protocol (DLEP) [RFC8175] from [MANET]
    can be leveraged at each hop to derive generic radio metrics
    (e.g., based on LQI, RSSI, queueing delays and ETX) on individual
    hops.

   *Operations, Administration and Maintenance (OAM) work at [DetNet]
    such as [DetNet-IP-OAM] for the case of the IP Data Plane
    observes the state of DetNet paths, typically MPLS and IPv6
    pseudowires [DetNet-DP-FW], in the direction of the traffic. RAW
    needs feedback that flows on the reverse path and gathers
    instantaneous values from the radio receivers at each hop to
    inform back the source and replicating relays so they can make
    optimized forwarding decisions. The work named ICAN may be
    related as well.

   *[BFD] detect faults in the path between an ingress and an egress
    forwarding engines, but is unaware of the complexity of a path
    with replication, and expects bidirectionality. BFD considers
    delivery as success whereas with RAW the bounded latency can be
    as important as the delivery itself.

   *[SPRING] and [BIER] define in-band signaling that influences the
    routing when decided at the head-end on the path. There's already
    one RAW-related draft at BIER [BIER-PREF] more may follow. RAW
    will need new in-band signaling when the decision is distributed,
    e.g., required chances of reliable delivery to destination within

latency. This signaling enables relays to tune retries and
replication to meet the required SLA.

*[CCAMP] defines protocol-independent metrics and parameters
(measurement attributes) for describing links and paths that are
required for routing and signaling in technology-specific
networks. RAW would be a source of requirements for CCAMP to
define metrics that are significant to the focus radios.

4.2. Prerequisites

A prerequisite to the RAW work is that an end-to-end routing
function computes a complex sub-topology along which forwarding can
happen between a source and one or more destinations. For 6TiSCH,
this is a Track. The concept of Track is specified in the 6TiSCH
Architecture [6TiSCH-ARCH]. Tracks provide a high degree of
redundancy and diversity and enable DetNet PREOF, end-to-end network
coding, and possibly radio-specific abstracted techniques such as
ARQ, overhearing, frequency diversity, time slotting, and possibly
others.

How the routing operation computes the Track is out of scope for
RAW. The scope of the RAW operation is one Track, and the goal of
the RAW operation is to optimize the use of the Track at the
forwarding timescale to maintain the expected service while
optimizing the usage of constrained resources such as energy and
spectrum.

Another prerequisite is that an IP link can be established over the
radio with some guarantees in terms of service reliability, e.g., it
can be relied upon to transmit a packet within a bounded latency and
provides a guaranteed BER/PDR outside rare but existing transient
outage windows that can last from split seconds to minutes. The
radio layer can be programmed with abstract parameters, and can
return an abstract view of the state of the Link to help forwarding
decision (think DLEP from MANET). In the layered approach, how the
radio manages its PHY layer is out of control and out of scope.
Whether it is single hop or meshed is also unknown and out of scope.

4.3. Routing Time Scale vs. Forwarding Time Scale

With DetNet, the end-to-end routing can be centralized and can
reside outside the network. In wireless, and in particular in a
wireless mesh, the path to the controller that performs the route
computation and maintenance expensive in terms of critical resources
such as air time and energy.

Reaching to the routing computation can also be slow in regards to
the speed of events that affect the forwarding operation at the
radio layer. Due to the cost and latency to perform a route

computation, the controller plane is not expected to be sensitive/
reactive to transient changes. The abstraction of a link at the
routing level is expected to use statistical operational metrics
that aggregate the behavior of a link over long periods of time, and
represent its availability as shades of gray as opposed to either up
or down.

```
                    +----------------+
                    |  Controller    |
                    |    (PCE)       |
                    |  [Routing ]    |
                    |  [Function]    |
                    +----------------+
                            ^
                            |
                          Slow
                            |
    _-._-._-._-._-._-._-.   |   ._-._-._-._-._-._-._-._-._-._-._-._-._-
     _-._-._-._-._-._-._-.  | _-._-._-._-._-._-._-._-._-._-._-._-._-
                            |
                        Expensive
                    ....    |  .......
                 ....     . | .       ....
              ....          v            ...
           ..    A-------B-------C---D      ..
         ...    / \   /       /     \      ..
        .      I ----M-------N--zzz-- E  ..
        ..       \       \   /         /     .
         ..      P--zzz--Q---------R    ..
          ..                          ..
           .......              ...
                .............
    zzz = flapping now
```
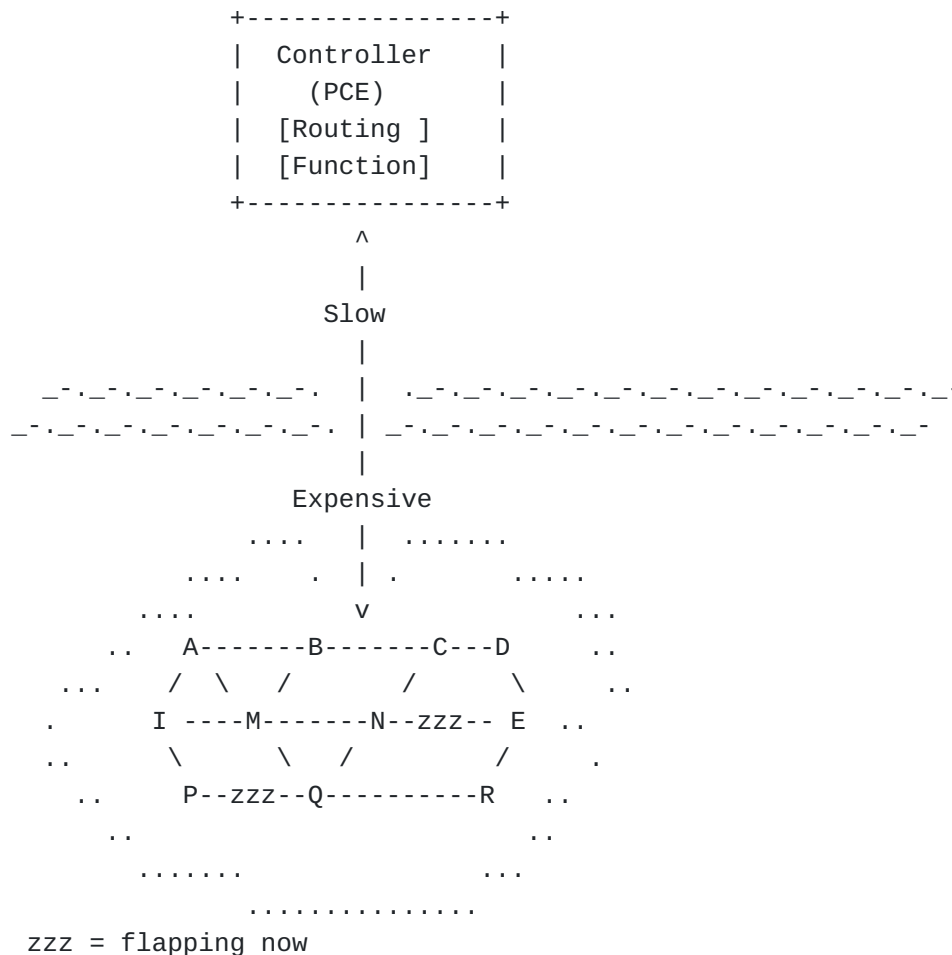
                    Figure 3: Time Scales

In the case of wireless, the changes that affect the forwarding
decision can happen frequently and often for short durations, e.g.,
a mobile object moves between a transmitter and a receiver, and will
cancel the line of sight transmission for a few seconds, or a radar
measures the depth of a pool and interferes on a particular channel
for a split second.

There is thus a desire to separate the long term computation of the
route and the short term forwarding decision. In such a model, the
routing operation computes a complex Track that enables multiple
Non-Equal Cost Multi-Path (N-ECMP) forwarding solutions, and leaves

it to the forwarding plane to make the per-packet decision of which
of these possibilities should be used.

In the case of wires, the concept is known in traffic engineering
where an alternate path can be used upon the detection of a failure
in the main path, e.g., using OAM in MPLS-TP or BFD over a
collection of SD-WAN tunnels. RAW formalizes a forwarding time scale
that is an order(s) of magnitude shorter than the controler plane
routing time scale, and separates the protocols and metrics that are
used at both scales. Routing can operate on long term statistics
such as delivery ratio over minutes to hours, but as a first
approximation can ignore flapping. On the other hand, the RAW
forwarding decision is made at packet speed, and uses information
that must be pertinent at the present time for the current
transmission.

4.4. Track Source-Routed vs. Distributed RAW Decision

Within a large routed topology, the route-over mesh operation builds
a particular complex Track with one source and one or more
destinations; within the Track, packets may follow different paths
and may be subject to RAW forwarding operations that include
replication, elimination, retries, overhearing and reordering.

The RAW forwarding decisions include the selection of points of
replication and elimination, how many retries can take place, and a
limit of validity for the packet beyond which the packet should be
destroyed rather than forwarded uselessly further down the Track.

The decision to apply the RAW techniques must be done quickly, and
depends on a very recent and precise knowledge of the forwarding
conditions within the complex Track. There is a need for an
observation method to provide the RAW forwarding plane with the
specific knowledge of the state of the Track for the type of flow of
interest (e.g., for a QoS level of interest). To observe the whole
Track in quasi real time, RAW will consider existing tools such as
L2-triggers, DLEP, BFD and in-band and out-of-band OAM.

One possible way of making the RAW forwarding decisions is to make
them all at the ingress and express them in-band in the packet,
which requires new loose or strict Hop-by-hop signaling. To control
the RAW forwarding operation along a Track for the individual
packets, RAW may leverage and extend known techniques such as DetNet
tagging, Segment Routing (SRv6) or BIER-TE such as done with [BIER-
PREF].

An alternate way is to enable each forwarding node to make the RAW
forwarding decisions for a packet on its own, based on its knowledge
of the expectation (timeliness and reliability) for that packet and

a recent observation of the rest of the way across the possible paths within the Track. Information about the service should be placed in the packet and matched with the forwarding node's capabilities and policies.

In either case, a per-flow state is installed in all intermediate nodes to recognize the flow and determine the forwarding policy to be applied.

5. RAW Architecture Components

5.1. Reliability and Availability

5.2. Wireless Tracks

The "6TiSCH Architecture" [6TiSCH-ARCH] introduces the concept of Track a a possibly complex path with the PAREO functions operated within.

A simple track is composed of a direct sequence of reserved hops to ensure the transmission of a single packet from a source node to a destination node across a multihop path.

A Complex Track is designed as a directed acyclic graph from a source node towards a destination node to support multi-path forwarding, as introduced in "6TiSCH Architecture" [6TiSCH-ARCH]. By employing Packet Replication and Elimination (PRE) functions [RFC8655], several paths may be computed, and these paths may be more or less independent. For example, a complex Track may branch off and rejoin over non-congruent paths (branches).

Some more details for Deterministic Network PRE techniques are presented in the following Section.

5.3. PAREO Functions

In a nutshell, Packet Replication and Elimination (PRE) establishes several paths in a network to provide redundancy and parallel transmissions to bound the end-to-end delay to traverse the network. Optionally, promiscuous listening between paths is possible, such that the nodes on one path may overhear transmissions along the other path. Considering the scenario shown in Figure 4, many different paths are possible for S to reach R. A simple way to benefit from this topology could be to use the two independent paths via nodes A, C, E and via B, D, F. But more complex paths are possible by interleaving transmissions from the lower level of the path to the upper level.

PRE may also take advantage of the shared properties of the wireless medium to compensate for the potential loss that is incurred with

radio transmissions. For instance, when the source sends to A, B may
listen also and get a second chance to receive the frame without an
additional transmission. Note that B would not have to listen if it
already received that particular frame at an earlier timeslot in a
dedicated transmission towards B.

                         (A)   (C)   (E)

        source (S)   (R) (root)

                         (B)   (D)   (F)
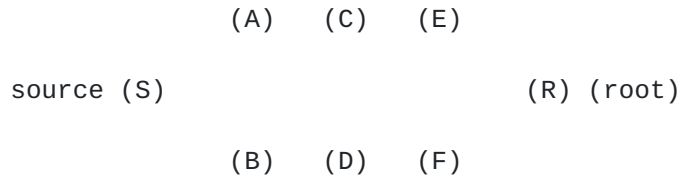

   Figure 4: A Typical Ladder Shape with Two Parallel Paths Toward the
                              Destination

   The PRE model can be implemented in both centralized and distributed
   scheduling approaches. In the centralized approach, a Path
   Computation Element (PCE) scheduler calculates the routes and
   schedules the communication among the nodes along a circuit such as
   a Label switched path. In the distributed approach, each node
   selects its route to the destination, typically using a source
   routing header. In both cases, at each node in the paths, a default
   parent and alternative parent(s) should be selected to set up
   complex tracks.

   In the following Subsections, all the required operations defined by
   PRE, namely, Alternative Path Selection, Packet Replication, Packet
   Elimination and Promiscuous Overhearing, are described.

5.3.1. Packet Replication

   The objective of PRE is to provide deterministic networking
   properties: high reliability and bounded latency. To achieve this
   goal, determinism in every hop of the forwarding paths MUST be
   guaranteed. By employing a Packet Replication procedure, each node
   forwards a copy of each data packet to multiple parents: its Default
   Parent (DP) and multiple Alternative Parents (APs). To do so, each
   node (i.e., source and intermediate node) transmits the data packet
   multiple times in unicast to each parent. For instance, in Figure 5,
   the source node S is transmitting the packet to both parents, nodes
   A and B, in two different timeslots within the same TSCH slotframe.
   An example TSCH schedule is shown in Figure 6. Thus, the packet
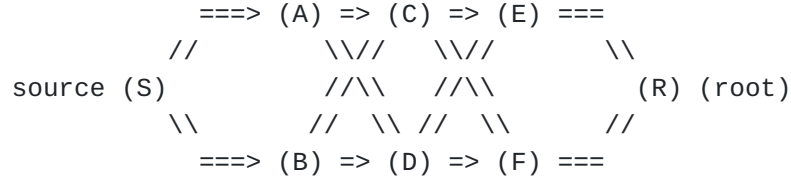   eventually obtains parallel paths to the destination.

```
                      ===> (A) => (C) => (E) ===
                //          \\//    \\//        \\
        source (S)          //\\    //\\         (R) (root)
                \\          //  \\ //  \\        //
                      ===> (B) => (D) => (F) ===
```

    Figure 5: Packet Replication: S transmits twice the same data packet,
                    to its DP (A) and to its AP (B).

```
                              Timeslot
    +---------++------+------+------+------+------+------+------+
    | Channel || 0    | 1    | 2    | 3    | 4    | 5    | 6    |
    +---------++======+======+======+======+======+======+======+
    | 0       || S->A | S->B | B->C | B->D | C->F | E->R | F->R |
    +---------++------+------+------+------+------+------+------+
    | 1       ||      | A->C | A->D | C->E | D->E | D->F |      |
    +---------++------+------+------+------+------+------+------+
```

            Figure 6: Packet Replication: Sample TSCH schedule

5.3.2. Packet Elimination

   The replication operation increases the traffic load in the network,
   due to packet duplications. Thus, a Packet Elimination operation
   SHOULD be applied at each RPL DODAG level to reduce the unnecessary
   traffic. To this aim, once a node receives the first copy of a data
   packet, it discards the subsequent copies. Because the first copy
   that reaches a node is the one that matters, it is the only copy
   that will be forwarded upward. Then, once a node performs the Packet
   Elimination operation, it will proceed with the Packet Replication
   operation to forward the packet toward the RPL DODAG Root.

5.3.3. Promiscuous Overhearing

   Considering that the wireless medium is broadcast by nature, any
   neighbor of a transmitter may overhear a transmission. By employing
   the Promiscuous Overhearing operation, a DP and some AP(s)
   eventually have more chances to receive the data packets. In Figure
   7, when node A is transmitting to its DP (node C), the AP (node D)
   and its sibling (node B) may decode this data packet as well. As a
   result, by employing corellated paths, a node may have multiple
   opportunities to receive a given data packet. This feature not only
   enhances the end-to-end reliability but also it reduces the end-to-
   end delay and increases energy efficiency.
```

```
           ===> (A) ====> (C) ====> (E) ====
         //       ^ | \\                      \\
    source (S)      | |    \\                     (R) (root)
         \\       | v       \\                  //
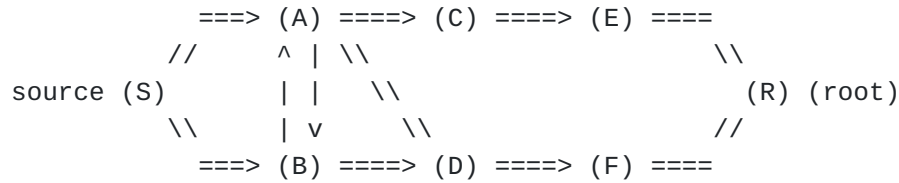           ===> (B) ====> (D) ====> (F) ====
```

Figure 7: Unicast to DP with Overhearing: by employing Promiscuous
Overhearing, DP, AP and the sibling nodes have more opportunities to
receive the same data packet.

6. Security Considerations

7. IANA Considerations

This document has no IANA actions.

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9. Acknowledgments

The authors wish to thank

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