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Extensions to IS-IS for Advertising Optional Router Capabilities

[draft-raggarwa-isis-cap-00.txt](#)

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2. Abstract

It is useful for routers in an IS-IS domain to know of the capabilities of their neighbors, and/or of other routers in the domain. This draft proposes extensions to IS-IS for advertising optional router capabilities. In particular it defines an optional Router Capability TLV for IS-IS.

[3.](#) Motivation

It is useful for routers in a domain to know of the capabilities of their IS-IS neighbors, and/or of other routers in the domain. This can be useful for various purposes:

- o In MPLS Traffic Engineering (TE) as a TE discovery mechanism [\[10\]](#) to announce a LSR's TE capabilities like Path Computation Server capability (Capability of a LSR to be a Path Computation Server for TE LSP path computation) or the intention of a LSR to be part of a particular MPLS TE mesh group.
- o For network management and troubleshooting. It gives operators a network wide view of IS-IS capabilities on different routers in the network. The presence of a capability on a given router implies that the software version supports the capability and the router is configured to support it. On the other hand the absence of an expected capability on a particular router can imply either mis-configuration or an incorrect software version. Hence this capability information can be used to track problems resulting from mis-configuration or an incorrect software version.

There is no existing mechanism in IS-IS to advertise optional router capabilities. We propose extensions to IS-IS for advertising these optional capabilities. For current IS-IS capabilities this advertisement will be used primarily for MPLS TE and informational purposes. Conceivably, future capability advertisements could be used for other purposes.

[4.](#) IS-IS Router Capability TLV

IS-IS [\[1\]](#) routers may optionally advertise their router capabilities in the TLV with code type 242. This TLV specifies the router ID of the router that originates the TLV, defines the flooding scope of the TLV, specifies the router capability bits in the first sub-TLV and certain capability related information in other sub-TLVs. This draft does not specify how an application may use the Router Capability TLV and such specification is outside the scope of this draft.

The router ID is a 32 bit unsigned integer to represent the router that originated this capability TLV. This is needed since this TLV can be flooded over the entire domain, hence the router ID of the originating router must be kept.

The capability bits are defined in a mandatory sub-TLV with code 1. It starts as a 32 bits flag, where each bit can represent a router capability. This flag can be expanded as needed to include more capabilities.

Some of the router capabilities may require more information than a single bit. The extra capability information can be encoded as sub-TLVs under this router capability TLV. The definition of these sub-TLVs is outside the scope of this draft.

If a router does not advertise this TLV, it does not imply that the router does not support one or more of the defined capabilities. If this TLV is included in the LSP, the router SHOULD set all the defined bits corresponding to the capabilities which the software supports, unless they are explicitly configured off.

4.1 Flooding Scope of the Router Capability TLV

There are three bits currently defined for this TLV in the information flag to control the flooding scope of the TLV. The Flooding bit, the Transit bit and the Down bit.

There are two flooding types defined for this router capability TLV's flooding scope. One is the domain wide flooding scope and the other is the intra-area flooding scope. The F bit if set indicates this TLV has the domain wide flooding scope.

The Transit bit can be used to signal the routers on the edge of the IGP routing domain to redistribute this TLV information into another routing process. How this is done is an application specific issue and is outside the scope of this document.

The L1/L2 routers MUST observe the Down bit to avoid TLV leak looping. This Down bit is not set when the router first originates this TLV and it MUST be set when leaking into a lower level or into another area of the same level. When the Down bit is set, this TLV can no longer be leaked to a higher level or into another area of the same level. This capability TLV MUST be preserved at the level boundary during TLV leaking. The L1/L2 router SHOULD NOT leak the TLV back into the same area which originated this TLV. It MAY be able to alter certain capability contents

during TLV leaking when specified by applications.

4.2 Encoding of the Router Capability TLV

The following figure depicts the structure of this IS-IS Router Capability TLV.

- ```
x CODE - 242
x LENGTH - total length of the value field in this TLV
x VALUE - 4-octet information flag, 4-octet router ID,
 1-octet sub-tlv length, the mandatory sub-TLV code 1
 for capability flags, and optional sub-TLVs for extra
 capability information, structured as follows:
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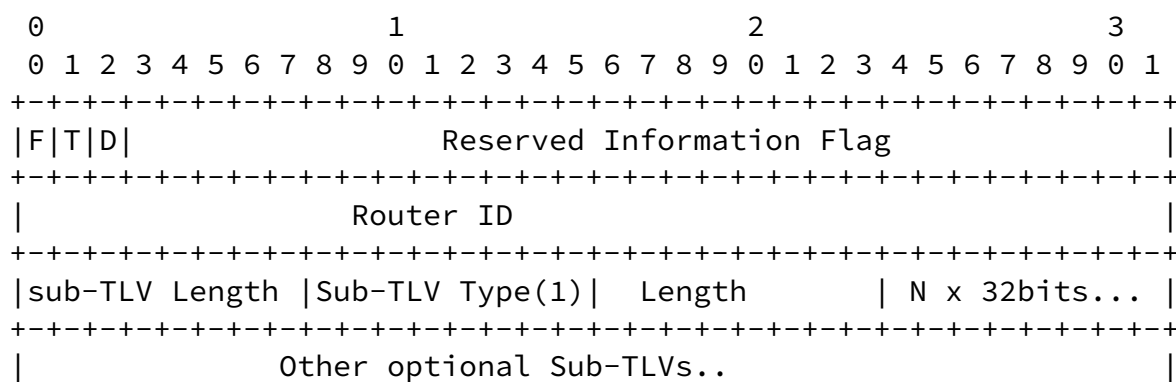


Figure 1. IS-IS Router Capability TLV

The first field is the 4-octet information flag, which consists of the F, T and D bits, the reserved information bits.

Bit F represents the Flooding scope of the TLV. If set, this TLV SHOULD be flooded to entire IGP domain. Otherwise, it SHOULD NOT be leaked into the other level or another area in the same level.

Bit T determines the Transit behavior into other routing domains. For example, if this bit is set, a router can leak this capability information into another routing protocol.

Bit D represents Down/Up behavior during the TLV leaking. When the capability is leaked from level 2 into level 1 or it is leaked into another area of the same level, this D bit **MUST** be set. Otherwise this bit **MUST** be cleared.

Router ID is an unsigned 32 bit number representing the router that originates this router capability TLV.

The next octet of the TLV is the total sub-TLV length of this router capability TLV. This sub-TLV length includes the first mandatory sub-TLV. The minimum value of this field is 6.

The first sub-TLV with code 1 is a mandatory sub-TLV, the router capability flag sub-TLV. The length is the length of this sub-TLV. Its set to  $N \times 4$  octets.  $N$  starts from 1 and can be increased when there is a need. Each 4 octets are referred to as a capability flag. For each capability flag the bits are indexed from the most significant to the least significant, where each bit represents one router capability.

There can be other sub-TLVs after the first sub-TLV to include extra information describing certain router capabilities. The description of those sub-TLVs is outside the scope of this draft.

The above data structure can be replicated within this TLV, but can not exceed the maximum length of 255 octets. If no other sub-TLVs are used and the capability flag is the minimum 4 octets, this encoding can contain up to 17 router capability TLVs where each have a minimum of 15 octets of data(4 byte information flag, 4 byte router-id, 1 byte total sub-tlv length, 6 byte capability flag).

#### [4.3](#) Reserved IS-IS Router Capability Bits

We have assigned some pre-determined bits to the first capability flag.

| Bit | Capabilities                                                  |
|-----|---------------------------------------------------------------|
| 0-3 | Reserved                                                      |
| 4   | IS-IS graceful restart capable <a href="#">[4]</a>            |
| 5   | IS-IS and BGP blackhole avoidance capable <a href="#">[6]</a> |
| 6   | IS-IS wide metric processing capable <a href="#">[3]</a>      |
| 7   | IS-IS short metric processing capable <a href="#">[1]</a>     |
| 8   | IS-IS hmac-md5 authentication capable <a href="#">[5]</a>     |
| 9   | IS-IS Traffic Engineering support <a href="#">[3]</a>         |

- 10 IS-IS point-to-point over LAN [7]
- 11 IS-IS Path Computation Server discovery [10]
- 12 M-ISIS capable [8]
- 13 IS-IS IPv6 capable [9]
- 14-31 For future assignments

## [6.](#) Security Consideration

This document does not introduce new security issues. The security considerations pertaining to the original IS-IS protocol remain relevant.

## [7.](#) Acknowledgments

The idea for this work grew out of a conversation with Andrew Partan and we would like to thank him for his contribution.

## [8.](#) References

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- [10] Vasseur et al, "RSVP Path computation request and reply " messages", [draft-vasseur-mpls-computation-rsvp-te-03.txt](#), work in progress

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