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**Device Enrollment in IETF protocols -- A Roadmap
draft-richardson-enrollment-roadmap-03**

Abstract

This document provides an overview of enrollment or imprinting mechanisms in current IETF protocols.

This document is being worked on in the ANIMA-WG, and on github at:
<https://github.com/anima-wg/enrollment-roadmap>

Status of This Memo

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1. Introduction

There are numerous mechanisms being proposed to solve the problem of securely introducing a new devices into an existing managed network.

This document provides an overview of the different mechanisms showing what technologies are common.

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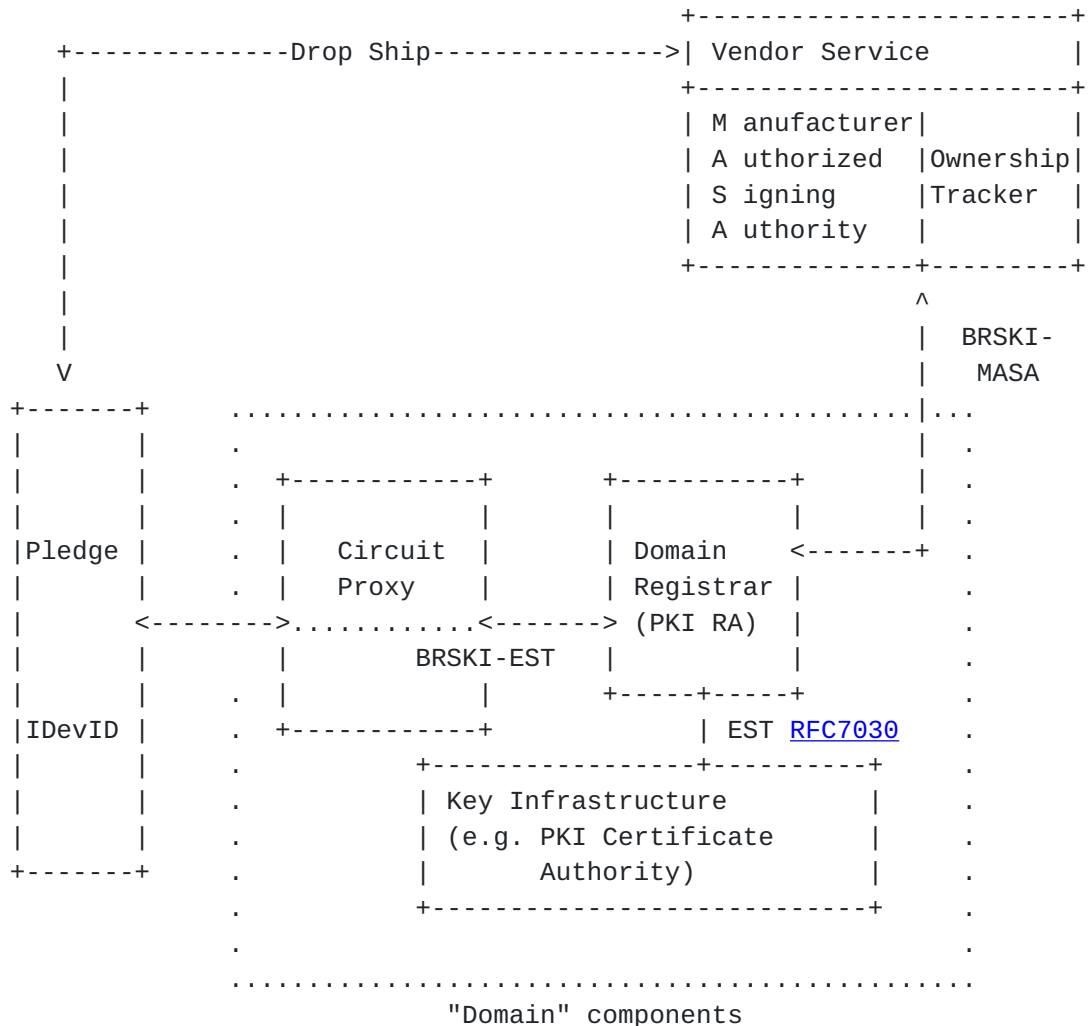
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The document starts with a diagram showing the various components and how they go together to form five enrollment scenarios.

2. Components of enrollment solutions

This diagram is taken from [\[I-D.ietf-anima-bootstrapping-keyinfra\]](#), which is where this work started.

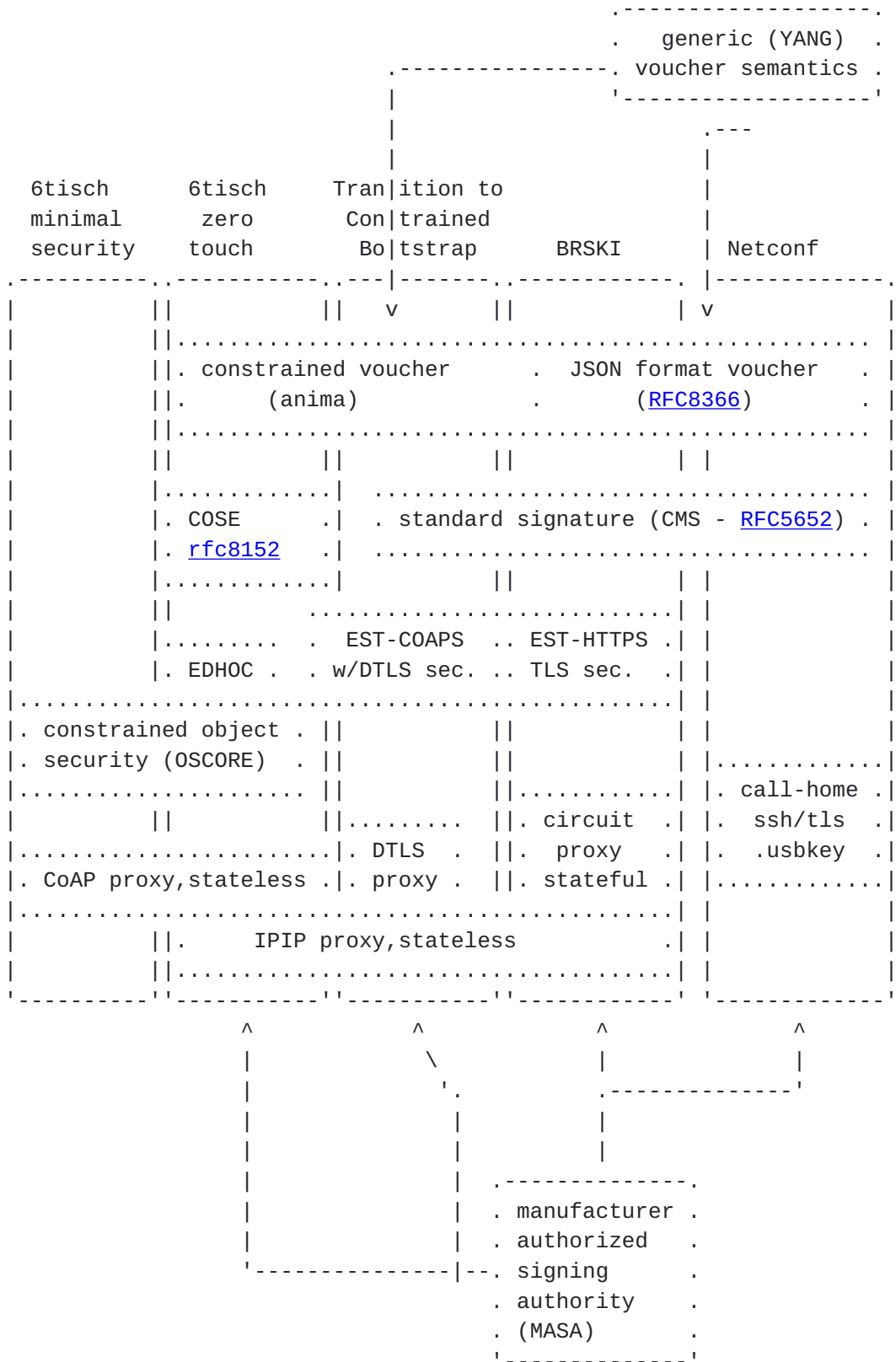


Five major components are described:

1. pledge: The node that is attempting to enroll.
2. proxy: A node that is within one layer-2 hop of the pledge that is helping.
3. domain registrar: the Join Registrar/Coordinator (JRC) that will determine eligibility of the pledge.

4. MASA: the representative of the manufacturer that has a pre-established trust relationship with the pledge.
5. the domain PKI (if any)

3. Map of Enrollment solution



4. Components

4.1. generic voucher semantics

The abstract semantics of the voucher, described in YANG, are in [[RFC8366](#)].

4.2. constrained voucher

The semantics of the constrained voucher, represented in CBOR, are described in [[I-D.ietf-anima-constrained-voucher](#)].

4.3. JSON format voucher

The semantics of the basic voucher, represented in JSON, are described in [[RFC8366](#)].

4.4. COSE-8152

In constrained systems the voucher is signed using the COSE mechanism described in [[RFC8152](#)]. This is to be replaced with [[I-D.ietf-cose-rfc8152bis-algs](#)] (now in RFC-EDITOR Q) and [[I-D.ietf-cose-rfc8152bis-struct](#)] (in WGLC, being revised to remove countersignatures).

4.5. standard signature (CMS)

In un-constrained systems the voucher is signed using the Cryptographic Message Syntax (CMS) described in [[RFC5652](#)].

4.6. EDHOC

On constrained and challenged networks, the session key management can be formed by [[I-D.ietf-lake-edhoc](#)].

This document has a WG called LAKE.

The CoAP-EST layer on top is described by [[I-D.ietf-ace-coap-est](#)]

4.7. EST-COAPS 2/DTLS sec(urity)

On unconstrained networks, the session key management is provided by [[RFC6347](#)]. The CoAP-EST layer on top is described by [[I-D.ietf-ace-coap-est](#)], which is in the RFC-EDITOR Q, waiting for DTLS 1.3.

The ACE WG adopted this document, and published it.

4.8. EST-HTTPS TLS sec(urity)

On unconstrained networks with unconstrained nodes, the EST layer and session key management is described by [\[RFC7030\]](#) as modified by [\[I-D.ietf-anima-bootstrapping-keyinfra\]](#) (BRSKI).

BRSKI is in the RFC-EDITOR Q, waiting for [\[I-D.ietf-anima-autonomic-control-plane\]](#), which is awaiting IESG go-ahead.

4.9. constrained object security (OSCORE)

On constained networks with constrained nodes, the CoAP transactions are secured by [\[I-D.ietf-core-object-security\]](#) using symmetric keys. The symmetric key may be pre-shared (for 6tisch minimal security), or MAY be derived using EDHOC.

4.10. ultra-constrained enrollment

The document [\[I-D.selander-ace-ake-authz\]](#). Version -01 expired in September 2020, but is expected to be reposted for discussion in the ACE WG.

4.11. Pledge traffic proxy mechanisms

Traffic between the Pledge and the JRC does not flow directly as the pledge does not typically have a globally reachable address, nor does it have any network access keys (whether WEP, WPA, 802.1x, or 802.15.4 keys).

Communication between the pledge and JRC is mediated by a proxy. This is primarily to protect the network against attacks. The proxy mechanism is provided by as many nodes as can afford to as a benefit to the network, and therefore MUST be as light weight as possible.

There are therefore stateless mechanisms and stateful mechanisms. The costs of the various methods is analyzed in [\[I-D.richardson-anima-state-for-joinrouter\]](#).

4.11.1. COAP proxy, stateless

The CoAP proxy mechanism uses the OSCORE Context Hint to statelessly store the address of the proxy within the CoAP structure. It is described in [\[I-D.ietf-6tisch-minimal-security\]](#).

4.12. DTLS proxy

A DTLS specific mechanism for COAPS, is described in [[I-D.vanderstok-anima-constrained-join-proxy](#)].

4.13. IPIP proxy, stateless

An IPIP proxy mechanism uses a layer of IP-in-IP header (protocol 98) to encapsulate the traffic between Join Proxy and JRC. It has some complexities to implement on typical POSIX platforms.

At one point, it was described in [[I-D.ietf-6tisch-dtsecurity-zerotouch-join](#)], in an Appendix. Another home for the text is desired.

4.14. circuit proxy stateful

The circuit proxy method utilizes either an application layer gateway (which in canonical 1990-era implementation requires a process per connection), or the use of NAT66. It maintains some state for each connection whether TCP or UDP.

It is this most expensive and most easily abused, but also the most widely available, code-wise.

5. call-home ssh/tls/usbkey

The NETCONF call-home mechanism assumes that the device can get basic connectivity, enough for an out "outgoing" TCP connection to the manufacturer.

6. manufacturer authorized signing authority (MASA)

The MASA is the manufacturers anchor of the manufacturer/pledge trust relationship that is established at the factory where the pledge is built.

7. Enrollment Mechanisms

7.1. NETCONF

The NETCONF WG is describing this in [[I-D.ietf-netconf-zerotouch](#)] document.

The NETCONF Zerotouch mechanism provides configuration and ownership information by having the pledge "call home" to a location determined by a mix of local hints (DHCPv4, DHCPv6, and mDNS), as well as built-

in anchors. Additionally, ownership vouchers can be alternatively distributed by portable storage such as USB key.

Upon reaching a validated call-home server, Zerotouch typically "reverses" the connection providing either an SSH or TLS connection to the pledge device such that it can be configured automatically.

Zerotouch relies upon either open or very easy access to network connectivity, along with the ability to make an outgoing TCP connection to the Internet, or to the provided local configuration agent.

Zerotouch is seen as an updated version of TR-69 by some, appropriate for configuration of residential appliances which are drop-shipped by ISPs or other service providers to homes. That is not the only targetted use.

7.2. BRSKI

The ANIMA WG is describing BRSKI in [\[I-D.ietf-anima-bootstrapping-keyinfra\]](#) document.

The ANIMA WG does enrollment with the aim of creating a secure channel with a public-key infrastructure (PKI) Registrar. The secured channel is used to perform Enrollment over Secure Transport (EST, [RFC7030](#)). The real goal is the enrollment a new device which was probably been drop-shipped into ANIMA's Autonomic Control Plane.

That is, after the pledge has been assigned a certificate within the (autonomic) domain, the device (no longer a pledge) will then form secure channels (typically using IKEv2 to key an IPsec channel). On top of that channel a routing protocol (RPL) is run to form the Autonomic Control Plane (ACP). The ACP is then used as a management network with which to configure the new device.

BRSKI is therefore step one of a number of steps, the ultimate goal of which is to bring the pledge into the ACP as a new device.

BRSKI itself does not provide for any direct keying of the network (802.11 WEP/WPA, or 802.15.4 security). The provision of a domain certificate at each node can, however, be used to do that kind of keying: for instance 802.15.9 provides for use of HIP and IKEv2 to key 802.15.4 networks.

7.3. Transition to Constrained Bootstrap

It is distinguished from BRSKI in that it uses DTLS (rather than TLS) and constrained (CBOR) vouchers. It is distinguished from 6tisch Zero Touch in that uses CMS to sign rather than COSE.

The ACE WG has adopted [[I-D.ietf-ace-coap-est](#)], but this is not a sufficient.

The constrained version of BRSKI is current described as part of [[I-D.ietf-anima-constrained-voucher](#)].

The use of this technology slice is attractive to IoT deployments where the devices are not battery powered (lighting for instance, AMI for electric meters). In such situations, the processors in each device have significantly more resources, and in particular far more code space available. The use of DTLS to secure application traffic (as described in the ACE documents) is already common, and so reuse of DTLS is desirable from a code point of view.

However, the network capacity is still limited so TCP and CBOR are still important. The network may also contain extremely constrained devices (kinetically powered light switches for n instance).

7.4. Cloud-based Registrars

BRSKI depends upon a local join proxy to direct the pledge device to the location of the Registrar. Not only does this require that a join proxy be widely available, but it also assumes that the owner is able to operate a Registrar.

[[I-D.friel-anima-brski-cloud](#)] deals with two use cases:

1. where there is a Registrar, but no join proxy. A redirection through a manufacturer or VAR operated register points the device to a Registrar located on the Internet.
2. where there is no Registrar (yet), but the network operator does have an [RFC7030](#) EST server. The cloud registrar provides a voucher connecting the pledge to the appropriate EST server.

7.5. BRSKY Asynchronous Enrollment

BRSKI depends the join proxy being available in the location where the devices are installed, and for that proxy to also be able to communicate to the Registrar (possibly over the Internet).

In a number of environments, such as new construction, a number of devices need to go through an onboarding process, but there is no connectivity available. (Possibly not even LTE/3G due to metal walls, lack of towers, etc.).

The [[I-D.ietf-anima-brski-async-enroll](#)] proposes two amendments to BRSKI:

1. the use of CMP rather than EST for enrollment, which eliminates the need to keep a live TLS connection between pledge and Registrar.
2. a discovery, store-and-forward mechanism for moving voucher requests and vouchers between the Registrar and Pledge, including a batch mechanism that aims to reduce the amount of walking that an installer might have to do.

7.6. 6tisch Zero Touch

The 6tisch WG is describing this in [[I-D.ietf-6tisch-dtsecurity-zerotouch-join](#)] document.

The 6tisch use case consists of very constrained devices with very constrained networks. Code space in the devices is larger than typical class 2, but the devices are typically battery powered and wish to sleep significantly.

The use of CBOR for vouchers, COSE to sign the vouchers saves significant network bandwidth and code space. Both CBOR, COSE and OSCORE are typically already in use for the application support. The addition of EDHOC to provide asymmetric bootstrap of OSCORE completes the suite of constrained security protocols.

7.7. 6tisch minimal security

The 6tisch WG is describing this in [[I-D.ietf-6tisch-minimal-security](#)] document. This mechanism does enrollment in a single request/response message, but requires at least one "touch" to pre-share symmetric keys.

The 6tisch WG felt that the number of round trips required to do EDHOC, and the size of the vouchers required an even simpler protocol. As existing 6tisch-type technology is typically deployed with network keys built-in at manufacturer time (no "drop-ship"), the switch from a static network key to a PSK for authentication is considered an incremental improvement.

All other methods are considered zero "touch".

[I-D.ietf-6tisch-minimal-security] is in the RFC-EDITOR Q, waiting for core-stateless on which it depends to leave the WG.

8. Discussion

A goal of this document is to provide some guidance in selecting which enrollment profile to use for a given scenario. This section tries to provide some contrasting comments between the various mechanisms.

(BUT, it does not yet do that..)

9. OPEN ISSUE

1. insufficient guidance if there is no internet connection.
2. classes of devices where the manufacturer will never have an HSM/PKI. so, it will only be a self-signed key.
3. how does a supply-chain manage BRSKI?

10. Security Considerations

This document includes a tradeoff of the security attributes of the different protocols, and so the entire document contains security advice.

11. IANA Considerations

This document does not define any new protocols, and therefore does not have any IANA Considerations.

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TBD

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