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Third Party Call Control in SIP

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Abstract

This document discusses the usage of SIP for third party call control. Third party call control refers to the ability of one entity to create a call in which communications is actually between other parties. We present a SIP mechanism for accomplishing third party call control that does not require any extensions or changes to SIP.

1 Introduction

In the traditional telephony context, third party call control refers to the ability of one entity (which we call the controller), to create, modify, or terminate calls between other participants. Third party call control is often used for operator services (where an operator creates a call that connects two participants together), and conferencing. On the Internet, a wider range of services are enabled through a third party session control mechanism. This is because other IP applications, such as web, email, presence, instant messaging, and chat can now be brought into the picture. An excellent example is click-to-dial. This service allows a user to click on a web page when they wish to speak to a customer service representative. The web server then creates a call between the user and a customer service representative. The call can be between two phones, a phone and an IP host, or two IP hosts.

In order to support third party call control applications, a mechanism is needed that allows a controller to create, modify, and terminate calls with other entities. In this document, we present a mechanism using SIP which allows a controller to execute third party services. The mechanism is not an extension to SIP. It is merely an application of the tools enabled through <u>RFC 2543</u>. A controller can create calls between any entity that contains a normal SIP user agent. After desribing the mechanism, we present three third party services which take advtantage of this mechanism. One is click-to-dial, the second is a feature that enables a mid-call announcement, and the third is a timed conference bridge initiation.

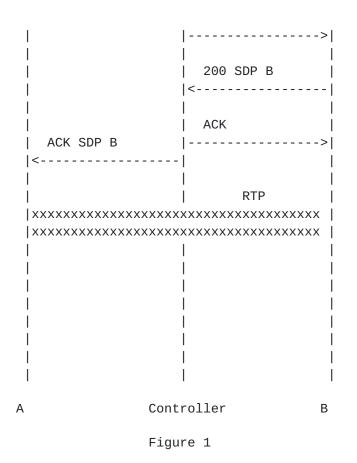
2 Third Party Control

The basic idea behind the third party mechanism is simple. Consider first the case of just connecting two users in a call. The controller first sends an INVITE to the first user whose phone is to ring. This is a standard INVITE, but it contains no SDP.[1] send an ACK. It generates a second INVITE. This INVITE is addressed to the second user to be connected in the call. This INVITE contains the SDP as received from the 200 OK of the first user. When the 200 OK to this second INVITE arrives, the controller ACK s it, takes the SDP, and includes that in the ACK for the first call. A flow diagram for this mechanism is given in Figure 1.

```
| INV no SDP |
|<-----|
| 200 SDP A |
|-----> | INV SDP A
```

^{[&}lt;u>1</u>] <u>RFC 2543</u> does allow for the initial INVITE to not contain a session description

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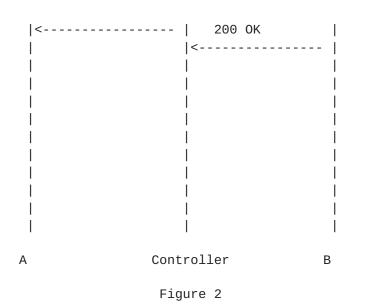


At this point, both participants believe they are in a single point to point call with some control system (assuming the controller identified itself as such in the From field of the INVITE). However, they are exchaning media directly with each other, rather than with the controller. The result is that the controller has set up a call between both participants.

Since the controller is still a central point for signaling, it now has complete control over the call. If it receives a BYE from one of the participants, it can create a new BYE and hang up with the other participant. This is shown in Figure 2.

1	
BYE From A	
	-> BYE From Cont.
200 OK	>

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As an alternative, when the controller receives a BYE from A, it can generate a new INVITE to a third party, C, using the SDP from B. When the 200 OK arrives from C, the controller sends a re-INVITE to B, using the SDP from C. If the 200 OK to the re-INVITE contains the same SDP as it used in the INVITE to C, the controller has successfully connected B to C, transparently to B. A call flow for this is shown in Figure 3.

 BYE From A > 200 OK <	 INV SDP B 	 200 SDP C
	АСК	
	INV SDP C	
	200 SDP B	
	< ACK	
	>	
		RTP

[Page 4]

İ	i	İ	i
A	Controller	В	С

From here, new parties can be added, removed, transferred, and so on, as the controller sees fit.

The general idea behind the mechanism is that there is a point to point SIP relationship between each participant and the controller. However, by passing the SDP it receives from one participant to another, it can causes users to actually communicate with each other rather than the controller.

3 Click to Dial

|

The first application of this capability we discuss is click to dial. In this service, a user is browsing the web page of an e-commerce site, and would like to speak to a customer service representative. They click on a link, and the phone on the desk (a normal telephone) rings. When the user picks up, the phone of the customer service representative (an IP phone) rings. When they pick up, the service representative is talking to the user.

We assume for purposes of this discussion that the web server is actually an applications server that contains an http interface. In this case, when the user clicks on the URL, the application server knows, through cookies or some other state mechanism, the addresses of the participants to be connected.

The call flow for this service is given in Figure 4. Note that it is identical to that of Figure 1, with the exception that the service is triggered through an http GET request when the user clicks on the link.

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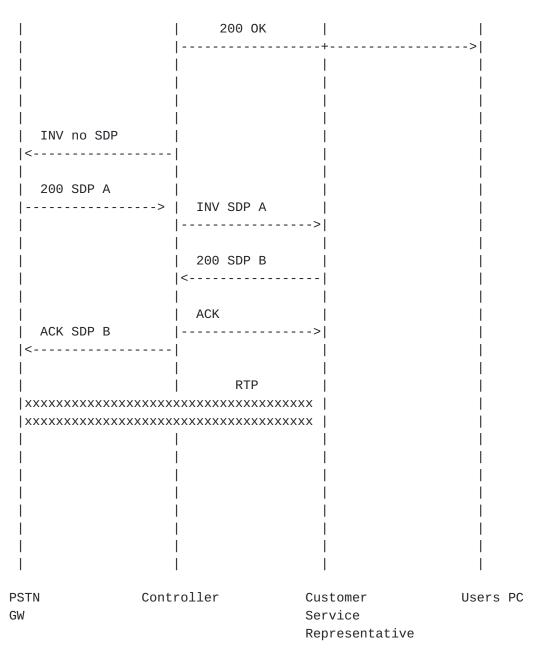


Figure 4

We note that this service can be provided through other mechanisms, namely PINT [1]. However, there are numerous differences between the way in which the service is provided by pint, and the way in which it is provided here:

o The pint solution enables calls only between two PSTN endpoints. The solution described here allows calls between PSTN phones (through SIP enabled gateways) and native IP

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Зрсс

phones.

- o When used for calls between two PSTN phones, the solution here may result in a portion of the call being routed over the Internet. In pint, the call is always routed only over the PSTN. This will probably result in better quality calls with the pint solution.
- o The PINT solution requires extensions to SIP (PINT is an extension to SIP), whereas the solution described here is done with baseline SIP.
- o The PINT solution allows the controller (acting as a PINT client) to "step out" once the call is established. The solution described here requires the controller to remain in the picture for the entire duration of the call.

<u>4</u> Mid-Call Announcement Capability

The third party call control mechanism described here can also be used to enable mid-call announcements. The call is set up by the controller as desribed above.[2] payphone, in which case the controller determines that the call is to be terminated after some amount of time if the user doesn't add more money to the phone. When this timer expires, the controller initiates places the called party on hold. It then sends an INVITE to the media server which will be collecting digits. It then sends a re-INVITE to the user on the payphone, connecting its media streams with the media server. The media server plays an announcement, and prompts the user to enter a credit card number, for example. After collecting the number and validating the card, if the call can continue, the media server hangs up. The controller takes this as a cue and reconnects the user to the original called party, and takes the original called party off hold.

A call flow for this service is shown in Figure 5

| RTP |

[2] It is actually not necessary for the controller to set up the call. However, if a participant initiates the call, the controller must step in as a virtual UAC/UAS, and act as a termination and re-initiation point

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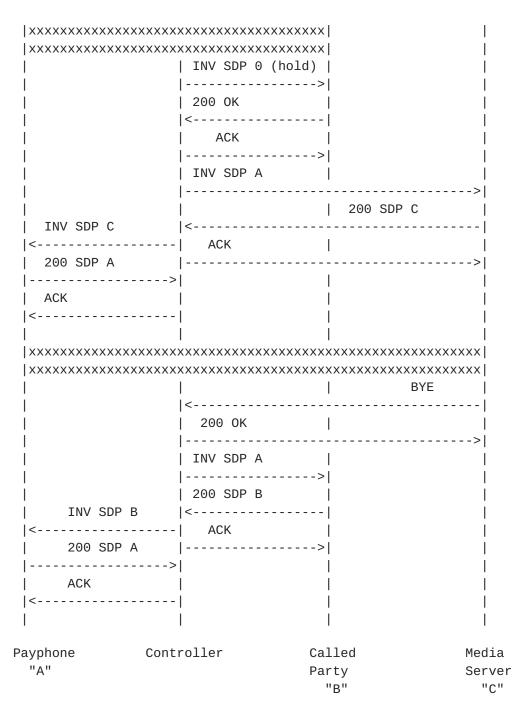


Figure 5

We have assumed that the media server and the controller have agreed, ahead of time, that a hangup implies that the desired service (extending the lifetime of the call) has succeeded. This is effectively allowing a call control interface between the controller

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and the media server. Parameters needed between the elements, such as the new expiration of the call, can be passed in the BYE. A separate draft, forthcoming, will discuss call control interfaces to media services in more detail.

<u>5</u> Timed Conference Intitation

In this service, a conference bridge is booked for some number of participants. In order to make sure the conference begins on time, the conference bridge will call each participant at the time of the call. If a participant doesn't answer, the bridge tries to contact them again (unless they call in) five minutes later.

In the call flow described here, we assume that the controller acts as the media bridge. This is not strictly necessary; some kind of control interface could be used to separate the media function from the controller.

The call flow, shown in Figure 6, is, not surprisingly, remarkably like that of Figure 1. The only difference is that the SDP listed in the INVITE s generated by the controller always contain SDP that points to the conference bridge, rather than one of the other participants. In the call flow diagram, user 1 is invited first, then user 2, and then user 3. User 3 is not available, but is called again five minutes later.

INV SDP X	1	I I
<	 INV SDP X	
200 SDP A >	> 	
İ	200 SDP B	i i
ACK	<	
<		
	ACK >	
	 INV SDP X	
	 	408 Timeout
	< ACK	
		>

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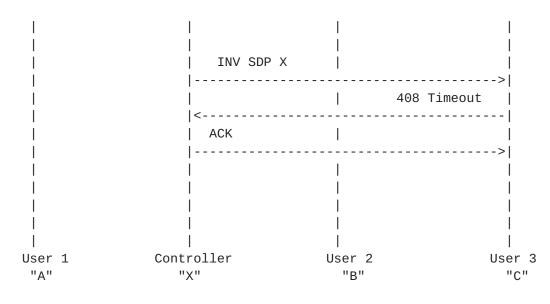


Figure 6

<u>6</u> Future Work

We plan on considering other mechanisms which might be used for third party call control, and discuss the pros and cons for each for providing numerous services.

7 Conclusions

We have presented a basic third party call control mechanism that uses SIP. This mechanism does not require any extensions to SIP and is completely backwards compatible.

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9 Bibliography

[1] S. Petrack and L. Conroy, "The PINT service protocol: Extensions to SIP and SDP for IP access to telephone call services," Internet Draft, Internet Engineering Task Force, June 1999. Work in progress.

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