

INTERNET-DRAFT

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February 2000

<[draft-skwan-gss-tsig-05.txt](#)>

Expires August 2000

## **GSS Algorithm for TSIG (GSS-TSIG)**

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Abstract

The TSIG protocol provides transaction level authentication for DNS. TSIG is extensible through the definition of new algorithms. This document specifies an algorithm based on the Generic Security Service Application Program Interface (GSS-API) [[RFC2078](#)].

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## [1. Introduction](#)

The Secret Key Transaction Signature for DNS [[TSIG](#)] protocol was developed to provide a lightweight alternative to full DNS security [[RFC2535](#)] and secure dynamic update [[RFC2137](#)], where full security is impractical due to implementation complexity, management overhead, or computational cost.

The [[TSIG](#)] protocol is extensible through the definition of new algorithms. This document specifies an algorithm based on the Generic Security Service Application Program Interface (GSS-API) [[RFC2078](#)]. GSS-API is a framework that provides an abstraction of security to the application protocol developer. The security services offered can include authentication, integrity, and confidentiality.

The GSS-API framework has several benefits:

- \* Mechanism and protocol independence. The underlying mechanisms that realize the security services can be negotiated on the fly and varied over time. For example, a client and server may use Kerberos for one transaction, whereas that same server may use TLS with a different client.

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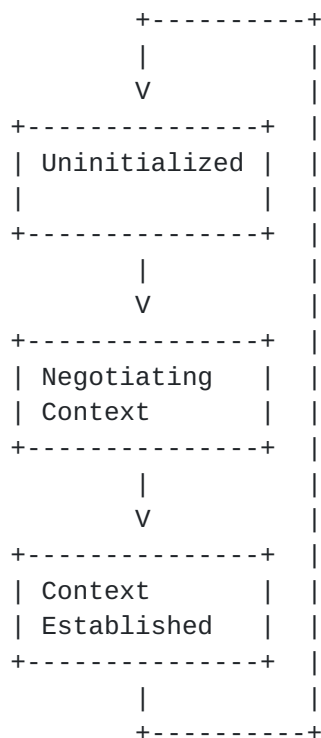
\* The protocol developer is removed from the responsibility of creating and managing a security infrastructure. For example, the developer does not need to create new key distribution or key management systems. Instead the developer relies on the security service mechanism to manage this on its behalf.

## 2. Protocol Overview

Readers that are unfamiliar with the GSS-API concepts are encouraged to read the characteristics and concepts section of [[RFC2078](#)] before examining this protocol in detail.

In GSS, client and server interact to create a "security context". The security context is used to create and verify transaction signatures on messages between the two parties. A unique security context is required for each unique connection between client and server.

Creating a security context involves a negotiation between client and server. Once a context has been established, it has a finite lifetime for which it can be used to secure messages. Thus there are three states of a context associated with a connection:



Every connection begins in the uninitialized state.

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## **2.1 GSS Details**

Client and server must be locally authenticated and have acquired default credentials per [[RFC2078](#)] before using this protocol.

Not all flags used in GSS-API interfaces are specified in this document. Where omitted, clients and servers may select the default or use a value based on local system policy.

Not all error return values from GSS-API interfaces are specified in this document. When errors are encountered, the caller should act appropriately.

## **3. Client Protocol Details**

A unique context is required for each server to which the client sends secure messages. A context is identified by a context handle. A client maintains a mapping of handles to servers,

(target\_server\_name, key\_name, context\_handle)

The value key\_name also identifies a context handle, and is used on the wire to indicate to a server which context should be used to process the current request.

### **3.1 Negotiating Context**

In GSS, establishing a security context involves the passing of opaque tokens between the client and the server. The client generates the initial token and sends it to the server. The server processes the token and if necessary, returns a subsequent token to the client. The client processes this token, and so on, until the negotiation is complete. The number of times the client and server exchange tokens depends on the underlying security mechanism. A completed negotiation results in a context handle.

The TKEY resource record [[TKEY](#)] is used as the vehicle to transfer tokens between client and server. The TKEY record is a general mechanism for establishing secret keys for use with TSIG. For more information, see [[TKEY](#)].

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### **3.1.1 Call GSS\_Init\_sec\_context**

The client obtains the first token by calling GSS\_Init\_sec\_context. The following input parameters are used. The outcome of the call is indicated with the output values below. Consult [[RFC2078](#)] for syntax definitions.

#### INPUTS

CONTEXT HANDLE	input_context_handle	= 0
INTERNAL NAME	targ_name	= DNS/<target_server_name>
OCTET STRING	input_token	= NULL
BOOLEAN	replay_det_req_flag	= TRUE
BOOLEAN	mutual_req_flag	= TRUE
BOOLEAN	deleg_req_flag	= TRUE (optional)

#### OUTPUTS

INTEGER	major_status
CONTEXT HANDLE	output_context_handle
OCTET STRING	output_token
BOOLEAN	replay_det_state
BOOLEAN	mutual_state

The values of replay\_det\_state and mutual\_state indicate if the security package can provide replay detection and mutual authentication, respectively. If one or both of these values are FALSE, the client must abandon this protocol.

The deleg\_req\_flag is optional, and can be used if the client wants the server to be able to call out to other services under the context of the client.

If major\_status indicates an error, the client must abandon the protocol. Success values of major\_status are GSS\_S\_CONTINUE\_NEEDED and GSS\_S\_COMPLETE. The exact success code is important during later processing.

The handle output\_context\_handle is unique to this negotiation and is stored in the client's mapping table as the context\_handle that maps to target\_server\_name.



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### **3.1.2 Send TKEY Query to Server**

The output\_token from GSS\_Init\_sec\_context is transmitted to the server in a query request for a TKEY record. The token itself will be placed in a TKEY record in the additional records section of the query. The domain-like name of the TKEY record set queried for and the name of the supplied TKEY record in the additional section will uniquely identify the security context to both the client and server, and thus the client should use a value which is globally unique.

#### **TKEY Record**

NAME = client-generated globally unique domain name string  
(see [[TKEY](#)])

#### **RDATA**

Algorithm Name	= gss-tsig.microsoft.com
Mode	= 3 (GSS-API negotiation - see [ <a href="#">TKEY</a> ])
Key Size	= size of output_token
Key	= output_token

Assign the remaining fields in the TKEY RDATA appropriate values per [[TKEY](#)].

If the last call to GSS\_Init\_sec\_context yielded a major\_status value of GSS\_S\_COMPLETE, then the message should be signed with a TSIG record before being sent to the server. See [section 5](#), Sending and Verifying Signed Messages, for the signing procedure.

The query is transmitted to the server.

### **3.1.3 Receive TKEY Query-Response from Server**

The server will return a standard TKEY query-response (see [[TKEY](#)]). The response may indicate that TKEY is not supported, or that the GSS-API mode and algorithm are not supported. If this is the case, the client must abandon this algorithm.

If the value of the Error field in the TKEY RDATA (TKEY.Error) is BADKEY, then the token provided by the client was invalid. The client must abandon this algorithm.

If TKEY.Error indicates success, then the client may continue. The next processing step depends on the value of major\_status from the most recent call to GSS\_Init\_sec\_context: either GSS\_S\_COMPLETE or GSS\_S\_CONTINUE.

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#### **3.1.3.1 Value of major\_status == GSS\_S\_COMPLETE**

If the last call to GSS\_Init\_sec\_context yielded a major\_status value of GSS\_S\_COMPLETE and a non-NULL output\_token was sent to the server, then the client side component of the negotiation is complete and the client is awaiting confirmation from the server.

Confirmation will be in the form of a NOERROR query response containing the last client supplied TKEY record in the answer section of the query. The response may also be signed with a TSIG record, and if present this signature must be verified using the procedure detailed in [section 5](#), Sending and Verifying Signed Messages.

If the message verification completes without an error, or if a TSIG signature was not included, the context state is advanced to Context Established. Proceed to [section 3.2](#) for usage of the security context.

#### **3.1.3.2 Value of major\_status == GSS\_S\_CONTINUE**

If the last call to GSS\_Init\_sec\_context yielded a major\_status value of GSS\_S\_CONTINUE, then the negotiation is not yet complete. The server will respond to the TKEY query with a NOERROR query response that contains a TKEY record in the answer section. The TKEY record contains a token that is passed to GSS\_Init\_sec\_context using the parameters from the previous call and the following modifications:

##### **INPUTS**

CONTEXT HANDLE	input_context_handle	= context_handle
OCTET STRING	input_token	= token from Key field of TKEY record

##### **OUTPUTS**

INTEGER	major_status
OCTET STRING	output_token

If major\_status indicates an error, the client must abandon this algorithm. Success values are GSS\_S\_CONTINUE and GSS\_S\_COMPLETE.

If major\_status is GSS\_S\_CONTINUE the negotiation is not yet finished. The token output\_token must again be passed to the server in a TKEY record. The negotiation sequence is repeated beginning with [section 3.1.2](#). The client should place a limit on the number of continuations in a context negotiation to prevent endless looping.

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If `major_status` is `GSS_S_COMPLETE` and `output_token` is non-NULL, the client-side component of the negotiation is complete but the token `output_token` must be passed to the server. The negotiation sequence is repeated beginning with [section 3.1.2](#).

If `major_status` is `GSS_S_COMPLETE` and `output_token` is NULL, context negotiation is complete. The response from the server may be signed with a TSIG record, and if present this signature must be verified using the procedure detailed in [section 5](#), Sending and Verifying Signed Messages.

If the message verification completes without an error, or if a TSIG signature was not included, the context state is advanced to Context Established. Proceed to [section 3.2](#) for usage of the security context.

### [3.2](#) Context Established

When context negotiation is complete, the handle `context_handle` is used for the generation and verification of transaction signatures.

The procedures for sending and receiving signed messages are given in [section 5](#), Sending and Verifying Signed Messages.

## [4](#). Server Protocol Details

As on the client-side, the result of a successful context negotiation is a context handle used in future processing.

A server may be managing several contexts with several clients. Clients identify their contexts by providing a key name in their request. The server maintains a mapping of key names to handles:

(`key_name`, `context_handle`)

### [4.1](#) Negotiating Context

A server recognizes TKEY queries as security context negotiation messages.



#### **4.1.1 Receive TKEY Query from Client**

Upon receiving a TKEY query, the server must examine the Mode and Algorithm Name fields to see if they match this algorithm. If they match, the (key\_name, context\_handle) mapping table is searched for the NAME value of the TKEY record. If the name is found in the table, the corresponding context\_handle is used in following GSS operations. If the name is not found a new negotiation is started.

#### **4.1.2 Call GSS\_Accept\_sec\_context**

The server performs its side of a context negotiation by calling GSS\_Accept\_sec\_context with the token provided by the client in the TKEY record.

The server calls GSS\_Accept\_sec\_context:

##### **INPUTS**

CONTEXT HANDLE	input_context_handle	= 0 if new negotiation, context_handle if ongoing
OCTET STRING	input_token	= Key field from TKEY RR

##### **OUTPUTS**

INTEGER	major_status
CONTEXT_HANDLE	output_context_handle
OCTET STRING	output_token

If this is the first call to GSS\_Accept\_sec\_context in a new negotiation, then output\_context\_handle is stored in the server's key-mapping table as the context\_handle that maps to the name of the TKEY record.

#### **4.1.3 Send TKEY Query-Response to Client**

If major\_status returns a GSS failure code, the negotiation has failed. The server must respond to the client with a standard TKEY query-response where the TKEY error field value is set to BADKEY.

Success values for major\_status are GSS\_S\_COMPLETE or GSS\_S\_CONTINUE.





If `major_status` is `GSS_S_COMPLETE` the server component of the negotiation is finished. The message from the client may be signed with a TSIG RR, and if present this signature must be verified using the procedure detailed in [section 5](#), Sending and Verifying Signed Messages. The server responds to the TKEY query using a standard query response. If `output_token` is non-NULL, then it must be returned to the client in a TKEY in the Answer section of the response. If `output_token` is NULL, then the TKEY record received from the client must be returned in the Answer section of the response. The answer should be signed with a TSIG record per the procedure given in [section 5](#), Sending and Verifying Signed Messages. The context state is advanced to Established. [Section 4.2](#) discusses the usage of the security context.

If `major_status` is `GSS_S_CONTINUE`, the server component of the negotiation is not yet finished. The server responds to the TKEY query with a standard query response, placing a TKEY record containing `output_token` in the answer section. The negotiation sequence then repeats beginning with [section 4.1.1](#). The server must limit the number of times that a given context is allowed to repeat, to prevent endless looping.

## **[4.2](#) Context Established**

When context negotiation is complete, the handle `context_handle` is used for the generation and verification of transaction signatures. The handle is valid for a finite amount of time determined by the underlying security mechanism. A server may unilaterally terminate a context at any time.

The procedures for sending and receiving signed messages are given in [section 5](#), Sending and Verifying Signed Messages.

### **[4.2.1](#) Terminating a Context**

A server can terminate any established context at any time. The server may hint to the client that the context is being deleted by including a TKEY RR in a response with the mode field set to "key deletion". See [[TKEY](#)] for more details.

An active context is deleted by calling `GSS_Delete_sec_context` providing the associated `context_handle`.

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## **5. Sending and Verifying Signed Messages**

### **5.1 Sending a Signed Message - Call GSS\_GetMIC**

The procedure for sending a signature-protected message is specified in [TSIG]. The data to be passed to the signature routine includes the whole DNS message with specific TSIG variables appended. For the exact format, see [TSIG]. For this protocol, use the following TSIG variable values:

```
TSIG Record
  NAME = key_name that identifies this context
  RDATA
    Algorithm Name = gss-tsig.microsoft.com
```

Assign the remaining fields in the TKEY RDATA appropriate values per [TKEY].

For the GSS algorithm, the signature is generated by calling GSS\_GetMIC:

```
INPUTS
  CONTEXT HANDLE context_handle = context_handle for key_name
  OCTET STRING   message       = outgoing message plus TSIG
                                variables (see [TSIG])

OUTPUTS
  INTEGER        major_status
  OCTET STRING   per_msg_token
```

If major\_status is GSS\_S\_COMPLETE, then signature generation succeeded. The signature in per\_msg\_token is inserted into the Signature field of the TSIG RR and the message is transmitted.

If major\_status is GSS\_S\_CONTEXT\_EXPIRED or GSS\_S\_CREDENTIALS\_EXPIRED, the caller needs to return to the uninitialized state and negotiate a new security context.

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## 5.2 Verifying a Signed Message - Call GSS\_VerifyMIC

The procedure for verifying a signature-protected message is specified in [TSIG].

The NAME of the TSIG record determines which context\_handle maps to this context. If the NAME does not map to a currently active context, the server must send a standard TSIG error response to the client indicating BADKEY in the TSIG error field (see [TSIG]).

For the GSS algorithm, a signature is verified by using GSS\_VerifyMIC:

### INPUTS

CONTEXT HANDLE context\_handle = context\_handle for key\_name  
OCTET STRING message = incoming message plus TSIG  
variables (see [TSIG])  
OCTET STRING per\_msg\_token = Signature field from TSIG RR

### OUTPUTS

INTEGER major\_status

If major\_status is GSS\_S\_COMPLETE, the signature is authentic and the message was delivered intact. Per [TSIG], the timer values of the TSIG record must also be valid before considering the message to be authentic. The caller must not act on the request or response in the message until these checks are verified.

If major\_status is GSS\_S\_CONTEXT\_EXPIRED, the negotiated context is no longer valid. If this failure occurs when a server is processing a client request, the server must send a standard TSIG error response to the client indicating BADKEY in the TSIG error field (see [TSIG]).

If major\_status is any other error code or if the timer values of the TSIG record are invalid, the message must not be considered authentic. If this error checking fails when a server is processing a client request, the appropriate error response must be sent to the client per [TSIG].

## 6. Security Considerations

This document describes a protocol for DNS security using GSS-API. The security provided by this protocol is only as effective as the security provided by the underlying GSS mechanisms.

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## **7. Acknowledgements**

The authors of this document would like to thank the following people for their contribution to this specification: Chuck Chan, Mike Swift, Ram Viswanathan and Donald E. Eastlake 3rd.

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