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INTERNET DRAFT

Expires: August 2, 2006

February 3 2006

The AES-CMAC-PRF-128 Algorithm for the Internet Key Exchange Protocol (IKE)

draft-songlee-aes-cmac-prf-128-03.txt

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Abstract

Some implementations of IP Security (IPsec) may want to use a pseudo-random function derived from the Advanced Encryption Standard (AES). This memo describes such an algorithm, called AES-CMAC-PRF-128.

1. Introduction

[AES-CMAC] describes a method to use the Advanced Encryption Standard (AES) as a message authentication code (MAC) whose output is 128 bits long. 128 bits output is useful as a long-lived pseudo-random function (PRF) in either IKE version 1 or version 2. This document specifies PRF that support fixed and variable key sizes for IKEv2 [IKEv2] Key Derivation Function (KDF) and authentication.

2. Basic definitions

VK	Variable length key for AES-CMAC-PRF-128, Denoted by VK.
0^n	The string that consists of n zero-bits. 0^3 means that 000 in binary format. 10^4 means that 10000 in binary format. 10^i means that 1 followed by i-times repeated

zero's.

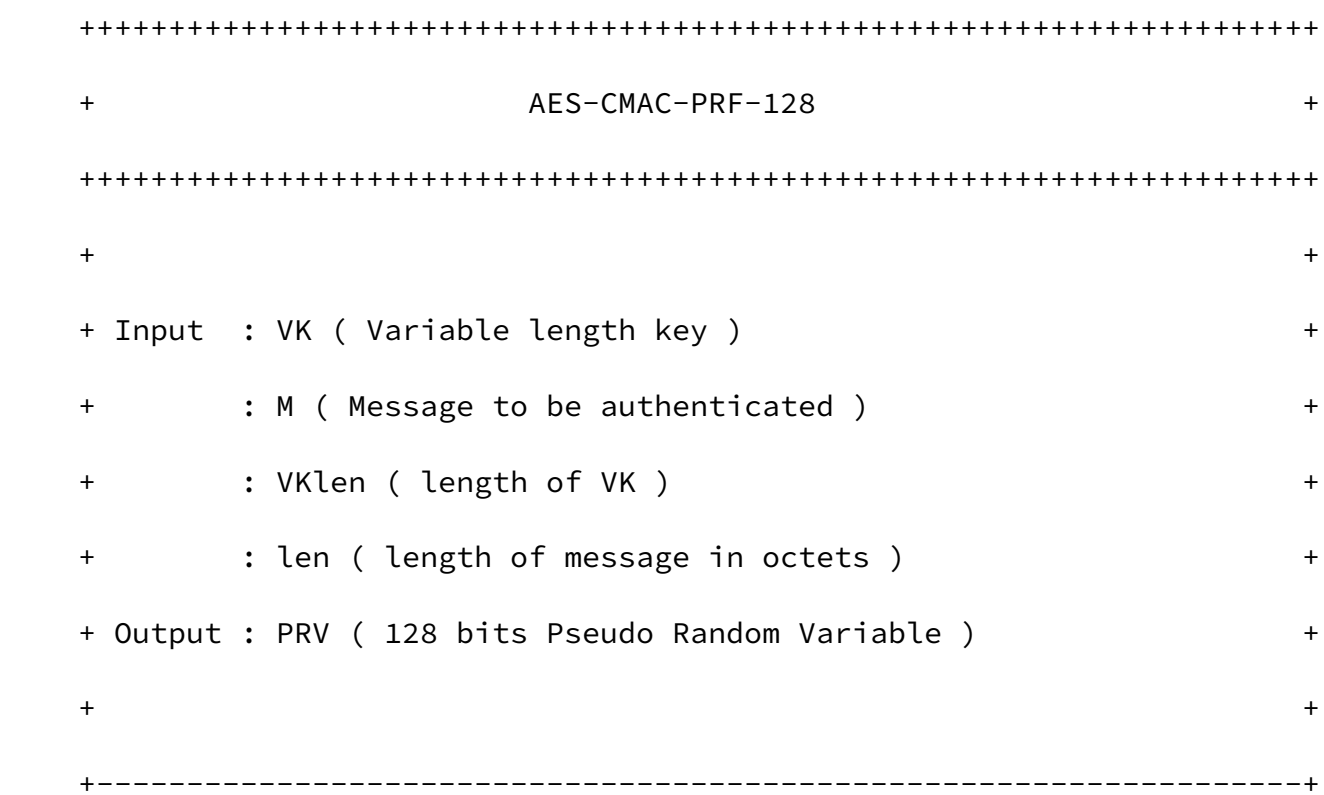
AES-CMAC AES-CMAC algorithm with 128 bits long key described in section 2.4 of [AES-CMAC].

3. The AES-CMAC-PRF-128 Algorithm

The AES-CMAC-PRF-128 algorithm is identical to AES-CMAC defined in [AES-CMAC] except that the 128 bits key length restriction is removed.

IKEv2 [IKEv2] uses PRFs for multiple purposes, most notably for generating keying material and authentication of the the IKE_SA. The IKEv2 specification differentiates between PRFs with fixed key sizes and those with variable key sizes

When using the PRF described in this document with IKEv2, the PRF is considered to be fixed-length for generating keying material but variable-length for authentication.


AES-CMAC-PRF-128
Input : VK (Variable length key)
: M (Message to be authenticated)
: VKlen (length of VK)
: len (length of message in octets)
Output : PRV (128 bits Pseudo Random Variable)
Variables: K (128-bits fixed key)
Step 1.

```
+           If VKlen is equal to 16 octets then                    +
+ Step 1a.      K := VK;                                           +
+           Else                                                   +
+ Step 1b.      K := AES-CMAC (0^128, VK, VKlen);                  +
+                                                                  +
+ Step 2.                                                         +
+           PRV := AES-CMAC (K,M,len);                             +
+           return PRV;                                            +
+                                                                  +
++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++
```

Figure 1. AES-CMAC-PRF-128 Algorithm

In step 1, the key for AES-CMAC-PRF-128 is created as follows:

o If the key is exactly 128 bits long, use it as-is.

o If the key is longer or shorter than 128 bits long, then we derive new key K by performing AES-CMAC algorithm using 128 bits all zero key and VK as the message. This step is described in step 1b.

In step 2, we perform AES-CMAC algorithm using K as the key and M as the message. The output of this algorithm is returned.

5. Test Vectors

```
Test Case AES-CMAC-PRF-128 with 20-octet input
Key        : 00010203 04050607 08090a0b 0c0d0e0f edcb
Key Length : 18
Message    : 00010203 04050607 08090a0b 0c0d0e0f 10111213
PRF Output : 84a348a4 a45d235b abfffc0d 2b4da09a

Test Case AES-CMAC-PRF-128 with 20-octet input
Key        : 00010203 04050607 08090a0b 0c0d0e0f
Key Length : 16
Message    : 00010203 04050607 08090a0b 0c0d0e0f 10111213
PRF Output : 980ae87b 5f4c9c52 14f5b6a8 455e4c2d
```

```
Test Case AES-CMAC-PRF-128 with 20-octet input
Key        : 00010203 04050607 0809
Key Length : 10
Message    : 00010203 04050607 08090a0b 0c0d0e0f 10111213
PRF Output : 290d9e11 2edb09ee 141fcf64 c0b72f3d

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```

6. Security Considerations

The security provided by AES-CMAC-PRF-128 is based upon the strength of AES and AES-CMAC. At the time of this writing, there are no known practical cryptographic attacks against AES or AES-CMAC. However as is true with any cryptographic algorithm, part of its strength lies in the secret key, 'K' and the correctness of the implementation in all of the participating systems. Keys need to be chosen at random based on RFC 4086 [RFC4086] and should be kept in safe and periodically refreshed.

Whenever keys larger than 128 bits are reduced to meet AES-128 key input size, some entropy might be lost. However, if using collision-resistant hash function such as AES-CMAC when generating new key for pseudo-random function, it preserves sufficient entropy as long as the pseudo-random function to be used requires 128 bits long input key.

7. IANA Consideration

IANA should allocate a value for IKEv2 Transform Type 2

(Pseudo-Random Function) to the PRF_AES128_CMAC algorithm when this document is published.

8. Acknowledgement

Portions of this text were borrowed from [AES-XCBC-PRF] and [AES-XCBC-PRF_bis], and many thanks to Russ Housley and Paul Hoffman for suggestions and guidance.

9. Reference

9.1 Normative References

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[RFC4086] Eastlake 3rd, D., Crocker, S., and J. Schiller, "Randomness Requirements for Security", RFC 4086 June 2005

9.2. Informative References

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Acknowledgment

Funding for the RFC Editor function is currently provided by the Internet Society.