Group OSCORE Profile of the Authentication and Authorization for Constrained Environments Framework

Abstract

This document specifies a profile for the Authentication and Authorization for Constrained Environments (ACE) framework. The profile uses Group OSCORE to provide communication security between a Client and a (set of) Resource Server(s) as members of an OSCORE Group. The profile securely binds an OAuth 2.0 Access Token with the public key of the Client associated with the private key used in the OSCORE group. The profile uses Group OSCORE to achieve server authentication, as well as proof-of-possession for the Client's public key. Also, it provides proof of the Client's membership to the correct OSCORE group, by binding the Access Token to information from the Group OSCORE Security Context, thus allowing the Resource Server(s) to verify the Client's membership upon receiving a message protected with Group OSCORE from the Client.

Discussion Venues

This note is to be removed before publishing as an RFC.

Discussion of this document takes place on the Constrained RESTful Environments Working Group mailing list (ace@ietf.org), which is archived at https://mailarchive.ietf.org/arch/browse/ace/.

Source for this draft and an issue tracker can be found at https://gitlab.com/crimson84/draft-tiloca-ace-group-oscore-profile.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

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1. Introduction

A number of applications rely on a group communication model, where a Client can access a resource shared by multiple Resource Servers at once, e.g., over IP multicast. Typical examples are switching of luminaries, actuators control, and distribution of software updates. Secure communication in the group can be achieved by sharing a set of keying material, which is typically provided upon joining the group.

For some of such applications, it may be just fine to enforce access control in a straightforward fashion. That is, any Client authorized to join the group, hence to get the group keying material, can be also implicitly authorized to perform any action at any resource of any Server in the group. An example of application where such implicit authorization might be used is a simple lighting scenario, where the lightbulbs are the Servers, while the user account on an app on the user's phone is the Client. In this case, it might be fine to not require additional authorization evidence from any user account, if it is acceptable that any current group member is also authorized to switch on and off any light, or to check their status.

However, in different instances of such applications, the approach above is not desirable, as different group members are intended to have different access rights to resources of other group members. That is, access control to the secure group communication channel and access control to the resource space provided by servers in the group should remain logically separated domains. For instance, a more fine-grained approach is required in the two following use cases.

As a first case, an application provides control of smart locks acting as Servers in the group, where: a first type of Client, e.g., a user account of a child, is allowed to only query the status of the smart locks; while a second type of Client, e.g., a user account of a parent, is allowed to both query and change the status of the smart locks. Further similar applications concern the enforcement of different sets of permissions in groups with sensor/actuator devices, e.g., thermostats, acting as Servers. Also, some group members may even be intended as Servers only. Hence, they must be prevented from acting as Clients altogether and from accessing resources at other Servers, especially when attempting to perform non-safe operations.

As a second case, building automation scenarios often rely on Servers that, under different circumstances, enforce different level of priority for processing received commands. For instance, BACnet deployments consider multiple classes of Clients, e.g., a normal light switch (C1) and an emergency fire panel (C2). Then, a C1 Client is not allowed to override a command from a C2 Client, until the latter relinquishes control at its higher priority. That is: i) only C2 Clients should be able to adjust the minimum required level of priority on the Servers, so rightly locking out C1 Clients if needed; and ii) when a Server is set to accept only high-priority commands, only C2 Clients should be able to perform such commands otherwise allowed also to C1 Clients. Given the different maximum authority of different Clients, fine-grained access control would effectively limit the execution of high- and emergency-priority commands only to devices that are in fact authorized to do so.

Besides, it would prevent a misconfigured or compromised device from
initiating a high-priority command and lock out normal control.

In the cases above, being a legitimate group member and storing the
group keying material is not supposed to imply any particular access
rights. Also, introducing a different security group for each
different set of access rights would result in additional keying
material to distribute and manage. In particular, if the access
rights for a single node change, this would require to evict that
node from the current group, followed by that node joining a
different group aligned with its new access rights. Moreover, the
keying material of both groups would have to be renewed for their
current members. Overall, this would have a non negligible impact on
operations and performance in the system.

A fine-grained access control model can be rather enforced within a
same group, by using the Authentication and Authorization for
Constrained Environments (ACE) framework [I-D.ietf-ace-oauth-authz].
That is, a Client has to first obtain authorization credentials in
the form of an Access Token, and post it to the Resource Server(s)
in the group before accessing the intended resources.

The ACE framework delegates to separate profile documents how to
secure communications between the Client and the Resource Server.
However each of the current profiles of ACE defined in [I-D.ietf-
ace-oscore-profile] [I-D.ietf-ace-dtls-authorize] [I-D.ietf-ace-
mqtt-tls-profile] admits a single security protocol that cannot be
used to protect group messages sent over IP multicast.

This document specifies the "coap_group_oscore" profile of the ACE
framework, where a Client uses CoAP [RFC7252] or CoAP over IP
multicast [I-D.ietf-core-groupcomm-bis] to communicate to one or
multiple Resource Servers, which are members of an application group
and share a common set of resources. This profile uses Group OSCORE
[I-D.ietf-core-oscore-groupcomm] as the security protocol to protect
messages exchanged between the Client and the Resource Servers.
Hence, it requires that both the Client and the Resource Servers
have previously joined the same OSCORE group.

That is, this profile describes how access control is enforced for a
Client after it has joined an OSCORE group, to access resources at
other members in that group. The process for joining the OSCORE
group through the respective Group Manager as defined in [I-D.ietf-
ace-key-groupcomm-oscore] takes place before the process described
in this document, and is out of the scope of this profile.

The Client proves its access to be authorized to the Resource Server
by using an Access Token, which is bound to a key (the proof-of-
possession key). This profile uses Group OSCORE to achieve server

authentication, as well as proof-of-possession for the Client's public key used in the OSCORE group in question. Note that the proof of possession is not done by a dedicated protocol element, but rather occurs after the first Group OSCORE exchange.

Furthermore, this profile provides proof of the Client's membership to the correct OSCORE group, by binding the Access Token to the Client's authentication credential used in the group and including the Client's public public key, as well as to information from the pre-established Group OSCORE Security Context. This allows the Resource Server to verify the Client's group membership upon reception of a message protected with Group OSCORE from that Client.

OSCORE [RFC8613] specifies how to use COSE [I-D.ietf-cose-rfc8152bis-struct][I-D.ietf-cose-rfc8152bis-algs] to secure CoAP messages. Group OSCORE builds on OSCORE to provide secure group communication, and ensures source authentication: by means of digital signatures embedded in protected messages (in group mode); or by protecting messages with pairwise keying material derived from the asymmetric keys of the two peers exchanging the message (in pairwise mode).

1.1. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

Readers are expected to be familiar with the terms and concepts related to CBOR [RFC8949], COSE [I-D.ietf-cose-rfc8152bis-struct][I-D.ietf-cose-rfc8152bis-algs], CoAP [RFC7252], OSCORE [RFC8613] and Group OSCORE [I-D.ietf-core-oscore-groupcomm]. These especially include:

  *Group Manager, as the entity responsible for a set of groups
   where communications among members are secured with Group OSCORE.

  *Authentication credential, as the set of information associated
   with an entity, including that entity's public key and parameters
   associated with the public key. Examples of authentication
   credentials are CBOR Web Tokens (CWTs) and CWT Claims Sets (CCSs)
   [RFC8392], X.509 certificates [RFC7925] and C509 certificates [I-D.ietf-cose-cbor-encoded-cert].

   Members of an OSCORE group have an associated authentication
   credential in the format used in the group. As per Section 2.3 of
   [I-D.ietf-core-oscore-groupcomm], an authentication credential
   provides the public key as well as the comprehensive set of

information related to the public key algorithm, including, e.g., the used elliptic curve (when applicable).

Readers are expected to be familiar with the terms and concepts described in the ACE framework for authentication and authorization [I-D.ietf-ace-oauth-authz], as well as in the OSCORE profile of ACE [I-D.ietf-ace-oscore-profile]. The terminology for entities in the considered architecture is defined in OAuth 2.0 [RFC6749]. In particular, this includes Client (C), Resource Server (RS), and Authorization Server (AS).

Note that, unless otherwise indicated, the term "endpoint" is used here following its OAuth definition, aimed at denoting resources such as /token and /introspect at the AS, and /authz-info at the RS. This document does not use the CoAP definition of "endpoint", which is "An entity participating in the CoAP protocol".

Additionally, this document makes use of the following terminology.

  *Equivalent COSE Key: a COSE Key built from an authentication credential used in the OSCORE group. The equivalent COSE Key preserves all the main information elements from the authentication credential, in particular the key coordinates and the comprehensive set of information related to the public key algorithm, including, e.g., the used elliptic curve (when applicable).

  *Pairwise-only group: an OSCORE group that uses only the pairwise mode (see Section 9 of [I-D.ietf-core-oscore-groupcomm]).

Examples throughout this document are expressed in CBOR diagnostic notation, without the tag and value abbreviations.

2. Protocol Overview

This section provides an overview of this profile, i.e., on how to use the ACE framework for authentication and authorization [I-D.ietf-ace-oauth-authz] to secure communications between a Client and a (set of) Resource Server(s) using Group OSCORE [I-D.ietf-core-oscore-groupcomm].

Note that this profile of ACE describes how access control can be enforced for a node after it has joined an OSCORE group, to access resources at other members in that group.

In particular, the process for joining the OSCORE group through the respective Group Manager as defined in [I-D.ietf-ace-key-groupcomm-oscore] must take place before the process described in this document, and is out of the scope of this profile.

An overview of the protocol flow for this profile is shown in Figure 1. In the figure, it is assumed that both RS1 and RS2 are associated with the same AS. It is also assumed that C, RS1 and RS2 have previously joined an OSCORE group with Group Identifier (gid) "abcd0000", and got assigned Sender ID (sid) "0", "1" and "2" in the group, respectively. The names of messages coincide with those of [I-D.ietf-ace-oauth-authz] when applicable.

```
C                            RS1            RS2                          AS
| [--- Resource Request -->] |              |                            |
|                            |              |                            |
| [<----  AS Request ------] |              |                            |
|        Creation Hints      |              |                            |
|                            |              |                            |
|-------- POST /token -------------------------------------------------->|
|   (aud: RS1, sid: 0, gid: abcd0000, ...) |                            |
|                            |              |                            |
|                            |              |                            |
|<----------------------------- Access Token + RS Information -----|
|                            | (aud: RS1, sid: 0, gid: abcd0000, ...) |
|                            |              |                            |
|---- POST /authz-info ----->|              |                            |
|       (access_token)       |              |                            |
|                            |              |                            |
|<--- 2.01 Created ----------|              |                            |
|                            |              |                            |
|-------- POST /token -------------------------------------------------->|
|   (aud: RS2, sid: 0, gid: abcd0000, ...) |                            |
|                            |              |                            |
|                            |              |                            |
|                            |              |                            |
|<----------------------------- Access Token + RS Information -----|
|                            | (aud: RS2, sid: 0, gid: abcd0000, ...) |
|                            |              |                            |
|----- POST /authz-info ----------------->|                            |
|        (access_token)      |              |                            |
|                            |              |                            |
|                            |              |                            |
|<--- 2.01 Created ---------------------|                            |
|                            |              |                            |
|-- Group OSCORE Request -+->|              |                            |
| (kid: 0, gid: abcd0000) \-------------->|                            |
|                            |              |                            |
|                         /proof-of-possession/                       |
|                            |              |                            |
|                            |              |                            |
|<-- Group OSCORE Response --|              |                            |
|        (kid: 1)            |              |                            |
|                            |              |                            |
/proof-of-possession/        |              |                            |
|                            |              |                            |
/Mutual authentication       |              |                            |
 between C and RS1/          |              |                            |
|                            |              |                            |
|                            |              |                            |
|<-- Group OSCORE Response --------------|                            |
|        (kid: 2)            |              |                            |
```

```
|                        |           |                   |
/proof-of-possession/     |           |                   |
|                        |           |                   |
/Mutual authentication    |           |                   |
 between C and RS2/        |           |                   |
|                        |           |                   |
|            ...           |           |                   |
```

Figure 1: Protocol Overview.

2.1. Pre-Conditions

Using Group OSCORE and this profile requires both the Client and the Resource Servers to have previously joined the same OSCORE group. This especially includes the derivation of the Group OSCORE Security Context and the assignment of unique Sender IDs to use in the group. Nodes may join the OSCORE group through the respective Group Manager by using the approach defined in [I-D.ietf-ace-key-groupcomm-oscore], which is also based on ACE.

After the Client and Resource Servers have joined the group, this profile provides access control for accessing resources on those Resource Servers, by securely communicating with Group OSCORE.

As a pre-requisite for this profile, the Client has to have successfully joined the OSCORE group where also the Resource Servers (RSs) are members. Depending on the limited information initially available, the Client may have to first discover the exact OSCORE group used by the RSs for the resources of interest, e.g., by using the approach defined in [I-D.tiloca-core-oscore-discovery].

2.2. Access Token Retrieval

This profile requires that the Client retrieves an Access Token from the AS for the resource(s) it wants to access on each of the RSs, using the /token endpoint, as specified in Section 5.8 of [I-D.ietf-ace-oauth-authz]. In a general case, it can be assumed that different RSs are associated with different ASs, even if the RSs are members of a same OSCORE group.

In the Access Token request to the AS, the Client MUST include the Group Identifier of the OSCORE group and its own Sender ID in that group. The AS MUST specify these pieces of information in the Access Token, included in the Access Token response to the Client.

Furthermore, in the Access Token request to the AS, the Client MUST also include: its own public key used in the OSCORE group; and a proof-of-possession (PoP) evidence to proof possession of the corresponding private key. The PoP evidence is computed over a PoP input uniquely related to the secure communication association between the Client and the AS. The AS MUST include also the public key indicated by the Client in the Access Token.

The Access Token request and response MUST be confidentiality-protected and ensure authenticity. This profile RECOMMENDS the use of OSCORE between the Client and the AS, to reduce the number of libraries the client has to support. Other protocols fulfilling the security requirements defined in Sections 5 and 6 of [I-D.ietf-ace-

oauth-authz] MAY alternatively be used, such as TLS [RFC8446] or
DTLS [RFC6347][I-D.ietf-tls-dtls13].

2.3. Access Token Posting

After having retrieved the Access Token from the AS, the Client
posts the Access Token to the RS, using the /authz-info endpoint and
mechanisms specified in Section 5.10 of [I-D.ietf-ace-oauth-authz],
as well as Content-Format = application/ace+cbor. When using this
profile, the communication with the /authz-info endpoint is not
protected.

If the Access Token is valid, the RS replies to this POST request
with a 2.01 (Created) response with Content-Format = application/
ace+cbor. Also, the RS associates the received Access Token with the
Group OSCORE Security Context identified by the Group Identifier
specified in the Access Token, following Section 3.2 of [RFC8613].
In practice, the RS maintains a collection of Security Contexts with
associated authorization information, for all the clients that it is
currently communicating with, and the authorization information is a
policy used as input when processing requests from those clients.

Finally, the RS stores the association between i) the authorization
information from the Access Token; and ii) the Group Identifier of
the OSCORE group together with the Sender ID and the authentication
credential of the Client in that group. This binds the Access Token
with the Group OSCORE Security Context of the OSCORE group.

Finally, when the Client communicates with the RS using the Group
OSCORE Security Context, the RS verifies that the Client is a
legitimate member of the OSCORE group and especially the exact group
member with the same Sender ID associated with the Access Token.
This occurs when verifying a request protected with Group OSCORE,
since the request includes the Client's Sender ID and either it
embeds a signature computed also over that Sender ID (if protected
with the group mode), or it is protected by means of pairwise
symmetric keying material derived from the asymmetric keys of the
two peers (if protected with the pairwise mode).

2.4. Secure Communication

The Client can send a request protected with Group OSCORE [I-D.ietf-
core-oscore-groupcomm] to the RS. This can be a unicast request
addressed to the RS, or a multicast request addressed to the OSCORE
group where the RS is also a member. To this end, the Client uses
the Group OSCORE Security Context already established upon joining
the OSCORE group, e.g., by using the approach defined in [I-D.ietf-
ace-key-groupcomm-oscore]. The RS may send a response back to the

Client, protecting it by means of the same Group OSCORE Security Context.

3. Client-AS Communication

This section details the Access Token POST Request that the Client sends to the /token endpoint of the AS, as well as the related Access Token response.

The Access Token MUST be bound to the public key of the client as proof-of-possession key (pop-key), by means of the 'cnf' claim.

3.1. C-to-AS: POST to Token Endpoint

The Client-to-AS request is specified in Section 5.8.1 of [I-D.ietf-ace-oauth-authz]. The Client MUST send this POST request to the /token endpoint over a secure channel that guarantees authentication, message integrity and confidentiality.

The POST request is formatted as the analogous Client-to-AS request in the OSCORE profile of ACE (see Section 3.1 of [I-D.ietf-ace-oscore-profile]), with the following additional parameters that MUST be included in the payload.

  * 'context_id', defined in Section 3.1.1 of this document. This parameter specifies the Group Identifier (GID), i.e., the Id Context of an OSCORE group where the Client and the RS are currently members. In particular, the Client wishes to communicate with the RS using the Group OSCORE Security Context associated with that OSCORE group.

  * 'salt_input', defined in Section 3.1.2 of this document. This parameter includes the Sender ID that the Client has in the OSCORE group whose GID is specified in the 'context_id' parameter above.

  * 'req_cnf', defined in Section 3.1 of [I-D.ietf-ace-oauth-params]. This parameter follows the syntax from Section 3.1 of [RFC8747] when including Value Type "COSE_Key" (1) and specifying an asymmetric key. In particular, the specified public key is the COSE Key equivalent to the authentication credential that the Client uses in the OSCORE group. The specified public key will be used as the pop-key bound to the Access Token.

    Alternative Value Types defined in future specifications are fine to consider, if indicating a non-encrypted asymmetric key or full-fledged autentication credential.

In addition, the Client computes its proof-of-possession (PoP) evidence, in order to prove possession of its own private key used

in the OSCORE group to the AS. This allows the AS to verify that the Client indeed owns the private key associated with that public key, as its alleged identity credential within the OSCORE group.

To this end, the Client MUST use as PoP input the byte representation of a quantity that uniquely represents the secure communication association between the Client and the AS. It is RECOMMENDED that the Client considers the following as PoP input.

  *If the Client and the AS communicate over (D)TLS, the PoP input is an exporter value computed as defined in Section 7.5 of [RFC8446]. In particular, the exporter label MUST be 'EXPORTER-ACE-Sign-Challenge-Client-AS' defined in Section 10.5 of this document, together with an empty 'context_value', and 32 bytes as 'key_length'.

  *If the Client and the AS communicate over OSCORE, the PoP input is the output PRK of a HKDF-Extract step [RFC5869], i.e., PRK = HMAC-Hash(salt, IKM). In particular, 'salt' takes (x1 | x2), where x1 is the ID Context of the OSCORE Security Context between the Client and the AS, x2 is the Sender ID of the Client in that Security Context, and | denotes byte string concatenation. Also, 'IKM' is the OSCORE Master Secret of the OSCORE Security Context between the Client and the AS.

  The HKDF MUST be one of the HMAC-based HKDF [RFC5869] algorithms defined for COSE [I-D.ietf-cose-rfc8152bis-algs]. The Client and AS may agree on the HKDF algorithm to use during the Client's registration at the AS. HKDF SHA-256 is mandatory to implement.

Then, the Client computes the PoP evidence as follows.

  *If the OSCORE group is not a pairwise-only group, the PoP evidence MUST be a signature. The Client computes the signature by using the same private key and signature algorithm it uses for signing messages in the OSCORE group. The private key corresponds to the authentication credential used in the OSCORE group, for which the equivalent COSE Key is specified in the 'req_cnf' parameter above.

  *If the OSCORE group is a pairwise-only group, the PoP evidence MUST be a MAC computed as follows, by using the HKDF Algorithm HKDF SHA-256, which consists of composing the HKDF-Extract and HKDF-Expand steps [RFC5869].

  MAC = HKDF(salt, IKM, info, L)

  The input parameters of HKDF are as follows.

    -salt takes as value the empty byte string.

-IKM is computed as a cofactor Diffie-Hellman shared secret, see Section 5.7.1.2 of [NIST-800-56A], using an ECDH algorithm pre-agreed between Client and AS. The Client uses its own Diffie-Hellman private key and the Diffie-Hellman public key of the AS. For X25519 and X448, the procedure is described in Section 5 of [RFC7748].

The Client's private key corresponds to the Client's authentication credential used in the OSCORE group, for which the equivalent COSE Key is specified in the 'req_cnf' parameter above. The Client may obtain the Diffie-Hellman public key of the AS during its registration process at the AS.

The Client and AS may agree on the ECDH algorithm to use during the Client's registration at the AS. The ECDH-SS + HKDF-256 algorithm specified in Section 6.3.1 of [I-D.ietf-cose-rfc8152bis-algs] is mandatory to implement.

-info takes as value the PoP input.

-L is equal to 8, i.e., the size of the MAC, in bytes.

Finally, the Client MUST include one of the two following parameters in the payload of the POST request to the AS.

*'client_cred_verify', defined in Section 3.1.3 of this document, specifying the Client's PoP evidence as a signature, which is computed as defined above. This parameter MUST be included if and only if the OSCORE group is not a pairwise-only group.

*'client_cred_verify_mac', defined in Section 3.1.4 of this document, specifying the Client's PoP evidence as a MAC, which is computed as defined above. This parameter MUST be included if and only if the OSCORE group is a pairwise-only group.

An example of such a request is shown in Figure 2.

```
   Header: POST (Code=0.02)
   Uri-Host: "as.example.com"
   Uri-Path: "token"
   Content-Format: "application/ace+cbor"
   Payload:
   {
     "audience" : "tempSensor4711",
     "scope" : "read",
     "context_id" : h'abcd0000',
     "salt_input" : h'00',
     "req_cnf" : {
       "COSE_Key" : {
         "kty" : EC2,
         "crv" : P-256,
         "x" : h'd7cc072de2205bdc1537a543d53c60a6acb62eccd890c7fa
                 27c9e354089bbe13',
         "y" : h'f95e1d4b851a2cc80fff87d8e23f22afb725d535e515d020
                 731e79a3b4e47120'
       }
     },
     "client_cred_verify" : h'...'
     (signature content omitted for brevity)
   }
```

Figure 2: Example C-to-AS POST /token request for an Access Token bound
to an asymmetric key.

3.1.1. 'context_id' Parameter

The 'context_id' parameter is an OPTIONAL parameter of the Access
Token request message defined in [Section 5.8.1](#) of [[I-D.ietf-ace-
oauth-authz](#)]. This parameter provides a value that the Client wishes
to use with the RS as a hint for a security context. Its exact
content is profile specific.

3.1.2. 'salt_input' Parameter

The 'salt_input' parameter is an OPTIONAL parameter of the Access
Token request message defined in [Section 5.8.1](#) of [[I-D.ietf-ace-
oauth-authz](#)]. This parameter provides a value that the Client wishes
to use as part of a salt with the RS, for deriving cryptographic
keying material. Its exact content is profile specific.

3.1.3. 'client_cred_verify' Parameter

The 'client_cred_verify' parameter is an OPTIONAL parameter of the
Access Token request message defined in [Section 5.8.1.](#) of [[I-D.ietf-
ace-oauth-authz](#)]. This parameter provides a signature computed by
the Client to prove the possession of its own private key.

3.1.4. 'client_cred_verify_mac' Parameter

The 'client_cred_verify_mac' parameter is an OPTIONAL parameter of
the Access Token request message defined in [Section 5.8.1.](#) of [[I-D.ietf-ace-oauth-authz](#)]. This parameter provides a Message
Authentication Code (MAC) computed by the Client to prove the
possession of its own private key.

3.2. AS-to-C: Access Token

After having verified the POST request to the /token endpoint and
that the Client is authorized to obtain an Access Token
corresponding to its Access Token request, the AS MUST verify the
proof-of-possession (PoP) evidence. In particular, the AS proceeds
as follows.

  *As PoP input, the AS uses the same value considered by the Client
   in [Section 3.1](#).

  *As public key of the Client, the AS uses the one specified in the
   'req_cnf' parameter of the Access Token request.

  *If the Access Token request includes the 'client_cred_verify'
   parameter, this specifies the PoP evidence as a signature. Then,
   the AS verifies the signature by using the public key of the
   Client.

  *If the Access Token request includes the 'client_cred_verify_mac'
   parameter, this specifies the PoP evidence as a Message
   Authentication Code (MAC).

   Then, the AS recomputes the MAC through the same process taken by
   the Client when preparing the value of the
   'client_cred_verify_mac' parameter for the Access Token (see
   [Section 3.1](#)), with the difference that the AS uses its own
   Diffie-Hellman private key and the Diffie-Hellman public key of
   the Client. The verification succeeds if and only if the
   recomputed MAC is equal to the MAC conveyed as PoP evidence in
   the Access Token request.

If both the 'client_cred_verify' and 'client_cred_verify_mac'
parameters are present, or if the verification of the PoP evidence
fails, the AS considers the Client request invalid.

If the Client request was invalid, or not authorized, the AS returns
an error response as described in [Section 5.8.3](#) of [[I-D.ietf-ace-oauth-authz](#)].

If all verifications are successful, the AS responds as defined in
Section 5.8.2 of [I-D.ietf-ace-oauth-authz]. In particular:

  *The AS can signal that the use of Group OSCORE is REQUIRED for a
   specific Access Token by including the 'ace_profile' parameter
   with the value "coap_group_oscore" in the Access Token response.
   The Client MUST use Group OSCORE towards all the Resource Servers
   for which this Access Token is valid. Usually, it is assumed that
   constrained devices will be pre-configured with the necessary
   profile, so that this kind of profile signaling can be omitted.

  *The AS MUST NOT include the 'rs_cnf' parameter defined in [I-
   D.ietf-ace-oauth-params]. In general, the AS may not be aware of
   the authentication credentials (and public keys included thereof)
   that the RSs use in the OSCORE group. Also, the Client is able to
   retrieve the authentication credentials of other group members
   from the responsible Group Manager, both upon joining the group
   or later on as a group member, as defined in [I-D.ietf-ace-key-
   groupcomm-oscore].

The AS MUST include the following information as metadata of the
issued Access Token. The use of CBOR web tokens (CWT) as specified
in [RFC8392] is RECOMMENDED.

  *The profile "coap_group_oscore". If the Access Token is a CWT,
   this is placed in the 'ace_profile' claim of the Access Token, as
   per Section 5.10 of [I-D.ietf-ace-oauth-authz].

  *The salt input specified in the 'salt_input' parameter of the
   Token Request. If the Access Token is a CWT, the content of the
   'salt_input' parameter MUST be placed in the 'salt_input' claim
   of the Access Token, defined in Section 3.2.1 of this document.

  *The Context Id input specified in the 'context_id' parameter of
   the Token Request. If the Access Token is a CWT, the content of
   the 'context_id' parameter MUST be placed in the
   'contextId_input' claim of the Access Token, defined in Section
   3.2.2 of this document.

  *The public key that the client uses in the OSCORE group and
   specified in the 'req_cnf' parameter of the Token request.

   If the Access Token is a CWT, the public key MUST be specified in
   the 'cnf' claim, which follows the syntax from Section 3.1 of
   [RFC8747] when including Value Type "COSE_Key" (1) and specifying
   an asymmetric key. In particular, the 'cnf' claim includes the
   same COSE Key specified in the 'req_cnf' parameter of the Token
   Request, i.e., the COSE Key equivalent to the authentication
   credential that the Client uses in the OSCORE group.

Alternative Value Types defined in future specifications are fine
to consider, if indicating a non-encrypted asymmetric key or
full-fledged autentication credential.

Figure 3 shows an example of such an AS response. The access token
has been truncated for readability.

```
Header: Created (Code=2.01)
Content-Type: "application/ace+cbor"
Payload:
{
  "access_token" : h'8343a1010aa2044c53 ...'
   (remainder of CWT omitted for brevity),
  "ace_profile" : "coap_group_oscore",
  "expires_in" : 3600
}
```

Figure 3: Example AS-to-C Access Token response with the Group OSCORE
profile.

Figure 4 shows an example CWT Claims Set, containing the Client's
public key in the group (as pop-key) in the 'cnf' claim.

```
{
  "aud" : "tempSensorInLivingRoom",
  "iat" : "1360189224",
  "exp" : "1360289224",
  "scope" :  "temperature_g firmware_p",
  "cnf" : {
    "COSE_Key" : {
      "kty" : EC2,
      "crv" : P-256,
      "x" : h'd7cc072de2205bdc1537a543d53c60a6acb62eccd890c7fa
              27c9e354089bbe13',
      "y" : h'f95e1d4b851a2cc80fff87d8e23f22afb725d535e515d020
              731e79a3b4e47120'
  },
  "salt_input" : h'00',
  "contextId_input" : h'abcd0000'
}
```

Figure 4: Example CWT Claims Set with OSCORE parameters.

The same CWT Claims Set as in Figure 4 and encoded in CBOR is shown
in Figure 5, using the value abbreviations defined in [I-D.ietf-ace-
oauth-authz] and [RFC8747]. The bytes in hexadecimal are reported in
the first column, while their corresponding CBOR meaning is reported

after the "#" sign on the second column, for easiness of
readability.

NOTE: it should be checked (and in case fixed) that the values used
below (which are not yet registered) are the final values registered
by IANA.

```
A7                                        # map(7)
   03                                     # unsigned(3)
   76                                     # text(22)
      74656D7053656E736F72496E4C6976696E67526F6F6D
   06                                     # unsigned(6)
   1A 5112D728                            # unsigned(1360189224)
   04                                     # unsigned(4)
   1A 51145DC8                            # unsigned(1360289224)
   09                                     # unsigned(9)
   78 18                                  # text(24)
      74656D70657261747572655F67206669726D776172655F70
   08                                     # unsigned(8)
   A1                                     # map(1)
      01                                  # unsigned(1)
      A4                                  # map(4)
         01                               # unsigned(1)
         02                               # unsigned(2)
         20                               # negative(0)
         01                               # unsigned(1)
         21                               # negative(1)
         58 20                            # bytes(32)
            D7CC072DE2205BDC1537A543D53C60A6ACB62ECCD890C7FA27C9
            E354089BBE13
         22                               # negative(2)
         58 20                            # bytes(32)
            F95E1D4B851A2CC80FFF87D8E23F22AFB725D535E515D020731E
            79A3B4E47120
   18 3C                                  # unsigned(60)
   41                                     # bytes(1)
      00
   18 3D                                  # unsigned(61)
   44                                     # bytes(4)
      ABCD0000
```

 Figure 5: Example CWT Claims Set with OSCORE parameters, CBOR encoded.

3.2.1. Salt Input Claim

The 'salt_input' claim provides a value that the Client requesting
the Access Token wishes to use as a part of a salt with the RS,
e.g., for deriving cryptographic material.

This parameter specifies the value of the salt input, encoded as a
CBOR byte string.

3.2.2. Context ID Input Claim

The 'contextId_input' claim provides a value that the Client
requesting the Access Token wishes to use with the RS, as a hint for
a security context.

This parameter specifies the value of the Context ID input, encoded
as a CBOR byte string.

4. Client-RS Communication

This section details the POST request and response to the /authz-
info endpoint between the Client and the RS.

The proof-of-possession required to bind the Access Token to the
Client is explicitly performed when the RS receives and verifies a
request from the Client protected with Group OSCORE, either with the
group mode (see Section 8 of [I-D.ietf-core-oscore-groupcomm]) or
with the pairwise mode (see Section 9 of [I-D.ietf-core-oscore-
groupcomm]).

In particular, the RS uses the Client's public key bound to the
Access Token, either when verifying the signature of the request (if
protected with the group mode), or when verifying the request as
integrity-protected with pairwise keying material derived from the
two peers' authentication credentials and asymmetric keys (if
protected with the pairwise mode). In either case, the RS also
authenticates the Client.

Similarly, when receiving a protected response from the RS, the
Client uses the RS's public key either when verifying the signature
of the response (if protected with the group mode), or when
verifying the response as integrity-protected with pairwise keying
material derived from the two peers' authentication credentials and
asymmetric keys (if protected with the pairwise mode). In either
case, the Client also authenticates the RS. Mutual authentication is
only achieved after the client has successfully verified the
protected response from the RS.

Therefore, an attacker using a stolen Access Token cannot generate a
valid Group OSCORE message as protected through the Client's private
key, and thus cannot prove possession of the pop-key bound to the
Access Token. Also, if a Client legitimately owns an Access Token
but has not joined the OSCORE group, it cannot generate a valid
Group OSCORE message, as it does not store the necessary keying
material shared among the group members.

Furthermore, a Client C1 is supposed to obtain a valid Access Token from the AS, as including the public key associated with its own private key used in the OSCORE group, together with its own Sender ID in that OSCORE group (see Section 3.1). This allows the RS receiving an Access Token to verify with the Group Manager of that OSCORE group whether such a Client has indeed that Sender ID and an authentication credential including that public key in the OSCORE group.

As a consequence, a different Client C2, also member of the same OSCORE group, is not able to impersonate C1, by: i) getting a valid Access Token, specifying the Sender ID of C1 and a different (made-up) public key; ii) successfully posting the Access Token to RS; and then iii) attempting to communicate using Group OSCORE impersonating C1, while blaming C1 for the consequences.

4.1. C-to-RS POST to authz-info Endpoint

The Client posts the Access Token to the /authz-info endpoint of the RS, as defined in Section 5.10.1 of [I-D.ietf-ace-oauth-authz].

4.2. RS-to-C: 2.01 (Created)

The RS MUST verify the validity of the Access Token as defined in Section 5.10.1 of [I-D.ietf-ace-oauth-authz], with the following additions.

  *The RS MUST check that the claims 'salt_input', 'contextId_input' and 'cnf' are included in the Access Token.

  *The RS considers: the content of the 'contextId_input' claim as the GID of the OSCORE group; the content of the 'salt_input' claim as the Sender ID that the Client has in the group; and the content of the 'cnf' claim as the COSE Key equivalent to the authentication credential that the Client uses in the group.

  The RS MUST check whether it already stores an authentication credential associated with the pair (GID, Sender ID) above, such that the COSE Key specified in the 'cnf' claim is its equivalent COSE Key.

  If this is not the case, the RS MUST request the Client's authentication credential to the Group Manager of the OSCORE group as described in Section 10 of [I-D.ietf-ace-key-groupcomm-oscore], specifying the Client's Sender ID in the OSCORE group, i.e., the value of the 'salt_input' claim. Then, the RS performs the following actions.

    -The RS MUST check whether the Client's authentication credential retrieved from the Group Manager is such that the

COSE Key specified in the 'cnf' claim of the Access Token is its equivalent COSE Key.

-The RS MUST check that the Client's Sender ID provided by the Group Manager together with the Client's authentication credential matches the one retrieved from the 'salt_input' claim of the Access Token.

If any of the checks above fails, the RS MUST consider the Access Token non valid, and MUST respond to the Client with an error response code equivalent to the CoAP code 4.00 (Bad Request).

If the Access Token is valid and further checks on its content are successful, the RS associates the authorization information from the Access Token with the Group OSCORE Security Context.

In particular, the RS associates the authorization information from the Access Token with the 3-tuple (GID, SaltInput, AuthCred), where GID is the Group Identifier of the OSCORE Group, while SaltInput and AuthCred are the Sender ID and the authentication credential that the Client uses in that OSCORE group, respectively.

The RS MUST keep this association up-to-date over time, as the 3-tuple (GID, SaltInput, AuthCred) associated with the Access Token might change. In particular:

  *If the OSCORE group is rekeyed (see Section 3.2 of [I-D.ietf-core-oscore-groupcomm] and Section 20 of [I-D.ietf-ace-key-groupcomm-oscore]), the Group Identifier also changes in the group, and the new one replaces the current 'GID' value in the 3-tuple.

  *If the Client requests and obtains a new OSCORE Sender ID from the Group Manager (see Section 2.5.3.1 of [I-D.ietf-core-oscore-groupcomm] and Section 9 of [I-D.ietf-ace-key-groupcomm-oscore]), the new Sender ID replaces the current 'SaltInput' value in the 3-tuple.

Finally, the RS MUST send a 2.01 (Created) response to the Client, as defined in Section 5.10.1 of [I-D.ietf-ace-oauth-authz].

4.3. Client-RS Secure Communication

When previously joining the OSCORE group, both the Client and RS have already established the related Group OSCORE Security Context to communicate as group members. Therefore, they can simply start to securely communicate using Group OSCORE, without deriving any additional keying material or security association.

4.3.1. Client Side

After having received the 2.01 (Created) response from the RS,
following the POST request to the authz-info endpoint, the Client
starts the communication with the RS, by sending a request protected
with Group OSCORE using the Group OSCORE Security Context [I-D.ietf-
core-oscore-groupcomm].

When communicating with the RS to access the resources as specified
by the authorization information, the Client MUST use the Group
OSCORE Security Context of the OSCORE group, whose GID was specified
in the 'context_id' parameter of the Token request.

4.3.2. Resource Server Side

After successful validation of the Access Token as defined in
Section 4.2 and after having sent the 2.01 (Created) response, the
RS can start to communicate with the Client using Group OSCORE [I-
D.ietf-core-oscore-groupcomm].

When processing an incoming request protected with Group OSCORE, the
RS MUST consider as valid Client's authentication credential only
the one associated to the stored Access Token. As defined in Section
4.5, a possible change of authentication credential requires the
Client to upload to the RS a new Access Token bound to the new
authentication credential.

Additionally, for every incoming request, if Group OSCORE
verification succeeds, the verification of access rights is
performed as described in Section 4.4.

After the expiration of the Access Token related to a Group OSCORE
Security Context, if the Client uses the Group OSCORE Security
Context to send a request for any resource intended for OSCORE group
members and that requires an active Access Token, the RS MUST
respond with a 4.01 (Unauthorized) error message protected with the
Group OSCORE Security Context.

4.4. Access Rights Verification

The RS MUST follow the procedures defined in Section 5.10.2 of [I-
D.ietf-ace-oauth-authz]. If an RS receives a Group OSCORE-protected
request from a Client, the RS processes it according to [I-D.ietf-
core-oscore-groupcomm].

If the Group OSCORE verification succeeds, and the target resource
requires authorization, the RS retrieves the authorization
information from the Access Token associated with the Group OSCORE
Security Context. Then, the RS MUST verify that the action requested
on the resource is authorized.

The response code MUST be 4.01 (Unauthorized) if the RS has no valid Access Token for the Client. If the RS has an Access Token for the Client but no actions are authorized on the target resource, the RS MUST reject the request with a 4.03 (Forbidden). If the RS has an Access Token for the Client but the requested action is not authorized, the RS MUST reject the request with a 4.05 (Method Not Allowed).

4.5. Change of Client's Authentication Credential in the Group

During its membership in the OSCORE group, the client might change the authentication credential it uses in the group. When this happens, the Client uploads the new authentication credential to the Group Manager, as defined in Section 11 of [I-D.ietf-ace-key-groupcomm-oscore].

After that, and in order to continue communicating with the RS, the Client MUST perform the following actions.

1. The Client requests a new Access Token to the AS, as defined in Section 3. In particular, when sending the POST request as defined in Section 3.1, the Client indicates:

   *The current Group Identifier of the OSCORE group, as value of the 'context_id' parameter.

   *The current Sender ID it has in the OSCORE group, as value of the 'salt_input' parameter.

   *The public key of the new authentication credential it uses in the OSCORE group, as value of the 'req_cnf' parameter. In particular, the specified public key is the COSE Key equivalent to the new authentication credential that the Client uses in the OSCORE group.

   *The proof-of-possession (PoP) evidence corresponding to the public key of the new authentication credential, as value of the 'client_cred_verify' or 'client_cred_verify_mac' parameter.

2. After receiving the response from the AS (see Section 3.2), the Client performs the same exchanges with the RS as defined in Section 4.

When receiving the new Access Token, the RS performs the same steps defined in Section 4.2, with the following addition, in case the new Access Token is successfully verified and stored. The RS also deletes the old Access Token, i.e., the one whose associated 3-tuple has the same GID and SaltInput values as in the 3-tuple including

the new authentication credential of the Client and associated with the new Access Token.

5. Secure Communication with the AS

As specified in the ACE framework (see Sections 5.8 and 5.9 of [I-D.ietf-ace-oauth-authz]), the requesting entity (RS and/or Client) and the AS communicate via the /token or /introspection endpoint. The use of CoAP and OSCORE [RFC8613] for this communication is RECOMMENDED in this profile. Other protocols fulfilling the security requirements defined in Sections 5 and 6 of [I-D.ietf-ace-oauth-authz] (such as HTTP and DTLS or TLS) MAY be used instead.

If OSCORE [RFC8613] is used, the requesting entity and the AS are expected to have a pre-established Security Context in place. How this Security Context is established is out of the scope of this profile. Furthermore, the requesting entity and the AS communicate using OSCORE through the /introspection endpoint as specified in Section 5.9 of [I-D.ietf-ace-oauth-authz], and through the /token endpoint as specified in Section 5.8 of [I-D.ietf-ace-oauth-authz].

6. Discarding the Security Context

As members of an OSCORE group, the Client and the RS may independently leave the group or be forced to, e.g., if compromised or suspected so. Upon leaving the OSCORE group, the Client or RS also discards the Group OSCORE Security Context, which may anyway be renewed by the Group Manager through a group rekeying process (see Section 3.2 of [I-D.ietf-core-oscore-groupcomm]).

The Client or RS can acquire a new Group OSCORE Security Context, by re-joining the OSCORE group, e.g., by using the approach defined in [I-D.ietf-ace-key-groupcomm-oscore]. In such a case, the Client SHOULD request a new Access Token and post it to the RS.

7. CBOR Mappings

The new parameters defined in this document MUST be mapped to CBOR types as specified in Figure 6, using the given integer abbreviation for the map key.

```
/------------------------+----------+------------\
| Parameter name         | CBOR Key | Value Type |
|------------------------+----------+------------|
| context_id             | TBD      | bstr       |
| salt_input             | TBD      | bstr       |
| client_cred_verify     | TBD      | bstr       |
| client_cred_verify_mac | TBD      | bstr       |
\------------------------+----------+------------/
```

Figure 6: CBOR mappings for new parameters.

The new claims defined in this document MUST be mapped to CBOR types as specified in Figure 7, using the given integer abbreviation for the map key.

```
/-----------------+----------+------------\
| Claim name      | CBOR Key | Value Type |
|-----------------+----------+------------|
| salt_input      | TBD      | bstr       |
| contextId_input | TBD      | bstr       |
\-----------------+----------+------------/
```

Figure 7: CBOR mappings for new claims.

8. Security Considerations

This document specifies a profile for the Authentication and Authorization for Constrained Environments (ACE) framework [I-D.ietf-ace-oauth-authz]. Thus, the general security considerations from the ACE framework also apply to this profile.

The proof-of-possession (PoP) key bound to an Access Token is always an asymmetric key, i.e., the public key that the Client uses in the OSCORE group. This means that there is never a same shared secret used as PoP key with possible multiple RSs. Therefore, it is possible and safe for the AS to issue an Access Token whose audience comprises multiple RSs.

In such a case, as per Section 6.1 of [I-D.ietf-ace-oauth-authz], the AS has to ensure the integrity protection of the Access Token by protecting it through an asymmetric signature. In addition, the used audience has to correctly identify all the RSs that are intended recipients of the Access Token. As a particular case, the audience can be the name of the OSCORE group, if the Access Token is intended to all the RSs in that group.

Furthermore, this document inherits the general security considerations about Group OSCORE [I-D.ietf-core-oscore-groupcomm], as to the specific use of Group OSCORE according to this profile.

Group OSCORE is designed to secure point-to-point as well as point-to-multipoint communications, providing a secure binding between a single request and multiple corresponding responses. In particular, Group OSCORE fulfills the same security requirements of OSCORE, for group requests and responses.

Group OSCORE ensures source authentication of messages both in group mode (see Section 8 of [I-D.ietf-core-oscore-groupcomm]) and in pairwise mode (see Section 9 of [I-D.ietf-core-oscore-groupcomm]).

When protecting an outgoing message in group mode, the sender uses its private key to compute a digital signature, which is embedded in the protected message. The group mode can be used to protect messages sent over multicast to multiple recipients, or sent over unicast to one recipient.

When protecting an outgoing message in pairwise mode, the sender uses a pairwise symmetric key, as derived from the asymmetric keys of the two peers exchanging the message. The pairwise mode can be used to protect only messages intended to one recipient.

9. Privacy Considerations

This document specifies a profile for the Authentication and Authorization for Constrained Environments (ACE) framework [I-D.ietf-ace-oauth-authz]. Thus the general privacy considerations from the ACE framework also apply to this profile.

As this profile uses Group OSCORE, the privacy considerations from [I-D.ietf-core-oscore-groupcomm] apply to this document as well.

An unprotected response to an unauthorized request may disclose information about the RS and/or its existing relationship with the Client. It is advisable to include as little information as possible in an unencrypted response. However, since both the Client and the RS share a Group OSCORE Security Context, unauthorized, yet protected requests are followed by protected responses, which can thus include more detailed information.

Although it may be encrypted, the Access Token is sent in the clear to the /authz-info endpoint at the RS. Thus, if the Client uses the same single Access Token from multiple locations with multiple Resource Servers, it can risk being tracked through the Access Token's value.

Note that, even though communications are protected with Group OSCORE, some information might still leak, due to the observable size, source address and destination address of exchanged messages.

10. IANA Considerations

This document has the following actions for IANA.

10.1. ACE Profile Registry

IANA is asked to add the following entry to the "ACE Profile"
registry defined in [Section 8.8](#) of [[I-D.ietf-ace-oauth-authz](#)].

   *Name: coap_group_oscore

   *Description: Profile to secure communications between constrained
    nodes using the Authentication and Authorization for Constrained
    Environments framework, by enabling authentication and fine-
    grained authorization of members of an OSCORE group, that use a
    pre-established Group OSCORE Security Context to communicate with
    Group OSCORE. Optionally, the dual mode defined in [Appendix A](#)
    additionally establishes a pairwise OSCORE Security Context, and
    thus also enables OSCORE communication between two members of the
    OSCORE group.

   *CBOR Value: TBD (value between 1 and 255)

   *Reference: [[this document]]

10.2. OAuth Parameters Registry

IANA is asked to add the following entries to the "OAuth Parameters"
registry.

   *Name: "context_id"

   *Parameter Usage Location: token request

   *Change Controller: IESG

   *Specification Document(s): [Section 3.1.1](#) of [[this document]]

   *Name: "salt_input"

   *Parameter Usage Location: token request

   *Change Controller: IESG

   *Specification Document(s): [Section 3.1.2](#) of [[this document]]

   *Name: "client_cred_verify"

   *Parameter Usage Location: token request

   *Change Controller: IESG

   *Specification Document(s): [Section 3.1.3](#) of [[this document]]

   *Name: "client_cred_verify_mac"

*Parameter Usage Location: token request

*Change Controller: IESG

*Specification Document(s): Section 3.1.4 of [[this document]]

*Name: "client_cred"

*Parameter Usage Location: token request

*Change Controller: IESG

*Specification Document(s): Appendix A.2.1.1 of [[this document]]

10.3. OAuth Parameters CBOR Mappings Registry

IANA is asked to add the following entries to the "OAuth Parameters CBOR Mappings" registry defined in Section 8.10 of [I-D.ietf-ace-oauth-authz].

*Name: "context_id"

*CBOR Key: TBD

*Value Type: bstr

*Reference: Section 3.1.1 of [[this document]]

*Name: "salt_input"

*CBOR Key: TBD

*Value Type: bstr

*Reference: Section 3.1.2 of [[this document]]

*Name: "client_cred_verify"

*CBOR Key: TBD

*Value Type: bstr

*Reference: Section 3.1.3 of [[this document]]

*Name: "client_cred_verify_mac"

*CBOR Key: TBD

*Value Type: bstr

*Reference: Section 3.1.4 of [[this document]]

*Name: "client_cred"

*CBOR Key: TBD

*Value Type: bstr

*Reference: Appendix A.2.1.1 of [[this document]]

10.4. CBOR Web Token Claims Registry

IANA is asked to add the following entries to the "CBOR Web Token Claims" registry.

*Claim Name: "salt_input"

*Claim Description: Client provided salt input

*JWT Claim Name: "N/A"

*Claim Key: TBD

*Claim Value Type(s): bstr

*Change Controller: IESG

*Specification Document(s): Section 3.2.1 of [[this document]]

*Claim Name: "contextId_input"

*Claim Description: Client context id input

*JWT Claim Name: "N/A"

*Claim Key: TBD

*Claim Value Type(s): bstr

*Change Controller: IESG

*Specification Document(s): Section 3.2.2 of [[this document]]

*Claim Name: "client_cred"

*Claim Description: Client Credential

*JWT Claim Name: "N/A"

*Claim Key: TBD

*Claim Value Type(s): map

*Change Controller: IESG

   *Specification Document(s): Appendix A.2.2.2 of [[this document]]

10.5. TLS Exporter Label Registry

IANA is asked to add the following entry to the "TLS Exporter Label"
registry defined in Section 6 of [RFC5705] and updated in Section 12
of [RFC8447].

   *Value: EXPORTER-ACE-Sign-Challenge-Client-AS

   *DTLS-OK: Y

   *Recommended: N

   *Reference: [[this document]] (Section 3.1)

11. References

11.1. Normative References

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Appendix A. Dual Mode (Group OSCORE & OSCORE)

This appendix defines the dual mode of this profile, which allows using both OSCORE [RFC8613] and Group OSCORE [I-D.ietf-core-oscore-groupcomm] as security protocols, by still relying on a single Access Token.

That is, the dual mode of this profile specifies how a Client uses CoAP [RFC7252] to communicate to a single Resource Server, or CoAP over IP multicast [I-D.ietf-core-groupcomm-bis] to communicate to multiple Resource Servers that are members of a group and share a common set of resources.

In particular, the dual mode of this profile uses two complementary security protocols to provide secure communication between the Client and the Resource Server(s). That is, it defines the use of either OSCORE or Group OSCORE to protect unicast requests addressed to a single Resource Server, as well as possible responses. Additionally, it defines the use of Group OSCORE to protect multicast requests sent to a group of Resource Servers, as well as possible individual responses. Like in the main mode of this profile, the Client and the Resource Servers need to have already joined the same OSCORE group, for instance by using the approach defined in [I-D.ietf-ace-key-groupcomm-oscore], which is also based on ACE.

The Client proves its access to be authorized to the Resource Server by using an Access Token, which is bound to a key (the proof-of-possession key). This profile mode uses OSCORE to achieve proof of possession, and OSCORE or Group OSCORE to achieve server authentication.

Unlike in the main mode of this profile, where a public key is used as pop-key, this dual mode uses OSCORE-related, symmetric keying material as pop-key instead. Furthermore, this dual mode provides proof of Client's membership to the correct OSCORE group, by securely binding the pre-established Group OSCORE Security Context to the pairwise OSCORE Security Context newly established between the Client and the Resource Server.

In addition to the terminology used for the main mode of this profile, the rest of this appendix refers also to "pairwise OSCORE Security Context" as to an OSCORE Security Context established

between only one Client and one Resource Server, and used to communicate with OSCORE [RFC8613].

A.1. Protocol Overview

This section provides an overview on how to use the ACE framework for authentication and authorization [I-D.ietf-ace-oauth-authz] to secure communications between a Client and a (set of) Resource Server(s) using OSCORE [RFC8613] and/or Group OSCORE [I-D.ietf-core-oscore-groupcomm].

Just as for main mode of this profile overviewed in Section 2, the process for joining the OSCORE group through the respective Group Manager as defined in [I-D.ietf-ace-key-groupcomm-oscore] must take place before the process described in the rest of this section, and is out of the scope of this profile.

An overview of the protocol flow for the dual mode of this profile is shown in Figure 8. In the figure, it is assumed that both RS1 and RS2 are associated with the same AS. It is also assumed that C, RS1 and RS2 have previously joined an OSCORE group with Group Identifier (gid) "abcd0000", and got assigned Sender ID (sid) "0", "1" and "2" in the group, respectively. The names of messages coincide with those of [I-D.ietf-ace-oauth-authz] when applicable.

```
C                         RS1          RS2                        AS
| [--- Resource Request -->] |           |                         |
|                          |            |                         |
| [<----  AS Request ------] |           |                         |
|        Creation Hints      |           |                         |
|                          |            |                         |
|-------- POST /token --------------------------------------------->|
| (aud: RS1, sid: 0, gid: abcd0000, ...)  |                         |
|                          |            |                         |
|<----------------------------- Access Token + RS Information -----|
|                          | (aud: RS1, sid: 0, gid: abcd0000, ...) |
|                          |            |                         |
|---- POST /authz-info ----->|           |                         |
|  (access_token, N1, ID1)  |           |                         |
|                          |            |                         |
|<-- 2.01 Created (N2, ID2) -|           |                         |
|                          |            |                         |
/Pairwise            /Pairwise           |                         |
 Security Context     Security Context   |                         |
 Derivation/          Derivation/        |                         |
|                          |            |                         |
|-------- POST /token --------------------------------------------->|
| (aud: RS2, sid: 0, gid: abcd0000, ...)  |                         |
|                          |            |                         |
|<----------------------------- Access Token + RS Information -----|
|                          | (aud: RS2, sid: 0, gid: abcd0000, ...) |
|                          |            |                         |
|---- POST /authz-info ------------------>|                         |
|  (access_token, N1', ID1') |           |                         |
|                          |            |                         |
|<-- 2.01 Created (N2', ID2')-------------|                         |
|                          |            |                         |
/Pairwise Security          |  /Pairwise Security                 |
 Context Derivation/        |   Context Derivation/               |
|                          |            |                         |
|----- OSCORE Request ------>|           |                         |
|        (kid: ID2)          |           |                         |
|                          |            |                         |
|                          |            |                         |
|                          |            |                         |
|                          |            |                         |
|              /Proof-of-possession;     |                         |
|                 Pairwise Security      |                         |
|                 Context storage/       |                         |
|                          |            |                         |
|<---- OSCORE Response ------|           |                         |
|                          |            |                         |
/Proof-of-possession;       |           |                         |
 Pairwise Security          |           |                         |
```

```
 Context storage/                |              |                              |
|                                |              |                              |
/Mutual authentication           |              |                              |
 between C and RS1               |              |                              |
 (as OSCORE peers)/              |              |                              |
|                                |              |                              |
|                                |              |                              |
|- Group OSCORE Request -+-->|    |              |                              |
| (kid: 0, gid: abcd0000) \-------------->|      |                              |
|                                |              |                              |
|<-- Group OSCORE Response --|    |              |                              |
|         (kid: 1)           |    |              |                              |
|                                |              |                              |
/Mutual authentication           |              |                              |
 between C and RS1               |              |                              |
 (as group members)/             |              |                              |
|                                |              |                              |
|<-- Group OSCORE Response ---------------|      |                              |
|         (kid: 2)           |    |              |                              |
|                                |              |                              |
/Mutual authentication           |              |                              |
 between C and RS2               |              |                              |
 (as group members)/             |              |                              |
|                                |              |                              |
|             ...            |    |              |                              |
```

Figure 8: Protocol Overview.

A.1.1. Pre-Conditions

The same pre-conditions for the main mode of this profile (see Section 2.1) hold for the dual mode described in this appendix.

A.1.2. Access Token Posting

After having retrieved the Access Token from the AS, the Client generates a nonce N1 and an identifier ID1 unique in the sets of its own Recipient IDs from its pairwise OSCORE Security Contexts. The client then posts both the Access Token, N1 and its chosen ID to the RS, using the /authz-info endpoint and mechanisms specified in Section 5.10 of [I-D.ietf-ace-oauth-authz] and Content-Format = application/ace+cbor.

When using the dual mode of this profile, the communication with the authz-info endpoint is not protected, except for update of access rights. Note that, when using the dual mode, this request can alternatively be protected with Group OSCORE, using the Group OSCORE Security Context paired with the pairwise OSCORE Security Context originally established with the first Access Token posting.

If the Access Token is valid, the RS replies to this POST request with a 2.01 (Created) response with Content-Format = application/ace+cbor, which in a CBOR map contains a nonce N2 and an identifier ID2 unique in the sets of its own Recipient IDs from its pairwise OSCORE Security Contexts.

A.1.3. Setup of the Pairwise OSCORE Security Context

After sending the 2.01 (Created) response, the RS sets the ID Context of the pairwise OSCORE Security Context (see Section 3 of [RFC8613]) to the Group Identifier of the OSCORE group specified in the Access Token, concatenated with N1, concatenated with N2, concatenated with the value in the contextId parameter of the OSCORE_Input_Material provided in the 'cnf' claim of the Access Token.

Then, the RS derives the complete pairwise OSCORE Security Context associated with the received Access Token, following Section 3.2 of [RFC8613]. In practice, the RS maintains a collection of Security Contexts with associated authorization information, for all the clients that it is currently communicating with, and the authorization information is a policy used as input when processing requests from those clients.

During the derivation process, the RS uses: the ID Context above; the exchanged nonces N1 and N2; the identifier ID1 received from the

Client, set as its own OSCORE Sender ID; the identifier ID2 provided
to the Client, set as its Recipient ID for the Client; and the
parameters in the Access Token. The derivation process uses also the
Master Secret of the OSCORE group, that the RS knows as a group
member, as well as the Sender ID of the Client in the OSCORE group,
which is specified in the Access Token. This ensures that the
pairwise OSCORE Security Context is securely bound to the Group
OSCORE Security Context of the OSCORE group.

Finally, the RS stores the association between i) the authorization
information from the Access Token; and ii) the Group Identifier of
the OSCORE group together with the Sender ID and the authentication
credential of the Client in that group.

After having received the nonce N2, the Client sets the ID Context
in its pairwise OSCORE Security Context (see Section 3 of [RFC8613])
to the Group Identifier of the OSCORE group, concatenated with N1,
concatenated with N2, concatenated with the value in the contextId
parameter of the OSCORE_Input_Material provided in the 'cnf'
parameter of the Access Token response from the AS. Then, the Client
derives the complete pairwise OSCORE Security Context, following
Section 3.2 of [RFC8613].

During the derivation process, the Client uses: the ID Context
above, the exchanged nonces N1 and N2; the identifier ID1 provided
to the RS, set as its own Recipient ID for the RS; the identifier
ID2 received from the RS, set as its own OSCORE Sender ID; and the
parameters received from the AS. The derivation process uses also
the Master Secret of the OSCORE group, that the Client knows as a
group member, as well as its own Sender ID in the OSCORE group.

When the Client communicates with the RS using the pairwise OSCORE
Security Context, the RS achieves proof-of-possession of the
credentials bound to the Access Token. Also, the RS verifies that
the Client is a legitimate member of the OSCORE group.

A.1.4. Secure Communication

Other than starting the communication with the RS using Group OSCORE
as described in Section 4.3, the Client can send to the RS a request
protected with OSCORE, using the pairwise OSCORE Security Context.

If the request is successfully verified, then the RS stores the
pairwise OSCORE Security Context, and uses it to protect the
possible response, as well as further communications with the
Client, until the Access Token is deleted, due to, for example,
expiration. This pairwise OSCORE Security Context is discarded when
an Access Token (whether the same or different) is used to
successfully derive a new pairwise OSCORE Security Context.

As discussed in [Section 7](#) of [[I-D.ietf-ace-oscore-profile](#)], the use
of random nonces N1 and N2 during the exchange between the Client
and the RS prevents the reuse of an Authenticated Encryption with
Associated Data (AEAD) nonce/key pair for two different messages.
Reuse might otherwise occur when Client and RS derive a new pairwise
OSCORE Security Context from an existing (non-expired) Access Token,
e.g., in case of reboot of either the Client or the RS, and might
lead to loss of both confidentiality and integrity.

Additionally, just as per the main mode of this profile (see [Section
4.3](#)), the Client and RS can also securely communicate by protecting
messages with Group OSCORE, using the Group OSCORE Security Context
already established upon joining the OSCORE group.

A.2. Client-AS Communication

This section details the Access Token POST Request that the Client
sends to the /token endpoint of the AS, as well as the related
Access Token response.

[Section 3.2](#) of [[RFC8613](#)] defines how to derive a pairwise OSCORE
Security Context based on a shared Master Secret and a set of other
parameters, established between the OSCORE client and server, which
the client receives from the AS in this exchange.

The proof-of-possession key (pop-key) received from the AS in this
exchange MUST be used to build the Master Secret in OSCORE (see
[Appendix A.3.3](#) and [Appendix A.3.4](#)).

A.2.1. C-to-AS: POST to Token Endpoint

The Client-to-AS request is specified in [Section 5.8.1](#) of [[I-D.ietf-
ace-oauth-authz](#)]. The Client MUST send this POST request to the /
token endpoint over a secure channel that guarantees authentication,
message integrity and confidentiality.

The POST request is formatted as the analogous Client-to-AS request
in the main mode of this profile (see [Section 3.1](#)), with the
following modifications.

  *The parameter 'req_cnf' MUST NOT be included in the payload.

  *The parameter 'client_cred', defined in [Appendix A.2.1.1](#) of this
   document, MUST be included in the payload. This parameter
   specifies the public key that the Client uses in the OSCORE
   group, whose identifier is indicated in the 'context_id'
   parameter. In particular, the specified public key is the COSE
   Key equivalent to the authentication credential that the Client
   uses in the OSCORE group.

*The proof-of-possession (PoP) evidence included in the
 'client_cred_verify' or 'client_cred_verify_mac' parameter is
 computed by using the Client's private key associated with the
 public key in the 'client_cred' parameter above.

An example of such a request is shown in Figure 9.

```
Header: POST (Code=0.02)
Uri-Host: "as.example.com"
Uri-Path: "token"
Content-Format: "application/ace+cbor"
Payload:
{
  "audience" : "tempSensor4711",
  "scope" : "read",
  "context_id" : h'abcd0000',
  "salt_input" : h'00',
  "client_cred" : {
    "COSE_Key" : {
      "kty" : EC2,
      "crv" : P-256,
      "x" : h'd7cc072de2205bdc1537a543d53c60a6acb62eccd890c7fa
              27c9e354089bbe13',
      "y" : h'f95e1d4b851a2cc80fff87d8e23f22afb725d535e515d020
              731e79a3b4e47120'
    }
  },
  "client_cred_verify" : h'...'
  (signature content omitted for brevity),
}
```

Figure 9: Example C-to-AS POST /token request for an Access Token bound
                        to a symmetric key.

Later on, the Client may want to update its current access rights,
without changing the existing pairwise OSCORE Security Context with
the RS. In this case, the Client MUST include in its POST request to
the /token endpoint a 'req_cnf' parameter, defined in Section 3.1 of
[I-D.ietf-ace-oauth-params], which MUST include a 'kid' field, as
defined in Section 3.1 of [RFC8747]. The 'kid' field has as value a
CBOR byte string containing the OSCORE_Input_Material Identifier
(assigned as discussed in Appendix A.2.2).

This identifier, together with other information such as audience,
can be used by the AS to determine the shared secret bound to the
proof-of-possession Access Token and therefore MUST identify a
symmetric key that was previously generated by the AS as a shared
secret for the communication between the Client and the RS. The AS

MUST verify that the received value identifies a proof-of-possession key that has previously been issued to the requesting Client. If that is not the case, the Client-to-AS request MUST be declined with the error code "invalid_request" as defined in Section 5.8.3 of [I-D.ietf-ace-oauth-authz].

This POST request for updating the access rights of an Access Token SHOULD NOT include the parameters 'salt_input', 'context_id', 'client_cred' and 'client_cred_verify'. An exception is the case defined in Appendix A.3.6, where the Client, following a change of authentication credential in the OSCORE group, requests a new Access Token associated with the public key of the new authentication credential, while still without changing the existing pairwise OSCORE Security Context with the RS.

An example of such a request is shown in Figure 10.

```
Header: POST (Code=0.02)
Uri-Host: "as.example.com"
Uri-Path: "token"
Content-Format: "application/ace+cbor"
Payload:
{
  "audience" : "tempSensor4711",
  "scope" : "read",
  "req_cnf" : {
    "kid" : h'01'
  }
}
```

Figure 10: Example C-to-AS POST /token request for updating rights to an Access Token bound to a symmetric key.

A.2.1.1. 'client_cred' Parameter

The 'client_cred' parameter is an OPTIONAL parameter of the Access Token request message defined in Section 5.8.1. of [I-D.ietf-ace-oauth-authz]. This parameter provides an asymmetric key that the Client wishes to use as its own public key, but which is not used as proof-of-possession key.

This parameter follows the syntax of the 'cnf' claim from Section 3.1 of [RFC8747] when including Value Type "COSE_Key" (1) and specifying an asymmetric key. Alternative Value Types defined in future specifications are fine to consider if indicating a non-encrypted asymmetric key.

A.2.2. AS-to-C: Access Token

After having verified the POST request to the /token endpoint and
that the Client is authorized to obtain an Access Token
corresponding to its Access Token request, the AS MUST verify the
proof-of-possession (PoP) evidence.

In particular, the AS proceeds as defined in Section 3.2, with the
difference that it uses the public key specified in the
'client_cred' parameter as public key of the Client.

If both the 'client_cred_verify' and 'client_cred_verify_mac'
parameters are present, or if the verification of the PoP evidence
fails, the AS considers the Client request invalid. The AS does not
perform this operation when asked to update a previously released
Access Token.

If all verifications are successful, the AS responds as defined in
Section 5.8.2 of [I-D.ietf-ace-oauth-authz]. If the Client request
was invalid, or not authorized, the AS returns an error response as
described in Section 5.8.3 of [I-D.ietf-ace-oauth-authz].

The AS can signal that the use of OSCORE and Group OSCORE is
REQUIRED for a specific Access Token by including the "ace_profile"
parameter with the value "coap_group_oscore" in the Access Token
response. This means that the Client MUST use OSCORE and/or Group
OSCORE towards all the Resource Servers for which this Access Token
is valid.

In particular, the Client MUST follow Appendix A.3.3 to derive the
pairwise OSCORE Security Context to use for communications with the
RS. Instead, the Client has already established the related Group
OSCORE Security Context to communicate with members of the OSCORE
group, upon previously joining that group.

Usually, it is assumed that constrained devices will be pre-
configured with the necessary profile, so that this kind of profile
signaling can be omitted.

In contrast with the main mode of this profile, the Access Token
response to the Client is analogous to the one in the OSCORE profile
of ACE, as described in Section 3.2 of [I-D.ietf-ace-oscore-
profile]. In particular, the AS provides an OSCORE_Input_Material
object, which is defined in Section 3.2.1 of [I-D.ietf-ace-oscore-
profile] and included in the 'cnf' parameter (see Section 3.2 of [I-
D.ietf-ace-oauth-params]) of the Access Token response.

The AS MUST send different OSCORE_Input_Material (and therefore
different Access Tokens) to different authorized clients, in order
for the RS to differentiate between clients.

In the issued Access Token, the AS MUST include as metadata the same information as defined in the main mode of this profile (see Section 3.2) with the following modifications.

   *The public key that the client uses in the OSCORE group and specified in the 'client_cred' parameter of the Token request (see Appendix A.2.1) MUST also be included in the Access Token.

   If the Access Token is a CWT, the AS MUST include it in the 'client_cred' claim of the Access Token, defined in Appendix A.2.2.2 of this document. In particular, the 'client_cred' claim includes the same COSE Key specified in the 'client_cred' parameter of the Token Request, i.e., the COSE Key equivalent to the authentication credential that the Client uses in the OSCORE group.

   *The OSCORE_Input_Material specified in the 'cnf' parameter of the Access Token response MUST also be included in the Access Token. If the Access Token is a CWT, the same OSCORE_Input_Material included in the 'cnf' parameter of the Access Token response MUST be included in the 'osc' field of the 'cnf' claim of the Access Token (see Section 3.2 of [I-D.ietf-ace-oscore-profile]).

Figure 11 shows an example of such an AS response. The access token has been truncated for readability.

```
Header: Created (Code=2.01)
Content-Type: "application/ace+cbor"
Payload:
{
   "access_token" : h'8343a1010aa2044c53 ...'
    (remainder of CWT omitted for brevity),
   "ace_profile" : "coap_group_oscore",
   "expires_in" : 3600,
   "cnf" : {
     "osc" : {
       "alg" : AES-CCM-16-64-128,
       "id"  : h'01',
       "ms" : h'f9af838368e353e78888e1426bd94e6f',
       "salt" : h'1122',
       "contextId" : h'99'
     }
   }
}
```

Figure 11: Example AS-to-C Access Token response with the Group OSCORE profile.

Figure 12 shows an example CWT, containing the necessary OSCORE
parameters in the 'cnf' claim.

```
  {
    "aud" : "tempSensorInLivingRoom",
    "iat" : 1360189224,
    "exp" : 1360289224,
    "scope" :  "temperature_g firmware_p",
    "cnf" : {
      "osc" : {
        "alg" : AES-CCM-16-64-128,
        "id"  : h'01',
        "ms" : h'f9af838368e353e78888e1426bd94e6f',
        "salt" : h'1122',
        "contextId" : h'99'
    },
    "salt_input" : h'00',
    "contextId_input" : h'abcd0000',
    "client_cred" : {
      "COSE_Key" : {
        "kty" : EC2,
        "crv" : P-256,
        "x" : h'd7cc072de2205bdc1537a543d53c60a6acb62eccd890c7fa
               27c9e354089bbe13',
        "y" : h'f95e1d4b851a2cc80fff87d8e23f22afb725d535e515d020
               731e79a3b4e47120'
      }
    }
  }
```

Figure 12: Example CWT with OSCORE parameters.

The same CWT as in Figure 12 and encoded in CBOR is shown in Figure
13, using the value abbreviations defined in [I-D.ietf-ace-oauth-
authz] and [RFC8747].

NOTE: it should be checked (and in case fixed) that the values used
below (which are not yet registered) are the final values registered
in IANA.

```
A8                                          # map(8)
   03                                       # unsigned(3)
   76                                       # text(22)
      74656D7053656E736F72496E4C6976696E67526F6F6D
   06                                       # unsigned(6)
   1A 5112D728                              # unsigned(1360189224)
   04                                       # unsigned(4)
   1A 51145DC8                              # unsigned(1360289224)
   09                                       # unsigned(9)
   78 18                                    # text(24)
      74656D70657261747572655F67206669726D776172655F70
   08                                       # unsigned(8)
   A1                                       # map(1)
      04                                    # unsigned(4)
      A5                                    # map(5)
         04                                 # unsigned(4)
         0A                                 # unsigned(10)
         00                                 # unsigned(0)
         41                                 # bytes(1)
            01                              # "\x01"
         02                                 # unsigned(2)
         50                                 # bytes(16)
            F9AF838368E353E78888E1426BD94E6F
         05                                 # unsigned(5)
         42                                 # bytes(2)
            1122                            # "\x11\""
         06                                 # unsigned(6)
         41                                 # bytes(1)
            99                              # "\x99"
   18 3C                                    # unsigned(60)
   41                                       # bytes(1)
      00
   18 3D                                    # unsigned(61)
   44                                       # bytes(4)
      ABCD0000
   18 3E                                    # unsigned(62)
   A1                                       # map(1)
      01                                    # unsigned(1)
      A4                                    # map(4)
         01                                 # unsigned(1)
         02                                 # unsigned(2)
         20                                 # negative(0)
         01                                 # unsigned(1)
         21                                 # negative(1)
         58 20                              # bytes(32)
            D7CC072DE2205BDC1537A543D53C60A6ACB62ECCD890C7FA27C9
            E354089BBE13
         22                                 # negative(2)
         58 20                              # bytes(32)
```

F95E1D4B851A2CC80FFF87D8E23F22AFB725D535E515D020731E
79A3B4E47120

Figure 13: Example CWT with OSCORE parameters.

   If the Client has requested an update to its access rights using the
   same pairwise OSCORE Security Context, which is valid and
   authorized, the AS MUST omit the 'cnf' parameter in the response to
   the client.

   Instead, the updated Access Token conveyed in the AS-to-C response
   MUST include a 'cnf' claim specifying a 'kid' field, as defined in
   Section 3.1 of [RFC8747]. The response from the AS MUST carry the
   OSCORE Input Material identifier in the 'kid' field within the 'cnf'
   claim of the Access Token. That is, the 'kid' field is a CBOR byte
   string, with value the same value of the 'kid' field of the
   'req_cnf' parameter from the C-to-AS request for updating rights to
   the Access Token (see Figure 10). This information needs to be
   included in the Access Token, in order for the RS to identify the
   previously generated pairwise OSCORE Security Context.

   Figure 14 shows an example of such an AS response. The Access Token
   has been truncated for readability.

     Header: Created (Code=2.01)
     Content-Type: "application/ace+cbor"
     Payload:
     {
       "access_token" : h'8343a1010aa2044c53 ...'
        (remainder of CWT omitted for brevity),
       "profile" : "coap_group_oscore",
       "expires_in" : 3600
     }

   Figure 14: Example AS-to-C Access Token response with the Group OSCORE
                  profile, for update of access rights.

   Figure 15 shows an example CWT, containing the necessary OSCORE
   parameters in the 'cnf' claim for update of access rights.

     {
       "aud" : "tempSensorInLivingRoom",
       "iat" : 1360189224,
       "exp" : 1360289224,
       "scope" : "temperature_h",
       "cnf" : {
         "kid" : h'01'
       }
     }

Figure 15: Example CWT with OSCORE parameters for update of access
rights.

A.2.2.1. Public Key Hash as Client Credential

As a possible optimization to limit the size of the Access Token,
the AS may specify as value of the 'client_cred' claim simply the
hash of the Client's public key, i.e., the hash of the COSE Key K
equivalent to the authentication credential that the Client uses in
the OSCORE group.

The specifically used hash-function MUST be collision-resistant on
byte-strings, and MUST be selected from the "Named Information Hash
Algorithm" Registry defined in Section 9.4 of [RFC6920].

In particular, the AS provides the Client with an Access Token as
defined in Appendix A.2.2, with the following differences.

The AS prepares INPUT_HASH as follows, with | denoting byte string
concatenation.

  *If the equivalent COSE Key K has COSE Key Type OKP, INPUT_HASH =
   i, where 'i' is the x-parameter of the COSE_Key specified in the
   'client_cred' parameter of the Token request, encoded as a CBOR
   byte string.

  *If the equivalent COSE Key K has COSE Key Type EC2, INPUT_HASH =
   (i_1 | i_2), where 'i_1' and 'i_2' are the x-parameter and y-
   parameter of the COSE_Key specified in the 'client_cred'
   parameter of the Token request, respectively, each encoded as a
   CBOR byte string.

  *If the equivalent COSE Key K has COSE Key Type RSA, INPUT_HASH =
   (i_1 | i_2), where 'i_1' and 'i_2' are the n-parameter and e-
   parameter of the COSE_Key specified in the 'client_cred'
   parameter of the Token request, respectively, each encoded as a
   CBOR byte string.

Then, the AS hashes INPUT_HASH according to the procedure described
in [RFC6920], with the output OUTPUT_HASH in binary format, as
described in Section 6 of [RFC6920].

Finally, the AS includes a single entry within the 'client_cred'
claim of the Access Token. This entry has label "kid" (3) defined in
Section 3.1 of [RFC8747], and value a CBOR byte string wrapping
OUTPUT_HASH.

Upon receiving the Access Token, the RS processes it according to
Appendix A.3.2, with the following differences.

The RS considers: the content of the 'contextId_input' claim as the GID of the OSCORE group; the content of the 'salt_input' claim as the Sender ID that the Client has in the group; and the content of the 'client_cred' claim as the hash RECEIVED_HASH of a COSE Key equivalent to the authentication credential that the Client uses in the group.

The RS MUST check whether it already stores an authentication credential associated with the pair (GID, Sender ID) above, such that the recomputed hash NEW_HASH of its equivalent COSE Key matches with RECEIVED_HASH from the 'client_cred' claim.

If this is not the case, the RS MUST request the Client's authentication credential to the Group Manager of the OSCORE group as described in Section 10 of [I-D.ietf-ace-key-groupcomm-oscore], specifying the Client's Sender ID in the OSCORE group, i.e., the value of the 'salt_input' claim. Then, the RS performs the following actions.

  *The RS MUST check whether RECEIVED_HASH matches with the recomputed hash NEW_HASH of a COSE Key equivalent to the Client's authentication credential retrieved from the Group Manager.

  *The RS MUST check that the Client's Sender ID provided by the Group Manager together with the Client's authentication credential matches the one retrieved from the 'salt_input' claim of the Access Token.

The RS MUST calculate NEW_HASH using the same method used by the AS described above, and using the same hash function. The hash function to use can be determined from the information conveyed in the 'client_cred' claim, as the procedure described in [RFC6920] also encodes the used hash function as metadata of the hash value.

A.2.2.2. Client Credential Claim

The 'client_cred' claim provides an asymmetric key that the Client owning the Access Token wishes to use as its own public key, but which is not used as proof-of-possession key.

This parameter follows the syntax of the 'cnf' claim from Section 3.1 of [RFC8747] when including Value Type "COSE_Key" (1) and specifying an asymmetric key. Alternative Value Types defined in future specifications are fine to consider, if indicating a non-encrypted asymmetric key or full-fledged autentication credential.

A.3. Client-RS Communication

This section details the POST request and response to the /authz-info endpoint between the Client and the RS. With respect to the

exchanged messages and their content, the Client and the RS perform as defined in the OSCORE profile of ACE (see Section 4 of [I-D.ietf-ace-oscore-profile]).

That is, the Client generates a nonce N1 and posts it to the RS, together with: an identifier ID1 unique in the sets of its own Recipient IDs from its pairwise OSCORE Security Contexts; and the Access Token that includes the material provisioned by the AS.

Then, the RS generates a nonce N2, and an identifier ID2 unique in the sets of its own Recipient IDs from its pairwise OSCORE Security Contexts. After that, the RS derives a pairwise OSCORE Security Context as described in Section 3.2 of [RFC8613]. In particular, it uses the two exchanged nonces and the two identifiers established with the Client, as well as two shared secrets together with additional pieces of information specified in the Access Token.

Both the client and the RS generate the pairwise OSCORE Security Context using the pop-key as part of the OSCORE Master Secret. In addition, the derivation of the pairwise OSCORE Security Context takes as input also information related to the OSCORE group, i.e., the Master Secret and Group Identifier of the group, as well as the Sender ID of the Client in the group. Hence, the derived pairwise OSCORE Security Context is also securely bound to the Group OSCORE Security Context of the OSCORE Group. Thus, the proof-of-possession required to bind the Access Token to the Client occurs after the first OSCORE message exchange.

Therefore, an attacker using a stolen Access Token cannot generate a valid pairwise OSCORE Security Context and thus cannot prove possession of the pop-key. Also, if a Client legitimately owns an Access Token but has not joined the OSCORE group, that Client cannot generate a valid pairwise OSCORE Security Context either, since it lacks the Master Secret used in the OSCORE group.

Besides, just as in the main mode (see Section 4), the RS is able to verify whether the Client has indeed the claimed Sender ID and authentication credential in the OSCORE group.

A.3.1. C-to-RS POST to authz-info Endpoint

The Client MUST generate a nonce N1, an OSCORE Recipient ID (ID1), and post them to the /authz-info endpoint of the RS together with the Access Token, as defined in the OSCORE profile of ACE (see Section 4.1 of [I-D.ietf-ace-oscore-profile]).

The same recommendations, considerations and behaviors defined in Section 4.1 of [I-D.ietf-ace-oscore-profile] hold.

If the Client has already posted a valid Access Token, has already established a pairwise OSCORE Security Context with the RS, and wants to update its access rights, the Client can do so by posting the new Access Token (retrieved from the AS and specifying the updated set of access rights) to the /authz-info endpoint.

The Client MUST protect the request using either the pairwise OSCORE Security Context established during the first Access Token exchange, or the Group OSCORE Security Context associated with that pairwise OSCORE Security Context.

In either case, the Client MUST only send the Access Token in the payload, i.e., no nonce or identifier are sent. After proper verification (see Section 4.2 of [I-D.ietf-ace-oscore-profile]), the new Access Token will supersede the old one at the RS, by replacing the corresponding authorization information. At the same time, the RS will maintain the same pairwise OSCORE Security Context and Group OSCORE Security Context, as now both associated with the new Access Token.

A.3.2. RS-to-C: 2.01 (Created)

The RS MUST verify the validity of the Access Token as defined in Section 4.2, with the following modifications.

*If the POST request to /authz-info is not protected, the RS checks that the 'cnf' claim is included in the Access Token and that it contains an OSCORE_Input_Material object. Also, the RS checks that the 'salt_input', 'client_cred' and 'contextId_input' claims are included in the Access Token.

*If the POST request to /authz-info is protected with the pairwise OSCORE Security Context shared with the Client or with the Group OSCORE Security Context of the OSCORE group, i.e., the Client is requesting an update of access rights, the RS checks that the 'cnf' claim is included in the Access Token and that it contains only the 'kid' field.

*If the 'salt_input', 'client_cred' and 'contextId_input' claims are included in the Access Token, the RS extracts the content of 'client_cred'. Then, the RS considers that content as the COSE Key equivalent to the authentication credential that the Client uses in the group, whose GID is specified in the 'contextId_input' claim. The RS can compare this public key with the actual COSE Key equivalent to the authentication credential of the claimed Client, retrieved from its local storage or from the Group Manager (see Section 4.2).

If any of the checks fails, the RS MUST consider the Access Token non valid, and MUST respond to the Client with an error response code equivalent to the CoAP code 4.00 (Bad Request).

If the Access Token is valid and further checks on its content are successful, the RS proceeds as follows.

In case the POST request to /authz-info was not protected, the RS MUST generate a nonce N2, an OSCORE Recipient ID (ID2), and include them in the 2.01 (Created) response to the Client, as defined in the OSCORE profile of ACE (see Section 4.2 of [I-D.ietf-ace-oscore-profile]).

Instead, in case the POST request to /authz-info was protected, the RS MUST ignore any nonce and identifiers in the request, if any was sent. Then, the RS MUST check that the 'kid' field of the 'cnf' claim in the new Access Token matches the identifier of the OSCORE Input Material of a pairwise OSCORE Security Context such that:

  *The pairwise OSCORE Security Context was used to protect the
   request, if this was protected with OSCORE; or

  *The pairwise OSCORE Security Context is bound to the Group OSCORE
   Security Context used to protect the request, if this was
   protected with Group OSCORE.

If either verification is successful, the new Access Token supersedes the old one at the RS. Besides, the RS associates the new Access Token to the same pairwise OSCORE Security Context identified above, as also bound to a Group OSCORE Security Context. The RS MUST respond with a 2.01 (Created) response with no payload, protected with the same Security Context used to protect the request. In particular, no new pairwise OSCORE Security Context is established between the Client and the RS. If any verification fails, the RS MUST respond with a 4.01 (Unauthorized) error response.

Further recommendations, considerations and behaviors defined in Section 4.2 of [I-D.ietf-ace-oscore-profile] hold for this document.

A.3.3. OSCORE Setup - Client Side

Once having received the 2.01 (Created) response from the RS, following an unprotected POST request to the /authz-info endpoint, the Client MUST extract the nonce N2 from the 'nonce2' parameter, and the Client identifier from the 'ace_server_recipientid' parameter in the CBOR map of the response payload. Note that this identifier is used by C as Sender ID in the pairwise OSCORE Security Context to be established with the RS, and is different as well as unrelated to the Sender ID of C in the OSCORE group.

Then, the Client performs the following actions, in order to set up and fully derive the pairwise OSCORE Security Context for communicating with the RS.

*The Client MUST set the ID Context of the pairwise OSCORE Security Context as the concatenation of: i) GID, i.e., the Group Identifier of the OSCORE group, as specified by the Client in the 'context_id' parameter of the Client-to-AS request; ii) the nonce N1; iii) the nonce N2; and iv) CID, i.e., the value in the contextId parameter of the OSCORE_Input_Material provided in the 'cnf' parameter of the Access Token response from the AS. The concatenation occurs in this order: ID Context = GID | N1 | N2 | CID, where | denotes byte string concatenation.

*The Client MUST set the updated Master Salt of the pairwise OSCORE Security Context as the concatenation of SaltInput, MSalt, the nonce N1, the nonce N2 and GMSalt, where: i) SaltInput is the Sender ID that the Client has in the OSCORE group, which is known to the Client as a member of the OSCORE group; ii) MSalt is the (optional) Master Salt in the pairwise OSCORE Security Context (received from the AS in the Token); and iii) GMSalt is the (optional) Master Salt in the Group OSCORE Security Context, which is known to the Client as a member of the OSCORE group. The concatenation occurs in this order: Master Salt = SaltInput | MSalt | N1 | N2 | GMSalt, where | denotes byte string concatenation. Optional values, if not specified, are not included in the concatenation. The five parameters SaltInput, MSalt, N1, N2 and GMSalt are to be concatenated as encoded CBOR byte strings. An example of Master Salt construction using CBOR encoding is given in Figure 16.

```
SaltInput, MSalt, N1, N2 and GMSalt, in CBOR diagnostic notation:
      SaltInput = h'00'
      MSalt = h'f9af838368e353e78888e1426bd94e6f'
      N1 = h'018a278f7faab55a'
      N2 = h'25a8991cd700ac01'
      GMSalt = h'99'

SaltInput, MSalt, N1, N2 and GMSalt, as CBOR encoded byte strings:
      SaltInput = 0x4100
      MSalt = 0x50f9af838368e353e78888e1426bd94e6f
      N1 = 0x48018a278f7faab55a
      N2 = 0x4825a8991cd700ac01
      GMSalt = 0x4199

Master Salt = 0x41 00
                 50 f9af838368e353e78888e1426bd94e6f
                 48 018a278f7faab55a
                 48 25a8991cd700ac01
                 41 99
```

Figure 16: Example of Master Salt construction using CBOR encoding.

*The Client MUST set the Master Secret of the pairwise OSCORE
 Security Context to the concatenation of MSec and GMSec, where:
 i) MSec is the value of the 'ms' parameter in the
 OSCORE_Input_Material of the 'cnf' parameter, received from the
 AS in Appendix A.2.2; while ii) GMSec is the Master Secret of the
 Group OSCORE Security Context, which is known to the Client as a
 member of the OSCORE group.

*The Client MUST set the Recipient ID as ace_client_recipientid,
 sent as described in Appendix A.3.1.

*The Client MUST set the Sender ID as ace_server_recipientid,
 received as described in Appendix A.3.1.

*The Client MUST set the AEAD Algorithm, ID Context, HKDF, and
 OSCORE Version as indicated in the corresponding parameters
 received from the AS in Appendix A.2.2, if present in the
 OSCORE_Input_Material of the 'cnf' parameter. In case these
 parameters are omitted, the default values SHALL be used as
 described in Sections 3.2 and 5.4 of [RFC8613].

Finally, the client MUST derive the complete pairwise OSCORE
Security Context following Section 3.2.1 of [RFC8613].

From then on, when communicating with the RS to access the resources
as specified by the authorization information, the Client MUST use
the newly established pairwise OSCORE Security Context or the Group

OSCORE Security Context of the OSCORE Group where both the Client
and the RS are members.

If any of the expected parameters is missing (e.g., any of the
mandatory parameters from the AS or the RS), or if
ace_client_recipientid equals ace_server_recipientid (and as a
consequence the Sender and Recipient Keys derived would be equal,
see Section 3.3 of [RFC8613]), then the client MUST stop the
exchange, and MUST NOT derive the pairwise OSCORE Security Context.
The Client MAY restart the exchange, to get the correct security
input material.

The Client can use this pairwise OSCORE Security Context to send
requests to the RS protected with OSCORE. Besides, the Client can
use the Group OSCORE Security Context for protecting unicast
requests to the RS, or multicast requests to the OSCORE group
including also the RS. Mutual authentication as group members is
only achieved after the client has successfully verified the Group
OSCORE protected response from the RS. Similarly, mutual
authentication as OSCORE peers is only achieved after the client has
successfully verified the OSCORE protected response from the RS,
using the pairwise OSCORE Security Context.

Note that the ID Context of the pairwise OSCORE Security Context can
be assigned by the AS, communicated and set in both the RS and
Client after the exchange specified in this profile is executed.
Subsequently, the Client and RS can update their ID Context by
running a mechanism such as the one defined in Appendix B.2 of
[RFC8613] if they both support it and are configured to do so. In
that case, the ID Context in the pairwise OSCORE Security Context
will not match the "contextId" parameter of the corresponding
OSCORE_Input_Material. Running the procedure in Appendix B.2 of
[RFC8613] results in the keying material in the pairwise OSCORE
Security Contexts of the Client and RS being updated. The Client can
achieve the same result by re-posting the Access Token to the
unprotected /authz-info endpoint at the RS, as described in
Section 4.1 of [I-D.ietf-ace-oscore-profile], although without
updating the ID Context.

A.3.4. OSCORE Setup - Resource Server Side

After validation of the Access Token as defined in Appendix A.3.2
and after sending the 2.01 (Created) response to an unprotected POST
request to the /authz-info endpoint, the RS performs the following
actions, in order to set up and fully derive the pairwise OSCORE
Security Context created to communicate with the Client.

   *The RS MUST set the ID Context of the pairwise OSCORE Security
    Context as the concatenation of: i) GID, i.e., the Group

Identifier of the OSCORE group, as specified in the 'contextId'
parameter of the OSCORE_Input_Material, in the 'cnf' claim of the
Access Token received from the Client (see Appendix A.3.1); ii)
the nonce N1; iii) the nonce N2; and iv) CID which is the value
in the contextId parameter of the OSCORE_Input_Material provided
in the 'cnf' claim of the Access Token. The concatenation occurs
in this order: ID Context = GID | N1 | N2 | CID, where | denotes
byte string concatenation.

*The RS MUST set the new Master Salt of the pairwise OSCORE
 Security Context as the concatenation of SaltInput, MSalt, the
 nonce N1, the nonce N2 and GMSalt, where: i) SaltInput is the
 Sender ID that the Client has in the OSCORE group, as specified
 in the 'salt_input' claim included in the Access Token received
 from the Client (see Appendix A.3.1); ii) MSalt is the (optional)
 Master Salt in the pairwise OSCORE Security Context as specified
 in the 'salt' parameter in the OSCORE_Input_Material of the 'cnf'
 claim, included in the Access Token received from the Client; and
 iii) GMSalt is the (optional) Master Salt in the Group OSCORE
 Security Context, which is known to the RS as a member of the
 OSCORE group. The concatenation occurs in this order: Master Salt
 = SaltInput | MSalt | N1 | N2 | GMSalt, where | denotes byte
 string concatenation. Optional values, if not specified, are not
 included in the concatenation. The same considerations for
 building the Master Salt, considering the inputs as encoded CBOR
 byte strings as in Figure 16, hold also for the RS.

*The RS MUST set the Master Secret of the pairwise OSCORE Security
 Context to the concatenation of MSec and GMSec, where: i) MSec is
 the value of the 'ms' parameter in the OSCORE_Input_Material of
 the 'cnf' claim, included in the Access Token received from the
 Client (see Appendix A.3.1); while ii) GMSec is the Master Secret
 of the Group OSCORE Security Context, which is known to the RS as
 a member of the OSCORE group.

*The RS MUST set the Recipient ID as ace_server_recipientid, sent
 as described in Appendix A.3.2.

*The RS MUST set the Sender ID as ace_client_recipientid, received
 as described in Appendix A.3.2.

*The RS MUST set the AEAD Algorithm, ID Context, HKDF, and OSCORE
 Version from the corresponding parameters received from the
 Client in the Access Token (see Appendix A.3.1), if present in
 the OSCORE_Input_Material of the 'cnf' claim. In case these
 parameters are omitted, the default values SHALL be used as
 described in Sections 3.2 and 5.4 of [RFC8613].

Finally, the RS MUST derive the complete pairwise OSCORE Security Context following Section 3.2.1 of [RFC8613].

Once having completed the derivation above, the RS MUST associate the authorization information from the Access Token with the just established pairwise OSCORE Security Context. Furthermore, as defined in Section 4.2, the RS MUST associate the authorization information from the Access Token with the Group OSCORE Security Context.

Then, the RS uses this pairwise OSCORE Security Context to verify requests from and send responses to the Client protected with OSCORE, when this Security Context is used. If OSCORE verification fails, error responses are used, as specified in Section 8 of [RFC8613].

Besides, the RS uses the Group OSCORE Security Context to verify (multicast) requests from and send responses to the Client protected with Group OSCORE. When processing an incoming request protected with Group OSCORE, the RS MUST consider as valid authentication credential of the Client only the authentication credential associated with the stored Access Token. As defined in Appendix A. 3.6, a change of authentication credential in the group requires the Client to upload to the RS a new Access Token, where the 'client_cred' claim specifies a COSE Key equivalent to the new authentication credential that the Client has in the group.

If Group OSCORE verification fails, error responses are used, as specified in Sections 8 and 9 of [I-D.ietf-core-oscore-groupcomm]. Additionally, for every incoming request, if OSCORE or Group OSCORE verification succeeds, the verification of access rights is performed as described in Appendix A.3.5.

After the deletion of the Access Token related to a pairwise OSCORE Security Context and to a Group OSCORE Security Context, due to, for example, expiration, the RS MUST NOT use the pairwise OSCORE Security Context. The RS MUST respond with an unprotected 4.01 (Unauthorized) error message to received requests that correspond to a pairwise OSCORE Security Context with a deleted Access Token. Also, if the Client uses the Group OSCORE Security Context to send a request for any resource intended for OSCORE group members and that requires an active Access Token, the RS MUST respond with a 4.01 (Unauthorized) error message protected with the Group OSCORE Security Context.

The same considerations, related to the value of the ID Context changing, as in Appendix A.3.3 hold also for the RS.

A.3.5. Access Rights Verification

The RS MUST follow the procedures defined in Section 4.4.

Additionally, if the RS receives an OSCORE-protected request from a Client, the RS processes it according to [RFC8613].

If the OSCORE verification succeeds, and the target resource requires authorization, the RS retrieves the authorization information from the Access Token associated with the pairwise OSCORE Security Context and to the Group OSCORE Security Context. Then, the RS MUST verify that the action requested on the resource is authorized.

The response code MUST be 4.01 (Unauthorized) if the RS has no valid Access Token for the Client.

A.3.6. Change of Client's Authentication Credential in the Group

During its membership in the OSCORE group, the client might change the authentication credential it uses in the group. When this happens, the Client uploads the new authentication credential to the Group Manager, as defined in Section 11 of [I-D.ietf-ace-key-groupcomm-oscore].

After that, the Client may still have an Access Token previously uploaded to the RS, which is not expired yet and still valid to the best of the Client's knowledge. Then, in order to continue communicating with the RS, the Client MUST perform the following actions.

1. The Client requests a new Access Token to the AS, as defined in Appendix A.2.1 for the update of access rights, i.e., with the 'req_cnf' parameter including only a 'kid' field. In particular, when sending the POST request to the AS, the Client indicates:

   *The current Group Identifier of the OSCORE group, as value of the 'context_id' parameter.

   *The current Sender ID it has in the OSCORE group, as value of the 'salt_input' parameter.

   *The public key of the new authentication credential it uses in the OSCORE group, as value of the 'client_cred' parameter. In particular, the specified public key is the COSE Key equivalent to the new authentication credential that the Client uses in the OSCORE group.

*The proof-of-possession (PoP) evidence corresponding to the
         public key of the new authentication credential, as value of
         the 'client_cred_verify' or 'client_cred_verify_mac'
         parameter.

        *The same current or instead new set of access rights, as
         value of the 'scope' parameter.

   2. After receiving the response from the AS (see [Appendix A.2.2](#)),
      the Client performs the same exchanges with the RS as defined
      in [Appendix A.3](#), with the following difference: the POST
      request to /authz-info for uploading the new Access Token MUST
      be protected with the pairwise OSCORE Security Context shared
      with the RS.

   When receiving the new Access Token, the RS performs the same steps
   defined in [Appendix A.3.2](#). In particular, no new pairwise OSCORE
   Security Context is established between the Client and the RS.

A.4. Secure Communication with the AS

   The same considerations for secure communication with the AS as
   defined in [Section 5](#) hold.

A.5. Discarding the Security Context

   The Client and the RS MUST follow what is defined in [Section 6](#) of
   [[I-D.ietf-ace-oscore-profile](#)] about discarding the pairwise OSCORE
   Security Context.

   Additionally, they MUST follow what is defined in the main mode of
   the profile (see [Section 6](#)), with respect to the Group OSCORE
   Security Context.

   The Client or RS can acquire a new Group OSCORE Security Context, by
   re-joining the OSCORE group, e.g., by using the approach defined in
   [[I-D.ietf-ace-key-groupcomm-oscore](#)]. In such a case, the Client
   SHOULD request a new Access Token and post it to the RS, in order to
   establish a new pairwise OSCORE Security Context and bind it to the
   Group OSCORE Security Context obtained upon re-joining the group.

A.6. CBOR Mappings

   The new parameters defined in this document MUST be mapped to CBOR
   types as specified in [Figure 6](#), with the following addition, using
   the given integer abbreviation for the map key.

```
/----------------+----------+------------\
| Parameter name | CBOR Key | Value Type |
|----------------+----------+------------|
| client_cred    | TBD      | map        |
\----------------+----------+------------/
```

Figure 17: CBOR mappings for new parameters.

The new claims defined in this document MUST be mapped to CBOR types
as specified in Figure 7, with the following addition, using the
given integer abbreviation for the map key.

```
/--------------+----------+------------\
| Claim name   | CBOR Key | Value Type |
|--------------+----------+------------|
| client_cred  | TBD      | map        |
\--------------+----------+------------/
```

Figure 18: CBOR mappings for new claims.

A.7. Security Considerations

The dual mode of this profile inherits the security considerations
from the main mode (see Section 8), as well as from the security
considerations of the OSCORE profile of ACE [I-D.ietf-ace-oscore-
profile]. Also, the security considerations about OSCORE [RFC8613]
hold for the dual mode of this profile, as to the specific use of
OSCORE.

Unlike the main mode and consistently with Section 6.1 of [I-D.ietf-
ace-oauth-authz], the dual mode of this profile cannot be used to
issue an Access Token for an audience that comprises multiple RSs.
This is because the proof-of-possession key bound to an Access Token
is the OSCORE Master Secret included in the OSCORE_Input_Material
object of the 'cnf' claim, and it has to be shared only between the
Client and one RS.

A.8. Privacy Considerations

The same privacy considerations as defined in the main mode of this
profile apply (see Section 9).

In addition, as this profile mode also uses OSCORE, the privacy
considerations from [RFC8613] apply as well, as to the specific use
of OSCORE.

Furthermore, this profile mode inherits the privacy considerations
from the OSCORE profile of ACE [I-D.ietf-ace-oscore-profile].

Appendix B. Profile Requirements

This appendix lists the specifications on this profile based on the requirements of the ACE framework, as requested in Appendix C of [I-D.ietf-ace-oauth-authz].

*(Optional) discovery process of how the Client finds the right AS for an RS it wants to send a request to: Not specified.

*Communication protocol the Client and the RS must use: CoAP.

*Security protocol(s) the Client and RS must use: Group OSCORE, i.e., exchange of secure messages by using a pre-established Group OSCORE Security Context. The optional dual mode defined in Appendix A additionally uses OSCORE, i.e., establishment of a pairwise OSCORE Security Context and exchange of secure messages.

*How the Client and the RS mutually authenticate: Explicitly, by possession of a common Group OSCORE Security Context, and by either: usage of digital signatures embedded in messages, if protected with the group mode of Group OSCORE; or protection of messages with the pairwise mode of Group OSCORE, by using pairwise symmetric keys, derived from the asymmetric keys of the two peers exchanging the message. Note that the mutual authentication is not completed before the Client has verified an OSCORE or a Group OSCORE response using the corresponding security context.

*Content-format of the protocol messages: "application/ace+cbor".

*Proof-of-Possession protocol(s) and how to select one; which key types (e.g., symmetric/asymmetric) supported: Group OSCORE algorithms; distributed and verified asymmetric keys. In the optional dual mode defined in Appendix A: OSCORE algorithms; pre-established symmetric keys.

*profile identifier: coap_group_oscore

*(Optional) how the RS talks to the AS for introspection: HTTP/ CoAP (+ TLS/DTLS/OSCORE).

*How the client talks to the AS for requesting a token: HTTP/CoAP (+ TLS/DTLS/OSCORE).

*How/if the authz-info endpoint is protected: Not protected.

*(Optional) other methods of token transport than the authz-info endpoint: Not specified.

Acknowledgments

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