

Internet Research Task Force (IRTF)
Internet-Draft
Intended status: Informational
Expires: December 16, 2017

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June 14, 2017

KangarooTwelve
draft-viguier-kangarootwelve-00

Abstract

This document defines the KangarooTwelve eXtendable Output Function (XOF), a hash function with arbitrary output length. It provides an efficient and secure hashing primitive, which is able to exploit the parallelism of the implementation in a scalable way. It uses tree hashing over a round-reduced version of SHAKE128 as underlying primitive.

This document builds up on the definitions of the permutations and of the sponge construction in [FIPS 202], and is meant to serve as a stable reference and an implementation guide.

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[1.](#) Introduction

This document defines the KangarooTwelve extendable Output Function (XOF) [[K12](#)], i.e. a generalization of a hash function that can return arbitrary output length. KangarooTwelve is based on a Keccak-p permutation specified in [[FIPS202](#)] and aims at higher speed than SHAKE and SHA-3.

The SHA-3 functions process data in a serial manner and unable to optimally exploit parallelism available in modern CPU architectures. KangarooTwelve splits the input message in fragments and applies an inner hash function F on each of them separately. It then applies F again on the concatenation of the digests. It makes use of Sakura coding for ensuring soundness of the tree hashing mode [[SAKURA](#)]. The inner hash function F is a sponge function and uses a round-reduced version of the permutation used in Keccak. Its security builds up on the scrutiny that Keccak has received since its publication [[KECCAK CRYPTANALYSIS](#)].

1.1. Conventions

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC 2119](#) [[RFC2119](#)].

The following notations are used throughout the document:

``...`` denotes a bit-string. For example, ``1010101``.

A 8 bit string ``b_0 b_1 b_2 b_3 b_4 b_5 b_6 b_7`` is a byte represented by an integer value v following the LSB 0 convention, i.e.

$$v = \text{sum for } i=0..7 \text{ of } 2^i * b_i$$

For example, ``11100000`` = 7. The following diagram represents the byte "07" with value 7 (decimal).

```

      Significance of Bits
    MSB 7 6 5 4 3 2 1 0 LSB
      +---+---+---+---+
      |0 0 0 0 0 1 1 1|
      +---+---+---+---+
    hex:  0       7

```

`"..."` denotes a string of bytes given in hexadecimal. For example, "0B 80", which can be also be seen as a bit-string : ``11010000 00000001``.

`|s|` denotes the length of a byte string "s". For example, `|"FF FF"|` = 2.

``0^b`` denotes the repetition of bit ``0`` b times. For example, ``0^4`` = ``0000``.

``0^0`` denotes the empty bit-string.

``1^b`` denotes the repetition of bit ``1`` b times. For example, ``1^3`` = ``111``.

`"00^b"` denotes the b times the repetition of byte "00". For example, `"00^7"` = "00 00 00 00 00 00 00".

`a||b` denotes the concatenation of two strings 'a' and 'b'. For example, ``10`||`01`` = ``1001``

`s[n:m]` denotes the selection of bytes from `n` to `m` exclusive of a string `'s'`. For example, for `s = "A5 C6 D7"`, `s[0:1] = "A5"` and `s[1:3] = "C6 D7"`.

2. Specifications

KangarooTwelve is an eXtendable Output Function (XOF). It takes as an input a pair of byte-strings (`M`, `C`) and a positive integer `L` where

`M` byte-string, is the Message and

`C` byte-string, is a Customization string and

`L` positive integer, the length of the output in bytes.

The Customization string serves as domain separation. It is typically a short string such as a name or an identifier (e.g. `URI`, `ODI...`)

2.1. Inner function: `F`

The inner function `F` makes use of the permutation Keccak-`p[1600,n_r=12]`, i.e., a version of the one used in SHAKE and SHA-3 instances reduced to `n_r=12` rounds and specified in FIPS 202 [[FIPS202](#)]. `F` is a sponge function calling this permutation, multi-rate padding `pad10*1` and with a rate of 168 bytes (= 1344 bits):

$$F = \text{Sponge}[\text{Keccak-p}[1600,n_r=12], \text{pad10*1}, r=1344]$$

It follows that `F` has a capacity of $1600 - 1344 = 256$ bits.

The sponge function `F` takes as an input a bit-string `S` and a positive integer `L` where

`S` bit-string, is the input String and

`L` positive integer, the Length of the output in bytes

The input string `S` SHOULD be represented as a pair (`Sbytes`, `dS`), where `Sbytes` contains only bytes and where `dS` is the delimited suffix representing the trailing bits.

First, let `S = Sbytes || Sbits`, where `Sbytes` contains only bytes and `Sbits` contains at most 7 bits. Then, convert `Sbits` into the delimited suffix `dS` by appending a bit `'1'` and as many bits `'0'` as necessary so that `dS` is a byte. The numerical value of `dS` is thus:

$$dS = 2^{|Sbits|} + \text{sum for } i=0..|Sbits|-1 \text{ of } 2^i * Sbits_i$$

Notice that the most significant bit `1` of dS coincides with the first bit of padding in the multi-rate padding rule pad10*1. The implementation of F therefore SHOULD add dS to the state and then the second bit of padding. [Appendix A.2](#) provides a pseudo code version.

In the table below, here are some examples of values, including those that are used in this document:

Sbits	bit-string	value (dec)	delimited Suffix (dS)
`	`10000000`	1	"01"
`01`	`01100000`	6	"06"
`11`	`11100000`	7	"07"
`110`	`11010000`	11	"0B"

2.2. Tree hashing over F

On top of the sponge function F, KangarooTwelve uses a Sakura-compatible tree hash mode [[SAKURA](#)]. First, merge M and C to a single input string S in a reversible way. `right_encode( |C| )` gives the length in bytes of C as a byte-string. See [Section 2.3](#).

$$S = M \parallel C \parallel \text{right_encode}(|C|)$$

Then, split S into n chunks of 8192 bytes.

$$\begin{aligned} S &= S_0 \parallel \dots \parallel S_{n-1} \\ |S_0| &= \dots = |S_{n-2}| = 8192 \text{ bytes} \\ |S_{n-1}| &\leq 8192 \text{ bytes} \end{aligned}$$

From $S_1 \dots S_{n-1}$, compute the 32-bytes hashes $CV_0 \dots CV_{n-2}$. This computation SHOULD exploit the parallelism available on the platform in order to be optimally efficient.

$$\begin{aligned} \text{Node}_i &= S_{i+1} \parallel \text{'110'} \\ CV_i &= F(\text{Node}_i, 32) \end{aligned}$$

Compute the final node: Node*.

- o If $|S| \leq 8192$ bytes, then $\text{Node}^* = S \parallel \text{'11'}$
- o Otherwise compute Node* as follow:


```

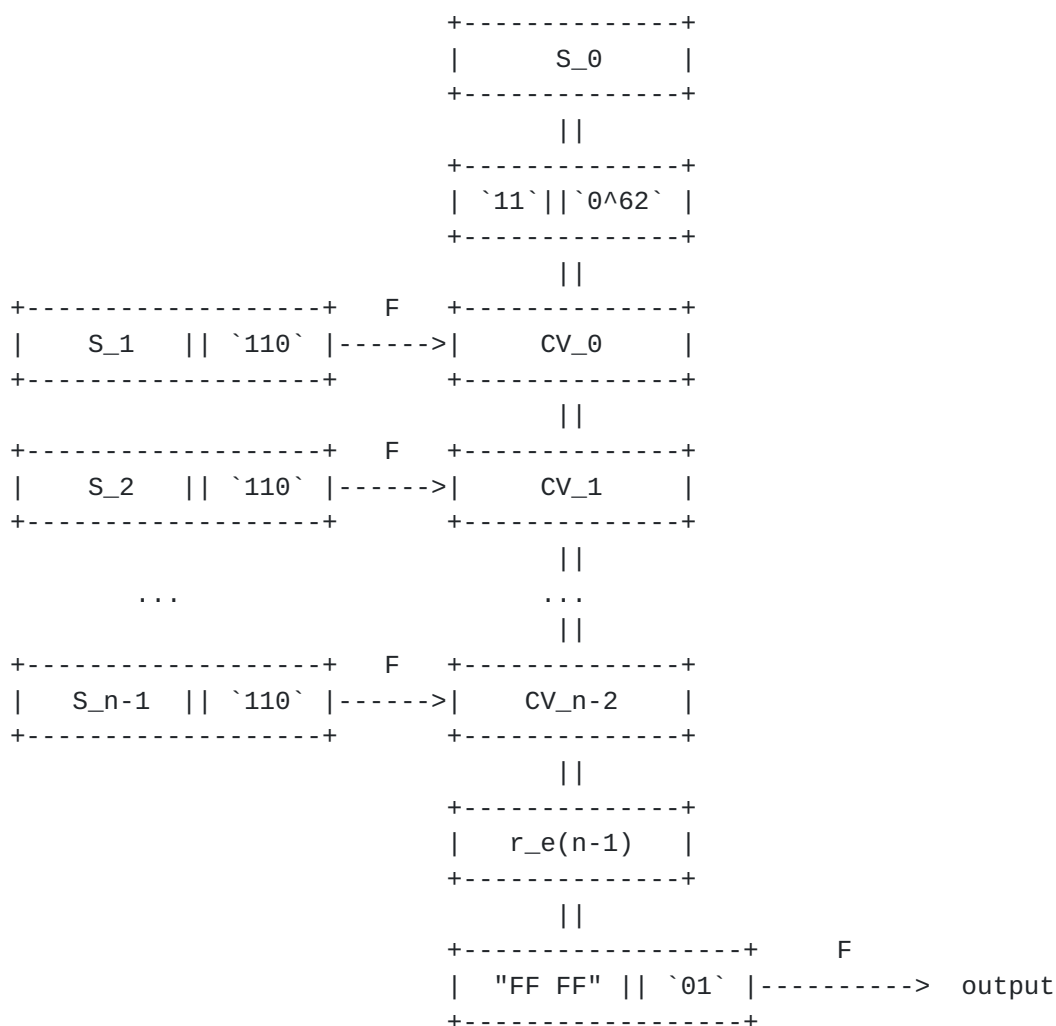
Node* = S_0 || "03 00 00 00 00 00 00 00"
Node* = Node* || CV_0
..
Node* = Node* || CV_n-2
Node* = Node* || right_encode(n-1)
Node* = Node* || "FF FF" || `01`

```

Finally, KangarooTwelve output is retrieved from $F(\text{Node}^*)$.

$\text{KangarooTwelve}(M, C, L) = F(\text{Node}^*, L)$

For $|S| > 8192$ bytes, KangarooTwelve computation flow is as follow:



We provide a pseudo code version in [Appendix A.3](#).

2.3. right_encode(x)

The function `right_encode` takes as inputs a non negative integer $x < 256^{255}$ and outputs a string of bytes `x_n || .. || x_0 || n` where

$$x = \text{sum from } i=0..n \text{ of } 256^i * x_i$$

A pseudo code version is as follow.

```
right_encode(x):  
  S = 0^0  
  
  while x > 0  
    S = x % 256 || S  
    x = x / 256  
  
  S = S || length(S)  
  
  return S  
end
```

3. Test vectors

Test vectors are based on the repetition of pattern the "00 01 .. FA" with a specific length. `ptn(n)` defines a string by repeating the pattern "00 01 .. FA" as many times as necessary and truncated to n bytes e.g.

Pattern for a length of 17 bytes:

`ptn(17) =`

"00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F 10"

Pattern for a length of 17^2 bytes:

`ptn(17^2) =`

```
"00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F
 10 11 12 13 14 15 16 17 18 19 1A 1B 1C 1D 1E 1F
 20 21 22 23 24 25 26 27 28 29 2A 2B 2C 2D 2E 2F
 30 31 32 33 34 35 36 37 38 39 3A 3B 3C 3D 3E 3F
 40 41 42 43 44 45 46 47 48 49 4A 4B 4C 4D 4E 4F
 50 51 52 53 54 55 56 57 58 59 5A 5B 5C 5D 5E 5F
 60 61 62 63 64 65 66 67 68 69 6A 6B 6C 6D 6E 6F
 70 71 72 73 74 75 76 77 78 79 7A 7B 7C 7D 7E 7F
 80 81 82 83 84 85 86 87 88 89 8A 8B 8C 8D 8E 8F
 90 91 92 93 94 95 96 97 98 99 9A 9B 9C 9D 9E 9F
 A0 A1 A2 A3 A4 A5 A6 A7 A8 A9 AA AB AC AD AE AF
 B0 B1 B2 B3 B4 B5 B6 B7 B8 B9 BA BB BC BD BE BF
 C0 C1 C2 C3 C4 C5 C6 C7 C8 C9 CA CB CC CD CE CF
 D0 D1 D2 D3 D4 D5 D6 D7 D8 D9 DA DB DC DD DE DF
 E0 E1 E2 E3 E4 E5 E6 E7 E8 E9 EA EB EC ED EE EF
 F0 F1 F2 F3 F4 F5 F6 F7 F8 F9 FA
 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F
 10 11 12 13 14 15 16 17 18 19 1A 1B 1C 1D 1E 1F
 20 21 22 23 24 25"
```

`KangarooTwelve(M=0^0, C=0^0, 32):`

```
"1A C2 D4 50 FC 3B 42 05 D1 9D A7 BF CA 1B 37 51
 3C 08 03 57 7A C7 16 7F 06 FE 2C E1 F0 EF 39 E5"
```

`KangarooTwelve(M=0^0, C=0^0, 64):`

```
"1A C2 D4 50 FC 3B 42 05 D1 9D A7 BF CA 1B 37 51
 3C 08 03 57 7A C7 16 7F 06 FE 2C E1 F0 EF 39 E5
 42 69 C0 56 B8 C8 2E 48 27 60 38 B6 D2 92 96 6C
 C0 7A 3D 46 45 27 2E 31 FF 38 50 81 39 EB 0A 71"
```

`KangarooTwelve(M=0^0, C=0^0, 10032), last 32 bytes:`

```
"E8 DC 56 36 42 F7 22 8C 84 68 4C 89 84 05 D3 A8
 34 79 91 58 C0 79 B1 28 80 27 7A 1D 28 E2 FF 6D"
```

`KangarooTwelve(M=ptn(1 bytes), C=0^0, 32):`

```
"2B DA 92 45 0E 8B 14 7F 8A 7C B6 29 E7 84 A0 58
 EF CA 7C F7 D8 21 8E 02 D3 45 DF AA 65 24 4A 1F"
```

`KangarooTwelve(M=ptn(17 bytes), C=0^0, 32):`

```
"6B F7 5F A2 23 91 98 DB 47 72 E3 64 78 F8 E1 9B
 0F 37 12 05 F6 A9 A9 3A 27 3F 51 DF 37 12 28 88"
```

`KangarooTwelve(M=ptn(17^2 bytes), C=0^0, 32):`

```
"0C 31 5E BC DE F6 14 26 DE 7D CF 8F B7 25 D1
 E7 46 75 D7 F5 32 7A 50 67 F3 67 B1 08 EC B6 7C"
```


KangarooTwelve(M=ptn(17³ bytes), C=0⁰, 32):

```
"CB 55 2E 2E C7 7D 99 10 70 1D 57 8B 45 7D DF 77
 2C 12 E3 22 E4 EE 7F E4 17 F9 2C 75 8F 0D 59 D0"
```

KangarooTwelve(M=ptn(17⁴ bytes), C=0⁰, 32):

```
"87 01 04 5E 22 20 53 45 FF 4D DA 05 55 5C BB 5C
 3A F1 A7 71 C2 B8 9B AE F3 7D B4 3D 99 98 B9 FE"
```

KangarooTwelve(M=ptn(17⁵ bytes), C=0⁰, 32):

```
"84 4D 61 09 33 B1 B9 96 3C BD EB 5A E3 B6 B0 5C
 C7 CB D6 7C EE DF 88 3E B6 78 A0 A8 E0 37 16 82"
```

KangarooTwelve(M=ptn(17⁶ bytes), C=0⁰, 32):

```
"3C 39 07 82 A8 A4 E8 9F A6 36 7F 72 FE AA F1 32
 55 C8 D9 58 78 48 1D 3C D8 CE 85 F5 8E 88 0A F8"
```

KangarooTwelve(M=0⁰, C=ptn(1 bytes), 32):

```
"FA B6 58 DB 63 E9 4A 24 61 88 BF 7A F6 9A 13 30
 45 F4 6E E9 84 C5 6E 3C 33 28 CA AF 1A A1 A5 83"
```

KangarooTwelve(M=0xff, C=ptn(41 bytes), 32):

```
"D8 48 C5 06 8C ED 73 6F 44 62 15 9B 98 67 FD 4C
 20 B8 08 AC C3 D5 BC 48 E0 B0 6B A0 A3 76 2E C4"
```

KangarooTwelve(M=0xff ff ff, C=ptn(41²), 32):

```
"C3 89 E5 00 9A E5 71 20 85 4C 2E 8C 64 67 0A C0
 13 58 CF 4C 1B AF 89 44 7A 72 42 34 DC 7C ED 74"
```

KangarooTwelve(M=0xff ff ff ff ff ff ff, C=ptn(41³ bytes), 32):

```
"75 D2 F8 6A 2E 64 45 66 72 6B 4F BC FC 56 57 B9
 DB CF 07 0C 7B 0D CA 06 45 0A B2 91 D7 44 3B CF"
```

4. IANA Considerations

None.

5. Security Considerations

This document is meant to serve as a stable reference and an implementation guide for the KangarooTwelve extendable Output Function. It makes no assertion to its security and relies on the cryptanalysis of Keccak [[KECCAK CRYPTANALYSIS](#)].

6. References

6.1. Normative References

- [FIPS202] National Institute of Standards and Technology, "FIPS PUB 202 - SHA-3 Standard: Permutation-Based Hash and Extendable-Output Functions",
WWW <http://dx.doi.org/10.6028/NIST.FIPS.202>, August 2015.
- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", [BCP 14](#), [RFC 2119](#),
DOI 10.17487/RFC2119, March 1997,
<<http://www.rfc-editor.org/info/rfc2119>>.

6.2. Informative References

- [K12] Bertoni, G., Daemen, J., Peeters, M., Van Assche, G., and R. Van Keer, "KangarooTwelve: fast hashing based on Keccak-p", WWW <http://eprint.iacr.org/2016/770.pdf>, August 2016.
- [KECCAK_CRYPTANALYSIS] Keccak Team, "Summary of Third-party cryptanalysis of Keccak", WWW https://www.keccak.team/third_party.html, 2017.
- [SAKURA] Bertoni, G., Daemen, J., Peeters, M., and G. Van Assche, "Sakura: a flexible coding for tree hashing",
WWW <http://eprint.iacr.org/2013/231.pdf>, April 2013.

[Appendix A.](#) Pseudo code

The sub-sections of this appendix contain pseudo code definitions of KangarooTwelve.

[A.1.](#) Keccak-p[1600] over 12 rounds

```

Keccak-p_1600_12(state):
    R = "D5"

    for x from 0 to 4
        for y from 0 to 4
            lanes[x][y] = state[8*(x+5*y):8*(x+5*y)+8]

    for round from 12 to 23
        # theta
        for x from 0 to 4
            C[x] = lanes[x][0]
            C[x] ^= lanes[x][1]
            C[x] ^= lanes[x][2]
            C[x] ^= lanes[x][3]
            C[x] ^= lanes[x][4]
        for x from 0 to 4
            D[x] = C[(x+4)%5] ^ ROL64(C[(x+1)%5], 1)
        for y from 0 to 4
            for x from 0 to 4
                lanes = lanes[x][y]^D[x]

        # rho and pi
        (x, y) = (1, 0)
        current = lanes[x][y]
        for t from 0 to 23
            (x, y) = (y, (2*x+3*y)%5)
            (current, lanes[x][y]) =
                (lanes[x][y], ROL64(current, (t+1)*(t+2)/2))

        # chi
        for y from 0 to 4
            for x from 0 to 4
                T[x] = lanes[x][y]
                for x from 0 to 4
                    lanes[x][y] = T[x] ^ ((not T[(x+1)%5]) & T[(x+2)%5])

        # iota
        for j from 0 to 6
            R = ((R << 1) ^ ((R >> 7)* "71")) % 256
            if (R & 2)
                lanes[0][0] = lanes[0][0] ^ (1 << ((1<<j)-1))

```



```
state = 0^0
for x from 0 to 4
  for y from 0 to 4
    state = state || lanes[x][y]

return state
end
```

where $\text{ROL64}(x, y)$ is a rotation of the 'x' 64-bit word toward the bits with higher indexes by 'y' bits.

[A.2.](#) Inner function F

```
F(inputBytes, Suffix, outputByteLen):
  state = "00^200"
  blockSize = 0
  offset = 0

  # === Absorb inputBytes ===
  while offset < |inputBytes|
    blockSize = min( |inputBytes| - offset, 168)
    state ^= inputBytes[offset : offset + blockSize]
    offset = offset + blockSize

    if blockSize = 168
      state = Keccak-p_1600_12(state)
      blockSize = 0

  # === Absorb Suffix ===
  state ^= "00^blockSize" || Suffix
  if (Suffix & "80") != 0 and blockSize == 167
    state = Keccak-p_1600_12(state)
  state ^= "00^167" || "80"

  state = Keccak-p_1600_12(state)

  # === Squeeze ===
  while outputByteLen > 0
    blockSize = min(outputByteLen, 168)
    outputBytes = outputBytes || state[0:blockSize]
    outputByteLen = outputByteLen - blockSize

    if outputByteLen > 0
      state = Keccak-p_1600_12(state)

  return outputBytes
end
```


[A.3.](#) KangarooTwelve

```
KangarooTwelve(inputMessage, customString, outputByteLen):
  S = inputMessage || customString
  S = S || right_encode( |customString| )

  if |S| <= 8192
    return F(S, "07", outputByteLen)
  else
    # === Kangaroo hopping ===
    Node* = S[0:8192] || "03 00^7"
    offset = 8192
    while offset < |inputBytes|
      blockSize = min( |inputBytes| - offset, 8192)
      CV = F(inputBytes[offset : offset + blockSize], "0B", 32)
      Node* = Node* || CV
      offset = offset + blockSize

    Node* = Node* || right_encode( |S| / 8192 ) || "FF FF"
    return F(Node*, "06", outputByteLen)
  end
```

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