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IP Delivery Delay Detection Protocol
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Abstract

This memo describes a method that indicates to the receiver of an IP Internet Protocol datagram the time at which it was sent. The receiver of the datagram can then judge how recently it was sent and discard packets deemed to be 'too old' according to their policy.

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1. Introduction

IP datagrams can be subject to a delivery delay attack, where a host or gateway receives datagrams that are not fresh. A fresh datagram is defined as one "Recently generated; not replayed from some earlier interaction of the protocol." [[RFC4949](#)]. Even when IP datagrams are validated legitimate using an integrity transform and are protected from replay, delivered datagrams may have been delayed for an unbounded period of time.

Delivery delay detection is useful when the value of data included in the IP datagram is time sensitive. Delivery delay protection is even more valuable when replay protection is not available. For example, while the IPsec [[RFC4301](#)] transforms IP Encapsulating Security Payload (ESP) [[RFC4303](#)] and IP Authentication Header [[RFC4302](#)] include replay protection, it must be disabled when multi-sender IPsec security associations are used to protect multicast and group communications. Such Many-to-Many Applications [[RFC3170](#)] often require the use of multi-sender security associations, where receivers cannot use the sequence number replay protection method. In such cases, a single counter cannot record responses from multiple senders, and no provision is made for multiple counters in the security association.

This memo describes an IP Delivery Delay Detection Protocol using timestamps to declare the age of an IP datagram, enabling receivers to make a judgement whether to accept an IP datagram as "fresh" [[RFC4949](#)].

1.1. Requirements notation

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [[RFC2119](#)].

2. IP Delivery Delay Detection Protocol (IP-D3P)

An IP-D3P datagram consists of a header and an IP payload. The IP-D3P header includes a timestamp, which receivers of the packet use to determine whether or not the packet has been recently generated. Receivers compare the timestamp delivered in the IP packet to their local time and make a determination as to whether it should be accepted. Freshness does not provide replay protection, as it makes no explicit judgement as to whether a receiver has received a particular packet before. Rather, it allows the receiver to determine whether the packet has been delayed in delivery beyond an acceptable point in time. A typical policy would be to choose a time larger than a reasonable delivery time, which delay indicates a possible packet delay attack. When combined with replay protection, IP-D3P provides the receiver assurance that the packet is not only unique, but also fresh.

Varying formats of timestamps are possible. This memo declares several formats with varying degrees of timestamp precision. It is important that both sender and receiver interpret the time in the same degree of accuracy, else the receiver determination of freshness may be in error. For example, if the sender reported a timestamp in seconds and the receiver interpreted the timestamp in microseconds it would likely reject many datagrams due to the precision error.

2.1. IP-D3P Header

The IP-D3P header consists of a type and a fixed-size timestamp, followed by a datagram, which may be either a full or partial IP datagram depending on the use case.

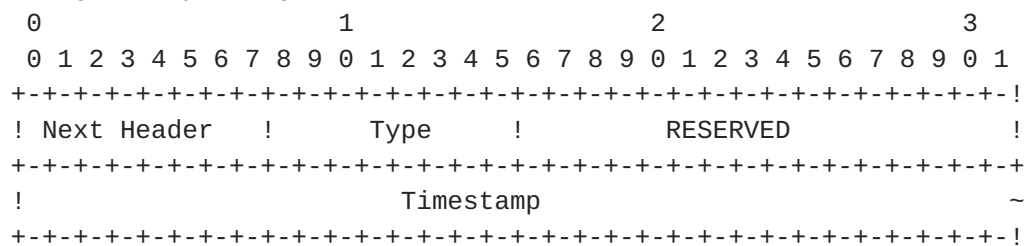


Figure 1: IP-D3P Header

The header fields are defined as follows:

- o Next Header (1 octet) -- Identifies the type of data in the Datagram following the IP-D3P header. Valid types are described in [[PROT-REG](#)].
- o Type (1 octet) -- Type of timestamp contained in the header. Valid values are defined in [Section 2.1.1](#).

- o RESERVED (2 octets) -- MUST be zero, and MUST be ignored on receipt.
- o Timestamp (variable) -- Timestamp declaring when the packet was originated. Length depends on type of timestamp.

2.1.1. Type

This memo defines the following types. A conforming implementation MUST implement POSIX-TIME-SEC. Each type indicates a fixed sized timestamp, which is valuable for packet evaluation efficiency. The length of the header can then also be trivially determined based on its Type.

- o POSIX-TIME-SEC. Specifies that the timestamp refers to current time as seconds since the POSIX Epoch [[POSIX.1](#)] (i.e., time as the number of seconds beginning January 1, 1970 not counting leap seconds. The timestamp has a length of 4 bytes.
- o POSIX-TIME-DECISEC. Specifies that timestamp refers to current time as seconds since the POSIX Epoch, specified in deciseconds (tenths of a second). The timestamp has a length of 6 bytes.
- o POSIX-TIME-MSEC. Specifies that the timestamp refers to current time as seconds since the POSIX Epoch, specified in milliseconds. The timestamp has a length of 8 bytes.

2.2. Outbound Packet Processing

An IP-D3P header is added to the packet as described in Figure 1. The sender inserts the type defined by policy into the Type field and places the current time value in the Timestamp field as specified by the Type. For example, when the Type is POSIX-TIME-SEC, the current value of seconds from the system clock is translated into POSIX time as defined in [Section 2.1.1](#) and placed into the header.

When IP-D3P is added to an IP packet, the IP-D3P Next Header field identifies the Protocol Id representing the next header (e.g., TCP). When IP-D3P is used with a IP-IP encapsulation (e.g., IPsec Tunnel Mode), the IP-D3P Next Header field represents the type of datagram following (e.g., 4 for IPv4, 41 for IPv6).

2.3. Inbound Packet Processing

A receiver first validates that the Type value in the IP-D3P header matches the expected type, as according to policy. It then extracts the timestamp and compares it to a sliding window maintained by the receiver that is based on the Type of time. A timestamp found within

the sliding window is accepted, and the packet is treated as a fresh packet (e.g., is processed further). If the timestamp is outside of the receiver's window, the packet is discarded.

Figure 2 shows the sliding window used by D3P. A receiver chooses a window size of W , which lies between W_s and W_e centered around the current time of the receiver (T_r). When a packet is received with a Timestamp T_s that lies below W_s , it is rejected as being too old. If T_s lies in between W_s and W_e it is accepted as a fresh packet. If T_s is in advance of W_e it is rejected as an invalid time. However, the reason for rejection (being too old or invalid) MUST NOT be discernible to any entity other than the receiver who rejected the packet.

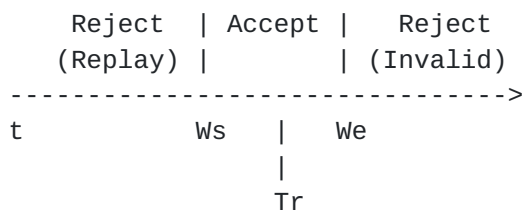


Figure 2: D3P Sliding Window

Decapsulation happens as follows. When IP-D3P follows an IP header (e.g., IPsec Transport Mode) the Next Header value in the IP-D3P header is placed in the IP header in place of TBD-1. The IP-D3P header is then removed from the packet. New Total Length and checksum fields of the IP header are also restored by the receiver. In the case of an IP-IP encapsulated packet, the IP-D3P header is simply discarded and processing continues on the encapsulated IP packet.

3. IP-D3P and IPsec

When IP-D3P is defined as part of the policy, it is applied before ESP or AH processing. Another instance of this approach is IPcomp [RFC3173]. Using IP-D3P with IPsec is particularly useful to protect group communications when the use of single-sender Security Associations (SAs) is not possible, and thus IPsec replay protection is not available (Section 2.2 of [RFC4303]). This has been a particular problem when applying IPsec to routing protocols using multicast communications (e.g., OSPF [RFC4552] and PIM [RFC4601]), where maintaining an IPsec SA for each sender is not always feasible. Additionally, for operational reasons IPsec SAs and keying material are often manually configured for these routing protocols, in which case the use of IPsec replay protection is not possible. The use of IP-D3P does not replace this missing replay protection, but can effectively bound the lifetime of a replayed packet.

When IP-D3P is applied as part of IPsec transport mode encapsulation, the IP-D3P encapsulation is added as the first Protocol following the IP header as shown in Figure 3 and Figure 4. Several IP header fields need to be adjusting when adding the IP-D3P header: Total Length is adjusted to the original length of the payload, the Protocol field (IPv4 header) or Next Header field (IPv6) is set to TBD-1 identifying IP-D3P as the next protocol, and the header checksum is adjusted.

[illegible]

Figure 3: IPv4 Transport Mode Encapsulation

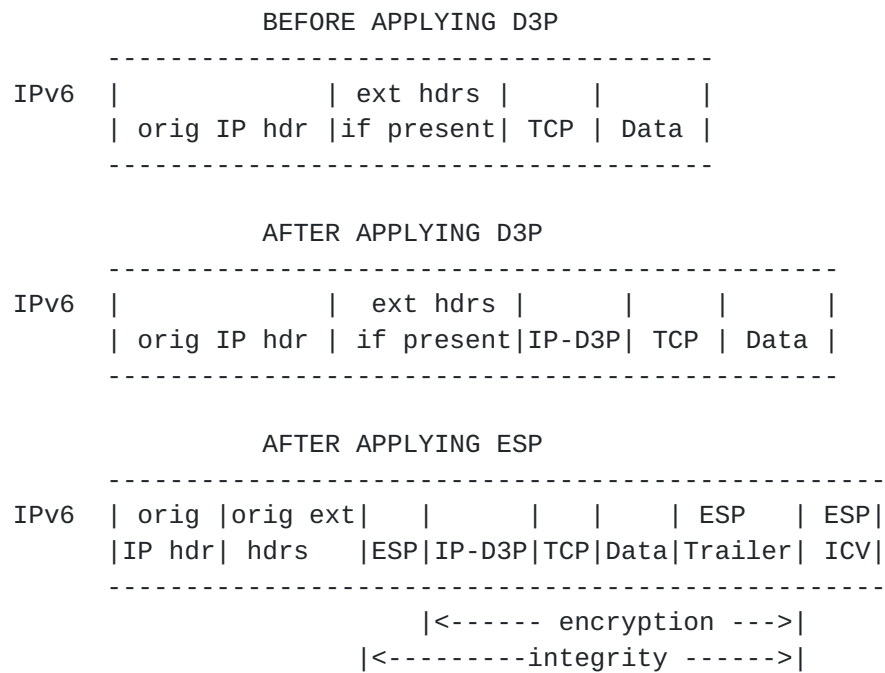


Figure 4: IPv6 Transport Mode Encapsulation

When IP-D3P is applied as part of IPsec tunnel mode encapsulation, the IP-D3P encapsulation is added immediately after the outer IP header, such that the inner IP header is fully encapsulated. This is shown in Figure 5 and Figure 6. The Protocol field in the outer IPv4 header or Next Header field in the outer IPv6 header is set to TBD-1, indicating that an IP-D3P datagram follows.

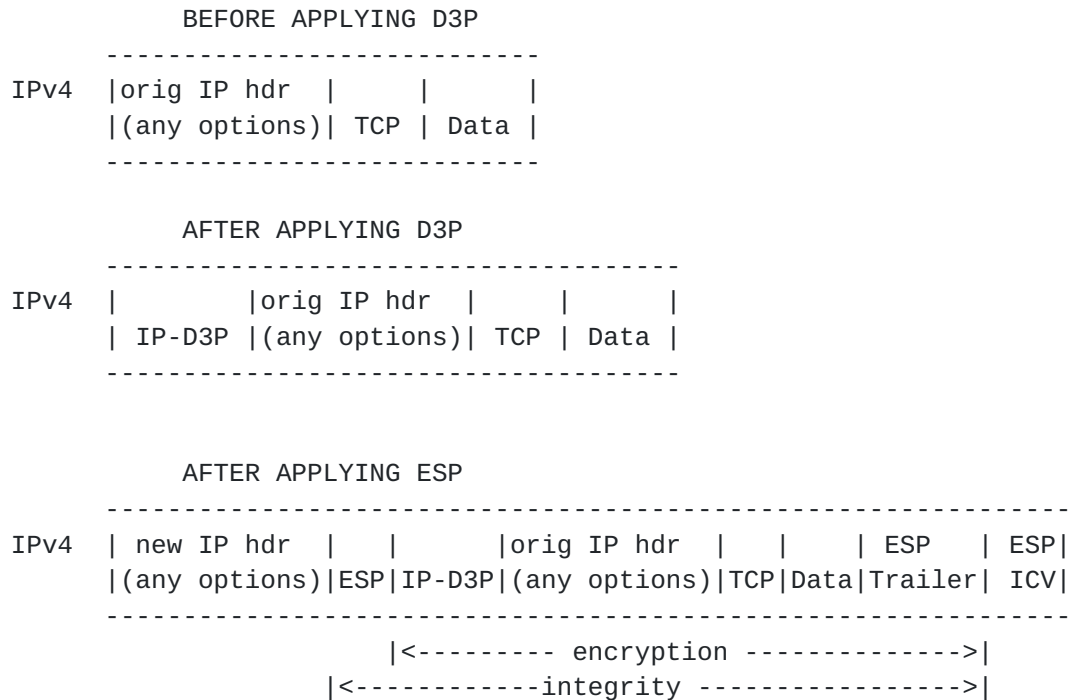


Figure 5: IPv4 Tunnel Mode Encapsulation

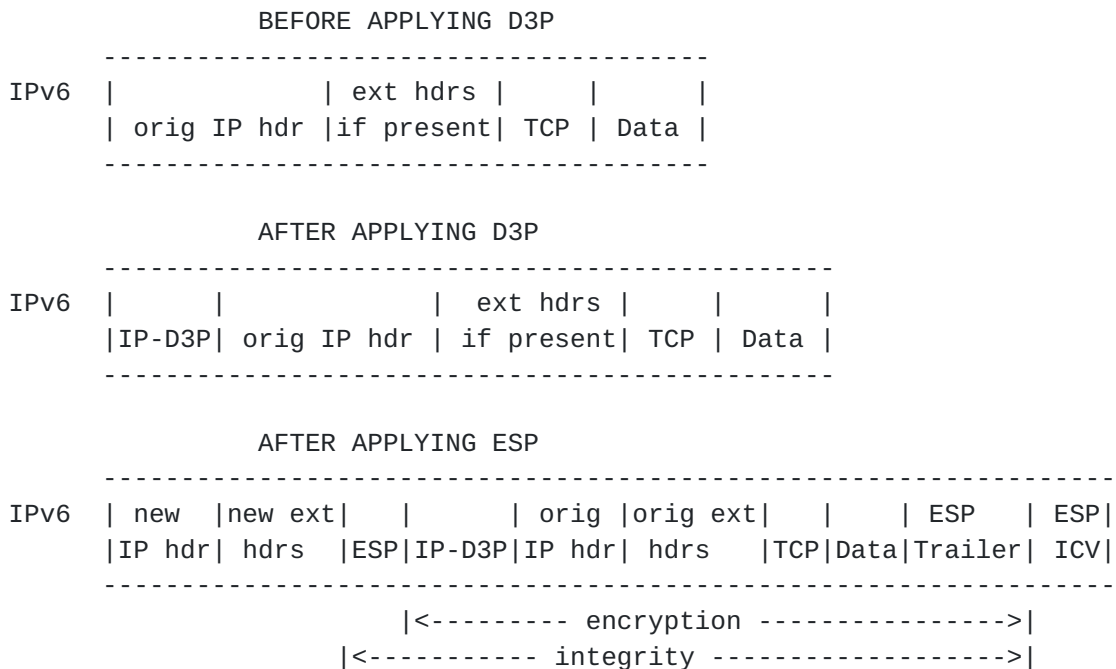


Figure 6: IPv6 Tunnel Mode Encapsulation

4. Policy Distribution using Group Domain of Interpretation

The Group Domain of Interpretation (GDOI) [[RFC6407](#)] is a group key management protocol used to distribute group policy and keying material to a set of Group Members (GMs). When groups using GDOI key management services require the use of IP-D3P, GDOI distributes this policy in a Group Associated Policy (GAP) payload using the D3P-TYPE attribute. This attribute indicates the use of IP-D3P for all SA TEKS distributed within the SA payload, and defines the Type of timestamp that is to be placed into datagrams matching those SA TEKS. Value for the D3P-TYPE attribute are taken from the IP-D3P Type Registry.

The GDOI KS policy can also define the size of the time window using the GAP payload D3P-WINDOWSIZE attribute. The value of the attribute is in the units defined by the Type. For example, for the POSIX-TIME-SEC attribute the value is in seconds. If the D3P-TYPE is sent without an accompanying D3P-WINDOWSIZE attribute then window size is chosen by the receiver.

5. Security Considerations

IP-D3P provides an indication of freshness an IP packet rather than an absolute detection of all replays.

In networks where active attackers are anticipated, packets including an IP-D3P header SHOULD be protected by an IPsec transform. IPsec integrity protection defends against active man in the middle attackers changing the timestamp, and making it appear to be stale or invalid.

The Types defined in this memo all use the system clock. In many cases, the system clock is set from an external protocol (e.g., NTP [[RFC5905](#)]), and indeed this will maximize the likelihood that the system clocks of both sender and receiver are synchronized. However, care should be taken that external protocol is resistant to a man in the middle attack. For example, NTP packets could be distributed within an independent well-protected network believed not available to active man in the middle attackers, or the NTP Message Digest could be used to provide integrity protection.

6. IANA Considerations

A new IP Protocol is defined for IP-D3P (value TBD-1).

A new IP-D3P Type registry is created. Values for this attribute are shown in the following table. The terms Reserved and Unassigned are to be applied as defined in [[RFC5226](#)].

| Value | Type |
|---------|--------------------|
| ----- | ---- |
| 0 | Reserved |
| 1 | POSIX-TIME-SEC |
| 2 | POSIX-TIME-DECISEC |
| 3 | POSIX-TIME-MSEC |
| 4-128 | Unassigned |
| 129-255 | Private Use |

The following additions are made to the GDOI Payloads [[GDOI-REG](#)] registry.

Two new attributes are added to the GAP Payload Policy Attributes registry.

- o Attribute type D3P-TYPE is a Basic attribute and takes the value of TBD-2. Valid values come from the IP-D3P Type registry.
- o Attribute type D3P-WINDOWSIZE is a Variable attribute and takes the value of TBD-3. There are no described set of valid values.

7. Acknowledgements

Some diagrams were adapted from similar diagrams in [[RFC4303](#)].

8. References

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