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Internet Cache Protocol (ICP), version 2

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Abstract

This draft document describes the Internet Cache Protocol (ICP) as currently implemented in a couple of World-Wide Web proxy cache packages. A companion document (RFCXXXX, <[draft-wessels-icp-v2-appl-00.txt](#)>) describes the application of ICP to Web caches. ICP was initially developed by Peter Danzig, et. al. at the University of Southern California. It evolved as an important part of hierarchical caching on the Harvest research project.

1. Introduction

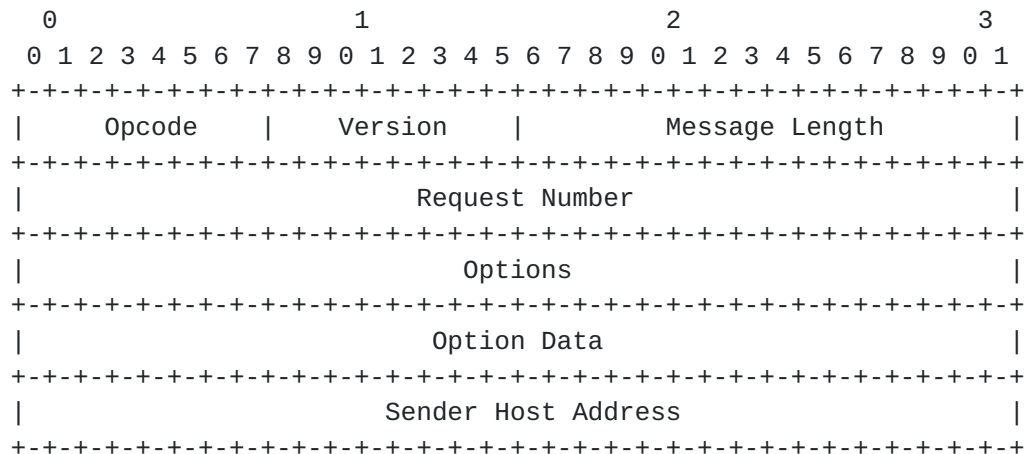
ICP is a message format used for communicating between Web caches. Although Web caches use HTTP[1] for the transfer of object data, caches benefit from a simpler, lighter communication protocol. ICP is primarily used in a cache mesh to locate specific Web objects in neighbor caches. One cache will send an ICP query to its neighbors. The neighbors will send back ICP replies indicating a ``HIT'' or a ``MISS.''

In current practice, ICP is implemented on top of UDP, but there is no requirement that it be limited to UDP. We feel that ICP over UDP offers features important to Web caching applications. An ICP query/reply exchange needs to occur quickly, typically within a second or two. A cache cannot wait longer than that before beginning to retrieve an object. Failure to receive a reply message most likely means the network path is either congested or broken. In either case we would not want to select that neighbor. As an indication of immediate network conditions between neighbor caches, ICP over a lightweight protocol such as UDP is better than one with the overhead of TCP.

In addition to its use as an object location protocol, ICP messages can be used for cache selection. Failure to receive a reply from a cache may indicate a network or system failure. The ICP reply may include information that could assist selection of the most appropriate source from which to retrieve an object.

ICP Message Format

The ICP message format consists of a 20-octet fixed header plus a variable sized payload:



NOTE: All fields must be represented in network byte order.

One of the opcodes defined below.

The ICP protocol version number. At the time of this writing, both versions two and three are in use. This document describes only version two. The version number field allows for future development of this protocol.

The total length (octets) of the ICP message. ICP messages MUST not exceed 16,384 octets in length.

An opaque identifier. When responding to a query, this value must be copied into the reply message.

A 32-bit field of option flags that allows extension of this version of the protocol in certain, limited ways. See ``ICP Option Flags'' below.

A four-octet field to support optional features. The following ICP features make use of this field:

The IPv4 address of the host sending the ICP message. This field should probably not be trusted over what is provided by `getpeername()`, `accept()`, and `recvfrom()`. There is some ambiguity over the original purpose of this field. In practice it is not used.

The contents of the Payload field vary depending on the Opcode, but most often it contains a null-terminated URL string.

The following table shows currently defined ICP opcodes:

ICP_OP_INVALID

ICP_OP_QUERY

ICP_OP_QUERY payload format:

0										1										2										3									
0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1								
+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-								
										Requester Host Address																													
+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-								
										Null-Terminated URL																													
+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-	+	-								

In response to an ICP_OP_QUERY, the recipient must return one of: ICP_OP_HIT, ICP_OP_MISS, ICP_OP_ERR, ICP_OP_MISS_NOFETCH, ICP_OP_DENIED, or ICP_OP_HIT_OBJ.

ICP_OP_SECHO

Similar to ICP_OP_QUERY, but for use in simulating a query to an origin server. When ICP is used to select the closest neighbor, the origin server can be included in the algorithm by bouncing an ICP_OP_SECHO message off it's echo port. The payload is simply the null-terminated URL.

NOTE: the echo server will not interpret the data (i.e. we could send it anything). This opcode is used to tell the difference between a legitimate query or response, random garbage, and an echo response.

ICP_OP_DECHO

Similar to ICP_OP_QUERY, but for use in simulating a query to a cache which does not use ICP. When ICP is used to choose the closest neighbor, a non-ICP cache can be included in the algorithm by bouncing an ICP_OP_DECHO message off it's echo port. The payload is simply the null-terminated URL.

NOTE: one problem with this approach is that while a system's echo port may be functioning perfectly, the cache software may not be running at all.

One of the following six ICP opcodes are sent in response to an ICP_OP_QUERY message. Unless otherwise noted, the payload must be the null-terminated URL string. Both the URL string and the Request Number field must be exactly the same as from the ICP_OP_QUERY message.

ICP_OP_HIT

An ICP_OP_HIT response indicates that the requested URL exists in this cache and that the requester is allowed to retrieve it.

ICP_OP_MISS

An ICP_OP_MISS response indicates that the requested URL does not exist in this cache. The querying cache may still choose to fetch the URL from the replying cache.

ICP_OP_ERR

An ICP_OP_ERR response indicates some kind of error in parsing or handling the query message (e.g. invalid URL).

ICP_OP_MISS_NOFETCH

An ICP_OP_MISS_NOFETCH response indicates that this cache is up, but is in a state where it does not want to handle cache misses. An example of such a state is during a startup phase where a cache might be rebuilding its object store. A cache in such a mode may wish to return ICP_OP_HIT for cache hits, but not ICP_OP_MISS for misses. ICP_OP_MISS_NOFETCH essentially means ``I am up and running, but please don't fetch this URL from me now.''

Note, ICP_OP_MISS_NOFETCH has a different meaning than ICP_OP_MISS. The ICP_OP_MISS reply is an invitation to fetch the URL from the replying cache (if their relationship allows it), but ICP_OP_MISS_NOFETCH is a request to NOT fetch the URL from the replying cache.

ICP_OP_DENIED

An ICP_OP_DENIED response indicates that the querying site is not allowed to retrieve the named object from this cache. Caches and proxies may implement complex access controls. This reply must be interpreted to mean ``you are not allowed to request this particular URL from me at this particular time.''

Caches receiving a high percentage of ICP_OP_DENIED replies are probably misconfigured. Caches should track percentage of all replies which are ICP_OP_DENIED and disable a neighbor which exceeds a certain threshold (e.g. 95% of 100 or more queries).

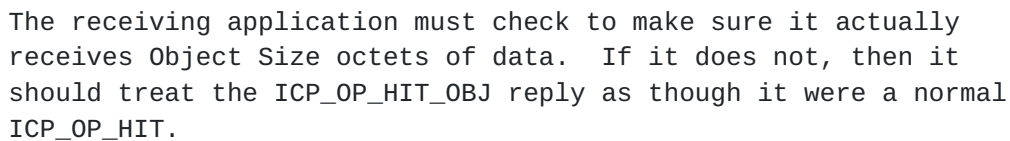
Similarly, a cache should track the percent of ICP_OP_DENIED messages that are sent to a given address. If the percent of denied messages exceeds a certain threshold (e.g. 95% of 100 or more), the cache may choose to ignore all subsequent ICP_OP_QUERY messages from that address until some sort of administrative intervention occurs.

ICP_OP_HIT_OBJ

Just like an ICP_OP_HIT response, but the actual object data has been included in this reply message. Many requested objects are small enough that it is possible to include them in the query response and avoid the need to make a subsequent HTTP request for the object.

CAVEAT: ICP_OP_HIT_OBJ has some negative side effects which make its use undesirable. It transfers object data without HTTP and therefore bypasses the standard HTTP processing, including authorization and age validation. Another negative side effect is that ICP_OP_HIT_OBJ messages will often be much larger than the path MTU, thereby causing fragmentation to occur on the UDP packet. For these reasons, use of ICP_OP_HIT_OBJ is NOT recommended.

ICP_OP_HIT_OBJ payload format:



UNRECOGNIZED OPCODES

3. ICP Option Flags

This flag is set in an ICP_OP_QUERY message indicating that it is okay to respond with an ICP_OP_HIT_OBJ message if the object data will fit in the reply.

0x40000000 ICP_FLAG_SRC_RTT

This flag is set in an ICP_OP_QUERY message indicating that the requester would like the ICP reply to include the responder's measured RTT to the origin server.

Upon receipt of an ICP_OP_QUERY with ICP_FLAG_SRC_RTT bit set, a cache should check an internal database of RTT measurements. If available, the RTT value MUST be expressed as a 16-bit integer, in units of milliseconds. If unavailable, the responder may either set the RTT value to zero, or clear the ICP_FLAG_SRC_RTT bit in the ICP reply. The ICP reply MUST not be delayed while waiting for the RTT measurement to occur.

This flag is set in an ICP reply message (ICP_OP_HIT, ICP_OP_MISS, ICP_OP_MISS_NOFETCH, or ICP_OP_HIT_OBJ) to indicate that the low 16-bits of the Option Data field contain the measured RTT to the host given in the requested URL. If ICP_FLAG_SRC_RTT is clear in the query then it MUST also be clear in the reply. If ICP_FLAG_SRC_RTT is set in the query, then it may or may not be set in the reply.

4. Security Considerations

The security issues relating to ICP are discussed in the companion document, RFCXXXX (<[draft-wessels-icp-v2-appl-00.txt](#)>).

5. References

- [1] Fielding, R., et. al, "Hypertext Transfer Protocol -- HTTP/1.1", [RFC 2068](#), UC Irvine, January 1997.
- [2] Berners-Lee, T., Masinter, L., and M. McCahill, "Uniform Resource Locators (URL)", [RFC 1738](#), CERN, Xerox PARC, University of Minnesota, December 1994.
- [3] Bowman M., Danzig P., Hardy D., Manber U., Schwartz M., and Wessels D., "The Harvest Information Discovery and Access System", Internet Research Task Force - Resource Discovery, <http://harvest.transarc.com/>.
- [4] Wessels D., Claffy K., "ICP and the Squid Web Cache", National

Laboratory for Applied Network Research, <http://www.nlanr.net/~wessels/Papers/icp-squid.ps.gz>

6. Acknowledgments

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