Network Working Group Internet-Draft Intended status: Standards Track Expires: April 1, 2018 X. Xu Huawei A. Bashandy Cisco H. Assarpour Broadcom S. Ma Juniper W. Henderickx Nokia J. Tantsura Individual September 28, 2017

Unified Source Routing Instructions using MPLS Label Stack draft-xu-mpls-unified-source-routing-instruction-04

Abstract

MPLS Segment Routing (SR-MPLS in short) is an MPLS data plane-based source routing paradigm in which a sender of a packet is allowed to partially or completely specify the route the packet takes through the network by imposing stacked MPLS labels to the packet. SR-MPLS could be leveraged to realize a unified source routing mechanism across MPLS, IPv4 and IPv6 data planes by using an MPLS label stack as a unified source routing instruction set while preserving backward compatibility with SR-MPLS.

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September 2017

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1. Introduction

MPLS Segment Routing (SR-MPLS in short)

[I-D.ietf-spring-segment-routing-mpls] is an MPLS data plane-based source routing paradigm in which a sender of a packet is allowed to partially or completely specify the route the packet takes through the network by imposing stacked MPLS labels to the packet. SR-MPLS could be leveraged to realize a unified source routing mechanism across MPLS, IPv4 and IPv6 data planes by using an MPLS label stack as a unified source routing instruction set while preserving backward compatibility with SR-MPLS. More specifically, the source routing instruction set information contained in a source routed packet could be uniformly encoded as an MPLS label stack no matter the underlay is IPv4, IPv6 or MPLS.

Although the source routing instructions are encoded as MPLS labels, this is a hardware convenience rather than an indication that the whole MPLS protocol stack and in particular the MPLS control protocols need to be deployed. Note that the complexity associated with the whole MPLS protocol stack is largely due to the complex control plane protocols.

<u>Section 3</u> describes various use cases for the unified source routing instruction mechanism and <u>Section 4</u> describes a typical application scenario and how the packet forwarding happens.

<u>1.1</u>. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in <u>RFC 2119</u> [<u>RFC2119</u>].

2. Terminology

This memo makes use of the terms defined in [<u>RFC3031</u>] and [<u>I-D.ietf-spring-segment-routing-mpls</u>].

3. Use Cases

The unified source routing mechanism across IPv4, IPv6 and MPLS is useful at least in the following use cases:

- o Incremental deployment of the SR-MPLS technology [<u>I-D.xu-mpls-spring-islands-connection-over-ip</u>]. Since there is no need to run any other label distribution protocol (e.g., LDP, see [<u>I-D.ietf-spring-segment-routing-ldp-interop</u>] for more details.) on those non-SR-MPLS routers for incremental deployment purposes, the network provisioning is greatly simplified, which is one of the major claimed benefits of the SR-MPLS technology (i.e., running a single protocol).
- o Overcome the load-balancing dilemma encountered by SR-MPLS. In fact, this unified source routing mechanism is even useful in a fully upgraded SR-MPLS network since the load-balancing dilemma encountered by SR-MPLS [I-D.ietf-mpls-spring-entropy-label] due to the maximum Readable Label-stack Depth (RLD) hardware limitation [I-D.ietf-ospf-mpls-elc] [I-D.ietf-isis-mpls-elc] [I-D.ietf-idr-bgp-ls-segment-routing-rld] and the Maximum SID Depth (MSD) hardware limitation [I-D.ietf-ospf-segment-routing-msd] [I-D.ietf-isis-segment-routing-msd] [I-D.ietf-idr-bgp-ls-segment-routing-msd]

encapsulation [RFC7510] where the source port of the UDP tunnel header is used as an entropy field.

- A poor man's light-weight alternative to SRv6
 [<u>I-D.ietf-6man-segment-routing-header</u>]. At least, it could be
 deployed as an interim until full featured SRv6 is available on
 more platforms. Since the Source Routing Header (SRH)
 [<u>I-D.ietf-6man-segment-routing-header</u>] consisting of an ordered
 list of 128-bit long IPv6 addresses is now replaced by an ordered
 list of 32-bit long label entries (i.e., label stack), the
 encapsulation overhead and forwarding performance issues
 associated with SRv6 are eliminated.
- A new IPv4 source routing mechanism which has overcome the security vulnerability issues associated with the traditional IPv4 source routing mechanism.
- o Traffic Engineering scenarios where only a few routers (e.g., the entry and exit nodes of each plane in the dual-plane network case or the egress node in the Egress Peer Engineering (EPE) case) are specified as segments of explicit paths. In this way, only a few routers are required to support the SR-MPLS capability while all the other routers just need to support IP forwarding capability, which would significantly reduce the deployment cost of the SR-MPLS technology.
- o MPLS-based Service Function Chaining (SFC) [I-D.xu-mpls-service-chaining]. Based on the unified source routing mechanism as described in this document, only SFC-related nodes including Service Function Forwarders (SFF), Service Functions (SF) and classifiers are required to recognize the SFC encapsulation header in the MPLS label stack form, while the intermediate routers just need to support vanilla IP forwarding (either IPv4 or IPv6). In other words, it undoubtedly complies with the transport-independence requirement for the SFC encapsulation header as listed in the SFC architecture document [RFC7665].

<u>4</u>. Packet Forwarding Procedures

The primary objective of this document is to describe how SR-MPLS capable routers and IP-only routers can seamlessly co-exist and interoperate. This section describes the forwarding information base (FIB) entry and the forwarding behavior that allow the deployment of SR-MPLS when some routers are IPv4 only or IPv6 only. Note that OSPF or ISIS is assumed to be enabled in the following examples as described in <u>Section 4.1</u> and 4.2, in fact, it's no doubt that BGP could be used as a replacement.

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<u>4.1</u>. Forwarding Entry Construction

This sub-section describes the how to construct the forwarding information base (FIB) entry on an SR-MPLS-capable router when some or all of the next-hops along the shortest path towards a prefix-SID are IPv4-only or IPv6-only routers. Consider the router "A" receiving a labeled packet whose top label L(E) corresponds to the prefix-SID is "SID(E)" of prefix "P(E)" advertised by the router "E". Suppose the ith next-hop router "NHi" along the shortest path from the router "A" towards the prefix-SID "SID(E)" is not SR-MPLS capable. That is both routers "A" and "E" are SR-MPLS capable but the next hop "NHi" along the shortest path from "A" to "E". The following applies:

- o It is assumed that the router "E" advertises the SR-Capabilities sub-TLV as described in and
 [<u>I-D.ietf-ospf-segment-routing-extensions</u>], which includes the SRGB because router "E" is SR-MPLS capabile.
- o The owning router "E" MUST advertise the encapsulation endpoint and the tunnel type using [<u>I-D.ietf-isis-encapsulation-cap</u>] and/or [<u>I-D.ietf-ospf-encapsulation-cap</u>].
- o If "A" and "E" are in different areas/levels, then
 - * The OSPF Tunnel Encapsulation TLV [<u>I-D.ietf-ospf-encapsulation-cap</u>] and/or the ISIS Tunnel Encapsulation sub-TLV [<u>I-D.ietf-isis-encapsulation-cap</u>] are flooded domain-wide.
 - * The OSPF SID/label range TLV [<u>I-D.ietf-ospf-segment-routing-extensions</u>] and the ISIS SR-Capabilities Sub-TLV [<u>I-D.ietf-isis-segment-routing-extensions</u>] are advertised domain-wide. This way router "A" knows the characteristics of the owning router "E".
 - * When the owning router "E" is running ISIS and advertises the prefix "P(E) ", the router "E" uses the extended reachability TLV (TLVs 135, 235, 236, 237) and associates the IPv4/IPv6 and/ or IPv4/IPv6 source router ID sub-TLV(s) [<u>RFC7794</u>].
 - * When the owning router "E" is running OSPF and advertises the prefix "P(E)", the router "E" uses the OSPFv2 Extended Prefix Opaque LSA [<u>RFC7684</u>] and sets the flooding scope to AS-wide.
 - * When the owning router "E" is running ISIS and advertises the ISIS capabilities TLV (TLV 242) [<u>RFC7981</u>], it must set the "router-ID" field to a valid value or include IPV6 TE router-

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ID sub-TLV (TLV 12), or do both. The "S" bit (flooding scope) of the ISIS capabilities TLV (TLV 242) MUST be set to "1".

- o Router "A" programs the FIB entry corresponding to the "SID(E)" as follows:
 - * If NP (OSPF) or P (ISIS) flag is clear,
 - *
- + pop the outer label.
- * If NP (OSPF) or P (ISIS) is set,
- *
- + the outer label is SID(E) plus the lower bound of the SRGB
 of "E".
- * Encapsulate the packet according to the encapsulation advertised in [<u>I-D.ietf-isis-encapsulation-cap</u>] or [<u>I-D.ietf-ospf-encapsulation-cap</u>].
- * Send the packet towards the next hop "NHi".

<u>4.2</u>. Packet Forwarding Procedures

++	++	++	++	++	
A ++ B ++ C ++ D ++ H					
++	++	+++	+++	· ++	
	++	++	+++		
E ++ F ++ G					
	++	++	++		
+	-+				
IP(A->E)				
+	-+	++			
L(G)		IP(E->G)			
+	-+	++		++	
L(H)		L(H)		IP(G->H)	
+	-+	++		++	
Packet	>	Packet	>	Packet	
+	-+	++		++	
Figure 1					

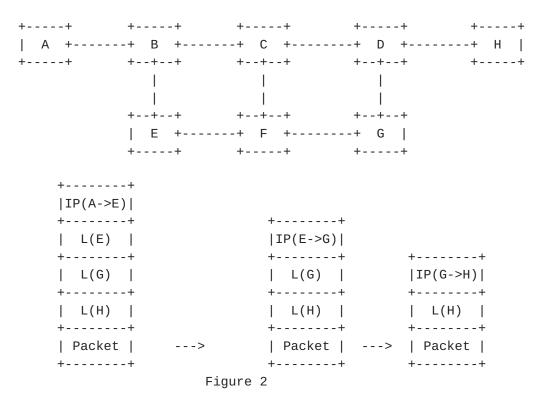
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As shown in Figure 1, Assume Router A, E, G and H are SR-MPLS-capable routers while the remaining are only capable of forwarding IP packets. Router A, E, G and H advertise their Segment Routing related information via IS-IS or OSPF. Now assume router A wants to send a given IP or MPLS packet via an explicit path of $\{E - >G - >H\}$, router A would impose an MPLS label stack corresponding to that explicit path on the received IP packet. Since there is no Label Switching Path (LSP) towards router E, router A would replace the top label indicating router E with an IP-based tunnel for MPLS (e.g., MPLS-over-UDP [RFC7510]) towards router E and then send it out. In other words, router A would pop the top label and then encapsulate the MPLS packet with an IP-based tunnel towards router E. When the IP-encapsulated MPLS packet arrives at router E, router E would strip the IP-based tunnel header and then process the decapsulated MPLS packet accordingly. Since there is no LSP towards router G which is indicated by the current top label of the decapsulated MPLS packet, router E would replace the current top label with an IP-based tunnel towards router G and send it out. When the packet arrives at router G, router G would strip the IP-based tunnel header and then process the decapsulated MPLS packet. Since there is no LSP towards router H, router G would replace the current top label with an IP-based tunnel towards router H. Now the packet encapsulated with the IPbased tunnel towards router H is exactly the original packet that router A had intended to send towards router H. If the packet is an MPLS packet, router G could use any IP-based tunnel for MPLS (e.g., MPLS-over-UDP [<u>RFC7510</u>]). If the packet is an IP packet, router G could use any IP tunnel for IP (e.g., IP-in-UDP [<u>I-D.xu-intarea-ip-in-udp</u>]). That original IP or MPLS packet would be forwarded towards router H via an IP-based tunnel. When the encapsulated packet arrives at router H, router H would decapsulate it into the original packet and then process it accordingly.

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Note that in the above description, it's assumed that the label associated with each prefix-SID advertised by the owner of the prefix-SID is a Penultimate Hop Popping (PHP) label (e.g., the NP-flag [<u>I-D.ietf-ospf-segment-routing-extensions</u>] associated with the corresponding prefix SID is not set).

Figure 2 demostrates the packet walk in the case where the label associated with each prefix-SID advertised by the owner of the prefix-SID is not a Penultimate Hop Popping (PHP) label (e.g., the NP-flag [<u>I-D.ietf-ospf-segment-routing-extensions</u>] associated with the corresponding prefix SID is set).



Although the above description is based on the use of prefix-SIDs, the unified source routing instruction approach is actually applicable to the use of adj-SIDs as well. For instance, when the top label of a received MPLS packet indicates an given adj-SID and the corresponding adjacent node to that adj-SID is not MPLS-capable, the top label would be replaced by an IP-based tunnel towards that adjacent node and then forwarded over the correponding link indicated by that adj-SID.

When encapsulating an MPLS packet with an IP-based tunnel header (e.g., a UDP header as per [<u>RFC7510</u>]), the corresponding entropy field (i.e., the source port in the MPLS-in-UDP case) should be filled with an entropy value that is generated by the encapsulator to uniquely identify a flow. However, what constitutes a flow is locally determined by the encapsulator. For instance, if the MPLS label stack contains at least one entropy label and the encapsulator is capable of reading that entropy label, the entropy label value could be directly copied to the entropy field (e.g., the source port of the UDP header). Otherwise, the encapsulator may have to perform a hash on the whole label stack or the five-tuple of the MPLS payload if the payload is determined as an IP packet. To avoid re-performing hash on the whole packet when re-encapsulating the packet with an IPbased tunnel header (e.g., a UDP tunnel header), especially when the encapsulator could not obtain at least one entropy label due to some reasons (e.g., 1) there is no EL at all in the label stack; 2) the encapsulator couldn't recognize the ELI; 3) the encapsulator could

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not read the EL due to the RLD limit), it's RECOMMENDED that the entropy value contained in the packet (e.g., the UDP source port value) is kept when stripping the IP-based tunnel header (e.g., the UDP tunnel header). As such, the entropy value could be directly copied to the entropy field (e.g., the source port of the UDP tunnel header) when re-encapsulating the packet with an IP-based tunnel header (e.g., a UDP tunnel header). As such, the load-balancing dilemma encountered by SR-MPLS as described in [I-D.ietf-mpls-spring-entropy-label] due to the maximum Readable Label-stack Depth (RLD) hardware limitation [I-D.ietf-ospf-mpls-elc] [I-D.ietf-isis-mpls-elc] and the Maximum SID Depth (MSD) hardware limitation [I-D.ietf-ospf-segment-routing-msd] [I-D.ietf-isis-segment-routing mechanism is even useful in a fully upgraded SR-MPLS network environment.

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6. Acknowledgements

Thanks Joel Halpern, Bruno Decraene, Loa Andersson and Stewart Bryant for their insightful comments on this document.

7. IANA Considerations

No IANA action is required.

8. Security Considerations

TBD.

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