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A Path-Aware Scheduling Scheme for Multipath Transport Protocols draft-zuo-mptcp-scheduler-01.txt

Abstract

Scheduling schemes play a key role in the performance of multipath transport protocols. A path-aware scheduling scheme for latencysensitive applications is proposed, which always schedules data in the path with the lowest One-Way Latency (OWL). The OWLs of all paths are obtained by employing the redundant transmission periodically.

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<u>1</u>. Introduction

For latency-sensitive applications, such as games, latency is the primary goal to make the interactive data exchanged in time. A suitable scheduling scheme will influence the latency performance of multipath transport protocols [RFC6824], since the packets with consecutive sequence numbers may arrive at the receiver out of order due to the different RTTs of multiple paths. Lots of scheduling schemes are proposed to select the path with the lowest latency to send data. RTT is commonly used as the condition for scheduling. However, it's found that the delays of the forward path and the reverse path are usually different in many cases, such as the asymmetric uplink delay and the downlink delay in wireless environment. Especially, the congestion may not happen in the forward path and the reverse path at the same time, resulting in the different queue delays. Therefore, scheduling data on the path with the lowest RTT cannot guarantee the lowest interactive latency required by the latency-sensitive applications.

A new path-aware scheduling scheme for multiple transport protocols

is proposed, which considers the One-Way Latency (OWL) as a scheduling condition. The data will be always sent in the path with the lowest OWL.

It is intended that the scheduling scheme presented in this draft can be applied to other multipath transport protocols, such as alternative multipath extensions to TCP [<u>RFC793</u>], UDP, QUIC, or indeed any other multipath protocols. However, for the purposes of example, this document will, where appropriate, refer to the MPTCP [<u>RFC6182</u>].

2. Acronyms and Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in <u>RFC 2119</u> [<u>RFC2119</u>].

3. A New Path-Aware Scheduling Scheme

A new path-aware scheduling scheme is proposed to always send data in the path with the lowest OWL.

3.1 The definition of OWL

OWL is defined as the latency of a data successfully delivered from the sender to the receiver, which is calculated as

OWL (i) = T_recv (i) - T_send (i),

where i means the i-th path, T_{recv} is the time when the data arrives at the receiver and T_{send} is the time when the data is sent from the sender.

Compared to the One-Way Delay (OWD) defined in [<u>RFC7679</u>], OWL can be deemed as

OWL (i) = OWD (i) + dT,

where dT means the time difference caused by the absolute clock time at the sender and the receiver. When the scheduler tries to select the path with the lowest latency to send data, it will compare the latencies of multiple paths. Take two paths for example,

dOWL = OWL (1) - OWL (2),

where dOWL indicates the OWL difference of path 1 and path 2. If dOWL > 0, the path 2 is selected or else the path 1 is selected. We could see that dT will not influence the value of dOWL, since it will be

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subtracted when path 1 and path 2 are connected between the same sender and the same receiver. Hence, OWL is more suitable than OWD as the scheduling condition in multipath environment, since it does not need to consider the issue of time synchronization.

3.2. The Operations of the Path-Aware Scheduling Scheme

<u>3.2.1</u>. Initialization

As in <u>RFC6824</u>, MPTCP first sets up the connection in the primary path and then subsequentially sets up the connections of other paths. After the handshake stage of the primary path, data is immediately sent in this path. Then when the connection of the second path is built up, the OWLs of these two paths could be obtained. The following data could be scheduled in the path with the lowest OWL. However, network is unstable at the moment, the value of the obtained OWL may change in the next round of data transmission, especially when there are other flows co-existed or packet loss is very high. For avoiding the frequent switching amongst the multiple paths due to the variant instant OWLs, the data are scheduled redundantly in the multiple paths for a period of time [CONEXT17]. The redundant scheduling means the same data are sent in all paths concurrently [MPTCP.org]. A Smoothed OWL (SOWL) is defined as the effective OWL of a path, as referred to the smoothed RTT, which is defined in [<u>RFC793</u>]:

SOWL = (ALPHA * SOWL) + ((1-ALPHA) * OWL).

Alternatively, other methods which could calculate the effective OWL during the period of time can be considered. The period of time could be set to a fixed number, e.g. 1s, or max{200ms, SRTT}.

<u>3.2.2</u>. Data scheduling

After the effective OWLs of all paths are obtained, the sender will schedule the following data from application to the path with the lowest OWL.

<u>3.2.3</u>. Periodically Redundant Transmission

Once the path with the lowest OWL is selected to send the data henceforth, and no data will be scheduled to the other paths. The result is that there is no chance to know the future OWLs of the other paths. For making the scheduling more responsive to the variant network environment and always sending data in the path with the lowest OWL, sending probe packets in other paths may be one of the potential solutions. However, it introduces extra packets which

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10s

consume extra bandwidth of the path. Therefore, the periodically redundant transmission is proposed, which sends data redundantly for obtaining the effective OWLs of all paths every a certain period of time (e.g. 10s). Figure 1 shows the process of the periodically redundant transmission.

1s 10s 1s

| transmission in | | transmission in | |-redundant -|-the path with -|-redundant -|-the path with -|-... | transmission| the lowest OWL | transmission| the lowest OWL |

Figure 1: Periodically redundant transmission.

On the one hand, the periodically redundant transmission could obtain the OWLs of all paths at the same time without introducing extra packets; on the other hand, the same data are transmitted concurrently in all the paths, which could guarantee the lowest delivering time as long as the data arrives at the receiver no matter in which path it is sent.

Once the OWLs of all paths are updated, the following data are scheduled as described in <u>Section 3.2.2</u>. During the data transmission in the path with the lowest OWL, the OWL of the selected path is continually updated. Once the updated OWL has increased so much that it may not the lowest OWL any more, the redundant transmission is immediately activated to obtain the updated OWLs of all paths, as shown in Figure 2.

 1s
 10s
 1s
 10s

 | transmission in |
 | transmission in |
 |

 |-redundant -|-the path with -|-redundant -|-the path with -|-...
 | transmission| the lowest OWL |

 | transmission| the lowest OWL | transmission| the lowest OWL |
 |

 | the OWL of the
 |

 selected path is not
 |

 the lowest any more
 |

Figure 2: Redundant transmission activated when the OWL of the selected path is not the lowest any more.

<u>4</u>. Implementation Considerations

4.1 Packet Format

As described in <u>Section 3.1</u>, the calculation of the OWL requires the time when the data is sent from the sender and the time when the data

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arrives at the receiver. Since the sender needs to compare the OWLs of all the paths, the sender is assigned executing the OWL calculation in this draft. A new MPTCP option, called as MP_OWL Option, is defined to carry the timestamp of data receiving at the receiver. The format of MP_OWL option is shown in Figure 3, where the value of 'Kind' is '30' (decimal) [RFC6824]. 1 2 3

0 1 2 3 4 5 6	78901234	6789012345678901				
+ Kind	+ L Lenath	++				
+	+	++				
Timestamp of receiving						

Figure 3: MP_OWL Option

The value of 'subtype' in MP_OWL option could be assigned as '0x8', since '0x0' through '0x7' have been used [RFC5226]. Table 1 gives the explanation of the value of '0x8'.

+	+ e +	Symbol		Name	++
0x8 +		MP_OWL		One-Way Latency	 ++

Table 1: The Subtype of MP_OWL option

4.2 Negotiation

The sender and the receiver need to negotiate with each other to see whether MP_OWL option is supported or not.

The details are to be defined.

4.3 The Steps for OWL Calculation

1) The sender sends each data and remember the sequence number as well as the sending time T_send of the data;

2) When the receiver receives the data, it responses an ACK with an MP_OWL option added, which carries the receiving time T_recv of the data;

3) When ACK gets back to the sender, the sender fetches the receiving time T_recv of the data from the ACK and subtracts the sending time T_send it has remembered to get the value of OWL, where $OWL = T_recv - T_send$.

<u>5</u>. Security Considerations

TBD.

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