

draft-ietf-6tisch-minimal-security

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6TiSCH - IETF 100 - Singapore



Status

- News
 - draft-ietf-6tisch-minimal-security-04
 - Published on 30 Oct 2017
- Relies completely on PSKs
- Summary of updates in -04

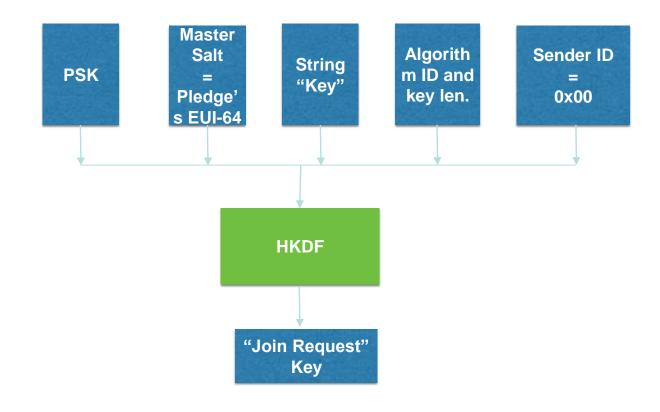


Update #1: Key/Nonce Derivation

- OSCORE-06 (formerly known as OSCOAP) updated the key/nonce derivation process
 - Same nonce used for both request and response but under a different key
- We could no longer use "EUI-64 | 0x00" and "EUI-64 | 0x01" as identifiers for the pledge and the JRC
- We now use EUI-64 of the pledge as Master Salt during key derivation and transport it as Context Hint
 - Sender ID of each pledge is 0x00; ID of the JRC is 0x01

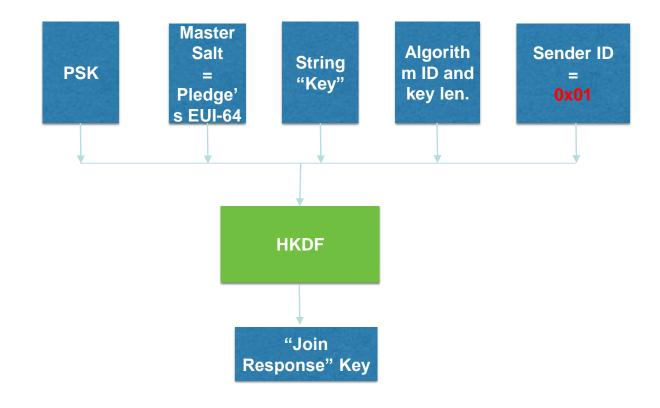
Derivation of Key Used to Protect Join Request





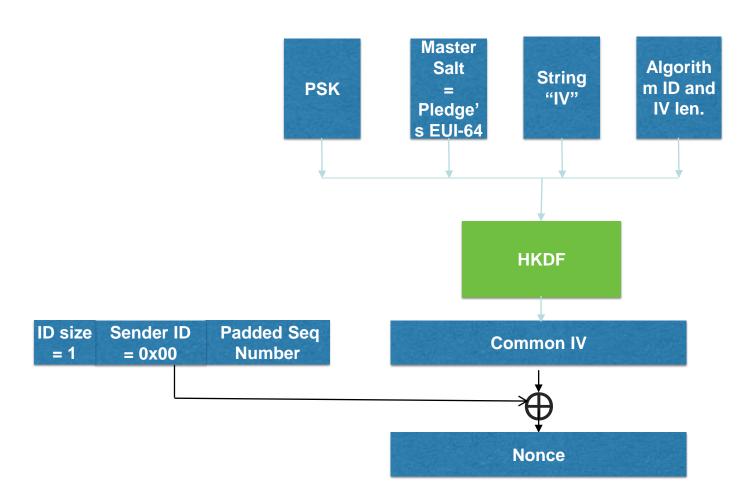
Derivation of Key Used to Protect Join Response







Nonce Derivation (used both for Join Request and Response)





Update #2: Error Handling

- Error handling in -03 opens the pledge to a DoS attack
 - Attacker could send (unprotected) error messages and force the pledge to attempt joining the next advertised network
- Solution in -04:
 - Using Non-Confirmable CoAP msg for Join Request will make OSCORE at JRC silently drop the request in case of failure (decryption, replay, unauthorized)
 - The pledge MUST silently discard any response not protected with OSCORE, including error codes.
 - Forces the pledge to implement a retransmission mechanism at the APP layer duplicating CoAP Confirmable msg functionality

Update #3: Join Request Retransmissions

- Binary exponential back-off mechanism to be implemented by the pledge at the APP layer specified in -04:
 - Super simple, inspired by the one in RFC7252 (CoAP)
 - Pledge keeps track of timeout and retransmission_counter
 - Parameters: TIMEOUT, TIMEOUT_RANDOM_FACTOR, MAX_RETRANSMIT
 - If the retransmission counter reaches MAX_RETRANSMIT on a timeout, the pledge SHOULD attempt to join the next advertised 6TiSCH network.

Name	Default Value	<u>+</u> 	1st attempt:	timeout in [10s, 15s]
TIMEOUT	10 s		2 nd attempt:	timeout in [20s, 30s]
TIMEOUT_RANDOM_FACTOR	1.5		3 rd attempt:	timeout in [40s, 60s]
MAX_RETRANSMIT	4		4th attempt:	timeout in [80s, 120s]



Misc updates

- Recommendation to store untrusted neighbor entries in a separate cache
- Join Request switched from GET -> POST to be more flexible with payload
- Added requirement on persistency of mutable OSCORE context parameters
 - Prevents nonce reuse and replay attacks across reboots
- Extensive editorial pass
 - Rewrote intro, clarifications on the PSK, etc...



Conclusion

- minimal-security-04 relies completely on PSKs
- Tracking of OSCORE, updates to error handling, editorial
- Open issue:
 - Join traffic, potentially controlled by the attacker, can influence SF to trigger 6P commands
 - In minimal-security, we recommend bandwidth cap at Join Proxy but this does not completely solve the problem
 - Should each SF specify how it handles the join traffic?
 - Proposed Resolution: Tag join packets at JP by using ToS bits in IPv6
- Reviews welcome!