

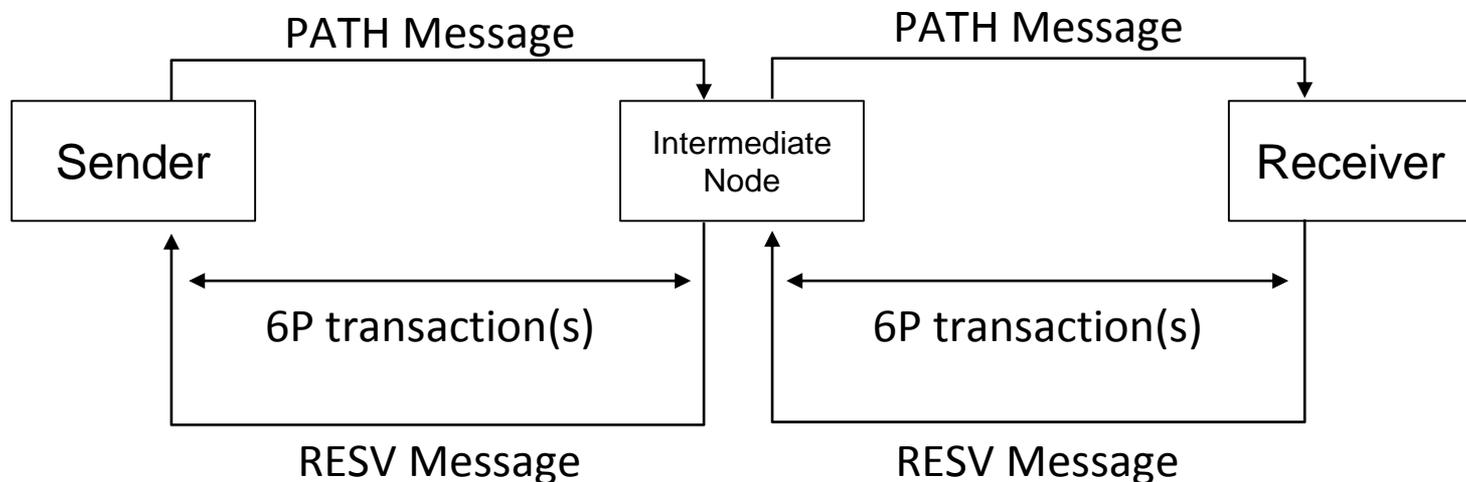
# draft-satish-6tisch-6top-sf1-04

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# Introduction to SF1

- Objectives
  - Reserve a track to a destination multiple hops away
  - Fulfill bandwidth and QoS (e.g., priority, time-critical) requirements
  - Hop-by-hop, distributed
- Combination of RSVP-TE and 6top protocol (6P)
  - RSVP-TE: end-to-end resource reservation
  - 6P: cell negotiation between two neighbors
- When to trigger SF1?
  - The Sender has an outgoing bandwidth requirement for a **new instance** to transmit data to Receiver.
  - The Sender has a **new outgoing bandwidth requirement** for an existing instance to transmit data to Receiver.

# End-to-end operation



- Step 1: The sender sends a PATH message downstream hop-by-hop to the receiver.
- Step 2: The receiver initiates a 6P transaction to verify if there are enough cells to fulfill the requirements. If so, the receiver sends the RESV message to reserve these cells.
- Step 3: Upon receiving a RESV message, an intermediate node reserves the cells between its upstream neighbor in the same way as in Step 2.
- Step 4: When RESV arrives at the sender before end-to-end timeout, a track from sender to receiver is built.

# PATH and RESV messages

```

<PATH Message> ::=  <Common Header> [ ... ]
                    <SESSION> <RSVP_HOP>
                    <TIME_VALUES>
                    [ ... ]
                    <LABEL_REQUEST>
                    [ <SF1 OPERATION REQUEST> ]
                    [ <6P OPERATION REQUEST> ]
                    [ ... ]
                    <sender descriptor>

<sender descriptor> ::= <SENDER_TEMPLATE> <SENDER_TSPEC>
                       [ <ADSPEC> ]
                       [ <RECORD_ROUTE> ]
  
```

- PATH: describe the flow (SENDER\_TSPEC) and collect path properties (ADSPEC).
- The three request objects: used to verify if the nodes along the route have the requested capabilities.
- The LABEL\_REQUEST is set to TSC (timeslot switching capability).

```

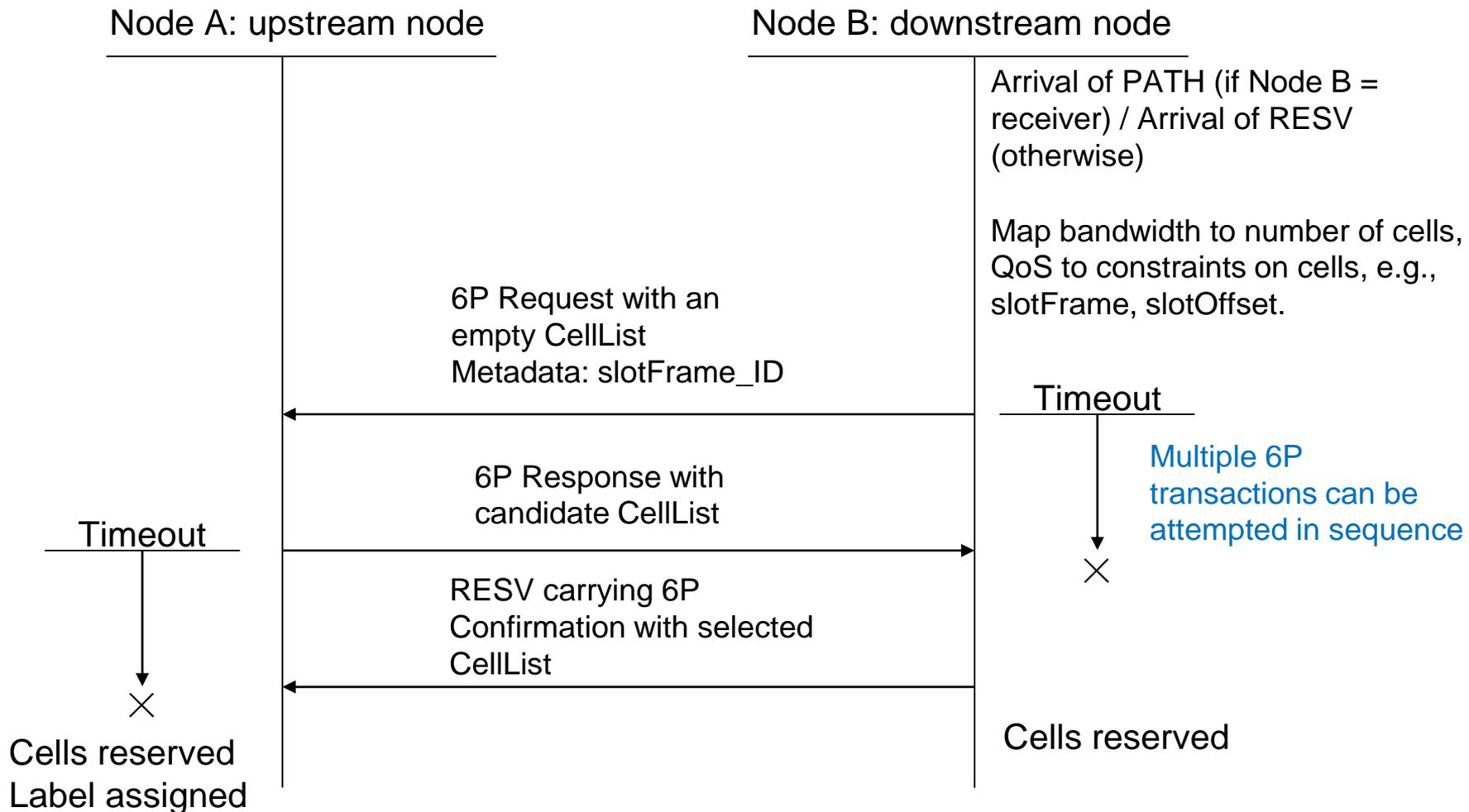
<RESV Message> ::= <Common Header> [ ... ]
                  <SESSION> <RSVP_HOP>
                  <TIME_VALUES>
                  [ <6P OPERATION> ]
                  [ ... ]
                  <STYLE> <flow descriptor list>

<flow descriptor list> ::= <FF flow descriptor>
<FF flow descriptor> ::= [ <FLOWSPEC> ] <FILTER_SPEC>
                       <LABEL> [ <RECORD_ROUTE> ]
  
```

- RESV: describe the bandwidth and QoS requirements (FLOWSPEC), assign label to the upstream node (LABEL).
- Parameters in FLOWSPEC are calculated according to SENDER\_TSPEC and ADSPEC.
- In case of 3-step transaction, 6P confirmation is encapsulated in 6P OPERATION.

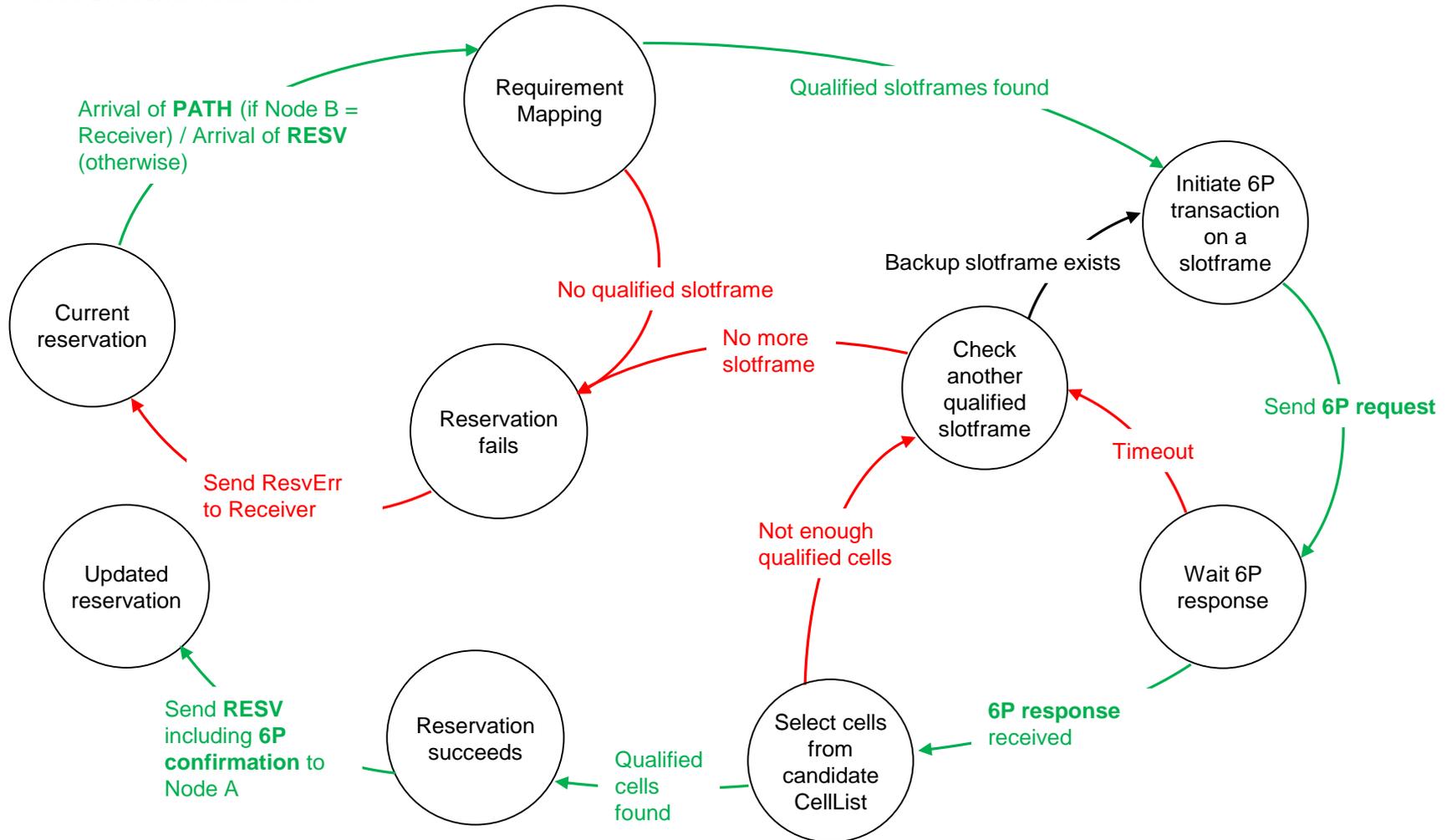
# One-hop operation using 3-step trans. (updated)

Note that the 2-step trans. can also be used.



# Node B State Transition Diagram

Node A: upstream node  
Node B: downstream node



# TrackID, Label and Bundle

- TrackID in SF1 (updated definition)
  - 16 bits identifier, assigned by the sender
  - Mapped from (source & destination IP address, RPLInstance)
  - Encapsulated in SENDER\_TEMPLATE of PATH and FILTER\_SPEC of RESV
- Label (updated definition)
  - 32 bits, mapped from a bundle between two neighbors
  - Encapsulated in the LABEL object of RESV
  - Locally valid between two neighbors, assigned by the downstream node
  - Associated to a track
- Bundle between two neighbors
  - A group of equivalent scheduled cells (slotFrame\_ID, CellList)



# Next steps

- Complete the definition
  - 6P request, SF1 request, 6P operation, Teardown message, etc.
- Mapping the traffic requirements to cells
- As suggested in 6top, the following requirements need to be covered
  - Error Handling (more detailed error code)
  - Specify the SF behavior of a node when it boots
  - Security considerations
  - Examples
- Implementation: simulation and hardware deployment

Thank you!

# Track Forwarding

- Not in the scope of SF1, just for the completeness of the story
- **A track can be seen as an LSP using bundles as implicit labels**
- The sender identifies which track a packet should follow based on “sender/receiver IP address, RPLInstance”.
- Then G-MPLS is used
  - The sender pushes the first label to the packet.
  - The label is swapped at each intermediate node
  - The label is popped out at the penultimate Hop or at the receiver
  - At each hop, the packet is forwarded using the bundle associated to the label.

