

Inter-domain DDoS mitigations: potentials, challenges, and solutions

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Large-scale *volumetric* attacks are common and often congest *more than one* networks

- Escalation in *volume* of DDoS attack traffic; e.g., 300 Gbps in 2013 – 1.2 Tbps in 2016.
- Volumetric attacks often flood *upstream* autonomous systems (ASes) [WISR'16]
- Advanced *link-flooding* attacks *congest multiple* ASes concurrently



[WISR'16] Arbor Networks, "Worldwide Infrastructure Security Report: Volume XI," Arbor Special Report, 2016.

Inter-domain DDoS mitigation becomes *necessary* for large-scale volumetric attacks

- Inter-domain packet filtering is often necessary:
 - e.g., inter-AS links are flooded, large portion of AS is flooded
- **Packet-filtering outsourcing**: an AS asks another AS for packet filtering
- State-of-the-art:
 - AT&T and CenturyLink*: *automated* packet-filtering outsourcing between two ASes for DDoS mitigation
 - **IETF DOTS**: standardization effort for common channels for inter-domain coordination

(*) N. Levy, D. Smith, and J. Schiel, "Operationalizing ISP cooperation during DDoS attacks," in NANOG 71, Oct 3, 2017.

Holy grail of inter-domain mitigation: source-end filtering



- Ideal DDoS mitigation: stops attacks earlier by outsourcing filtering to source AS (e.g., D-WARD, StopIt)
- Advantages
 - *Reduction of bandwidth* waste (thus *cost* saving)
 - Resource-demanding filtering operations
 - Local contexts may be utilized (e.g., list of usual suspects)

Yet, source-end filtering has **not** been deployed due to the **lack of trust** between ASes



- *Mutually untrusted* ASes may launch attacks
 - Source AS can *modify* or *leak* the requested policy
- No strong incentives for source AS
 - Filtering incurs *non-negligible cost*.
- *Risk* of dropping packets
 - Source AS may be *blamed* for dropping its own customer packets

From *outsourcing* to *collaboration*: source and destination *collaboratively* determine filtering policy



- *Source* AS can also *express its own policy* for its customers (e.g., preference, black/white lists)
- Security concerns:
 - 1) how to **guarantee fair policy** composition?
 - 2) how to protect the **sensitive filtering policies**?
 - 3) what if source AS **bypasses** the packet filtering?

Desired property 1)

Fair composition must be *verified by both* source and destination ASes

- Collaboration platform requires fair policy composition and enforcement
- *Fair*: we find a *middle ground* of two policies, favoring neither of the policies
 - Example: ASes express flow preferences



 Two ASes should be able to *verify* the fair composition and enforcement of their policies

Desired property 2)

Filtering policy of each AS must be *protected* from each other with privacy guarantees

- Filtering policies are inherently *sensitive*; e.g.,
 - Preference due to *private contracts, proprietary algorithms*
 - Attack information, vulnerable points of the network
 - Internal white/black policies
- Two ASes should be able to <u>negotiate</u> and determine the <u>guaranteed</u> degree of <u>privacy</u>

Desired property 3)

Filtering operations must *not* be *bypassed* for *any packet* from source to destination

- Source AS can evade the filtering when it wishes to ignore the composed filtering policy
- <u>Non-bypassability</u>: filtering operation *must be invoked* for all packets from source AS to destination AS

Middlebox-based filtering: a practical design choice for *rapid* and *widespread* deployment



- Commodity hardware for multi-Giga-bps throughput
- Trusted execution environment (TEE) capabilities (e.g., memory isolation, remote attestation) from commodity CPUs
 > verifiable control-/data-plane operations

TEE-based **middlebox** can satisfy three desired **security properties**

- Policy composition and enforcement are isolated and verifiable via remote attestation
 - fair packet filtering is guaranteed
- Two ASes can *negotiate* the desired level of *privacy* and policy *fairness*
 - policy inference attack: ~O(log(N))
 - tradeoff: degree of anonymity vs. fairness
- Bypass is immediately detected by efficient sketch and MAC operations
 - only 5-tuple information copied to TEE
 => multi-Giga-bps performance







Preliminary results:

multi-Giga-bps filtering with Intel SGX platform

- Filtering up to ~1.8 Gbps with 3 Intel SGX cores
 - CPU: Intel[®] Core[™] i5-6400
 - Memory: 8 GB (128 MB reserved for EPC)
 - NIC: Intel[®] 10-Gigabit X540-AT2
- *TCB* 1,369 SLoC, 1.9 MB binary.
- SGX-integration of **DPDK**
 - libraries ported to SGX: mempool, mbuf, ring, sched
- Plan: *scale out* with load balanced parallel middleboxes for *10 Gbps* or higher throughput

Conclusion

Collaborative source-end packet filtering

Potentials:

Collaborative source-end filtering is an ideal defense for everincreasing volumetric attacks

Challenges:

Lack of trust between ASes makes existing solutions impractical

Solutions:

TEE-based middlebox solution can offer **three security properties** (i.e., verifiable fairness, privacy, non-bypassability) necessary for secure and practical collaborative DDoS solution



We are open for *feedback* and *collaboration*

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