



KANGAROOTWELVE

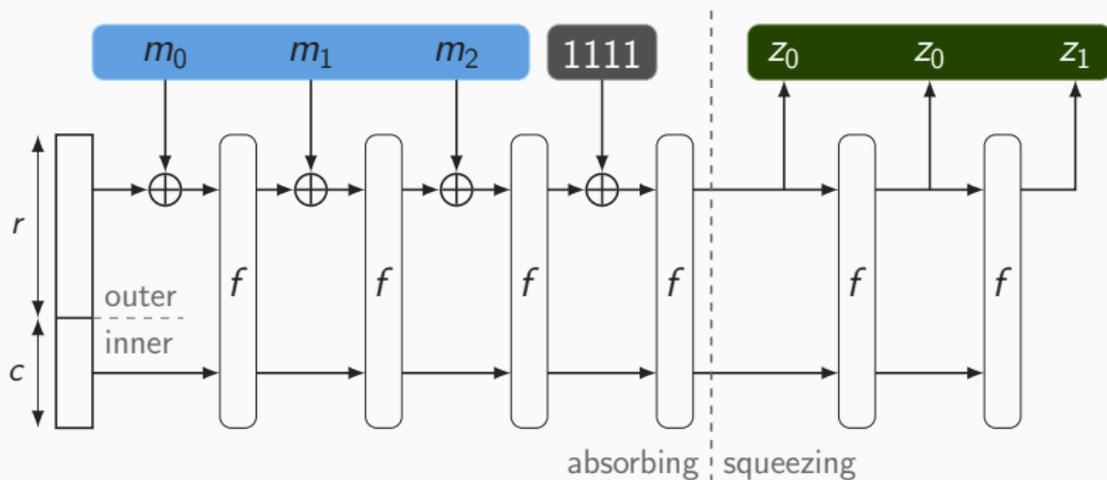
draft-viguiet-kangarootwelve-01

Benoît Viguiet¹

CFRG Meeting, March 19, 2018

¹Radboud University, Nijmegen, The Netherlands

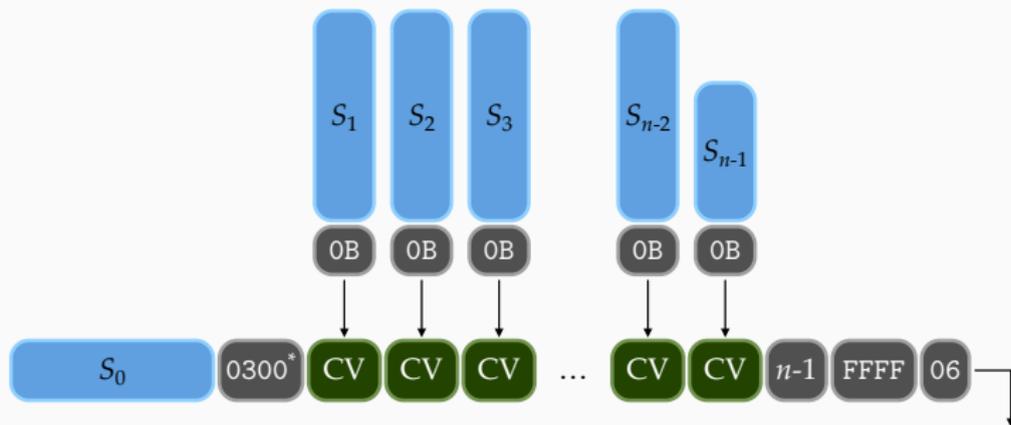
What is KANGAROOTWELVE?



► SHAKE128

- eXtendable Output Function
- Sponge construction
- Uses KECCAK- $p[1600, n_r = 24]$
- BUT no parallelism

What is KANGAROOTWELVE?



► KangarooTwelve

- eXtendable Output Function
- Tree on top of sponge construction
- KECCAK- p reduced from 24 to 12 rounds
- Parallelism grows automatically with input size
- No penalty for short messages

How secure is KANGAROOTWELVE?

- ▶ Same security claim as SHAKE128: 128 bits of security
- ▶ Sponge generic security

[EuroCrypt 2008] – On the Indifferentiability of the Sponge Construction

- ▶ Parallel mode with proven generic security

[IJIS 2014] – Sufficient conditions for sound tree and sequential hashing modes
[ACNS 2014] – Sakura: A Flexible Coding for Tree Hashing

- ▶ Sponge function on top of KECCAK- p [1600, $n_T = 12$]
 - Round function unchanged
 - ⇒ cryptanalysis since 2008 still valid
 - Safety margin: from *rock-solid* to *comfortable*

Status of KECCAK cryptanalysis

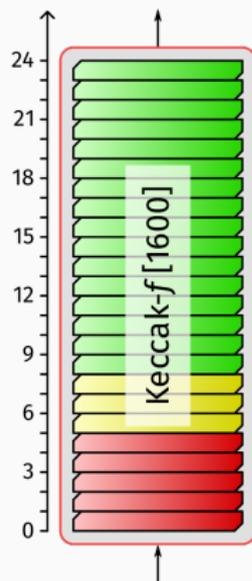
- ▶ Collision attacks up to 5 rounds
 - Also up to 6 rounds, but for non-standard parameters ($c = 160$)

[Song, Liao, Guo, CRYPTO 2017]

- ▶ Stream prediction
 - in 8 rounds (2^{128} time, prob. 1)
 - in 9 rounds (2^{256} time, prob. 1)

[Dinur, Morawiecki, Pieprzyk, Srebrny, Straus, EUROCRYPT 2015]

- ▶ Lots of third party cryptanalysis available at:
https://keccak.team/third_party.html



How fast is KANGAROOTWELVE?

- ▶ At least twice as fast as SHAKE128 on short inputs
- ▶ Much faster when parallelism is exploited on long inputs

	Short input	Long input
Intel Core i5-4570 (Haswell)	3.68 c/b	1.44 c/b
Intel Core i5-6500 (Skylake)	2.89 c/b	1.22 c/b
Intel Core i7-7800X (Skylake-X)	2.35 c/b	0.55 c/b

Single core only.

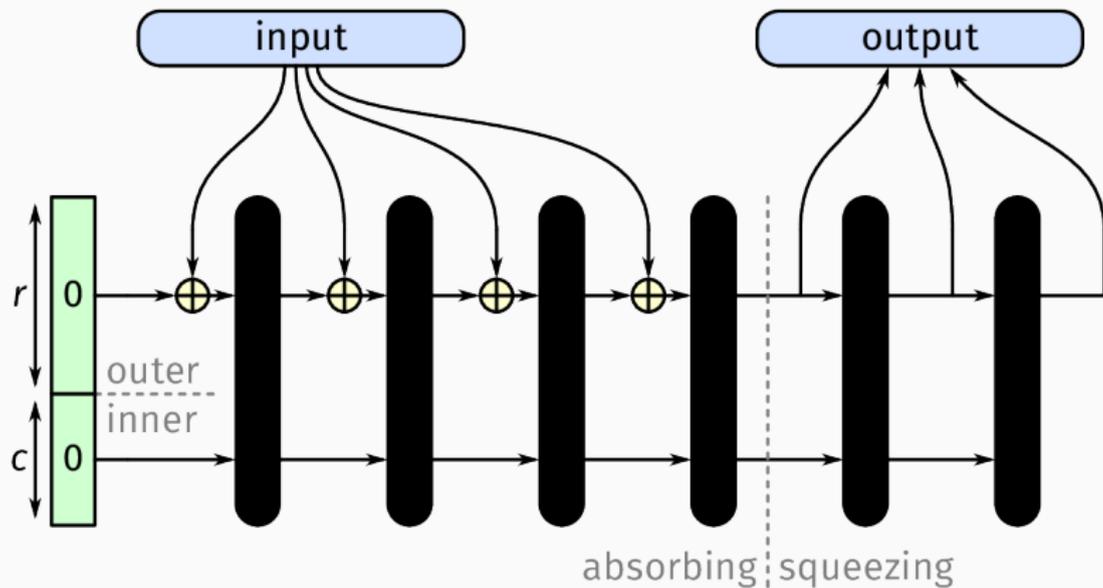


Why is it interesting for the IETF?

- ▶ KECCAK/KANGAROOTWELVE is an open design
 - Public design rationale
 - Result of an open international competition
 - Long-standing active scrutiny from the crypto community
- ▶ Best security/speed trade-off
 - Speed-up w/o wasting cryptanalysis resources (no tweaks)
 - Proven generic security
- ▶ Scalable parallelism
 - As much parallelism as the implementation can exploit
 - Without parameter

[https://tools.ietf.org/html/
draft-viguier-kangarootwelve-01](https://tools.ietf.org/html/draft-viguier-kangarootwelve-01)

Analyzing the sponge construction



Generic security of the sponge construction

Theorem 2. *A padded sponge construction calling a random permutation, $\mathcal{S}'[\mathcal{F}]$, is (t_D, t_S, N, ϵ) -indistinguishable from a random oracle, for any $t_D, t_S = O(N^2)$, $N < 2^c$ and for any ϵ with $\epsilon > f_P(N)$.*

If N is significantly smaller than 2^c , $f_P(N)$ can be approximated closely by:

$$f_P(N) \approx 1 - e^{-\frac{(1-2^{-r})N^2 + (1+2^{-r})N}{2^{c+1}}} < \frac{(1-2^{-r})N^2 + (1+2^{-r})N}{2^{c+1}}. \quad (6)$$

[EuroCrypt 2008]

<http://sponge.noekeon.org/SpongeIndifferentiability.pdf>

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Theorem, explained

$$\Pr[\text{attack}] \leq \frac{N^2}{2^{c+1}} \text{ (or so)}$$

⇒ if $N \ll 2^{c/2}$, then the probability is negligible

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 - Strong mathematical proofs

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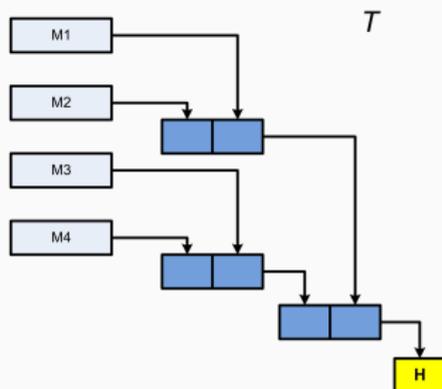
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 - ⇒ scope of cryptanalysis reduced to primitive
- ▶ Security of the primitive
 - No proof!
 - ⇒ open design rationale
 - ⇒ lots of third-party **cryptanalysis!**
 - Confidence
 - ⇐ sustained cryptanalysis activity and no break
 - ⇐ proven properties

Impact of parallelism

$\text{KECCAK-}f[1600] \times 1$	1070 cycles
$\text{KECCAK-}f[1600] \times 2$	1360 cycles
$\text{KECCAK-}f[1600] \times 4$	1410 cycles

CPU: Intel Core i5-6500 (Skylake) with AVX2 256-bit SIMD

Tree hashing



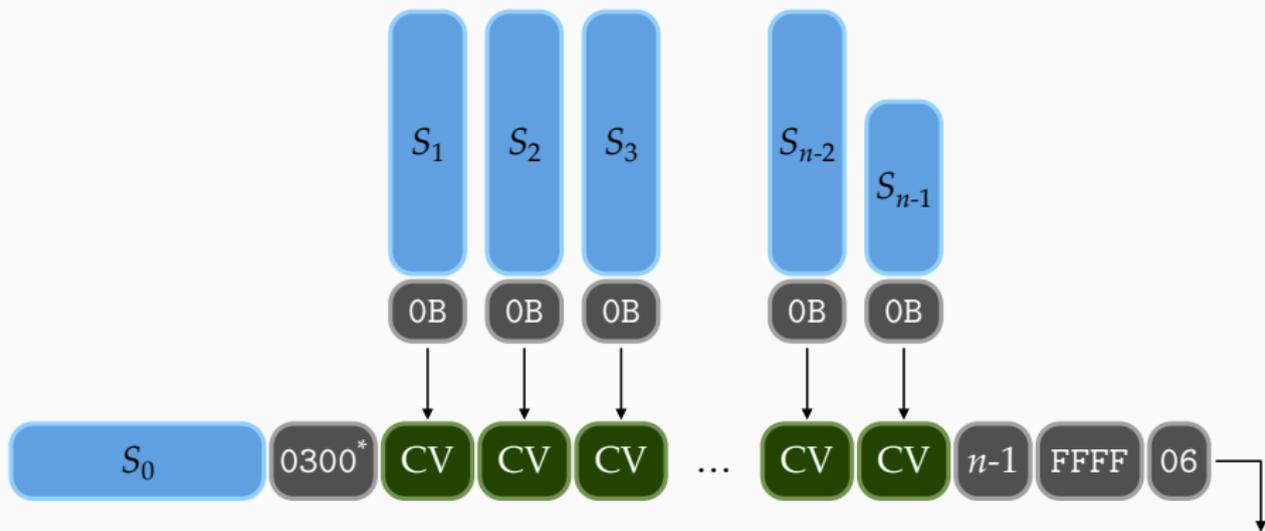
Example: **ParallelHash** [SP 800-185]

function	instruction set	cycles/byte ¹
$\text{KECCAK}[c = 256] \times 1$	x86_64	6.29
$\text{KECCAK}[c = 256] \times 2$	AVX2	4.32
$\text{KECCAK}[c = 256] \times 4$	AVX2	2.31

CPU: Intel Core i5-6500 (Skylake) with AVX2 256-bit SIMD

¹for long messages.

KANGAROOTWELVE's mode



Final node growing with kangaroo hopping and SAKURA coding

[ACNS 2014]