

Auxiliary Exchange in IKEv2 Protocol

`draft-smyslov-ipsecme-ikev2-aux`

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IETF 101

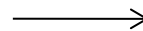
Initial IKEv2 Exchanges

Initiator

Responder

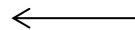
IKE_SA_INIT

HDR (MID=0) , SAi1, KEi, Ni



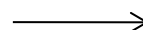
IKE_SA_INIT

HDR (MID=0) , SAr1, KEr, Nr



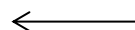
IKE_AUTH

HDR (MID=1) , SK{IDi, AUTH, SAi2, TSi, TSr}



IKE_AUTH

HDR (MID=1) , SK{AUTH, SAR2, TSi, TSr}



- IKE_SA_INIT messages are usually less than MTU – no IP fragmentation
- IKE_AUTH messages can be large, so IP fragmentation is possible
 - IP fragmentation interacts badly with some middleboxes like NAT and firewalls
- RFC7383 defines a way to avoid IP fragmentation by fragmenting messages in IKE
 - can only be used on encrypted messages, so IKE_SA_INIT is out of scope

The Problem

- Some recent proposals for IKEv2 protocol may lead to the situation when IKE_SA_INIT messages grow above MTU
 - Quantum Safe Key Exchange (QSKE) proposal defines additional Key Exchange payloads to be included into IKE_SA_INIT
 - something else?
- As result IKE_SA_INIT messages become subject for IP fragmentation with all aftermath
- Adding IKE fragmentation to IKE_SA_INIT is cumbersome and may lead to vulnerability to DoS attacks
 - IKE_SA_INIT messages have no protection, so an attacker who is able to see them and to inject bogus fragments can easily mount a reassembly queue poisoning attack

Proposed Solution

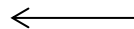
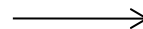
New auxiliary (**IKE_AUX**) exchange is added between IKE_SA_INIT and IKE_AUTH:

Initiator

Responder

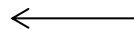
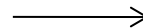
IKE_SA_INIT

HDR(MID=0), SAi1, KEi, Ni,
N(AUX_EXCHANGE_SUPPORTED)



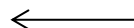
IKE_AUX

HDR(MID=1), SK{...}



IKE_AUTH

HDR(MID=2), SK{IDi, AUTH, SAi2, TSi, TSr}



IKE_SA_INIT

HDR(MID=0), SAR1, KEr, Nr
N(AUX_EXCHANGE_SUPPORTED)

IKE_AUX

HDR(MID=1), SK{...}

IKE_AUTH

HDR(MID=2), SK{AUTH, SAR2, TSi, TSr}

IKE_AUX Exchange

- New large payloads are placed in IKE_AUX, keeping IKE_SA_INIT messages small
- IKE_AUX messages are encrypted and MACed, so standard IKE fragmentation can be used
- IKE_AUX messages are authenticated by including their ICVs in signature calculation in IKE_AUTH:

```
InitiatorSignedOctets = RealMessage1 | AUX_I | NonceRData | MACedIDForI  
AUX_I = ICV_INIT_1 [ | ICV_INIT_2 [ | ICV_INIT_3 ... ] ]  
ResponderSignedOctets = RealMessage2 | AUX_R | NonceIData | MACedIDForR  
AUX_R = ICV_RESP_1 [ | ICV_RESP_2 [ | ICV_RESP_3 ... ] ]
```

Using IKE_AUX with QSKE

- Additional QSKE payload(s) are transferred using IKE_AUX
- IKE_AUX messages are protected using keys derived from key exchange performed in IKE_SA_INIT
 - IKE_SA_INIT messages must always contain KE payload
 - this KE payload may either contain classic (EC)DH public key or public key for some QSKE method, but it must be small enough not to cause IP fragmentation
- Keys for IKE_AUTH and for subsequent exchanges can be calculated as modification of standard IKE SA re-keying:

```
SKEYSEED(final) = prf(SK_d(initial), QSKE1 [| QSKE2 [| QSKE3 ...]] | Ni | Nr)
```

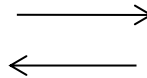
Keys in case of QSKE (example)

Initiator

Responder

IKE_SA_INIT

HDR(MID=0), SA_{i1}, KE_i, Ni,
N(AUX_EXCHANGE_SUPPORTED)



IKE_SA_INIT

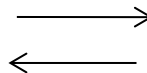
HDR(MID=0), SA_{r1}, KE_r, Nr
N(AUX_EXCHANGE_SUPPORTED)

$SKEYSEED(initial) = \text{prf}(Ni \mid Nr, g^{ir})$

IKE_AUX is protected using SK_e/SK_a keys derived from SKEYSEED(initial)

IKE_AUX

HDR(MID=1), SK{QSKE_{1i}, QSKE_{2i}}



IKE_AUX

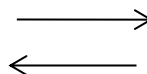
HDR(MID=1), SK{QSKE_{1r}, QSKE_{2r}}

$SKEYSEED(final) = \text{prf}(SK_d(initial), QSKE1 \mid QSKE2 \mid Ni \mid Nr)$

IKE_AUTH (and subsequent exchanges) is protected using SK_e/SK_a keys derived from SKEYSEED(final)

IKE_AUTH

HDR(MID=2), SK{ID_i, AUTH, SA_{i2}, TS_i, TS_r}



IKE_AUTH

HDR(MID=2), SK{AUTH, SA_{r2}, TS_i, TS_r}

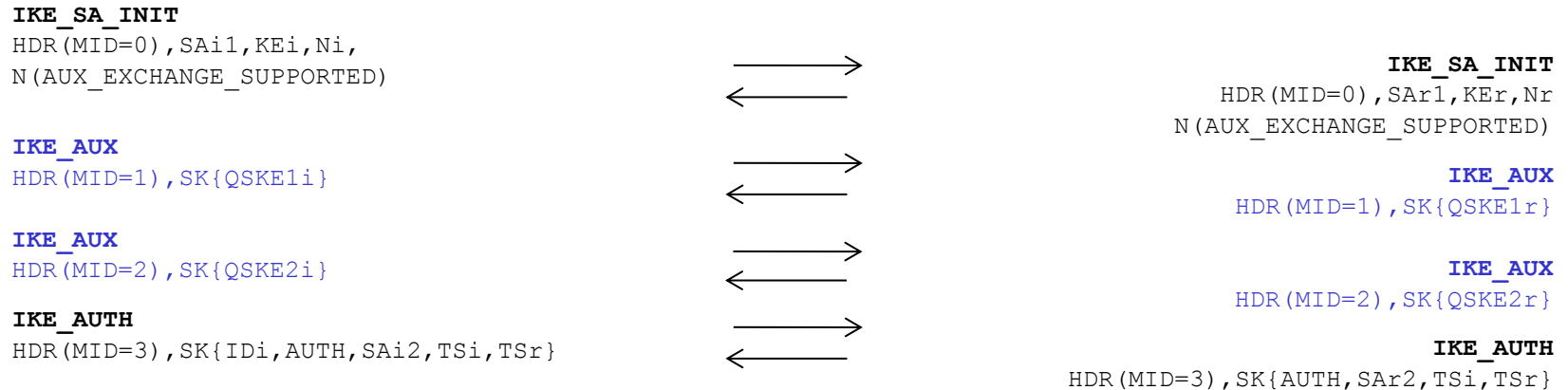
IKE_AUX Properties

- Complexity
 - a simple standard IKEv2 exchange
 - minimal influence on IKE_SA_INIT and IKE_AUTH
 - IKE_AUTH would start with Message ID > 1
 - uses standard IKEv2 fragmentation
 - some (small) impact on IKE state machine
 - modification of AUTH payload calculation
- Modularity
 - IKE_AUX is not tied to QSKE and can be used in other situations when large amount of data needs to be transferred prior to IKE_AUTH
- Security
 - DoS attacks surface in case of fragmentation is smaller than it would be if fragmentation were done in unprotected IKE_SA_INIT

IKE_AUX Properties (continued)

- Reliability

- if IKE_AUX is used with QSKE and several QSKE methods are employed, then each QSKE method can optionally be done in a separate IKE_AUX exchange:



- This would increase probability of IKE SA successful setup on congested or lossy networks in case IKE_AUX messages got fragmented using IKE fragmentation.

- Performance

- adds extra round trip(s)
- with QSKE re-calculation of SKEYSEED and derived keys is required

Thanks

- Comments? Questions?
- More details in the draft
- Please review and send feedback to author
- WG adoption?