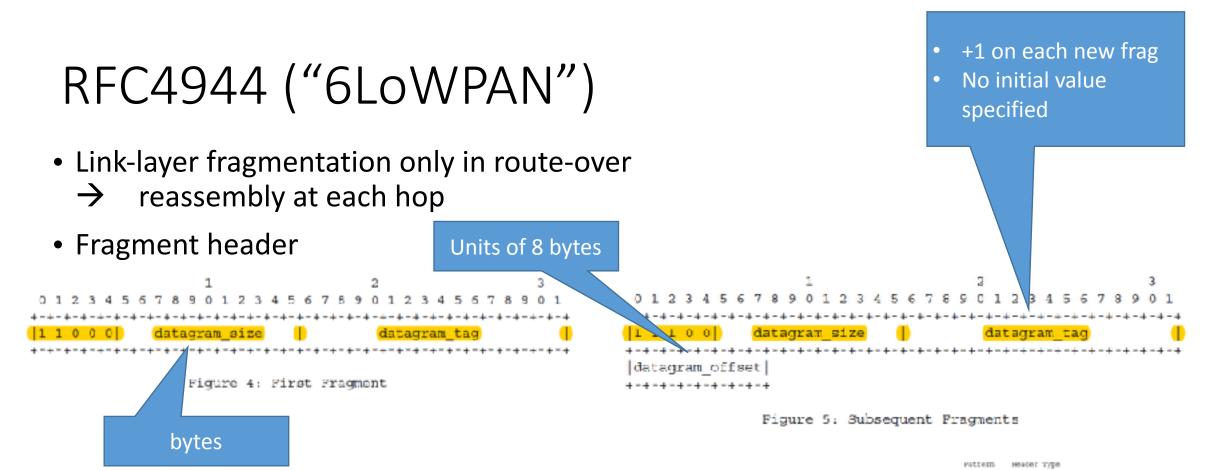
Memory-efficient fragment forwarding with Virtual Reassembly Buffers

draft-bormann-lwig-6lowpan-virtual-reassembly Carsten Bormann <<u>cabo@tzi.org</u>> with help from the 6lo Fragmentation design team Slides mostly provided by Thomas Watteyne — thanks!

Outline

- Review of standards
 - RFC4944
 - RFC6282
- Problem statement
- Reviewing Virtual Reassembly Buffers (VRB)
- Documents and Working Groups



- Reassembly timer:
 - Starts when node receives first fragment
 - Timeout value MUST be <60s
 - When times out, buffer cleared, packet dropped

0.0	NORTH AND	871.7		Bot a LoNTAN frame
D1	100001	-8V9	-	uncompressed this vograses
01	600610	LONDEN BOT	-	LONDAN HC: componented 1976
01	000011	reserved	-	Reserved for future use
		reserved	-	Reserved for future use
Drt.	001111	TERATION	-	beserved for future use
P1	030400	LOWING BOD	-	LOVENN_BC# browleast
D1	010001	reserved	-	Reserved for future use
		reserved	-	Reserved for future use
D1	111110	reseaved	-	Reseaved for future use
D1	1111111	ISC	-	Additional Dispatch byte follows
2.0	XXXXXXXXXX	#842.4		Nosh Header
	0001388	P RANGE I		fragmentation Header (first)
11	001000	reserved	-	Reserved for future use
		reserved		Reserved for future use
2.1	911111	Lensined	-	Reserved for future use
81	1003306	PEA-IN	-	Tragmentation Header (subsequent)
1.1	1014-00	recorved		Desceved for future use
	101+00	reserved		Reserved for future use
-		reserved		

RFC6282 ("6LoWPAN-HC")

Section 5.3 of [RFC4944] also defines how to fragment compressed IPV6 datagrams that do not fit within a single link frame. Section 5.3 of [RFC4944] defines the fragment header's datagram_size and datagram_offset values as the size and offset of the IPv6 datagram before compression. As a result, all fragment payload outside the first fragment must carry their respective portions of the IPv6 datagram before compression. This document does not change that requirement. When using the fragmentation mechanism described in Section 5.3 of [RFC4944], any header that cannot fit within the first fragment MUST NOT be compressed.

Problem statement

- Per-hop fragmentation and reassembly has 2 issues:
 - Latency:
 - Increases end-to-end latency as you need to wait for each fragment at each hop
 - Reliability:
 - Limited memory → limited number of buffers (1-2?) → packet dropped when new frag received and old not fully reassembled yet
 - No frag recovery: 1 frag loss == packet dropped
- Proposed solution:
 - Fragment forwarding:
 - Source fragments
 - Intermediate nodes relays
 - LBR reassembles

6LoWPAN: The Wireless Embedded Internet a.k.a. "Zach and Carsten's book", Wiley 2009

2.5.2 L3 routing ("Route-Over")

Layer-3 Route-Over forwarding is illustrated in Figure 2.6. In contrast to layer-2 mesh forwarding, layer-3 Route-Over forwarding does not require any special support from the adaptation layer format. Before the layer-3 forwarding engine sees the packet, the adaptation layer has done its work and decapsulated the packet – at least conceptually (implementations may be able to perform some optimizations by keeping the encapsulated form if they know how to rewrite it into the proper encapsulated form for the next layer-3 hop).

Note that this in particular means that fragmentation and reassembly are performed at each hop in Route-Over forwarding – it is hard to imagine otherwise, as the layer-3 addresses are part of the initial bytes of the IPv6 header, which is present only in the first fragment of a larger packet. Again, implementations may be able to optimize this process by keeping virtual reassembly buffers that remember just the IPv6 header including the relevant addresses (and the contents of any fragments that arrived out of order before the addresses).

Context: typical fragmentation implementation

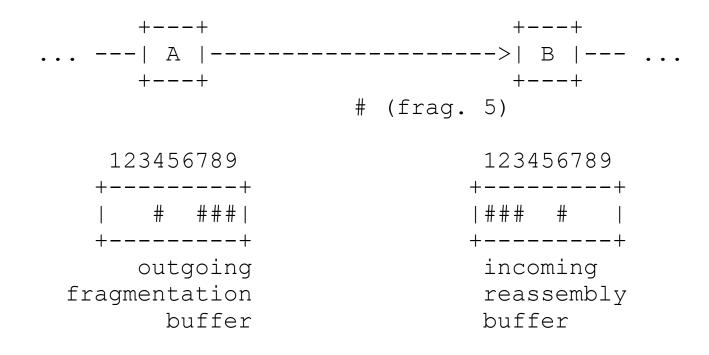
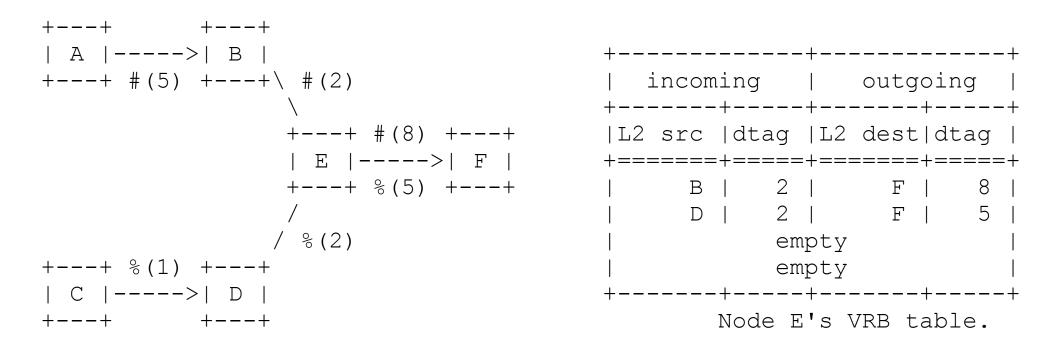


Figure 1: Fragmentation at node A, reassembly at node B.

→ Limits: latency, end-to-end reliability (see preliminary simulation results)

Virtual Reassembly Buffer (VRB) Implementation



VRB: Gotchas and Limits

- Gotchas
 - Long headers (e.g. source routing) \rightarrow receive multiple fragments before forwarding
 - Out-of-order fragments \rightarrow receive multiple fragments before forwarding
 - Changing header length \rightarrow put all slack in the frame sizes into the **first** fragment
- Security
 - DoS attack still possible by "fragment 1" flood to overflow VRB
- This is an implementation technique, not a protocol. Limits:
 - Non-zero Packet Drop Probability
 - No Fragment Recovery
 - No Per-Fragment Routing

→ If limits are not acceptable you need an actual protocol (such as draft-thubert-6lo-fragment-recovery) — out of scope for LWIG

6lo fragmentation design team

- Produce 2 documents (to be submitted to 6lo WG):
 - informational document
 - summarize fragmentation as standardized now
 - describes Carsten's virtual reassembly buffer implementation
 - discusses its limits
 - standards-track document
 - builds upon the first one
 - adds fragment recovery
- Philosophy
 - keep activity as swift as possible
 - small DT, but regular information to WGs
 - ideally close the DT after London (to be discussed)

draft-watteyne-6lominimal-fragment-00

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6lo vs. lwig

- VRB discussion was first published as draft-bormann-lwig-6lowpan-virtual-reassembly
- Original intention was to have everything in draft-watteyne-6lominimal-fragment
- Proposal now is to:
 - Keep a focused document just on the VRB implementation technique in LWIG (draft-...-lwig-6lowpan-virtual-reassembly)
 - Put the analysis and overview over new protocol proposals in 6Lo (draft-...-6lo-minimal-fragment)