

dmm
Internet-Draft
Intended status: Standards Track
Expires: October 25, 2019

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April 23, 2019

Co-existence of 3GPP 5GS and Identifier-Locator Separation Architecture
draft-homma-dmm-5gs-id-loc-coexistence-03

Abstract

This document describes an approach to introduce Identifier Locator Separation architecture into 3GPP 5GS with low-impact on its specifications, and shows the features and considerations of this approach.

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1. Introduction

Identifier-Locator Separation (ID-LOC) architectures aim to simplify management of network, devices, and sessions by employing two namespaces: Identifier for device's identity, and Locator for its location in the network.

An ID-LOC architecture can be implemented by a dedicated protocol such as LISP [I-D.ietf-lisp-rfc6833bis], ILA [I-D.herbert-intarea-ila], ILNP [RFC6740], etc. The control plane of such ID-LOC protocols can be combined with one of different encapsulation techniques such as GTP-U [TS.29281], SRv6 [I-D.filsfils-spring-srv6-network-programming], MPLS [RFC3031], VXLAN [RFC7348], etc. at data plane to provide a customized solution. Furthermore, regarding control plane of ID-LOC, it can optionally even take advantage of enhanced PUB/SUB capable distributed databases to store ID-LOC mappings.

ID-LOC protocols are also expected to be used for optimizing user-plane of mobile network [I-D.bogineni-dmm-optimized-mobile-user-plane]. Different alternatives to introduce ID-LOC architecture into 3GPP 5GS (5th Generation System), are under consideration in related IETF WG such as DMM WG.

Introducing ID-LOC architecture into mobile network can involve modifications to 5GS architecture and specifications that might span over several 5GS releases.

Therefore, an approach that enables the introduction of ID-LOC architecture into 5GS without change of its specifications and supports migration path toward a native ID-LOC network can be useful to operators. Here, ID-LOC native network refers to a network that employs the ID-LOC architecture as only mechanism for packet forwarding.

The document aims to describe one such approach and clarify different features, and benefits.

2. Definition of Terms

This section describes general terms of ID-LOC architecture. This document also refers definitions of 3GPP 5GS [TS.23.501-3GPP], and some of such terms which are used in this document are listed in this section.

The LISP terms are described in [I-D.ietf-lisp-rfc6833bis].

The ILA terms are described in [I-D.herbert-intarea-ila].

The SRv6 terms are described in [I-D.ietf-6man-segment-routing-header].

2.1. Terms of ID-LOC Protocols

Device Identifier (ID): An ID is an identifier of host or end point such as UE or network function including VM instance, container, etc. In ID-LOC architectures, an IP or MAC address is generally assigned to an end device as identifier. In this case, IDs are used as values for the source and destination IP/MAC address fields of packets sent from end points. Alternatively, other attributes of the end point, such as its Fully Qualified Domain Name (FQDN), can also be used as IDs.

Locator (LOC): A LOC is generally an address (e.g. IPv4, IPv6, MAC, etc) of the ID-LOC node. In the case of SRv6 it can be the ID-LOC node's local SID representing the segment for which the ID-LOC node is the segment termination node.

ID-LOC node; An ID-LOC node is a node that has at least one unique locator within a network domain, and has functionalities to obtain destination locator and to forward packets to the ID-LOC node which has the destination locator. This node has an ID-LOC mapping cache and obtains destination locator by looking up destination ID (destination address of a data packet) from the mapping cache. If ID of the received packet is not present in its own mapping cache, an ID-LOC node requests mapping information of the ID and the assigned locator to ID-to-LOC mapping system. Also an ID-LOC node forwards packet to a peered LOC node by encapsulation or conversion of the IP header field such as IP address field, and decapsulates or reconverts packets received from another ID-LOC node. Different implementations of ID-LOC architecture use different forwarding mechanisms. LISP data-plane, for example, uses IPv4/v6 header and LISP header for encapsulation, whereas ILA tightly couples itself with IPv6, and SRv6 uses IPv6 and SIDs (Segment Identifiers).

ID-to-LOC Mapping System: An ID-to-LOC mapping system is a database which contains all known ID-to-LOC mappings within an ID-LOC domain. The mapping information is updated when an end point moves to under another ID-LOC node. This database can be logically centralized, distributed across the ID-LOC nodes, or a combination of both. If the database is logically centralized, each ID-LOC node has an interface to the system to send a request and receive mapping information.

ID-to-LOC Mapping Cache: An ID-LOC mapping cache is a table in an ID-LOC node that stores ID-to-LOC mapping information and it is used for obtaining destination LOC from ID of received packet. ID-to-LOC mapping cache typically contains a small piece of database. The cache is updated when the ID-LOC node receives a new ID-to-LOC mapping information from ID-to-LOC mapping system.

2.2. Terms of 5GS

User Plane Function (UPF): An UPF handles the user plane paths. An UPF is connected to SMF with N4 interface. More detailed information is described in [TS.23.501-3GPP]. This document defines two types of UPF, Central UPF (cUPF) and Distributed UPF (dUPF). Their features are described in Section 3

Uplink Classifier (ULCL): An ULCL is an UPF functionality that aims at diverting Uplink traffic, based on filter rules provided by SMF, towards Data Network (DN).

Data Network (DN): A DN is a network where network functions and entities, including operator or 3rd party services, are deployed. This document defines two types of DN, Central DN (cDN) and Distributed DN (dDN). Their features are described in Section 3.

Radio Access Network (RAN): A RAN is an access network where radio bearer sent by UEs traverse. A RAN encapsulate users' packets with GTP-U.

Session Management Function (SMF): An SMF is a function which provides control plane functionalities for handling user traffic.

Application Function (AF): An AF is a control plane functionality and connected to SMF with Naf interfaces.

3. Mechanism on Data Plane

This approach achieves traffic forwarding with optimized path and session continuity by using ID-LOC and ULCL for particular communication including UE-to-UE or MEC (Mobile Edge Computing) communication. ULCL is one of fundamental functions of 5GC Rel.15 and it provides functionalities of packet filtering and divert for uplink packets sent by UEs.

The overview of the assumed 5GC architecture of data plane where the proposal approach works is shown in Figure 1. The details of numbered interfaces in the figure are described in [TS.23.501-3GPP].

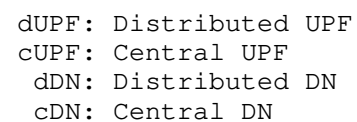


Figure 1: Assumed 5GC Network Architecture

This network has following features;

- o A Central UPF (cUPF) is deployed at a connecting point to Central DN (cDN). A cUPF becomes anchor point for UEs and it assigns IP

addresses (IDs) for each UE. The traffic transmitted from UEs are basically sent to the cUPF.

- o Distributed UPFs (dUPFs) and Distributed DN (dDN) are deployed and geographically distributed at user edge side. A unique address space (it's not necessarily globally unique) is assigned to dDN. When a dUPF forwards a UE's uplink packet, and if the subnet of the destination address is the same as the one assigned to dDN at proximity, then dUPF, with the help of ULCL, may divert the packet to that dDN. Here, the ULCL identifies each encapsulated uplink packet to be diverted, by checking if the destination of the inner packet is one of IP addresses assigned to the dDN. A dUPF removes GTP-U header from the packets, and sends them to dDN via N6. When dUPF receives packets from dDN, dUPF encapsulates them with GTP-U header, and merges them into downlink packets from cUPF. An overview of behaviors of dUPF and ULCL is shown in Figure 2.
- o Network topology between RAN and dUPF/cUPF adopts tree structure and the section between RAN and dUPF and the section between dUPF and cUPF are connected with GTP-U.

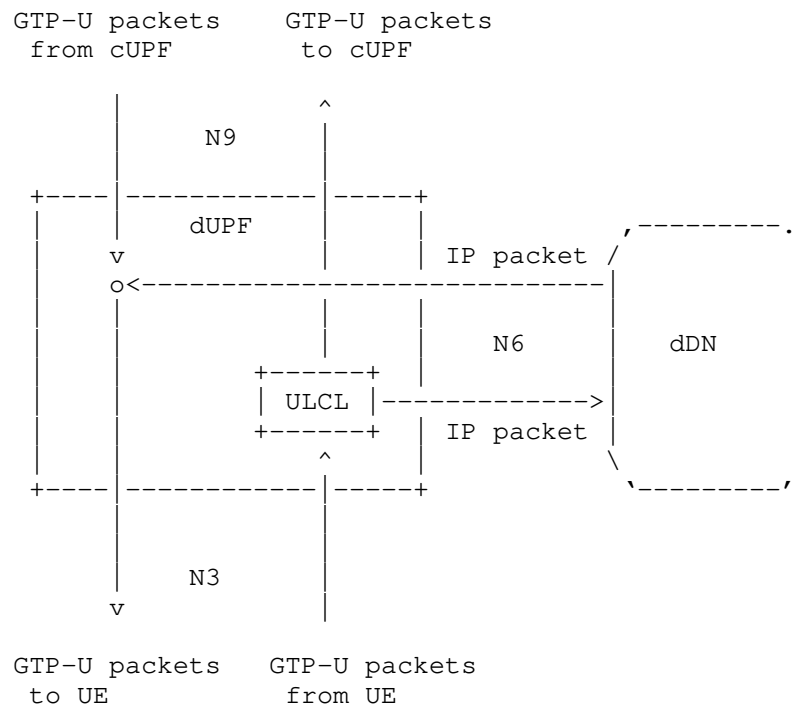


Figure 2: Behaviors of dUPF and ULCL

In the proposed approach, IDs are assumed to be IP addresses and an ID-LOC node is installed between dUPF and dDN. ID-LOC nodes are connected with a IP mechanism such as IP tunnels or translation of destination IP field. As examples of such data plane protocols for providing connectivity between ID-LOC nodes, IPv4/v6 header with LISP header or SRv6 ([I-D.ietf-6man-segment-routing-header]) can be used. In addition, each ID-LOC node has connectivity with one or more Mapping Systems. The overview is shown in Figure 3.

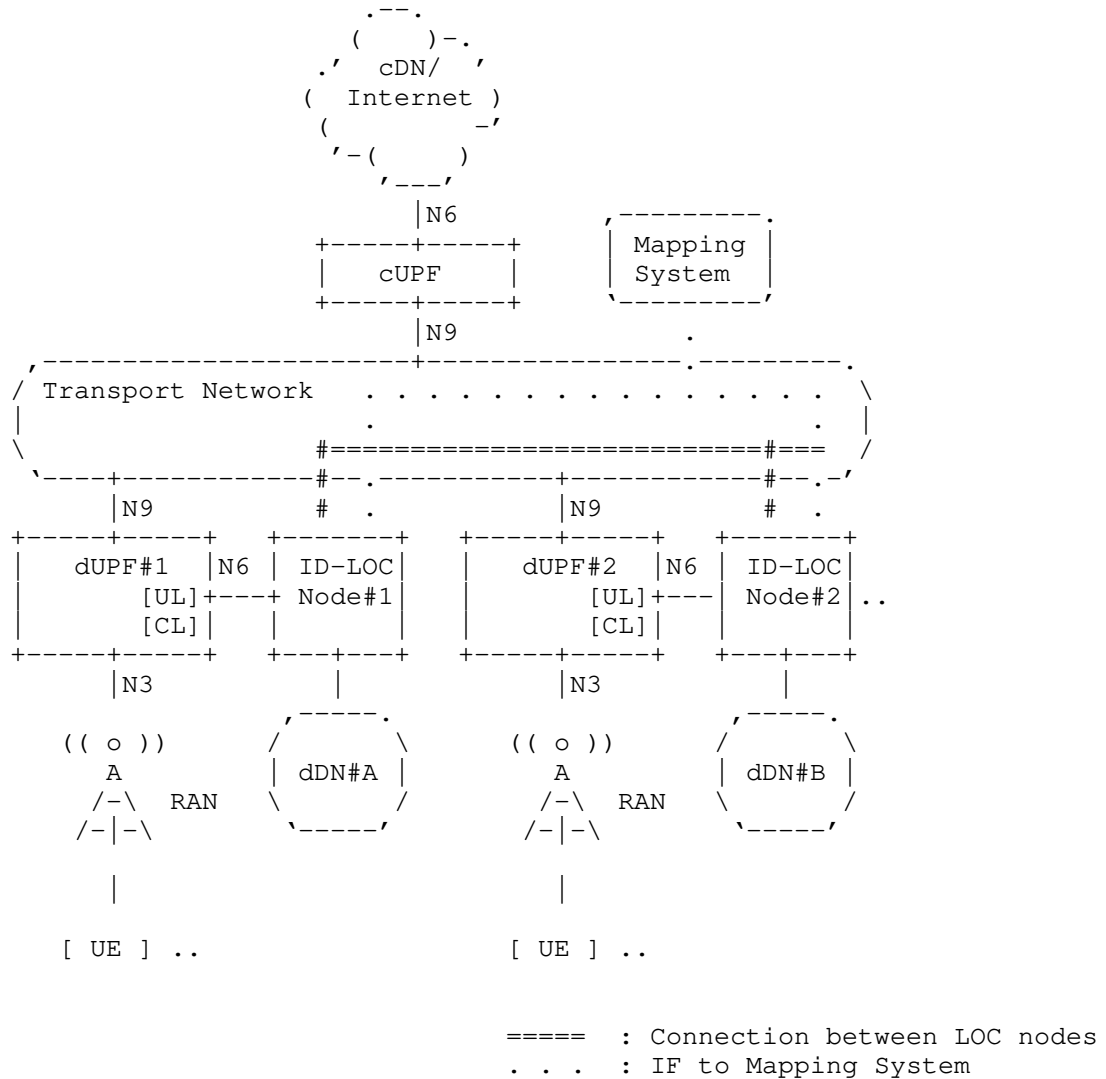


Figure 3: Proposal Network Architecture

Each dUPF has a filter table of ULCL. Each filter table is configured to match the addresses of UEs within the network domain (i.e., addresses for UEs assigned by the cUPF). Filter tables can also be configured to match the address corresponding to the address space (or part of) corresponding with the dDNs in the network domain. UPFs monitor each uplink GTP-U packet with its ULCL and divert it to the connected ID-LOC node with decapsulation of GTP-U if the

destination address of the inner packet (payload) matches the filtering table. When ID-LOC node receives a packet from the dUPF, it obtains LOC which the destination of the packet (ID) belongs to by looking up its own ID-to-LOC mapping table or querying it from the Mapping System according ID-LOC mechanism. Then it sends the packet to peered ID-LOC node indicated by the LOC. The peered ID-LOC node converts the received packet to appropriate form and forwards them the destination by following its own forwarding table.

From such processes, forwarding paths of user traffic diverted by ULCL from 5GC to ID-LOC node are optimized.

A cUPF is connected with dUPFs via N9 interface and packets are forwarded with GTP-U encapsulation between cUPF and dUPF.

Some case studies of ID-LOC protocols are described in Appendix A and Appendix B.

4. Mechanisms on Control Plane

For ID-LOC mechanism in mobile networks, a control plane mechanism is required to manage location information of UEs and NFs in each dDN. There are mainly three models to realize control plane mechanism for ID-LOC as follows:

Model 1: Independent Control Planes

Model 2: Interworking Control Planes

Model 3: Integrated Control Planes

Some of models may require to use 5GS interfaces or add some functionalities to functions of 5GC. 5GS architecture and the service-based interfaces are shown in Figure 4. The details of functions and interfaces are described in [TS.23.501-3GPP].

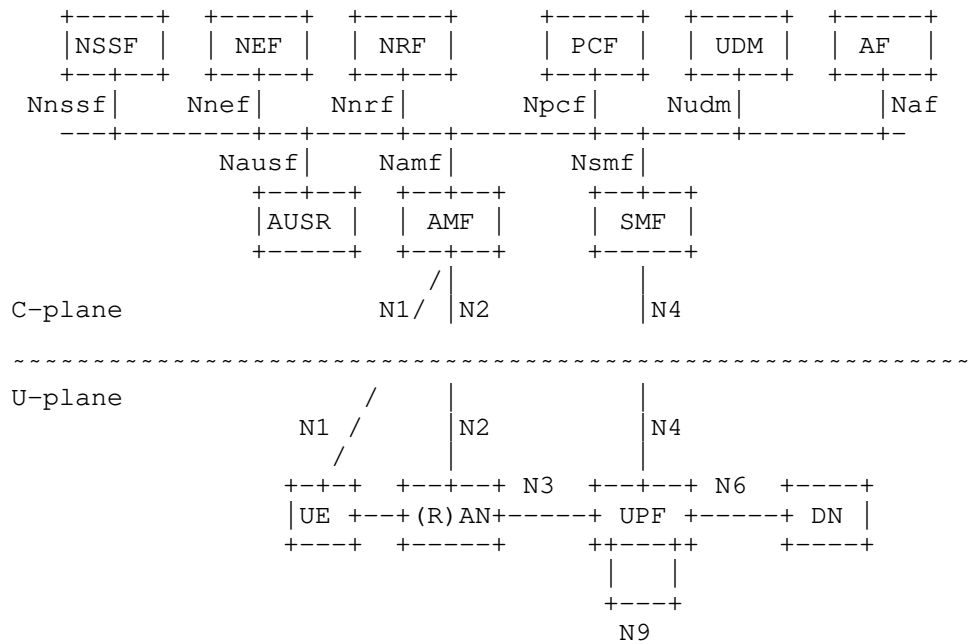


Figure 4: 5GS Architecture and Service-based Interfaces

4.1. Model 1: Independent Control Planes

In this model, control plane of 5GC and ID-to-LOC mapping mechanism are completely separated. Information of a UE and an ID-LOC node which the UE is attached is sent to a mapping system and registered in the mapping database only when the ID-LOC node receives a packet from the UE and the UE is not registered yet.

This model does not cause any impacts on 5GC architecture. However, in this model, a UE cannot be accessed from other UEs within the same network domain until a packet from the UE is diverted to the ID-LOC node by the UPF which the UE is located and the ID and LOC are registered to the Mapping System.

4.2. Model 2: Interworking Control Planes

In this model, a mapping system interworks with an SMF which manages sessions of each UE. A scheme to inform, that a UE moves and is relocated to another UPF, from SMF to AF via Naf interface is defined in 5GS ([TS.23.502-3GPP])* . A Mapping System is installed as an AF and obtains mobility information of UEs with the above scheme.

* The stage 3 of discussion of 5GS has not been fixed yet and the specification may be changed.

This model would not cause any impacts on 5GS architecture, and a mapping system can always keep the current mobility information of each UE.

4.3. Model 3: Integrated Control Planes

In this model, SMF functionalities are integrated into a mapping system. In other words, the mapping system becomes a part of 5GS. In 5GS architecture, an SMF has a role of session management of UEs, and it updates its own mapping database depending on movement of a UE.

This approach enables to always keep mapping databases the latest status, however, it obviously requires extension or replacement of SMF actually deployed in 5GS network.

5. Features Analysis

5.1. Benefits

- o This approach provides a mechanism for introducing ID-LOC architecture into 5GS with no or nominal impact, and achieves optimized forwarding with session continuity in the assumed use cases such as UE-to-UE or UE-to-dDN communications.
- o Regarding communication to the cDN, this approach can keep scalability because it does not change the current mechanism of 5GS.

5.2. Issues

- o dUPF and ID-LOC node are separated, and thus an extra hop may occur against the optimized forwarding. However, it can be resolved by implementing dUPF and ID-LOC node within a same box or application.

6. Security Considerations

TBD

7. IANA Considerations

This memo includes no request to IANA.

8. Acknowledgement

The authors would like to thank Ryosuke Kurebayashi, Koji Tsubouchi, Toru Okugawa, and Dino Farinacci for their kind reviews and technical feedback.

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Appendix A. Case Studies on Use of LISP

This Appendix describes detailed processes of the proposal approach with LISP mechanism in the following types of communications.

1. UE-to-UE Communication
2. UE-to-dDN Communication
3. UE-to-cDN/Internet Communication

In the following description of case studies, ID and Locator are called EID (End-point Identifier) and RLOC (Routing Locator) in LISP terms. Mapping Server has the master of EID-to-RLOC mapping database, and each xTR (Ingress/Egress Tunnel Router) has EID-to-RLOC mapping cache. An xTR obtains the destination RLOC from its own

cache by looking up the destination EID of received packet. They obtain mappings from the mapping system if an EID looked up is not registered in the cache. Packets are passed between xTRs with some tunnel protocols.

A.1. UE-to-UE Communication

In the current architecture, a cUPF becomes an anchor point for UEs, and all packets between UEs even those which are located to the same dUPF are transferred through the anchor point. This may cause communication delay and inefficient resource usage. In the proposed procedure, packets can be transferred without going through an anchor point, and low latency and efficient resource usage can be achieved.

The UE-to-UE communications include communications between UEs located to different dUPFs (Case 1), and communication between UEs located to the same dUPF (Case 2). In this section, the detailed procedures of the cases are described.

Moreover, in a mobile network, a UE may move during communications. This section describes considerations about UE's handover in such case.

A.1.1. Case A-1: UEs allocated different dUPF

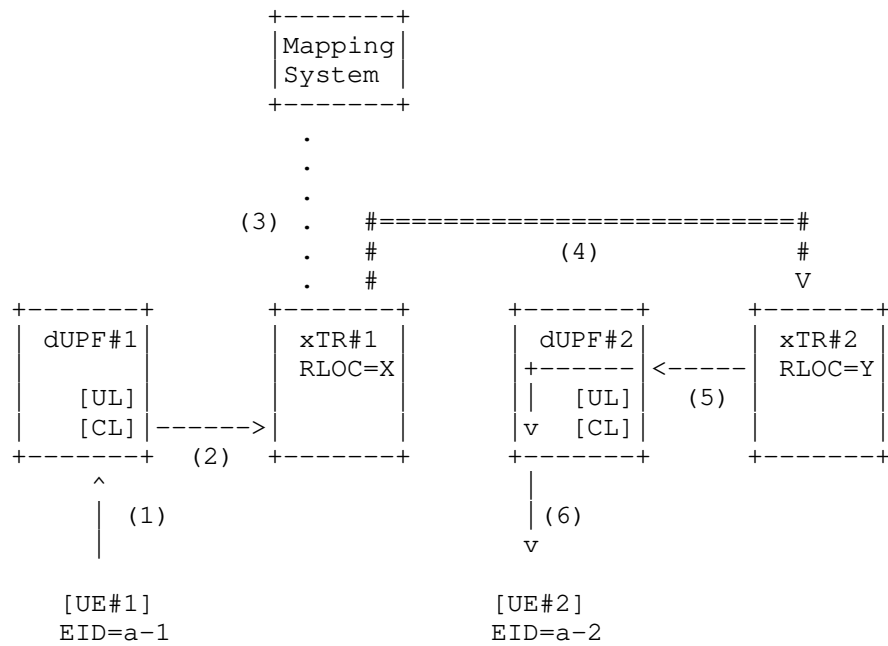


Figure 5: Procedure in Case A-1

- (0) Within this network, addresses are assigned to UEs from a address space [A]. These addresses are described as a-n (n=1,2,..). EID=a-1 and a-2 are assigned to UE#1 and UE#2.
- (1) UE#1 sends packets to UE#2 with setting EID=a-2 as the destination IP address.
- (2) dUPF#1 monitors inner packet of received GTP-U packet and divert it to xTR#1 with decapsulation if the destination address is one of address space [A].
- (3) xTR#1 updates own EID-to-RLOC mapping cache by interaction with Mapping System (if needed).
- (4) xTR#1 obtains the RLOC(=Y) of EID=a-2 from the EID-to-RLOC mapping cache, and sends the packets to the xTR#2 with a tunnel with RLOC=Y as the destination address.
- (5) xTR#2 decapsulate the packets, and sends them to dUPF#2.
- (6) dUPF#2 encapsulate packets with GTP-U header, and sends them to UE#2.

A.1.2. Case A-2: UEs allocated the same xTR

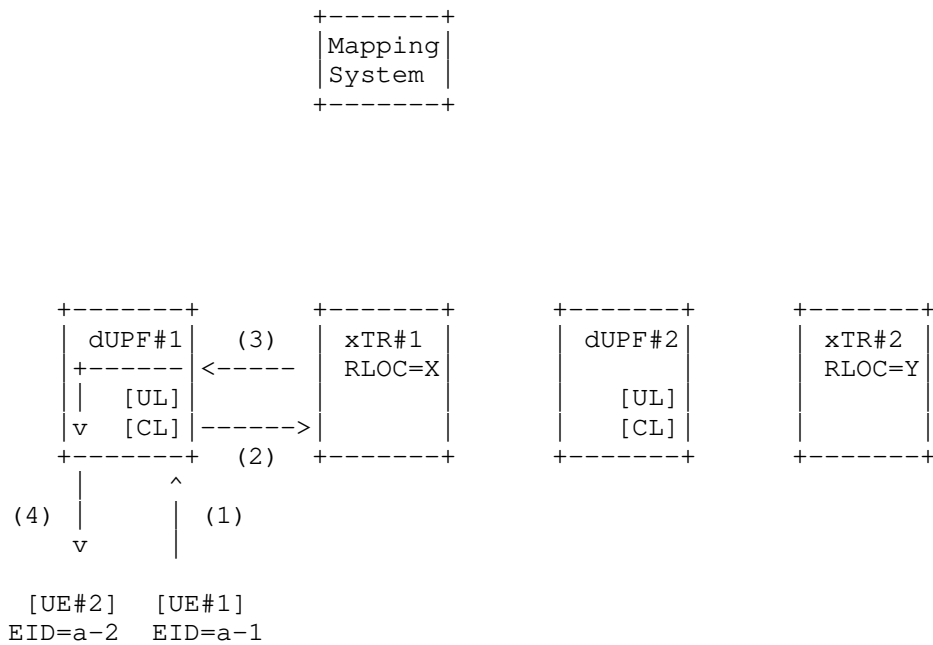


Figure 6: Procedure in Case A-2

- (0) Within this network, addresses are assigned to UEs from a address space [A]. These addresses are described as a-n (n=1,2,..). EID=a-1 and a-2 are assigned to UE#1 and UE#2.
- (1) UE#1 sends packets to UE#2 with setting EID=a-2 as the destination IP address.
- (2) dUPF#1 monitors inner packets of received GTP-U traffic and divert it to xTR#1 with decapsulation if the destination address is one of address space [A].
- (3) Since xTR#1 serves UE#2, it locally routes the traffic for EID=a-2. xTR#1 sends the received packets back to dUPF#1.
- (4) dUPF#1 encapsulate packets with GTP-U, and sends them to UE#2.

A.1.3. Consideration of Case that UE Moves to under Another xTR

When a UE moves to a serving area of another dUPF during communication with another UE, EID-to-RLOC mapping database of a Mapping System and the tables of the xTR and the peered xTR must be updated. Unless some of the mechanism described below are in place, the xTRs can't send packets to the appropriate xTR during the updating, and thus packet drop or stalling may occur.

For example, a mechanism that immediately advertise the update of location of UEs to Mapping System and the appropriate xTRs depending on movement of each UE might be required. Some documents (e.g., [I-D.ietf-lisp-eid-mobility], [I-D.ietf-lisp-pubsub]) discuss such mechanisms. Alternatively, a mechanism that replicates packets to both the old and new location while the UE is in transit could also be used. This approach is discussed in detail in [I-D.ietf-lisp-predictive-rlocs].

A.2. UE-to-dDN Communication

The UE-to-dDN communications basically correspond the communication between a UE and neighbor dDN (Case3). On the other hand, if a UE moved under another dUPF during usage of a stateful application, or the application is not uniformly deployed in every dDN, the UE needs to continue to communicate with the previous dDN (Case4).

In such cases, in the current architecture, all packets are needed to go through the anchor point or dynamic GTP tunnel reconfiguration between dUPF is required. The former solution causes additional communication delay and inefficient resource usage. The latter solution increase the cost of 5GS control plane to dynamically update the GTP tunnel with multiple UPFs and their ULCL filter tables along with the movement of the UE. The proposed approach achieves appropriate packet transfer in such cases.

In this section, the detailed procedures of communications between a UE and neighbor dDN and communications between a UE and non-neighbor dDN

A.2.1. Case A-3: UE communicates with neighbor dDN

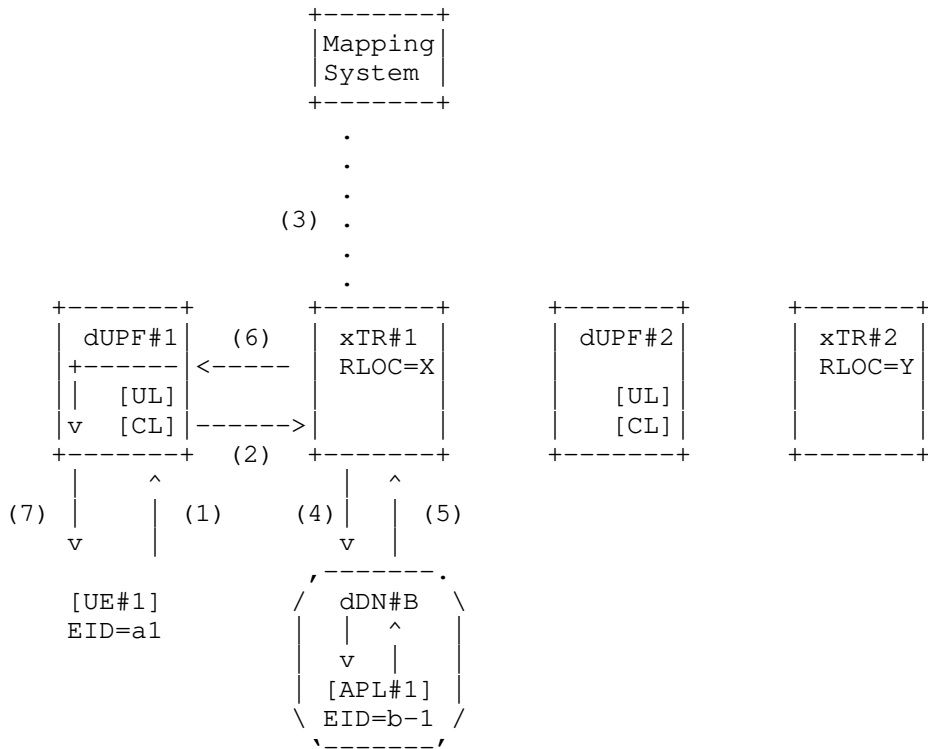


Figure 7: Procedure in Case A-3

- (0) Within this network, UEs are assigned their addresses from an address space [A]. These addresses are described as a-n (n=1,2,...). Also, applications in dDN#B are assigned their addresses from a address space [B]. These addresses are described as b-n (n=1,2,...). EID=a-1 and b-1 assigned to UE#1 and APL#1 which is located in dDN#B.

[Uplink Processes]

- (1) UE#1 sends packets to dDN#B with setting EID=b-1 as the destination IP address.
- (2) dUPF#1 monitors inner of received GTP-U packets and divert it to xTR#1 with decapsulation if the destination IP address is one of address space [B].

- (3) xTR#1 updates own EID-to-RLOC mapping cache by interaction with Mapping System (if needed). Or xTR#1 may update its own cache by a Map-Notify message when an APL is deployed or deleted in dDB#B.
- (4) Since xTR#1 serves dDN#B, it locally routes the traffic for EID=b-1. xTR#1 sends the packets to the dDN#B.

[Downlink Processes]

- (5) APL#1 in dDN#B sends packets to UE#1 with setting EID=a-1 as the destination IP address.
- (6) Since xTR#1 serves UE#1, it locally routes the traffic for EID=a-1. xTR#1 sends packets to dUPF#1.
- (7) dUPF#2 encapsulates packets with GTP-U, and sends them to UE#1.

A.2.2. Case A-4: UE communicates with non-neighbor dDN

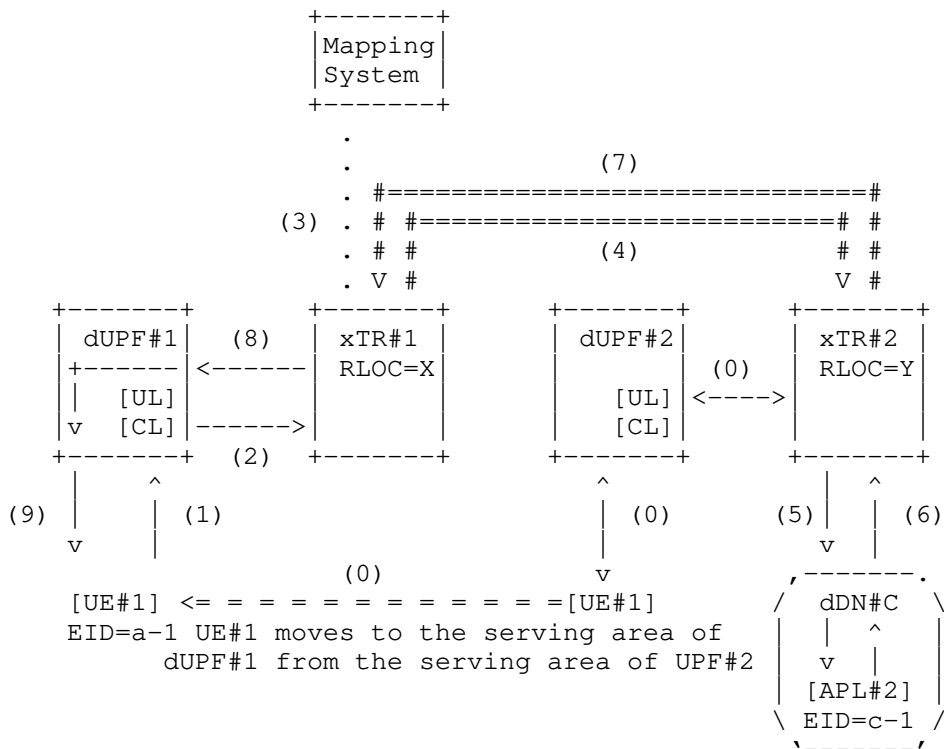


Figure 8: Procedure in Case A-4

- (0) Within this network, UEs are assigned their addresses from an address space [A]. These addresses are described as a-n (n=1,2,..). And applications in dDN#C are assigned their addresses from an address space [C]. These addresses are described as c-n (n=1,2,..). EID=a-1 and c-1 assigned to UE#1 and APL#2 which is located in dDN#C. UE#1 has moved to the serving area of dUPF#1 from the serving area of UPF#2 while communicating to APL#2.

[Uplink Processes]

- (1) UE#1 sends packets to APL#2 with setting EID=c-1 as the destination IP address.
- (2) dUPF#1 monitors each inner packet of received GTP-U traffic and divert it to xTR#1 with decapsulation if the destination address is one of address space [C].
- (3) xTR#1 updates own EID-to-RLOC mapping cache by interaction with Mapping System (if needed).
- (4) xTR#1 obtains RLOC(=Y) of EID=c-1 from the EID-to-RLOC mapping cache, and sends the packet to the xTR#2 with a tunnel with RLOC=Y as the destination address.
- (5) xTR#2 decapsulates the packets received from xTR#1, and sends them to dDN#C depending on its forwarding table.

[Downlink Processes]

- (6) APL#2 sends packets to UE#1 with setting EID=a-1 as the destination IP address.
- (7) xTR#2 obtains RLOC(=X) of EID=a-1 from the EID-to-RLOC mapping cache, and sends the packets to the xTR#1 with a tunnel with RLOC=X as the destination address.
- (8) xTR#1 decapsulates the packets received from xTR#2m and sends them to the dUPF#1 depending on its forwarding table.
- (9) dUPF#1 encapsulates the packets with GTP-U and sends packets to UE#1.

A.3. UE-to-cDN/Internet Communication

UE-to-cDN/Internet communication is achieved by GTP-U mechanism originally equipped in 3GPP 5GS architecture. In this section, we

describe processes of UE-to-cDN communication in the proposal architecture as an example.

A.3.1. Case A-5: UE communicates with cDN

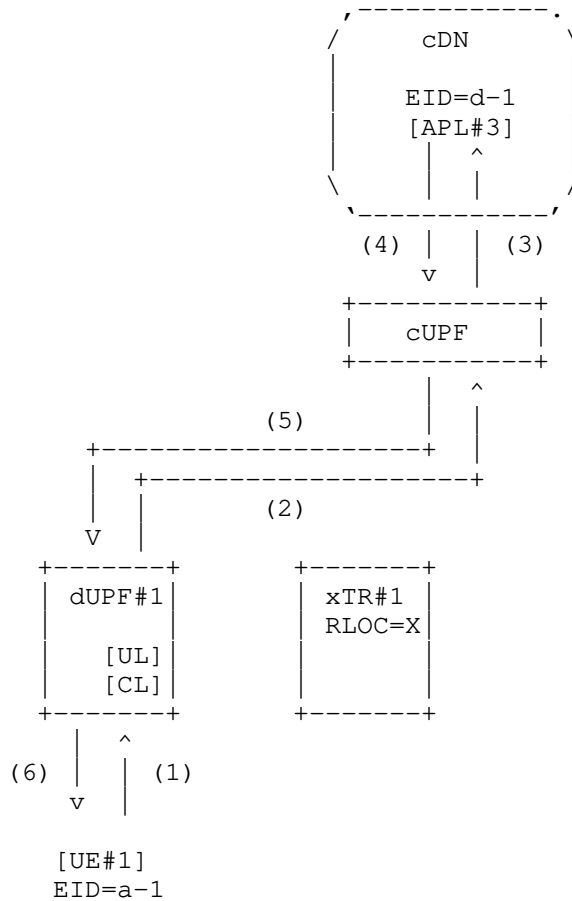


Figure 9: Procedure in Case A-5

- (0) Within this network, UEs are assigned their addresses from an address space [A]. These addresses are described as a-n (n=1,2,..). And applications in cDN are assigned their addresses from an address space [D]. These addresses are described as d-n (n=1,2,..). EID=a-1 and d-1 assigned to UE#1 and APL#3 which is located in cDN.

[Uplink Processes]

- (1) UE#1 sends packets to cDN with setting EID=d-1 as the destination IP address.
- (2) dUPF#1 monitors inner of received GTP-U packets. Since the destination IP address (EID=d-1) does not hit the filter of ULCL, dUPF#1 re-encapsulates the packet to another GTP-U connecting to cUPF and forwards to cUPF.
- (3) cUPF decapsulates GTP-U packets and forwards them to APL#3 in cDN depending on its own forwarding table.

[Downlink Processes]

- (4) APL#3 in cDN sends packets to UE#1 with setting EID=a-1 as the destination IP address.
- (5) cUPF encapsulates the packets received from APL#3 and forwards them to dUPF#1 depending on its own forwarding table.
- (6) dUPF re-encapsulates the packets to another GTP-U and forwards to UE#1.

Appendix B. Case Studies on Use of ILA

This Appendix describes detailed processes of the proposal approach with ILA mechanism in the following types of communications.

1. UE-to-UE Communication
2. UE-to-dDN Communication
3. UE-to-cDN/Internet Communication

Each ILA node has ID-to-LOC mapping table. Mappings are propagated amongst ILA routers or hosts in a network using mapping propagation protocols.

In the following description of case studies, a mapping system, called ILA resolver in ILA terms, has the master of ID-to-LOC mapping database, and each ILA node obtains mappings from the mapping system. In some cases, each ILA node has an ID-to-LOC mapping database.

In ILA, an SIR address expressed by composition of SIR prefix and identifier is assigned to each UE or VM instance. An SIR prefix and an identifier are described SIR_prefix_n and id_m (n=1,2,..., m=1,2,...), and an SIR address is expressed as SIR_addr_x =[n,m]

(x=1,2,...) in the following description. Also, each ILA-Nodes are assigned unique Locators, which is a network prefix that routes to a host. Locators are described as loc_n (n=1,2,...).

B.1. UE-to-UE Communications

The overview of this communication type is described in A.1.

B.1.1. Case B-1: UEs allocated different dUPF

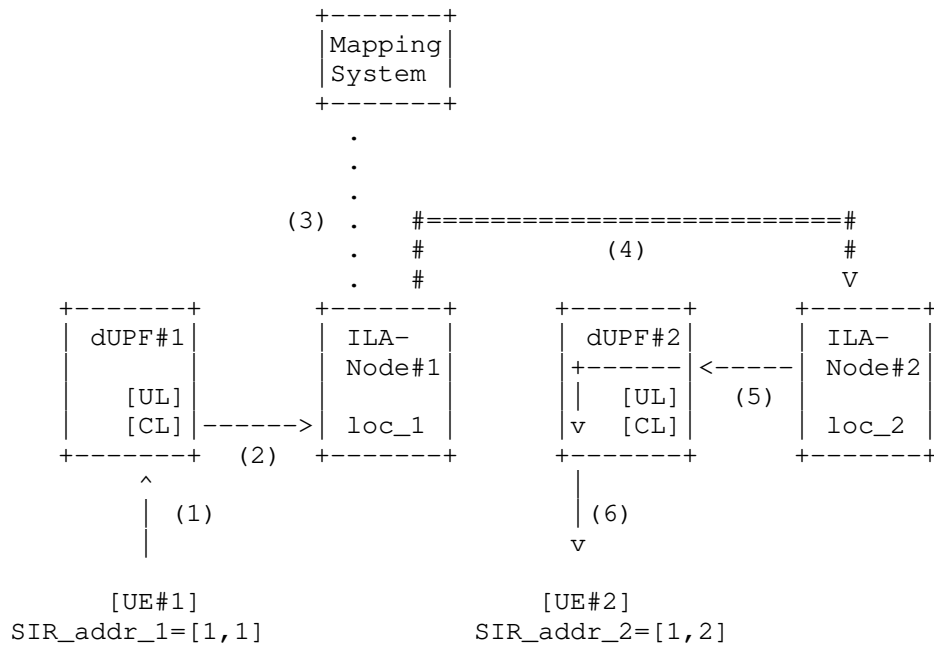
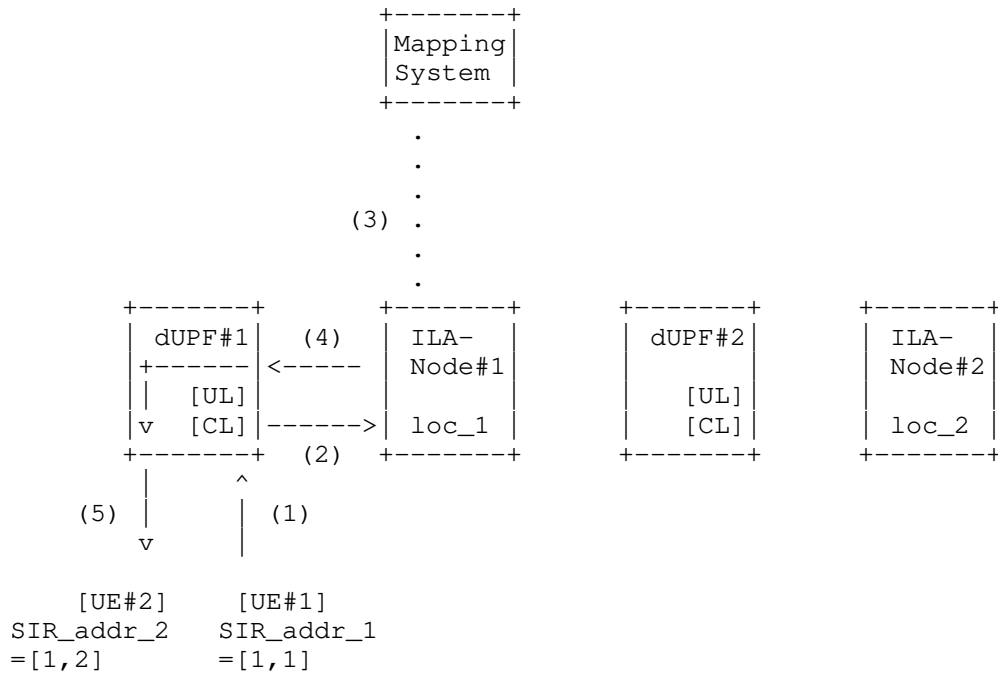


Figure 10: Procedure in Case B-1

- (0) Within this network, UEs are belonged to the same ILA domain, and the same SIR prefix is assigned to UEs. SIR_addr_1=[1,1] and SIR_addr_2=[1,2] are assigned to UE#1 and UE#2.
- (1) UE#1 sends packets to UE#2 with setting SIR_addr_2 as the destination IP address.

- (2) dUPF#1 monitors inner packet of received GTP-U packet and diverts it to ILA-Node#1 with decapsulation if the prefix of the destination address is SIR_prefix_1.
- (3) ILA-Node#1 updates own ID-to-LOC mapping table by interaction with the mapping system (if needed).
- (4) ILA-Node#1 obtains loc_2 as Locator of the ILA node#2 from the ID-to-LOC mapping table. ILA-Node#1 converts the prefixes of the source and destination addresses to loc_1 (Locator of id_1) and loc_2 (Locator of id_2). ILA-Node#1 sends the packet to the ILA-Node#2.
- (5) ILA-Node#2 receives the packet and converts the prefixes of the source and destination addresses to SIR_prefix_1, and then sends packets to dUPF#2.
- (6) dUPF#2 encapsulate packets with GTP-U header, and sends them to UE#2.

B.1.2. Case B-2: UEs allocated the same ILA node



Legend: $SIR_addr_x = [(SIR_Prefix), (Identifier)]$

Figure 11: Procedure in Case B-2

- (0) Within this network, UEs are belonged to the same ILA domain, and the same SIR prefix is assigned to UEs. $SIR_addr_1 = [1,1]$ and $SIR_addr_2 = [1,2]$ are assigned to UE#1 and UE#2.
- (1) UE#1 sends packets to UE#2 with setting SIR_addr_2 as the destination IP address.
- (2) dUPF#1 monitors inner packet of received GTP-U packet and diverts it to ILA-Node#1 with decapsulation if the prefix of the destination address is SIR_prefix_1 .
- (3) ILA-node#1 updates own ID-to-LOC mapping table by interaction with Mapping System (if needed).

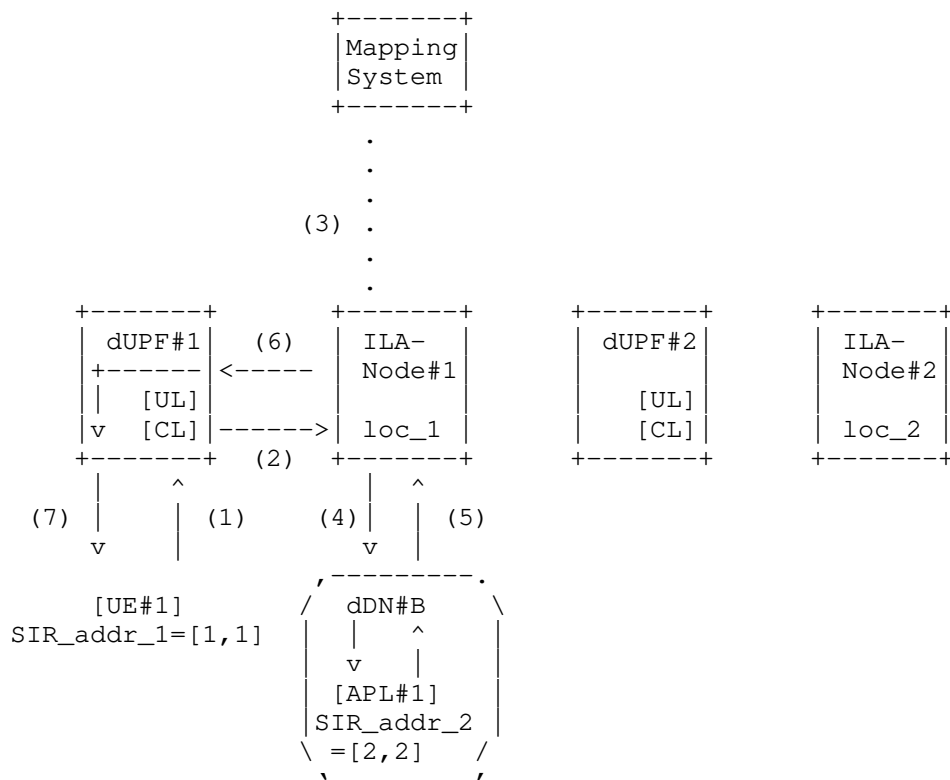
(4) ILA-Node#1 obtains loc_1 as Locator of ILA node#2 from the ID-to-LOC mapping table. Since loc_1 indicates itself, ILA-Node#1 sends the packets back to dUPF#1.

(5) dUPF#1 encapsulate packets with GTP-U, and sends them to UE#2.

B.2. UE-to-dDN Communication

The overview of this communication type is described in A.2.

B.2.1. Case B-3: UE communicates with neighbor dDN



Legend: SIR_addr_x=[(SIR_Prefix), (Identifier)]

Figure 12: Procedure in Case B-3

- (0) Within this network, UEs are belonged to the same ILA domain, and the same SIR prefix (SIR_prefix_1) are assigned to UEs. Applications in dDN#B are belonged to different ILA domain. and different SIR prefix (SIR_prefix_2) is assigned to these applications. SIR_addr_1=[1,1] and SIR_addr_2=[2,2] are assigned to UE#1 and APL#1. APL#1 is located in dDN#B.

Uplink Processes

- (1) UE#1 sends packets to APL#1 with setting SIR_addr_2 as the destination IP address.

- (2) dUPF#1 monitors inner packet of received GTP-U packet and diverts it to ILA-Node#1 with decapsulation if the prefix of the destination address is SIR_prefix_2.
- (3) ILA-Node#1 updates own ID-to-LOC mapping table by interaction with a mapping system (if needed). Or ILA-Node#1 may update its own table by a Map-Notify message when an APL is deployed or deleted in dDB#B.
- (4) ILA-Node#1 obtains loc_1 as Locator of id_2 from the ID-to-LOC mapping table. Since loc_1 indicates itself, ILA-Node#1 sends the packets to the dDN#B.

Downlink Processes

- (5) APL#1 in dDN#B sends packets to UE#1 with setting SIR_address_1 as the destination IP address.
- (6) ILA-Node#1 obtains loc_1 as Locator of id_1 from the ID-to-LOC mapping table. Since loc=1 indicates itself, ILA-Node#1 sends packets to dUPF#1.
- (7) dUPF#2 encapsulates packets with GTP-U, and sends them to UE#1.

B.2.2. Case B-4: UE communicates with non-neighbor dDN

- (2) dUPF#1 monitors inner packet of received GTP-U packet and diverts it to ILA-Node#1 with decapsulation if the prefix of the destination address is SIR_prefix_3.
- (3) ILA-Node#1 updates own ID-to-LOC mapping table by interaction with Mapping System (if needed).
- (4) ILA-Node#1 obtains loc_2 as Locator of id_3 from the ID-to-LOC mapping table. ILA-Node#1 converts the prefix of the source address to loc_1 (Locator of id_1), and the prefix of the destination address to loc_2 (Locator of id_3). ILA-Node#1 sends the packet to the ILA-Node#2.
- (5) ILA-Node#2 converts the prefix of the source address to SIR_prefix_1, and the prefix of the destination address to SIR_prefix_3, and then sends packets to dDN#C depending on its forwarding table.

Downlink Processes

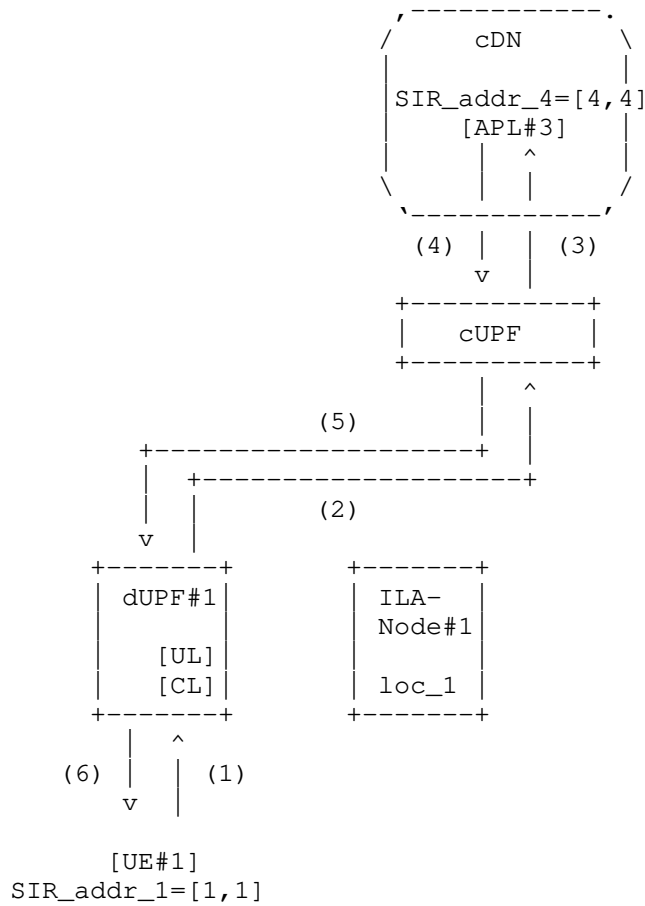
- (6) APL#2 sends packets to UE#1 with setting SIR_address_1 as the destination IP address.
- (7) ILA-Node#2 obtains loc_1 as Locator of id_1 from the ID-to-LOC mapping table. ILA-Node#2 converts the prefix of the source address to loc_2 (Locator of id_3), and the prefix of the destination address to loc_1 (Locator of id_1). ILA-Node#1 sends the packet to the ILA-Node#1.
- (8) ILA-Node#1 converts the prefix of the source address to SIR_prefix_3, and the prefix of the destination address to SIR_prefix_1, and then sends packets to d#UPF1 depending on its forwarding table.
- (9) dUPF#1 encapsulates the packets with GTP-U and sends packets to UE#1.

B.3. UE-to-cDN/Internet Communication

UE-to-cDN/Internet communication are basically achieved by GTP-U mechanism originally equipped in 3GPP 5GS architecture. ILA causes some limitation on IP addressing to UEs (e.g., all UEs in an ILA domain have the same SIR prefix), and thus some IP translation node such as NAT (Network Address Translation) may be required to enable UEs to access to external network. In this section, we describe processes of UE-to-cDN/Internet communication in the proposal architecture. In Internet communication, from aspect of privacy or routing with external network, SIR addresses assigned to UEs are

translated by NAT function deployed between dUPF and connection point.

B.3.1. Case B-5: Internet Communication



Legend: $SIR_addr_x=[(SIR_Prefix), (Identifier)]$

Figure 14: Procedure in Case B-5

- (0) Within this network, UEs are belonged to the same ILA domain, and the same SIR prefix (SIR_prefix_1) are assigned to UEs. Applications in cDN are belonged to different ILA domain. and

different SIR prefix (SIR_prefix_4) is assigned to these applications. SIR_addr_1=[1,1] and SIR_addr_4=[4,4] are assigned to UE#1 and APL#3. APL#3 is located in cDN.

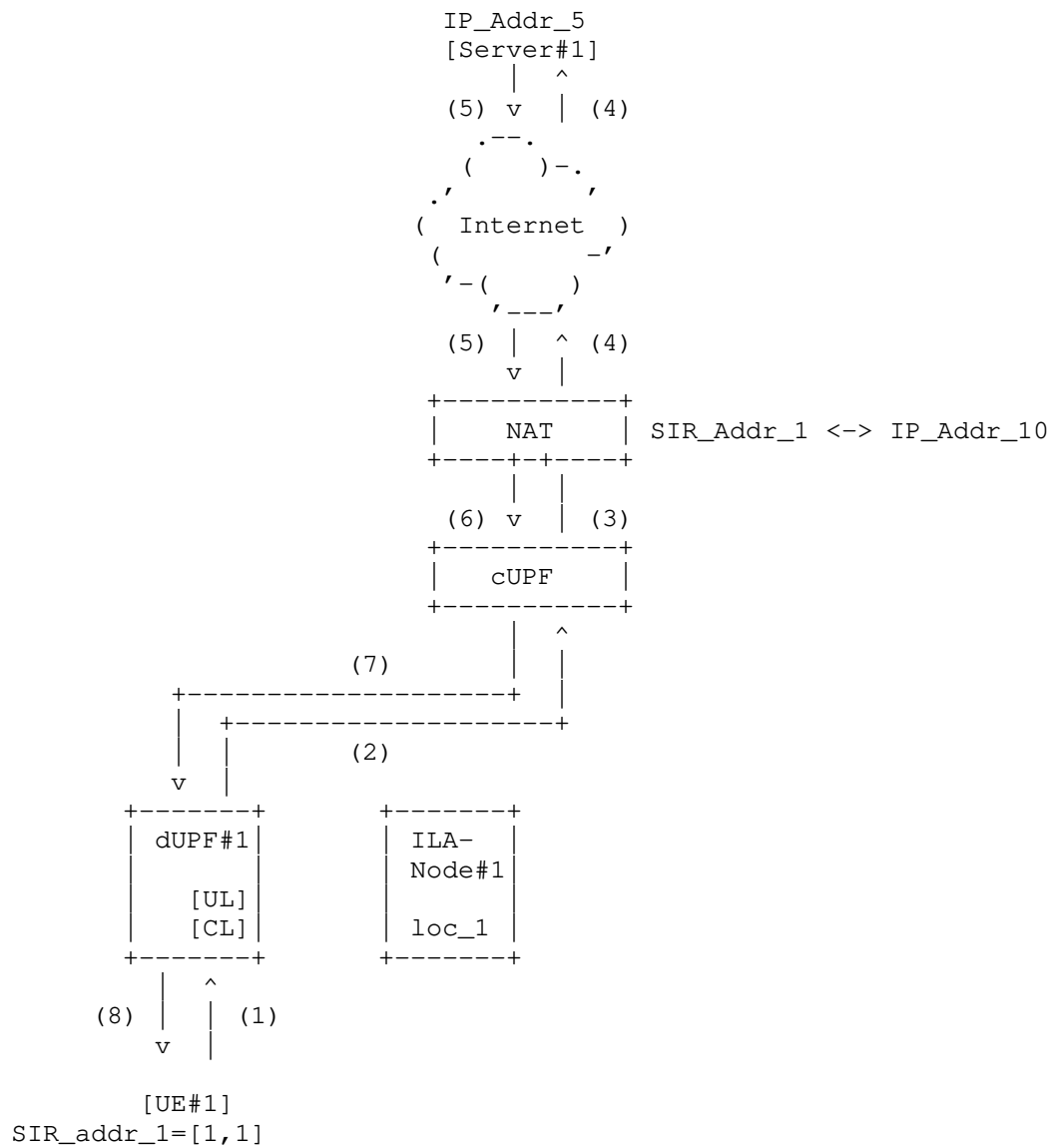
Uplink Processes

- (1) UE#1 sends packets to APL#3 with setting SIR_adrr_4 as the destination IP address.
- (2) dUPF#1 monitors inner of received GTP-U packets. Since the destination IP address (SIR_adder_4) does not hit the filter of ULCL, dUPF#1 re-encapsulates the packet to another GTP-U connecting to cUPF and forwards to cUPF.
- (3) cUPF decapusalates GTP-U packets and forwards them to APL#3 in cDN depending on its own forwarding table.

Downlink Processes

- (4) APL#3 in cDN sends packets to UE#1 with setting SIR_addr_1 as the destination IP address.
- (5) cUPF encapsulates the packets received from APL#3 and forwards them to dUPF#1 with GTP-U encapsulation depending on its own forwarding table.
- (6) dUPF re-encapsulates the packets to another GTP-U and forwards to UE#1.

B.3.2. Case B-6: Internet Communication



Legend: `SIR_addr_x=[(SIR_Prefix), (Identifier)]`

Figure 15: Procedure in Case B-6

- (0) Within this network, UEs are belonged to the same ILA domain, and the same SIR prefix (SIR_prefix_1) are assigned to UEs. SIR_addr_1=[1,1] assigned to UE#1 and server#1 has IP_addr_5. UE#1 communicate with server#1 over Internet.

Uplink Processes

- (1) UE#1 sends packets to server#1 with setting IP_adrr_5 as the destination IP address.
- (2) dUPF#1 monitors inner of received GTP-U packets. Since the destination IP address (SIR_adddr_4) does not hit the filter of ULCL, dUPF#1 re-encapsulates the packet to another GTP-U connecting to cUPF and forwards to cUPF.
- (3) cUPF decapusulates GTP-U packets and forwards them to Internet depending on its own forwarding table.
- (4) NAT translates SIR_adddr_1 of received packets to IP_addr_10 and the packets are forwarded to server#1 over Internet.

Downlink Processes

- (5) Server#1 sends packets to UE#1 with setting IP_adddr_1 as the destination IP address.
- (6) NAT translates IP_addr_10 of received packets to SIR_adddr_1, and packets are sent to cUPF.
- (7) cUPF encapsulates the packets with GTP-U and sends them to dUPF#1 depending on its own forwarding table.
- (8) dUPF re-encapsulates the packets to another GTP-U and forwards to UE#1.

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