QUIC Loss Detection & Congestion Control

draft-ietf-quic-recovery

TCPM, **IETF 103**

Goals

Discuss QUIC loss recovery mechanisms

No slides on congestion control, but can discuss (it's just NewReno)

Learn about egregious errors and blind spots TCPM has the right experts

Increase engagement with TCPM can do an update again at the next IETF



Non-Goals (for the next hour)

Re-design mechanisms

Re-litigate constants

Re-litigate QUIC's use of TCP standards 6298 and 5681 are non-normative references

... these things can be done, just don't do them right now



Overview

Some relevant QUIC details

Recovery mechanisms

Potential improvements / Open questions







QUIC Packet Numbers

Monotonically increasing 62-bit packet numbers (caveat: multiple PN spaces during connection setup)

Packet number DOES NOT indicate delivery ordering



QUIC Acknowledgements

ACK frame is encrypted and carried within QUIC packets

ACK frame contains:

largest acked one or more ack ranges "ack delay": T(ack send) - T(largest acked packet received) 3 ECN counts: #ECT(0), #ECT(1), #CE



Generating ACKs

SHOULD ACK every other packet subject to 25ms delayed ack timer

SHOULD ACK immediately if: Received packet number != largest received + 1 CE codepoint received

MAY process more packets before ACK allows less frequent acking



Notation

Packet: PN X packet with Packet Number X

Ack Frame: A X(K-L)(M-N) largest acked of X ack ranges K-L and M-N (Note: X > K, L, M, N)



Same, but different

Loss Detection

fast retransmit, early retransmit tail loss probe, RTO spurious RTO detection

Congestion Control

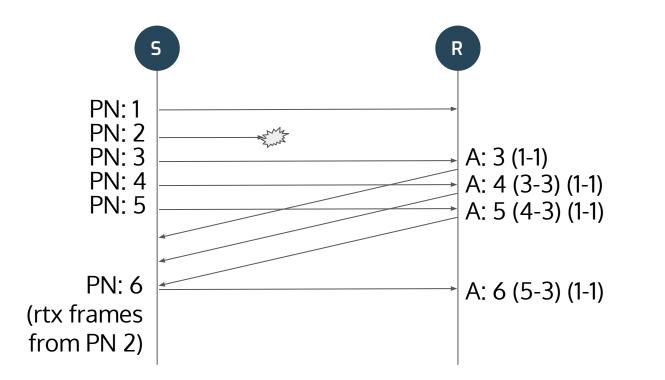
NewReno, but largest_acked ends recovery period



Recovery Mechanisms

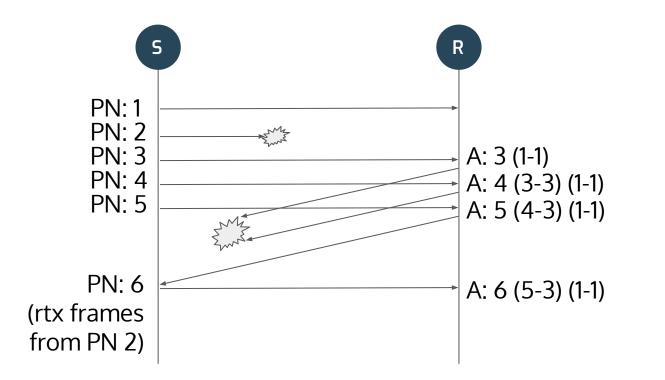


Fast Retransmit (Packet threshold)



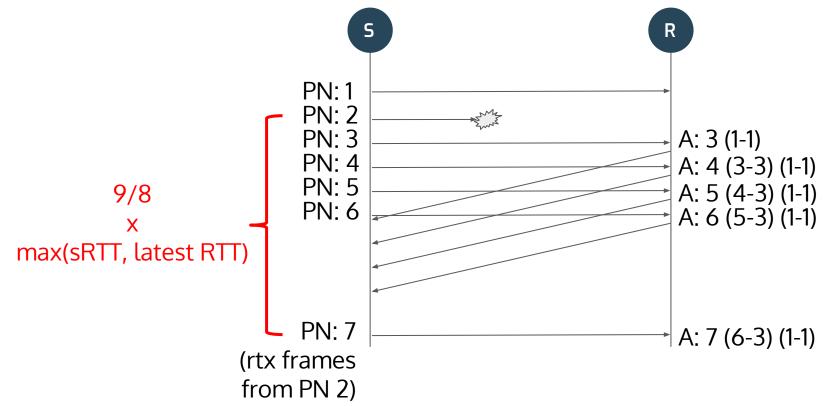


Fast Retransmit (FACK)



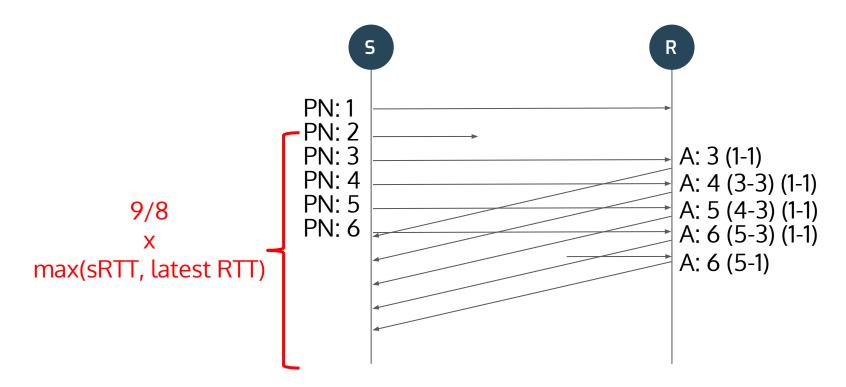


Fast Retransmit (Time threshold)



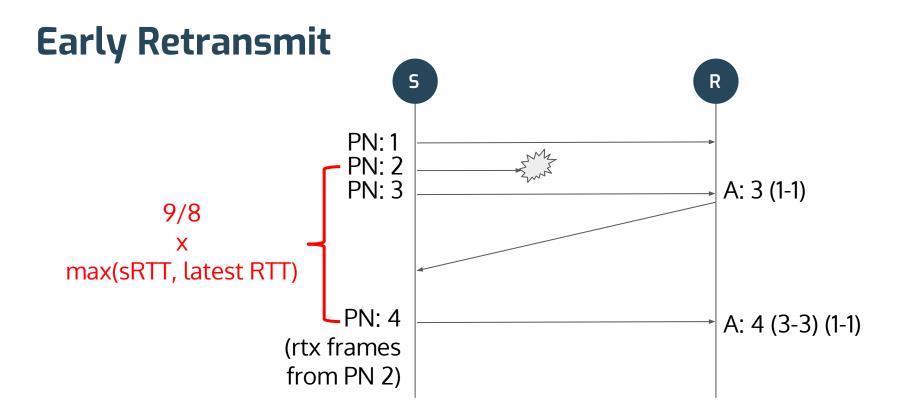


Fast Retransmit (Time threshold)



Time threshold allows reordering tolerance in packet space





Small delay allows for some reordering



RTT and Timeouts

RTT is RFC 6298, except for RTT sample:

rtt = now - largest_acked.sent_time - ack.ack_delay

max_ack_delay declared by both endpoints during handshake

Timeouts:

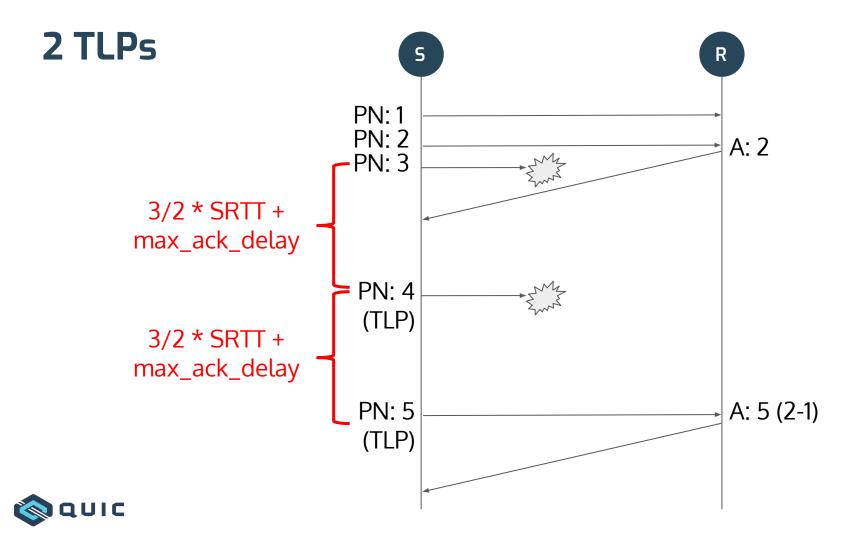
RTO = srtt + 4 * rttvar + max_ack_delay (min: 200ms) TLP = 1.5 * srtt + max_ack_delay (min: 10ms)

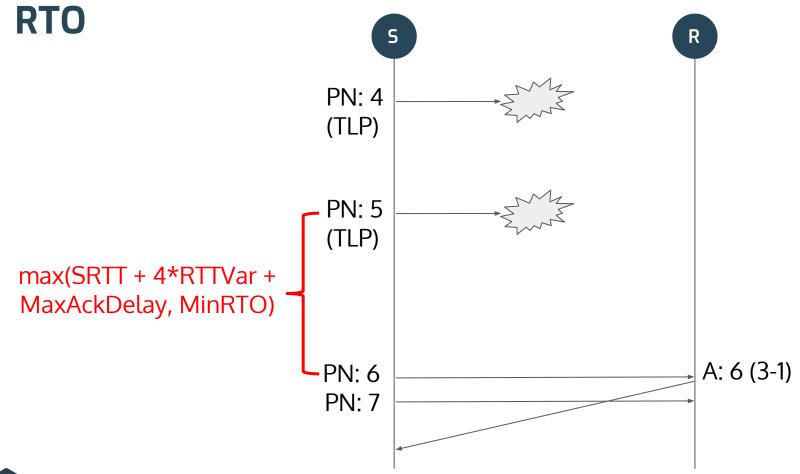


TLP S R PN: 1 PN: 2 And a PN: 3 A: 2 3/2 * SRTT + max_ack_delay PN: 4 A: 4 (2-1) (TLP: send new data or retx)

TLP always includes max_ack_delay

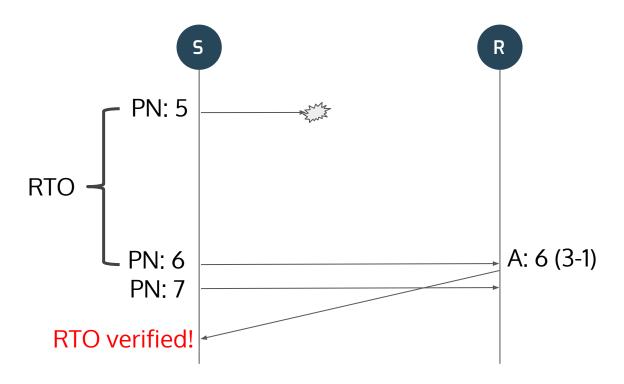






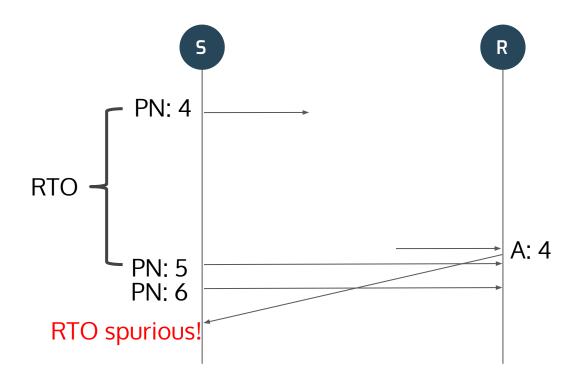


Spurious RTO Detection





Spurious RTO Detection





Spurious RTO

No congestion control actions on RTO

If any packet sent prior to RTO is newly acked declare RTO as spurious nothing more to be done

If all packets acked are ones sent after RTO declare RTO as verified congestion control actions (open issue: <u>#1966</u>)



Crypto Timeout

Set aggressively

before RTT sample: 200 ms after RTT sample: 2 x smoothed RTT set to max(timeout, kMinTLPTimeout)

Exponential backoff on consequent timeouts

Retransmit all outstanding crypto packets on timeout



Potential Improvements (NOT IN DRAFT!)



Generating fewer ACKs

SHOULD be sent immediately upon receipt of a second packet wireless drivers, middleboxes compress TCP acks should QUIC generate acks less frequently *by default?*



Removing MinRTO (#1017)

MaxAckDelay is explicitly communicated in the handshake

TCP's minRTO was to avoid spurious RTOs (RFC 6298) primary cost is bandwidth collapse when timer fires spurious RTO detection eliminates this cost

QUIC could remove the MinRTO since spurious RTOs have substantially lower cost



Potential Timeout Simplification

Combine TLP and RTO

both are similar, but no practical difference in QUIC

Issue TLP is commonly spurious

Why different than TCP cost of spurious RTO and TLP is low in QUIC



Fast retransmit

Should we do adaptive time thresholding?

How do we best use both packet and time thresholds? working on this now

