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Abstract

The applications of EVPN-based solutions ([RFC7432] and [RFC8365]) have become pervasive in Data Center, Service Provider, and Enterprise segments. It is being used for fabric overlays and inter-site connectivity in the Data Center market segment, for Layer-2, Layer-3, and IRB VPN services in the Service Provider market segment, and for fabric overlay and WAN connectivity in Enterprise networks. For Data Center and Enterprise applications, there is a need to provide inter-site and WAN connectivity over public Internet in a secured manner with same level of privacy, integrity, and authentication for tenant's traffic as IPsec tunneling using IKEv2. This document presents a solution where BGP point-to-multipoint signaling is leveraged for key and policy exchange among PE devices to create private pair-wise IPsec Security Associations without IKEv2 point-to-point signaling or any other direct peer-to-peer session establishment messages.

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Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

AC: Attachment Circuit.

ARP: Address Resolution Protocol.

BD: Broadcast Domain. As per [RFC7432], an EVI consists of a single or multiple BDs. In case of VLAN-bundle and VLAN-based service models (see [RFC7432]), a BD is equivalent to an EVI. In case of VLAN-aware bundle service model, an EVI contains multiple BDs. Also, in this document, BD and subnet are equivalent terms.

BD Route Target: refers to the Broadcast Domain assigned Route Target [RFC4364]. In case of VLAN-aware bundle service model, all the BD instances in the MAC-VRF share the same Route Target.

BT: Bridge Table. The instantiation of a BD in a MAC-VRF, as per [RFC7432].

DGW: Data Center Gateway.

Ethernet A-D route: Ethernet Auto-Discovery (A-D) route, as per [RFC7432].

Ethernet NVO tunnel: refers to Network Virtualization Overlay tunnels with Ethernet payload. Examples of this type of tunnels are VXLAN or GENEVE.

EVI: EVPN Instance spanning the NVE/PE devices that are participating

on that EVPN, as per [RFC7432].

EVPN: Ethernet Virtual Private Networks, as per [RFC7432].

GRE: Generic Routing Encapsulation.

GW IP: Gateway IP Address.

IPL: IP Prefix Length.

IP NVO tunnel: it refers to Network Virtualization Overlay tunnels with IP payload (no MAC header in the payload).

IP-VRF: A VPN Routing and Forwarding table for IP routes on an NVE/PE. The IP routes could be populated by EVPN and IP-VPN address families. An IP-VRF is also an instantiation of a layer 3 VPN in an NVE/PE.

IRB: Integrated Routing and Bridging interface. It connects an IP-VRF to a BD (or subnet).

MAC-VRF: A Virtual Routing and Forwarding table for Media Access Control (MAC) addresses on an NVE/PE, as per [RFC7432]. A MAC-VRF is also an instantiation of an EVI in an NVE/PE.

ML: MAC address length.

ND: Neighbor Discovery Protocol.

NVE: Network Virtualization Edge.

GENEVE: Generic Network Virtualization Encapsulation, [GENEVE].

NVO: Network Virtualization Overlays.

RT-2: EVPN route type 2, i.e., MAC/IP advertisement route, as defined in [RFC7432].

RT-5: EVPN route type 5, i.e., IP Prefix route. As defined in Section 3 of [EVPN-PREFIX].

SBD: Supplementary Broadcast Domain. A BD that does not have any ACs, only IRB interfaces, and it is used to provide connectivity among all the IP-VRFs of the tenant. The SBD is only required in IP-VRF- to-IP-VRF use-cases (see Section 4.4.).

SN: Subnet.

TS: Tenant System.

VA: Virtual Appliance.

VNI: Virtual Network Identifier. As in [RFC8365], the term is used as a representation of a 24-bit NVO instance identifier, with the understanding that VNI will refer to a VXLAN Network Identifier in VXLAN, or Virtual Network Identifier in GENEVE, etc. unless it is stated otherwise.

VTEP: VXLAN Termination End Point, as in [RFC7348].

VXLAN: Virtual Extensible LAN, as in [RFC7348].

This document also assumes familiarity with the terminology of [RFC7432], [RFC8365] and [RFC7365].

1 Introduction

The applications of EVPN-based solutions have become pervasive in Data Center, Service Provider, and Enterprise segments. It is being used for fabric overlays and inter-site connectivity in the Data Center market segment, for Layer-2, Layer-3, and IRB VPN services in the Service Provider market segment, and for fabric overlay and WAN connectivity in the Enterprise networks. For Data Center and Enterprise applications, there is a need to provide inter-site and WAN connectivity over public Internet in a secured manner with the same level of privacy, integrity, and authentication for tenant's traffic as used in IPsec tunneling using IKEv2. This document presents a solution where BGP point-to-multipoint signaling is leveraged for key and policy exchange among PE devices to create private pair-wise IPsec Security Associations without IKEv2 point-to-point signaling or any other direct peer-to-peer session establishment messages.

EVPN uses BGP as control-plane protocol for distribution of information needed for discovery of PEs participating in a VPN, discovery of PEs participating in a redundancy group, customer MAC addresses and IP prefixes/addresses, aliasing information, tunnel encapsulation types, multicast tunnel types, multicast group memberships, and other info. The advantages of using BGP control plane in EVPN are well understood including the following:

- 1) A full mesh of BGP sessions among PE devices can be avoided by using Route Reflector (RR) where a PE only needs to setup a single BGP session between itself and the RR as opposed to setting up N BGP sessions to N other remote PEs; therefore, reducing number of BGP sessions from $O(N^2)$ to $O(N)$ in the network. Furthermore, RR hierarchy can be leveraged to scale the number of BGP routes on the RR.
- 2) MP-BGP route filtering and constrained route distribution can be leveraged to ensure that the control-plane traffic for a given VPN is only distributed to the PEs participating in that VPN.

For setting up point-to-point security association (i.e., IPsec tunnel) between a pair of EVPN PEs, it is important to leverage BGP point-to-multipoint signaling architecture using the RR along with its route filtering and constrain mechanisms to achieve the performance and the scale needed for large number of security associations (IPsec tunnels) along with their frequent re-keying requirements. Using BGP signaling along with the RR (instead of peer-to-peer protocol such as IKEv2) reduces number of message exchanges needed for SAs establishment and maintenance from $O(N^2)$ to $O(N)$ in the network.

2 Requirements

The requirements for secured EVPN are captured in the following subsections.

2.1 Tenant's Layer-2 and Layer-3 data & control traffic

Tenant's layer-2 and layer-3 data and control traffic must be protected by IPsec cryptographic methods. This implies not only tenant's data traffic must be protected by IPsec but also tenant's control and routing information that are advertised in BGP must also be protected by IPsec. This in turn implies that BGP session must be protected by IPsec.

2.2 Tenant's Unicast & Multicast Data Protection

Tenant's layer-2 and layer-3 unicast traffic must be protected by IPsec. In addition to that, tenant's layer-2 broadcast, unknown unicast, and multicast traffic as well as tenant's layer-3 multicast traffic must be protected by IPsec when ingress replication or assisted replication are used. The use of BGP P2MP signaling for setting up P2MP SAs in P2MP multicast tunnels is for future study.

2.3 P2MP Signaling for SA setup and Maintenance

BGP P2MP signaling must be used for IPsec SAs setup and maintenance. The BGP signaling must follow P2MP signaling framework per [CONTROLLER-IKE] for IPsec SAs setup and maintenance in order to reduce the number of message exchanges from $O(N^2)$ to $O(N)$ among the participant PE devices.

2.4 Granularity of Security Association Tunnels

The solution must support the setup and maintenance of IPsec SAs at the following level of granularities:

- 1) Per PE: A single IPsec tunnel between a pair of PEs to be used for all tenants' traffic supported by the pair of PEs.
- 2) Per tenant: A single IPsec tunnel per tenant per pair of PEs. For example, if there are 1000 tenants supported on a pair of PEs, then 1000 IPsec tunnels are required between that pair of PEs.
- 3) Per subnet: A single IPsec tunnel per subnet (e.g., per VLAN/EVI) of a tenant on a pair of PEs.
- 4) Per IP address: A single IPsec tunnel per pair of IP addresses of a tenant on a pair of PEs.

5) Per MAC address: A single IPsec tunnel per pair of MAC addresses of a tenant on a pair of PEs.

6) Per Attachment Circuit: A single IPsec tunnel per pair of Attachment Circuits between a pair of PEs.

2.5 Support for Policy and DH-Group List

The solution must support a single policy and DH group for all SAs as well as supporting multiple policies and DH groups among the SAs.

3 Solution Description

This solution uses BGP P2MP signaling where an originating PE only send a message to the Route Reflector (RR) and then the RR reflects that message to the interested recipient PEs. The framework for such signaling is described in [CONTROLLER-IKE] and it is referred to as device-to-controller trust model. This trust model is significantly different than the traditional peer-to-peer trust model where a P2P signaling protocol such as IKEv2 [RFC7296] is used in which the PE devices directly authenticate each other and agree upon security policy and keying material to protect communications between themselves. The device-to-controller trust model leverages P2MP signaling via the controller (e.g., the RR) to achieve much better scale and performance for establishment and maintenance of large number of pair-wise Security Associations (SAs) among the PEs.

This device-to-controller trust model first secures the control channel between each device and the controller using peer-to-peer protocol such as IKEv2 [RFC7296] to establish P2P SAs between each PE and the RR. It then uses this secured control channel for P2MP signaling in establishment of P2P SAs between each pair of PE devices.

Each PE advertises to other PEs via the RR the information needed in establishment of pair-wise SAs between itself and every other remote PEs. These pieces of information are sent as Sub-TLVs of IPsec tunnel type in BGP Tunnel Encapsulation attribute. These Sub-TLVs are detailed in section 5 and are based on the DIM message components from [CONTROLLER-IKE] and the IKEv2 specification [RFC7296]. The IPsec tunnel TLVs along with its Sub-TLVs are sent along with the BGP route (NLRI) for a given level of granularity.

If only a single SA is required per pair of PE devices to multiplex user traffic for all tenants, then IPsec tunnel TLV is advertised along with IPv4 or IPv6 NLRI representing loopback address of the

originating PE. It should be noted that this is not a VPN route but rather an IPv4 or IPv6 route.

If a SA is required per tenant between a pair of PE devices, then IPsec tunnel TLV can be advertised along with EVPN IMET route representing the tenant or can be advertised along with a new EVPN route representing the tenant.

If a SA is required per tenant's subnet (e.g., per VLAN) between a pair of PE devices, then IPsec tunnel TLV is advertised along with EVPN IMET route.

If a SA is required between a pair of tenant's devices represented by a pair of IP addresses, then IPsec tunnel TLV is advertised along with EVPN IP Prefix Advertisement Route or EVPN MAC/IP Advertisement route.

If a SA is required between a pair of tenant's devices represented by a pair of MAC addresses, then IPsec tunnel TLV is advertised along with EVPN MAC/IP Advertisement route.

If a SA is required between a pair of Attachment Circuits (ACs) on two PE devices (where an AC can be represented by <VLAN, port>), then IPsec tunnel TLV is advertised along with EVPN Ethernet AD route.

3.1 Inheritance of Security Policies

Operationally, it is easy to configure a security association between a pair of PEs using BGP signaling. This is the default security association that is used for traffic that flows between peers. However, in the event more finer granularity of security association is desired on the traffic flows, it is possible to set up SAs between a pair of tenants, a pair of subnets within a tenant, a pair of IPs between a subnet, and a pair of MACs between a subnet using the appropriate EVPN routes as described above. In the event, there are no security TLVs associated with an EVPN route, there is a strict order in the manner security associations are inherited for such a route. This results in an EVPN route inheriting the security associations of the parent in a hierarchical fashion. For example, traffic between an IP pair is protected using security TLVs announced along with the EVPN IP Prefix Advertisement Route or EVPN MAC/IP Advertisement route as a first choice. If such TLVs are missing with the associated route, then one checks to see if the subnets the IPs are associated with has security TLVs with the EVPN IMET route. If they are present, those associations are used in securing the traffic. In the absence of them, the peer security associations are used. The order in which security associations are inherited are from the granular to the coarser, namely, IP/MAC associated TLVs with the

EVPN route being the first preference, and the subnet, the tenant, and the peer associations preferred in that fashion.

It should be noted that when a security association is made it is possible for it to be re-used by a large number of traffic flows. For example, a tenant security association may be associated with a number of child subnet routes. Clearly it is mandatory to keep a tenant security association alive, if there are one or more subnet routes that want to use that association. Logically, the security associations between a pair of entities creates a single secure tunnel. It is thus possible to classify the incoming traffic in the most granular sense {IP/MAC, subnet, tenant, peer} to a particular secure tunnel that falls within its route hierarchy. The policy that is applied to such traffic is independent from its use of an existing or a new secure tunnel. It is clear that since any number of classified traffic flows can use a security association, such a security association will not be torn down, if at least there is one policy using such a secure tunnel.

3.2 Distribution of Public Keys and Policies

One of the requirements for this solution is to support a single DH group and a single policy for all SAs as well as to support multiple DH groups and policies among the SAs. The following subsections describe what pieces of information (what Sub-TLVs) are needed to be exchanged to support a single DH group and a single policy versus multiple DH groups and multiple policies.

3.2.1 Minimal DIM

For SA establishment, at the minimum, a PE needs to advertise to other PEs, its DIM values as specified in [CONTROLLER-IKE]. These include:

ID	Tunnel ID
N	Nonce
RC	Rekey Counter
I	Indication of initial policy distribution
KE	DH public value.

When this minimal set of DIM values is sent, then it is assumed that all peer PEs share the same policy for which DH group to use, as well as which IPsec SA policy to employ. Section 5.1 defines the Minimal DIM sub-TLV as part of IPsec tunnel TLV in BGP Tunnel Encapsulation Attribute.

3.2.2 Multiple Policies

There can be scenarios for which there is a need to have multiple policy options. This can happen when there is a need for policy change and smooth migration among all PE devices to the new policy is required. It can also happen if different PE devices have different capabilities within the network. In these scenarios, PE devices need to be able to choose the correct policy to use for each other. This multi-policy scheme is described in section 6 of [CONTROLLER-IKE]. In order to support this multi-policy feature, a PE device MUST distribute a policy list. This list consists of multiple distinct policies in order of preference, where the first policy is the most preferred one. The receiving PE selects the policy by taking the received list (starting with the first policy) and comparing that against its own list and choosing the first one found in common. If there is no match, this indicates a configuration error and the PEs MUST NOT establish new SAs until a message is received that does produce a match.

3.2.2.1 Multiple DH-groups

It can be the case that not all peers use the same DH group. When multiple DH groups are supported, the peer may include multiple KE Sub-TLVs. The order of the KE Sub-TLVs determines the preference. The preference and selection methods are specified in Section 6 of [CONTROLLER-IKE].

3.2.2.2 Multiple or Single ESP SA policies

In order to specify an ESP SA Policy, a DIM may include one or more SA Sub-TLVs. When all peers are configured by a controller with the same ESP SA policy, they MAY leave the SA out of the DIM. This minimizes messaging when group configuration is static and known. However, it may also be desirable to include the SA. If a single SA is included, the peer is indicating what ESP SA policy it uses, but is not willing to negotiate. If multiple SA Sub-TLVs are included, the peer is indicating that it is willing to negotiate. The order of the SA Sub-TLVs determines the preference. The preference and selection methods are specified in Section 6 of [CONTROLLER-IKE].

3.3 Initial IPsec SAs Generation

The procedure for generation of initial IPsec SAs is described in section 3 of [CONTROLLER-IKE]. This section gives a summary of it in context of BGP signaling. When a PE device first comes up and wants to setup an IPsec SA between itself and each of the interested remote PEs, it generates a DH pair along for each [what word here? "tenant"?] using an algorithm defined in the IKEv2 Diffie-Hellman

Group Transform IDs [IKEv2-IANA]. The originating PE distributes the DH public value along with the other values in the DIM (using IPsec Tunnel TLV in Tunnel Encapsulation Attribute) to other remote PEs via the RR. Each receiving PE uses this DH public number and the corresponding nonce in creation of IPsec SA pair to the originating PE - i.e., an outbound SA and an inbound SA. The detail procedures are described in section 5.2 of [CONTROLLER-IKE].

3.4 Re-Keying

A PE can initiate re-keying at any time due to local time or volume based policy or due to the result of cipher counter nearing its final value. The rekey process is performed individually for each remote PE. If rekeying is performed with multiple PEs simultaneously, then the decision process and rules described in this rekey are performed independently for each PE. Section 4 of [CONTROLLER-IKE] describes this rekeying process in details and gives examples for a single IPsec device (e.g., a single PE) rekey versus multiple PE devices rekey simultaneously.

3.5 IPsec Databases

The Peer Authorization Database (PAD), the Security Policy Database (SPD), and the Security Association Database (SAD) all need to be setup as defined in the IPsec Security Architecture [RFC4301]. Section 5 of [CONTROLLER-IKE] gives a summary description of how these databases are setup for the controller-based model where key is exchanged via P2MP signaling via the controller (i.e., the RR) and the policy can be either signaled via the RR (in case of multiple policies) or configured by the management station (in case of single policy).

4 Encapsulation

Vast majority of Encapsulation for Network Virtualization Overlay (NVO) networks in deployment are based on UDP/IP with UDP destination port ID indicating the type of NVO encapsulation (e.g., VxLAN, GPE, GENEVE, GUE) and UDP source port ID representing flow entropy for load-balancing of the traffic within the fabric based on n-tuple that includes UDP header. When encrypting NVO encapsulated packets using IP Encapsulating Security Payload (ESP), the following two options can be used: a) adding a UDP header before ESP header (e.g., UDP header in clear) and b) no UDP header before ESP header (e.g., standard ESP encapsulation). The following subsection describe these encapsulation in further details.

4.1 Standard ESP Encapsulation

When standard IP Encapsulating Security Payload (ESP) is used (without outer UDP header) for encryption of NVO packets, it is used in transport mode as depicted below. When such encapsulation is used, for BGP signaling, the Tunnel Type of Tunnel Encapsulation TLV is set to ESP-Transport and the Tunnel Type of Encapsulation Extended Community is set to NVO encapsulation type (e.g., VxLAN, GENEVE, GPE, etc.). This implies that the customer packets are first encapsulated using NVO encapsulation type and then it is further encapsulated & encrypted using ESP-Transport mode.

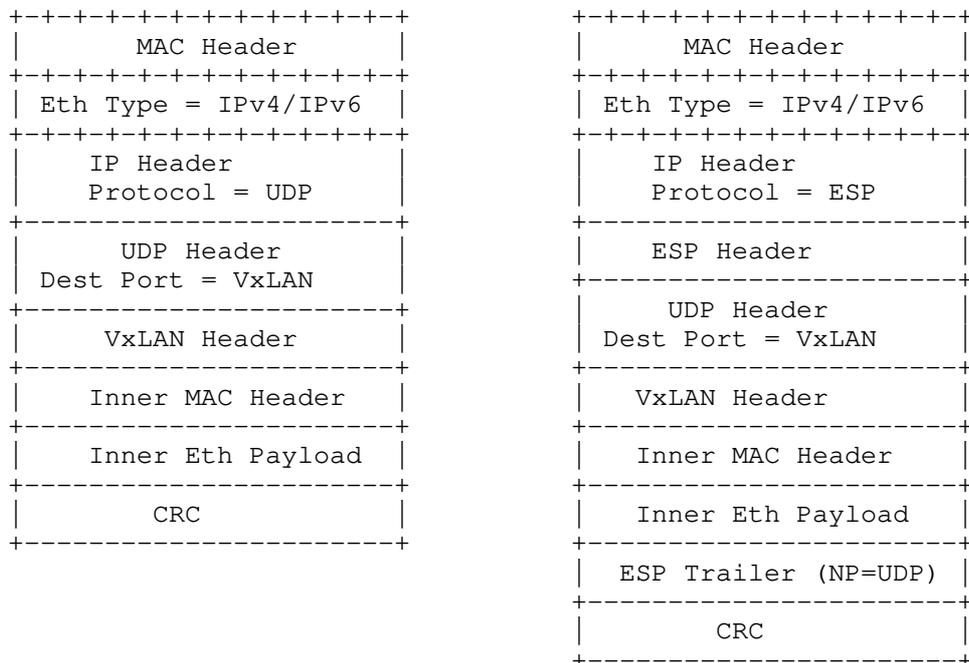


Figure 3: VxLAN Encapsulation within ESP

4.2 ESP Encapsulation within UDP packet

In scenarios where NAT traversal is required ([RFC3948]) or where load balancing using UDP header is required, then ESP encapsulation within UDP packet as depicted in the following figure is used. The ESP for NVO applications is in transport mode. The outer UDP header

(before the ESP header) has its source port set to flow entropy and its destination port set to 4500 (indicating ESP header follows). A non-zero SPI value in ESP header implies that this is a data packet (i.e., it is not an IKE packet). The Next Protocol field in the ESP trailer indicates what follows the ESP header, is a UDP header. This inner UDP header has a destination port ID that identifies NVO encapsulation type (e.g., VxLAN). Optimization of this packet format where only a single UDP header is used (only the outer UDP header) is for future study.

When such encapsulation is used, for BGP signaling, the Tunnel Type of Tunnel Encapsulation TLV is set to ESP-in-UDP-Transport and the Tunnel Type of Encapsulation Extended Community is set to NVO encapsulation type (e.g., VxLAN, GENEVE, GPE, etc.). This implies that the customer packets are first encapsulated using NVO encapsulation type and then it is further encapsulated & encrypted using ESP-in-UDP with Transport mode.

[RFC3948]

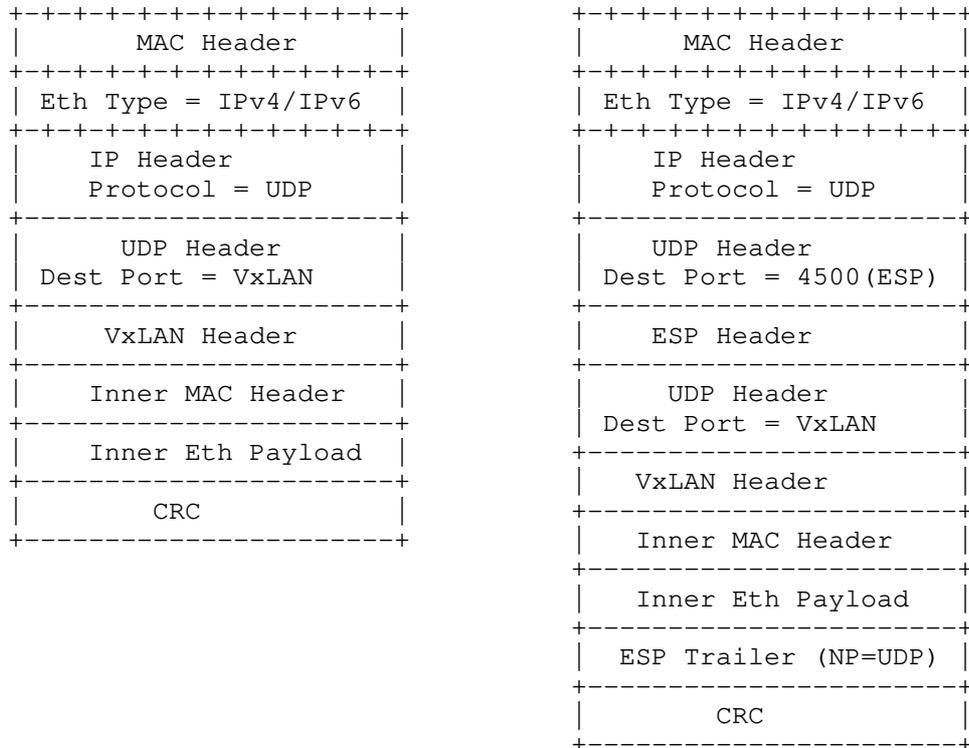


Figure 4: VxLAN Encapsulation within ESP Within UDP

5 BGP Encoding

This document defines two new Tunnel Types along with its associated sub-TLVs for The Tunnel Encapsulation Attribute [TUNNEL-ENCAP]. These tunnel types correspond to ESP-Transport and ESP-in-UDP-Transport as described in section 4. The following sub-TLVs apply to both tunnel types unless stated otherwise.

5.1 The Base (Minimal Set) DIM Sub-TLV

The Base DIM is described in 3.2.1. One and only one Base DIM may be sent in the IPsec Tunnel TLV.

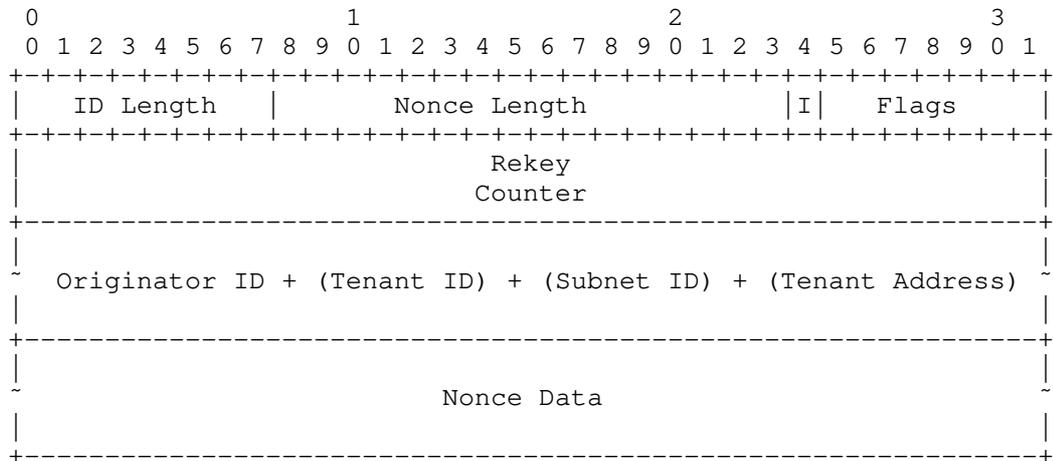


Figure 5: The Base DIM Sub-TLV

ID Length (16 bits) is the length of the Originator ID + (Tenant ID) + (Subnet ID) + (Tenant Address) in bytes.

Nonce Length (8 bits) is the length of the Nonce Data in bytes

I (1 bit) is the initial contact flag from [CONTROLLER-IKE]

Flags (7 bits) are reserved and MUST be set to zero on transmit and ignored on receipt.

The Rekey Counter is a 64 bit rekey counter as specified in [CONTROLLER-IKE]

The Originator ID + (Tenant ID) + (Subnet ID) + (Tenant Address) is the tunnel identifier and uniquely identifies the tunnel. Depending on the granularity of the tunnel, the fields in () may not be used - i.e., for a tunnel at the PE level of granularity, only Originator ID is required.

The Nonce Data is the nonce described in [CONTROLLER-IKE]. Its length is a multiple of 32 bits. Nonce lengths should be chosen to meet minimum requirements described in IKEv2 [RFC7296].

5.2 Key Exchange Sub-TLV

The KE Sub-TLV is described in 3.2.1 and 3.2.2.1. A KE is always required. One or more KE Sub-TLVs may be included in the IPSec Tunnel TLV.

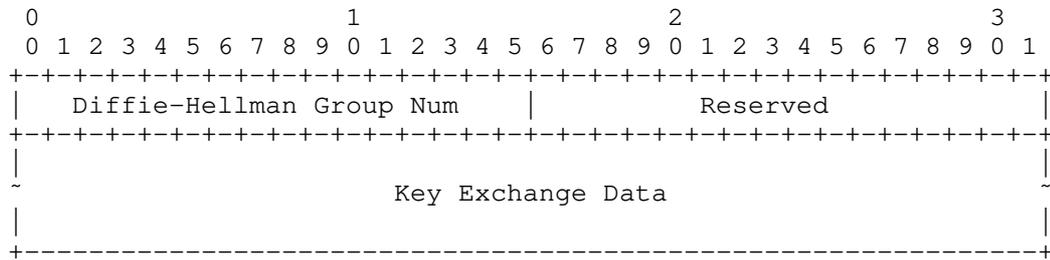


Figure 6: Key Exchange Sub-TLV

Diffie-Hellman Group Num 916 bits) identifies the Diffie-Hellman group in the Key Exchange Data was computed. Diffie-Hellman group numbers are discussed in IKEv2 [RFC7296] Appendix B and [RFC5114].

The Key Exchange payload is constructed by copying one's Diffie-Hellman public value into the "Key Exchange Data" portion of the payload. The length of the Diffie-Hellman public value is described for MOPD groups in [RFC7296] and for ECP groups in [RFC4753].

5.3 ESP SA Proposals Sub-TLV

The SA Sub-TLV is described in 3.2.2.2. Zero or more SA Sub-TLVs may be included in the IPSec Tunnel TLV.

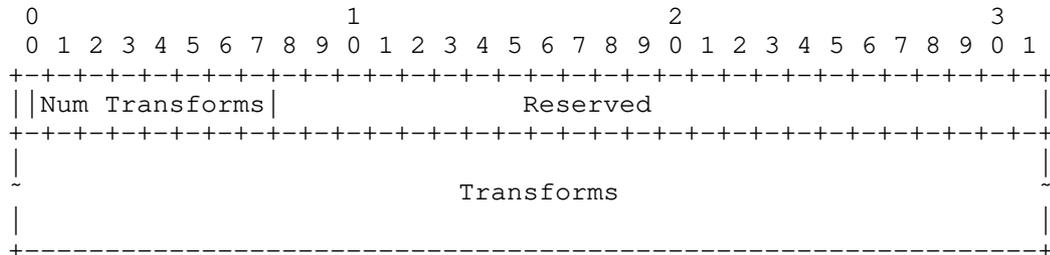


Figure 8: ESP SA Proposals Sub-TLV

Num Transforms is the number of transforms included.

Reserved is not used and MUST be set to zero on transmit and MUST be ignored on receipt.

5.3.1 Transform Substructure

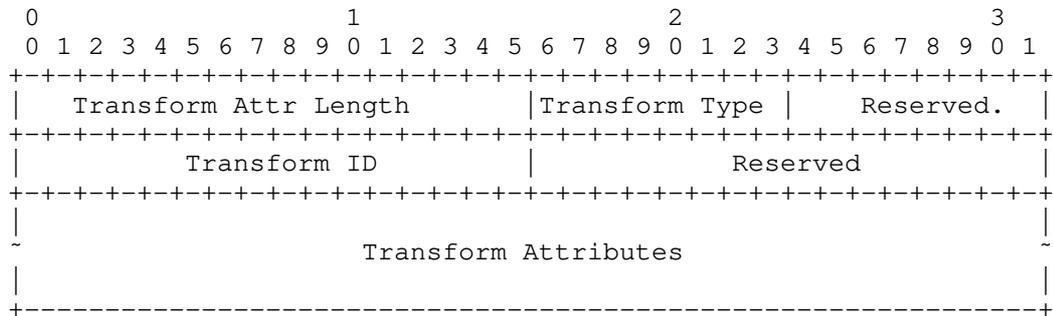


Figure 9: Transform Substructure Sub-TLV

The Transform Attr Length is the length of the Transform Attributes field.

The Transform Type is from Section 3.3.2 of [RFC7296] and [IKEV2IANA]. Only the values ENCR, INTEG, and ESN are allowed.

The Transform ID specifies the transform identification value from [IKEV2IANA].

Reserved is unused and MUST be zero on transmit and MUST be ignored on receipt.

The Transform Attributes are taken directly from 3.3.5 of [RFC7296].

6 Applicability to other VPN types

Although P2MP BGP signaling for establishment and maintenance of SAs among PE devices is described in this document in context of EVPN, there is no reason why it cannot be extended to other VPN technologies such as IP-VPN [RFC4364], VPLS [RFC4761] & [RFC4762], and MVPN [RFC6513] & [RFC6514] with ingress replication. The reason EVPN has been chosen is because of its pervasiveness in DC, SP, and Enterprise applications and because of its ability to support SA establishment at different granularity levels such as: per PE, Per tenant, per subnet, per Ethernet Segment, per IP address, and per MAC. For other VPN technology types, a much smaller granularity levels can be supported. For example for VPLS, only the granularity of per PE and per subnet can be supported. For per-PE granularity level, the mechanism is the same among all the VPN technologies as IPsec tunnel type (and its associated TLV and sub-TLVs) are sent along with the PE's loopback IPv4 (or IPv6) address. For VPLS, if per-subnet (per bridge domain) granularity level needs to be supported, then the IPsec tunnel type and TLV are sent along with

VPLS AD route.

The following table lists what level of granularity can be supported by a given VPN technology and with what BGP route.

Functionality	EVPN	IP-VPN	MVPN	VPLS
per PE	IPv4/v6 route	IPv4/v6 route	IPv4/v6 rte	IPv4/v6
per tenant	IMET (or new)	lpbk (or new)	I-PMSI	N/A
per subnet	IMET	N/A	N/A	VPLS AD
per IP	EVPN RT2/RT5	VPN IP rt	*,G or S,G	N/A
per MAC	EVPN RT2	N/A	N/A	N/A

7 Acknowledgements

8 Security Considerations

9 IANA Considerations

A new transitive extended community Type of 0x06 and Sub-Type of TBD for EVPN Attachment Circuit Extended Community needs to be allocated by IANA.

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