

Packet Loss Signaling for Encrypted Protocols

draft-ferrieuxhamchaoui-tsvwg-lossbits

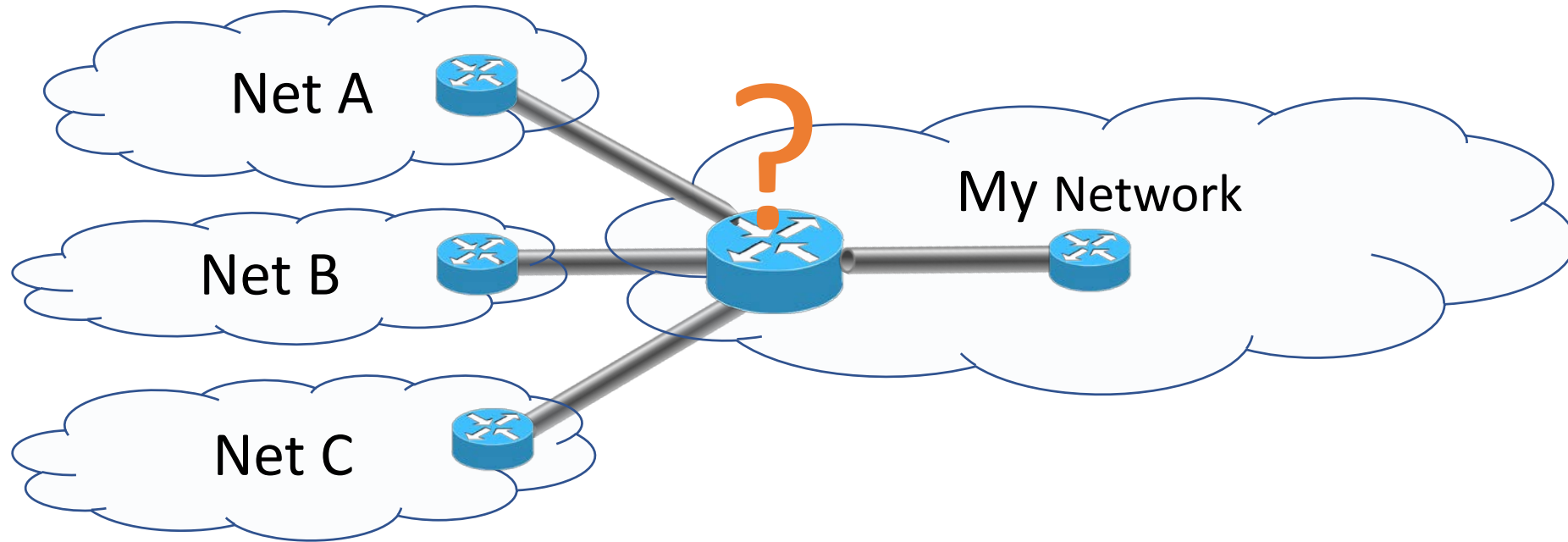
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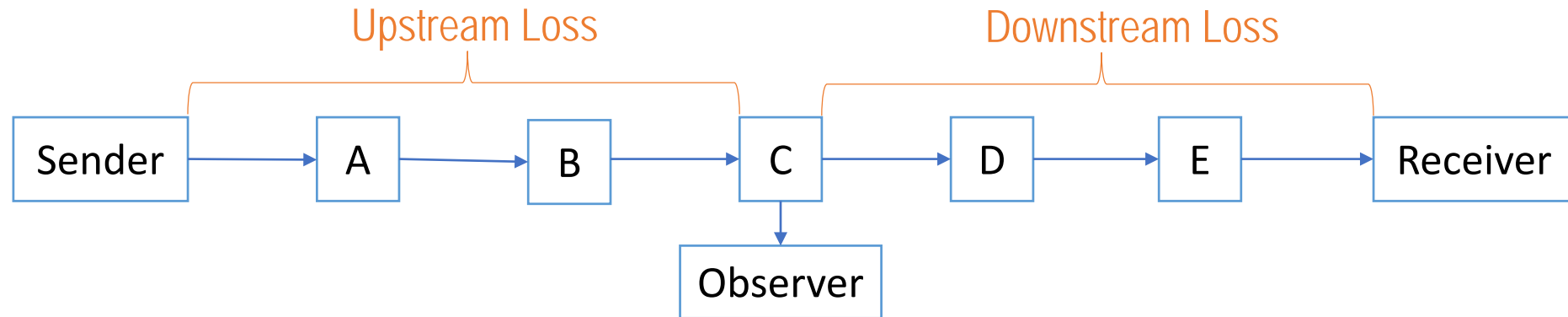
Igor Lubashev – Akamai

Motivation: Loss Detection/Localization Matters

*Networks can look like dumb pipes,
only if a plumber can find leaks and patch them quickly*



Motivation: Loss Detection/Localization Matters



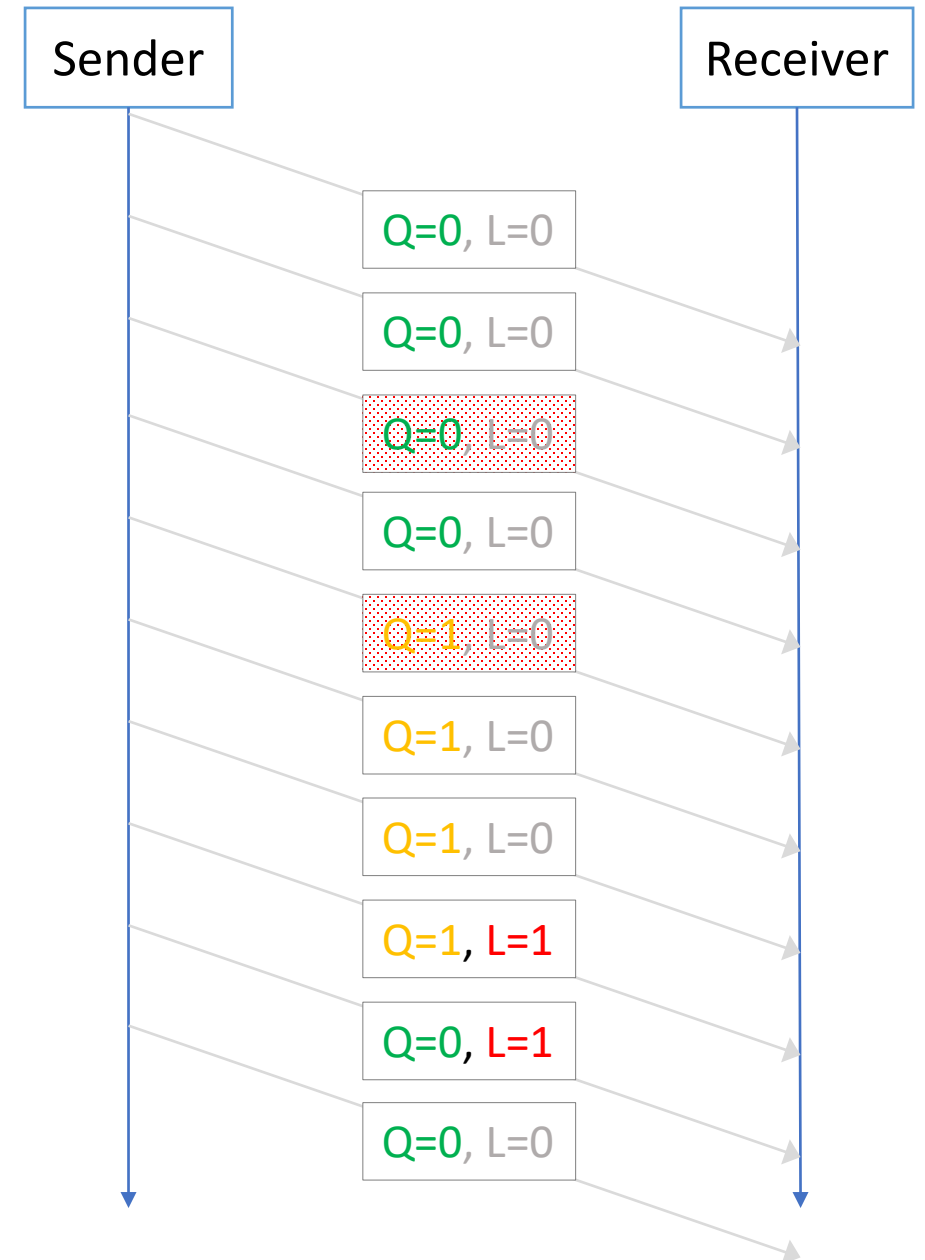
- TCP: observe seq# (and ack#/sack#, if path is symmetric)
- Transport with encrypted headers: ☹
 - QUIC has a “latency Spin bit”, so you may get an RTT estimate but *not* loss
- “*Just observe similar TCP flows*” is not a good answer

I am not a Network Operator. Why Should I Care?

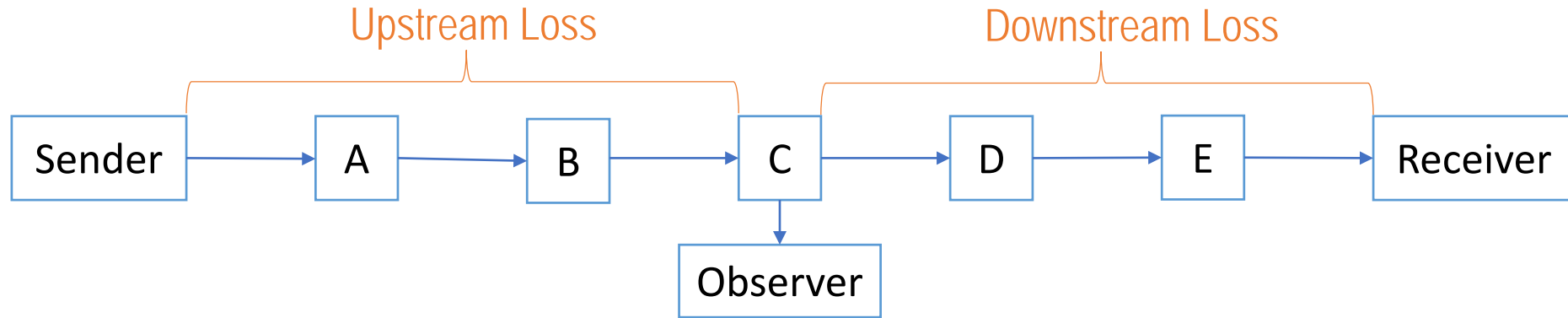
- If you publish content or services, you derive some benefit from those sites being available and fast.
- If you are a CDN, your customers pay you to ensure their sites are available and fast and to take care of “those Internet issues”.

Proposal: Two “Loss bits”

- **Q**: The “sSquare signal” bit is toggled every N outgoing packets (akin to *color* in RFC 8321)
- **L**: The “Loss event” bit is 1 when “Unreported Loss Counter” (ULC) > 0
 - ULC is incremented for each packet deemed lost by the protocol
 - ULC is decremented for each packet sent with L=1



Loss Calculation



- End-to-End loss (e)

e = fraction of packets with $L=1$

- Upstream loss (u)

$$u = 1 - \frac{\text{average \# of observed packets in a block (same } Q)}{\text{size of the block}}$$

- Downstream loss (d)

$$(1 - u)(1 - d) = 1 - e$$

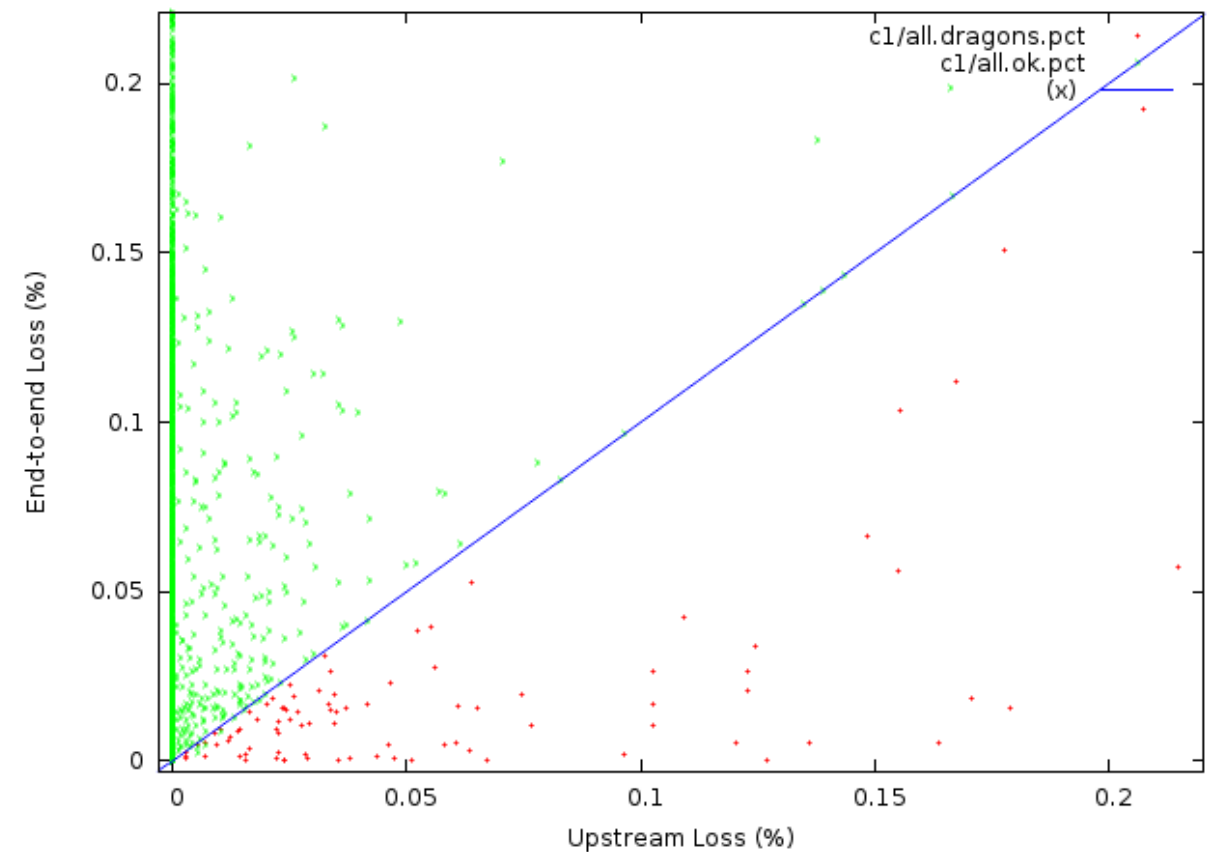
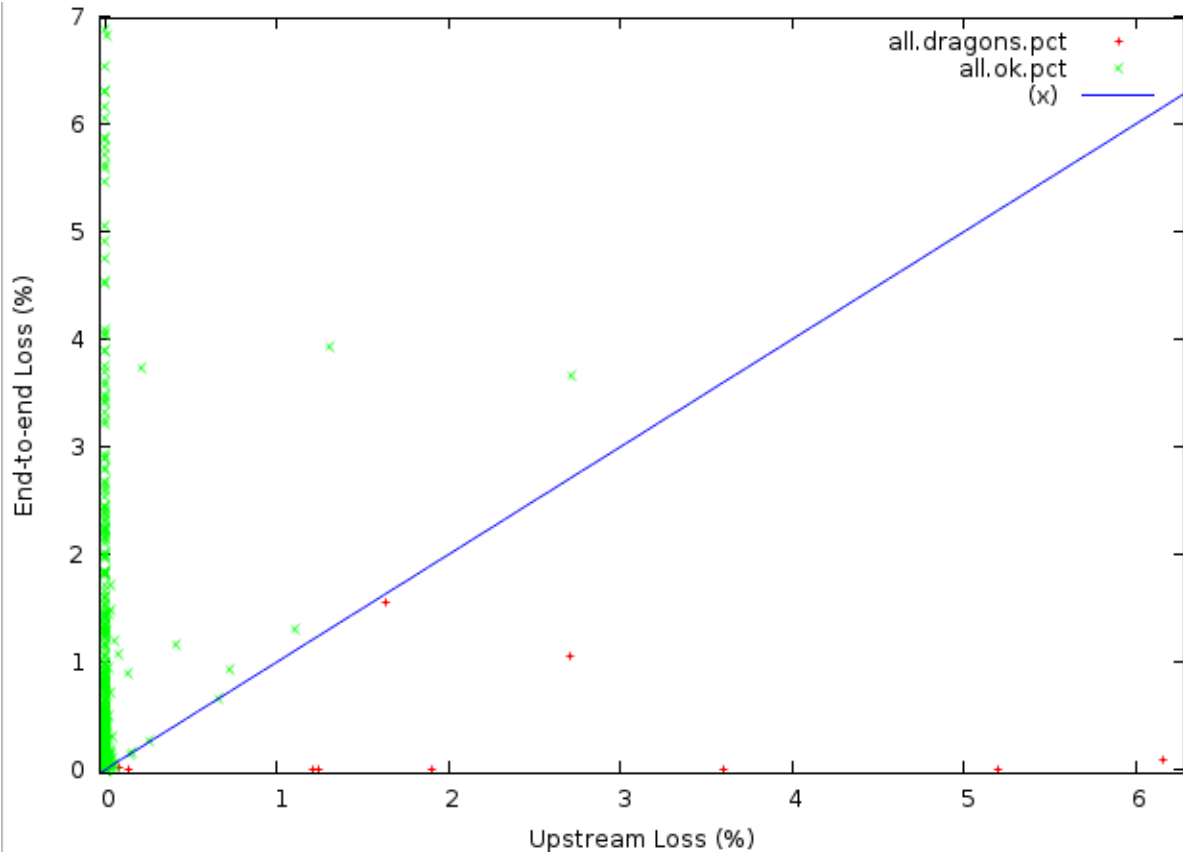
$$d = \frac{e - u}{1 - u} \approx e - u$$

Which Protocol Header?

- This draft requires answers to:
 - Question 1: “Do we need loss detection by non-endpoints?”
 - Question 2: “If we do, are Q & L bits fit for the purpose?”
- If “Yes” to both of the above, we can find a home for the bits in a subsequent draft (possibly in a different WG):
 - IPv4/IPv6 header?
 - IPv4 options / IPv6 HBH option?
 - UDP trailer?
 - QUIC header?

Experimental Data – Akamai to Orange (4 countries)

- Q&L bits are in *ip.ttl* \gg 6 (and *ip6.hoplimit* \gg 6)
- A lot more data and discussion in maprg tomorrow at 10am



Privacy and Ossification

- Protecting Privacy

- Explicit signal means less information leakage (RFC 8558)
- Separate counters for separate flows, subflows, paths, QUIC connection IDs, ... to prevent loss signals used to link multiple connections to the same device

- Ossification Resistance

- Loss signals are not integral protocol bits, so they can be greased, if desired
- QUIC latency spin bit is an example:
 - can mandate random-looking values for Q&L bits if unused
 - can mandate to not using for a certain portion of connections

Getting in Touch

- Mailing List: ietf-loss-bits@googlegroups.com
- Data Discussion on Friday at 10am (mapgr)
- draft-ferrieuxhamchaoui-tsvwg-lossbits