lispers.net NAT Implementation
draft-farinacci-lisp-simple-nat-00

IETF Madrid (Virtual) LISP WG
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Purpose of Draft

- To describe a simple version of NAT-traversal

- Based on a subset of procedures and message formats from:

    ```
    draft-ermagan-lisp-nat-traversal
    ```

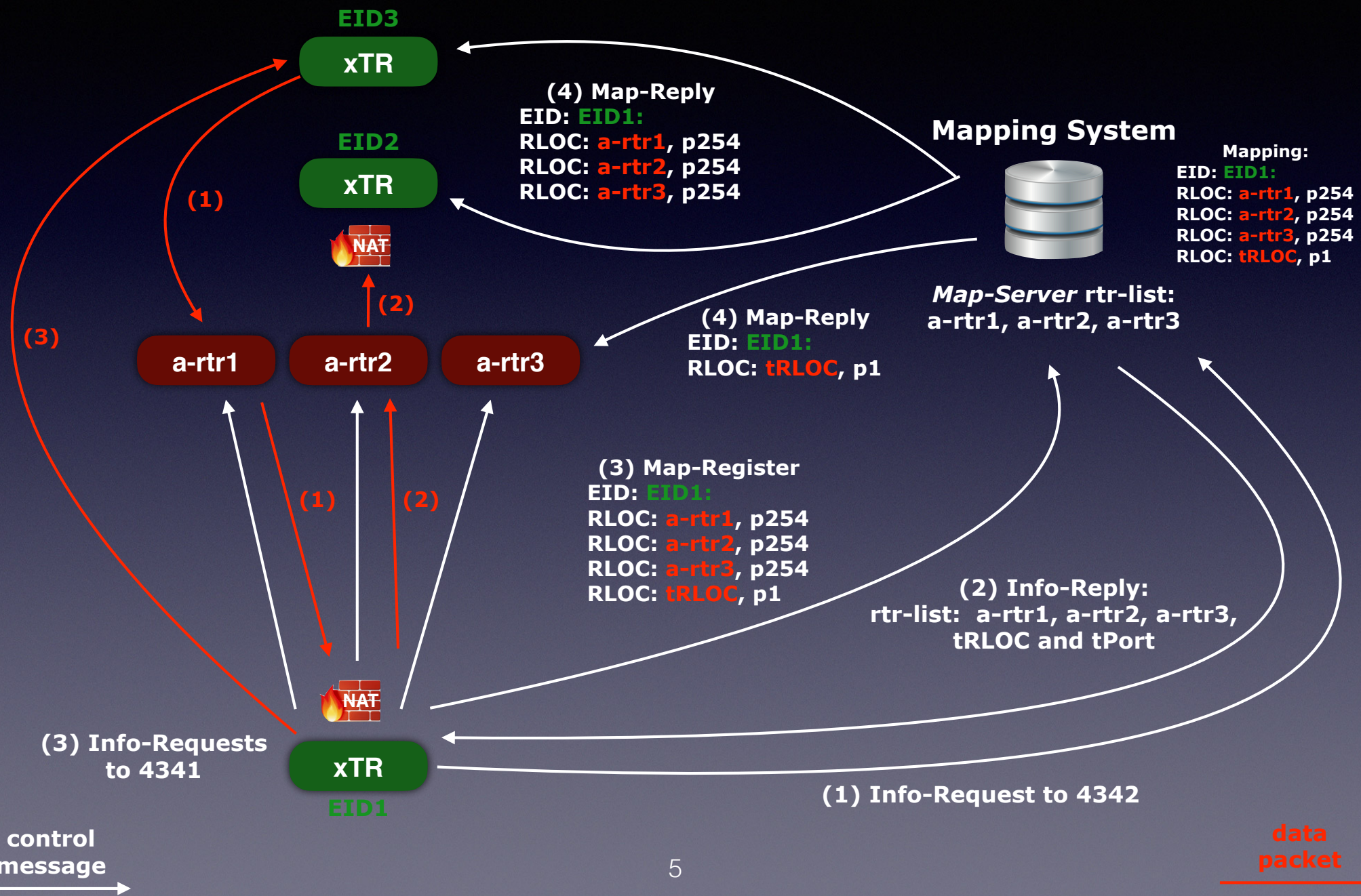- There has been *lispers.net* deployments since 2014

Requirements

- The xTR must function when it resides 1 or more hops away from a NAT device

- The xTR must function when it runs as a LISP-MN, in a VM, or a container (logically or physically)

- The xTR must function when it resides behind multiple NAT devices between itself and an RTR

- The xTR must function when it is multi-homed behind different NAT devices

- RTRs and PxTRs must function behind NATs

- *lig* must function when requests are sent from behind NATs

- IPv4 unicast/multicast and IPv6 unicast/multicast EIDs must be supported

- IPv6 RLOCs are not supported and don't need to be NATed

Protocol Exchange Sequence

- xTR sends Info-Request port **4342** to Map-Server

- Map-Server replies with list of RTRs in Info-Reply

- xTR sends Info-Request port **4341** to each RTR

- NATs install entries so RTRs can send from port **4341** to xTR

- Remote ITRs encapsulate to RTRs

- RTRs encapsulate to translated address and port

- xTRs Map-Register list of RTR RLOCs and its translated RLOC address

- xTRs have default map-cache entries to RTRs

- xTRs can short-cut map-cache entries to "non-NAT xTR" RLOCs

Mapping System Interface

EID3
xTR

(1)

EID2
xTR

NAT

(2)

a-rtr1   **a-rtr2**   **a-rtr3**

(3)

(1)   **(2)**

(4) Map-Reply
EID: EID1:
RLOC: a-rtr1, p254
RLOC: a-rtr2, p254
RLOC: a-rtr3, p254

Mapping System

Mapping:
EID: EID1:
RLOC: a-rtr1, p254
RLOC: a-rtr2, p254
RLOC: a-rtr3, p254
RLOC: tRLOC, p1

Map-Server rtr-list:
a-rtr1, a-rtr2, a-rtr3

(4) Map-Reply
EID: EID1:
RLOC: tRLOC, p1

(3) Map-Register
EID: EID1:
RLOC: a-rtr1, p254
RLOC: a-rtr2, p254
RLOC: a-rtr3, p254
RLOC: tRLOC, p1

(2) Info-Reply:
rtr-list: a-rtr1, a-rtr2, a-rtr3,
tRLOC and tPort

(3) Info-Requests
to 4341

NAT

xTR

EID1

(1) Info-Request to 4342

control
message

5

data
packet

Design Observations

9. Design Observations

    The following benefits and observations can be attributed to this
    design:

    o   An ITR behind a NAT virtually runs no control-plane and a very
        simple data-plane. All it does is RLOC-probe the RTRs in the
        common RLOC-set for each default map-cache entry. And maintains a
        very small map-cache of 4 entries per instance-ID it supports.

    o   An xTR behind a NAT can tell if another xTR is behind the same set
        of NAT devices and use Private RLOCs to reach each other on a
        short-cut path. It does this by comparing the Global RLOC
        returned from RTRs in Info-Reply messages.

    o   An xTR behind a NAT is free to send a Map-Request to the mapping
        system for any EID to test to see if there is a direct path to the
        LISP site versus potentially using a sub-optimal path through an
        RTR. This happens when xTRs exist that are not behind NAT devices
        where their RLOCs are global RLOCs.

    o   By sending Info-Requests to Map-Servers, an xTR behind a NAT can
        tell if they are reachable and if those Map-Servers also act as
        Map-Resolvers, the xTR behind the NAT can avoid sending Map-
        Requests to unreachable Map-Resolvers.

    o   Enhanced data-plane security can be used via LISP-Crypto
        mechanisms detailed in [RFC8061] using this NAT-Traversal design
        so both unicast and multicast traffic are encrypted.

    o   This design allows for the minimum amount of NAT device state
        since only RTRs are encapsulating to ETRs behind NAT devices.
        Therefore, the number of ITRs sending packets to EIDs behind NAT
        devices is aggregated out for scale. Scale is also achieved when
        xTRs behind NATs roam and RLOC-set changes need to update only RTR
        map-caches.

    o   The protocol procedures in this document can be used when a LISP
        site has multiple xTRs each connected via separate NAT devices to
        the public network. Each xTR registers their Global RLOCs and
        RTRs with merge semantics to the mapping system so remote ITRs can
        load-split traffic across a merged RTR set as well as RTRs across
        each xTR behind different NAT devices.

Draft Destiny

Ask if WG okay with Informational RFC?

Questions/Reactions/Tomatoes?