ristretto2558 decaf448 draft-irtf-cfrg-ristretto255-decaf448

what problem are we solving?

2.4 Notation Let \mathbb{G} denote a cyclic group of prime order p,



weierstrass e.g. secp256k1

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prime order





weierstrass e.g. secp256k1

prime order

fastest formulas



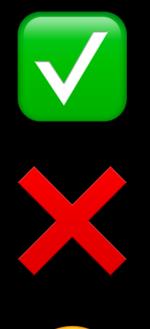


weierstrass e.g. secp256k1

prime order

fastest formulas

complete formulas







weierstrass e.g. secp256k1

prime order

fastest formulas

complete formulas

easy in constant time









okay, it's only a small conceptual mismatch...

...so what's wrong with a small cofactor?



WEII.

security analysis for the abstract protocol does not apply to its concrete implementation

We

security analysis for the abstract protocol does not apply to its concrete implementation subgroup validation is expensive, negating any speedup

We

to its concrete implementation

ad-hoc protocol tweaks are specific to each protocol

We

- security analysis for the abstract protocol does not apply
- subgroup validation is expensive, negating any speedup

tweaks cause subtle "features"

to agree on whether a signature is valid

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e.g., rfc8032 does not require ed25519 implementations

to agree on whether a signature is valid

- different behaviour between batch, single verification
- very bad for applications involving consensus
- also incompatible with hierarchical key derivation

tweaks cause subtle "features"

e.g., rfc8032 does not require ed25519 implementations

...or catastrophic failures

how do we fix this mismatch?

decaf & ristretto

construction, by mike hamburg, of a prime-order group

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uses a non-prime-order curve internally, no overhead

construction, by mike hamburg, of a prime-order group uses a non-prime-order curve internally, no overhead canonical, non-malleable encoding of group elements

uses a non-prime-order curve internally, no overhead "batteries included": a single hash-to-group method

- construction, by mike hamburg, of a prime-order group
- canonical, non-malleable encoding of group elements

how does this work?

three families of curves

montgomery curves

montgomery curves



$\mathcal{M}_{B,A}: \quad B_{u^2} = v(v^2 + Av + 1)$

montgomery curves $\mathcal{M}_{B,A}: \quad B_{u^2} = v(v^2 + A_v + 1)$

support fast "pseudomultiplication" $v(P) \mapsto v(CkJP)$

montgomery curves $\mathcal{M}_{B,A}: \quad Bu^2 = v(v^2 + Av + 1)$

require few constraints in circuits

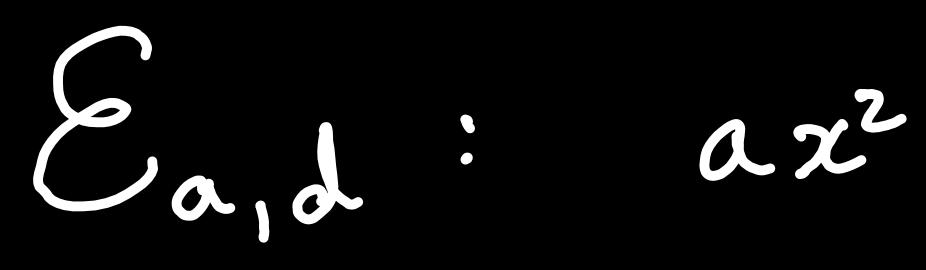
support fast "pseudomultiplication" $v(P) \mapsto v(CkJP)$







$\mathcal{C}_{a,d}: \quad ax^2 + y^2 = | + dx^2 y^2$

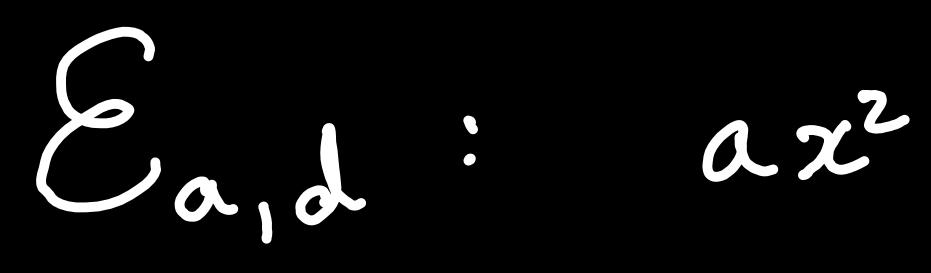


birationally equivalent to montgomery curves



$\mathcal{E}_{a,d}: \quad ax^2 + y^2 = \left| + dx^2 y^2 \right|$

edwards curves

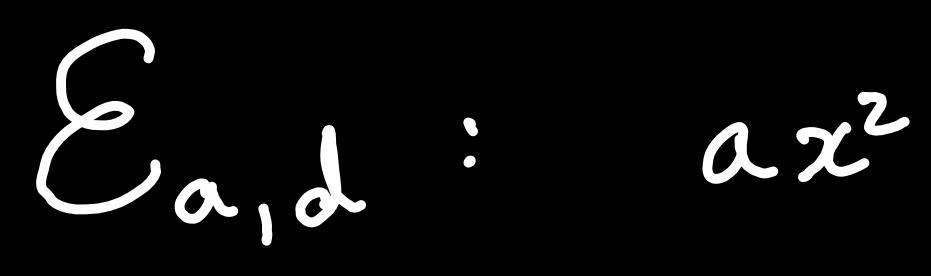


birationally equivalent to montgomery curves

fastest known formulas for curve operations

$\mathcal{E}_{a,d}: \quad ax^2 + y^2 = | + dx^2 y^2$

edwards curves



birationally equivalent to montgomery curves

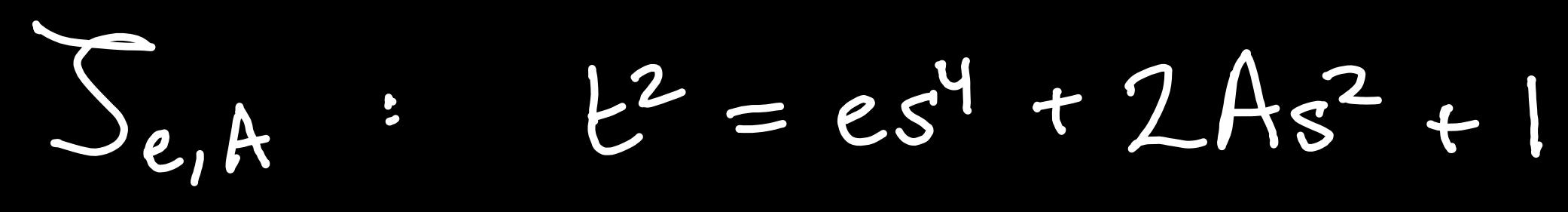
fastest known formulas for curve operations

formulas allow parallelism inside a curve operation

$\mathcal{E}_{a,d}: \quad ax^2 + y^2 = \left| + dx^2 y^2 \right|$

jacobi quartic curves

jacobi quartic curves



jacobi quartic curves $S_{e,A}: E^2 = es^4 + 2As^2 + 1$

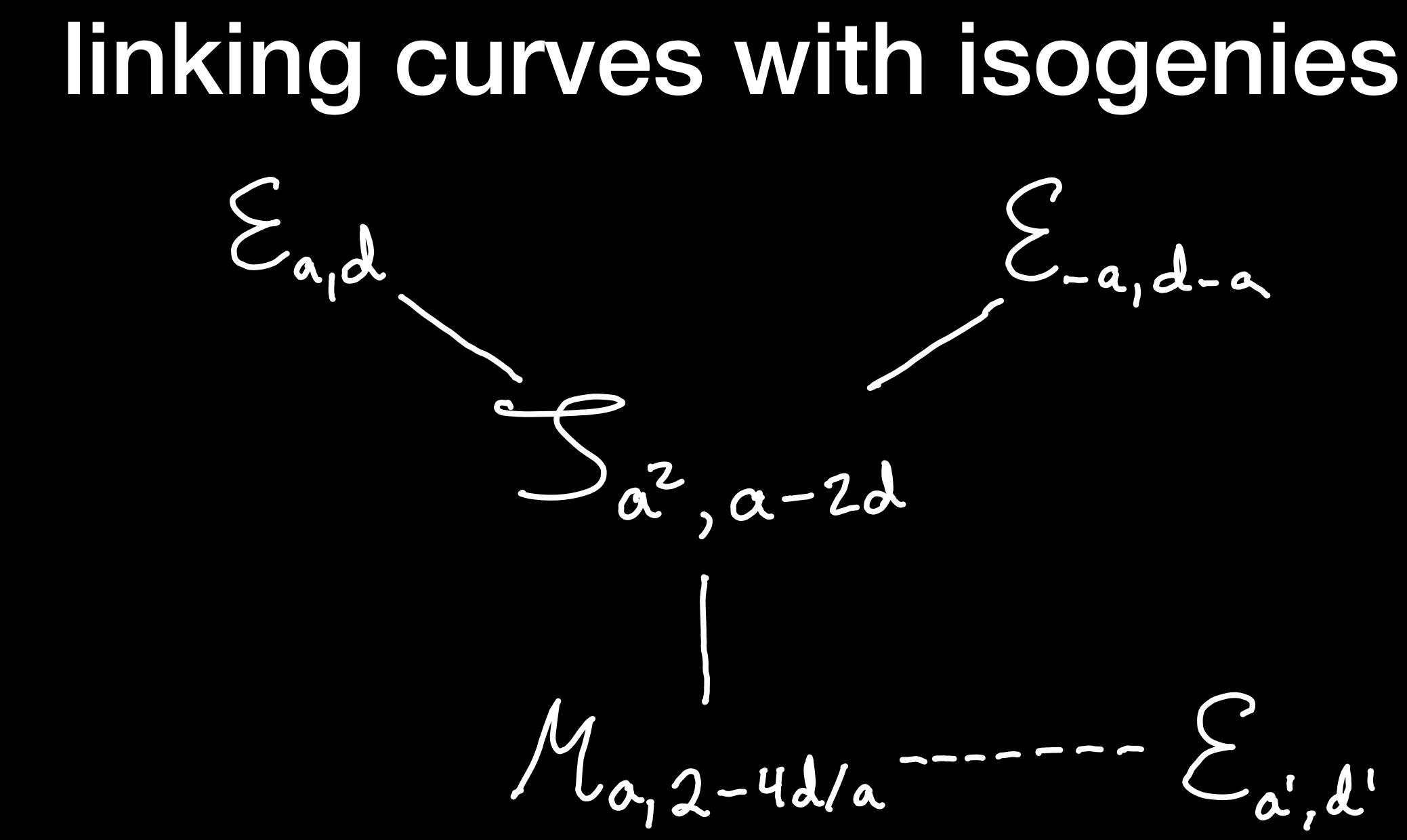
easy to write down 4 points of order 2

jacobi quartic curves $S_{e,A}: E^2 = es^4 + 2As^2 + 1$

easy to write down 4 points of order 2

we can efficiently encode $(s,t) \mod S[z]$

linking curves with isogenies



-a, d-a

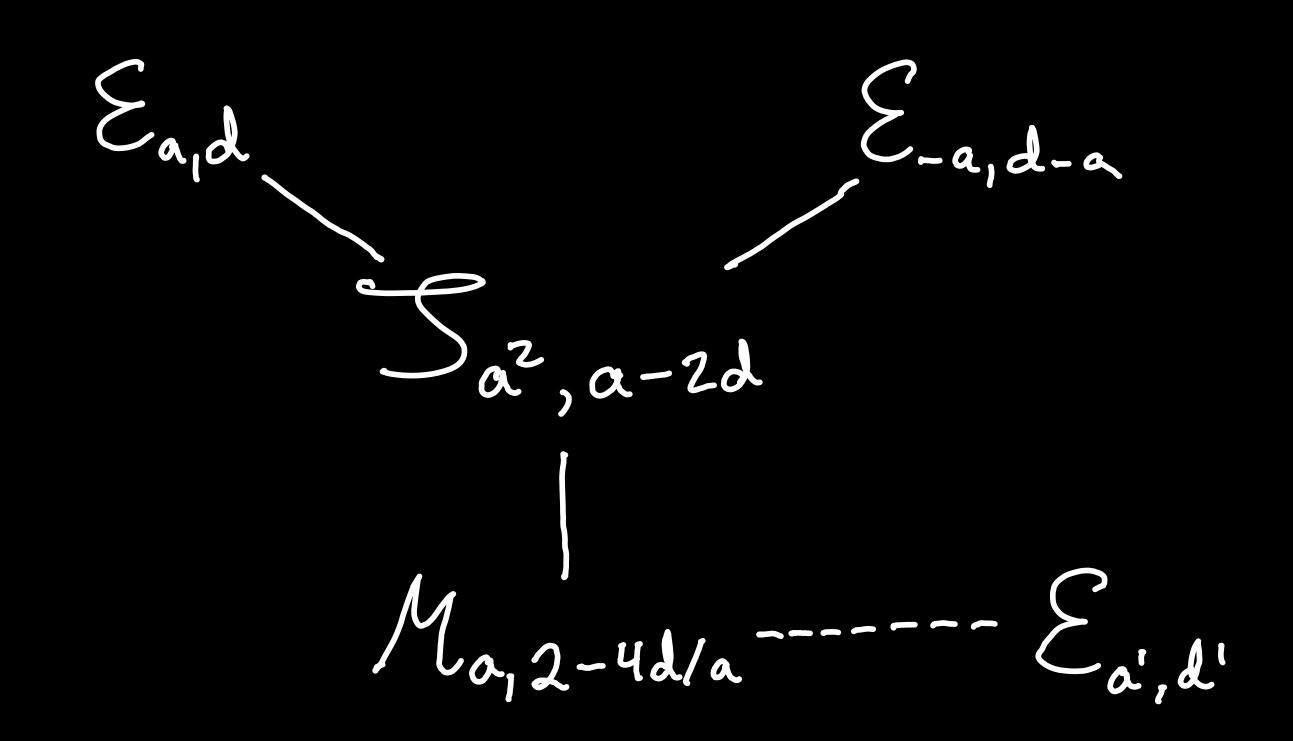
-a', d'

specify an encoding on the jacobi quartic

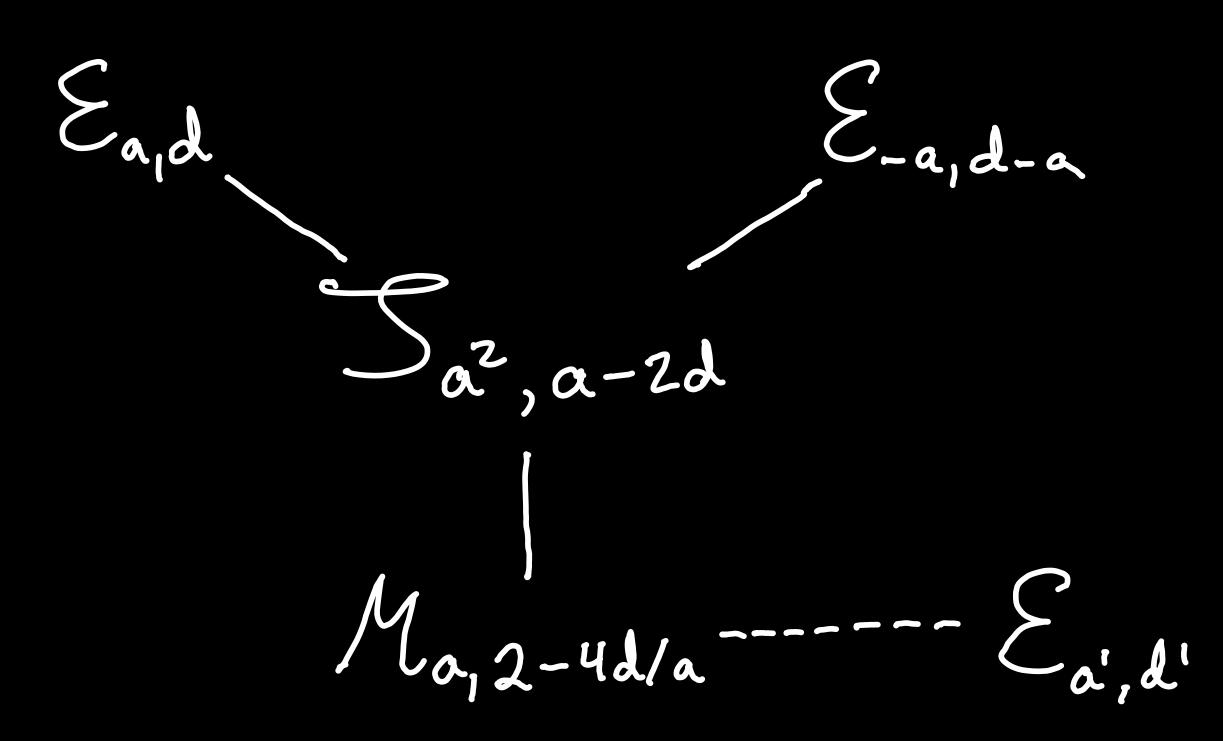
specify an encoding on the jacobi quartic isogenies "transport" the encoding to other curve shapes

specify an encoding on the jacobi quartic extra step to handle cofactor 8 instead of cofactor 4

- isogenies "transport" the encoding to other curve shapes

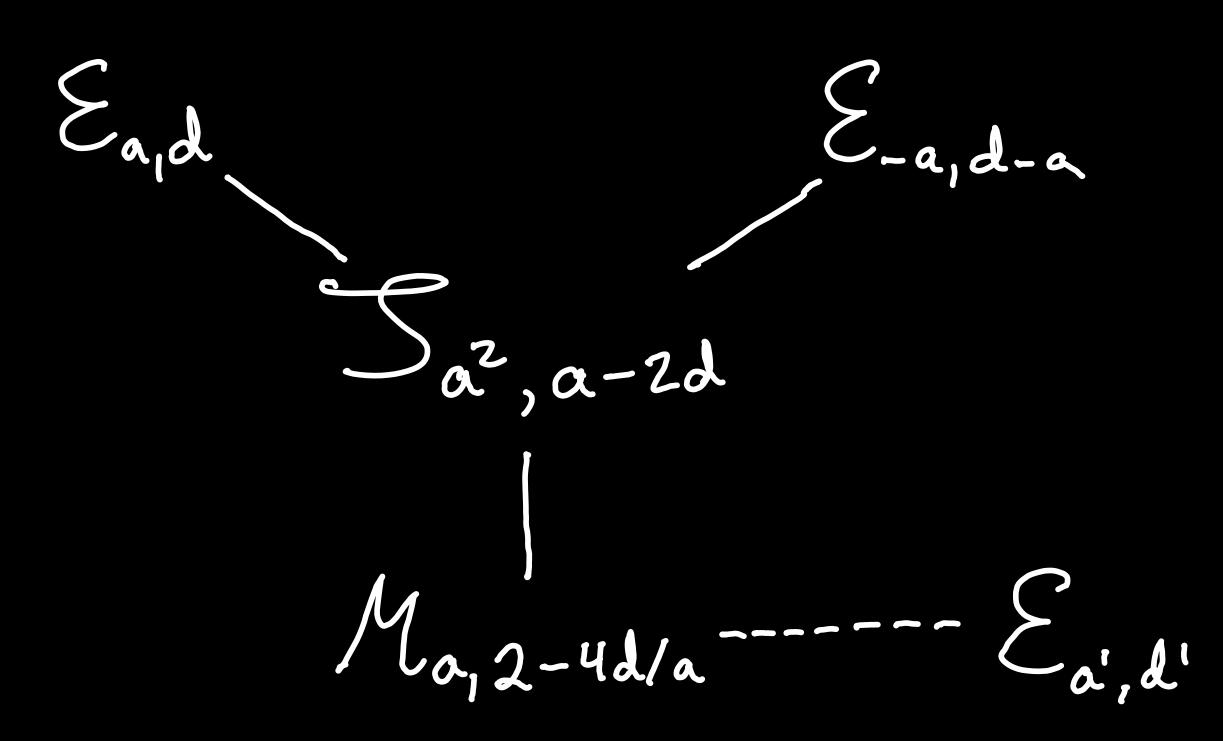


decaf transports the encoding to edwards directly



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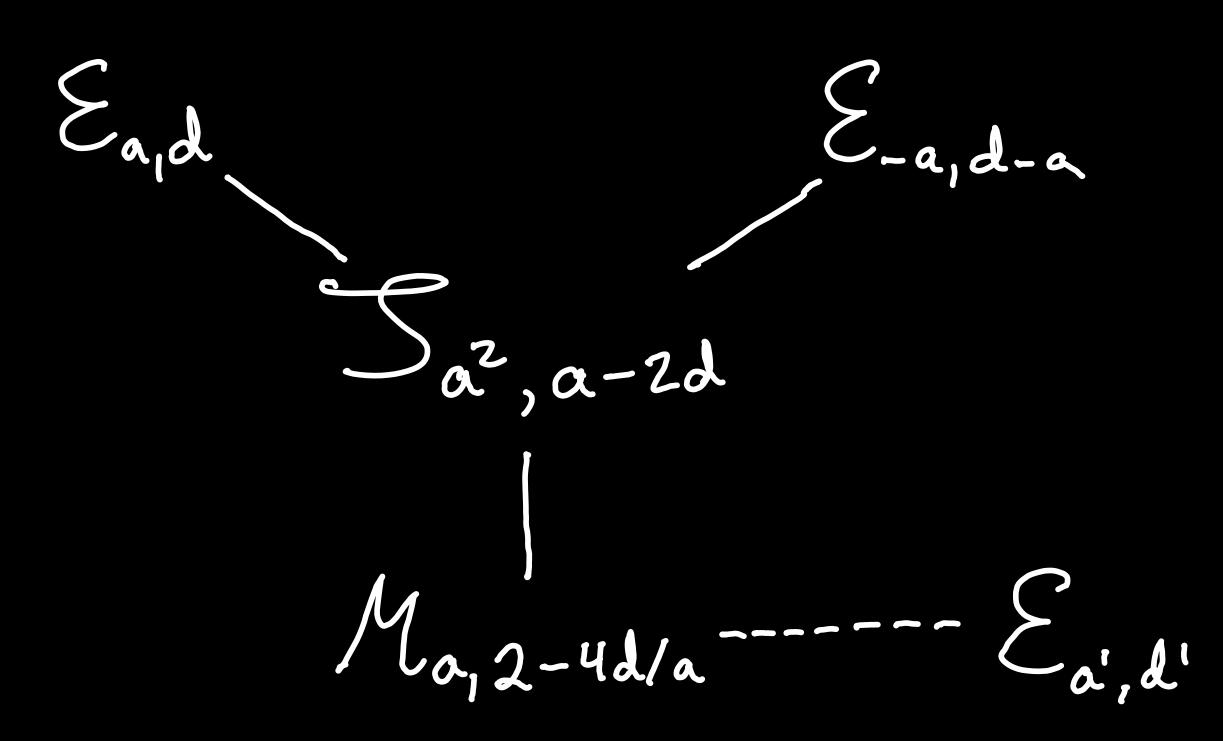
ristretto transports the encoding to edwards via montgomery + cofactor 8



decaf transports the encoding to edwards directly

ristretto transports the encoding to edwards via montgomery + cofactor 8

either can use any curve internally



concrete parameterizations



can be implemented using curve25519



can be implemented using curve25519 ~128-bit security level



can be implemented using curve25519 ~128-bit security level increasing adoption: zk proofs, psi, pake, etc



decaf448



can be implemented using edwards448

decaf448



can be implemented using edwards448 ~224-bit security level

decaf448

decaf448

can be implemented using edwards448 ~224-bit security level suitable where ed448 would be used

test vectors achieve "coverage" of all edge cases

test vectors achieve "coverage" of all edge cases language was rewritten last year to address feedback

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test vectors achieve "coverage" of all edge cases language was rewritten last year to address feedback added decaf448 parameters no outstanding issues, ready to go