Deploying LISP in a Campus Network

Jordi Paillisse, Marc Portoles, Albert Lopez, Alberto Rodriguez-Natal, David Iacobacci, Johnson Leong, Victor Moreno, Albert Cabellos, Fabio Maino, and Sanjay Hooda

LISP WG, IETF 109, Nov, 19th

Based on the paper:

SD-Access: Practical Experiences in Designing and Deploying Software Defined Enterprise Networks.
ACM CoNEXT '20, December, 2020,
https://arxiv.org/abs/2010.15236

Summary

1. Overview of Software-Defined Access (SDA)
2. LISP use-cases in SDA
3. Architecture and Design
4. Evaluation

Introduction: SDA

- SDN-based solution for Campus and Access Networks
- VXLAN data plane
- LISP control plane
- Unified Wired + Wireless
- Endpoint Mobility
- L3 Segmentation:
  - VXLAN VNI
  - Group-based polices
- L2 stretching

https://www.nordu.net/article/infoshare-campus-network-service

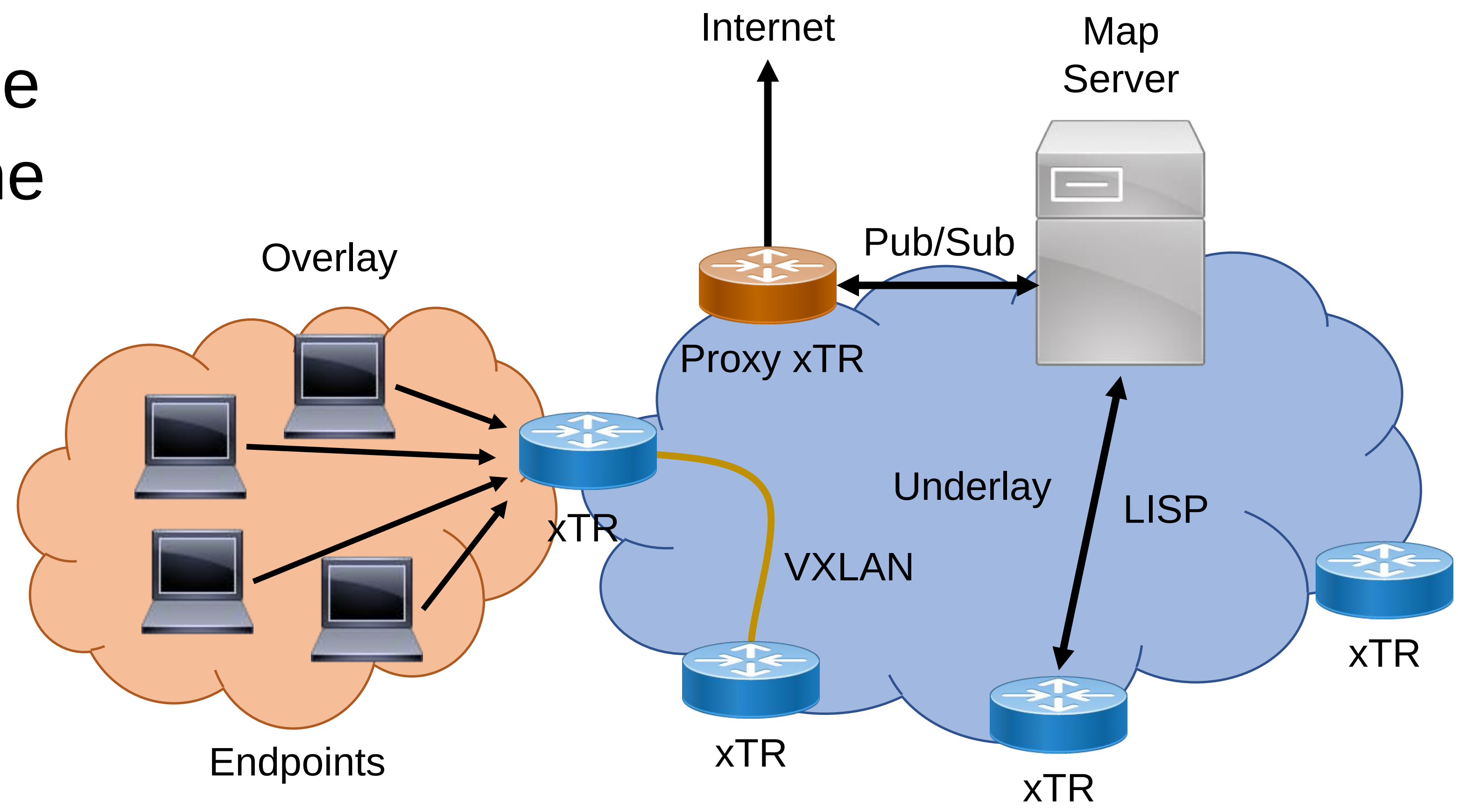
LISP use-cases in SDA

- L2/L3 EID Mobility
- Reduce and distribute data plane state
- Minimize CAPEX via providing routes on-demand
  - Less data plane entries → Smaller FIB → Less memory → Reduced cost
- Incremental deployment
  - Keep existing underlay, with standard OSPF or IS-IS

Architecture & Design

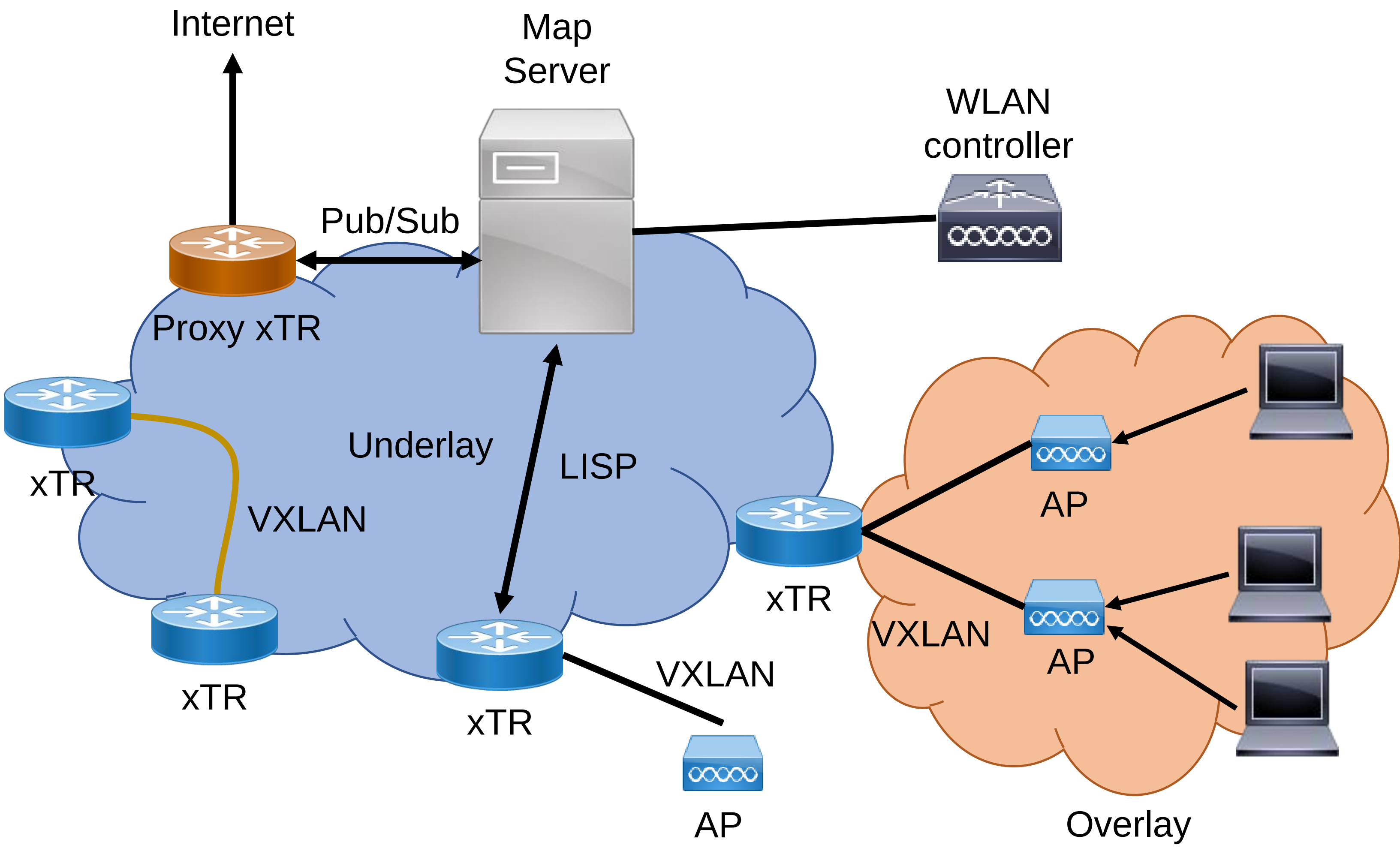
Architecture

- LISP control plane
- VXLAN data plane
- EIDs are individual hosts
- Three mappings per endpoint: IPv4, IPv6, and MAC to RLOC
- Plus MAC EID to IP EID



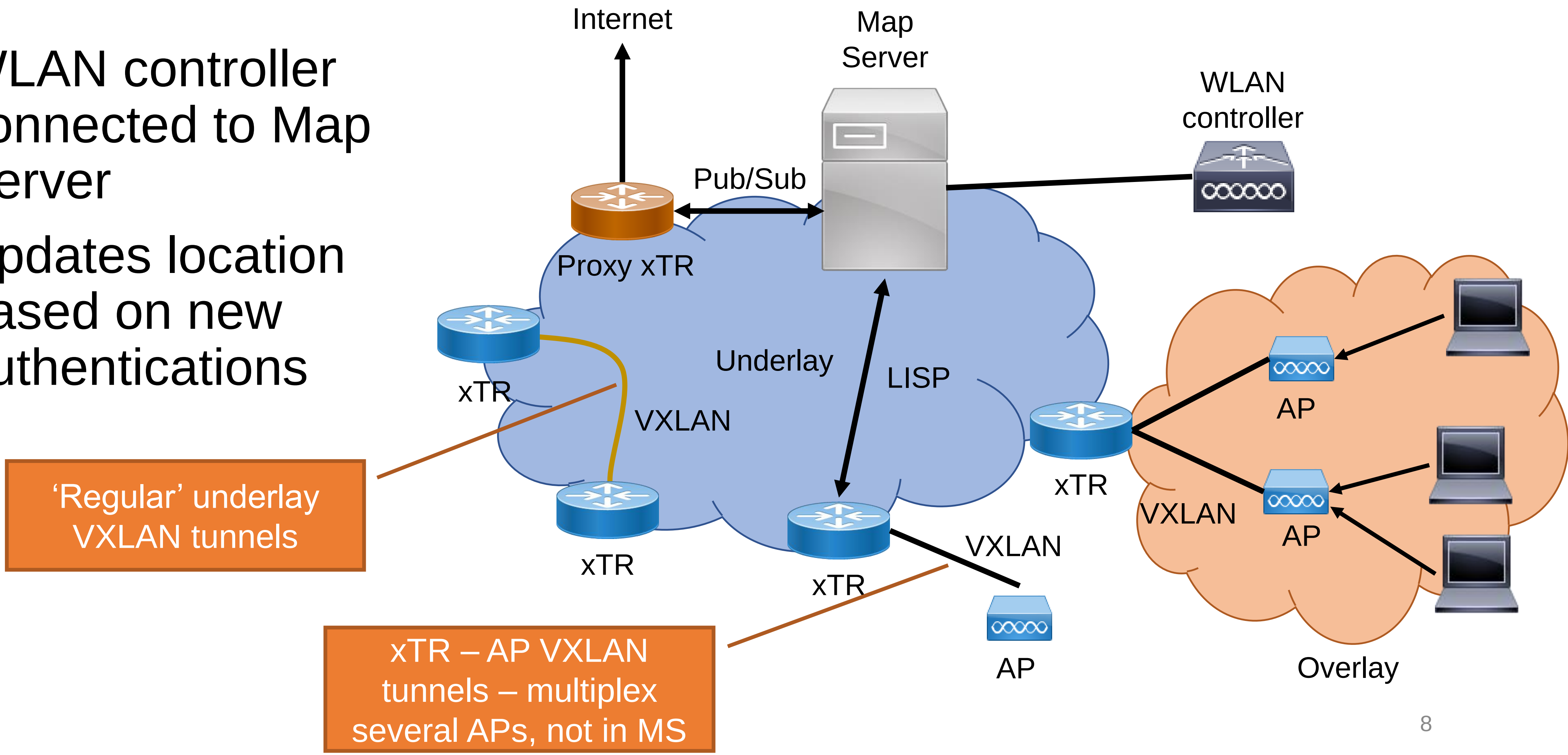
Architecture - Wireless

- WLAN controller connected to Map Server
- Updates location based on new authentications



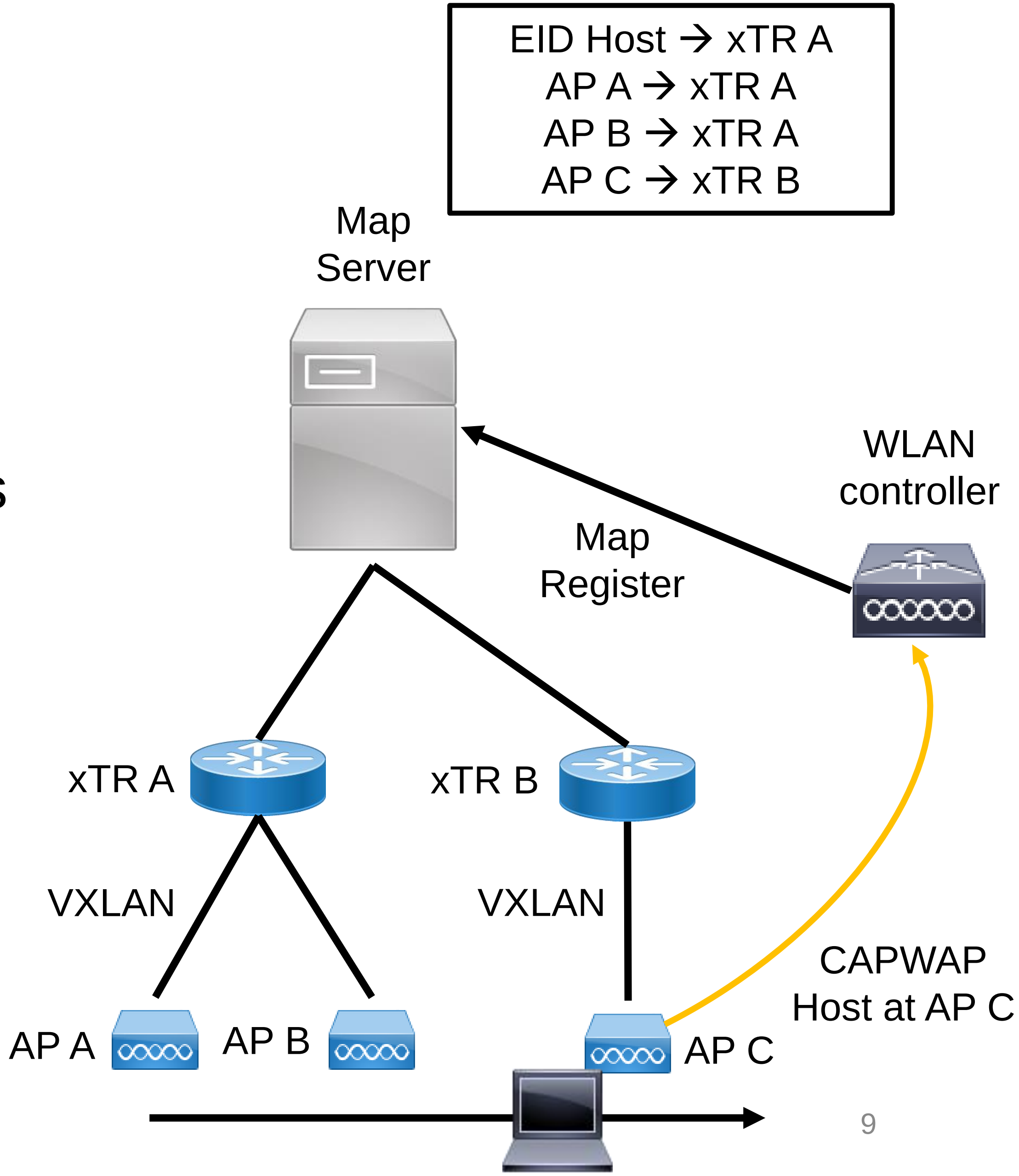
Architecture - Wireless

- WLAN controller connected to Map Server
- Updates location based on new authentications



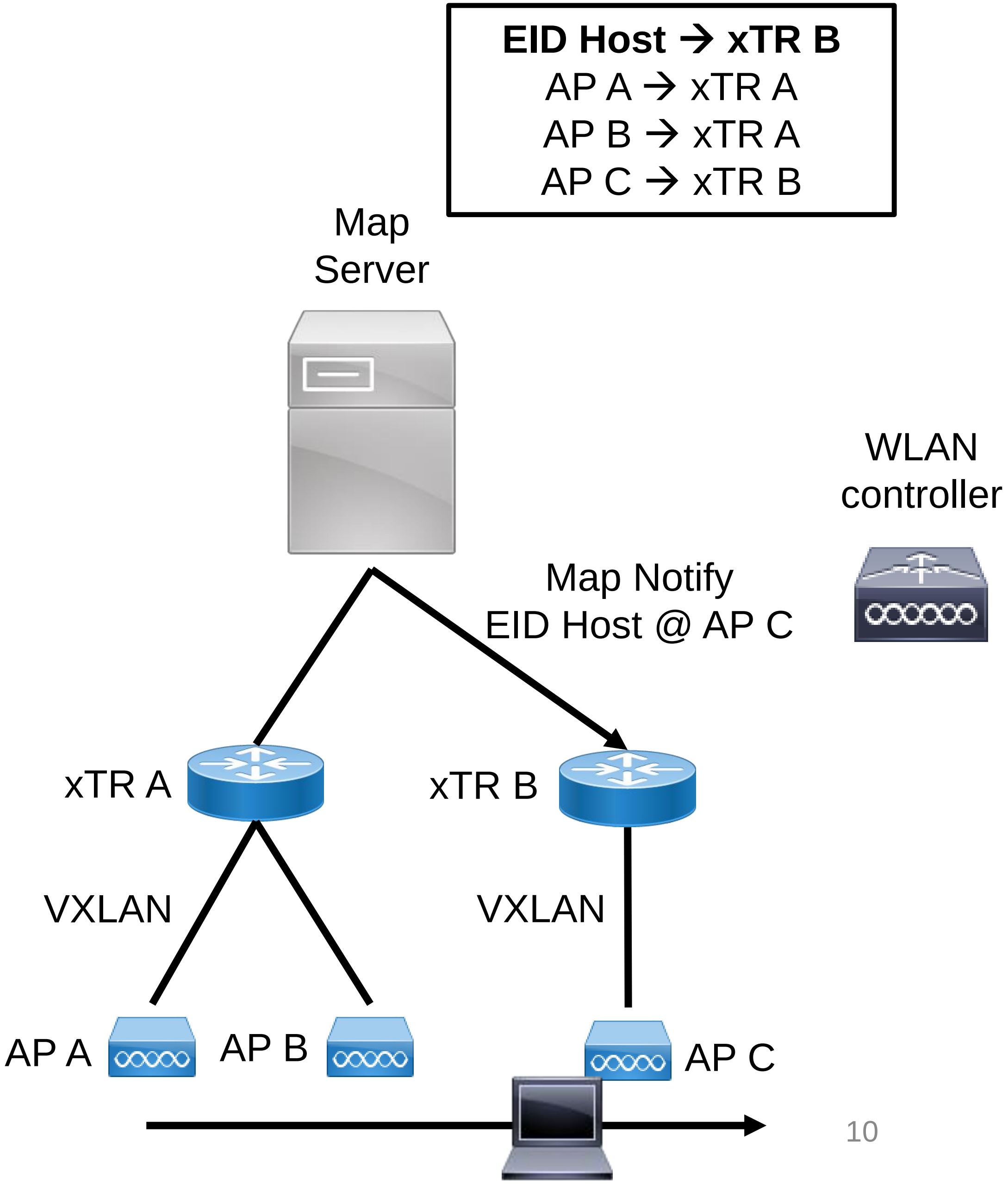
Design – Mobility (I)

- Static VXLAN tunnel between AP and xTR
  - Multiple APs per xTR
  - Store (AP IP → xTR RLOC) mappings
- WLAN controller detects host movement
  - Map Request: AP C? → xTR B
- Registers new location
  - EID Host → xTR B



Design – Mobility (II)

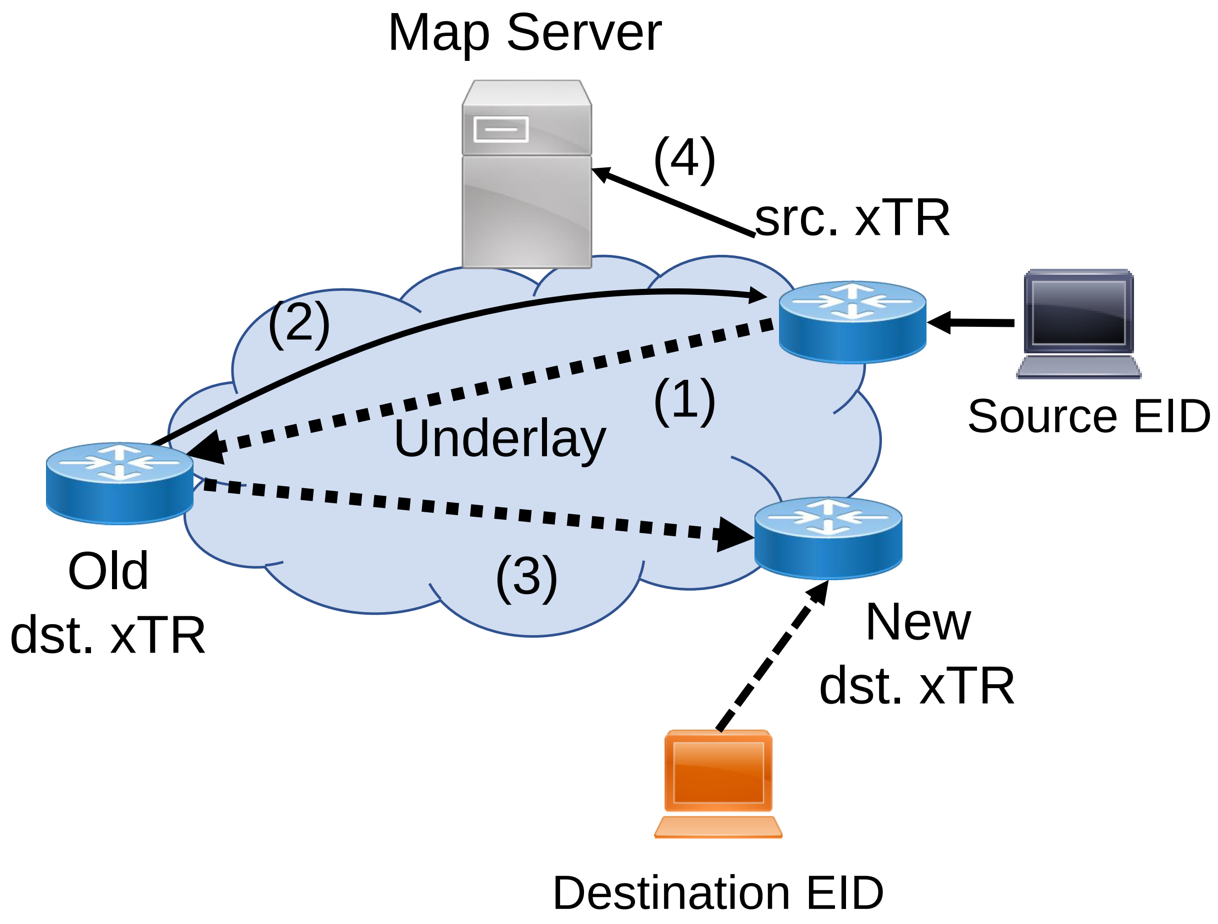
- xTR B receives Map Notify and updates local state → EID Host behind AP C
- Additional mechanisms to improve mobility:
  - Away entry (to not drop packets)
  - SMR (to notify outdated remote xTR)



Design – Mobility (and III)

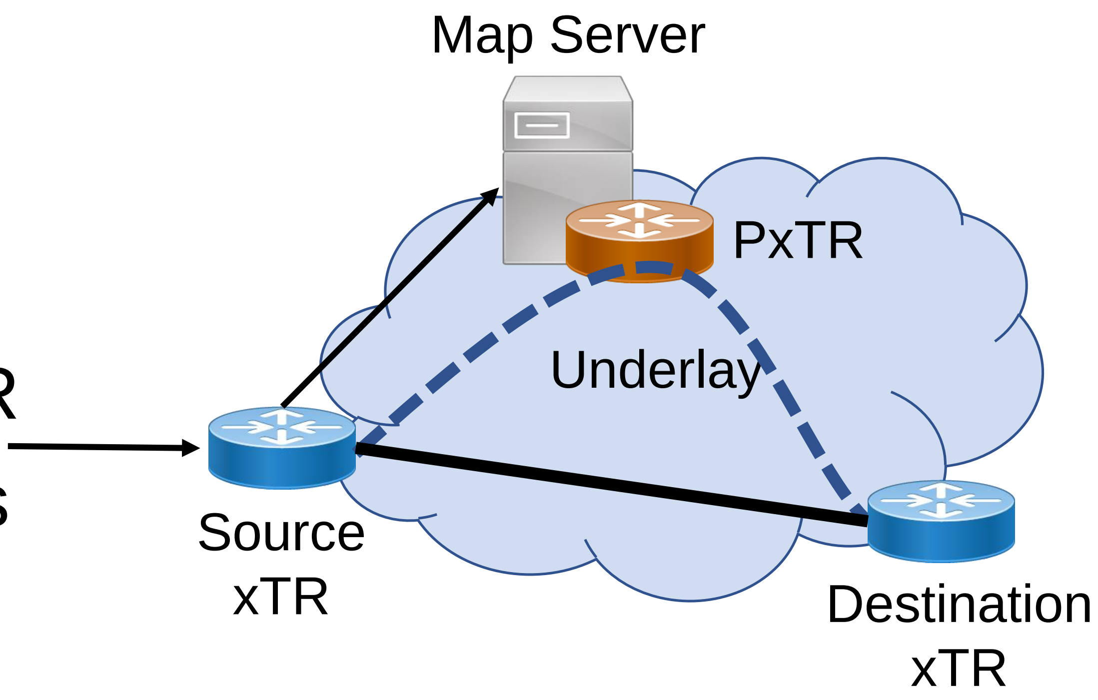
- SMR
  (1) If receive traffic for EID no longer in xTR
  (2) Send SMR to xTR
  (4) xTR updates map cache

- Away entry
  - Remember EIDs no longer in xTR
  - Forward traffic to new xTR (3)



Design – Avoid resolution delay (aka No First Packet Drop)

- Add default route in xTRs
- Points to PxTR (aka Border)
- Proxy has Pub/Sub functionality
- Forwards traffic on behalf of xTR
- Until more specific map cache is installed in the xTR



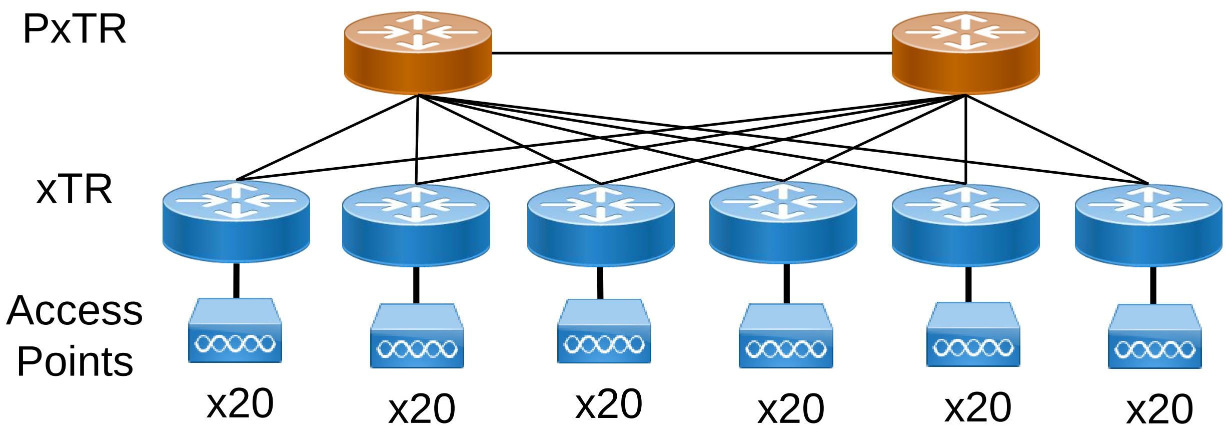
Design – Scalable L2 stretching

- Use cases:
  - Convert ARP broadcast to unicast
  - L2 protocols (eg. Apple Bonjour)
  - Legacy IoT devices (that do not use IP)
- Src. xTR encaps L2 frames to dst. xTR on VXLAN
- Resolve in Map Server missing info:
  - Use inner dst. MAC to locate dst. RLOC
  - For ARP: use MAC to EID mapping
- Forward ARP requests (instead of creating them) for coherence with IPv6 NDP

Evaluation

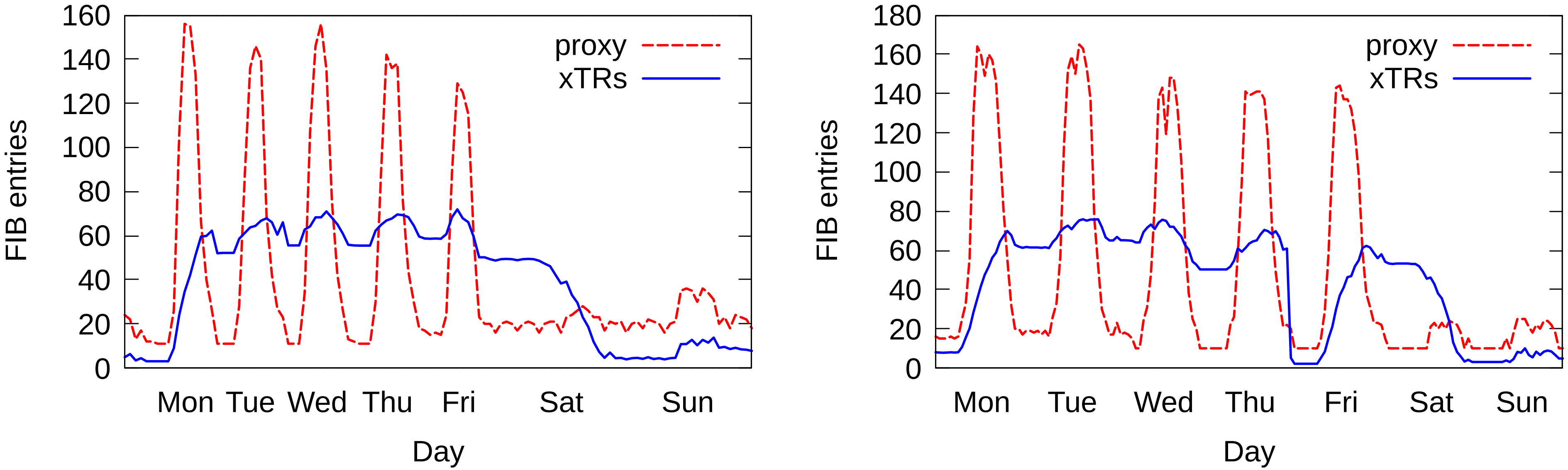
Data Plane State Reduction

- Count map cache entries in
  - PxTR
  - xTRs
- PxTR has all MS data due to Pub/Sub → fraction of mappings in xTRs
- Two different deployments:
  - Depl. A 150 hosts
  - Depl. B 450 hosts
  - Routers:
    - 1-2 PxTR
    - 7-6 xTR
    - 120 AP (20 per xTR)

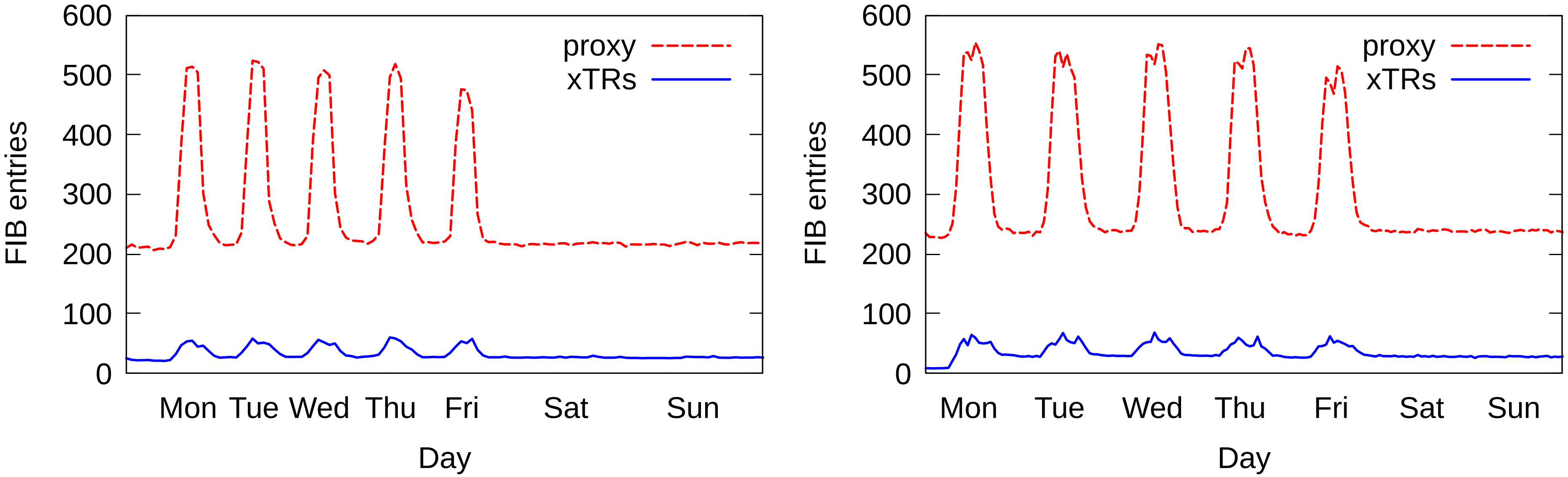


Data Plane State Reduction

A

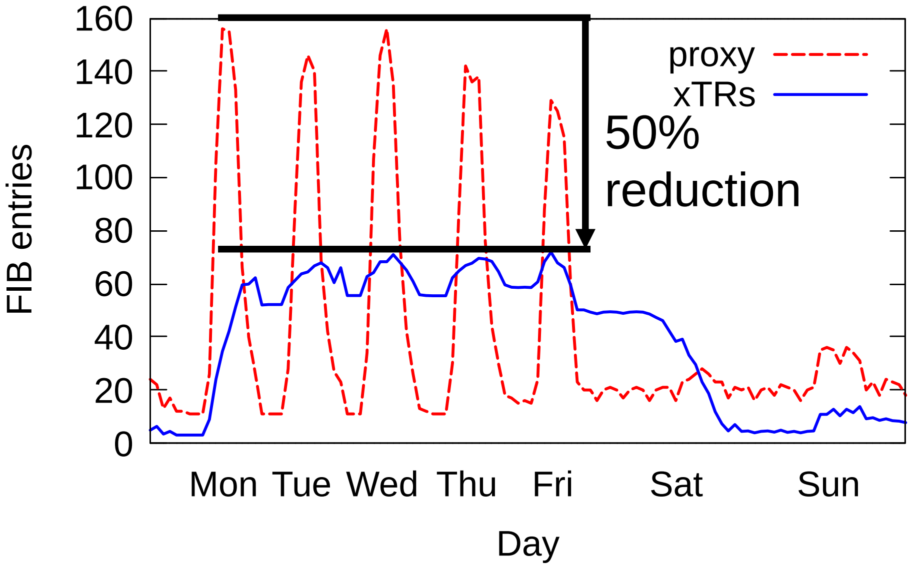


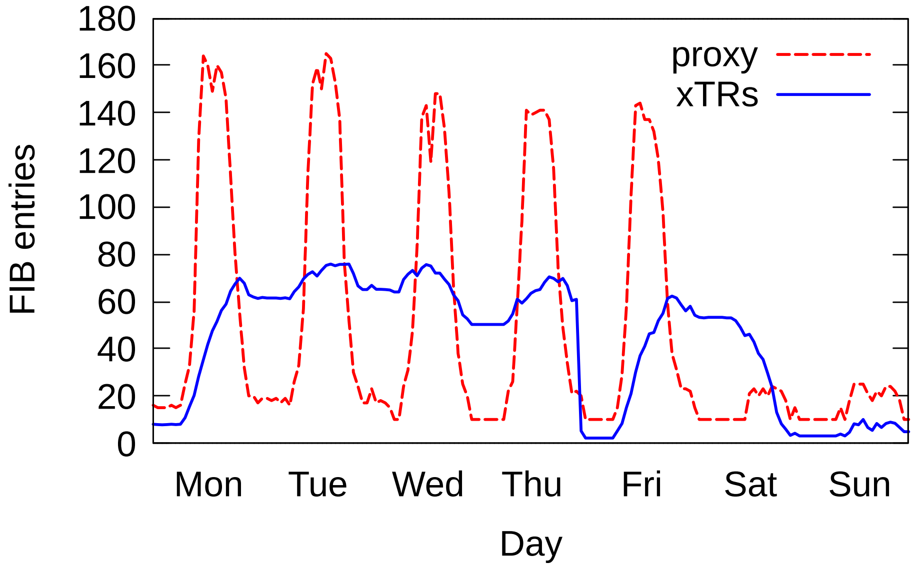
B



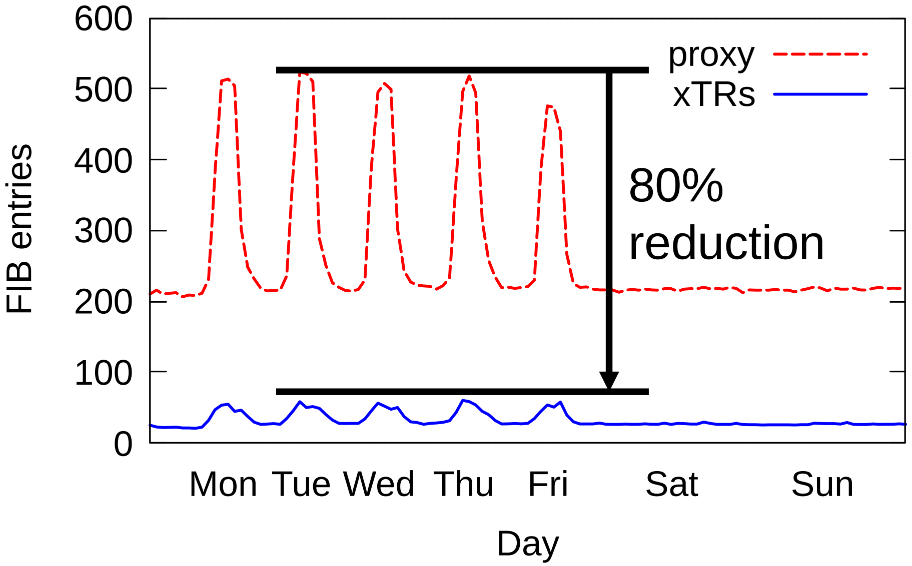
Data Plane State Reduction

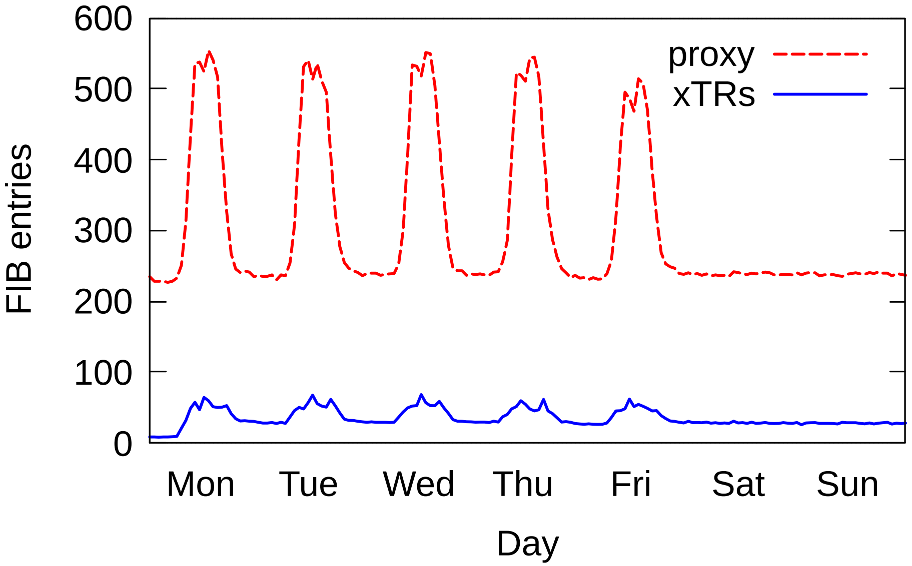
A



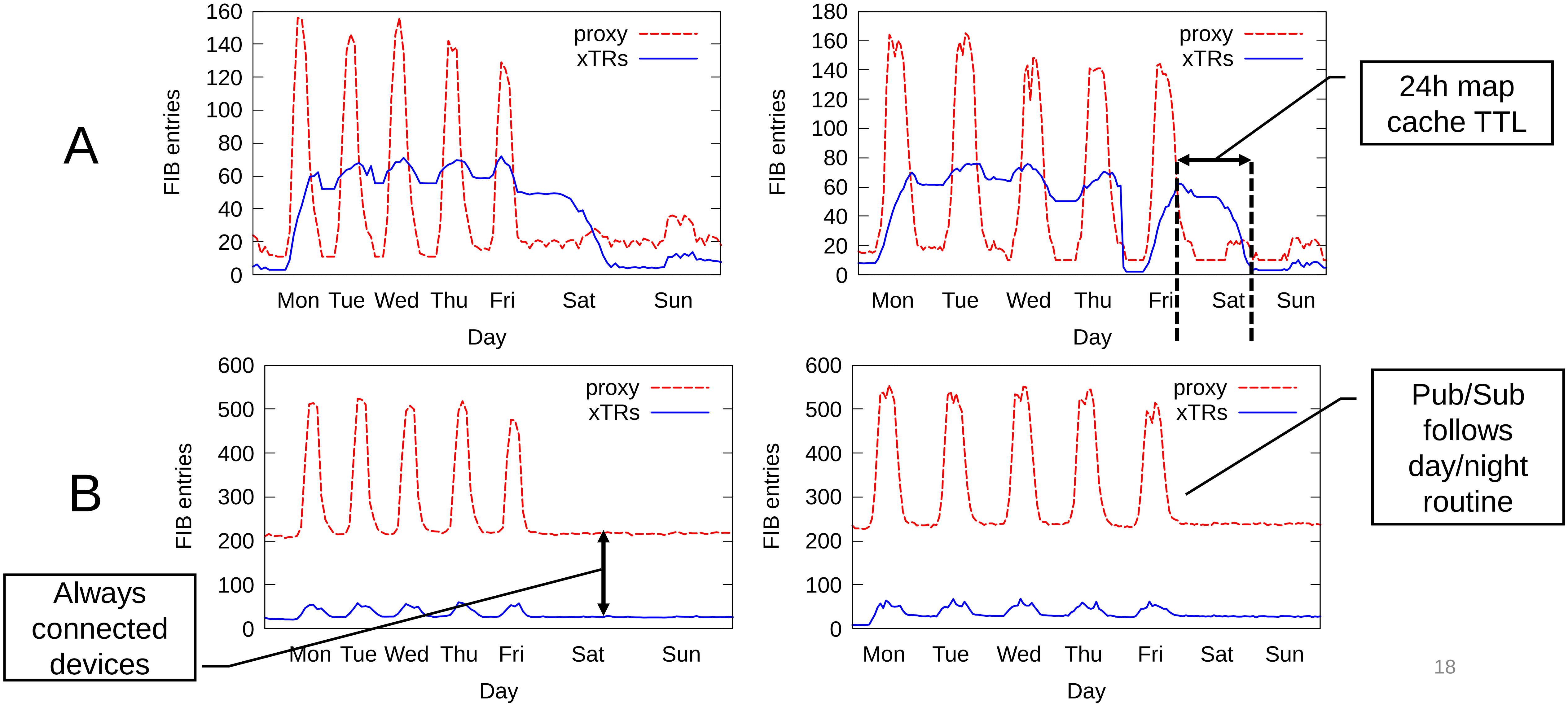


B



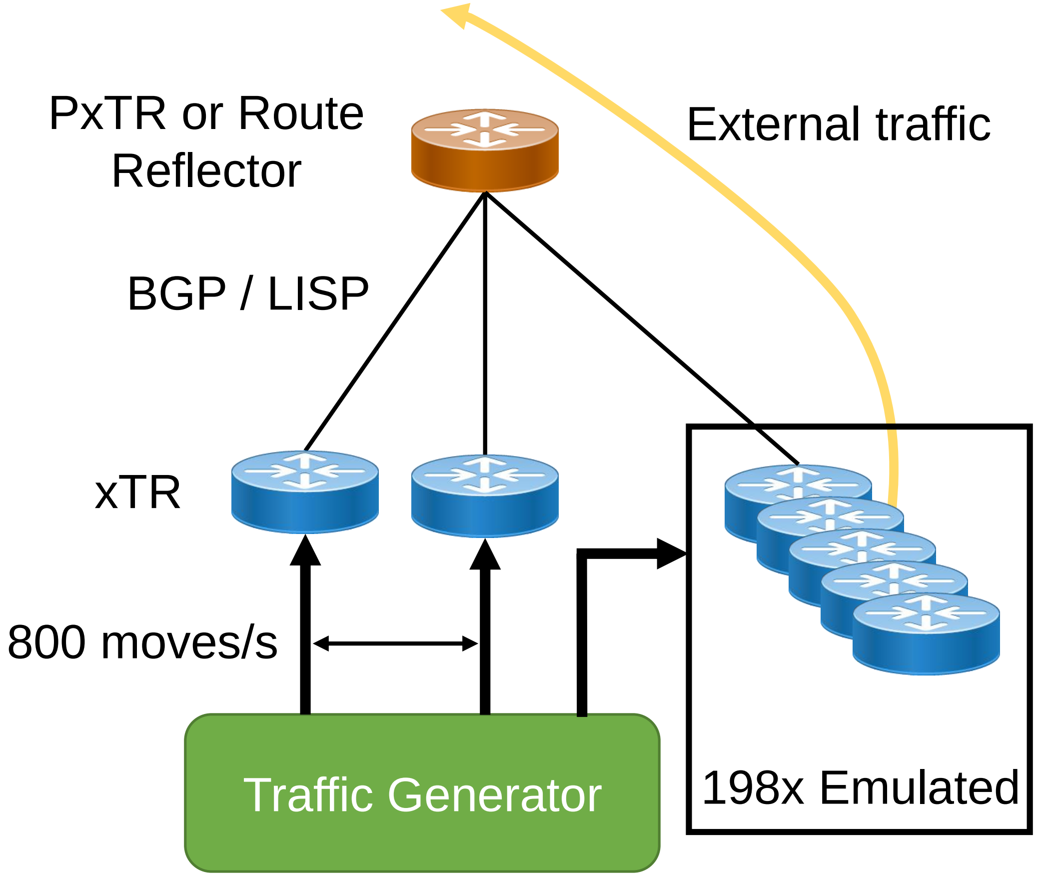


Data Plane State Reduction



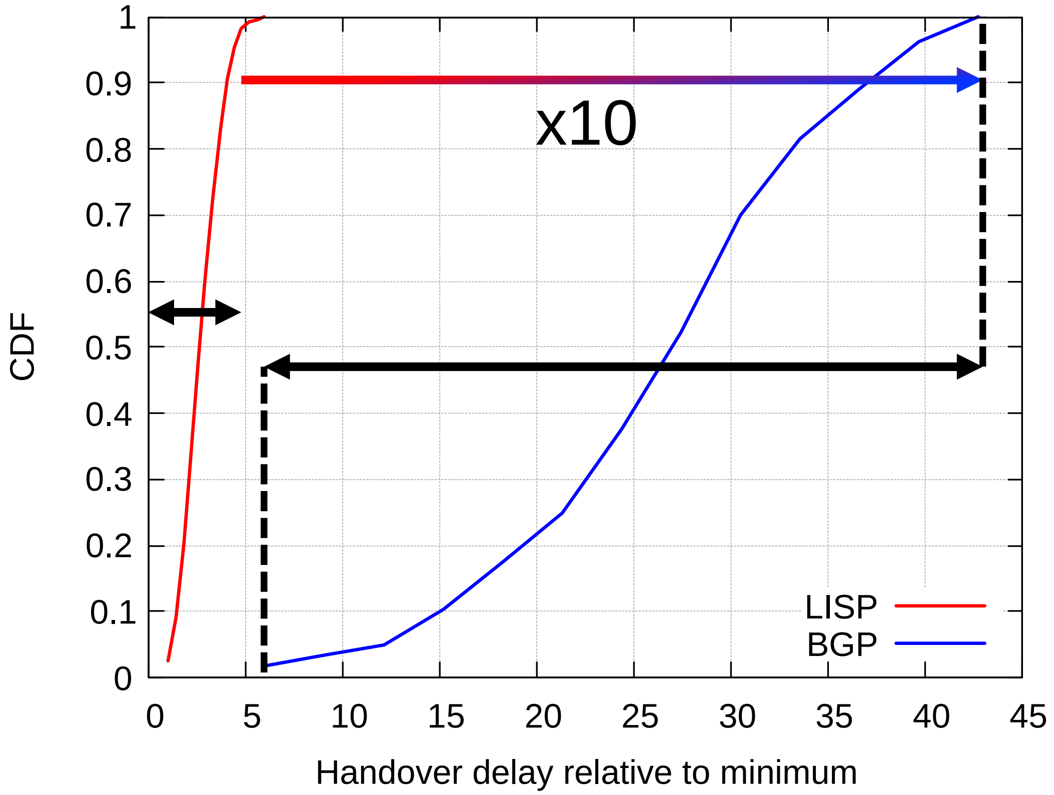
Handover Delay

- Massive mobility events
  - Eg. warehouse with mobile robots
- Calculate handover delay
- Lab setup
- 3 physical routers
- 198 emulated xTR
- Simulate handovers with a traffic generator



Handover Delay

- Compare LISP and BGP control planes
- Difference of aprox. one order of magnitude
- Notify only affected routers vs. all of them
- Less variability



Conclusions

- Example of a LISP deployment in an enterprise setting
- Reduced data plane state
- Distributed mobility data plane, with centralized control
- Versus classical WLAN controllers:
  - Improved routing (no triangular routing)
  - More scalable
- Reduce mobility handover

https://www.nordu.net/article/infoshare-campus-network-service

Thanks for listening!

You can find the paper at:

https://arxiv.org/abs/2010.15236