DSGLUE

DPRIVE, IETF 112, November 2021 Ben Schwartz Slides v01

Background: Authenticated ADoT

- "Authenticated" -> An active network adversary cannot ever gain access to the DNS query or response.
- Authenticated ADoT (A2DoT) is possible without any modification to parents ... if resolvers are very patient:
 - Always do NS revalidation before using a nameserver.
 - Also send a SVCB query for _dns.\$NSNAME (in parallel).
 - Use DNSSEC to authenticate the answers (requires signed child).
- This slows down resolution of all domains, not just those that use ADoT. Not likely to be deployed at scale (?)

Background: ADoT Parent Signals

- A signal in the parent domain is purely a performance optimization
- Most resolvers are impatient, so enabling optimized performance may be a prerequisite for wide deployment.

Background: Design Space

- Can we slow down resolution of existing domains?
- Do we care about A2DoT under non-A2DoT parents?
 - \circ ~ i.e. protecting label N+1 after label N has leaked
 - Can we require that non-A2DoT parents are signed?
- Can we add new RR types to the glue?
- Can we add new digest types to the DS record?
- Do we care about the latency of A2DoT-enabled domains?
- Can the child atomically update NS/DS/glue RRSets together in the parent?

Design assumptions for DSGLUE

- Can we slow down resolution of existing domains? NO
- Do we care about A2DoT under non-A2DoT parents? **YES**
 - \circ i.e. protecting label N+1 after label N has leaked
 - Can we require that non-A2DoT parents are signed? **YES**
- Can we add new RR types to the glue? NO
- Can we add new digest types to the DS record? **YES**
- Do we care about the latency of A2DoT-enabled domains?
 YES
- Can the child atomically update NS/DS/glue RRSets together in the parent? **NO**

DSGLUE Structure

• A DSGLUE record is a DS record with

- Algorithm: DSGLUE (TBD1)
- Digest type: VERBATIM (TBD2)
 - See draft-vandijk-dnsop-ds-digest-verbatim
- Contents: One arbitrary **RRSet** in a compact TLV encoding
 - Name must be below the zone cut, so only the prefix is encoded.
- Nonexistence is indicated by encoding an empty RRSet

Interpretation

- DSGLUE records are DS records.
 - Covered by the usual parent RRSIG, Bogus if tampered or removed

• RRSets in DSGLUE are glue.

- Only for delegation-following, not authoritative for the child zone.
- DSGLUE can repeat ordinary glue to secure it.
 - If they disagree, DSGLUE overrides the unsigned glue RRSet.
- Like glue, each DSGLUE record's RRSet SHOULD actually exist.
- Any RR Type is expressible, but not all are allowed (yet).
 - We can add RR Types as we figure out what they mean.

Simple Example

;; Child zone	
\$ORIGIN	child.example.
secret	AAAA 2001:db8::1
a	NS ns
ns	AAAA 2001:db8::2
@	DNSKEY
@	CDS <real ds=""></real>

;; Parent zone
\$ORIGIN example.
child NS ns.child
ns.child AAAA 2001:db8::2
child DS <real DS>

Both zones are fully signed. RRSIGs and TTLs are omitted for brevity.

Simple Example with slow A2DoT

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;; Parent zone
\$ORIGIN example.
child NS ns.child
ns.child AAAA 2001:db8::2
child DS <real DS>

Simple Example with DSGLUE

;; Child zone	
\$ORIGIN cl	nild.example.
secret A	AAA 2001:db8::1
@ NS	5 ns
ns A	AAA 2001:db8::2
_dns.ns SV	/CB 1 ns alpn=dot
@ DI	NSKEY
@ CI	DS <real ds=""></real>
CI	DS \$DSGLUE(., NS,
<mark>[ns.child</mark>	.example.])
CI	DS \$DSGLUE(_dns.ns.,
ا SVCB, [1	ns.child.example.
<mark>alpn=dot])</mark>	

;; Parent zone
\$ORIGIN example.
child NS ns.child
ns.child AAAA 2001:db8::2
child DS <real DS>
 DS \$DSGLUE(., NS,
[ns.child.example.])
 DS \$DSGLUE(_dns.ns.,
SVCB, [1 ns.child.example.
alpn=dot])

Closing thoughts

- DSGLUE shows that A2DoT is achievable even under very challenging assumptions.
 - No slowdown for non-participating zones
 - Minimum additional latency for participating zones
 - The child zone doesn't have to be signed.
- This is a very large design space with a lot of options to consider.
- We are not blocked. We can start testing slow A2DoT while we figure out what we want from parent signals.