BGP MVPN in IPv6 Infrastructure Networks: Problems and Solution Approaches

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Abstract

MVPN deployment faces some problems while used in provider’s IPv6 infrastructure networks. This document describes these problems, and the solutions to solve these problems.

Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

Status of This Memo

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1. Introduction

BGP MVPN procedure is defined in [RFC6514]. As a mature MVPN technology, it has been accepted by most operators and vendors. In [RFC6515], BGP MVPN is updated for IPv6 infrastructure networks. However, the deployment of BGP MVPN in IPv6 network still faces some problems. This document describes these problems, and the solutions to solve these problems.

2. Terminology

Readers of this document are assumed to be familiar with the terminology and concepts of the documents listed as Normative References.

3. Problems and Solutions

3.1. Problems

In [RFC6514], the following processes are not suitable if IPv6 infrastructure addresses are used in provider’s networks:
1. In paragraph 7 of section 11.1.3, it described that "To support non-segmented inter-AS tunnels, the Source AS field in the C-multicast route is set to value of the Originating Router’s IP Address field of the found Intra-AS I-PMSI A-D route". In NLRI of C-multicast route, Source AS field is 4 bytes length, while the Originating Router’s IP Address field of Intra-AS I-PMSI A-D route is 16 bytes length in provider’s IPv6 networks.

2. In paragraph 2 of section 11.2, it described that "To support non-segmented inter-AS tunnels, instead of matching the RD and Source AS carried in the C-multicast route against the RD and Source AS of an Inter-AS I-PMSI A-D route, the ASBR should match it against the RD and the Originating Router’s IP Address of the Intra-AS I-PMSI A-D routes". Source AS field in NLRI of C-multicast route cannot be translated to the Originating Router’s IP Address of the Intra-AS I-PMSI A-D routes in provider’s IPv6 networks, because of the mismatch of their field length.

In [RFC7716], zero RD is introduced in BGP MVPN NLRIs to enable Global Table Multicast service in provider’s networks. In IPv6 infrastructure networks, Leaf PEs cannot send two distinct C-multicast route to two individual upstream root PEs for selective forwarding, because the RD of the two roots is the same.

In the process of evolution to IPv6, IPv4 and IPv6 infrastructure addresses will coexist in the provider’s network. The following figure is an example of BGP MVPN evolution to IPv6.

![Figure 1: BGP MVPN Evolution to IPv6 Infrastructure](image)

During the evolution process, IPv4 and IPv6 parallel BGP sessions are established between Provider Edge routers and Reflector routers, if the BGP MVPN route send to all IPv4 and IPv6 BGP peer without any
control, the number of the PATHs of these routes will be doubled with each reflection while BGP ADD-PATH [RFC7911] procedure is enabled on Reflector routers.

3.2. Modification of C-Multicast route NLRI

To support non-segmented inter-AS tunnels in IPv6 infrastructure network, the C-Multicast route NLRI is redefined as following:

```
+-----------------------------------+
|      RD   (8 octets)              |
+-----------------------------------+
|   Root Distinguisher (4 octets)   |
+-----------------------------------+
| Multicast Source Length (1 octet) |
+-----------------------------------+
|   Multicast Source (variable)     |
+-----------------------------------+
| Multicast Group Length (1 octet)  |
+-----------------------------------+
|  Multicast Group   (variable)     |
+-----------------------------------+
```

In the above figure, the Root Distinguisher field replaces the Source As field defined in [RFC6514]. When constructing a C-Multicast route, leaf PE follows the following specification:

1. If the Originating Router’s IP Address field of the found Intra-AS I-PMSI A-D route is an IPv4 address, the Root Distinguisher field MUST be treated as Source AS field and section 11.1.3 of [RFC6514] MUST be followed.

2. If the Originating Router’s IP Address field of the found Intra-AS I-PMSI A-D route is an IPv6 address and the root PE and leaf PE are in the same AS, the Root Distinguisher field MUST be treated as Source AS field and section 11.1.3 of [RFC6514] MUST be followed.

3. If the Originating Router’s IP Address field of the found Intra-AS I-PMSI A-D route is an IPv6 address and the root PE and leaf PE are in the different ASs, a four bytes distinct value MUST be assigned by leaf PE for each root PE, the Root Distinguisher field in C-Multicast NLRI is filled with this value and a distinct C-multicast route will be send to individual upstream root PE.
When receiving a C-Multicast route from E-BGP neighbors, the ASBR checks whether an IPv6 VRF Route Import Extended Community is included in this route and takes following actions:

1. If the IPv6 VRF Route Import Extended Community does not exist in the C-Multicast route, the ASBR treats the Root Distinguisher field as Source AS field and follows the description in section 11.2 of [RFC6514].

2. If the IPv6 VRF Route Import Extended Community does exist in the C-Multicast route, the ASBR match the IPv6 address carried in this extended community and the RD in C-Multicast route NLRI against the Originating Router’s IP Address and the RD of the Intra-AS I-PMSI A-D routes. If the corresponding Intra-AS I-PMSI A-D route exists, the ASBR propagates the C-Multicast route in its local AS.

3.3. Route reflection control

To reduce BGP MVPN routes in Parallel IPv4 and IPv6 BGP sessions scenario, the following actions should be taken by sender PEs:

1. For Intra-AS I-PMSI A-D Route, S-PMSI A-D Route and Leaf A-D Route, if the Originating Router’s IP Address field in the route is filled with an IPv6 address, it is sent to the IPv6 BGP neighbors; otherwise, it is sent to the IPv4 BGP neighbors.

2. For Inter-AS I-PMSI A-D Route and Source Active A-D Route, it is sent to both IPv6 BGP neighbors and IPv4 BGP neighbors.

3. For C-Multicast Route, if the IPv6 VRF Route Import Extended Community exists in the route, it is sent to the IPv6 BGP neighbors; otherwise, it is sent to the IPv4 BGP neighbors.

In the reflector routers, the part of routes which are received from IPv6 BGP neighbors will be reflected to other IPv6 BGP neighbors and the other part of routes which are received from IPv4 BGP neighbors will be reflected to other IPv4 BGP neighbors.

4. Security Considerations

This document introduces no new security considerations beyond those already specified in [RFC6514] and [RFC6515].
5. IANA Considerations

This document contains no actions for IANA.

6. Acknowledgements

Your name here

7. Normative References


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IPv6-Only PE Design for IPv4-NLRI with IPv6-NH
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Abstract

As Enterprises and Service Providers upgrade their brown field or green field MPLS/SR core to an IPv6 transport, Multiprotocol BGP (MP-BGP) now plays an important role in the transition of their Provider (P) core network as well as Provider Edge (PE) Edge network from IPv4 to IPv6. Operators must be able to continue to support IPv4 customers when both the Core and Edge networks are IPv6-Only.

This document details an important External BGP (eBGP) PE-CE Edge and Inter-AS IPv6-Only peering design that leverages the MP-BGP capability exchange by using IPv6 peering as pure transport, allowing both IPv4 Network Layer Reachability Information (NLRI) and IPv6 Network Layer Reachability Information (NLRI) to be carried over the same (Border Gateway Protocol) BGP TCP session. The design change provides the same Dual Stacking functionality that exists today with separate IPv4 and IPv6 BGP sessions as we have today. With this design change from a control plane perspective a single IPv6 is required for both IPv4 and IPv6 routing updates and from a data plane forwarding perspective an IPv6 address need only be configured on the PE and CE interface for both IPv4 and IPv6 packet forwarding.

This document provides a much needed solution for Internet Exchange Point (IXP) that are facing IPv4 address depletion at large peering points. With this design, IXP can now deploy PE-CE IPv6-Only eBGP Edge or Inter-AS peering design to eliminate IPv4 provisioning at the Edge. This core and edge IPv6-Only peering design paradigm change can apply to any eBGP peering, public internet or private, which can
be either Core networks, Data Center networks, Access networks or can be any eBGP peering scenario. This document provides vendor specific test cases for the IPv6-Only peering design as well as test results for the five major vendors stakeholders in the routing and switching industry, Cisco, Juniper, Arista, Nokia and Huawei. With the test results provided for the IPv6-Only Edge peering design, the goal is that all other vendors around the world that have not been tested will begin to adopt and implement this new Best Current Practice for eBGP IPv6-Only Edge peering.

As this issue with IXP IPv4 address depletion is a critical issue around the world, it is imperative for an immediate solution that can be implemented quickly. This Best Current Practice IPv6-only eBGP peering design specification will help proliferate IPv6-Only deployments at the eBGP Edge network peering points to starting immediately at a minimum with operators around the world using Cisco, Juniper, Arista, Nokia and Huawei. As other vendors start to implement this Best Current Practice, the IXP IPv4 address depletion gap will eventually be eliminated.

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1. Introduction

As Enterprises and Service Providers upgrade their brown field or green field MPLS/SR core to an IPv6 transport such as MPLS LDPv6, SR-MPLSv6 or SRv6, Multiprotocol BGP (MP-BGP) now plays an important role in the transition of the Provider (P) core networks and Provider Edge (PE) edge networks from IPv4 to IPv6. Operators have a requirement to support IPv4 customers and must be able to support IPv4 address family and Sub-Address-Family Virtual Private Network (VPN)-IPv4, and Multicast VPN IPv4 customers.

IXP are also facing IPv4 address depletion at their peering points, which are large Layer 2 transit backbones that service providers peer and exchange IPv4 and IPv6 Network Layer Reachability Information (NLRI). Today, these transit exchange points are Dual Stacked. With this IPv6-only BGP peering design, only IPv6 is configured on the PE-CE interface, the Provider Edge (PE) - Customer Edge (CE), or Inter-AS ASBR (Autonomous System Boundary Router) to ASBR (Autonomous System Boundary Router) PE-PE Provider Edge (PE) - Provider Edge (PE), the IPv6 BGP peer is now used to carry IPv4 (Network Layer Reachability Information) NLRI over an IPv6 next hop using IPv6 next hop encoding defined in [RFC8950], while continuing to forward both IPv4 and IPv6 packets. In the framework of this design the PE is no longer Dual Stacked. However in the case of the CE, PE-CE link CE side of the link is no longer Dual Stacked, however all other internal links within the CE domain may or may not be Dual stacked. In the Inter-AS case the ASBR-ASBR PE-PE peering all peerings would be now IPv6-Only for all Inter-AS options Peering Option-A, Option-B, Option-AB and Option-C per [RFC4364]. We now refer to this PE as an "IPv6-Only PE" using the IPv6-Only PE Design framework.

MP-BGP specifies that the set of usable next-hop address families is determined by the Address Family Identifier (AFI) and the Subsequent Address Family Identifier (SAFI). Historically the AFI/SAFI definitions for the IPv4 address family only have provisions for
advertising a Next Hop address that belongs to the IPv4 protocol when advertising IPv4 or VPN-IPv4. [RFC8950] specifies the extensions necessary to allow advertising IPv4 NLRI, Virtual Private Network Unicast (VPN-IPv4) NLRI, Multicast Virtual Private Network (MVPN-IPv4) NLRI with a Next Hop address that belongs to the IPv6 protocol. This comprises of an extended next hop encoding MP-REACH BGP capability exchange to allow the address of the Next Hop for IPv4 NLRI, VPN-IPv4 NLRI and MVPN-IPv4 NLRI to also belong to the IPv6 Protocol. [RFC8950] defines the encoding of the Next Hop to determine which of the protocols the address actually belongs to, and a new BGP Capability allowing MP-BGP Peers to discover dynamically whether they can exchange IPv4 NLRI and VPN-IPv4 NLRI with an IPv6 Next Hop.

The current specification for carrying IPv4 NLRI of a given address family via a Next Hop of a different address family is now defined in [RFC8950], and specifies the extended next hop encoding MP-REACH capability extension necessary to do so. This comprises an extension of the AFI/SAFI definitions to allow the address of the Next Hop for IPv4 NLRI or VPN-IPv4 NLRI to belong to either the IPv4 or the IPv6 protocol, the encoding of the Next Hop information to determine which of the protocols the address belongs to, and a new BGP Capability allowing MP-BGP peers to dynamically discover whether they can exchange IPv4 NLRI and VPN-IPv4 NLRI with an IPv6 Next Hop.

With the new extensions defined in [RFC8950] supporting NLRI and next hop address family mismatch, the BGP peer session can now be treated as a pure TCP transport and carry both IPv4 and IPv6 NLRI at the Provider Edge (PE) - Customer Edge (CE) or Inter-AS ASBR to ASBR PE-PE over a single IPv6 TCP session. This allows for the elimination of dual stack from the PE-CE and Inter-AS ASBR-ASBR PE-PE peering point, and now enable the peering to be IPv6-ONLY. The elimination of IPv4 on the PE-CE and Inter-AS ASBR-ASBR PE-PE peering points translates into OPEX expenditure savings of point-to-point infrastructure links as well as /31 address space savings and administration and network management of both IPv4 and IPv6 BGP peers. This reduction decreases the number of PE-CE BGP peers by fifty percent, which is a tremendous cost savings for operators.

While the savings exists at the Edge eBGP PE-CE peering, on the core side PE to Route Reflector (RR) peering carrying <AFI/SAFI> IPv4 <1/1>, VPN-IPv4 <1/128>, and Multicasat VPN <1/129>, there is no savings as the Provider (P) Core is IPv6 Only and thus can only have an IPv6 peer and must use [RFC8950] extended next hop encoding to carrying IPv4 NLRI IPv4 <2/1>, VPN-IPv4 <2/128>, and Multicasat VPN <2/129> over an IPv6 next hop.
This IPv6-Only PE design is applicable to both PE-CE Edge over a IPv4-Only Core, IPv6-Only Core as well as Global table or VPN overlay scenario as well as all Inter-AS Options Option-A, Option-B, Option-AB and Option-C The following Address Family (AFI) / Subsequent Address Family (SAFI) will be tested with both IPv4-Only Core, IPv6-Only Core and Global Routing Table (GRT) and IP Virtual Private Network (VPN) [RFC4364]. <AFI/SAFI> IPv4 <1/1>, VPN-IPV4 <1/128>, and Multicasat VPN <1/129>.

This document provides a much needed solution for Internet Exchange Point (IXP) that are facing IPv4 address depletion at large peering points. With this design, IXP can now use deploy PE-CE IPv6-Only eBGP Edge and Inter-AS peering design to eliminate IPv4 provisioning at the PE Edge as well as PE Inter-AS. This core and edge IPv6-Only peering design paradigm change can apply to any eBGP peering, public internet or private, which can be either Core networks, Data Center networks, Access networks or can be any eBGP peering scenario. This document provides detailed vendor specific test cases and test results for the IPv6-Only peering design as well as successful test results between five major vendors stakeholders in the routing and switching industry, Cisco, Juniper, Arista, Nokia and Huawei. With the test results provided for the IPv6-Only Edge peering design, the goal is that all other vendors around the world that have not been tested will begin to adopt and implement this new best practice for eBGP IPv6-Only Edge peering. This will give confidence to operators to start the proliferation of this IPv6-Only PE design.

As this issue with IXP address depletion is a critical issue around the world, it is imperative for an immediate solution that can be implemented quickly. This best practice IPv6-only eBGP peering design specification will help proliferate IPv6-Only deployments at the eBGP Edge and Inter-AS network peering points starting immediately at a minimum with operators around the world using Cisco, Juniper, Arista, Nokia and Huawei.

2. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

3. Terminology

Terminolgy used in defining the IPv6-Only Edge specification.

AFBR: Address Family Border Router Provider Edge (PE).
Edge: PE-CE Edge Network Provider Edge - Customer Edge

Core: P Core Network Provider (P)

4to6 Softwire: IPv4 edge over an IPv6-Only core

6to4 Softwire: IPv6 edge over an IPv4-Only core

E2E: End to End

4. IPv6-Only Edge Peering Architecture

4.1. Problem Statement

This specification addresses a real issue that has been discussed at many operator groups around the world related to IXP major peering points where hundreds of AS’s have both IPv4 and IPv6 dual stacked peering. IPv4 address depletion have been a major issue issue for many years now. Operators around the world are clamoring for a solution that can help solve issues related to IPv4 address depletion at these large IXP peering points. With this solution IXPs as well as all infrastructure networks such as Core networks, DC networks, Access networks as well as any PE-CE public or private network can now utilize this IPv6-Only Edge solution and reap the benefits immediately on IPv4 address space saving.

IXP Problem Statement

Dual Stacked

CE

+--------+ IPv4 BGP Peer +--------+

CE

+--------+ IPv6 BGP Peer +--------+

CE

+--------+ IPv4 forwarding +--------+

IPV4 forwarding

IPv6 forwarding

IPv6 forwarding

IPv4 forwarding

Figure 1: Problem Statement - IXP Dual Stack Peering
4.2. IPv6-Only PE-CE Design Solution

The IPv6-Only Edge design solution provides a means of E2E single protocol design solution extension of [RFC5565] Softwire Mesh framework from the PE-CE Edge to the Core from ingress so egress through the entire operators domain. This solution eliminates all IPv4 addressing from end to end while still providing the same Dual Stack functionality of IPv4 and IPv6 packet forwarding from a data plane perspective by leveraging the [RFC8950] extended next hop encoding so that IPv4 NLRI can be advertised over a single IPv6 pure transport TCP session. This IPv6-Only E2E architecture eliminates all IPv4 peering and IPv4 addressing E2E from the ingress CE to ingress PE to egress PE to egress CE and all hops along the operator E2E path.

Solution applicable to any Edge peering scenario - IXP, Core, DC, Access, etc

Figure 3: IPv6-Only Solution Applicability
4.3. IPv6-Only Edge Peering Design

4.3.1. IPv6-Only Edge Peering Packet Walk

The IPv6-Only Edge Peering design utilizes two key E2E Softwire Mesh Framework scenario’s, 4to6 softwire and 6to4 softwire. The Softwire mesh framework concept is based on the overlay and underlay MPLS or SR based technology framework, where the underlay is the transport layer and the overlay is a Virtual Private Network (VPN) layer, and is the the tunneled virtualization layer containing the customer payload. The concept of a 6to4 Softwire is based on transmission of IPv6 packets at the edge of the network by tunneling the IPv6 packets over an IPv4-Only Core. The concept of a 4to6 Softwire is also based on transmission of IPv4 packets at the edge of the network by tunneling the IPv4 packets over an IPv6-Only Core.

This document describes End to End (E2E) test scenarios that follow a packet flow from IPv6-Only attachment circuit from ingress PE-CE to egress PE-CE tracing the routing protocol control plane and data plane forwarding of IPv4 packets in a 4to6 softwire or 6to4 softwire within the IPv4-Only or IPv6-Only Core network. In both scenario we are focusing on IPv4 packets and the control plane and data plane forwarding aspects of IPv4 packets from the PE-CE Edge network over an IPv6-Only P (Provider) core network or IPv4-Only P (Provider) core network. With this IPv6-Only Edge peering design, the Softwire Mesh Framework is not extended beyond the Provider Edge (PE) and continues to terminate on the PE router.
4.3.2. 6to4 Softwire IPv4-Only Core packet walk

6to4 softwire where IPv6-Edge eBGP IPv6 peering where IPv4 packets at network Edge traverse a IPv4-Only Core

In the scenario where IPv4 packets originating from a PE-CE edge are tunneled over an MPLS or Segment Routing IPv4 underlay core network, the PE and CE only have an IPv6 address configured on the interface. In this scenario the IPv4 packets that ingress the CE from within the CE AS are over an IPv6-Only interface and are forwarded to an IPv4 NLRI destination prefix learned from the Pure Transport Single IPv6 BGP Peer. In the IPv6-Only Edge peering architecture the PE is IPv6-Only as all PE-CE interfaces are IPv6-Only. However, on the CE, the PE-CE interface is the only interface that is IPv6-Only and all other interfaces may or may not be IPv6-Only. Following the data plane packet flow, IPv4 packets are forwarded from the ingress CE to the IPv6-Only ingress PE where the VPN label imposition push per prefix, per-vrf, per-CE occurs and the labeled packet is forwarded over a 6to4 softwire IPv4-Only core, to the egress PE where the VPN label disposition pop occurs and the native IPv4 packet is forwarded to the egress CE. In the reverse direction IPv4 packets are forwarded from the egress CE to egress PE where the VPN label imposition per prefix, per-vrf, per-CE push occurs and the labeled packet is forwarded back over the 6to4 softwire IPv4-Only core, to the ingress PE where the VPN label disposition pop occurs and the native IPv4 packet is forwarded to the ingress CE. The functionality of the IPv4 forwarding plane in this scenario is identical from a data plane forwarding perspective to Dual Stack IPv4 forwarding scenario.
4.3.3. 4to6 Softwire IPv6-Only Core packet walk

4to6 softwire where IPv6-Edge eBGP IPv6 peering where IPv4 packets at network Edge traverse a IPv6-Only Core
In the scenario where IPv4 packets originating from a PE-CE edge are tunneled over an MPLS or Segment Routing IPv4 underlay core network, the PE and CE only have an IPv6 address configured on the interface. In this scenario the IPv4 packets that ingress the CE from within the CE AS are over an IPv6-Only interface and are forwarded to an IPv4 NLRI destination prefix learned from the Pure Transport Single IPv6 BGP Peer. In the IPv6-Only Edge peering architecture the PE is IPv6-Only as all PE-CE interfaces are IPv6-Only. However, on the CE, the PE-CE interface is the only interface that is IPv6-Only and all other interfaces may or may not be IPv6-Only. Following the data plane packet flow, IPv4 packets are forwarded from the ingress CE to the IPv6-Only ingress PE where the VPN label imposition push per prefix, per-vrf, per-CE occurs and the labeled packet is forwarded over a 4to6 softwire IPv6-Only core, to the egress PE where the VPN label disposition pop occurs and the native IPv4 packet is forwarded to the egress CE. In the reverse direction IPv4 packets are forwarded from the egress CE to egress PE where the VPN label imposition per prefix, per-vrf, per-CE push occurs and the labeled packet is forwarded back over the 4to6 softwire IPv6-Only core, to the ingress PE where the VPN label disposition pop occurs and the native IPv4 packet is forwarded to the ingress CE. The functionality of the IPv4 forwarding plane in this scenario is identical from a data plane forwarding perspective to Dual Stack IPv4 forwarding scenario.
Figure 6: 4to6 Softwire - IPv4 Edge over an IPv6-Only Core

4.4. RFC5549 and RFC8950 Applicability
4.4.1. IPv6-Only Edge Peering design next-hop encoding

This section describes [RFC8950] next hop encoding updates to [RFC5549] applicability to this specification. IPv6-only eBGP Edge PE-CE peering to carry IPv4 Unicast NLRI <AFI/SAFI> IPv4 <1/1> over an IPv6 next hop BGP capability extended hop encoding IANA capability codepoint value 5 defined is applicable to both [RFC5549] and [RFC8950] as IPv4 Unicast NLRI <AFI/SAFI> IPv4 <1/1> does not change in the RFC updates.

IPv4 packets over an IPv6-Only core 4to6 Softwire E2E packet flow is part of the IPv6-Only design vendor interoperability test cases and in that respect is applicable as [RFC8950] updates [RFC5549] for <AFI/SAFI> VPN-IPv4 <1/128>, and Multicasat VPN <1/129>

4.4.2. RFC8950 updates to RFC5549 applicability

This section describes the [RFC8950] next hop encoding updates to [RFC5549]

In [RFC5549] when AFI/SAFI 1/128 is used, the next-hop address is encoded as an IPv6 address with a length of 16 or 32 bytes. This document modifies how the next-hop address is encoded to accommodate all existing implementations and bring consistency with VPNv4oIPv4 and VPNv6oIPv6. The next-hop address is now encoded as a VPN-IPv6 address with a length of 24 or 48 bytes [RFC8950] (see Sections 3 and 6.2 of this document). This change addresses Erratum ID 5253 (Err5253). As all known and deployed implementations are interoperable today and use the new proposed encoding, the change does not break existing interoperability. Updates to [RFC8950] is applicable to the IPv6-Only PE-CE edge design for the IPv6 next hop encoding E2E test case of IPv4 packets over and IPv6-Only core 4to6 Softwire. In this test case IPv4 Unicast NLRI <AFI/SAFI> IPv4 <1/1> is advertised over the PE to RR core peering 4to6 softwire in <AFI/SAFI> VPN-IPv4 <1/128>. In this test case label allocation mode comes into play which is discussed in section 8.9.

[RFC5549] next hop encoding of MP_REACH_NLRI with:

* NLRI= NLRI as per current AFI/SAFI definition

Advertising with [RFC4760] MP_REACH_NLRI with:

* AFI = 1
* SAFI = 128 or 129
* Length of Next Hop Address = 16 or 32
* NLRI= NLRI as per current AFI/SAFI definition

[RFC8950] next hop encoding of MP_REACH_NLRI with:

* NLRI= NLRI as per current AFI/SAFI definition

Advertising with [RFC4760] MP_REACH_NLRI with:

* AFI = 1
* SAFI = 128 or 129
* Length of Next Hop Address = 24 or 48
* Next Hop Address = VPN-IPv6 address of next hop with an 8-octet RD set to zero (potentially followed by the link-local VPN-IPv6 address of the next hop with an 8-octet RD is set to zero).

* NLRI= NLRI as per current AFI/SAFI definition

5. IPv6-Only PE Design Edge and Inter-AS Options E2E Test Cases

Proof of concept interoperability testing of the 4 test cases between the 5 vendors Cisco, Juniper, Arista, Nokia and Huawei.

Cisco, Juniper, Arista, Nokia, Huawei, platform, code revision and test results for all use cases

Cisco: Edge Router- XR ASR 9910 IOS XR 7.4.1, Core Router- NCS 6000 7.2.2, CRS-X 6.7.4

Juniper: Edge Router- MX platform MX480, MX960, Core Router- PTX Platform PFX5000, PTC10K8 (JUNOS and EVO) Release 20.4R2

Tested v4 edge over v6 core in a virtual setup using vMX platform and 20.4R2 and LDPv6 as underlay, but there were some data plane forwarding issues. Tested same setup on latest release 21.4 and it worked. Investigating what the minimum version is for this setup to work.
Tested on above Juniper platforms. Completed IPv6-Only PE design functionality test with PE-CE IPv6 peer carrying IPv4 and IPv6 prefixes control plane validation and data plane forwarding plane validation and verified end to end reachability CE to CE forwarding plane with Default Per-CE label allocation mode. Tested with IPv4-Only Core and IPv6-Only Core and proved that the IPv6-Only PE design solution works. Both IPv4 and IPv6 packets were forwarded identical functionality of Dual Stack without having IPv4 address configured.

Nokia: Edge and Core-7750 Service Router, Release R21
Huawei: Edge and Core-VRPv8, Release VRP-V800R020C10
Arista:

Intra-AS tests PE-CE Edge Peering IPv4-Only Core, IPv6-Only Core, Global Table (GRT) and IP VPN
AFI/SAFI IPv4-Unicast SAFI IPv6-Unicast SAFI

IPv4 Core:
Test-1 Global table (6PE)
Test-2 IP VPN
Global table IPv6

IPv6 Core:
Test-3 Global table
Test-4 IP VPN

Inter-AS Options tests IPv4-Only Core, IPv6-Only Core, Global Table (GRT) and IP VPN
AFI/SAFI VPN and MVPN
IPv4-Only Core
Test-5 Global table 6PE Option-B (Segmented LSP stitched IPv4 Core - Inter-AS Link IPv6-Only PE - IPv4 Core)
Test-6 Global table 6PE Option-C (Redistribute IPv4 Loopbacks into BGP-LU AFI/SAFI 2/6)
Test-7 IP VPN Inter AS Option-B (Segmented LSP stitched IPv4 Core – Inter-AS Link IPv6-Only PE - IPv4 Core)

Test-8 IP VPN Inter AS Option-C (Redistribute IPv4 Loopbacks into BGP-LU AFI/SAFI 2/6)

IPv6-Only Core

Test-9 Global table Option-B

Test-10 Global table Option-C

Test-11 IP VPN Inter AS Option-B

Test-12 IP VPN Inter AS Option-C

5.1. Test-1 E2E IPv6-Only PE-CE, Global Table over IPv4-Only Core (6PE), 6to4 softwire

Figure 7: Test-1 E2E IPv6-Only PE-CE, Global Table over IPv4-Only Core (6PE)

5.2. Test-2 E2E IPv6-Only PE-CE, VPN over IPv4-Only Core, 6to4 Softwire
5.3. Test-3 E2E IPv6-Only PE-CE, Global Table over IPv6-Only Core, 4to6 Softwire

Figure 9: Test-3 E2E IPv6-Only PE-CE, Global Table over IPv6-Only Core

5.4. Test-4 E2E IPv6-Only PE-CE, VPN over IPv6-Only Core, 4to6 Softwire
5.5. Test-5 E2E IPv6-Only PE-CE, Global Table over IPv4-Only Core(6PE), 6to4 softwire -Inter-AS Option-B

5.6. Test-6 E2E IPv6-Only PE-CE, Global Table over IPv4-Only Core(6PE), 6to4 softwire -Inter-AS Option-C
5.7. Test-7 E2E IPv6-Only PE-CE, VPN over IPv4-Only, 6to4 softwire - Inter-AS Option-B

Figure 13: Test-7 E2E IPv6-Only PE-CE, VPN over IPv4-Only Core - Inter-AS Option-B

5.8. Test-8 E2E IPv6-Only PE-CE, VPN over IPv4-Only Core, 6to4 softwire - Inter-AS Option-C
5.9. Test-9 E2E IPv6-Only PE-CE, Global Table over IPv6-Only Core, 4to6 softwire - Inter-AS Option-B

5.10. Test-10 E2E IPv6-Only PE-CE, Global Table over IPv6-Only Core, 4to6 softwire - Inter-AS Option-C
Inter-AS ASBR-ASBR link is IPv6-Only PE

IPv6-Only
PE / CE
--- +----+  AS 1 IPv6
- CE - PE -- IPv6-Only Core 0=Underlay==0 --- IPv6-Only Core 0==Underlay==0 -- PE - CE--- +----+
IPv6 BGP peer MPLS/SR IPv6 BGP peer
IPv4 forwarding IPv4 forwarding
IPv6 forwarding IPv6 forwarding

Figure 16: Test-10 E2E IPv6-Only PE-CE, Global Table over IPv6-Only Core - Inter-AS Option-C

5.11. Test-11 E2E IPv6-Only PE-CE, VPN over IPv6-Only Core, 4to6 softwire - Inter-AS Option-B

Inter-AS ASBR-ASBR link is IPv6-Only PE

IPv6-Only
PE / CE
--- +----+  AS 1 IPv6
- CE - PE -- IPv6-Only Core 0=Overlay===0 --- IPv6-Only Core 0==Overlay===0 -- PE - CE--- +----+
IPv6 BGP peer MPLS/SR IPv6 BGP peer
IPv4 forwarding IPv4 forwarding
IPv6 forwarding IPv6 forwarding

Figure 17: Test-11 E2E IPv6-Only PE-CE, VPN over IPv6-Only Core - Inter-AS Option-B

5.12. Test-12 E2E IPv6-Only PE-CE, VPN over IPv6-Only Core, 4to6 softwire - Inter-AS Option-C
5.13. IPv6-Only PE-CE Operational Considerations Testing

Ping CE to PE when destination prefix is withdrawn
Traceroute CE to PE and test all ICMPv4 and ICMPv6 type codes

6. Operational Considerations

With a single IPv6 Peer carrying both IPv4 and IPv6 NLRI there are some operational considerations in terms of what changes and what does not change.

What does not change with a single IPv6 transport peer carrying IPv4 NLRI and IPv6 NLRI below:

Routing Policy configuration is still separate for IPv4 and IPv6 configured by capability as previously.
Layer 1, Layer 2 issues such as one-way fiber or fiber cut will impact both IPv4 and IPv6 as previously.

If the interface is in the Admin Down state, the IPv6 peer would go down, and IPv4 NLRI and IPv6 NLRI would be withdrawn as previously.

Changes resulting from a single IPv6 transport peer carrying IPv4 NLRI and IPv6 NLRI below:

Physical interface is no longer dual stacked.

Any change in IPv6 address or DAD state will impact both IPv4 and IPv6 NLRI exchange.

Single BFD session for both IPv4 and IPv6 NLRI fate sharing as the session is now tied to the transport, which now is only IPv6 address family.

Both IPv4 and IPv6 peer now exists under the IPv6 address family configuration.

Fate sharing of IPv4 and IPv6 address family from a logical perspective now carried over a single physical IPv6 peer.

From an operations perspective, prior to elimination of IPv4 peers, an audit is recommended to identify and IPv4 and IPv6 peering incongruencies that may exist and to rectify them. No operational impacts or issues are expected with this change.

With MPLS VPN overlay, per-CE next-hop label allocation mode where both IPv4 and IPv6 prefixes have the same label in no table lookup pop-n-forward mode should be taken into consideration.

7. IANA Considerations

There are not any IANA considerations.

8. Security Considerations

The extensions defined in this document allow BGP to propagate reachability information about IPv4 prefixes over an MPLS or SR IPv6-Only core network. As such, no new security issues are raised beyond those that already exist in BGP-4 and the use of MP-BGP for IPv6. Both IPv4 and IPv6 peers exist under the IPv6 address family configuration. The security features of BGP and corresponding security policy defined in the ISP domain are applicable. For the inter-AS distribution of IPv6 routes according to case (a) of Section 4 of this document, no new security issues are raised beyond
those that already exist in the use of eBGP for IPv6 [RFC2545].

9.  Acknowledgments

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10.  Contributors

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11.  References

11.1.  Normative References


11.2. Informative References


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A YANG Model for SRv6 Mobile User Plane
draft-mahesh-spring-srv6-mobile-yang-00

Abstract

This document defines a YANG data model for configuration and management of SRv6 for the mobile network.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

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1. Introduction

In mobile networks, mobility systems provide connectivity over a wireless link to stationary and non-stationary nodes. The user-plane establishes a tunnel between the mobile node and its anchor node over IP-based backhaul and core networks.

When SRv6 is applied to mobile networks, it enables a source routing architecture, where operators get to explicitly specify a route for the packets to traverse both to and from a mobile node. The SRv6 Endpoint nodes serve as mobile user-plane anchors.

For example, in an Enhanced mode topology, the intermediate waypoints, SIDs, can be used for Traffic Engineering. For more details, see Segment Routing IPv6 for Mobile User Plane [I-D.ietf-dmm-srv6-mobile-uplane], Section 5.2. The gNB and UPF are SR-aware, and there are two service segments, one for traffic engineering to support a low latency path, and the other for service programming. In such a topology the operator routes the traffic through these SRv6 nodes, so they can perform their Endpoint functionality and forward the packet. Further, in the uplink direction, when the gNB receives a packet from a UE, it adds the segments of the SR policy to route the traffic through those two segments, while doing something similar in the downlink direction.

This document describes a YANG 1.1 [RFC7950] data model for the Segment Routing IPv6 (SRv6) user plane of mobile networks.

The model conforms to the NMDA [RFC8342] architecture.
1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

2. Terminology

This document references terms defined in other documents. In particular, it imports definitions for the following terms from Segment Routing Architecture [RFC8402], and IPv6 Segment Routing Header (SRH) [RFC8754].

* Active Segment
* BGP-Prefix Segment
* Prefix SID
* Segment
* SID
* SRH
* SRv6
* SRv6 Endpoint nodes
* SRv6 SID
* Segment Routing domain (SR domain)
* SR Global Block (SRGB)
* SR Local Block (SRLB)

2.1. Acronyms

This document uses a few acronyms. Some of them are defined here for reference.
<table>
<thead>
<tr>
<th>Acronym</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>gNB</td>
<td>gNodeB, a 5G Base Station using New Radio technology</td>
</tr>
<tr>
<td>MUP</td>
<td>Mobile User Plane</td>
</tr>
<tr>
<td>SR</td>
<td>Segment Routing</td>
</tr>
<tr>
<td>SRv6</td>
<td>Segment Routing over v6</td>
</tr>
<tr>
<td>UE</td>
<td>User Equipment</td>
</tr>
<tr>
<td>UPF</td>
<td>User Plane Function</td>
</tr>
</tbody>
</table>

Table 1: Acronyms

3. Tree Diagram

An abridged version of the tree diagram is shown here. Annotations used in the diagram are defined in YANG Tree Diagrams [RFC8340].
module: ietf-srv6-mobile

augment /rt:routing/rt:control-plane-protocols
    /rt:control-plane-protocol/bgp:bgp/bgp:global:
        +++rw route-distinguisher?     rt-types:route-distinguisher
        +++rw label-allocation-mode?   identityref
        +++rw sid-allocation-mode?     identityref
        +++rw srv6
            +++rw locator?                 leafref
            +++ro sid-manager-connected?   boolean
            +++ro locator-registered?      boolean
            +++ro micro-segment-enabled?   boolean
        +++rw mobile
            +++rw encapsulation
            |     ...
            +++rw decapsulations
            |     ...
            +++rw decapsulation-source-prefix?   inet:ipv6-prefix

augment /rt:routing/rt:control-plane-protocols
    /rt:control-plane-protocol/bgp:bgp/bgp:global
        /bgp:route-selection-options:
            +++rw selection-deferral-time?   uint16
            +++rw med-missing-as-worst?      boolean
            +++rw multipath-as-path-relax?   boolean
            +++rw multipath-next-hop-relax?  boolean

augment /rt-pol:routing-policy/rt-pol:defined-sets
    /bp:bgp-defined-sets:
        +++rw n4-interface-sets
            +++rw interface-set* [name]
                +++rw name      string
                +++rw member*   identityref

augment /rt-pol:routing-policy/rt-pol:policy-definitions
    /rt-pol:policy-definition/rt-pol:statements
        /rt-pol:statement/rt-pol:conditions/bp:bgp-conditions:
            +++rw match-n4-network-interface-set
                +++rw n4-network-instance-set?   leafref
                +++rw match-set-options?         match-set-options-type

augment /rt-pol:routing-policy/rt-pol:policy-definitions
    /rt-pol:policy-definition/rt-pol:statements
        /rt-pol:statement/rt-pol:actions/bp:bgp-actions:
            +++rw set-network-interface
                +++rw apply-policy
                    +++rw import-policy*           leafref
                    +++rw default-import-policy?   default-policy-type
                    +++rw export-policy*           leafref
                    +++rw default-export-policy?   default-policy-type
4. YANG Model

The YANG model is divided into two parts. The first part of the model augments the BGP model in BGP Model for Service Provider Network [I-D.ietf-idr-bgp-model] for the BGP configuration, while the second part augments the BGP Routing Policy model in BGP Model for Service Provider Network [I-D.ietf-idr-bgp-model].

The BGP model is augmented both at a global level to add SRv6 configuration, and at the route selection option. The BGP policy model is augmented to add a defined set, a set of match options, and a set of actions.

The model imports Common YANG Data Types [RFC6991], A YANG Data Model for Routing Management (NMDA Version) [RFC8349], A YANG Data Model for Routing Policy [RFC9067], YANG Data Model for Segment Routing [RFC9020], YANG Data Model for SRv6 Base and Static [I-D.ietf-spring-srv6-yang], and BGP Model for Service Provider Network [I-D.ietf-idr-bgp-model].

```yang
<CODE BEGINS> file "ietf-srv6-mobile@2022-03-03.yang"
module ietf-srv6-mobile {
  yang-version "1.1";
  namespace "urn:ietf:params:xml:ns:yang:ietf-srv6-mobile";
  prefix "srv6-mob";

  import ietf-inet-types {
    prefix "inet";
    reference
      "RFC 6991: Common YANG Data Types.";
  }
  import ietf-routing {
    prefix rt;
    reference
      "RFC 8349, A YANG Data Model for Routing Management (NMDA Version).";
  }
  import ietf-routing-types {
    prefix rt-types;
    reference
      "RFC 8294: Common YANG Types for the Routing Area.";
  }
  import ietf-routing-policy {
    prefix rt-pol;
    reference
      "RFC 9067: A YANG Data Model for Routing Policy.";
  }
  import ietf-bgp {
```
prefix bgp;
reference
"I-D.ietf-idr-bgp-model: BGP Model for Service Provider
Network.";
}
import ietf-bgp-policy {
  prefix bp;
  reference
  "I-D.ietf-idr-bgp-model: BGP Model for Service Provider
  Network.";
}
import ietf-bgp-types {
  prefix bt;
  reference
  "I-D.ietf-idr-bgp-model: BGP Model for Service Provider
  Network.";
}
import ietf-segment-routing {
  prefix sr;
  reference
  "RFC 9020: YANG Data Model for Segment Routing.";
}
import ietf-srv6-base {
  prefix srv6;
  reference
  "I-D.ietf-spring-srv6-yang: YANG Data Model for SRv6 Base
  and Static.";
}
import ietf-srv6-types {
  prefix srv6-types;
  reference
  "RFC 9020: YANG Data Model for Segment Routing.";
}
organization
"IETF SPRING Working Group";
contact
"WG Web: <https://datatracker.ietf.org/wg/spring/about>
WG List: <spring@ietf.org>
Editor: Mahesh Jethanandani (mjethanandani at gmail dot com)
Author: Tetsuya Murakami (tetsuya at arrcus dot com)";
description
"This module augments the BGP YANG model to add support for
configuration in mobile networks."
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This version of this YANG module is part of RFC XXXX (https://www.rfc-editor.org/info/rfcXXXX); see the RFC itself for full legal notices.


revision "2022-03-03" {
  description
    "Initial Version.";
  reference
    "RFC XXXX, A YANG Model for BGP configuration in mobile networks.";
}

/*
 * Features
 */

/*
 * Typedefs
 */
typedef srv6-mobile-function-type {
  type union {
    type identityref {
      base "srv6-types:srv6-endpoint-type";
    }
    type uint16;
  }
  description
    "Type definition for SRV6 mobile function. Unknown values are represented as the numeric value.";
  reference
    "draft-ietf-dmm-srv6-mobile-uplane";
}
/ * Identities
 */

identity srv6-mup {
    base bt:afi-safi-type;
    description
        "Segment Routing for Mobile User Plane (AFI,SAFI = ?,?)";
    reference
        "RFC XXXX: A YANG Model for BGP configuration in mobile networks.";
}

identity label-allocation-mode {
    description
        "Base identity to be used to express types of label allocation strategies to be used within a network instance";
}

identity label-per-prefix {
    base label-allocation-mode;
    description
        "A label is to be allocated per prefix entry in the RIB for the network instance";
}

identity label-per-nexthop {
    base label-allocation-mode;
    description
        "A label is to be allocated per nexthop entry in the RIB for the network instance";
}

identity label-instance {
    base label-allocation-mode;
    description
        "A single label is to be used for the instance";
}

identity sid-allocation-mode {
    description
        "Base identity to be used to express types of SRv6 segment ID allocation strategies to be used within a network instance.";
}

identity sid-per-nexthop {
    base sid-allocation-mode;
    description
        "A segment ID is to be allocated per nexthop entry in the RIB";
identity sid-instance {  
    base sid-allocation-mode;  
    description  
        "A single segment ID is to be used for the instance";
}

identity sid-per-nexthop-no-transportation {  
    base sid-allocation-mode;  
    description  
        "A segment ID is to be allocated per nexthop entry in the  
         RIB for the network instance";
}

identity sid-instance-no-transportation {  
    base sid-allocation-mode;  
    description  
        "A single segment ID is to be used for the instance";
}

identity mobile-interface-type {  
    description  
        "Base identity for different mobile interfaces.";
}

identity n4 {  
    base mobile-interface-type;  
    description  
        "N4 interface.";
}

/*
 * Groupings
 */

/*
 * BGP configuration
 */
augment "/rt:routing/rt:control-plane-protocols" +  
    "/rt:control-plane-protocol/bgp:bgp/bgp:global" {
    description  
        "Augmentation of the BGP global configuration to add srv6  
         mobile configuration.";

    leaf route-distinguisher {

type rt-types:route-distinguisher;

description
"The route distinguisher that should be used for the local VRF or VSI instance when it is signalled via BGP."
}

leaf label-allocation-mode {
  type identityref {
    base label-allocation-mode;
  }
  must "not(../sid-allocation-mode)" {
    error-message "label-allocation-mode and sid-allocation-mode cannot co-exist";
  }
    error-message "SRv6 configurations must be removed first";
  }
  description
"The label allocation mode to be used for L3 entries in the network instance";
}

leaf sid-allocation-mode {
  type identityref {
    base sid-allocation-mode;
  }
  must "not(../label-allocation-mode)" {
    error-message "label-allocation-mode and sid-allocation-mode cannot co-exist";
  }
    error-message "SRv6 locator name must be configured";
  }
  description
"The route distinguisher that should be used for the local VRF or VSI instance when it is signalled via BGP."
}
"The segment ID allocation mode to be used for L3 entries in the network instance";
}

container srv6 {
  description
  "SRv6 mobile container.";

  leaf locator {
    type leafref {
      path "/rt:routing/sr:segment-routing/" +
      "srv6:srv6/srv6:locators/srv6:locator/srv6:name";
    }
    description
    "Locator configuration.";
  }

  leaf sid-manager-connected {
    type boolean;
    config false;
    description
    "Connection with segment ID manager is active";
  }

  leaf locator-registered {
    type boolean;
    config false;
    description
    "Locator name is registered";
  }

  leaf micro-segment-enabled {
    type boolean;
    config false;
    description
    "Locator has enabled micro-segment behavior";
  }

  container mobile {
    when ":derived-from-or-self(/rt:routing" +
    "/rt:control-plane-protocols" +
    "/rt:control-plane-protocol/bgp:bgp/bgp:global" +
    "/bgp:afi-safis/bgp:afi-safi/bgp:name, 'srv6-mup')" {
      description
      "This augmentation is valid only for a MUP SAFI.";
    }
  }
  description
  "This augmentation is valid only for a MUP SAFI.";
}
"Mobile configuration of SRv6."

container encapsulation {
    description "Encapsulation configuration.";

    leaf locator {
        type leafref {
            path "/rt:routing/sr:segment-routing/" +
            "srv6:srv6:locators/srv6:locator/srv6:name";
        }
        description "Reference to SRv6 locator key";
    }

    leaf function {
        type srv6-mobile-function-type;
        must "boolean(current()/../locator)" {
            error-message
            "SRv6 Mobile Locator name must be configured";
        }
        description "One of the SRv6 function types.";
    }

    leaf source-adress {
        type inet:ipv4-address;
        description "GTP source IP address";
    }

    leaf source-position {
        type uint8;
        description "Bit position of GTP source IP address";
    }

    container n4-network-instance {
        description "Definitions for the N4 interface.";

        leaf routing-policy {
            type leafref {
                path "/rt-pol:routing-policy/" +
                "rt-pol:policy-definitions/" +
                "rt-pol:policy-definition/rt-pol:name";
            }
            must "boolean(current()/../access)" {
            }
error-message
  "SRv6 Mobile access instance name must be " +
  "configured";
}
description
  "Reference to routing-policy";
}

leaf access {
  type string;
  description
    "Mobile access instance.";
}
}
}

container decapsulations {
  description
    "SRv6 mobile decapsulation configuration.";
}

list decapsulation {
  key "id";
  description
    "SRv6 mobile Decapsulation config";
  leaf id {
    type uint16;
    description
      "SRv6 mobile decapsulation entry id";
  }
  leaf locator {
    type leafref {
      path "/rt:routing/sr:segment-routing/" +
        "srv6:srv6/srv6:locators/srv6:locator/srv6:name";
    }
    description
      "Reference to SRv6 locater key";
  }
  leaf function {
    type srv6-mobile-function-type;
    must "boolean(current()/../locator)" { 
      error-message
        "SRv6 Mobile Locator name must be configured";
    }
    description
      "One of SRv6 function types.";
  }
container n4-network-instance {
    description "Definitions for the N4 interface.";
    leaf core {
        type string;
        description "Core instance";
    }
}

leaf decapsulation-source-prefix {
    type inet:ipv6-prefix;
    description "IPv6 prefix for GTP source address";
}

    description "Augmentation of the BGP global configuration for route selection options to add srv6 mobile configuration.";
    leaf selection-deferral-time {
        type uint16 { 
            range 1..3600;
        }
        default 300;
        description "An upper-bound on the time (in seconds) that the best-path selection is deferred";
    }
    leaf med-missing-as-worst {
        type boolean;
        description "A route without MED is treated as with highest MED value";
    }
    leaf multipath-as-path-relax {
type boolean;
default true;
description
"Paths with different AS-Path but of same length can form ECMP";
}

leaf multipath-nexthop-relax {
    type boolean;
default false;
description
"Enable BGP multi-path for paths with same next-hop";
}

augment "/rt-pol:routing-policy/rt-pol:defined-sets" +
"/bp:bgp-defined-sets" {
    description
"Augmentation of the Routing Policy module to add mobile interface defined sets.";
}

container n4-interface-sets {
    description
"Enclosing container for list of n4 interface sets.";

    list interface-set {
        key "name";
        description
"List of defined interface sets.";

        leaf name {
            type string;
            description
"Name of interface set. This is used to reference the set in match conditions.";
        }

        leaf-list member {
            type identityref {
                base "mobile-interface-type";
            }
            description
"Members of interface set.";
        }
    }
}
augment "/rt-pol:routing-policy/rt-pol:policy-definitions" + 
"/rt-pol:policy-definition/rt-pol:statements" + 
"/rt-pol:statement/rt-pol:conditions/bp:bgp-conditions" { 
  description 
  "Augmentation of the Routing Policy module to add conditions.";

container match-n4-network-interface-set {
  description 
  "Match a referenced network instance.";
  leaf n4-network-instance-set {
    type leafref {
      path "/rt-pol:routing-policy/rt-pol:defined-sets/" + 
      "bp:bgp-defined-sets/n4-interface-sets/" + 
      "interface-set/name";
    }
    description 
    "References a defined community set.";
  }
  uses rt-pol:match-set-options-group;
}
}

augment "/rt-pol:routing-policy/rt-pol:policy-definitions" + 
"/rt-pol:policy-definition/rt-pol:statements" + 
"/rt-pol:statement/rt-pol:actions/bp:bgp-actions" { 
  description 
  "Augmentation of the Routing Policy module to add actions.";

container set-network-interface {
  description 
  "Set a referenced network instance.";
  uses rt-pol:apply-policy-group;
}
}

<CODE ENDS>

Figure 2: SRv6 YANG Model for Mobile User Plane

5. IANA Considerations

This memo registers the following namespace URIs in the IETF XML in the "IETF XML Registry" [RFC3688]:

Registrant Contact:  The IESG.
XML:  N/A; the requested URI is an XML namespace.
This document registers the following YANG modules in the "YANG Module Names" registry [RFC6020]:

Name:  ietf-srv6-mobile
Prefix:  srv6-mob
Reference:  RFC XXXX

6. Security Considerations

The YANG module specified in this document defines a schema for data that is designed to be accessed via network management protocols such as NETCONF [RFC6241] or RESTCONF [RFC8040]. The lowest NETCONF layer is the secure transport layer, and the mandatory-to-implement secure transport is Secure Shell (SSH) [RFC6242]. The lowest RESTCONF layer is HTTPS, and the mandatory-to-implement secure transport is TLS [RFC8446].

The Network Configuration Access Control Model (NACM) [RFC8341] provides the means to restrict access for particular NETCONF or RESTCONF users to a preconfigured subset of all available NETCONF or RESTCONF protocol operations and content.

There are a number of data nodes defined in this YANG module that are writable/creatable/deletable (i.e., config true, which is the default). These data nodes may be considered sensitive or vulnerable in some network environments. Write operations (e.g., edit-config) to these data nodes without proper protection can have a negative effect on network operations. These are the subtrees and data nodes and their sensitivity/vulnerability:

Some of the readable data nodes in this YANG module may be considered sensitive or vulnerable in some network environments. It is thus important to control read access (e.g., via get, get-config, or notification) to these data nodes. These are the subtrees and data nodes and their sensitivity/vulnerability:

Some of the RPC operations in this YANG module may be considered sensitive or vulnerable in some network environments. It is thus important to control access to these operations. These are the operations and their sensitivity/vulnerability:

7. References

7.1. Normative References


Appendix A. Appendix 1 Complete Tree Diagram

Here is a complete tree diagram for the configuration and operational part of the model.

```yang
module: ietf-srv6-mobile

augment /rt:routing/rt:control-plane-protocols
    /rt:control-plane-protocol/bgp:bgp/bgp:global:
        +++-rw route-distinguisher? rt-types:route-distinguisher
        +++-rw label-allocation-mode? identityref
        +++-rw sid-allocation-mode? identityref
        +++-rw srv6
            +++-rw locator? leafref
            +++-ro sid-manager-connected? boolean
            +++-ro locator-registered? boolean
            +++-ro micro-segment-enabled? boolean
            +++-rw mobile
                +++-rw encapsulation
                    +++-rw locator? leafref
                    +++-rw function? srv6-mobile-function-type
                    +++-rw source-addr? inet:ipv4-address
                    +++-rw source-position? uint8
                    +++-rw n4-network-instance
                        +++-rw routing-policy? leafref
                        +++-rw access? string
                +++-rw decapsulations
                    +++-rw decapsulation* [id]
                        +++-rw id uint16
                        +++-rw locator? leafref
                        +++-rw function? srv6-mobile-function-type
                        +++-rw n4-network-instance
                            +++-rw core? string
                    +++-rw decapsulation-source-prefix? inet:ipv6-prefix

augment /rt:routing/rt:control-plane-protocols
    /rt:control-plane-protocol/bgp:bgp/bgp:global
    /bgp:route-selection-options:
```
Figure 3: Complete tree diagram

Acknowledgements

TBA

Contributors

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IPv6-Only PE Design for IPv4-NLRI All SAFI over IPv6-NH
draft-mishra-bess-ipv4nlri-all-safi-ipv6nh-01

Abstract

As Enterprises and Service Providers upgrade their brown field or green field MPLS/SR core to an IPv6 transport, Multiprotocol BGP (MP-BGP) now plays an important role in the transition of their Provider (P) core network as well as Provider Edge (PE) Inter-AS peering network from IPv4 to IPv6. Operators must be able to continue to support IPv4 customers when both the Core and Edge networks are IPv6-Only.

This document details an important External BGP (eBGP) PE-PE Inter-AS IPv6-Only peering design that leverages the MP-BGP capability exchange by using IPv6 peering as pure transport, allowing both IPv4 Network Layer Reachability Information (NLRI) and IPv6 Network Layer Reachability Information (NLRI) to be carried over the same (Border Gateway Protocol) BGP TCP session for all Address Family Identifiers (AFI) and Subsequent Address Family Identifiers (SAFI). The design change provides the same Dual Stacking functionality that exists today with separate IPv4 and IPv6 BGP sessions as we have today. With this design change from a control plane perspective a single IPv6-Only peer is required for both IPv4 and IPv6 routing updates and from a data plane forwardindg perspective an IPv6 address need only be configured on the PE to PE Inter-AS peering interface for both IPv4 and IPv6 packet forwarding. This document extends the IPv6-Only PE-CE peering architecture defined in [I-D.ietf-bess-ipv6-only-pe-design] to PE-PE inter-as peering architecture where the 4to6 softwire is now extended to Inter-AS L3 VPN options Option-A, Option-AB and Option-C and now applies to all
AFI/SAFI ubiquitously. As service providers migrate to Segment Routing architecture SR-MPLS and SRv6, VPN overlay exists as well, and thus Inter-AS options Option-A, Option-AB and Option-C are still applicable and thus this extension of IPv6-Only peering architecture extension to Inter-AS peering is very relevant to Segment Routing as well.

Status of This Memo

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1. Introduction

As Enterprises and Service Providers upgrade their brown field or green field MPLS/SR core to an IPv6 transport such as MPLS LDPv6, SR-MPLSv6 or SRv6, Multiprotocol BGP (MP-BGP) now plays an important role in the transition of the Provider (P) core networks and Provider Edge (PE) edge networks from IPv4 to IPv6. Operators have a requirement to support IPv4 customers and must be able to support IPv4 address family and Sub-Address-Family Virtual Private Network (VPN)-IPv4, and Multicast VPN IPv4 customers.

IXP are also facing IPv4 address depletion at their peering points, which are large Layer 2 transit backbones that service providers peer and exchange IPv4 and IPv6 Network Layer Reachability Information (NLRI). Today, these transit exchange points are Dual Stacked. With this IPv6-only BGP peering design, only IPv6 is configured on the PE-PE inter-as peering interface, the Inter-AS Provider Edge (PE) – Provider Edge (PE), the IPv6 BGP peer is now used to carry IPv4 (Network Layer Reachability Information) NLRI over an IPv6 next hop using IPv6 next hop encoding defined in [RFC8950], while continuing
to forward both IPv4 and IPv6 packets. In the framework of this design the ASBRs providing Inter-AS options peering PE to PE extending L3 VPN services is now no longer Dual Stacked.

MP-BGP specifies that the set of usable next-hop address families is determined by the Address Family Identifier (AFI) and the Subsequent Address Family Identifier (SAFI). Historically the AFI/SAFI definitions for the IPv4 address family only have provisions for advertising a Next Hop address that belongs to the IPv4 protocol when advertising IPv4 or VPN-IPv4. [RFC8950] specifies the extensions necessary to allow advertising IPv4 NLRI, Virtual Private Network Unicast (VPN-IPv4) NLRI, Multicast Virtual Private Network (MVPN-IPv4) NLRI with a Next Hop address that belongs to the IPv6 protocol. This comprises of an extended next hop encoding MP-REACH BGP capability exchange to allow the address of the Next Hop for IPv4 NLRI, VPN-IPv4 NLRI and MVPN-IPv4 NLRI to also belong to the IPv6 Protocol. [RFC8950] defines the encoding of the Next Hop to determine which of the protocols the address actually belongs to, and a new BGP Capability allowing MP-BGP Peers to discover dynamically whether they can exchange IPv4 NLRI and VPN-IPv4 NLRI with an IPv6 Next Hop.

The current specification for carrying IPv4 NLRI of a given address family via a Next Hop of a different address family is now defined in [RFC8950], and specifies the extended next hop encoding MP-REACH capability extension necessary to do so. This comprises an extension of the AFI/SAFI definitions to allow the address of the Next Hop for IPv4 NLRI or VPN-IPv4 NLRI to belong to either the IPv4 or the IPv6 protocol, the encoding of the Next Hop information to determine which of the protocols the address belongs to, and a new BGP Capability allowing MP-BGP peers to dynamically discover whether they can exchange IPv4 NLRI and VPN-IPv4 NLRI with an IPv6 Next Hop.

With the new extensions defined in [RFC8950] supporting NLRI and next hop address family mismatch, the BGP peer session can now be treated as a pure TCP transport and carry both IPv4 and IPv6 NLRI at the Provider Edge (PE) - Customer Edge (CE) over a single IPv6 TCP session. This allows for the elimination of dual stack from the PE-PE Inter-AS peering point, and now enable the Inter-AS peering to be IPv6-ONLY. The elimination of IPv4 Inter Provider ASBR tie point, PE-PE Inter-AS peering points translates into OPEX expenditure savings of point-to-point infrastructure links as well as /31 address space savings and administration and network management of both IPv4 and IPv6 BGP peers. This reduction decreases the number of PE-PE Inter-AS options BGP peers by fifty percent, which is a tremendous cost savings for operators.
While the savings exists at the Edge eBGP PE-PE Inter-AS peering, on the core side PE to Route Reflector (RR) peering carrying <AFI/SAFI> IPv4 <1/1>, VPN-IPV4 <1/128>, and Multicasat VPN <1/129>, there is no savings as the Provider (P) Core is IPv6 Only and thus can only have an IPv6 peer and must use [RFC8950] extended next hop encoding to carrying IPv4 NLRI IPV4 <2/1>, VPN-IPV4 <2/128>, and Multicast VPN <2/129> over an IPv6 next hop.

This document details an important External BGP (eBGP) PE-PE Inter-AS IPv6-Only peering design that leverages the MP-BGP capability exchange by using IPv6 peering as pure transport, allowing both IPv4 Network Layer Reachability Information (NLRI) and IPv6 Network Layer Reachability Information (NLRI)to be carried over the same (Border Gateway Protocol) BGP TCP session for all Address Family Identifiers (AFI) and Subsequent Address Family Identifiers(SAFI). The design change provides the same Dual Stacking functionality that exists today with separate IPv4 and IPv6 BGP sessions as we have today. With this design change from a control plane perspective a single IPv6-Only peer is required for both IPv4 and IPv6 routing updates and from a data plane forwarding perspective an IPv6 address need only be configured on the PE to PE Inter-AS peering interface for both IPv4 and IPv6 packet forwarding. This document extends the IPv6-Only PE-CE peering architecture defined in [I-D.ietf-bess-ipv6-only-pe-design] to PE-PE inter-as peering architecture where the 4to6 software is now extended to Inter-AS L3 VPN options Option-A, Option-AB and Option-C and now applies to all AFI/SAFI ubiquitously. As service providers migrate to Segment Routing architecture SR-MPLS and SRv6, VPN overlay exits as well, and thus Inter-AS options Option-A, Option-AB and Option-C are still applicable and thus this extension of IPv6-Only peering architecture extension to Inter-AS peering is relevant to Segment Routing.

This document details an important External BGP (eBGP) PE-PE Inter-AS IPv6-Only peering design that leverages the MP-BGP capability exchange by using IPv6 peering as pure transport, allowing both IPv4 Network Layer Reachability Information (NLRI) and IPv6 Network Layer Reachability Information (NLRI)to be carried over the same (Border Gateway Protocol) BGP TCP session for the following Address Family Identifiers (AFI) and Subsequent Address Family Identifiers(SAFI) to be carried over IPv6-Only Inter-AS peerings described in detail in this document: <AFI/SAFI> IPv4 Unicast <1/1>, IPv4 Multicast <1/2>, VPN-IPV4 <1/128>, Multicasat VPN <1/129>, BGP-LU IPV4 (4PE) <1/4>, BGP-LU IPV4 <1/4>

This document details an important External BGP (eBGP) PE-PE Inter-AS IPv6-Only peering design that leverages the MP-BGP capability exchange by using IPv6 peering as pure transport, allowing both IPv4 Network Layer Reachability Information (NLRI) and IPv6 Network Layer
Reachability Information (NLRI) to be carried over the same (Border Gateway Protocol) BGP TCP session for all remaining Address Family Identifiers (AFI) and Subsequent Address Family Identifiers (SAFI) below as well that can be carried over IPv6-Only Inter-AS peerings:


2. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

3. Terminology

Terminology used in defining the IPv6-Only Edge specification.

AFBR: Address Family Border Router Provider Edge (PE).

Edge: PE-CE Edge Network Provider Edge - Customer Edge

Core: P Core Network Provider (P)

4to6 Softwire: IPv4 edge over an IPv6-Only core

6to4 Softwire: IPv6 edge over an IPv4-Only core

E2E: End to End

4. IPv6-Only Edge Peering Architecture
4.1. Problem Statement

This specification addresses a real issue that has been discussed at many operator groups around the world related to IXP major peering points where hundreds of AS’s have both IPv4 and IPv6 dual stacked peering. IPv4 address depletion have been a major issue issue for many years now. Operators around the world are clamoring for a solution that can help solve issues related to IPv4 address depletion at these large IXP peering points. With this solution IXPs as well as all infrastructure networks such as Core networks, DC networks, Access networks as well as any PE-CE public or private network can now utilize this IPv6-Only Edge solution and reap the benefits immediately on IPv4 address space saving.

![Figure 1: Problem Statement - IXP Dual Stack Peering](image-url)
4.2. IPv6-Only PE-CE Design Solution

The IPv6-Only Edge design solution provides a means of E2E single protocol design solution extension of [RFC5565] Softwire Mesh framework from the PE-CE Edge to the Core from ingress so egress through the entire operators domain. This solution eliminates all IPv4 addressing from end to end while still providing the same Dual Stack functionality of IPv4 and IPv6 packet forwarding from a data plane perspective by leveraging the [RFC8950] extended next hop encoding so that IPv4 NLRI can be advertised over a single IPv6 pure transport TCP session. This IPv6-Only E2E architecture eliminates all IPv4 peering and IPv4 addressing E2E from the ingress CE to ingress PE to egress PE to egress CE and all hops along the operator E2E path.

Solution applicable to any Edge peering scenario - IXP, Core, DC, Access, etc

Figure 3: IPv6-Only Solution Applicability
4.3. IPv6-Only Edge Peering Design

4.3.1. IPv6-Only Edge Peering Packet Walk

The IPv6-Only Edge Peering design utilizes two key E2E Softwire Mesh Framework scenario’s, 4to6 softwire and 6to4 softwire. The Softwire mesh framework concept is based on the overlay and underlay MPLS or SR based technology framework, where the underlay is the transport layer and the overlay is a Virtual Private Network (VPN) layer, and is the the tunneled virtualization layer containing the customer payload. The concept of a 6to4 Softwire is based on transmission of IPv6 packets at the edge of the network by tunneling the IPv6 packets over an IPv4-Only Core. The concept of a 4to6 Softwire is also based on transmission of IPv4 packets at the edge of the network by tunneling the IPv4 packets over an IPv6-Only Core.

This document describes End to End (E2E) test scenarios that follow a packet flow from IPv6-Only attachment circuit from ingress PE-CE to egress PE-CE tracing the routing protocol control plane and data plane forwarding of IPv4 packets in a 4to6 softwire or 6to4 softwire within the IPv4-Only or IPv6-Only Core network. In both scenario we are focusing on IPv4 packets and the control plane and data plane forwarding aspects of IPv4 packets from the PE-CE Edge network over an IPv6-Only P (Provider) core network or IPv4-Only P (Provider) core network. With this IPv6-Only Edge peering design, the Softwire Mesh Framework is not extended beyond the Provider Edge (PE) and continues to terminate on the PE router.

4.3.2. 6to4 Softwire IPv4-Only Core packet walk

6to4 softwire where IPv6-Edge eBGP IPv6 peering where IPv4 packets at network Edge traverse a IPv4-Only Core

In the scenario where IPv4 packets originating from a PE-CE edge are tunneled over an MPLS or Segment Routing IPv4 underlay core network, the PE and CE only have an IPv6 address configured on the interface. In this scenario the IPv4 packets that ingress the CE from within the CE AS are over an IPv6-Only interface and are forwarded to an IPv4 NLRI destination prefix learned from the Pure Transport Single IPv6 BGP Peer. In the IPv6-Only Edge peering architecture the PE is IPv6-Only as all PE-CE interfaces are IPv6-Only. However, on the CE, the PE-CE interface is the only interface that is IPv6-Only and all other interfaces may or may not be IPv6-Only. Following the data plane packet flow, IPv4 packets are forwarded from the ingress CE to the IPv6-Only ingress PE where the VPN label imposition push per prefix, per-vrf, per-CE occurs and the labeled packet is forwarded.
over a 6to4 softwire IPv4-Only core, to the egress PE where the VPN label disposition pop occurs and the native IPv4 packet is forwarded to the egress CE. In the reverse direction IPv4 packets are forwarded from the egress CE to egress PE where the VPN label imposition per prefix, per-vrf, per-CE push occurs and the labeled packet is forwarded back over the 6to4 softwire IPv4-Only core, to the ingress PE where the VPN label disposition pop occurs and the native IPv4 packet is forwarded to the ingress CE. The functionality of the IPv4 forwarding plane in this scenario is identical from a data plane forwarding perspective to Dual Stack IPv4 forwarding scenario.
Figure 5: 6to4 Softwire - IPv6 Edge over an IPv4-Only Core

4.3.3. 4to6 Softwire IPv6-Only Core packet walk

4to6 softwire where IPv6-Edge eBGP IPv6 peering where IPv4 packets at network Edge traverse a IPv6-Only Core
In the scenario where IPv4 packets originating from a PE-CE edge are tunneled over an MPLS or Segment Routing IPv4 underlay core network, the PE and CE only have an IPv6 address configured on the interface. In this scenario the IPv4 packets that ingress the CE from within the CE AS are over an IPv6-Only interface and are forwarded to an IPv4 NLRI destination prefix learned from the Pure Transport Single IPv6 BGP Peer. In the IPv6-Only Edge peering architecture the PE is IPv6-Only as all PE-CE interfaces are IPv6-Only. However, on the CE, the PE-CE interface is the only interface that is IPv6-Only and all other interfaces may or may not be IPv6-Only. Following the data plane packet flow, IPv4 packets are forwarded from the ingress CE to the IPv6-Only ingress PE where the VPN label imposition push per prefix, per-vrf, per-CE occurs and the labeled packet is forwarded over a 4to6 softwire IPv6-Only core, to the egress PE where the VPN label disposition pop occurs and the native IPv4 packet is forwarded to the egress CE. In the reverse direction IPv4 packets are forwarded from the egress CE to egress PE where the VPN label imposition per prefix, per-vrf, per-CE push occurs and the labeled packet is forwarded back over the 4to6 softwire IPv6-Only core, to the ingress PE where the VPN label disposition pop occurs and the native IPv4 packet is forwarded to the ingress CE. The functionality of the IPv4 forwarding plane in this scenario is identical from a data plane forwarding perspective to Dual Stack IPv4 forwarding scenario.
Figure 6: 4to6 Softwire - IPv4 Edge over an IPv6-Only Core

4.4. RFC5549 and RFC8950 Applicability
4.4.1. IPv6-Only Edge Peering design next-hop encoding

This section describes [RFC8950] next hop encoding updates to [RFC5549] applicability to this specification. IPv6-only eBGP Edge PE-CE peering to carry IPv4 Unicast NLRI <AFI/SAFI> IPv4 <1/1> over an IPv6 next hop BGP capability extended hop encoding IANA capability codepoint value 5 defined is applicable to both [RFC5549] and [RFC8950] as IPv4 Unicast NLRI <AFI/SAFI> IPv4 <1/1> does not change in the RFC updates.

IPv4 packets over an IPv6-Only core 4to6 Softwire E2E packet flow is part of the IPv6-Only design vendor interoperability test cases and in that respect is applicable as [RFC8950] updates [RFC5549] for <AFI/SAFI> VPN-IPV4 <1/128>, and Multicasat VPN <1/129>.

4.4.2. RFC8950 updates to RFC5549 applicability

This section describes the [RFC8950] next hop encoding updates to [RFC5549]

In [RFC5549] when AFI/SAFI 1/128 is used, the next-hop address is encoded as an IPv6 address with a length of 16 or 32 bytes. This document modifies how the next-hop address is encoded to accommodate all existing implementations and bring consistency with VPNv4oIPv4 and VPNv6oIPv6. The next-hop address is now encoded as a VPN-IPv6 address with a length of 24 or 48 bytes [RFC8950] (see Sections 3 and 6.2 of this document). This change addresses Erratum ID 5253 (Err5253). As all known and deployed implementations are interoperable today and use the new proposed encoding, the change does not break existing interoperability. Updates to [RFC8950] is applicable to the IPv6-Only PE-CE edge design for the IPv6 next hop encoding E2E test case of IPv4 packets over and IPv6-Only core 4to6 Softwire. In this test case IPv4 Unicast NLRI <AFI/SAFI> IPv4 <1/1> is advertised over the PE to RR core peering 4to6 softwire in <AFI/SAFI> VPN-IPV4 <1/128>. In this test case label allocation mode comes into play which is discussed in section 8.9.

[RFC5549] next hop encoding of MP_REACH_NLRI with:

* NLRI= NLRI as per current AFI/SAFI definition

Advertising with [RFC4760] MP_REACH_NLRI with:

* AFI = 1
* SAFI = 128 or 129
* Length of Next Hop Address = 16 or 32
* NLRI= NLRI as per current AFI/SAFI definition

[RFC8950] next hop encoding of MP_REACH_NLRI with:

* NLRI= NLRI as per current AFI/SAFI definition

Advertising with [RFC4760] MP_REACH_NLRI with:

* AFI = 1
* SAFI = 128 or 129
* Length of Next Hop Address = 24 or 48
* Next Hop Address = VPN-IPv6 address of next hop with an 8-octet RD set to zero (potentially followed by the link-local VPN-IPv6 address of the next hop with an 8-octet RD is set to zero).

* NLRI= NLRI as per current AFI/SAFI definition

5. IPv6-Only PE Design Edge E2E Design for all AFI/SAFI

Proof of concept interoperability testing of the 4 test cases bet

5.1. Design Solution-1 E2E IPv6-Only PE-CE, Global Table over IPv4-Only Core(6PE), 6to4 softwire

---

![Diagram](image.png)

Figure 7: Design Solution-1 E2E IPv6-Only PE-CE, Global Table over IPv4-Only Core (6PE)
Cisco, Juniper, Arista, Nokia, Huawei code and platform and test results.

Cisco: Edge Router- XR ASR 9910 IOS XR 7.4.1, Core Router- NCS 6000 7.2.2, CRS-X 6.7.4

Juniper: Edge Router- MX platform MX480, MX960, Core Router- PTX Platform PTX5000, PTC10K8 (JUNOS and EVO) Release 20.4R2

Tested v4 edge over v6 core in a virtual setup using vMX platform and 20.4R2 and LDPv6 as underlay, but there were some data plane forwarding issues. Tested same setup on latest release 21.4 and it worked. Investigating what the minimum version is for this setup to work.

Arista:

Nokia: Edge and Core-7750 Service Router, Release R21

Huawei: Edge and Core-VRPv8, Release VRP-V800R020C10

5.2. Design Solution-2 E2E IPv6-Only PE-CE, VPN over IPv4-Only Core, 6to4 Softwire

![Diagram of Design Solution-2 E2E IPv6-Only PE-CE, VPN over IPv4-Only Core](image)

5.3. Design Solution-3 E2E IPv6-Only PE-CE, Global Table over IPv6-Only Core (4PE), 4to6 Softwire

Figure 9: Design Solution-3 E2E IPv6-Only PE-CE, Global Table over IPv6-Only Core (4PE)

5.4. Design Solution-4 E2E IPv6-Only PE-CE, VPN over IPv6-Only Core, 4to6 Softwire

Figure 10: Design Solution-4 E2E IPv6-Only PE-CE, VPN over IPv6-Only Core

5.5. Design Solution-5 E2E Inter-AS Option B and AB, 4to6 Softwire
5.6. IPv6-Only PE-CE Operational Considerations Testing

Ping CE to PE when destination prefix is withdrawn
Traceroute CE to PE and test all ICMPv4 and ICMPv6 type codes

6. IPv6-Only PE ALL AFI/SFI Operational Considerations

With a single IPv6 Peer carrying both IPv4 and IPv6 NLRI there are some operational considerations in terms of what changes and what does not change.

What does not change with a single IPv6 transport peer carrying IPv4 NLRI and IPv6 NLRI below:

Routing Policy configuration is still separate for IPv4 and IPv6 configured by capability as previously.
Layer 1, Layer 2 issues such as one-way fiber or fiber cut will impact both IPv4 and IPv6 as previously.

If the interface is in the Admin Down state, the IPv6 peer would go down, and IPv4 NLRI and IPv6 NLRI would be withdrawn as previously.

Changes resulting from a single IPv6 transport peer carrying IPv4 NLRI and IPv6 NLRI below:

Physical interface is no longer dual stacked.

Any change in IPv6 address or DAD state will impact both IPv4 and IPv6 NLRI exchange.

Single BFD session for both IPv4 and IPv6 NLRI fate sharing as the session is now tied to the transport, which now is only IPv6 address family.

Both IPv4 and IPv6 peer now exists under the IPv6 address family configuration.

Fate sharing of IPv4 and IPv6 address family from a logical perspective now carried over a single physical IPv6 peer.

From an operations perspective, prior to elimination of IPv4 peers, an audit is recommended to identify and IPv4 and IPv6 peering incongruencies that may exist and to rectify them. No operational impacts or issues are expected with this change.

With MPLS VPN overlay, per-CE next-hop label allocation mode where both IPv4 and IPv6 prefixes have the same label in no table lookup pop-n-forward mode should be taken into consideration.

7. Vendor Implementations and Operator Deployments

Vendor implementations are with Cisco, Juniper, Nokia, Arista and Huawei

8. IANA Considerations

There are not any IANA considerations.
9. Security Considerations

The extensions defined in this document allow BGP to propagate reachability information about IPv4 prefixes over an MPLS or SR IPv6-Only core network. As such, no new security issues are raised beyond those that already exist in BGP-4 and the use of MP-BGP for IPv6. Both IPv4 and IPv6 peers exist under the IPv6 address family configuration. The security features of BGP and corresponding security policy defined in the ISP domain are applicable. For the inter-AS distribution of IPv6 routes according to case (a) of Section 4 of this document, no new security issues are raised beyond those that already exist in the use of eBGP for IPv6 [RFC2545].

10. Acknowledgments

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12. References

12.1. Normative References

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12.2. Informative References


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BGP Extensions for the Mobile User Plane (MUP) SAFI
draft-mpmz-bess-mup-safi-00.txt

Abstract

This document defines a new SAFI known as a BGP Mobile User Plane (BGP-MUP) SAFI to support MUP Extensions and a extended community for BGP. This document also provides BGP signaling and procedures for the new SAFI to convert mobile session information into appropriate IP forwarding information. These extensions can be used by operators between MUP PE, MUP GW and MUP Controller for integrating mobile user plane into BGP MUP network using the IP based routing.

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1. Introduction

[I-D.mhkk-dmm-srv6mup-architecture] defines the Segment Routing IPv6 Mobile User Plane (SRv6 MUP) architecture for Distributed Mobility Management. As part of the architecture, the document defines a new
SRv6 segment type called as a MUP Segment, new routing information that can carried within BGP, and 3 new network nodes; MUP PE, MUP GW and a MUP Controller.

This document defines a new SAFI known as a BGP Mobile User Plane (BGP-MUP) SAFI to support MUP Extensions for BGP. This draft also provides BGP signaling and procedures for the new SAFI to convert mobile session information into appropriate IP routing information. These extensions can be used by operators between the MUP PE, MUP GW and MUP Controller for integrating mobile user plane into BGP MUP network using the IP based routing. These extensions also works with routing instances accommodating two new well-known segment types known as Interwork and Direct [I-D.mhkk-dmm-srv6mup-architecture]. Finally, the BGP encoding and procedures defined in this document that uses SRv6 as the forwarding fabric follows the SRv6 MUP architecture defined in [I-D.mhkk-dmm-srv6mup-architecture]. The BGP extensions to build networks that use forwarding mechanisms other than SRv6 (SRv6 MUP) are outside the scope of this document.

1.1. Requirements Notation

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] when, and only when, they appear in all capitals, as shown here.

2. Terminology

MUP: Mobile User Plane

MUP Segment: Representation of mobile user plane segment

MUP PE: Provider Edge node in a MUP network

MUP GW: Gateway node to interwork with another mobile user plane networks

MUP Controller: Controller node for a MUP network

UE: User Equipment, as per [TS.23501]

gNodeB: 3GPP-compliant implementation of 5G-NR base station, as per [TS.23501]

UPF: 3GPP-compliant implementation of User Plane Function, as per [TS.23501]
3. BGP MUP Extensions

3.1. BGP MUP SAFI

This draft defines a new BGP SAFI known as a BGP-MUP SAFI. The value of this SAFI is to be assigned by IANA from the Subsequent Address Family Identifiers (SAFI) registry.

This document also defines a new BGP NLRI known as the BGP-MUP NLRI. The following is the format of the BGP-MUP NLRI:

```
+-------------------------------+
|    Architecture Type (1 octet) |
+-------------------------------+
|       Route Type (2 octets)    |
+-------------------------------+
|         Length (1 octet)       |
+-------------------------------+
|  Route Type specific (variable) |
+-------------------------------+
```

The Architecture Type field defines the encoding of the rest of BGP-MUP NLRI for a given Mobile User Plane architecture. This draft defines the following architecture type for BGP-MUP NLRI:

+ 1 - 3gpp-5g;

The Route Type field defines the encoding of the rest of BGP-MUP NLRI (Route Type specific BGP-MUP NLRI) for a given architecture type.

The Length field indicates the length in octets of the Route Type specific field of the BGP-MUP NLRI.

This document defines the following Route Types for BGP-MUP NLRI.

+ 1 - Interwork Segment Discovery route;
+ 2 - Direct Segment Discovery route;
+ 3 - Type 1 Session Transformed (ST) route;
+ 4 - Type 2 Session Transformed (ST) route;

These Route Types are applicable for the 3gpp-5G architecture type as per [I-D.mhkk-dmm-srv6mup-architecture]. Other new architectures can share them if it is applicable as well.
The BGP-MUP NLRI is carried in BGP [RFC4271] using BGP Multiprotocol Extensions [RFC4760] with an Address Family Identifier (AFI) of 1 or 2 and a Subsequent AFI (SAFI) of BGP-MUP. The NLRI field in the MP_REACH_NLRI/MP_UNREACH_NLRI attribute contains the BGP-MUP NLRI (encoded as specified above). The value of the AFI field in the MP_REACH_NLRI/MP_UNREACH_NLRI attribute that carries the BGP-MUP NLRI determines whether the addresses carried in the routes are IPv4 or IPv6 addresses (AFI 1 indicates IPv4 addresses, AFI 2 indicates IPv6 addresses).

In order for two BGP speakers to exchange BGP-MUP NLRIs, they must use a BGP Capabilities Advertisement to ensure that they both are capable of properly processing such an NLRI. This is done as specified in [RFC4760], by using capability code 1 (multiprotocol BGP) with an AFI of 1 or 2 and an SAFI of BGP-MUP.

This document defines 4 Route Types for 3gpp-5G architecture type. Any other Route Types MUST be silently ignored upon a receipt if a BGP speaker supports only 3gpp-5G architecture type. An implementation MAY log an error when such Route Types are ignored. An implementation MAY consider retrieving any discarded Route Types by simply resetting the session or issuing a Route-REFRESH message [RFC2918] if the Route Refresh Capability is successfully negotiated.

The following sections describes the format of the Route Type specific BGP-MUP NLRI for various Route Types defined in this document.

3.1.1. BGP Interwork Segment Discovery route

A BGP Interwork Segment Discovery route Type specific BGP-MUP NLRI consist of the following:

```
+-----------------------------------+
|           RD  (8 octets)          |
+-----------------------------------+
|       Prefix Length (1 octet)     |
+-----------------------------------+
|        Prefix (variable)          |
+-----------------------------------+
```

The Interwork Segment Discovery route Type NLRI consist of RD which is encoded as described in [RFC4364]. It also has a prefix associated with Interwork segment connected locally. For the purpose of BGP route key processing, only the RD, Prefix Length and Prefix are considered to be part of the prefix in the NLRI.
In 3GPP 5G specific case, a prefix used for a N3 interface of the gNodeB connected locally MAY be used as this prefix. The prefix length of one octet indicating length of the prefix. If the AFI is IPv4, then the maximum value of the Prefix Length is 32 bits otherwise it is considered as a malformed NLRI. If the AFI is IPv6, then the maximum value of of the Prefix length is 128 bits otherwise it is considered as a malformed NLRI. A BGP speaker MUST handle such a malformed NLRI as a "Treat-as-withdraw" [RFC7606]. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.

3.1.2. BGP Direct Segment Discovery route

A BGP Direct Segment Discovery route Type specific BGP-MUP NLRI consist of the following:

```
+-----------------------------------+
|           RD  (8 octets)          |
+-----------------------------------+
|        Address (4 or 16 octets)   |
+-----------------------------------+
```

The Direct Segment Discovery route Type NLRI consist of RD which is encoded as described in [RFC4364]. It also has an Address of originating BGP speaker. For the purpose of BGP route key processing, only the RD and Address are considered to be part of the prefix in the NLRI.

If the AFI is IPv4 then the address length is 4 octets otherwise it is considered as a malformed NLRI. If the AFI is IPv6 then the address length is 16 octets otherwise it is considered as a malformed NLRI. A BGP speaker MUST handle such a malformed NLRI as a "Treat-as-withdraw" [RFC7606]. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.

3.1.3. BGP Type 1 Session Transformed (ST) Route

A BGP Type 1 ST Route Type specific BGP-MUP NLRI consist of the following:
The Type 1 ST Route Type NLRI consist of RD which is encoded as described in [RFC4364]. It also has Prefix length of one octet indicating length of the Prefix. For the purpose of BGP route key processing, only the RD, Prefix Length and Prefix are considered to be part of the prefix in the NLRI.

In 3GPP 5G specific case, Prefix is the prefix allocated to a UE. If the AFI is IPv4, then the maximum value of the Prefix Length field is 32. If the AFI is IPv6, then the maximum value of the Prefix Length field is 128. Any other length field is considered a malformed NLRI. A BGP speaker MUST handle such a malformed NLRI as a "Treat-as-withdraw" [RFC7606]. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.

The architecture specific encoding values will follow after the variable length Prefix.

3.1.3.1. 3gpp-5g Specific BGP Type 1 ST Route

A BGP 3gpp-5g Type 1 ST Route Type specific BGP-MUP NLRI consist of the following:

The TEID has a fixed length of 4 octets. The TEID value of 0 is considered as an invalid and a malformed TEID. A BGP speaker MUST handle such a malformed NLRI as a "Treat-as-withdraw" [RFC7606]. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.
The QFI has a fixed length of 1 octet.

The Endpoint Address Length has a fixed length of 1 octet. Endpoint Address field contains of an IPv4 address, then the value of the Endpoint Address Length field is 32. If the Endpoint Address field contains of an IPv6 Address, then the value of the Endpoint Address Length field is 128. Any other value is considered as an invalid and a malformed Endpoint Address. A BGP speaker MUST handle such a malformed NLRI as a "Treat-as-withdraw" [RFC7606]. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.

The NLRI architecture field MUST be encoded as shown above if a BGP speaker receives 3gpp-5g specific BGP Type 1 ST route. Otherwise the NLRI is considered as a malformed. A BGP speaker MUST handle such a malformed NLRI as a "Treat-as-withdraw" [RFC7606]. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.

3.1.4. BGP Type 2 Session Transformed (ST) Route

A BGP Type 2 ST Route Type specific BGP-MUP NLRI consist of the following:

```
+----------------+                        +----------------+                        +----------------+                        +----------------+                        +----------------+
| RD (8 octets)  |                        | Endpoint Length (1 octet) |                        | Endpoint Address (variable) |                        | Architecture specific Endpoint Identifier (variable) |
+----------------+                        +----------------+                        +----------------+                        +----------------+                        +----------------+
```

The Type 2 ST Route Type NLRI consist of RD which is encoded as described in [RFC4364]. It also has Endpoint length of one octet indicating length of the Endpoint Address and the Architecture specific Endpoint Identifier as per [I-D.mhkk-dmm-srv6mup-architecture] defines aggregation capability by the Type2 ST Route. If the AFI is IPv4 and the Endpoint Length is longer than 32 then the Architecture specific endpoint Identifier field exists with the IPv4 Endpoint Address. If the AFI is IPv6 and the Endpoint Length is longer than 128 then the Architecture specific endpoint Identifier field exists with then the IPv6 Endpoint Address. For the purpose of BGP route key processing, only the RD, Endpoint Address and Architecture specific Endpoint Identifier are considered to be part of the prefix in the NLRI.
In 3GPP 5G specific case, the Endpoint Address is a N3 interface address of the UPF. If the AFI is IPv4, then the maximum Endpoint length is 64 otherwise it is considered as a malformed NLRI. If the AFI is IPv6, then the maximum Endpoint length is 160 otherwise it is considered as a malformed NLRI. A BGP speaker MUST handle such a malformed NLRI as a "Treat-as-withdraw" [RFC7606]. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.

The architecture specific encoding values will follow after the variable length Prefix.

3.1.4.1. 3gpp-5g Specific BGP Type 2 ST Route

A BGP 3gpp-5g Type 2 ST Route Type specific BGP-MUP NLRI consist of the following:

```
+-----------------------------------+
|          TEID (0-4 octets)        |
+-----------------------------------+
```

The maximum length of TEID is 4 octets. The TEID value of 0 is considered as an invalid and a malformed TEID. A BGP speaker MUST handle such a malformed NLRI as a "Treat-as-withdraw" [RFC7606]. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.

The NLRI architecture field MUST be encoded as shown above if a BGP speaker receives 3gpp-5g specific BGP Type 2 ST route. Otherwise the NLRI is considered as a malformed. A BGP speaker MUST handle such a malformed NLRI as a "Treat-as-withdraw" [RFC7606]. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.

3.2. BGP MUP Extended Community

This document defines a new BGP Extended community known as BGP MUP Extended community as per [I-D.mhkk-dmm-srv6mup-architecture]. This is a BGP MUP specific Extended Community, of an extended type, and is transitive across AS boundaries [RFC4360].

The Type value of this Extended community is set to MUP type, to be assigned by IANA from BGP Extended community transitive registry. The Sub-Type value is set to Direct-Type Segment Identifier type, to be assigned by IANA from BGP Extended community transitive registry. The value field of the community is set to the 6 bytes of configurable segment identifier value.
The usage of this Extended community is described in the Section 3.3.12 and Section 3.3.10

3.3. Operation

BGP speakers acting as a MUP PE, MUP GW and a MUP Controller MUST establish a BGP session to exchange BGP-MUP NLRIs for both, IPv4 and IPv6 AFIs. Once the session is established successfully, BGP speakers can exchange the Discovery routes as well as Session Transformed routes. This information is specific to a given routing instance. BGP-MUP SAFI is expected to work with routing instances accommodating MUP segments. In 3GPP 5G specific case, the routing instances are depicted as N3RAN and N6DN instances defined in [I-D.mhkk-dmm-srv6mup-architecture]. The subsequent sections describes procedures of generating and processing of each route types.

3.3.1. Generation of the Interwork Segment Discovery route

The Interwork Segment Discovery route is generated by the MUP GW when a routing instance accommodates an Interwork type MUP Segment, e.g., N3 interfaces or routes on RAN side in 3GPP 5G specific case. It generates route per each N3RAN IP prefix and stores the route in the routing instance of N3RAN. The IP prefix MAY include a gNodeB address which is connecting to the MUP GW. The BGP API for BGP MP_REACH_NLRI attribute to carry the Discovery route is decided based on the API of the prefix.

When advertising the Interwork Segment Discovery route, a MUP GW MUST attach the export BGP Route Target Extended Community of the associated routing instance.

When advertising the Interwork Segment Discovery route, a MUP GW MUST use the IPv6 address of the MUP GW as the nexthop address in the MP_REACH_NLRI attribute.

The Interwork Segment Discovery route update MUST have a prefix SID attribute which the SID consists of the MUP GW locator followed by a function. In 3GPP 5G specific case, if the BGP API is IPv4, the function MUST be GTP4.E [I-D.ietf-dmm-srv6-mobile-uplane], or MUST be GTP6.E [I-D.ietf-dmm-srv6-mobile-uplane] if the BGP API is IPv6.

3.3.2. Withdrawal of the Interwork Segment Discovery route

The Interwork Segment Discovery route is withdrawn by the MUP GW when it detects that the MUP Segment no longer present for the N3RAN. The BGP API for BGP MP_UNREACH_NLRI attribute to carry the Discovery route is decided based on the API of the prefix.
When withdrawing the Interwork Segment Discovery route, a MUP GW MUST attach the export BGP Route Target Extended Community of the associated routing instance.

### 3.3.3. Processing of the Interwork Segment Discovery routes

Both, the MUP GW and the MUP PE MAY receive the Discovery Interwork routes from other MUP GWs in the BGP MUP network. A BGP speaker acting as a MUP PE or a MUP GW MAY keep the received MUP Interwork Segment Discovery routes advertised from other MUP GWs. The receiving BGP speaker will perform the importing of the received MUP Interwork Segment Discovery routes in the configured routing instance based on the Route Target extended communities. The IP prefixes for the received segments are imported into the configured routing instance table. Thereby the receiving BGP speaker can provide network connectivity between the nodes that exist in the segments. A BGP speaker MAY discard the received Interwork Segment Discovery route if the speaker fails to import the route based on the Route Target extended communities.

The BGP speaker receiving the Interwork Segment Discovery routes SHOULD ignore the nexthop in the MP_REACH_NLRI attribute. However, the receiving BGP speaker MUST ensure that the value of Address filed in the NLRI is an address of the originator of the locator value in the prefix SID attribute. The originator of the locator value can be resolved from the IPv6 IGP table. If the result of the match is not identical then the receiving BGP speaker MUST consider it as a malformed NLRI and the "Treat-as-withdraw procedure of [RFC7606] is applied. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.

When a BGP speaker receives a MP_REACH_NLRI attribute update message with a Discovery Internetwork NLRI without a prefix SID attribute, than it MUST be treated as if it contained a malformed prefix SID attribute and the "Treat-as-withdraw procedure of [RFC7606] is applied. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.

When a BGP speaker receives an MP_UNREACH_NLRI attribute update message it MUST delete the withdrawn Interwork Segment Discovery route from the routing instance table where it was created.

### 3.3.4. Generation of the Direct Segment Discovery route

The Direct Segment Discovery route is generated by the MUP PE when a routing instance accommodates a Direct type MUP Segment, e.g., N6 interface or routes on DN side in 3GPP 5G specific case. It generates the Direct Segment Discovery route per each routing...
instance for the MUP Segment. The address in the BGP-MUP NLRI MUST be a unique MUP PE identifier. The BGP AFI for BGP MP_REACH_NLRI attribute to carry the Direct Segment Discovery route is decided based on the AFI of the MUP PE identifier [I-D.mhkk-dmm-srv6mup-architecture].

When announcing the Direct Segment Discovery route, a MUP PE MUST attach a BGP MUP Extended community of the associated routing instance. The sub-type of the Extended community is Direct-Type Segment Identifier.

When advertising the Direct Segment Discovery route, a MUP PE MUST use the IPv6 address of the MUP PE as the nexthop address in the MP_REACH_NLRI attribute.

The Direct Segment Discovery route update MUST have a prefix SID attribute which the SID consists of the MUP PE locator followed by a function. The function MAY be End.DT4/6 or End.DX4/6.

3.3.5. Withdrawal of the Direct Segment Discovery route

The Direct Segment Discovery route is withdrawn by the MUP PE when it detects that the MUP Segment no longer present for the routing instance. The BGP AFI for BGP MP_UNREACH_NLRI attribute to carry the Discovery route is decided based on the AFI of the MUP PE identifier.

When withdrawing the Direct Segment Discovery route, a BGP speaker MUST attach a BGP MUP Extended community of the associated routing instance.

3.3.6. Processing of the Direct Segment Discovery routes

Both, the MUP GW and the MUP PE MAY receive the Discovery Direct routes from other MUP PEs in the BGP MUP network. A BGP speaker acting as a MUP PE or a MUP GW MAY keep the received MUP Direct Segment Discovery routes advertised from other MUP PEs. The receiving BGP speaker will perform the importing of the received MUP Direct Segment Discovery routes in the configured routing instance based on the Route Target extended communities. The IP prefixes for the received segments are imported into the configured routing instance table. Thereby the receiving BGP speaker can provide network connectivity between the nodes that exist in the segments. A BGP speaker MAY discard the received Direct Segment Discovery route if the speaker fails to import the route based on the Route Target extended communities.

The BGP speaker receiving the Direct Segment Discovery routes SHOULD ignore the nexthop in the MP_REACH_NLRI attribute. However, the
receiving BGP speaker MUST ensure that the received nexthop value in the MP_REACH_NLRI attribute is identical to the originator of the locator value in the prefix SID attribute. The originator of the locator value can be resolved from the IPv6 IGP table. If the result of the match is not identical then the receiving BGP speaker MUST consider it as a malformed NLRI and the "Treat-as-withdraw procedure of [RFC7606] is applied. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.

When a BGP speaker receives a MP_REACH_NLRI attribute update message with a Direct Segment Discovery route without a prefix SID attribute, than it MUST be treated as if it contained a malformed prefix SID attribute and the "Treat-as-withdraw procedure of [RFC7606] is applied. A BGP speaker MUST skip such NLRIs and continue processing of rest of the Update message.

When a BGP speaker receives an MP_UNREACH_NLRI attribute update message it MUST delete the withdrawn Direct Segment Discovery route from the routing instance table where it was created.

3.3.7. Generation of the Type 1 ST Route

A BGP speaker acting as a MUP controller generates Type 1 ST route from corresponding session information through a northbound API as per [I-D.mhkk-dmm-srv6mup-architecture]. The northbound API definition for ST1 route creation is out of scope of this document.

In 3GPP 5G specific case, to compose a Type 1 ST route the MUP controller acquires UE address or prefix, Tunnel Endpoint address of GTP [TS.29281] tunnel, TEID, and QFI for access side as input parameters from the northbound API. The MUP controller decides the RD of the Type 1 ST route based on operator policy.

When advertising the Type 1 ST route, the MUP controller SHOULD attach a Route Target Extended community which the MUP PEs are importing into the routing instance for the corresponding Direct segment.

The MUP controller MUST set the nexthop of the route to the address of the MUP Controller.

The MUP controller MUST announce this route using a AFI of the route and the SAFI of BGP-MUP to all other BGP speakers within the SRv6 domain.
3.3.8. Withdrawing of the Type 1 ST Route

A BGP speaker acting as a MUP controller withdraws Type 1 ST route based on deletion of corresponding session information through a northbound API as per [I-D.mhkk-dmm-srv6mup-architecture]. The northbound API definition for ST1 route withdrawal is out of scope of this document.

In 3GPP 5G specific case, to withdraw a Type 1 ST route the MUP controller acquires the UE address or prefix, Tunnel Endpoint address of GTP[TS.29281] tunnel, TEID, and QFI for access side as input parameters from the northbound API. The MUP controller MUST advertise the withdrawals of the Type 1 ST route.

When withdrawing the Type 1 ST route, the MUP controller SHOULD attach the Route Target Extended community which the MUP PEs are importing into the routing instance accommodating the corresponding Direct segment to the Route Target Extended community.

3.3.9. Processing of the Type 1 ST Route

Both, the MUP GW and the MUP PE MAY receive the Type 1 ST routes from the MUP Controller in the BGP MUP network. A BGP speaker acting as a MUP PE or a MUP GW MAY keep the received MUP Type 1 ST routes advertised from the MUP Controller. The receiving BGP speaker will perform the importing of the received MUP Type 1 ST routes in the configured N6DN routing instance based on the Route Target extended communities. A BGP speaker MAY discard the received Type 1 ST route if the speaker fails to import the route based on the Route Target extended communities.

In case of a BGP speaker receiving a Type 1 ST routes is a MUP PE, the MUP PE SHOULD use the received Tunnel Endpoint Address in this NLRI as a key to lookup the associated Interwork Segment Discovery route and extract the locator and the function in the prefix SID attribute of the Interwork route.

In 3GPP 5G specific case, the MUP PE extracts TEID, QFI and Tunnel Endpoint address from the NLRI. Then the MUP PE SHOULD generate the forwarding SID for GTP4/6.E based on the procedures mentioned in the [I-D.ietf-dmm-srv6-mobile-uplane]. If the MUP PE cannot generate the prefix SID, then it SHOULD mark the received Type 1 ST route as an invalid route. The MUP PE MAY hold such an invalid route until the route as a valid route upon successful generation of prefix SID.

The MUP PE receiving Type 1 ST routes SHOULD ignore the received nexthop in the MP_REACH_NLRI attribute.
The MUP PE receiving Type 1 ST routes in MP_UNREACH_NLRI attribute MUST delete all the routes from the associated routing instance.

### 3.3.10. Generation of the Type 2 ST Route

A BGP speaker acting as a MUP controller generates Type 2 ST route from corresponding session information through a northbound API, or pre-defined configuration as per [I-D.mhkk-dmm-srv6mup-architecture]. The northbound API definition for ST2 route creation is out of scope of this document.

In 3GPP 5G specific case, to compose a Type 2 ST route the MUP controller acquires the Endpoint consists of Endpoint address of GTP [TS.29281] tunnel and TEID for core side with the effective length of the Endpoint as input parameters. The MUP controller decides the RD of the Type 2 ST route based on operator policy.

When advertising the Type 2 ST route, the MUP controller SHOULD attach a BGP MUP Extended community corresponding to the Direct segment. The sub-type of the Extended community is Direct-Type Segment Identifier. This Segment Identifier is generated from the information received through a northbound API, or a pre-defined configuration as per [I-D.mhkk-dmm-srv6mup-architecture]. The northbound API definition for receiving this information is out of scope of this document.

The MUP controller MUST also attach a Route Target Extended community of the routing instances in the MUP GW accommodating the corresponding Interwork segment.

The MUP controller MUST set the next hop of the route to the address of the MUP Controller.

### 3.3.11. Withdrawing of the Type 2 ST Route

A BGP speaker acting as a MUP controller withdraws Type 2 ST route based on deletion of corresponding session information through a northbound API as per [I-D.mhkk-dmm-srv6mup-architecture]. The northbound API definition for ST2 route withdraw is out of scope of this document.

In 3GPP 5G specific case, to withdraw a Type 2 ST route the MUP controller acquires the Endpoint consists of Endpoint address of GTP [TS.29281] tunnel and TEID for core side with the effective length of the Endpoint as input parameters. The MUP controller MUST advertise the withdraws of the Type 2 ST route.
When withdrawing the Type 2 ST route, the MUP controller SHOULD attach the BGP MUP Extended community corresponding to the Direct segment, and the Route Target Extended community which the MUP GWs are importing into the routing instance accomodating the corresponding Interwork segment to the Route Target Extended community.

3.3.12. Processing of the Type 2 ST Route

Both, the MUP GW and the MUP PE MAY receive the Type 2 ST routes from the MUP Controller in the BGP MUP network. A BGP speaker acting as a MUP PE or a MUP GW MAY keep the received MUP Type 2 ST routes advertised from the MUP Controller. The receiving BGP speaker will perform the importing of the received MUP Type 2 ST routes in the configured N3RAN routing instance based on the Route Target extended communities. A BGP speaker MAY discard the received Type 2 ST route if the speaker fails to import the route based on the Route Target extended communities.

The BGP speaker receiving the Type 2 ST routes SHOULD ignore the received nexthop in the MP_REACH_NLRI attribute.

A MUP GW receiving the Type 2 ST routes in a MP_REACH_NLRI attribute without a BGP MUP Extended community SHOULD consider the route as a malformed route. The MUP GW MUST handle such a malformed NLRI as a "Treat-as-withdraw" [RFC7606].

The MUP GW receiving Type 2 ST route with a BGP MUP Extended Community should extract the received segment identifier from the community. The segment identifier is used to resolve an appropriate Direct segment routing instance.

4. Security Considerations

The mechanisms described in this document could reuse the existing BGP security mechanisms [RFC4271] [RFC4272]. The security model and threats specific to Provider Provisioned VPNs, including L3VPNs, are discussed in [RFC4111]. The method defined in [RFC5925] SHOULD be used where authentication of BGP control packets is needed.

This document defines 3 new network nodes; MUP PE, MUP GW and a MUP Controller. These MUP BGP speakers SHOULD NOT establish BGP sessions with other BGP speakers in the domains which are not trusted without any explicit configuration or an operator intervention. Usage of procedures defined in [RFC5925] SHOULD be enforced at such boundaries to ensure the proper authentication of BGP control packets.
Furthermore, [RFC5925] will not help to keep the BGP messages private. To protect the BGP messages exchanged between BGP speakers from eavesdrop, establishing BGP sessions over encrypted paths SHOULD be considered.

MUP PEs and GWs SHOULD impose an upper bound on number of routes they should store to protect their control plane load. For example, BGP implementations MAY provide a configuration knob to impose an upper bound on Type 1 ST Routes.

5. IANA Considerations

This document defines a new BGP SAFI known as a BGP-MUP SAFI. IANA is requested to assign the value for the new SAFI from the Subsequent Address Family Identifiers (SAFI) registry.

This document defines a new Architecture Type for a BGP-MUP SAFI. IANA is requested to create a new Architecture Type NLRI registry for BGP-MUP SAFI. Furthermore, IANA is also requested to assign values for the following Architecture Types from the newly created BGP-MUP Architecture Type NLRI registry:

+ 1 - 3gpp-5g;

This document defines new NLRIs for a BGP-MUP SAFI. IANA is requested to create a new NLRI registry for BGP-MUP SAFI. Furthermore, IANA is also requested to assign values for the following NLRIs from the newly created BGP-MUP NLRI registry:

+ 1 - Discovery Internetwork route;
+ 2 - Direct Segment Discovery route;
+ 3 - Type 1 Session Transformed (ST) route;
+ 4 - Type 2 Session Transformed (ST) route;

This document defines a new BGP Extended Community called "SRv6 MUP Extended Community". This Community is of an extended type and is transitive. IANA is requested to assign the Type and the Sub-Type value for this Community from the Transitive Extended Community registry.

6. Contributors

In addition to the authors listed on the front page, the following individuals have also made significant contributions to the draft:

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EVPN Support for L3 Fast Convergence and Aliasing/Backup Path
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Abstract

This document proposes an EVPN extension to allow several of its multihoming functions, fast convergence and aliasing/backup path, to be used in conjunction with inter-subnet forwarding.

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1. Introduction

This document proposes an EVPN extension to allow several of its multihoming functions, fast convergence and aliasing/backup path, to be used in conjunction with inter-subnet forwarding. It re-uses the existing EVPN routes, the Ethernet A-D per ES and the Ethernet A-D per EVI routes, which are used for these multihoming functions. In particular, there are three use-cases that could benefit from the use of these multihoming functions:

a. Inter-subnet forwarding for host routes in symmetric IRB [RFC9135].

b. Inter-subnet forwarding for prefix routes in the interface-less IP-VRF-to-IP-VRF model [RFC9136].

c. Inter-subnet forwarding for prefix routes when the ESI is used exclusively as an L3 construct [RFC9136].

1.1. Ethernet Segments for Host Routes in Symmetric IRB

Consider a pair of multi-homing PEs, PE1 and PE2, as illustrated in Figure 1. Let there be a host H1 attached to them. Consider PE3 and a host H3 attached to it.

![Figure 1: Inter-subnet traffic between Multihoming PEs and Remote PE](image)

With Asymmetric IRB [RFC9135], if H3 sends inter-subnet traffic to H1, routing will happen at PE3. PE3 will be attached to the destination IRB interface and will trigger ARP/ND requests if it does not have an ARP/ND adjacency to H1. A subsequent routing lookup will resolve the destination MAC to H1’s MAC address. Furthermore, H1’s
MAC will point to an ECMP EVPN destination on PE1 and PE2, either due to host route advertisement from both PE1 and PE2, or due to Ethernet Segment MAC Aliasing as detailed in [RFC7432].

With Symmetric IRB [RFC9135], if H3 sends inter-subnet traffic to H1, a routing lookup will happen at PE3’s IP-VRF and this routing lookup will not yield the destination IRB interface and therefore MAC Aliasing is not possible. In order to have per-flow load balancing for H3’s routed traffic to H1, an IP ECMP list (to PE1/PE2) needs to be associated to H1’s host route in the IP-VRF route-table. If H1 is locally learned only at one of the multi-homing PEs, PE1 or PE2, due to LAG hashing, PE3 will not be able to build an IP ECMP list for the H1 host route.

With the extension described in this document, PE3’s IP-VRF becomes Ethernet-Segment-aware and builds an IP ECMP list for H1 based on the advertisement of ES1 along with H1 in a MAC/IP route and the availability of ES1 on PE1 and PE2.

1.2. Inter-subnet Forwarding for Prefix Routes in the Interface-less IP-VRF-to-IP-VRF Model

In the Interface-less IP-VRF-to-IP-VRF model described in [RFC9136] there is no Overlay Index and hence no recursive resolution of the IP Prefix route to either a MAC/IP Advertisement or an Ethernet A-D per ES/EVI route, which means that the fast convergence and aliasing/backup path functions are disabled. The recursive resolution of an IP Prefix route to an Ethernet A-D per ES/EVI route is already described in [RFC9136].

The scenario illustrated in Figure 2 will be used to explain the procedures.
Consider PE1 and PE2 are multi-homed to CE1 (in an All-Active Ethernet Segment ES1), and PE1, PE2 and PE3 are attached to an IP-VRF of the same tenant. Suppose H1’s host route is learned (via ARP or ND snooping) on PE1 only, and PE1 advertises an EVPN IP Prefix route for H1’s host route. If H3 sends inter-subnet traffic to H1, a routing lookup on PE3 would normally yield a single next-hop, i.e., PE1.

This document proposes the use of the ESI in the IP Prefix route and the recursive resolution to A-D per ES/EVI routes advertised from PE1 and PE2, so that H1’s host route in PE3 can be associated to an IP ECMP list (to PE1/PE2) for aliasing purposes.

1.3. Ethernet Segments for Prefix routes in IP-VRF-to-IP-VRF use-cases

This document also enables fast convergence and aliasing/backup path to be used even when the ESI is used exclusively as an L3 construct, in an Interface-less IP-VRF-to-IP-VRF scenario [RFC9136]. There are two use cases analyzed and supported by this document:

* IP Aliasing for EVPN IP Prefix routes
* Centralized Routing Model

1.3.1. IP Aliasing for EVPN IP Prefix routes

As an example, consider the scenario in Figure 3 in which PE1 and PE2 are multi-homed to CE1. However, and contrary to CE1 in Figure 2, in this case the links between CE1 and PE1/PE2 are used exclusively for L3 protocols and L3 forwarding in different BDs, and a BGP session established between CE1’s loopback address and PE1’s IRB address.
In these use-cases, sometimes the CE supports a single BGP session to one of the PEs (through which it advertises a number of IP prefixes seating behind itself) and yet, it is desired that remote PEs can build an IP ECMP list or backup IP list including all the PEs multi-homed to the same CE. For example, in Figure 3, CE1 has a single eBGP neighbor, i.e., PE1. Load-balancing for traffic from CE1 to H4 can be accomplished by a default route with next-hops PE1 and PE2, however, load-balancing from H4 to any of the prefixes attached to CE1 would not be possible since only PE1 would advertise EVPN IP Prefix routes for CE1’s prefixes. This document provides a solution so that PE3 considers PE2 as a next-hop in the IP ECMP list for CE1’s prefixes, even if PE2 did not advertise the IP Prefix routes for those prefixes in the first place.

1.3.2. Centralized Routing Model

Figure 4 illustrates a model in which multiple CEs establish an eBGP PE-CE session with a Centralized PE.
The CEs in this case are usually VNFs (Virtual Network Function entities) or CNFs (Containerized Network Function entities) and by provisioning the same network parameters on all of them, the operation gets significantly simplified. The configuration on the PEs also gets simplified, since the PE-CE eBGP sessions to the CEs are only configured on a centralized PE. In the diagram, CE1 is one of these VNF/CNFs that sets up a multi-hop eBGP session to the centralized PEC. As an example, CE1 advertises prefix 50.0.0.0/24 with Next Hop 10.0.0.1 (to PEC) via the multi-hop eBGP session. PEC then exports the prefix into a RT-5 route, following the Interface-less IP-VRF-to-IP-VRF model [RFC9136], with Next Hop PEC. When H4 sends traffic to an IP address of the subnet 50.0.0.0/24, the traffic will be forwarded to PEC first, and PEC will then forward to PE1 (or PE2). In other words, this model simplifies the configuration and
operation of the CEs, however, it introduces an inefficiency since traffic needs to go through the Centralized PE (PEC) instead of going directly to the PE(s) attached to the destination CE. The IP Aliasing solution specified in this document overcomes this inefficiency and allows traffic from PE3 to be forwarded directly to PE1 or PE2, without going through PEC.

1.4. Terminology and Conventions

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

* IRB: Integrated Routing and Bridging

* IRB Interface: Integrated Bridging and Routing Interface. A virtual interface that connects the Bridge Table and the IP-VRF on an NVE.

* BD: Broadcast Domain. An EVI may be comprised of one BD (VLAN-based or VLAN Bundle services) or multiple BDs (VLAN-aware Bundle services).

* Bridge Table: An instantiation of a broadcast domain on a MAC-VRF.

* CE: Customer Edge device, e.g., a host, router, or switch.s

* EVI: An EVPN instance spanning the Provider Edge (PE) devices participating in that EVPN.

* MAC-VRF: A Virtual Routing and Forwarding table for Media Access Control (MAC) addresses on a PE.

* Ethernet Segment (ES): When a customer site (device or network) is connected to one or more PEs via a set of Ethernet links, then that set of links is referred to as an 'Ethernet segment'.

* Ethernet Segment Identifier (ESI): A unique non-zero identifier that identifies an Ethernet segment is called an 'Ethernet Segment Identifier'.

* IP-VRF: A VPN Routing and Forwarding table for IP routes on an NVE/PE. The IP routes could be populated by any routing protocol, e.g., EVPN, IP-VPN and BGP PE-CE IP address families. An IP-VRF is also an instantiation of a layer 3 VPN in an NVE/PE.
* EVPN IP route: An EVPN IP Prefix route or an EVPN MAC/IP Advertisement route.

* LACP: Link Aggregation Control Protocol.

* PE: Provider Edge device.

* Single-Active Redundancy Mode: When only a single PE, among all the PEs attached to an Ethernet segment, is allowed to forward traffic to/from that Ethernet segment for a given VLAN, then the Ethernet segment is defined to be operating in Single-Active redundancy mode.

* All-Active Redundancy Mode: When all PEs attached to an Ethernet segment are allowed to forward known unicast traffic to/from that Ethernet segment for a given VLAN, then the Ethernet segment is defined to be operating in All-Active redundancy mode.

* RT-2: EVPN MAC/IP Advertisement route, as specified in [RFC7432].

* RT-4: EVPN Ethernet Segment route, as specified in [RFC7432].

* RT-5: EVPN IP Prefix route, as specified in [RFC9136].

2. Ethernet Segments for L3 Aliasing/Backup Path and Fast Convergence

The first two use cases described in Section 1 do not require any extensions to the Ethernet Segment definition and both cases support Ethernet Segments as a set of Ethernet links and specified in [RFC7432], or virtual Ethernet Segments as a set of logical links specified in [I-D.ietf-bess-evpn-virtual-eth-segment].

The third use case in Section 1 requires an extension to the way Ethernet Segments are defined and associated. In this case, the Ethernet Segment is a Layer-3 construct characterized as follows:

1. The ES is defined as a set of Layer-3 links to the multi-homed CE and its state MUST be linked to the layer-3 reachability from each multi-homed PE to the CE’s loopback address via a non-EVPN route in the PE’s IP-VRF.

2. The ESI SHOULD be of type 4 [RFC7432] and set to the router ID of the multi-homed CE.

3. All-active or single-active multi-homing redundancy modes are supported, however, the redundancy mode only affects the procedures in Section 3.
4. PEs attached to the same Layer-3 ES discover each other through the exchange of RT-4 routes (Ethernet Segment routes). DF Election procedures [RFC8584] MAY be used for single-active multi-homing mode.

5. The routes advertised from the multi-homed CE’s and installed in the PE’s IP-VRF table with the CE’s loopback as the next-hop MUST be re-advertised by the PE in EVPN IP Prefix routes with the ESI of the CE. The rest of the EVPN IP Prefix routes fields are set as per the Interface-less model in [RFC9136]. Note that the BGP PE-CE routes advertised by the multi-homed CE are installed in the IP-VRF normally irrespective of the Next Hop being resolved to an EVPN or a non-EVPN route, and they are exported as a RT-5 with the ESI.

In the example depicted in Figure 3, ES1 is defined as the set of layer-3 links that connects PE1 and PE2 to CE1. Its ESI, e.g., ESI-1, is derived as a type 4 ESI using the CE’s router ID. ES-1 will be operationally up in the PE as long as CE1’s loopback route is installed in the PE’s IP-VRF and learned via any routing protocol except for an EVPN route. E.g., an active static route to 1.1.1.1 via next-hop 10.0.0.2 would make the ES operationally up in PE1, and the eBGP routes received from CE1 with next-hop 1.1.1.1 will be re-advertised as RT-5 routes with ESI-1.

In the example illustrated in Figure 4, ES1 is a set of layer-3 links connecting PE1, PE2 and PEC to CE1. ESI-1 is derived as a type 4 ESI using the CE’s router ID, as in the previous example. CE1’s loopback route (which is associated to ES1) is installed in PE1 and PE2 via non-EVPN route, hence ES1 is operationally up in PE1 and PE2. On PE-C though, CE1’s loopback is installed via EVPN IP Prefix route, therefore, as per point 1 in the current section, ES1 is operationally down in PEC. As per point 5, this does not prevent PEC from exporting CE1’s prefixes into RT-5 routes with ESI-1. However, since ES-1 is operationally down in PEC, no IP A-D per EVI routes (Section 3) and no IP A-D per ES routes Section 4 for ESI-1 will be advertised from PEC, preventing PEC from attracting traffic destined to CE1.

3. IP Aliasing and Backup Path

In order to address the use-cases described in Section 1, above, this document proposes that:
1. A PE that is attached to a given ES will advertise a set of one or more Ethernet A-D per ES routes for that ES. Each is termed an 'IP A-D per ES' route and is tagged with the route targets (RTs) for one or more of the IP-VRFs defined on it for that ES; the complete set of IP A-D per ES routes contains the RTs for all of the IP-VRFs defined on it for that ES.

A remote PE imports an IP A-D per ES route into the IP-VRFs corresponding to the RTs with which the route is tagged. When the complete set of IP A-D per ES routes has been processed, a remote PE will have imported an IP A-D per ES route into each of the IP-VRFs defined on it for that ES; this enables fast convergence for each of these IP-VRFs.

2. A PE advertises for this ES, an Ethernet A-D Per EVI route for each of the IP-VRFs defined on it. Each is termed an 'IP A-D per EVI' route and is tagged with the RT for a given IP-VRF, and conveys a label that identifies that IP-VRF.

A remote PE imports an IP A-D per EVI route into the IP-VRF corresponding to the RT with which the route is tagged. The label contained in the route enables aliasing/backup path for the routes in that IP-VRF.

To address the third use-case described in Section 1, where the links between a CE and its multihomed PEs are used exclusively for L3 protocols and L3 forwarding, a PE uses the procedures described in 1) and 2), above.

The processing of the IP A-D per ES and the IP A-D per EVI routes is as defined in [RFC7432] and [RFC8365] except that the fast convergence and aliasing/backup path functions apply to the routes contained in an IP-VRF. In particular, a remote PE that receives an EVPN MAC/IP Advertisement route or an IP Prefix route with a non-reserved ESI and the RT of a particular IP-VRF SHOULD consider it reachable by every PE that has advertised an IP A-D per ES and IP A-D per EVI route for that ESI and IP-VRF.

3.1. Constructing the IP A-D per EVI Route

The construction of the IP A-D per EVI route is the same as that of the Ethernet A-D per EVI route, as described in [RFC7432], with the following exceptions:

* The Route-Distinguisher is for the corresponding IP-VRF.

* The Ethernet Tag should be set to 0.
* The route SHOULD carry the Route Target of the corresponding IP-VRF.

* The route MUST carry the MPLS label, VNI (VXLAN Network Identifier [RFC8365]) or Segment Routing IPv6 SID (Segment Identifier [I-D.ietf-bess-srv6-services]) that identifies the corresponding IP-VRF.

* The route MUST carry the PE’s MAC Extended Community if the encapsulation used between the PEs for inter-subnet forwarding is an Ethernet NVO tunnel [RFC9136].

* The route SHOULD carry the EVPN Layer 2 Extended Community [I-D.ietf-bess-rfc7432bis]. For all-active multihoming, all PEs attached to the specified ES will advertise P=1. For backup path, the Primary PE will advertise P=1 and the Backup PE will advertise P=0, B=1.

  - The Primary PE SHOULD be a PE with a routing adjacency to the attached CE.
  - The Primary PE MAY be determined by policy or MAY be elected by a DF Election as in [RFC8584] as described in Section 2.

4. Fast Convergence for Routed Traffic

Host or Prefix reachability is learned via the BGP-EVPN control plane over the MPLS/NVO network. EVPN IP routes for a given ES are advertised by one or more of the PEs attached to that ES. When one of these PEs fails, a remote PE needs to quickly invalidate the EVPN IP routes received from it.

To accomplish this, EVPN defined the fast convergence function specified in [RFC7432]. This document extends fast convergence to inter-subnet forwarding by having each PE advertise a set of one or more IP A-D per ES routes for each locally attached Ethernet segment (refer to Section 4.1 below for details on how these routes are constructed). A PE may need to advertise more than one IP A-D per ES route for a given ES because the ES may be in a multiplicity of IP-VRFs and the Route Targets for all of these IP-VRFs may not fit into a single route. Advertising a set of IP A-D per ES routes for the ES allows each route to contain a subset of the complete set of Route Targets. Each IP A-D per ES route is differentiated from the other routes in the set by a different Route Distinguisher (RD).

Upon failure in connectivity to the attached ES, the PE withdraws the corresponding set of IP A-D per ES routes. This triggers all PEs that receive the withdrawal to update their next-hop adjacencies for
all IP addresses associated with the Ethernet Segment in question, across IP-VRFs. If no other PE has advertised an IP A-D per ES route for the same Ethernet Segment, then the PE that received the withdrawal simply invalidates the IP entries for that segment. Otherwise, the PE updates its next-hop adjacencies accordingly.

These routes should be processed with higher priority than EVPN IP route withdrawals upon failure. Similar priority processing is needed even on the intermediate Route Reflectors.

4.1. Constructing IP A-D per Ethernet Segment Route

This section describes the procedures used to construct the IP A-D per ES route, which is used for fast convergence (as discussed in Section 4). The usage/construction of this route remains similar to that described in section 8.2.1. of [RFC7432] with a few notable exceptions as explained in following sections.

4.1.1. IP A-D per ES Route Targets

Each IP A-D per ES route MUST carry one or more Route Targets. The set of IP A-D per ES routes MUST carry the entire set of IP-VRF Route Targets for all the IP-VRFs defined on that ES.

4.2. Avoiding convergence issues by synchronizing IP prefixes

Consider a pair of multi-homing PEs, PE1 and PE2. Let there be a host H1 attached to them. Consider PE3 and a host H3 attached to it.

If the host H1 is learned on both the PEs, the ECMP path list is formed on PE3 pointing to (PE1/PE2). Traffic from H3 to H1 is not impacted even if one of the PEs fails as the path list gets corrected upon receiving the withdrawal of the fast convergence route(s) (IP A-D per ES routes).

In a case where H1 is locally learned only on PE1 due to LAG hashing or a single routing protocol adjacency to PE1, at PE3, H1 has ECMP path list (PE1/PE2) using Aliasing as described in this document. Traffic from H3 can reach H1 via either PE1 or PE2.

PE2 should install local forwarding state for EVPN IP routes advertised by other PEs attached to the same ES (i.e., PE1) but not advertise them as local routes. When the traffic from H3 reaches PE2, PE2 will be able forward the traffic to H1 without any convergence delay (caused by triggering ARP/ND to H1 or to the next-hop to reach H1). The synchronization of the EVPN IP routes across all PEs of the same Ethernet Segment is important to solve convergence issues.
4.3. Handling Silent Host MAC/IP route for IP Aliasing

Consider the example of Figure 1 for IP aliasing. If PE1 fails, PE3 will receive the withdrawal of the fast convergence route(s) and update the ECMP list for H1 to be just PE2. When the EVPN IP route for H1 is also withdrawn, neither PE2 nor PE3 will have a route to H1, and traffic from H3 to H1 is blackholed until PE2 learns H1 and advertises an EVPN IP route for it.

This blackholing can be much worse if the H1 behaves like a silent host. IP address of H1 will not be re-learned on PE2 till H1 ARP/ND messages or some traffic triggers ARP/ND for H1.

PE2 can detect the failure of PE1’s reachability in different ways:

a. When PE1 fails, the next hop tracking to PE1 in the underlay routing protocols can help detect the failure.

b. Upon the failure of its link to CE1, PE1 will withdraw its IP A-D route(s) and PE2 can use this as a trigger to detect failure.

Thus to avoid blackholing, when PE2 detects loss of reachability to PE1, it should trigger ARP/ND requests for all remote IP prefixes received from PE1 across all affected IP-VRFs. This will force host H1 to reply to the solicited ARP/ND messages from PE2 and refresh both MAC and IP for the corresponding host in its tables.

Even in core failure scenario on PE1, PE1 must withdraw all its local layer-2 connectivity, as Layer-2 traffic should not be received by PE1. So when ARP/ND is triggered from PE2 the replies from host H1 can only be received by PE2. Thus H1 will be learned as local route and also advertised from PE2.

It is recommended to have a staggered or delayed deletion of the EVPN IP routes from PE1, so that ARP/ND refresh can happen on PE2 before the deletion.

4.4. MAC Aging

In the same example as in Section 4.3, PE1 would do ARP/ND refresh for H1 before it ages out. During this process, H1 can age out genuinely or due to the ARP/ND reply landing on PE2. PE1 must withdraw the local entry from BGP when H1 entry ages out. PE1 deletes the entry from the local forwarding only when there are no remote synced entries.

5. Determining Reachability to Unicast IP Addresses
5.1. Local Learning

The procedures for local learning do not change from [RFC7432] or [RFC9136].

5.2. Remote Learning

The procedures for remote learning do not change from [RFC7432] or [RFC9136].

5.3. Constructing the EVPN IP Routes

The procedures for constructing MAC/IP Address or IP Prefix Advertisements do not change from [RFC7432] or [RFC9136].

5.3.1. Route Resolution

If the ESI field is set to reserved values of 0 or MAX-ESI, the EVPN IP route resolution MUST be based on the EVPN IP route alone.

If the ESI field is set to a non-reserved ESI, the EVPN IP route resolution MUST happen only when both the EVPN IP route and the associated set of IP A-D per ES routes have been received. To illustrate this with an example, consider a pair of multi-homed PEs, PE1 and PE2, connected to an all-active Ethernet Segment. A given host with IP address H1 is learned by PE1 but not by PE2. When the EVPN IP route from PE1 and a set of IP A-D per ES and IP A-D per EVI routes from PE1 and PE2 are received, then (1) PE3 can forward traffic destined to H1 to both PE1 and PE2.

If after (1) PE1 withdraws the IP A-D per ES route, then PE3 will forward the traffic to PE2 only.

If after (1) PE2 withdraws the IP A-D per ES route, then PE3 will forward the traffic to PE1 only.

If after (1) PE1 withdraws the EVPN IP route, then PE3 will do delayed deletion of H1, as described in Section 4.3.

If after (1) PE2 advertised the EVPN IP route, but PE1 withdraws it, PE3 will continue forwarding to both PE1 and PE2 as long as it has the IP A-D per ES and the IP A-D per EVI route from both.

6. Forwarding Unicast Packets

Refer to Section 5 in [RFC9135] and [RFC9136].
7. Load Balancing of Unicast Packets

The procedures for load balancing of Unicast Packets do not change from [RFC7432]

8. IP Aliasing and Unequal ECMP for IP Prefix Routes

[I-D.ietf-bess-evpn-unequal-lb] specifies the use of the EVPN Link bandwidth extended community to achieve weighted load balancing to an ES or Virtual ES for unicast traffic. The procedures in [I-D.ietf-bess-evpn-unequal-lb] MAY be used along with the procedures described in this document for any of the three cases described in Section 1, with the following considerations:

* The ES weight is signaled by the multi-homed PEs in the IP A-D per ES routes.

* The remote ingress PE learning an EVPN IP Route to prefix/host P that is associated to a weighted load balancing ES, will follow the procedures in [I-D.ietf-bess-evpn-unequal-lb] to influence the load balancing for traffic to P.

* [I-D.ietf-bess-evpn-unequal-lb] also allows the use of the EVPN Link Bandwidth Extended Community along with RT-5s. If the ingress PE learns a prefix P via a non-reserved ESI RT-5 route with a weight (for which IP A-D per ES routes also signal a weight) and a zero ESI RT-5 that includes a weight, the ingress PE will consider all the PEs attached to the ES as a single PE when normalizing weights.

As an example, consider PE1 and PE2 are attached to ES-1 and PE1 advertises an RT-5 for prefix P with ESI-1 (and EVPN Link Bandwidth of 1). Consider PE3 advertises an RT-5 for P with ESI=0 and EVPN Link Bandwidth of 2. If PE1 and PE2 advertise an EVPN Link Bandwidth of 1 and 2, respectively, in the IP A-D per ES routes for ES-1, an ingress PE4 SHOULD assign a normalized weight of 1 to ES-1 and a normalized weight of 2 to PE3. When PE4 sprays the flows to P, it will send twice as many flows to PE3. For the flows sent to ES-1, the individual PE EVPN Link Bandwidths advertised in the IP A-D per ES routes will be considered.
9. Security Considerations

The mechanisms in this document use EVPN control plane as defined in [RFC7432]. Security considerations described in [RFC7432] are equally applicable. This document uses MPLS and IP-based tunnel technologies to support data plane transport. Security considerations described in [RFC7432] and in [RFC8365] are equally applicable.

10. IANA Considerations

No IANA considerations.

11. Contributors

12. Acknowledgments

13. References

13.1. Normative References


13.2. Informative References

[I-D.ietf-bess-evpn-virtual-eth-segment]

[I-D.ietf-bess-evpn-unequal-lb]

[I-D.ietf-bess-srv6-services]

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EVPN Mac Dampening Back-off
draft-saum-bess-dampening-backoff-02

Abstract

MAC move handling in EVPN deployments is discussed in detail in
[RFC7432]. There are few optimizations which can be done in existing
way of handling the mac duplication. This document describes few of
the potential techniques to do so.

Status of This Memo

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1. Important Terms

MDAS: Mac Dampening Attribute Set:

   MDT: Mac Dampening Timer

   MDC: Mac Dampening Count

   MFT: MAC Freezing Timer

Mac Dampening: Process of stalling the mobility of MAC as define in
[RFC7432].

VTEP: Virtual Tunnel End Point or Vxlan Tunnel End Point

DT: Damptened Time: Actual time taken to dampen the contentious MAC
2. Introduction

The host mobility solution described in [RFC7432] elaborates on few use-cases related to dual mac discovery which leads to dampening logic coming into play. The host move handling logic addresses the problem of frequent mac-moves and culminates by freezing the MACs against further moves. If there is no mellowing down of the issue, then it leads to unending cycle of mac dampening and freezing. Hence this problem needs organic measures for arriving at MAC freezing point, sooner than later.

The events that can lead to never ending duplication are as follows:

(a) Misconfiguration of hosts with identical configuration, in the same bridge-domain, across ESIs and across NVEs.

(b) Looping of traffic due to layer 2 loops created in the bridge domain in the tenant network behind the NVEs.

2.1. Misconfiguration of Hosts

Consider the following figure wherein two hosts, Host-1 and Host-2, are misconfigured with same mac-address MAC-1. These hosts are placed behind two different Ethernet segments, ES12 and ES3 respectively and hooked to the same bridge-domain (BD-1). PE1, PE2 and PE3 will get into a never ending loop of learning the MAC-1 locally and also from the remote Vtep. Thus entering into a control-plane BGP-EVPN cycle of bumping up the sequence number in the MACMobility Extended Community till the maximum MAC move count is hit with the stipulated time. The MAC published to other Vteps like PE4 also changes accordingly based on the latest update with highest sequence number.
2.2. Loopy Traffic in Tenant Network

Consider the following case of a loopy tenant network, leading to MAC duplicity in the network. Let's say, Host-1 generates a BUM traffic like GARP (Gratuitous ARP) and sends it over the VLAN which is part of BD-1 and mapped to a configured EVI on the PEs. PE1 sprays the BUM over the EVPN fabric tying it with the mapped EVI. The BUM packet arrives at PE1 (assuming it's the elected DF) over the EVPN fabric. PE1 sprays the traffic towards the directly attached tenant network attached, tagging it with VLAN that maps to the bridge domain, BD-1, which in turn maps to the MAC-VRF pointed to by the EVI. If the layer-2 network on tenant side is loopy due to STP network not...
converging or STP not configured at all, or for some other unknown reason (not under the purview of this document); then the BUM traffic may loop back to PE1, thus creating a duplicate MAC learning for MAC-1. Till the tenant network is curtailed or put to order via admin intervention or otherwise, continuous MAC moves for MAC-1 can be observed between PEs attached to ethernet segment ES12 (PE1) and ES3 (PE2).

![Diagram of network topology]

**Figure 2**: Loopy traffic in Tenant Network

**LEGEND:**

- PE1, PE2, PE3: Vxlan/overlay gateways
- HOST-1: Hosts behind PE3
- MAC_1: MAC address of Host-1
- ES12: Ethernet segment between PE1 and PE2 for BD-1
- ES3: Ethernet segment attached to PE3 for BD-1
- BD-1: Bridge Domain 1

3. Requirements

3.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].
When used in lowercase, these words convey their typical use in common language, and they are not to be interpreted as described in [RFC2119].

4. Problem Description

The mac dampening procedure mentioned in [RFC7432], suggests that a Overlay Tunnel Endpoint that detects the mac mobility event upon local learning, should start a ‘M’ seconds timer and track the MAC for ‘N’ moves before the timer expires. Hence forth, concluding that it is a MAC Duplication issue and freezing the MAC while also raising the alarm, for the admin to take corrective action. It is observed in few vendor implementations, that involves defreezing the MAC in deterministic time (configurable or derived) after freezing it, with a positive assumption that admin shall take corrective action meanwhile. Else, the subsequent unfreeze shall end up in the same cycle of MAC Duplication detection and freezing of the MAC. In case of lazy, none or inaccurate intervention by the admin, this can potentially result in a prolong state of network disarray:

(1) Unnecessary and periodic control-plane protocol churn

(2) Exchange of control plane states which are transient and inaccurate

(3) Reachability to end device remains in the realms of ambiguity for prolonged duration

(4) Traffic destined to the Duplicate MAC case, panning across fabrics, sites or across geographies, ends up hogging the precious WAN bandwidth.

Potential solutions are discussed subsequent sections.

5. Solution(s)

The potential solutions are as follows:

5.1. Mac Freeze

The eventual solution is to FREEZE the MAC forever till admin does the clearing of the MAC. The unfreeze and clearing actions are not organic in nature and can be accompanied by unwarranted impact like clearing of other MACs in the bridge-domain. The way out may be resetting the layer-2 port and thus impacting all tenant bridge-domains hosted on the port. This solution, hence, does not always solves or mitigate the situation, or, it may create a situation from which the eventual bail-out is expensive and not restricted to the
5.2. Backing Off MAC Mobility Timer and Count

The best-bet to organically mellow down the never ending MAC-mobility (indicating Duplicate MAC), is to freeze the MAC temporarily, for lets say, the same time as MAC Dampening Timer (MDT). Lets term this timer as MAC Freeze Timer (MFT). MFT is the time span for which the contentious MAC is frozen, i.e., no further control plane and data flow is allowed for this MAC. The duplicity/un-ending-mobility is expected to be addressed by the admin. In case the problem is not addressed within MAC Freeze Timer, the MAC duplicity is again identified based on the MAC mobility count within the MAC Dampening timer. The best way forward MAY be:

1) to get to the duplicity conclusion faster than the earlier iteration
2) and freeze the MAC for a longer duration than earlier iteration
3) , With the assumption that the problem shall be resolved in that time frame.

The MAC Dampening Attribute Set (MDAS), comprises of following three parameters:

1) MAC Dampening Timer (MDT): Defined in [RFC7432]
2) MAC Dampening Count (MDC): Defined in [RFC7432]
3) MAC Freeze Timer (MFT): Time for which the MAC is frozen after MAC duplicity is detected

For example, let the first iteration of MDAS_iter_1 {MDT=180 seconds, MDC=5, MFT=180 seconds}. The default values of MDT and MDC are picked from [RFC7432], while lets define the default value of MFT same as MDT. In case admin fails to intervene, the MAC is unfrozen after MFT expires.

For second iteration of the MDAS for the problem-MAC, i.e. MDAS_iter_2 = function (MDAS_iter_1). The MDT and MDC values in second iteration are derived by backing off the MDT and MDC values by a pre-defined delta, i.e.

1) MDAS_iter_2 (MDT) = MDAS_iter_1 (MDT) decrement_timer_delta
2) MDAS_iter_2 (MDC) = MDAS_iter_1 (MDC) decrement_count_delta
Thus reducing the time and moves to conclude on duplicity of the MAC. The values of decrement_timer_delta and decrement_count_delta can be configured or derived on a case to case basis. [TBD: Elaborate on the case]. The next step is to freeze the MAC for some more time as compared to the previous iteration set of MDAS, thus increasing the probability of the admin, correcting the issue:

\[(1)\text{ MDAS}_{iter\_2} (\text{MFT}) = \text{MDAS}_{iter\_1} (\text{MFT}) + \text{increment\_timer\_delta} \]

\[(2)\text{ The value of increment\_timer\_delta is also configurable in nature.} \]

5.2.1. MDAS Derivation

The following formulae generalizes the derivation of MDAS attributes in the Nth iteration of Duplicate MAC detection on a PE:

\[(1)\text{ MDAS}_{iter\_\text{(N)}} (\text{MDT}) = (\text{MDAS}_{iter\_\text{(N-1)}} (\text{MDT})) - \text{decrement\_mdt\_delta} \]

\[(2)\text{ MDAS}_{iter\_\text{(N)}} (\text{MDC}) = (\text{MDAS}_{iter\_\text{(N-1)}} (\text{MDC})) - \text{decrement\_mdc\_delta} \]

\[(3)\text{ MDAS}_{iter\_\text{(N)}} (\text{MFT}) = \text{MDAS}_{iter\_\text{(N-1)}} (\text{MFT}) + \text{increment\_mft\_delta} \]

Where in, the following values for 1st iteration can be define as follows:

\[
\begin{align*}
\text{MDAS}_{iter\_1} (\text{MDT}) & = 180 \text{ seconds} \\
\text{MDAS}_{iter\_1} (\text{MDC}) & = 5 \\
\text{MDAS}_{iter\_1} (\text{MFT}) & = 180 \text{ seconds. Many implementations keep the MDT and MFT values as same.}
\end{align*}
\]

The derivation of MDAS perimeters can be exponential in nature. The delta values can be exponentially increased or decreased after certain iterations, thus triggering a exponential backing off the delta values.

5.2.1.1. MDAS Boundry Values

The new MDT value SHOULD not be less than the time taken to Dampen the MAC movement in last set of MDAS iteration. On the same lines, the new MDC count SHOULD not go below ‘2’, as count below 2, the MAC Dampening procedure does not gets triggered.
5.2.2. Delta Values Calculation

Following bullets give an overview of potential ways the delta values, i.e., decrement_mdt_delta, increment_mdc_delta and decrement_mft_delta:

(a) Delta values should be such that they SHOULD not infringe the time or count taken to reach Dampening state in the last set

(b) Delta values are static all through the sets

(c) Delta variable gets incremented/decremented based on the reduction in time (proportionally) to achieve the ‘Dampened state’ in the last ‘MDAS set’ as compared to the time to reach the ‘Dampened state’ in the MDAS set previous to the last one. For the same, the time taken to reach the Dampened State should be cached so that comparisons can be made in subsequent sets. In case, it is the first ‘MDAS Set’, the delta values MAY be either default or configured ones. For the second ‘MDAS set’, the value MAY be cross-checked against the Dampened time for the first set.

(d) Delta values are always inherited from admin configuration.

As mentioned in the Section 5.2.1, the derivation of new delta values can be done by exponentially backing them off in subsequent MDAS set(s).

5.3. Backing Off Example

This section describes the example of MDAS calculation with respect to the use-case defined in Section 2.1. Though it’s equally applicable to the case described in Section 2.2. This example explains the logic in perspective of PE1. Let’s say PE1 learns the MAC-1 locally and publishes it over EVPN control plane before PE2 does the same. PE1 publishes it over control plane before PE2 learns it locally (ignoring the case where both learn in tandem and publish it over control plane). Subsequently, PE2 learns it and publishes it over control plane with sequence number 1. PE1 starts the dampening logic by incrementing the local count by 1 and starting the dampening timer. If this jiggle goes on for 5 counts at PE1, MAC Dampening logic described in [RFC7432] shall freeze the MAC. PE1 SHOULD cache the time it took to dampen the MAC. Let’s say it’s 30 seconds.

Assuming admin does not take any action, before MAC freeze timer expires and PE1 defreezes the MAC, it will start moving again. PE1 shall reduce the MDT value by decrement_mdt_delta = 30 seconds to 150 seconds. The MDC counts are reduced by decrement_mdc_count = 1 to 4
and the MFT is incremented by increment_mft_delta = 20 seconds to 170 seconds. Thus PE1 shall wait for 150 seconds for concluding the dampening logic and tracks the MAC for 4 moves. Once dampening is hit, MAC is rendered as frozen for 170 seconds for admin to take action thus giving some more time for admin to take action.

The whole intention is to gradually move towards a permanent freeze of the MAC if no admin does not do the needful in the stipulated time frame.

6. Backward Compatibility

The backward comptability is a no-op for MDAS derivation and recalculation, as MAC Dampening logic is very local to the Vtep. Even if the remote Vtep does not conforms to the logic presented in this literature, it will still work towards the dampening the frequent mac-mobility with the same parameters of MDT and MDS. The instant freezing or temporary freezing of the dampened MAC is implementation dependent and should not impact or get impacted by the MDAS derivations presented in this document.

7. Security Considerations

This document inherits all the security considerations discussed in [RFC7432].

8. IANA Considerations

This document inherits all the IANA considerations discussed in [RFC7432].

9. Acknowledgements

10. References

10.1. Normative References


10.2. Informative References


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Abstract

The Designated forwarder concept is leveraged to prevent looping of BUM traffic into tenant network sourced across NVO fabric for multihoming deployments. [RFC7432] defines a prelimn approach to select the DF for an ES,VLAN or ES,Vlan Group panning across multiple NVE’s. [RFC8584] makes the election logic more robust and fine grained inculcating fair election of DF handling most of the prevalent use-cases. This document presents a deployment problem and a corresponding solution which cannot be easily resolve by rules mentioned in [RFC7432] and [RFC8584].

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6 March 2022
1. **Important Terms**

DF: Designated Forwarder as defined in [RFC7432].

VTEP: Virtual Tunnel End Point or Vxlan Tunnel End Point

2. **Introduction**

The Designated forwarder concept is leveraged to prevent looping of BUM traffic into tenant network sourced across NVO fabric for multihoming deployments. [RFC7432] defines a prelimm approach to select the DF for an ES,VLAN or ES,Vlan Group panning across multiple NVE’s. [RFC8584] makes the election logic more robust and fine grained inculcating fair election of DF handling most of the prevalent use-cases. This document presents a deployment problem and a corresponding solution which cannot be easily resolve by rules mentioned in [RFC7432] and [RFC8584].

3. **Requirements Language**

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

Dikshit & Joshi

Expires 7 September 2022
4. Problem Description

It’s a typical case of Firewall devices also configured as default gateway for the NVO fabric or default gateways in turn redirect traffic to firewalls over shared vlan. This example for simplicity assumes the former case wherein firewall is also configured as default gateway for all VLANs in the site (SITE-1 and SITE-2).

All PEs (Vtep1 and Vtep2 in below example) in the diagram are attached to same ES and both intend to act act as DF for the broadcast domain (BD-1) for their respective sites. As already mentioned, this is a typical case of firewall-gateways (active/active) across fabrics (sites), Where in, the preferred firewall-gateway is the one local to the site, whereas, upon failure, packets need to be redirected (over WAN, via DCI/VPN) towards the remote site firewall. The firewall-device is connected to it’s first-hop vtep over the same bridge-domain and same ESI. All in all, it’s an emulated multi-homing scenario. This is a scenario of firewall devices hosting same (IP and MAC) credentials.

Simplistic example: There are two sites, SITE-1 and SITE-2 in the below diagram. Traffic (including BUM) generated by Host1 (in SITE-1) (for a bridge-domain) should run through site-local firewall instance (firewall_1) preferably. Only in case of local-outage, the traffic should be send across over WAN to the remote firewall (firewall_2). Same should apply to traffic generated by Host2 (in SITE-2), wherein, it should preferably run through the local firewall (firewall_2) and over a failure should go over the WAN towards firewall_1.

Vtep1/2 learn the firewall MAC (MAC_F) as local learning and also from the remote Vtep2/1. But since both the learnings are over the same ESI, it should not lead to MAC move. Cometh the local firewall failure, Vteps (1 or 2) should start redirecting the traffic to remote SITE. Any ARP request (BUM traffic) for firewall credentials landing at either Vtep1 or Vtep2 should be flooded to network towards the local firewall.
5. Solution(s)

The control plane part of the solution can be leveraged from the ‘DF Election Extended Community’ described in [RFC8584]. Since the requirement is to ensure all the PEs attached to ESI forward the BUM traffic arriving from NVO fabric towards the Attachment circuits (ACs) configured over the ES for a BD (broadcast domain) mapped to VLAN or bundle of VLANs. As explained in the above section that this is a case where PEs are in disparate networks and the ACs behind them are not connected to each other to a common physical device.

This document proposes a new mode of DF-election, ALL-PEs-DF where-in all of the participating PEs intend to play DF role for a VLAN(s) enabled on an ESI. This requires "DF Election Extended Community" to carry this information with the ES route to indicate it to remote PEs. This ensures all PEs receiving BUM traffic over NVO fabric destined to ESI, BD, SHOULD flood it on the associated ES on the access/tenant side. PE MAY be explicitly configured to choose the ALL-PEs-DF mode.

5.1. Sending All PEs are DF mode

The All-PEs-DF mode is used as follows:

1. PEs configured to use ALL-PEs-DF mode SHOULD set "DF Alg" algorithm field in ‘DF Election Extended Community’ to appropriate value.

2. This document proposes value ‘2’ for All-PEs-DF mode, as values ‘0’ and ‘1’ are already defined for usage in [RFC8584].

Figure 1: Active-Active Firewall Across Sites
This algorithm is agnostic to the values carried in 'Bitmap' but does not discounts any use-case(s) in future which may need extra information carried in 'Bitmap' along with All-PEs-DF mode.

5.2. Receive All PEs are DF mode

When a PE receives the ES routes from all the other PEs for the ES in question carrying the ALL-PEs-DF mode set in 'DF Election Extended Community', it SHOULD checks to see if all the advertisements have the Extended Community with 'All-DF-mode' set as 'DF Alg'. If yes, then SHOULD ignore the 'Bitmap' and 'Rsvd' field in the extended community. As also mentioned in [RFC8584], if even a single advertisement for Route Type 4 is received without the locally configured DF Alg and capability, the default DF election algorithm MUST be used as prescribed in [RFC7432].

5.3. Example of algorithm

The BGP-EVPN control plane support prescribed in this document helps in resolving the problem described in Section 4. If PEs Vtep1 and Vtep2 are configured to use ALL-PEs-DF mode, then any BUM traffic from respective hosts Host1/Host2 over the EVPN fabric, should get broadcasted towards the AC for the ESI, Vlan to which the firewall_1/firewall_2 (respectively) is attached. For example the arp-request for the Firewall IP will be honored by the Firewall_1 behind the Vtep1 which receives the ARP-request, whereas, when Vtep2 receives the arp-request it will be honored by Firewall_2. Vtep1 and Vtep2 will publish the arp-request in their respective ACs attached to the firewall on which Vlan, ESI is enabled.

6. Interoperability with other Algos

Since All-DF-algo is special mode and not exactly an algorithm, which requires the participation of all PEs for an ESI, VLAN. Hence, even if one PE publishes an algo which is NOT "All-DF-mode", other PEs SHOULD revert back to default algorithm. The reason being that, if there are PE1, PE2, PE3 and PE4 in contention. PE1 and PE2 publishes DF Algo 'ALL-PEs-DF', PE3 publishes '0' and PE4 publishes '1'. Once this mismatch is perceived, all PEs SHOULD try and converge towards the default mode. An admin intervention may be required to achieve the same or to converge on any other supported 'DF Algo'.
7. Backward Compatibility

As prescribed in [RFC8584], PEs not supporting (hence not publishing) 'ALL-PEs-DF', SHOULD ignore the processing of the 'DF Election Extended Community' and SHOULD indulge in DF-election using the default algorithm mentioned in [RFC7432]. The PEs configured with this new algorithm (hence publishing it), if receive Route Type 4 without 'DF Election Extended Community', SHOULD also revert back to default algorithm. If PEs receive Route Type 4 with another algorithm published in 'DF Election Extended Community', then it should follow procedures prescribed in Section 6.

8. Impact on Local Bias

There is no impact on the local-bias handling, as the PE receiving the BUM from access side over {ESI, VLAN} and relays it to other PEs that published {ESI, VLAN} in Route Type 4; the receiving side PEs will not relay it to EVPN fabric nor will they redirect it to same ESI configured with same VLAN on the access/tenant side.

9. Security Considerations

This document inherits all the security considerations discussed in [RFC7432] and [RFC8584].

10. IANA Considerations

This document inherits all the IANA considerations discussed in [RFC7432] and [RFC8584].

11. Acknowledgements

12. References

12.1. Normative References


12.2. Informative References


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Domain Path (D-PATH) for Ethernet VPN (EVPN) Interconnect Networks
draft-sr-bess-evpn-dpath-01

Abstract

The BGP Domain PATH (D-PATH) attribute is defined for Inter-Subnet Forwarding (ISF) BGP Sub-Address Families that advertise IP prefixes. When used along with EVPN IP Prefix routes or IP-VPN routes, it identifies the domain(s) through which the routes have passed and that information can be used by the receiver BGP speakers to detect routing loops or influence the BGP best path selection. This document extends the use of D-PATH so that it can also be used along with other EVPN route types.

Status of This Memo

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1. Introduction

The BGP Domain PATH (D-PATH) attribute [I-D.ietf-bess-evpn-ipvpn-interworking] is defined for Inter-Subnet Forwarding (ISF) BGP Sub-Address Families that advertise IP prefixes. When used along with EVPN IP Prefix routes or IP-VPN routes, it identifies the domain(s) through which the routes have passed and that information can be used by the receiver BGP speakers to detect routing loops or influence the BGP best path selection. This document extends the use of D-PATH so that it can also be used along with other EVPN route types.

The D-PATH attribute can be used to prevent control plane loops for EVPN routes, or to provide full path visibility of all the EVPN Interconnect Gateways through which a route has gone and modify the best path selection based on it. Some use cases in which D-PATH can be used along with (non-IP Prefix) EVPN routes follow, but the use cases are not limited to the ones described in this section.

1.1. D-PATH to Prevent Loops for EVPN Routes

Figure 1 illustrates an EVPN Interconnect case where EVPN MAC/IP Advertisement routes can be looped indefinitely. The three Gateways (GW1, GW2 and GW3) and PE1 in the diagram are attached to the same EVPN Broadcast Domain (BD1). However, BD1 is extended throughout three different domains that are interconnected by the Gateways, which follow [RFC9014] procedures. Suppose a host with MAC address M1 is learned on GW1 and GW1 advertises an EVPN MAC/IP Advertisement route for M1 into Domain-1 and Domain-2. When the route gets imported by GW2 and GW3 and later exported into Domain-3, GW2 and GW3 may redistribute each other’s route for M1 back into Domain-1 and Domain-2, respectively, creating a loop. D-PATH can be used by the Gateways when redistributing the route between Domains, to identify the Domains through which the route for M1 has gone. When GW1 receives an EVPN MAC/IP Advertisement route for M1 that contains a D-PATH with a domain-id locally assigned, GW1 identifies the route as "looped".

![Diagram of EVPN Interconnect case](image)

Figure 1: Loops for EVPN routes
Similar examples are possible with EVPN VPWS services on the Gateways and PEs, where loop prevention for the redistributed A-D per EVI routes is needed. D-PATH provides the end to end path visibility that is required to prevent the loop.

1.2. Add Path Visibility and Influence Best Path Selection for EVPN Routes

Figure 2 illustrates another [RFC9014] EVPN Interconnect case where, in addition to using D-PATH to prevent EVPN MAC/IP Advertisement route loops when redistributing routes between domains, the D-PATH attribute can also influence the best path selection for the routes. For example, if all the Gateways in the diagram are attached to the same BD1, an EVPN MAC/IP Advertisement route for MAC address M1 advertised by GW1 is advertised into Domain-1 and Domain-4. Two routes for M1 will arrive at GW3 with different Route Distinguishers and BGP Next Hops. If D-PATH is used by all the Gateways, the two routes arriving at GW3 will have a different sequence of domain-ids in the D-PATH attribute. GW3 can use the length of the D-PATH as a way of influencing the selection (i.e., the shortest D-PATH route is selected). D-PATH improves the path visibility of the route since it provides information about all the Domains through which the route has passed.

![Figure 2: Influence Best Path Selection for EVPN routes](image-url)
2. Conventions used in this document

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

3. Terminology

This section summarizes the terminology that is used throughout the rest of the document.

* AC: Attachment Circuit or logical interface associated to a given BT. To determine the AC on which a packet arrived, the PE examines the combination of a physical port and VLAN tags (where the VLAN tags can be individual c-tags, s-tags or ranges of both).

* BD and BT: a Broadcast Domain and Bridge Table, as defined in [I-D.ietf-bess-rfc7432bis]. A BD is a group of PEs attached to the same EVPN layer-2 multipoint service. A BT is the instantiation of a Broadcast Domain in a PE. When there is a single Broadcast Domain in a given EVI, the MAC-VRF in each PE contains a single BT. When there are multiple BTs within the same MAC-VRF, each BT is associated to a different Ethernet Tag. The EVPN routes specific to a BT indicate which Ethernet Tag the route corresponds to.

* ES and ESI: Ethernet Segment and Ethernet Segment Identifier, as defined in [I-D.ietf-bess-rfc7432bis].

* EVPN Domain: two PEs are in the same EVPN Domain if they are attached to the same service and the packets between them do not require a data path lookup of the inner frame (e.g., in the BT of a MAC-VRF) in any intermediate router. An EVPN Domain Gateway PE is always configured with multiple Domain identifiers (EVPN Domain-ID) in the MAC-VRF or VPWS that connects those EVPN Domains, each EVPN Domain-ID representing an EVPN Domain.

Example: Figure 3 illustrates an example where PE1 and PE2 belong to different EVPN Domains since packets between them (for flows between hosts with MAC addresses M1 and M2) require a MAC lookup in two of the gateways that are connecting the three EVPN Domains. E.g., if frames from M1 to M2 go through PE1, GW1, GW3 and PE2, a MAC lookup is performed at GW1 and GW3.
* EVPN Domain Gateway: a PE where a service (BD or VPWS instance) is instantiated and is attached to two or more EVPN Domains. An example of EVPN Domain Gateway is a PE following the procedures in section 4.4 or section 4.6 of [RFC9014]. In the example in Figure 3, GW1 and GW2 connect EVPN Domains 1 and 2, whereas GW3 and GW4 connect EVPN Domains 2 and 3. GW1 and GW2 import the MAC/IP Advertisement route for M1 coming from the EVPN Domain 1 into the MAC-VRF for BD1, and redistribute it into EVPN Domain 2. Likewise, GW3 and GW4 import the route into their MAC-VRF and re-advertise it into EVPN Domain 3.

* MAC-VRF: a MAC Virtual Routing and Forwarding table, as defined in [I-D.ietf-bess-rfc7432bis]. It is also the instantiation of an EVI (EVPN Instance) in a PE.

4. Use of Domain Path Attribute (D-PATH) with EVPN routes

This document extends the use of the D-PATH attribute specified in [I-D.ietf-bess-evpn-ipvpn-interworking] so that D-PATH can be advertised and processed along with the following EVPN route types:

* EVPN MAC/IP Advertisement routes that are not used for Inter-Subnet Forwarding (ISF). Note that if the EVPN MAC/IP Advertisement route is used for Inter-Subnet Forwarding as in [RFC9135], the procedures for the D-PATH advertisement and processing are described in [I-D.ietf-bess-evpn-ipvpn-interworking].

* EVPN A-D per EVI routes that are used for EVPN VPWS [RFC8214]. Advertised A-D per EVI routes not used for EVPN VPWS SHOULD NOT contain D-PATH.
As discussed, the use of D-PATH with EVPN IP Prefix routes is specified in [I-D.ietf-bess-evpn-ipvpn-interworking]. When used along with EVPN routes other than IP Prefix routes, the D-PATH attribute is characterized as follows:

1. D-PATH is composed of a sequence of Domain segments following the format specified in [I-D.ietf-bess-evpn-ipvpn-interworking] where each Domain is represented as <DOMAIN-ID:ISF_SAFI_TYPE>. In this specification, DOMAIN-ID is an EVPN Domain identifier configured in an EVPN Domain Gateway and ISF_SAFI_TYPE is set to either 70 (EVPN) or 0 (local route). To simplify the explanation, this document represents the domains for EVPN routes as <Domain-ID:TYPE>.

2. D-PATH identifies the sequence of EVPN Domains the route has gone through, with the last <Domain-ID:TYPE> entry (rightmost) identifying the first PE or EVPN Domain Gateway that added the D-PATH attribute.

3. For non-ISF MAC/IP Advertisement routes or A-D per EVI routes, D-PATH SHOULD be added/modified by a EVPN Domain Gateway that redistributes the route between EVPN Domains and MAY be added by a PE or EVPN Domain Gateway that originates the route, as follows:

   a. An EVPN Domain Gateway that connects two EVPN Domains "X" and "Y", and receives a route on a EVPN Domain identified by a Domain-ID "X" SHOULD add a domain <X:EVPN> to the existing (or newly added) D-PATH attribute when redistributing the route to EVPN Domain "Y". The route is redistributed if it is first imported in a MAC-VRF (or VPWS instance), the MAC (or Ethernet Tag) is active, and policy allows the re-export of the route to a BGP neighbor.

   b. An EVPN Domain Gateway MAY also add the D-PATH attribute for locally learned MACs or MAC/IP pairs. In this case, the domain added would be <A:0>, where "A" is the Domain-ID configured on the Gateway's MAC-VRF that is specific to local routes (MAC/IP learned via local AC), and "0" is the TYPE of the EVPN Domain and indicates that the route is locally originated and not redistributed after receiving it from a BGP-EVPN neighbor. The EVPN Domain-ID for local routes MAY be shared by all the EVPN Domain Gateways of the same redundancy group for local routes, or each EVPN Domain Gateway of the redundancy group can use its own Domain-ID.
c. A PE that is connected to a single EVPN Domain (therefore the PE is not an EVPN Domain Gateway) MAY add the D-PATH with a domain <B:0>, where "B" is the Domain-ID configured on the PE’s MAC-VRF (or VPWS) for locally learned MAC/IPs (or Ethernet Tag IDs for VPWS). "0" is the TYPE that indicates the route is not re-advertised, but originated in the PE.

4. For EVPN Inclusive Multicast Ethernet Tag routes, an EVPN Domain Gateway must not redistribute routes between Domains as specified in [RFC9014]. An EVPN Domain Gateway originates an EVPN Inclusive Multicast Ethernet Tag route per Domain to which the Gateway is attached, so that BUM traffic can be attracted from one Domain to the rest. Therefore, only the above point 3.b. applies to EVPN Domain Gateways. That is, an EVPN Domain Gateway MAY add a <A:0> D-PATH attribute for the Inclusive Multicast Ethernet Tag routes generated for the EVPN Domains, where "A" is the configured local Domain-ID, and "0" is the TYPE that indicates the route is locally originated and not redistributed across EVPN Domains. When two EVPN Domain Gateways of the same redundancy group connect two EVPN Domains "X" and "Y" and D-PATH is used for EVPN Inclusive Multicast Ethernet Tag routes, it is RECOMMENDED to add the D-PATH attribute with the same local Domain-ID and only on "X" or "Y" but not on both Domains.

5. On received EVPN routes, D-PATH is processed and used for loop detection (Section 4.4) as well as to influence the best path selection (Section 4.1, Section 4.2 and Section 4.3).

4.1. D-PATH and Best Path Selection for MAC/IP Advertisement routes

When two (or more) MAC/IP Advertisement routes with the same route key (and same or different RDs) are received, a best path selection algorithm is used to select and install only one route. This section summarizes the best path selection for MAC/IP Advertisement routes, which extends the rules in [I-D.ietf-bess-rfc7432bis] by including D-PATH in the tie-breaking algorithm. While the algorithm may be implemented in different ways, the selection result SHOULD be the same as the result of the rules that follow.

The tie-breaking algorithm begins by considering all EVPN MAC/IP Advertisement routes equally preferable routes to the same destination, and then selects routes to be removed from consideration. The process terminates as soon as only one route remains in consideration.
1. When the Default Gateway extended community is present in some of the routes, the MAC/IP Advertisement routes without the Default Gateway indication are removed from consideration, as defined in [I-D.ietf-bess-rfc7432bis].

2. Then the routes with the Static bit set in the MAC Mobility extended community are preferred, and the routes without the Static bit set are removed from consideration, as defined in [I-D.ietf-bess-rfc7432bis]. Note that this rule does not apply to routes with the Default Gateway extended community, since these routes SHALL NOT convey the MAC Mobility extended community [I-D.ietf-bess-rfc7432bis]. Therefore if two or more routes with the Default Gateway extended community remain in consideration, the selection process skips this step.

3. Then the routes with the highest MAC Mobility Sequence number are preferred, hence the routes that are not tied for having the highest Sequence number are removed from consideration, as defined in [I-D.ietf-bess-rfc7432bis]. If two or more routes with the Default Gateway extended community remain in consideration, the selection process skips this step (for the same reason as in step 2).

4. Then routes with the highest Local Preference are preferred, hence routes that are not tied for having the highest Local Preference are removed from consideration, as defined in [RFC4271].

5. Then routes with the shortest D-PATH are preferred, hence routes not tied for the shortest D-PATH are removed. Routes without D-PATH are considered zero-length D-PATH.

6. Then routes with the numerically lowest left-most Domain-ID are preferred, hence routes not tied for the numerically lowest left-most Domain-ID are removed from consideration.

7. If the steps above do not produce a single route, the rest of the rules after the highest Local Preference in [RFC4271] apply after step 6.

The above selection criteria is followed irrespective of the ESI value in the routes. EVPN Multi-Homing procedures for Aliasing or Backup paths in [I-D.ietf-bess-rfc7432bis] are applied to the selected MAC/IP Advertisement route.
4.2. D-PATH and Best Path Selection for Ethernet A-D per EVI routes

When two (or more) EVVPN A-D per EVI routes with the same route key (and same or different RDs) are received for a VPWS, a best path selection algorithm is used. The selection algorithm follows the same steps as in Section 4.1 except for steps 1-3 which do not apply since the Default Gateway and MAC Mobility extended community are irrelevant to the EVVPN A-D per EVI routes.

The above selection is followed for A-D per EVI routes with ESI=0. For non-zero ESI routes, the EVVPN Multi-Homing procedures in [RFC8214] for Aliasing and Backup path are followed to select the routes (P and B flags are considered for the selection of the routes when sending traffic to a remote Ethernet Segment). If the mentioned Multi-Homing procedures do not produce a single route to each of the remote PEs attached to the same ES, steps 4-7 in Section 4.1 are followed.

4.3. D-PATH and Best Path Selection for Inclusive Multicast Ethernet Tag routes

When two (or more) EVVPN Inclusive Multicast Ethernet Tag routes with the same route key (and same or different RDs) are received for a MAC-VRF, a best path selection algorithm is used. The selection algorithm follows the same steps as in Section 4.1 except for steps 1-3 which do not apply.

4.4. Loop Detection

An EVVPN route received by a PE with a D-PATH attribute that contains one or more of its locally associated Domain-IDs for the MAC-VRF or VPWS instance is considered to be a looped route. A looped route MUST NOT be redistributed to a different domain and SHOULD be flagged as "looped".

EVVPN A-D per EVI looped routes and Inclusive Multicast Ethernet Tag looped routes MUST NOT be installed, where "install" means "create forwarding state" in this document. An EVVPN MAC/IP Advertisement looped route MAY be installed if selected as the best route.

For instance, in the example of Figure 3, assuming PE1 advertises M1’s MAC/IP and does not add the D-PATH attribute, the EVVPN Domain Gateway GW1 receives two MAC/IP Advertisement routes for M1’s MAC/IP:

* A MAC/IP Advertisement route with next-hop PE1 and no D-PATH.
* A MAC/IP Advertisement route with next-hop GW2 and
  D-PATH={length=1,<6500:1:EVPN>}, assuming that the Domain-ID for
  EVPN Domain 1 is 6500:1.

In this case, EVPN Domain Gateway GW1 flags the MAC/IP Advertisement
route with D-PATH as "looped", and does not install the MAC in the BT
of the MAC-VRF and does not redistribute the route back to EVPN
Domain 1 (since the route includes one of GW1's Domain-IDs). In case
the MAC/IP Advertisement route with next-hop PE1 is withdrawn, GW1
may install the route with next-hop GW2 and D-PATH <6500:1:EVPN>;
this may help speed up convergence in case of failures.

4.5. Error Handling

The error handling for the D-PATH attribute is described in
[I-D.ietf-bess-evpn-ipvpn-interworking]. This document extends the
use of D-PATH to non-ISF EVPN routes.

5. Use-Case Examples

This section illustrates the use of D-PATH in EVPN routes with
examples.

Figure 4 and Figure 5 illustrate an integrated interconnect solution
for an EVPN BD, as described in section 4.4 and section 4.6 of
[RFC9014]. GW1 and GW2 are EVPN Domain Gateways connecting two EVPN
Domains identified by D-PATH domain {1:1:EVPN} and {1:2:EVPN},
respectively. Received Ethernet A-D routes, ES routes, and Inclusive
Multicast routes from the routers in one EVPN Domain are consumed and
processed by GW1 and GW2, but not redistributed to the other EVPN
Domain. However, MAC/IP Advertisement routes received by GW1 and GW2
in one EVPN Domain are processed and, if installed, redistributed
into the other EVPN Domain.
Consider the example of Figure 4, where PE1 advertises a MAC/IP Advertisement route for M1/IP1. The route is processed and installed by GW1 and GW2 in BD1, and both redistribute the routes into the EVPN Domain-2. By using D-PATH in GW1 and GW2, when the route is received on PE2, PE2 has the visibility of the EVPN Domains through which the route has gone, and can also use the D-PATH for best path selection. In addition, GW1 and GW2 can compare the D-PATH of the incoming routes with their local list of EVPN Domain-IDs, and detect looped routes if any of the local EVPN Domain-IDs matches a domain in the received D-PATH. This procedure prevents the redistribution of the route back into EVPN Domain-1. For example, when GW1 receives the MAC/IP Advertisement route for M1/IP1 with D-PATH <1:1:EVPN>, GW1 identifies the route as looped and it does not redistribute it back to Domain-1. The M1/IP1 route with Next Hop PE1 is installed. If M1/IP1 with Next Hop PE1 is withdrawn, GW1 MAY install the route M1/IP1 with Next Hop GW2, as specified in Section 4.4.

The example of Figure 5 illustrates how GW1 and GW2 can also have local ACs in BD1 and learn local MAC (or MAC/IP) addresses from devices connected to the ACs.
Assuming GW2 learns M3/IP3 via local AC, GW2 advertises a MAC/IP Advertisement route for M3/IP3 into each of the EVPN Domains that GW2 is connected to. As described in Section 4, GW2 can advertise these two MAC/IP Advertisement routes with a configured EVPN Domain-ID for local MAC/IPs routes that can be shared with GW1. Consider this EVPN Domain-ID is 1:3 and it is configured on both, GW1 and GW2. When GW2 advertises the route into each EVPN Domain, it adds the D-PATH attribute with a domain {1:3:0}. These routes are flagged by GW1 as "looped" since 1:3 is configured as a local EVPN Domain-ID in GW1. In addition, PE1 and PE2 receive the routes with the D-PATH and they have the visibility of the origin of the routes, in this case local EVPN Domain Gateway routes. This information can be used to influence the best path selection in case of multiple routes for M3/IP3 are received on PE1 or PE2 for BD1.

As an alternative solution to configuring the same EVPN Domain-ID for local routes on both EVPN Domain Gateways, GW2 can be configured with EVPN Domain-ID 1:3 for local routes, and GW1 can use a different EVPN Domain-ID, e.g., 1:4. In this case, GW2 advertises the route for M3/IP3 into each EVPN Domain as before, but now GW1 does not flag the route as "looped" since 1:3 is not on the list of GW1’s local EVPN Domain-IDs. GW1 receives the routes from both EVPN Domains, and GW1 selects the route from e.g., EVPN Domain-1. GW1 then installs the route in its BT and redistributes the route into EVPN Domain-2 with
D-PATH \{1:1:EVVPN, 1:3:0\}. When PE2 receives two routes for M3/IP3, one from GW2 with D-PATH \{1:3:0\} and another from GW1 with D-PATH \{1:1:EVVPN, 1:3:0\}, PE2 uses best path selection and choose to send its traffic to GW2. Also GW2 receives the route for M3/IP3 from GW1 and mark it as "looped" since that route conveys its own EVVPN Domain-IDs 1:1 and 1:3.

In a nutshell, the use of D-PATH in MAC/IP Advertisement routes helps prevent loops and influences the best path selection so that PEs choose the shortest paths to the destination PEs.

6. Security Considerations

Most of the considerations included in [I-D.ietf-bess-evpn-ipvpn-interworking] apply to this document.

7. IANA Considerations

None.

8. Acknowledgments

9. Contributors

10. References

10.1. Normative References


10.2. Informative References


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Secure EVPN MAC Signaling
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Abstract

This specification adds attributes to EVPN to carry IPv6 address metadata learned from RFC 8505 and RFC 8928 so as to maintain a synchronized copy of the 6LoWPAN ND registrar at each EVPN router and perform locally a unicast IPv6 ND service for address lookup and duplicate address detection.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

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1. Introduction

"Registration Extensions for IPv6 over 6LoWPAN Neighbor Discovery" [RFC8505] (ND) provides a zeroconf routing-agnostic Host-to-Router Link-Local interface for Stateful Address Autoconfiguration. "Address-Protected Neighbor Discovery for Low-Power and Lossy Networks" [RFC8928] (AP-ND) adds a zeroconf anti-theft protection that protects the ownership of the autoconfigured address with autoconfigured proof of ownership called a Registration Ownership Verifier (ROVR).

[RFC8505] enables the host to claim an IPv6 address and obtain reachability services for that address. It is already used to inject host routes in RPL [RFC9010] and RIFT "Routing in Fat Trees" [RIFT], and to maintain a proxy-ND state in a backbone router [RFC8929]; this specification extends its applicability to the case of Ethernet Virtual Private Network (EVPN).

[RFC8505] specifies a unicast address registration mechanism that enables the host called a 6LoWPAN Node (6LN) to install a ND binding state in the 6LoWPAN Router (6LR) that can serve as Neighbor Cache Entry (NCE), though it is not operated as a cache. The protocol provides the means to reject the registration in case of address duplication. It also enables to discriminate mobility from multihoming. [RFC8928] adds the capability to verify the ownership of the address and prevent an attacker from stealing and/or impersonating an address.

[RFC8505] defines the 6LoWPAN Border Router (6LBR) as an abstract address registrar that provides authoritative service for Address Registration and duplicate detection. The 6LBR stores address metadata that is obtained during the Address Registration, including an owner ID and a sequence counter. As part of the process of a new Address Registration, the 6LR queries the 6LBR for existing metadata related to the address being registered. This enables in particular to detect a duplication and reject the registration. This specification extends the 6LBR abstract data model to store the Link Layer Address (LLA) of the Registering Node. This enables the 6LBR to perform locally, and using unicast communication, the IPv6 ND services of address lookup and duplicate address detection.

The [RFC8505] address registrar can be centralized, but it can also be distributed and maintained synchronized using a routing protocol. This specification adds attributes to EVPN to carry the IPv6 address metadata learned from [RFC8505] so as to maintain a synchronized copy of the 6LBR abstract data at each EVPN router.
2. Terminology

2.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

2.2. Glossary

This document uses the following acronyms:

- 6CIO: Capability Indication Option [RFC7400]
- 6LN: 6LoWPAN Node (the Host) [RFC6775]
- 6LR: 6LoWPAN router (the router) [RFC6775]
- 6LBR: 6LoWPAN Border router [RFC6775]
- AMC: Address Mapping Confirmation [UNICAST-LOOKUP]
- AMR: Address Mapping Request [UNICAST-LOOKUP]
- ARO: Address Registration Option [RFC6775]
- CIP0: Crypto-ID Parameters Option
- DAD: Duplicate Address Detection [RFC4862]
- ICMPv6: Internet Control Message Protocol for IPv6
- DAC: Duplicate Address Confirmation [RFC6775]
- DAR: Duplicate Address Request [RFC6775]
- EDAC: Extended Duplicate Address Confirmation [RFC8505]
- EDAR: Extended Duplicate Address Request [RFC8505]
- EARO: Extended Address Registration Option [RFC8505]
- EVPVN: Ethernet VPN [RFC7432]
- LLA: Link-Layer Address (the MAC address on Ethernet)
- LLN: Low-Power and Lossy Network [RFC6550]
- NA: Neighbor Advertisement [RFC4861]
- NCE: Neighbor Cache Entry [RFC4861]
- ND: Neighbor Discovery [RFC4861]
- NDPSO: Neighbor Discovery Protocol Signature Option
- NS: Neighbor Solicitation [RFC4861]
- RA: Router Advertisement [RFC4861]
- ROVR: Registration Ownership Verifier [RFC8505]
- TID: Transaction ID (a sequence counter in the EARO) [RFC8505]
- SLAAC: Stateless Address Autoconfiguration [RFC4862]
- SLLAO: Source Link-Layer Address Option [RFC4861]
- TLLAO: Target Link-Layer Address Option [RFC4861]
- ROVR MAC: MAC obtained from a host meeting requirements in Section 5
- Validated ROVR MAC: ROVR MAC validated by procedures specified in [RFC8928]
- ROVR Node: EVPN node capable of advertising ROVR MACs
- non-ROVR Node: EVPN node not supporting extensions defined in this
2.3. References

This document uses the terms Clos fabric and Fat Tree interchangeably, to refer to a folded spine-and-leaf topology as defined in the terminology section of "RIFT: Routing in Fat Trees" [RIFT].

The term "leaf" represents the access switch that connects the servers to the Fat Tree. The leaf is typically a Top-of-Rack (ToR) switch.

This specification uses the terms 6LN, 6LR and 6LBR to refer specifically to nodes that implement the said roles in [RFC8505] and does not expect other functionality such as 6LoWPAN Header Compression:

* In the context of this document, the 6LN is a server that advertises an address mapping using [RFC8505], and optionally protects its ownership with [RFC8928].

* The 6LR and 6LBR function are collapsed at the leaf and its state is synchronized with that of the EVPN functional support using an internal interface that is out of scope. That interface could be "pull" meaning that the 6LBR fetches the EVPN information when it needs it, or "push", meaning that any information that EVPN distributes is immediately fed in all the 6LBRs in all the leaves. Note that this is pure control plane and is not subject to abbreviating optimization as the FIB may be.

In this document, readers will encounter terms and concepts that are discussed in the following documents:

EVPN: "BGP MPLS-Based Ethernet VPN" [RFC7432] and "Network Virtualization Overlay Solution" [RFC8365],

Classical IPv6 ND: "Neighbor Discovery for IP version 6" [RFC4861] and "IPv6 Stateless Address Autoconfiguration" [RFC4862],

6LoWPAN ND: Neighbor Discovery Optimization for Low-Power and Lossy Networks [RFC6775], "Registration Extensions for 6LoWPAN Neighbor Discovery" [RFC8505], "Address Protected Neighbor Discovery for Low-power and Lossy Networks" [RFC8928], and "IPv6 Backbone Router" [RFC8929].
3. 6LoWPAN Neighbor Discovery

6LoWPAN Neighbor Discovery defines a stateful address autoconfiguration mechanism for IPv6. 6LoWPAN ND enables to divorce the L3 abstractions for link and subnet from the characteristics of the L2 link and broadcast domain. It is applicable beyond its original field of IoT to any environment where the broadcast nature of the underlaying network should not be exploited, e.g., in the case of a wireless link where broadcast uses an excessive amount of spectrum, and a distributed cloud, where it may span too widely.

In contrast to Stateless Address Autoconfiguration (SLAAC) [RFC4862] which relies on broadcast for duplicate address detection (DAD) and address lookup, 6LoWPAN ND installs and maintains a state in the neighbors for the duration of their interaction. Though it is also called a Neighbor Cache Entry (BCE) in [RFC6775], and in contrast with the the BCE in SLAAC, that state is not a cache that can be casually flushed and rebuilt. It must be installed proactively and refreshed periodically to maintain the connectivity and enable unicast-only operations.

This section goes through the 6LoWPAN ND network abstractions and mechanisms that this specification leverages, as a non-normative reference to the reader. The relevant normative text is to be found in [RFC6775], [RFC8505], and [RFC8928].

3.1. IPv6 Interface, Link, and Subnet

The typical abstraction for an IP Link with 6LoWPAN ND is a logical point-to-point (P2P) link between a node (a host or a router) and a router, regardless of the physical medium between the node and the router, which may or may not be shared with other nodes.

A Subnet is deployed over a mesh of nodes connected with those logical P2P Links, where routers form a connected dominating set as represented in Figure 1; the resulting aggregate is called a multilink subnet (MLSN). An MLSN may be only partially meshed, and the underlaying network is not expected to provide a multicast or a broadcast service across the subnet.
Consequently, the subnet model is not-on-link, meaning that the any-to-any connectivity across the subnet is ensured through L3 operations (routing or proxy) as opposed to transitive (any-to-any) reachability from L2. It also means that hosts do not lookup other nodes using IPv6 Neighbor Discovery but forward all their traffic via their connected routers. Which in turn means that only routers need to be discovered, which is done by sending Router Advertisement (RA) messages to all directly reachable nodes in the subnet, e.g., using a radio broadcast.

As illustrated in Figure 2, an IP interface bundles multiple sub-interfaces to connect the IP links between this node and peers in the same subnet, which is known as a point-to-multipoint (P2MP) interface.
In the case of a 6LoWPAN radio, the IP Interface may be physical, and the P2P IP links are virtual based on discovered neighbor routers; the same model can apply when the node is connected via a switch to one or more routers.

In the case of a multihomed NIC card in a datacenter, the NIC is connected to several Top-of-Rack (ToR) switches acting as leaves in the fabric, over as many Ethernet physical interfaces. If the NIC provides a L2 virtual switch, then the stack can apply the same model as above, modeling the virtual port to the virtual switch as a P2MP interface.
On the other hand, if the NIC provides a virtual router, then Ethernet ports are L3 ports and the physical link to the leaf is modeled as P2P. To form the P2MP interface, the router bundles (aggregates) the physical interfaces as the sub-interfaces of a single logical P2MP Link, as shown in Figure 3.

```
+---------------------
| NIC Aggregate interface
| -----------------------
| virtual MAC Address
| IPv6 global addresses

+------------------------
| Ethernet Port 1
| ------------------
| Physical MAC Address
| IPv6 Link-Local address(es)
| Leaf MAC Address
| Leaf IP addresses
+------------------------

+------------------------
| Ethernet Port 2
| ------------------
| Physical MAC Address
| IPv6 Link-Local address(es)
| Leaf MAC Address
| Leaf IP addresses
+------------------------

....

+---------------------
```

Figure 3: Logical P2MP Interface

To conserve the same model, it makes sense to configure the same (virtual) MAC address on all the physical interfaces, and use it for the purpose of IPv6 ND. In that case, the same MAC address is exposed as Link-layer Address (LLA) to both leaves for the NIC IP addresses, and the IPv6 address still appears as unicast. Note that the Link-Local addresses used to register the global IPv6 addresses to the leaf may be different but that does not affect the exposed mapping.
When that is not possible, then the same IP address is advertised with the physical MAC address of each port as the LLA over that port. In that case, the global IPv6 address appears as anycast, and SHOULD be advertised as such, more in Section 3.3.5.

3.2. RFC 6775 Address Registration

The classical "IPv6 Neighbor Discovery (IPv6 ND) Protocol" [RFC4861] [RFC4862] was defined for serial links and transit media such as Ethernet. It is a reactive protocol that relies heavily on multicast operations for Address Discovery (aka Lookup) and Duplicate Address Detection (DAD).

"Neighbor Discovery Optimizations for 6LoWPAN networks" [RFC6775] adapts IPv6 ND for operations over energy-constrained LLNs. The main functions of [RFC6775] are to proactively establish the Neighbor Cache Entry (NCE) in the 6LR and to prevent address duplication. To that effect, [RFC6775] introduces a new unicast Address Registration mechanism that contributes to reducing the use of multicast Registration messages compared to the classical IPv6 ND protocol.

[RFC6775] defines a new Address Registration Option (ARO) that is carried in the unicast Neighbor Solicitation (NS) and Neighbor Advertisement (NA) messages between the 6LoWPAN Node (6LN) and the 6LoWPAN router (6LR). It also defines the Duplicate Address Request (DAR) and Duplicate Address Confirmation (DAC) messages between the 6LR and the 6LBR. In a Low-Power and Lossy Network (LLN), the 6LBR is the central repository of all the Registered Addresses in its domain and the authoritative source of truth for uniqueness and ownership.

3.3. RFC 8505 Extended Address Registration

"Registration Extensions for 6LoWPAN Neighbor Discovery" [RFC8505] updates RFC 6775 into a generic Address Registration mechanism that can be used to access services such as routing and ND proxy. To that effect, [RFC8505] defines the Extended Address Registration Option (EARO), shown in Figure 4:
3.3.1. R Flag

[RFC8505] introduces the R Flag in the EARO. The Registering Node sets the R Flag to indicate whether the 6LR should ensure reachability for the Registered Address. If the R Flag is set to 0, then the Registering Node handles the reachability of the Registered Address by other means. In an EVPN network, this means that either it is a RAN that injects the route by itself or that it uses another EVPN router for reachability services.

This document specifies how the R Flag is used in the context of EVPN. An EVPN Host that implements the 6LN functionality from [RFC8505] requires reachability services for an IPv6 address if and only if it sets the R Flag in the NS(EARO) used to register the address to a 6LR acting as an EVPN border router. Upon receiving the NS(EARO), the EVPN router generates a BGP advertisement for the Registered Address if and only if the R flag is set to 1.

[RFC9010] specifies that the 'R' flag is set in the responded NA messages if and only if the route was installed. This specification echoes that behavior.

3.3.2. TID, "I" Field and Opaque Fields

When the T Flag is set to 1, the EARO includes a sequence counter called Transaction ID (TID), that is needed to format the MAC Mobility Extended Community. This is the reason why the support of [RFC8505] by the Host, as opposed to only [RFC6775], is a prerequisite for this specification; this requirement is fully explained in Section 5.1. The EARO also transports an Opaque field and an associated "I" field that describes what the Opaque field transports and how to use it.
This document specifies the use of the "I" field and the Opaque field by a Host.

3.3.3. Status

The values of the EARO status are maintained by IANA in the Address Registration Option Status Values subregistry [IANA-EARO-STATUS] of the Internet Control Message Protocol version 6 (ICMPv6) Parameters registry.

[RFC6775] and [RFC8505] defined the original values whereas [RFC9010] reduced range to 64 values and reformatted the octet field to enable to transport an external error, e.g., coming from a routing protocol.

This specification uses the format expressed in [RFC9010]. The value of 0 denotes an unqualified success, 1 indicates an address duplication, 3 a TID value that is outdated, and 4 is used in an asynchronous NA to indicate that 6LN should remove that address and possibly form new ones.

3.3.4. Route Ownership Verifier

Section 5.3 of [RFC8505] introduces the Registration Ownership Verifier (ROVR) field of variable length from 64 to 256 bits. The ROVR is a replacement of the EUI-64 in the ARO [RFC6775] that was used to identify uniquely an Address Registration with the Link-Layer address of the owner but provided no protection against spoofing.

"Address Protected Neighbor Discovery for Low-power and Lossy Networks" [RFC8928] leverages the ROVR field as a cryptographic proof of ownership to prevent a rogue third party from registering an address that is already owned. The use of ROVR field enables the 6LR to block traffic that is not sourced at an owned address.

This specification does not address how the protection by [RFC8928] could be extended for use in EVPN. On the other hand, it adds the ROVR to the BGP advertisement to share the state with the other routers via the Reflector (see Section 6.1), which means that the routers that are aware of the Host route are also aware of the ROVR associated to the Target Address, whether it is cryptographic and should be verified.
3.3.5. Anycast and Multicast Addresses

"IPv6 Neighbor Discovery Multicast Address Registration" [I-D.thubert-6lo-multicast-registration] updates [RFC8505] to enable the address registration of IPv6 anycast and multicast addresses. From the host perspective, the registration is very similar to that of unicast addresses, but for a flag in the EARO that signals that the address is multicast or anycast.

This procedure can be used as a replacement to "Multicast Listener Discovery Version 2 (MLDv2) for IPv6" [RFC3810] for source-independant multicast operation. As for unicast, the method saves the need for the host to listen to pollings from the router, and allows the host to sleep for periods of time.

3.4. RFC 8505 Extended DAR/DAC

[RFC8505] updates the DAR/DAC messages to EDAR/EDAC messages to carry the ROVR field. The EDAR/EDAC exchange takes place between the 6LR to which the node registers an address, and the abstract 6LBR that stores the reference value for the ROVR and the TID associated to that address. It is triggered by an NS(EARO) message from a 6LN to the 6LR, to create, refresh, compare and delete the corresponding state in the 6LBR.

In the status returned with the EDAC message, the 6LBR indicates if the registration is accepted, should be challenged, or is duplicate. The status of 0 (success) indicates that the address is either new or that the current registration matches, and in particular that the ROVR at the 6LBR and the one in the EDAR message are identical.
The EDAR/EDAC exchange is protected by the retry mechanism specified in Section 8.2.6 of [RFC6775], though in a data center, a duration significantly shorter than the default value of the Retransmission Timer [RFC4861] of 1 second may be sufficient to cover the round-trip delay between the 6R and the 6LBR.

With this specification, the 6LBR is distributed across the leaves, and all the leaves where an address is currently registered maintain a full 6LBR state for the address, aka local state in the following text. The specification leverages the EDAR/EDAC exchange to ensure that a leaf (acting as a 6LR) that needs to create a 6LBR state for a new registration has the same value for the ROVR as any 6LBR already serving that address on another leaf. At the same time, the specification avoids placing full ROVR information in BGP so 1) it is not observable by a potential attacker and 2) the new attributes remain reasonably small.

3.5. RFC 7400 Capability Indication Option

"6LoWPAN-GHC: Generic Header Compression for IPv6 over Low-Power Wireless Personal Area Networks (6LoWPANs)" [RFC7400] defines the 6LoWPAN Capability Indication Option (6CIO) that enables a node to expose its capabilities in router Advertisement (RA) messages.

[RFC8505] defines a number of bits in the 6CIO, in particular:

L: Node is a 6LR.
E: Node is an IPv6 ND Registrar -- i.e., it supports registrations
based on EARO.

P: Node is a Routing Registrar, -- i.e., an IPv6 ND Registrar that also provides reachability services for the Registered Address.

```
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|     Type      |   Length = 1  |     Reserved      |D|L|B|P|E|G|
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|                           Reserved                            |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
+-----------------------------------------------+---
|                                      Reserved |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
```

Figure 6: 6CIO flags

A 6LR that provides reachability services for a Host in an EVPN network as specified in this document includes a 6CIO in its RA messages and set the L, P and E flags to 1 as prescribed by [RFC8505].

4. Extending 6LoWPAN ND

4.1. Use of the R flag in NA

This document extends [RFC8928] and [RFC8505] as follows

This document also updates the behavior of a 6LR acting as EVPN router and of a 6LN acting as Host in the 6LoWPAN ND Address Registration as follows:

* The use of the R Flag is extended to the NA(EARO) to confirm whether the route was installed.

4.2. Distributing the 6LBR

This specification enables to distribute the 6LBR at the edge of the EVPN network and collapse the 6LBR function with that of the EVPN support. In that model, the EVPN to 6LBR interaction becomes an internal interface, where each side informs the other in case of new information concerning an IP to Link-Layer Address (LLA) mapping. Since this is an internal interface, this specification makes no assumption on whether the 6LBR stores its own representation of the full EVPN state, which means that the EVPN support informs the 6LBR in case of any change on the EVPN side (this is called the push model, see Figure 13), or if the 6LBR queries the EVPN support when it does not have a mapping to satisfy a request (pull model, see Figure 12).
This specification leverages [RFC8929] that augments the abstract data model of the 6LBR to store the LLA associated with the registered address. Based on that additional state, the 6LBR in a leaf can communicate the mapping to the collocated EVPN function and respond to unicast address mapping lookups from the server side.

In an environment where the server ranges from a classical host to a more complex platform that runs a collection of virtual hosts interconnected by a virtual switch, but where the host-to-leaf interface remains at layer 2, the 6LR and the 6LBR functions can be collapsed in the leaf. The 6LR to 6LBR interaction also becomes an internal interface, and there is no need for EDAR/EDAC messages.

In that case, the MAC address associated to the Registered Address is indicated in the Target Link-Layer Address Option (TLLAO) in the NS message used for the registration, as shown in Figure 7. In the case of a pull model, if the 6LBR does not have a local state for the mapping, it queries the EVPN support to obtain the EVPN state if any. If a mapping is known then the 6LR/6LBR evaluates the registration for address duplication and other possible issues per [RFC8505]. Else (this is for a new mapping), if the registration is accepted, then the 6LBR notifies the EVPN support to inject a route type 2 in the fabric.
In another type of deployment, the 6LR may be a virtual router in the server whereas the 6LBR runs in the leaf node. To address that case, the EDAR/EDAC may be used to communicate as shown in figure 5 of [RFC8505]. This draft leverages the capability to insert IPv6 ND options in the EDAR and EDAC messages introduced in [RFC8929] to place a TLLAO that carries the MAC address associated to the Registered address in the EDAR and EDAC messages as shown in Figure 8:
[RFC8505] updates the DAR/DAC messages into the Extended DAR/DAC to carry the ROVR field. With this specification, the abstract 6LBR is distributed in all the Leaf nodes and synchronized with EVPN. When a server successfully registers an address to a leaf, the 6LR on that leaf becomes 6LBR for that address. It stores the full state for that address including the ROVR and the TID. When the address registration moves to another leaf, an EDAR/EDAC flow between the 6LR in the new leaf and the 6LBR in the old leaf confirms that the ROVR in the NS(EARO) received at the new leaf is correct, in which case the 6LR in the new leaf becomes 6LBR.

When the address is already registered to the local leaf, the EDAR/EDAC exchange is either local between a virtual router in the server and the leaf, or internal to the leaf between a collapsed 6LR and

Figure 8: leveraging EDAR
6LBR. Based on its local state, the 6LBR in the leaf checks whether the proposed address/route is new and legit, and can reject it otherwise.

It results that duplicate addresses and address impersonation attacks can be filtered at the level of IPv6 ND by the 6LBR before the information reaches EVPN.

4.3. Unicast Address Lookup with the 6LBR

A classical IPv6 ND stack in the server that treats the subnet prefix as on-link (more in section 4.6.2. of [RFC4861]), will resolve an unknown LLA mapping with a multicast NS(lookup) message addressed to the solicited node multicast address (SNMA) associated with the destination address being resolved. The RECOMMENDED operation in that case is for the 6LBR that has a mapping state to forward the packet as a unicast MAC to the LLA that is stored for the IPv6 address as expected by [RFC6085]. The actual owner of the address can then answer unicast with a NA message, setting the override (O) bit to 1, as shown in Figure 9.

Figure 9: Forwarding legacy NS (Lookup)
Section 3.1. of [RFC8929] adds the capability to insert IPv6 ND options in the EDAR and EDAC messages. This enables the 6LBR to store the link-layer address associated with the Registered Address and to serve as a mapping server. [UNICAST-LOOKUP] leverages that state to define a new unicast address lookup operation, extending the EDAR and EDAC messages as the Address Mapping Request (AMR) and Confirmation (AMC) with a different Code Prefix [RFC8505].

In that model, the router advertises the subnet prefix as not on-link by setting the L flag to 0 in the Prefix Information Option (PIO), more in section 4.6.2. of [RFC4861]. The expected behavior is that the host that communicates with a peer in the same subnet refrains from resolving the address mapping and passes the packets directly to the router.

In the case where the router is a virtual 6LR running in the server, and the source and destination are in the same subnet served by EVPN, the router then resolves the address mapping on behalf of the host. To that effect, the router sends a unicast AMR message to the 6LBR. The message contains the SLLAO of the router to which the 6LBR will reply. If the binding is found, the 6LBR replies with an AMC message that contains the TLLOA with the requested MAC address, as shown in Figure 10.
If it is not found, [UNICAST-LOOKUP] provides the capability to indicate immediately that the mapping is not known with a "not found" status in the AMC, as opposed to waiting for an NS(lookup) and retries to time out per [RFC4861].

In a fully stateful subnet where all nodes register all their addresses with [RFC8505], this means that the looked up address is not present in the network; in that case the packet is dropped and an ICMP error type 1 "Destination Unreachable" code 3 "Address unreachable" [RFC4443] is returned as shown in Figure 11.
Figure 11: Unicast Lookup failure

Note that the figures above make no assumption on the pull vs. push model. In the case of pull model, the 6LBR queries the EVPN support when it does not have the mapping information to satisfy a request. Figure 12 illustrates a successful pull model lookup flow, when the route type 2 for the mapping is already known on the EVPN side.
In the case of push model, the EVPN support synchronizes its state upon a route type 2 with the 6LBR, and the 6LBR maintains an abstract data structure for all information known to EVPN. This way, the 6LBR already has the mapping information to satisfy any request for an existing mapping and it can answer right away. Figure 13 illustrates a successful push model lookup flow, when the 6LBR is already in possession of the mapping.
In a mixed environment, a lookup failure (the mapping is not found though the address is present in the network) may be caused by a legacy node that was not discovered (aka a silent node). In that case, it is an administrative decision for the 6LR to broadcast an NS(lookup) or to return an error as shown in Figure 11.
5. Requirements on the EVPN-Unaware Host

This document describes how EVPN routing can be extended to reach a Host. This section specifies the minimal EVPN-independent functionality that the Host needs to implement to obtain routing services for its addresses.

5.1. Support of 6LoWPAN ND

A host sees a prefix as not on-link (e.g., it learned that prefix in a PIO in a RA with the L flag not set) should not attempt to resolve an address within that prefix using a multicast NS(lookup). Instead, it must pass its packets to a router, preferably one that advertises that prefix in a PIO; it must register the address that it uses as source to that router to enable source address validation using [RFC8505]. It is recommended that the Host also implements [RFC8928] to prove its ownership of its addresses.

The Host is expected to request routing services from a router only if that router originates RA messages with a 6CIO that has the L, P, and E flags all set to 1 as discussed in Section 3.5, unless configured to do so. To obtain routing services for one of its addresses, the host must register the address to a router that advertises the prefix, setting the "R" and "T" flags in the EARO to 1 as discussed in Section 3.3.1 and Section 3.3.2, respectively.

This document echoes the behavior specified in [RFC9010] whereby, when the R Flag set to 1 in a NS(EARO) is not echoed in the NA(EARO), the host must understand that the route injection failed, and if the R flag is reset later in an asynchronous NA(EARO), the host must understand that routing service has failed.

The host may attach to multiple 6LRs and is expected to prefer those that provide routing services. The abstract model for this is a P2MP interface that wraps together as many P2P IP Links the host has adjacencies to 6LRs over that interface. The IPv6 address and the subnet are associated to that interface. The interface may be virtual and it may bundle multiple physical Ethernet interfaces that connect to the individual 6LRs over point to point wires, possibly via a software switch. It can also be associated to one physical interface to an external switch, either way the PI Links can be associated to sub-interface of the interface.

The Host needs to register to all the 6LRs from which it desires routing services. The multiple Address Registrations to several 6LRs should be performed in a rapid sequence, using the same EARO for the same Address. Gaps between the Address Registrations will invalidate some of the routes till the Address Registration finally shows on
those routes. The routers recognize the same (ROVR, TID) as the signal of a multihomed address and maintain all the routes. In the case of EVPN, the Ethernet Segment must also be the same. The flow for a successful multihomed registration is illustrated in Figure 17.

[RFC8505] introduces error Status values in the NA(EARO) which can be received synchronously upon an NS(EARO) or asynchronously. The Host needs to support both cases and refrain from using the address when the Status value indicates a rejection.

This specification can be used to register Anycast and Multicast IPv6 addresses as discussed in Section 3.3.5 and replace MLDv2. To benefit from that capability, both the host and the 6LR MUST support the "A" and "M" flags that indicate Anycast and Multicast Addresses respectively. Those flags are defined in [I-D.thubert-6lo-multicast-registration] for use in the EARO and in the EDAR and EDAC messages.

6. Enhancements to EVPN

This section addresses the necessary changes to EVPN formats and behavior to support address registration security per [RFC8928] and mobility per [RFC8505] while retaining interoperability with traditional nodes. Basically the MAC Mobility Extended Community [RFC7432] and the ARP/ND Extended Community [RFC9047] are extended to advertise the sequencing and ownership validation information received in the EARO. With 6LRs injecting not only MACs via packet sources and TLLAO options but also ROVR into MAC Mobility and ARP/ND Extended Community, their semantics must be extended. Specifically following issues have to be addressed:

* The ROVR extends the semantics of the type-2 MAC advertisement via changes in ARP/ND and MAC Mobility Extended Community in the sense that the MAC must be aligned with the ROVR and under normal circumstances only the validity of ROVR guarantees that the type-2 MAC can be allocated to the requester. A MAC validated by ROVR should take precedence over MAC addresses allocated without using it given it presents a much more trustworthy topological information (it will be called ROVR MAC in further text). EVPN nodes not supporting extensions introduced by this document will need to be led to believe that a ROVR MAC is to be preferred over any advertisement they see as long a ROVR MAC route is present. The primary key of NRLI is still the (IP, MAC, Ethernet Tag ID) tuple as defined in [RFC7432], Section 7.2 and 7.7. This implies that the same MAC (and consequently ROVR MAC) can be assigned multiple IP addresses with different ROVRs and those represent independent NLRIs.
* The TID field in the EARO is smaller than the mobility sequence number in [RFC7432]. To allow a ROVR MAC mobility to "win" over legacy MACs in every circumstance, signaling must be introduced that enables to distinguish a TID-generated sequence number from a legacy sequence number.

* [RFC8505] supports IP multihoming, but does not differentiate multihoming from anycast or MAC address rotation. If an anycast IP address is registered with a different ROVR it will be rejected as duplicate. If it is registered with a different TID, the older sequence will be withdrawn. So the basic expectation with [RFC8505] is that the advertisement of an anycast address is coordinated, with the same keypair known to all parties, and the same value of the TID used by all nodes (and possibly never increasing), in other words, with no concept of mobility.

* [I-D.thubert-6lo-multicast-registration] adds new flags in the EARO to signal that an address is anycast or multicast, respectively. This specification injects that information in the ARP/ND extended community using matching flags as follows:

  - This specification uses the "O" flag in the ARP/ND extended community to signal that the IP address is anycast, and requires the local 6LBR to ignore the duplication if the same IP address is registered locally, and then to inject the NLRI with the "O" flag set on the ARP/ND Extended Community as well.

  - This specification adds the new "M" flag in the ARP/ND extended community to signal that the IP address is multicast, and requires the local 6LBR to ignore the duplication if the same IP address is registered locally, and then to inject the NLRI with the "M" flag set on the ARP/ND Extended Community as well.

* [RFC8928] needs the full ROVR to validate the address ownership, but the full ROVR can be too large to advertise through BGP. When an IP address is advertised through EVPN, it is REQUIRED that the EVPN Next Hop represents the address of the 6LBR of the leaf where the address was registered as well. This way, if the address is registered later on a second leaf, the 6LR in second leaf can leverage an out-of-band, i.e. via EVPN traffic carrying tunnels, EDAR/EDAC exchange with that 6LBR to validate that the ROVR in the registration is indeed the same. When that is the case, it can continue with the registration procedure and if successful, become a 6LBR for that IP address, either as a mobility event or as a multihomed registration.
* [RFC8928] expects nodes to autoconfigure the keypair that is used to form the ROVR, in which case the IPv6 address can be locally autoconfigured with no central coordination; in that case, the ROVR protects the address ownership and allows the fabric to enforce first-come first-serve and source address validation (SAVI). But it is also possible to provision the ROVR in the 6LBR in advance and later configure the keypair in the node, e.g., in the case of a trusted server. To enable that capability in EVPN, this specification adds a flag (U) to signal that the 6LBR that injects the address in EVPN does not provide reachability to the address. When that flag is set, the value of the TID is ignored in the mobility computation, the mapping to the MAC address is ignored, and the route to the IP address is not injected in the RIB on ROVR nodes. Non-ROVR nodes will consider the node a "honey-pot". Once the address is registered by a 6LN in the network and the according validation with the node advertising the U-bit version of the route is performed, the owner will inject the route without the U-bit. A node advertising the NLRI with U-bit in its ARP/ND Extended Community MUST withdraw the U-bit route once it sees a validated NLRI without the U-bit and it MAY reinject the route with the U-bit once all routes without the U-bit are withdrawn to protect the address again.

6.1. Updated ARP/ND Extended Community

The ARP/ND Extended Community defined in [RFC9047] is a transitive EVPN Extended Community (Type field value of 0x06) with a Sub-Type of 0x08. It is advertised along with EVPN MAC/IP Advertisement routes that carry an IPv4 or IPv6 address. This the ARP/ND Extended Community to transport fields from the EARO is natural since the EARO is an extension to IPv6 ND.

EVPN signaling is not used to carry the full ROVR since without challenge per [RFC8928] they do not represent any difference over using the IP/MAC combination. A Hash of the ROVR is still passed in the ARP/ND Extended Community to immediately detect a duplication, with 2 different values of the hash for the same address. The full ROVR is verified upon a movement or a multi-homed advertisement using an EDAR/EDAC exchange with the 6LBR that advertised the address first.

Additionally, backwards compatibility could not be preserved given comparing routes based on ROVR would present a change in primary key of NLRIs which non-ROVR routers do not implement. An indication from a ROVR node that a MAC has been validated by proof of ownership is enough to convey the necessary information.
ROVR nodes MUST set the "H" flag in the ARP/ND Extended Community to indicate that the advertisement carries a TID and a ROVR Hash in case the host followed the according procedures.

ROVR nodes MUST set the "V" flag if the address assignment passed proof of ownership per [RFC8928]. Such "validated" ROVR MAC addresses will be preferred by ROVR nodes over non-validated ROVR MACs.

In case a ROVR node configures the address as "sticky" (since the sticky bit semantics have been changed to the point a ROVR cannot tell whether address is really sticky unless advertised as such by non-ROVR node) a new "X" flag called "super sticky" is introduced.

```
+---------------+---------------+-------------------+
| Type=0x06     | Sub-Type=0x08 | Flags (2 octet)   |
+---------------+---------------+-------------------+
| TID           | ROVR Hash    |                   |
+---------------+---------------+-------------------+
```

Figure 14: Updated ARP/ND Extended Community

The "I", "O", and "R" flags are defined in [RFC9047]. The following new fields are defined:

U: 1-bit flag. "Unreachable", indicating that the IP address is not reachable via that EVPN next hop, but is advertised for the purpose of protecting the value of the ROVR until a first 6LBR that can reach the address becomes available.

M: 1-bit flag. "Multicast", indicating that the IP address duplication should be ignored. When this bit is set, TID should be ignored in comparison of EVPN advertisements, i.e. all ROVR MACs at same level of validation MUST be considered having same TID.

V: 1-bit flag. "ROVR Validated" indicates that the MAC passed proof of ownership per [RFC8928]. Presence of this bit implies the "R" bit being set regardless of its value.

H: 1-bit flag. "ROVR Capable" indicates that the advertisement is
originated after processing signaling from host meeting the requirements in Section 5, and that the advertised MAC address is a ROVR MAC. This also indicates that the TID and ROVR Hash fields are populated based on information from the most recent EARO [RFC8505] from the host.


ROVR Hash: 16-bits hash of ROVR [RFC8505] from the Registering Node. The Hash is built by XOR’ing ROVR bytes in network order into the least significant byte and rotating the two bytes result after every byte by one bit to the left.

6.2. Updated Mobility Extended Community

This specification extends the MAC Mobility Extended Community to transport the TID instead of increasing the normal sequence number. The TID is placed in the high bits of the sequence number field to "override" any normal MAC advertisement (further considerations will be provided in Section 6.3). This allows to design a solution that, while backwards compatible, allows to introduce ROVR MAC as "more trusted" entities. Figure 15 presents the according extensions that will however necessitate some further explanation.

To introduce a "precedence" of ROVR MACs over normal EVPN MACs ROVR MACs are advertised to look like "sticky" MACs for non-ROVR nodes. As defined in the glossary, for simplicity reasons such nodes will be called non-ROVR nodes vs. ROVR nodes. The "sticky" bit will force non-ROVR nodes to disregard the sequence number and accept any IP/MAC route provided.

```
+-------------------------------------------+-------------------------------------------+--
<p>| Type=0x06       | Sub-Type=0x00 | Flags = 0 |X|S|  Reserved = 0 |</p>
<table>
<thead>
<tr>
<th>-------------------------------------------</th>
<th>-------------------------------------------</th>
</tr>
</thead>
<tbody>
<tr>
<td>-------------------------------------------</td>
<td>-------------------------------------------</td>
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<td>-------------------------------------------</td>
<td>-------------------------------------------</td>
</tr>
<tr>
<td>-------------------------------------------</td>
<td>-------------------------------------------</td>
</tr>
</tbody>
</table>
```

Figure 15: Updating the MAC Mobility Extended Community for ROVR MACs

This specification updates the Sequence Number field defined in [RFC7432]. The field is split in 3 parts, one 8-bit flags field, the TID, and a reserved 2-byte field. The unspecified flags and the reserved fields MUST be set to 0 by the sender and ignored by the receiver.

The "S" flag is defined in [RFC7432]. The following new fields are defined:
X: 1-bit flag. "Super Sticky" indicates that the ROVR MAC is sticky and should follow procedures of sticky per [RFC7432].

T: 1-bit flag. "TID Present", MUST be set to 1 when this specification is used. This ensures that the TID always wins vs. the sequence counter defined in [RFC7432]

TID: The TID copied from the most recent EARO [RFC8505] from the Registering Node.

6.3. Extended ROVR MAC Procedures

In case a non-ROVR node advertises a sticky MAC by setting the "S" bit and a ROVR node sees an ROVR address registration for the same MAC it MUST follow procedures per [RFC7432].

In case a non-ROVR node advertises a sequence number larger than the one generated by TID on a ROVR node, the ROVR node SHOULD advertise a Sequence Number consisting of all bits being set to force a "roll-over" on all nodes and then fall back to advertising the TID generated sequence number again. In case a non-ROVR node persists in increasing the sequence number after that it is indication of violation of [RFC7432] on its part.

A ROVR node advertising a ROVR MAC that has not been validated and receiving same type-2 route that has been validated MUST immediately withdraw its advertisement.

A ROVR node advertising a ROVR MAC and receiving an equivalent ROVR MAC from other node with a higher TID MUST immediately withdraw its advertisement. This will allow the non-ROVR nodes to correctly interpret the sequence as MAC move despite ignoring the sequence number due to presence of "S" bit.

A ROVR node that receives a ROVR MAC with "super sticky" indication and seeing the MAC locally MUST follow analogous procedures to [RFC7432].

Multi-homing a MAC on mix of ROVR and non-ROVR nodes will lead to operational notifications since per [RFC7432] the non-ROVR node will interpret the situation as a sticky MAC that has shown up on its local interface unless an implementation is somewhat clever and understands that the presence of the same ESI on all the routes indicates that this situation does not represent a sticky MAC being moved.
7. Protocol Operations

Following section illustrates several situations and resulting signaling in EVPN from the point of view of a ROVR node.

Figure 16 illustrates the registration flow of a new address protected by [RFC8928]. The ROVR in the EARO is a Crypto-ID that derives from a public address through hashing with some other terms. The router challenges the host with a status of 5 (validation requested).

The host performs the NS again, passing the parameters that enable to build the Crypto-ID in a Crypto-ID Parameters Option (CIPO), and signing that set of parameters together with a pair of Nonce values, one from each side, in a resulting Neighbor Discovery Protocol Signature Option (NDPSO). The 6LR first verifies that the Crypto-ID can be rebuilt based on the public key, then verifies that the signature in the NDPSO was effectively performed with the associated public key. When that is the case, the registration flow can continue, else the registration is rejected with a status of 10 (Validation Failed) in the NA(EARO).

With this specification, the 6LBR communicates internally with the collocated EVPN router to inject the route in EVPN. Since the [RFC8928] validation was performed, the V flag is set. Once this is done, the local 6LBR installs a local state associated to the NCE and becomes owner of the registration, whereas the remote leaves optionally install a remote state for the address with the indication of the 6LBR that owns the registration. The local 6LBR MUST be signalled as EVPN Next Hop for the route.
Figure 16: Host Registration

Figure 17 presents the same flow but for a multihomed address; here and in the following flows, the proof of ownership section is not shown, but its use is RECOMMENDED. The interesting piece is that
when the node registers to the second 6LBR, that second 6LBR finds
that there is a first 6LBR that already owns the registration. Using
and EDAR / EDAC flow, the second 6LBR validates that the ROVR and TID
are identical, in which case it accepts the registration and becomes
another 6LBR owner of the registration. The result is that the 2
6LBRs are synchronized and any of the 2 can now be used, e.g., if the
address is registered a third time.
The registration is associated with a lifetime, and it must be renewed with an incremented TID. The new TID is propagated in EVPN as illustrated in Figure 18.

Figure 18: Host Registration Renewal

Figure 19 illustrates the case where a second host registers the same address, creating a potential address duplication situation. In most cases, the ROVR hash will be different, and the local 6LBR can reject the registration in a status of 1 (duplicate) right away.
Figure 20 illustrates the case of an address duplication situation where by chance, the ROVR hashes are the same. In that case, the local 6LR checks with the 6LBR that owns the registration using an EDAR/EDAC message exchange. As opposed to the ROVR hash, the full ROVRs do not collide and the registration is also rejected.
Figure 20: Duplicate Addresses, ROVR Hash Collision

Figure 21 shows a rare case where the registration has already moved elsewhere with an incremented TID when the local registration is received after being delayed in the network. In that case, the registration is rejected with a status of 3 (moved).
Address move differs from multi-homing by the ESI being different as visualized by Figure 22. In case of different ESI BGP signalling happens immediately, in case of multi-homing we can reasonably expect for the signalling to catch up on the other leg with a new, higher TID. However, since ESI matches TID doesn't matter strictly speaking and the new remote state can be installed as is. However, if 6LN is not refreshing it registration we can expect elapsed lifetime to create scenario Figure 25 over time.
The host that registered the address may cancel the registration at any time, e.g., if the address is removed from its own interface. This is done by registering with a lifetime of 0 as shown in Figure 23. The Leaf may respond with a status of 0 to indicate success, but a status of 4 (removed) is preferred for this situation.

The host that registered the address may withdraw the route but maintain the NCE, e.g., in the case where it is multihomed but does not want to use one interface for the traffic back as this time. This is done by registering with the R flag set to 0 as shown in Figure 24.
Figure 24: Route Type 2 Removal

When the lifetime elapses, the 6LBR requires the collocated EVPN router to withdraw the route.
8. Security Considerations

TBD

9. IANA Considerations

9.1. MAC Mobility Extended Community Flags

This document creates the "MAC Mobility Extended Community Flags" registry based on one flag initially defined in [RFC9047] and more flags defined in Section 6.2 as shown in Table 1:

```
+-----------------+------------------+--------+
| Flag Position   | Name              | Reference |
+-----------------+------------------+--------+
| 0..5            | Unassigned        | N/A    |
| 6               | Super-sticky (X)  | THIS RFC |
| 7               | Sticky (S)        | RFC 7432 |
+-----------------+------------------+--------+
```

Table 1: MAC Mobility Extended Community Flags
9.2. ARP/ND Extended Community Flags

This document modifies the "ARP/ND Extended Community Flags" registry initially created in Section 5 of [RFC9047]. Section 6.1 suggests to assign new entries in the Registry as indicated in Table 2:

<table>
<thead>
<tr>
<th>Flag Position</th>
<th>Name</th>
<th>Reference</th>
</tr>
</thead>
<tbody>
<tr>
<td>0 (Suggested)</td>
<td>Unreachable(U)</td>
<td>THIS RFC</td>
</tr>
<tr>
<td>1 (Suggested)</td>
<td>Multicast (M)</td>
<td>THIS RFC</td>
</tr>
<tr>
<td>2 (Suggested)</td>
<td>ROVR Validated (V)</td>
<td>THIS RFC</td>
</tr>
<tr>
<td>3 (Suggested)</td>
<td>ROVR Capable (H)</td>
<td>THIS RFC</td>
</tr>
</tbody>
</table>

Table 2: New ARP/ND Extended Community Flags

10. Acknowledgments

The authors wish to thank you for reading that far. We acknowledge and express gratitude to Wen Lin, Stephane Litkowski, Eric Levy-Abegnoli, Lukas Krattiger, Jerome Tollet, and Ali Sajassi, for their help and support.

11. Normative References


12. Informative References


IANA, "Address Registration Option Status Values", <https://www.iana.org/assignments/icmpv6-parameters/icmpv6-parameters.xhtml#address-registration>.

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Abstract

This document defines Multicast Virtual Private Network (VPN) extensions and procedures that allow fast failover for upstream failures by allowing upstream Provider Edges (PEs) to determine a single forwarder for a VPN multicast flow, without the downstream PEs' duplication prevention. The fast failover is accomplished by using Virtual Router Redundancy Protocol (VRRP) [RFC5798] or similar technologies for the upstream PEs to determine a single designated forwarder. Also, this document introduces a new BGP Extended Community called "Upstream Forwarder Selection", carried by BGP VPN route so that the upstream PEs can inform downstream PEs the election behavior. The downstream PEs, accordingly, send C-multicast routes to both the primary and standby upstream PEs and forward the multicast flow coming from both sides to the CEs.

Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

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Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."
1.  Introduction

MVPN [RFC6513] and [RFC6514] defines the MVPN architecture and MVPN
protocol specification which include the basic procedures for
selecting the Upstream Multicast Hop. Further [RFC9026] defines some
extension that allow fast failover for upstream failures by allowing
downstream PEs to consider the status of Provider-Tunnels (P-tunnels)
when selecting the Upstream PE for a VPN multicast flow. However,
there are some problems when deploying the "hot root standby"
mechanism described in [RFC9026].

First, all the ingress PEs, regardless of the primary or standby
role, forward (C-S,C-G) flow to other PEs though a P-tunnel, forcing
the egress PEs to discard all but one, which will cause the steady
traffic redundancy throughout the backbone network.
Second, an efficient and accurate method for the downstream PEs to determine the "status" of a P-tunnel is required, which is somewhat complicated in some cases, as mentioned in Section 3.1.8 of [RFC9026]. This document proposes a different "warm root standby" procedure mentioned in Section 4.2 of [RFC9026]. The procedures include a) an upstream designated forwarder election between multi homing ingress PEs, and b) the downstream PEs’ advertising Primary and Standby BGP C-multicast route and accepting traffic from any of both sides.

Section 3 describes procedures allowing multi homing ingress PEs to determine "locally" a single forwarder to avoid duplicate packets sending through the backbone, without the egress PEs’ primary or standby UMH selection.

Section 4 describes an optional BGP Extended Community called "Upstream Forwarder Selection", which is carried by BGP VPN routes (SAFI 128 or 129), to inform the downstream PEs the selection behavior describes in Section 3.

Section 5 describes the downstream PEs’ behavior in this case. The downstream PEs advertise C-multicast Source Tree Join route to both the primary and secondary Upstream PEs (carrying, as Route Target extended communities, the values of the VRF Route Import Extended Community of each VPN route from each Upstream PE). The Upstream Forwarder Selection Extended Community indicates that the packet duplication prevention will be accomplished by the upstream PEs and that any of the traffic from both the primary and secondary upstream PEs would be acceptable to be forwarded to the CEs.

2. Terminology

Readers of this document are assumed to be familiar with the terminology and concepts of the documents listed as Normative References.

3. Upstream Designated Forwarder Selection

Section 9.1.2 of [RFC6513] describes a "single forwarder selection" to ensure that duplicate packets not sending through the backbone. This document proposes a deployment of VRRP or some similar technology to enable dual or multi homing ingress PEs to determine a designated forwarder.
3.1. Upstream Designated Forwarder Selection by VRRP

VRRP specifies an election protocol that dynamically assigns responsibility for a virtual router to one of the VRRP routers on a LAN. The VRRP router controlling the IPv4 or IPv6 address(es) associated with a virtual router is called the Master, and it forwards packets sent to these IPv4 or IPv6 addresses. Similarly, the role of the VRRP routers associated with a virtual router can also be that of the upstream PEs in MVPN dual homing upstream PEs deployment.

Virtual Router -- pair of dual homing upstream PEs
Virtual Router Master -- the primary upstream PE
Virtual Router Backup -- the standby upstream PE

The method of mapping the role of a VRRP router to that of a MVPN upstream PE is more likely an administrative measure and could be implemented as configurable policies. Both the primary and standby PEs install VRF PIM state corresponding to BGP Source Tree Join route and send C-Join messages to the CE toward C-S. Whereas only the primary upstream PE (Virtual Router Master according to VRRP) forwards (C-S,C-G) flow to downstream PEs through a P-tunnel.

3.2. Other Feasible Selection Technologies

VRRP is just an example of the feasible choices for the dual homing upstream PEs’ single forwarder selection. Other private implementations or similar designated forwarder selection technologies could also be optional for further study. However, a feasible technology should have the ability of being deployed per VRF and being associated with one Multicast VPN instance.

4. Upstream Forwarder Selection Extended Community

This document defines a new BGP Extended Community called "Upstream Forwarder Selection".

The Upstream Forwarder Selection is an IP-address-specific Extended Community, of an extended type, and is transitive across AS boundaries [RFC4360].
An upstream PE constructs Upstream Forwarder Selection as follows, regardless of the role of the selection result:

The Global Administrator field of the Upstream Forwarder Selection SHOULD be set to a virtual IP address (or similar identity) of the upstream PEs (such as the VRRP Virtual IP address when using VRRP), which is identical between primary and standby PEs.

The Local Administrator field of the Upstream Forwarder Selection SHOULD be set to a master or backup status determined by the election which is different between primary and standby PEs.

Similar with the carrying of the VRF Route Import Extended Community imposed in Section 7 of [RFC6514], the multi homing PEs MUST also include in the BGP Updates message that carries the (unicast) VPN route the Upstream Forwarder Selection Extended Community that has the value of DF election result associated with this VRF.

5. Downstream PE Behavior

5.1. Standby C-multicast Route Advertisement

The Standby BGP C-multicast route advertisement described in Section 4 of [RFC9026] is still necessary. One downstream PE needs to determine a secondary UMH, originates and sends C-multicast routes with RTs that identify both the Primary and Standby Upstream PEs. However, because of the duplication prevention being accomplished by the upstream DF selection described above, carrying the new Standby PE BGP Communities with C-multicast routes is no longer a indispensable requirement.

5.2. Anycast Reverse Path Forwarding Checking

Multicast VPN specifications [RFC6513] impose that a downstream PE only forwards to CEs the packets coming from the expected Upstream PE (Section 9.1.1 of [RFC6513]).

When performing the UMH selection, if a route in the set of VPN-IP eligible UMH routes carries the Upstream Forwarder Selection Extended Community, the Upstream PE determined from the route should be considered a potentially valid Upstream PE. In most cases, there should be two of that routes for one (C-S,C-G) flow, indicating the primary and standby upstream PEs. As a result, the downstream PE
accepts the (C-S,C-G) flow from any of both sides and forward it to CEIs. It is a kind of "anycast" reverse path forwarding (RPF) checking. Eventually, it is the upstream single forwarder selection mechanism that ensures the duplicate packets not passing through the backbone network, as described in Section 3.

6. Security Considerations

This document introduces no new security considerations beyond those already specified in [RFC6513] and [RFC6514].

7. IANA Considerations

This document contains no actions for IANA.

8. Acknowledgements

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9. Normative References


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Advertising S-BFD Discriminators in BGP  
draft-wang-bess-sbfd-discriminator-02

Abstract

Seamless Bidirectional Forwarding Detection (S-BFD) is a simplified BFD mechanism. It eliminates most negotiation aspects and provides advantages such as fast configuration injection. S-BFD is especially useful in multi-homing PE scenarios and reduces resource overheads on the dual-homing PEs. Although S-BFD is simpler than BFD, a large number of manual configurations are required when there are a large number of PEs.

This document provides the mechanism of distributing S-BFD discriminators with VPN service routes, which simplifies S-BFD deployment for VPN services.

Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

Status of This Memo

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1. Introduction

[RFC7880] defines the Seamless Bidirectional Forwarding Detection (S-BFD) mechanism. S-BFD is a simplified mechanism for using BFD with a large proportion of negotiation aspects eliminated, thus providing benefits such as quick provisioning, as well as improved control and flexibility for network nodes initiating path monitoring. Currently, S-BFD can be used to simplify the service deployment.
During network construction, carriers usually deploy active and standby nodes to improve network reliability. In this way, when a single node is faulty, a protection switchover can be performed quickly. To accelerate fault detection, BFD is generally used. BFD sessions must be deployed on both ends of the BFD session, which occupies resources on both ends of the PE.

[RFC7880] defines Seamless Bidirectional Forwarding Detection (S-BFD), a simplified mechanism for using BFD with a large proportion of negotiation aspects eliminated, thus providing benefits such as quick provisioning, as well as improved control and flexibility for network nodes initiating path monitoring. This mechanism is useful for asymmetric scenarios, such as 3PE scenarios. In dual-homing scenarios, BFD does not need to be deployed to detect single-homing nodes. In this scenario, S-BFD greatly saves resources on the dual-homing side.

To deploy S-BFD, the initiator needs to know the reflector’s endpoint and identifier. When a large number of PEs need to be deployed, the deployment is complicated. [RFC7883] and [RFC7884] introduced an IGP-based S-BFD discriminator advertisement mechanism to simplify S-BFD deployment. VPN service may be deployed across inter-area or inter-AS. In this case, the IGP flooding mechanism does not work.

It is recommended to use BGP to distribute BFD discriminator information. BGP can transmit routes across domains, and service routes can drive to generate the end-to-end S-BFD sessions on demand.

2. Terminology

BFD : Bidirectional Forwarding Detection
S-BFD : Seamless Bidirectional Forwarding Detection
APE : Access PE, used to access users
SPE: Service PE, used to support service for users
UCE: User CE
SCE: Service CE
3. Scenarios

In some EVPN deployments, for example, when it spans over multiple domains, only one of a pair of interconnected PEs benefits from monitoring the status of the connection. In such a case, using S-BFD [RFC7880] is advantageous as it reduces the load on one of the PEs while providing the benefit of faster convergence. The following sections provide examples of EVPNs that would benefit from using S-BFD.

For SRv6 services, there are two different service types. One is service over SRv6 BE, the other is service over SRv6 Policy. For the service over SRv6 BE, it will use the VPNSID to resolve the forwarding information. Thus we must generate an S-BFD session to detect the VPNSID’s reachability. This is an IP-routed S-BFD. We may use the remote VPNSID’s locator as the destination of the S-BFD session. For the service over SRv6 Policy, it will use <nexthop, color> of the service route to resolve an SRv6 Policy. Then we must generate an S-BFD session to detect the reachability of the SR Policy.

3.1. EVPN Layer 3 Service Over SRv6 BE Use Case

Figure 1 shows a SRv6 BE based seamless scenario. UCE is single-homed to APE, and SCE is dual-homed to SPE1 and SPE2. The service is across multiple ASes.

SCE1 accesses SPE1 and SPE2 through Layer 3 and advertises its private network routes to them. SPE1 and SPE2 encapsulate the routes into Type 5 routes in the EVPN format and sends them to APE1. After receiving Type 5 routes advertised by SPE1 and SPE2, APE1 generates
primary and backup entries for the routes to speed up service switchover. In this scenario, the SRv6 BE service mode is used. APE1 will resolve SPE1’s VPN routes reachability through the VPNSID. To ensure that APE1 can properly route to PE1, PE1 needs to advertise its own locator route. The advertisement of the locator route is not in the scope of this document.

To speed up fault detection, we may configure an S-BFD session on APE1 to detect SPE1 or SPE2’s reachability. In traditional mode, a discriminator needs to be assigned by SPE1 and SPE2, and two S-BFD sessions need to be configured on APE1 to detect the VPN SID’s reachability of SPE1 and SPE2. It needs to generate an S-BFD session with the destination set to the VPN SID. To reduce the number of S-BFD sessions, locator-based S-BFD sessions can be used instead of S-BFD sessions for VPNSIDs.

There are a large number of such APEs that exist on the network. Each APE is configured with several S-BFD sessions to detect PE1 and PE2, which increases the deployment complexity.

3.2. EVPN Layer 3 Service Over SRv6 Policy Use Case

Figure 2: EVPN Layer 3 Service Over SRv6 Policy

Figure 2 shows a SRv6 Policy scenario. SCE1 is dual-homed to SPE1 and SPE2, and UCE1 is accessed to APE1. SPE1, SPE2, and APE1 are cross BGP ASes.

SCE1 accesses SPE1 and SPE2 through Layer 3 and advertises its private network routes to APE1. SPE1 and SPE2 encapsulate the routes into Type 5 routes in the EVPN format and sends them to APE1.
After receiving Type 5 routes advertised by SPE1 and SPE2, APE1 generates primary and backup entries for the routes, speeding up service switchover. APE1 parses the tunnel based on the <nexthop, color> of the service routes advertised by SPE1 and SPE2, and matches an SRv6 Policy. After receiving the traffic from UCE1 to SCE1, APE1 encapsulates and forwards the traffic based on the SRv6 Policy.

An S-BFD session needs to be established for these SRv6 Policy-based forwarding paths to swiftly detect the availability of the paths. When detecting a fault on the SRv6 Policy path of the primary service route, services can be swiftly switched to the backup path, providing more reliable protection for services.

There are a large number of such PEs that exist on the network. Each PE is configured with several S-BFD sessions to detect PE1 and PE2, which increases the deployment complexity.

Certainly, this scenario may also be implemented in other methods. For example, when delivering an SRv6 policy, specify a tunnel to generate an S-BFD session.

4. Procedure

4.1. BGP Encoding

[RFC9026] specifies the "BFD Discriminators" (38) attribute, which is an optional transitive BGP attribute that conveys the Discriminators and other optional attributes used to establish BFD sessions.

The attribute defined in [RFC9026] is used to transmit P2MP BFD session creation information through the BFD Discriminator attribute in MVPN scenarios. For non-multicast services, such as L3VPN services, L2VPN services, and native IP services, BFD discriminators are also required to create an S-BFD session.

The S-BFD Discriminator attribute introduced in this document is defined as follows:
The BFD Mode field is 1 octet. [RFC9026] defines only the P2MP BFD session for MVPN. This document defines two new types of S-BFD session types based on the preceding scenarios.

As described in the preceding scenario. There are two types of S-BFD sessions for SRv6 services. For service over SRv6 BE, an IP-routed S-BFD session needs to be created to detect the locator route. For service over SRv6 Policy, an S-BFD session for SRv6 Policy path needs to be created to detect the SRv6 Policy path. So two new BFD modes should be introduced here.

S-BFD for SRv6 Locator Session Mode, which is dedicated to detecting the locator. The temporary type is 176, and is to be allocated by IANA.

S-BFD for Common Session Mode, which is for general S-BFD session. The temporary type is 177, and is to be allocated by IANA. This mode is not only for SRv6, but also can be used for other scenarios.

BFD Discriminators:

The field length is 4 octets. Used to specify the discriminator for S-BFD session.

Optional TLVs:

Variable-length fields are optional. Indicates the additional information required for creating a S-BFD session. The format is as follows:
If a transit node changes the next hop or reassigns a VPN SID when forwarding a route, the transit node needs to use the locally allocated S-BFD discriminator to advertise the S-BFD discriminator attribute. If the transit node does not recognize the S-BFD Discriminator attribute in the learned route and continues to advertise the route to the remote PE, the receiver may use incorrect information when creating an S-BFD session. Therefore, the advertised S-BFD Discriminator attribute needs to carry the IP address for receiver verification.

In this document, S-BFD for SRv6 Locator Session and S-BFD for Common Session must carry IP addresses except discriminators, which reuse the Source IP Address TLV defined in [RFC9026].

If the mode is set to S-BFD for SRv6 Locator Session, the SRv6 Locator address used for the service is carried.

If the mode is set to S-BFD for Common Session, the next-hop address used for the service is carried.

For details about the error handling, see section "Error Handling".

4.2. Router Procedure

In BGP address families, such as L3VPN or EVPN, routes can carry the S-BFD Discriminator attribute as required so that S-BFD sessions can be established based on the attribute. The following uses S-BFD for SRv6 Locator as an example. If mode is set to S-BFD for Common Session, the processing method is similar.

4.2.1. Egress Node Process

As shown in figure 1, the S-BFD discriminator is configured on PE1. After obtaining the information, BGP encapsulates the attribute into the EVPN route and sets the BFD Mode to S-BFD for Locator Session, when advertising the EVPN route. The Discriminator value is local discriminator value. The optional TLV carries the local PE’s locator address used by the VPN.
4.2.2. Transit Node Process

Here is the end-to-end SRv6 BE scenario. The ASBR does not re-allocate the VPN SID. Thus, the ASBR does not require to modify the VPN SID, and not to alter the BFD discriminator attribute.

4.2.3. Ingress Node Process

After receiving the EVPN Type 5 routes from PE1 and PE2, PE3 imports the routes to the VRF of PE3 based on the route targets. Routes triggers establish the S-BFD sessions based on <S-BFD discriminator, locator ip>.

Then, routes with the same prefix from PE1 and PE2 form primary and backup paths. When the primary path or the egress node is in fault, S-BFD detects that fault and forms switch to backup path quickly.

To avoid the waste of redundant resources, assume that the ASBR re-assigns the SID in Option B and the ASBR does not recognize the attribute. In this case, the SID and locator carried in the route received by PE3 do not match the Source IP carried in the Optional TLV in the BFD attribute. Therefore, PE3 does not need to establish an S-BFD session to remote PE, which can avoid resource waste.

5. Error handling

Error handling complies with [RFC7606]. In this document, the BFD discriminator information is used only to establish an S-BFD session. Therefore, if the BFD discriminator information is invalid, the BFD attribute will be discard and not transmit to other devices.

For BFD discriminator attribute, the following case will be processed:

- The BFD Discriminator value in receiving BFD Discriminator attribute is 0, the attribute is invalid.

For BFD mode type is S-BFD for SRv6 Locator Session, the following case will be processed:

- The BFD discriminator attribute doesn’t contain optional TLV with type set to 1, the attribute is invalid.

- The optional TLV type is 1 but the length is not 16, the attribute is invalid.

- The optional TLV type is 1 but the value is all 0, the attribute is invalid.
o If multiple Source IP Optional TLVs are carried, the first source IP address should be used as the destination to establish an S-BFD session. For EVPN type 2 MAC-IP routes may use the first and the second IP address because it may carry two SRv6 SIDs with different locators. Other source IP addresses should be ignored.

o If a non-Source IP Optional TLV is carried, the Optional TLV will be ignored.

For BFD mode type is S-BFD for Common Session, the following case will be processed:

- The BFD discriminator attribute doesn’t contain optional TLV with type set to 1, the attribute is invalid.
- The optional TLV type is 1 but the length is not 4 or 16, the attribute is invalid.
- The optional TLV type is 1 but the value is all 0, the attribute is invalid.
- If multiple Source IP Optional TLVs are carried, only the first source IP address should be used as the destination to establish an S-BFD session. Other source IP addresses should be ignored.
- If a non-Source IP Optional TLV is carried, the Optional TLV will be ignored.

6. IANA Considerations

This document defines two new BFD modes in the BFD Discriminator attribute. The following values are recommended to be assigned by IANA:

<table>
<thead>
<tr>
<th>Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>176</td>
<td>S-BFD for SRv6 Locator Session</td>
</tr>
<tr>
<td>177</td>
<td>S-BFD for Common Session</td>
</tr>
</tbody>
</table>

7. Security Considerations

The new S-BFD Discriminators sub-TLV does not introduce any new security risks for BGP.

When creating an S-BFD session, the initiator verifies the S-BFD session based on routing information. This reduces the number of invalid S-BFD sessions and avoid attribute attack.
8. Acknowledgements

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9. References

9.1. Normative References


9.2. References


Authors’ Addresses

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