Considerations for Benchmarking Network Performance in Containerized Infrastructures
draft-dcn-bmwg-containerizedinfra-08

Abstract

This draft describes considerations for benchmarking network performance in containerized infrastructures. In the containerized infrastructure, Virtualized Network Functions (VNFs) are deployed on an operating-system-level virtualization platform by abstracting the user namespace as opposed to virtualization using a hypervisor. Hence, the system configurations and networking scenarios for benchmarking will be partially changed by how the resource allocation and network technologies are specified for containerized VNFs. This draft compares the state of the art in the container networking architecture with VM-based virtualized systems networking architecture and provides several test scenarios for benchmarking network performance in containerized infrastructures.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at https://datatracker.ietf.org/drafts/current/.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on 2 September 2022.
Copyright Notice

Copyright (c) 2022 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust's Legal Provisions Relating to IETF Documents (https://trustee.ietf.org/license-info) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Revised BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Revised BSD License.

Table of Contents

1. Introduction ................................................. 3
2. Terminology .................................................. 4
3. Containerized Infrastructure Overview ....................... 4
4. Networking Models in Containerized Infrastructure .......... 8
   4.1. Kernel-space vSwitch Model .............................. 9
   4.2. User-space vSwitch Model ............................... 10
   4.3. eBPF Acceleration Model ............................... 10
   4.4. Smart-NIC Acceleration Model ......................... 12
   4.5. Model Combination ..................................... 13
5. Performance Impacts .......................................... 14
   5.1. CPU Isolation / NUMA Affinity ......................... 14
   5.2. Hugepages ............................................. 15
   5.3. Service Function Chaining ............................. 15
   5.4. Additional Considerations ............................. 16
6. Security Considerations ...................................... 16
7. References .................................................... 16
   7.1. Informative References .................................. 16
Appendix A. Benchmarking Experience (Contiv-VPP) ............. 18
   A.1. Benchmarking Environment ............................... 18
   A.2. Trouble shooting and Result ........................... 22
Appendix B. Benchmarking Experience (SR-IOV with DPDK) ....... 23
   B.1. Benchmarking Environment ............................... 24
   B.2. Trouble shooting and Results ........................... 27
Appendix C. Benchmarking Experience (Multi-pod Test) ........... 27
   C.1. Benchmarking Overview ................................. 27
   C.2. Hardware Configurations ............................... 28
   C.3. NUMA Allocation Scenario ............................. 30
   C.4. Traffic Generator Configurations ..................... 30
   C.5. Benchmark Results and Trouble-shootings ............... 30
Authors’ Addresses .............................................. 31
1. Introduction

The Benchmarking Methodology Working Group (BMWG) has recently expanded its benchmarking scope from Physical Network Function (PNF) running on a dedicated hardware system to Network Function Virtualization (NFV) infrastructure and Virtualized Network Function (VNF). [RFC8172] described considerations for configuring NFV infrastructure and benchmarking metrics, and [RFC8204] gives guidelines for benchmarking virtual switch which connects VNFs in Open Platform for NFV (OPNFV).

Recently NFV infrastructure has evolved to include a lightweight virtualized platform called the containerized infrastructure, where VNFs share the same host Operating System (OS) and are logically isolated by using a different namespace. While previous NFV infrastructure uses a hypervisor to allocate resources for Virtual Machine (VMs) and instantiate VNFs, the containerized infrastructure virtualizes resources without a hypervisor, making containers very lightweight and more efficient in infrastructure resource utilization compared to the VM-based NFV infrastructure. When we consider benchmarking for VNFs in the containerized infrastructure, it may have a different System Under Test (SUT) and Device Under Test (DUT) configuration compared with both black-box benchmarking and VM-based NFV infrastructure as described in [RFC8172]. Accordingly, additional configuration parameters and testing strategies may be required.

In the containerized infrastructure, a VNF network is implemented by running both switch and router functions in the host system. For example, the internal communication between VNFs in the same host uses the L2 bridge function, while communication with external node(s) uses the L3 router function. For container networking, the host system may use a virtual switch (vSwitch), but other options exist. In the [ETSI-TST-009], they describe differences in networking structure between the VM-based and the containerized infrastructure. Occasioned by these differences, deployment scenarios for testing network performance described in [RFC8204] may be partially applied to the containerized infrastructure, but other scenarios may be required.

This draft aims to distinguish benchmarking of containerized infrastructure from the previous benchmarking methodology of common NFV infrastructure. Considering the point in [RFC8204] that virtual switch (vSwitch) is the networking principle of containerized infrastructure, this draft investigates different network models based on vSwitch location and acceleration technologies. At the same time, it is essential to uncover the impact of different deployment configurations on containerized infrastructure, such as resource
isolation, hugepages, service function chaining. The benchmark experiences of various combinations of these mentioned configurations and networking models are also presented in this draft as the references to set up and benchmark containerized infrastructure. Note that, although the detailed configurations of both infrastructures differ, the new benchmarks and metrics defined in [RFC8172] can be equally applied in containerized infrastructure from a generic-NFV point of view, and therefore defining additional metrics or methodologies are out of scope.

2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document is to be interpreted as described in [RFC2119]. This document uses the terminology described in [RFC8172], [RFC8204], [ETSI-TST-009].

3. Containerized Infrastructure Overview

For benchmarking of the containerized infrastructure, as mentioned in [RFC8172], the basic approach is to reuse existing benchmarking methods developed within the BMWG. Various network function specifications defined in BMWG should still be applied to containerized VNF(C-VNF)s for the performance comparison with physical network functions and VM-based VNFs. A major distinction of the containerized infrastructure from the VM-based infrastructure is the absence of a hypervisor. Without hypervisor, all C-VNFs share the same host resources, including but not limited to computing, storage, and networking resources, as well as the host Operating System(OS), kernel, and libraries. These architectural differences bring additional considerations of resource management impacts for benchmarking.

In a common containerized infrastructure, thanks to the proliferation of Kubernetes, the pod is defined as a basic unit for orchestration and management that can host multiple containers. Based on that, [ETSI-TST-009] defined two test scenario for container infrastructure as follows.

- Container2Container: Communication between containers running in the same pod. It can be done by shared volumes or Inter-process communication (IPC).
- Pod2Pod: Communication between containers running in the different pods.
As mentioned in [RFC8204], vSwitch is also an important aspect of the containerized infrastructure. For Pod2Pod communication, every pod has only one virtual Ethernet (vETH) interface. This interface is connected to the vSwitch via vETH pair for each container. Not only Pod2Pod but also Pod2External scenario that communicates with an external node is also required. In this case, vSwitch SHOULD support gateway and Network Address Translation (NAT) functionalities.

Figure 1 shows briefly differences of network architectures based on container deployment models. Basically, on bare metal, C-VNFs can be deployed as a cluster called POD by Kubernetes. Otherwise, each C-VNF can be deployed separately using Docker. In the former case, there is only one external network interface, even a POD containing more than one C-VNF. An additional deployment model considers a scenario where C-VNFs or PODs are running on VM. In our draft, we define new terminologies; BMP, which is Pod on bare metal, and VMP, which is Pod on VM.
In [ETSI-TST-009], they described data plane test scenarios in a single host. In that document, there are two scenarios for containerized infrastructure; Container2Container, which is internal communication between two containers in the same Pod, and the Pod2Pod model, which is communication between two containers running in different Pods. According to our new terminologies, we can call the Pod2Pod model the BMP2BMP scenario. When we consider container running on VM as an additional deployment option, there can be more single host test scenarios as follows:

- BMP2VMP scenario
Figure 2: Single Host Test Scenario - BMP2VMP

- VMP2VMP scenario
4. Networking Models in Containerized Infrastructure

Container networking services are provided as network plugins. Basically, by using them, network services are deployed as an isolation environment from container runtime through the host namespace, creating a virtual interface, allocating interface and IP address to C-VNF. Since the containerized infrastructure has different network architecture depending on its using plugins, it is necessary to specify the plugin used in the infrastructure. Especially for Kubernetes infrastructure, several Container Networking Interface (CNI) plugins are developed, which describes network configuration files in JSON format, and plugins are instantiated as new namespaces. When the CNI plugin is initiated, it pushes forwarding rules and networking policies to the existing vSwitch (i.e., Linux bridge, Open vSwitch) or creates its own switch functions to provide networking service.
The container network model can be classified according to the location of the vSwitch component. There are some CNI plugins that provide networking without the vSwitch components; however, this draft focuses on plugins using vSwitch components.

4.1. Kernel-space vSwitch Model

![Kernel-space vSwitch Model Diagram]

Figure 4: Examples of Kernel-Space vSwitch Model

Figure 4 shows kernel-space vSwitch model. In this model, because the vSwitch component is running on kernel space, data packets should be processed in-network stack of host kernel before transferring packets to the C-VNF running in user-space. Not only pod2External but also pod2pod traffic should be processed in the kernel space. For dynamic networking configuration, the Forwarding policy can be pushed by the controller/agent located in the user-space. In the case of Open vSwitch (OVS) [OVS], the first packet of flow can be sent to the user space agent (ovs-switchd) for forwarding decision. Kernel-space vSwitch models are listed below;

- Docker Network [Docker-network], Flannel Network [Flannel], OVS (OpenvSwitch) [OVS], OVN (Open Virtual Network) [OVN]
4.2. User-space vSwitch Model

Figure 5: Examples of User-Space vSwitch Model

Figure 5 shows user-space vSwitch model, in which data packets from physical network port are bypassed kernel processing and delivered directly to the vSwitch running on user-space. This model is commonly considered as Data Plane Acceleration (DPA) technology since it can achieve high-rate packet processing than a kernel-space network with limited packet throughput. For bypassing kernel and directly transferring the packet to vSwitch, Data Plane Development Kit (DPDK) is essentially required. With DPDK, an additional driver called Pull-Mode Driver (PMD) is created on vSwitch. PMD driver must be created for each NIC separately. User-space vSwitch models are listed below;

- ovs-dpdk[ovs-dpdk], vpp[vpp]

4.3. eBPF Acceleration Model
Figure 6: Examples of eBPF Acceleration Model

Figure 6 shows eBPF Acceleration model, which leverages extended Berkeley Packet Filter (eBPF) technology to achieve high-performance packet processing. It enables execution of sandboxed programs inside abstract virtual machines within the Linux kernel without changing the kernel source code or loading the kernel module. To accelerate data plane performance, eBPF programs are attached to different BPF hooks inside the Linux kernel stack.

One type of BPF hook is the eXpress Data Path (XDP) at the networking driver. It is the first hook that triggers eBPF program upon packet reception from external network. The other type of BPF hook is Traffic Control Ingress/Egress eBPF hook (tc eBPF). These hooks are attached to the vETH pair of the pod and the XDP hook. The tc Egress
eBPF hooks at the vETH pair enforce policy on all traffic exit the pod, while the tc Ingress eBPF hook at the end of the kernel networking runs after initial packet processing from XDP hook.

On the egress datapath side, whenever a packet exits the pod, it goes through vETH pair then is picked up by the tc egress eBPF hook. These hooks trigger eBPF programs to forward the packet directly to the external facing network interface, bypassing all of the kernel network layer processing such as iptables. On the ingress datapath side, eBPF programs at the XDP and tc ingress eBPF hook pick up packets from the network device and directly deliver it to the vETH interface pair, or bypassing context-switching process to the pod network namespace in the case of Cilium project [Cilium].

Notable eBPF Acceleration models are 2 CNI plugin projects: Calico[Calico], Cilium[Cilium]. In the case of Cilium, eBPF/XDP program can be offloaded directly on the smart NIC card, which allows data plane acceleration without using the CPU. Container network performance of these eBPF-based project is reported in [cilium-benchmark].

4.4. Smart-NIC Acceleration Model

![Smart-NIC Acceleration Model Diagram](attachment:fig7.png)

**Figure 7: Examples of Smart-NIC Acceleration Model**
Figure 7 shows Smart-NIC acceleration model, which does not use vSwitch component. This model can be separated into two technologies.

One is Single-Root I/O Virtualization (SR-IOV) [SR-IOV], which is an extension of PCIe specifications to enable multiple partitions running simultaneously within a system to share PCIe devices. In the NIC, there are virtual replicas of PCI functions known as virtual functions (VF), and each of them is directly connected to each container’s network interfaces. Using SR-IOV, data packets from external bypass both kernel and user space and are directly forwarded to container’s virtual network interface.

The other technology is eBPF/XDP programs offloading to Smart-NIC card as mentioned in the previous section. It enables general acceleration of eBPF. eBPF programs are attached to XDP and run at the Smart-NIC card, which allows server CPUs to perform more application-level work. However, not all Smart-NIC cards provide eBPF/XDP offloading support.

4.5. Model Combination
Figure 8 shows the networking model when combining user-space vSwitch model and Smart-NIC acceleration model. This model is frequently considered in service function chain scenarios when two different types of traffic flows are present. These two types are North/South traffic and East/West traffic.

North/South traffic is the type that packets are received from other servers and routed through VNF. For this traffic type, Smart-NIC model such as SR-IOV is preferred because packets always have to pass the NIC. User-space vSwitch involvement in north-south traffic will create more bottlenecks. On the other hand, East/West traffic is a form of sending and receiving data between containers deployed in the same server and can pass through multiple containers. For this type, user-space vSwitch models such as OVS-DPDK and VPP are preferred because packets are routed within the user space only and not through the NIC.

The throughput advantages of these different networking models with different traffic direction cases are reported in [Intel-SRIOV-NFV].

5. Performance Impacts

5.1. CPU Isolation / NUMA Affinity

CPU pinning enables benefits such as maximizing cache utilization, eliminating operating system thread scheduling overhead as well as coordinating network I/O by guaranteeing resources. This technology is very effective in avoiding the "noisy neighbor" problem, and it is already proved in existing experience [Intel-EPA].

Using NUMA, performance will be increasing not CPU and memory but also network since that network interface connected PCIe slot of specific NUMA node have locality. Using NUMA requires a strong understanding of VNF’s memory requirements. If VNF uses more memory than a single NUMA node contains, the overhead will occur due to being spilled to another NUMA node. Network performance can be changed depending on the location of the NUMA node whether it is the same NUMA node where the physical network interface and CNF are attached to. There is benchmarking experience for cross-NUMA performance impacts [ViNePERF]. In that tests, they consist of cross-NUMA performance with 3 scenarios depending on the location of the traffic generator and traffic endpoint. As the results, it was verified as below:

- A single NUMA Node serving multiple interfaces is worse than Cross-NUMA Node performance degradation
5.2. Hugepages

Hugepages configure a large page size of memory to reduce Translation Lookaside Buffer (TLB) miss rate and increase the application performance. This increases the performance of logical/virtual to physical address lookups performed by a CPU’s memory management unit, and overall system performance. In the containerized infrastructure, the container is isolated at the application level, and administrators can set huge pages more granular level (e.g., Kubernetes allows to use of 512M bytes huge pages for the container as default values). Moreover, this page is dedicated to the application but another process, so the application uses the page more efficiently way. From a network benchmark point of view, however, the impact on general packet processing can be relatively negligible, and it may be necessary to consider the application level to measure the impact together. In the case of using the DPDK application, as reported in [Intel-EPA], it was verified to improve network performance because packet handling processes are running in the application together.

5.3. Service Function Chaining

When we consider benchmarking for containerized and VM-based infrastructure and network functions, benchmarking scenarios may contain various operational use cases. Traditional black-box benchmarking focuses on measuring the in-out performance of packets from physical network ports since the hardware is tightly coupled with its function and only a single function is running on its dedicated hardware. However, in the NFV environment, the physical network port commonly will be connected to multiple VNFs (i.e., Multiple PVP test setup architectures were described in [ETSI-TST-009]) rather than dedicated to a single VNF. This scenario is called Service Function Chaining. Therefore, benchmarking scenarios should reflect operational considerations such as the number of VNFs or network services defined by a set of VNFs in a single host. [service-density] proposed a way for measuring the performance of multiple NFV service instances at a varied service density on a single host, which is one example of these operational benchmarking aspects. Another aspect in benchmarking service function chaining scenario should be considered is different network acceleration technologies. Network performance differences may occur because of different traffic patterns based on the provided acceleration method.
5.4. Additional Considerations

Apart from the single-host test scenario, the multi-hosts scenario should also be considered in container network benchmarking, where container services are deployed across different servers. To provide network connectivity for container-based VNFs between different server nodes, inter-node networking is required. According to [ETSI-NFV-IFA-038], there are several technologies to enable inter-node network: overlay technologies using a tunnel endpoint (e.g. VXLAN, IP in IP), routing using Border Gateway Protocol (BGP), layer 2 underlay, direct network using dedicated NIC for each pod, or load balancer using LoadBalancer service type in Kubernetes. Different protocols from these technologies may cause performance differences in container networking.

6. Security Considerations

TBD

7. References

7.1. Informative References


[ETSI-NFV-IFA-038]

[ETSI-TST-009]


[Intel-EPA]

[Intel-SRIOV-NFV]

[OVN]

[OVS]


Appendix A. Benchmarking Experience (Contiv-VPP)

A.1. Benchmarking Environment

In this test, our purpose is to test the performance of user-space based model for container infrastructure and figure out the relationship between resource allocation and network performance. With respect to this, we set up Contiv-VPP, one of the user-space based network solutions in container infrastructure and tested like below.

- Three physical server for benchmarking
<table>
<thead>
<tr>
<th>Node Name</th>
<th>Specification</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Container Control for Master</td>
<td>- Intel(R) Xeon(R)</td>
<td>Container Deployment and Network Allocation</td>
</tr>
<tr>
<td></td>
<td>CPU E5-2690</td>
<td>- ubuntu 18.04</td>
</tr>
<tr>
<td></td>
<td>(2Socket X 12Core)</td>
<td>- Kubernetes Master</td>
</tr>
<tr>
<td></td>
<td>- MEM 128G</td>
<td>- CNI Controller</td>
</tr>
<tr>
<td></td>
<td>- DISK 2T</td>
<td>.. Contive VPP Controller</td>
</tr>
<tr>
<td></td>
<td>- Control plane : 1G</td>
<td>.. Contive VPP Agent</td>
</tr>
<tr>
<td>Container Service for Worker</td>
<td>- Intel(R) Xeon(R)</td>
<td>Container Service</td>
</tr>
<tr>
<td></td>
<td>Gold 6148</td>
<td>- ubuntu 18.04</td>
</tr>
<tr>
<td></td>
<td>(2socket X 20Core)</td>
<td>- Kubernetes Worker</td>
</tr>
<tr>
<td></td>
<td>- MEM 128G</td>
<td>- CNI Agent</td>
</tr>
<tr>
<td></td>
<td>- DISK 2T</td>
<td>.. Contive VPP Agent</td>
</tr>
<tr>
<td></td>
<td>- Control plane : 1G</td>
<td></td>
</tr>
<tr>
<td></td>
<td>- Data plane : MLX 10G (1NIC 2PORT)</td>
<td></td>
</tr>
<tr>
<td>Packet Generator</td>
<td>- Intel(R) Xeon(R)</td>
<td>Packet Generator</td>
</tr>
<tr>
<td></td>
<td>CPU E5-2690</td>
<td>- CentOS 7</td>
</tr>
<tr>
<td></td>
<td>(2Socket X 12Core)</td>
<td>- installed Trex 2.4</td>
</tr>
<tr>
<td></td>
<td>- MEM 128G</td>
<td></td>
</tr>
<tr>
<td></td>
<td>- DISK 2T</td>
<td></td>
</tr>
<tr>
<td></td>
<td>- Control plane : 1G</td>
<td></td>
</tr>
<tr>
<td></td>
<td>- Data plane : MLX 10G (1NIC 2PORT)</td>
<td></td>
</tr>
</tbody>
</table>

Figure 9: Test Environment-Server Specification

- The architecture of benchmarking
Figure 10: Test Environment-Architecture

- Network model of Containerized Infrastructure (User space Model)
We set up a Contive-VPP network to benchmark the user space container network model in the containerized infrastructure worker node. We set up network interface at NUMA0, and we created different network subnets VRF1, VRF2 to classify input and output data traffic,

Figure 11: Test Environment-Network Architecture
respectively. And then, we assigned two interfaces which connected
to VRF1, VRF2 and, we setup routing table to route Trex packet from
eth1 interface to eth2 interface in POD.

A.2. Trouble shooting and Result

In this environment, we confirmed that the routing table doesn’t work
when we send packets using Trex packet generator. The reason is that
when kernel space based network configured, ip forwarding rule is
processed to kernel stack level while ‘ip packet forwarding rule’ is
processed only in vrf0, which is the default virtual routing and
forwarding (VRF0) in VPP. The above testing architecture makes
problem since vrf1 and vrf2 interface couldn’t route packet.
According to above result, we assigned vrf0 and vrf1 to POD and, data
flow is like below.

```
+---------------------------------------------+---------------------+
|                   NUMA 0                    |        NUMA 0       |
+---------------------------------------------|---------------------+
|  Containerized Infrastructure Worker Node   |                     |
|        +---------------------------+        |  +----------------+ |
|        |      POD1                 |        |  |     POD2       | |
|        |      +-------------+      |        |  |   +-------+    | |
|        |   +--v----+    +---v--+   |        |  | +-v--+  +-v--+ | |
|        |   | eth1 |     | eth2 |   |        |  | |eth1|  |eth2| | |
|        |   +--^---+     +---^--+   |        |  | +-^--+  +-^--+ | |
|        +------|-------------|------+        |  +---|-------|----+ |
|       +-------+             |               |      |       |      |
|       |       +-------------|---------------|------+       |      |
|       |       |             |        +------|--------------+      |
| +-----|-------|-------------|--------|----+ |                     |
| |     |       |             v        v    | |                     |
| |     |       |          +-tap10--tap11-+ | |                     |
| |     |       |          |  ^        ^  | | |                     |
| |     |       |          |  |  VRF1  |  | | |                     |
| |     |       |          +--|--------|--+ | |                     |
| |     |       |             |    +---+    | |                     |
| | +-*tap00--*tap01----------|----|---+    | | User Space          |
| |   +-----^-----+      +------^----+        | User Space          |
| |   +-----v-----+      +------v----+      | |                     |
| +---|*DPDK PMD0 |------|*DPDK PMD1 |------+ |                     |
+-----------|-------------------|-------------|---------------------+
|-----------|-------------------|-------------|---------------------+

*- CPU pinning interface
```
We conducted benchmarking with three conditions. The test environments are as follows. - Basic VPP switch - General kubernetes (No CPU Pining) - Shared Mode / Exclusive mode. In the basic Kubernetes environment, all PODs share a host’s CPU. Shared mode is that some POD share a pool of CPU assigned to specific PODs. Exclusive mode is that a specific POD dedicates a specific CPU to use. In shared mode, we assigned two CPUs for several PODs, in exclusive mode, we dedicated one CPU for one POD, independently. The result is like Figure 13. First, the test was conducted to figure out the line rate of the VPP switch, and the basic Kubernetes performance. After that, we applied NUMA to the network interface using Shared Mode and Exclusive Mode in the same node and different node. In Exclusive and Shared mode tests, we confirmed that Exclusive mode showed better performance than Shared mode when same NUMA CPU was assigned, respectively. However, we confirmed that performance is reduced at the section between the vpp switch and the POD, affecting the total result.

<table>
<thead>
<tr>
<th>Model</th>
<th>NUMA Mode (pinning)</th>
<th>Result (Gbps)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Maximum Line Rate</td>
<td>N/A</td>
<td>3.1</td>
</tr>
<tr>
<td></td>
<td>same NUMA</td>
<td>9.8</td>
</tr>
<tr>
<td>Without CMK</td>
<td>N/A</td>
<td>1.5</td>
</tr>
<tr>
<td>CMK-Exclusive Mode</td>
<td>same NUMA</td>
<td>4.7</td>
</tr>
<tr>
<td></td>
<td>Different NUMA</td>
<td>3.1</td>
</tr>
<tr>
<td>CMK-shared Mode</td>
<td>same NUMA</td>
<td>3.5</td>
</tr>
<tr>
<td></td>
<td>Different NUMA</td>
<td>2.3</td>
</tr>
</tbody>
</table>

Figure 13: Test Results
B.1. Benchmarking Environment

In this test, our purpose is to test the performance of Smart-NIC acceleration model for container infrastructure and figure out relationship between resource allocation and network performance. With respect to this, we setup SRIOV combining with DPDK to bypass the Kernel space in container infrastructure and tested based on that.

- Three physical server for benchmarking

<table>
<thead>
<tr>
<th>Node Name</th>
<th>Specification</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Conatiner Control for Master</td>
<td>- Intel(R) Core(TM) i5-6200U CPU (1socket x 4Core) - MEM 8G - DISK 500GB - Control plane : 1G</td>
<td>Container Deployment and Network Allocation - ubuntu 18.04 - Kubernetes Master - CNI Controller MULTUS CNI SRIOV plugin with DPDK</td>
</tr>
<tr>
<td>Conatiner Service for Worker</td>
<td>- Intel(R) Xeon(R) E5-2620 v3 @ 2.4Ghz (1socket X 6Core) - MEM 128G - DISK 2T - Control plane : 1G - Data plane : XL710-qda2 (1NIC 2PORT- 40Gb)</td>
<td>Container Service - Centos 7.7 - Kubernetes Worker - CNI Agent MULTUS CNI SRIOV plugin with DPDK</td>
</tr>
<tr>
<td>Packet Generator</td>
<td>- Intel(R) Xeon(R) Gold 6148 @ 2.4Ghz (2Socket X 20Core) - MEM 128G - DISK 2T - Control plane : 1G - Data plane : XL710-qda2 (1NIC 2PORT- 40Gb)</td>
<td>Packet Generator - CentOS 7.7 - installed Trex 2.4</td>
</tr>
</tbody>
</table>

Figure 14: Test Environment-Server Specification

- The architecture of benchmarking
Figure 15: Test Environment-Architecture

- Network model of Containerized Infrastructure (User space Model)
We set up a Multus CNI, SRIOV CNI with DPDK to benchmark the user-space container network model in the containerized infrastructure worker node. The Multus CNI support creates multiple interfaces for a container. The traffic is bypassed the Kernel space by SRIOV with DPDK. We established two modes of CMK: shared core and exclusive core. We created VFs for each network interface of a container. Then, we set up TREX to route packet from eth1 to eth2 in a POD.
B.2. Trouble shooting and Results

Figure 17 shows the test results when using 1518 bytes packet traffic from the T-Rex traffic generator. First, we get the maximum line rate of the system using SR-IOV as the packet acceleration technique. Then we measured throughput when applying the CMK feature. We observed similar results as VPP CPU Pinning test. The default Kubernetes system without CMK feature enabled had the worst performance as the CPU resources are shared without any isolation. When the CMK feature is enabled, Exclusive Mode performed better than Shared Mode because each pod had its own dedicated CPU.

<table>
<thead>
<tr>
<th>Model</th>
<th>Result (Gbps)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Maximum Line Rate</td>
<td>39.3</td>
</tr>
<tr>
<td>Without CMK</td>
<td>11.5</td>
</tr>
<tr>
<td>CMK-Exclusive Mode</td>
<td>39.2</td>
</tr>
<tr>
<td>CMK-shared Mode</td>
<td>29.6</td>
</tr>
</tbody>
</table>

Figure 17: SR-IOV CPU Pinning Test Results

Appendix C. Benchmarking Experience (Multi-pod Test)

C.1. Benchmarking Overview

The main goal of this experience was to benchmark the multi-pod scenario, in which packets are traversed through two pods. To create additional interfaces for forwarding packets between two pods, Multus CNI was used. We compared two userspace-vSwitch model network technologies: OVS/DPDK and VPP-memif. Since vpp-memif has a different packet forwarding mechanism by using shared memory interface, it is expected that vpp-memif may provide higher performance than OVS-DPDK. Also, we consider NUMA impact for both cases, and made 6 scenarios depending on CPU location of vSwitch and two pods. Figure 18 is packet forwarding scenario in this test, where two pods run on the same host and vSwitch delivers packets between two pods.
C.2. Hardware Configurations
<table>
<thead>
<tr>
<th>Node Name</th>
<th>Specification</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Container Control</td>
<td>Intel(R) Core(TM) E5-2620v3 @ 2.40GHz (1socket x 12Cores)</td>
<td>Container Deployment and Network Allocation - ubuntu 18.04 - Kubernetes Master - CNI Controller - MULTUS CNI - DPDK-OVS/VPP-memif</td>
</tr>
<tr>
<td>for Master</td>
<td>MEM 32GB</td>
<td></td>
</tr>
<tr>
<td></td>
<td>DISK 1TB</td>
<td></td>
</tr>
<tr>
<td></td>
<td>NIC: Control plane: 1G</td>
<td></td>
</tr>
<tr>
<td></td>
<td>OS: CentOS Linux7.9</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Container Service</td>
<td>Intel(R) Xeon(R) Gold 6148 @ 2.40GHz (2socket X 40Cores)</td>
<td>Container dpdk-L2fwd - Kubernetes Worker - CNI Agent - Multus CNI - DPDK-OVS/VPP-memif</td>
</tr>
<tr>
<td>for Worker</td>
<td>MEM 256GB</td>
<td></td>
</tr>
<tr>
<td></td>
<td>DISK 2TB</td>
<td></td>
</tr>
<tr>
<td></td>
<td>NIC</td>
<td></td>
</tr>
<tr>
<td></td>
<td>Control plane: 1G</td>
<td></td>
</tr>
<tr>
<td></td>
<td>Data plane: XL710-qda2 (1NIC 2PORT- 40Gb)</td>
<td></td>
</tr>
<tr>
<td></td>
<td>OS: CentOS Linux 7.9</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Packet Generator</td>
<td>Intel(R) Xeon(R) Gold 6148 @ 2.4Ghz (2Socket X 40Core)</td>
<td>Packet Generator - Installed Trex v2.92</td>
</tr>
<tr>
<td></td>
<td>MEM 256GB</td>
<td></td>
</tr>
<tr>
<td></td>
<td>DISK 2TB</td>
<td></td>
</tr>
<tr>
<td></td>
<td>NIC</td>
<td></td>
</tr>
<tr>
<td></td>
<td>Data plane: XL710-qda2 (1NIC 2PORT - 40Gb)</td>
<td></td>
</tr>
<tr>
<td></td>
<td>OS: CentOS Lunix 7.9</td>
<td></td>
</tr>
</tbody>
</table>

Figure 19: Hardware Configurations for Multi-pod Benchmarking

For installations and configurations of CNIs, we used userspace-cni network plugin. Among this CNI, multus provides to create multiple interfaces for each pod. Both OVS-DPDK and VPP-memif bypass kernel with DPDK PMD driver. For CPU isolation and NUMA allocation, we used Intel CMK with exclusive mode. Since Trex generator is upgraded to the new version, we used the latest version of Trex.
C.3. NUMA Allocation Scenario

To analyze benchmarking impacts of different NUMA allocation, we set 6 scenarios depending on CPU location allocating to two pods and vSwitch. For this scenario, we did not consider cross-NUMA case, which allocates CPUs to pod or switch in a manner that two cores are located in different NUMA nodes. 6 scenarios we considered are listed in Table 1. Note that, NIC is attached to the NUMA1.

+============+=========+=======+=======+
| Scenario # | vSwitich |  pod1 |  pod2 |
|-------------+---------+-------+-------|
|     S1      |  NUMA1  | NUMA0 | NUMA0 |
|-------------+---------+-------+-------|
|     S2      |  NUMA1  | NUMA1 | NUMA1 |
|-------------+---------+-------+-------|
|     S3      |  NUMA0  | NUMA0 | NUMA0 |
|-------------+---------+-------+-------|
|     S4      |  NUMA0  | NUMA1 | NUMA1 |
|-------------+---------+-------+-------|
|     S5      |  NUMA1  | NUMA1 | NUMA0 |
|-------------+---------+-------+-------|
|     S6      |  NUMA0  | NUMA0 | NUMA1 |
|-------------+---------+-------+-------|

Table 1: NUMA Allocation Scenarios

C.4. Traffic Generator Configurations

For multi-pod benchmarking, we discovered Non Drop Rate (NDR) with binary search algorithm. In Trex, it supports command to discover NDR for each testing. Also, we test for different ethernet frame sizes from 64bytes to 1518bytes. For running Trex, we used command as follows;

```
./ndr --stl --port 0 1 -v --profile stl/bench.py --prof-tun size=x --opt-bin-search
```

C.5. Benchmark Results and Trouble-shootings

As the benchmarking results, Table 2 shows packet loss ratio using 1518 bytes packet in OVS-DPDK/vpp-memif. From that result, we can say that the vpp-memif has better performance than OVS-DPDK, which is came from the difference in the way to forward packets between vSwitch and pod. Also, the impact of NUMA is bigger when vSwitch and both pods are located in the same node than when allocating CPU to the node where NIC is attached.
<table>
<thead>
<tr>
<th>Networking Model</th>
<th>S1</th>
<th>S2</th>
<th>S3</th>
<th>S4</th>
<th>S5</th>
<th>S6</th>
</tr>
</thead>
<tbody>
<tr>
<td>vpp-memif</td>
<td>59.96</td>
<td>34.17</td>
<td>45.13</td>
<td>57.1</td>
<td>33.47</td>
<td>44.92</td>
</tr>
</tbody>
</table>

Table 2: Multi-pod Benchmarking Results (% of Line Rate)

Authors’ Addresses

Kyoungjae Sun  
ETRI  
218, Gajeong-ro, Yuseung-gu  
Dajeon  
34065  
Republic of Korea  
Phone: +82 10 3643 5627  
Email: kjsun@etri.re.kr

Hyunsik Yang  
KT  
KT Research Center 151  
Taebong-ro, Seocho-gu  
Seoul  
06763  
Republic of Korea  
Phone: +82 10 9005 7439  
Email: yangun@dcn.ssu.ac.kr

Jangwon Lee  
Soongsil University  
369, Sangdo-ro, Dongjak-gu  
Seoul  
06978  
Republic of Korea  
Phone: +82 10 7448 4664  
Email: jangwon.lee@dcn.ssu.ac.kr
Tran Minh Ngoc
Soongsil University
369, Sangdo-ro, Dongjak-gu
Seoul
06978
Republic of Korea
Phone: +82 2 820 0841
Email: mipearlskal307@dcn.ssu.ac.kr

Younghan Kim
Soongsil University
369, Sangdo-ro, Dongjak-gu
Seoul
06978
Republic of Korea
Phone: +82 10 2691 0904
Email: younghak@ssu.ac.kr
Abstract

This document defines methodologies for benchmarking EVPN and PBB-EVPN performance. EVPN is defined in RFC 7432, and is being deployed in Service Provider networks. Specifically, this document defines the methodologies for benchmarking EVPN/PBB-EVPN convergence, data plane performance, and control plane performance.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at https://datatracker.ietf.org/drafts/current/.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on December 20, 2021.

Copyright Notice

Copyright (c) 2021 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust’s Legal Provisions Relating to IETF Documents (https://trustee.ietf.org/license-info) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of
the Trust Legal Provisions and are provided without warranty as described in the Simplified BSD License.

Table of Contents

1. Introduction ........................................... 3
   1.1. Requirements Language ................................ 3
   1.2. Terminologies ...................................... 3
2. Test Topology .......................................... 4
3. Test Cases for EVPN Benchmarking ......................... 7
   3.1. Data Plane MAC Learning ............................. 7
   3.2. Control Plane MAC Learning .......................... 8
   3.3. MAC Flush-Local Link Failure and Relearning .......... 9
   3.4. MAC Flush-Remote Link Failure and Relearning ........ 10
   3.5. MAC Aging ........................................ 11
   3.6. Remote MAC Aging .................................... 12
   3.7. Control and Data plane MAC Learning ................ 12
   3.8. High Availability .................................. 13
   3.9. ARP/ND Scale ...................................... 14
   3.10. Scaling of Services ................................ 15
   3.11. Scale Convergence ................................ 16
   3.12. SOAK Test ......................................... 17
4. Test Cases for PBB-EVPN Benchmarking ..................... 18
   4.1. Data Plane Local MAC Learning ........................ 18
   4.2. Data Plane Remote MAC Learning ........................ 18
   4.3. MAC Flush-Local Link Failure ........................ 19
   4.4. MAC Flush-Remote Link Failure ........................ 20
   4.5. MAC Aging .......................................... 21
   4.6. Remote MAC Aging .................................... 22
   4.7. Local and Remote MAC Learning ........................ 23
   4.8. High Availability .................................. 23
   4.9. Scale ............................................. 24
   4.10. Scale Convergence ................................ 25
   4.11. Soak Test ......................................... 26
5. Acknowledgments ......................................... 27
6. IANA Considerations ...................................... 27
7. Security Considerations .................................. 27
8. References ............................................... 27
   8.1. Normative References ................................ 27
   8.2. Informative References ............................... 28
Appendix A. Appendix ....................................... 28
Authors’ Addresses .......................................... 28
1. Introduction

EVPN is defined in RFC 7432, and describes BGP MPLS based Ethernet VPNS (EVPN). PBB-EVPN is defined in RFC 7623, discusses how Ethernet Provider backbone Bridging can be combined with EVPNs to provide a new/combined solution. This draft defines methodologies that can be used to benchmark both RFC 7432 and RFC 7623 solutions. Further, this draft provides methodologies for benchmarking the performance of EVPN data and control planes, MAC learning, MAC flushing, MAC aging, convergence, high availability, and scale.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 8174 [RFC8174].

1.2. Terminologies

Most of the terminology used in this documents comes from [RFC7432] and [RFC7632].

All-Active Redundancy Mode: When all PEs attached to an Ethernet segment are allowed to forward known unicast traffic to/from that Ethernet segment for a given VLAN, then the Ethernet segment is defined to be operating in All-Active redundancy mode.

AA: All Active mode

CE: Customer Router/Devices/Switch.

DF: Designated Forwarder

DUT: Device under test.

Ethernet Segment (ES): When a customer site (device or network) is connected to one or more PEs via a set of Ethernet links, then that set of links is referred to as an ‘Ethernet segment’.

EVI: An EVPN instance spanning the Provider Edge (PE) devices participating in that EVPN.

Ethernet Segment Identifier (ESI): A unique non-zero identifier that identifies an Ethernet segment is called an ‘Ethernet Segment Identifier’.
Ethernet Tag: An Ethernet tag identifies a particular broadcast domain, e.g., a VLAN. An EVPN instance consists of one or more broadcast domains.

Interface: Physical interface of a router/switch.

IRB: Integrated routing and bridging interface

MAC: Media Access Control addresses on a PE.

MHPE2: Multi homed Provider Edge router 2.

MHPE1: Multi homed Provider Edge router 1.

SHPE3: Single homed Provider Edge Router 3.

PE: Provider Edge device.

P: Provider Router.

RR: Route Reflector.

RT: Traffic Generator.

Sub Interface: Each physical Interfaces is subdivided into a set of Logical units.

SA: Single Active

Single-Active Redundancy Mode: When a single PE (among all the PEs attached to an Ethernet segment) is the only PE allowed to forward traffic to/from a given Ethernet segment for a given VLAN, then that Ethernet segment is defined to be operating in Single-Active redundancy mode.

2. Test Topology

There are five routers in the Test setup. SHPE3, RR/P, MHPE1 and MHPE2 emulating a service provider network. CE is a customer device connected to MHPE1 and MHPE2. It is configured with bridge domains in multiple VLANS. The traffic generator is connected to the CE and SHPE3. The MHPE1 acts as DUT. The traffic generator will be used as sender and receiver of traffic. The test measurements are taken from the DUT. MHPE1 and MHPE2 are multi-homed routers connected to CE running single active mode. The traffic generator will be generating traffic at 10% of the line rate.
SHPE3  +----------------+  Traffic Generator sender/receiver of layer 2 traffic with multiple Vlans
        |                |                           |
        |                |                | Core Link
        |                |                | RR/P  +-------------------+ Core link
        |                |                | core link
        |                |                +-------------------+
MHPE1(DUT) | MHPE2
        +-------------------+-------------------+
        | PE-CE link       | PE-CE link
        +-------------------+-------------------+

-----+  |  |  +---------------------+  Traffic Generator sender/receiver of layer 2 traffic with multiple Vlans
er/ |  |  | Ce/Layer 2 bridge
ffic|  |  | +---------------------+

Topology 1

Test Setup

Figure 1
Table showing the traffic directions of various EVPN/PBB-EVPN benchmarking test cases. Depending on the test scenario, the traffic can be uni-directional or bi-directional (configured in the traffic generator).

**Figure 2**

Test Setup Configurations:

SHPE3 is configured with Interior Gateway protocols like OPSF or IS-
IS for underlay, LDP for MPLS support, Interior Border Gateway with EVPN address family for overlay support. This router must be configured with N EVPN/PBB-EVPN instances for testing. Traffic generator is connected to this router for sending and receiving traffic.

RR is configured with Interior Gateway protocols like OSPF or IS-IS for underlay, LDP for MPLS support, Interior Border Gateway with EVPN address family for overlay support. This router function as both provider router and a route reflector.
MHPE1 is configured with Interior Gateway protocols like OSPF or IS-IS for underlay, LDP for MPLS support, Interior Border Gateway with EVPN address family for overlay support. This router must be configured with N EVPN/PBB-EVPN instances for testing. This router is configured with ESI per vlan or ESI per interface. It is functioning as multi homing PE working on Single Active EVPN mode. This router serves as the DUT and it is connected to CE. MHPE1 is acting as DUT for all the test cases.

MHPE2 is configured with Interior Gateway protocols like OSPF or IS-IS for underlay, LDP for MPLS support, Interior Border Gateway with EVPN address family for overlay support. This router must be configured with N EVPN/PBB-EVPN instances for testing. This router is configured with ESI per vlan or ESI per interface. It is functioning as multi homing PE working on Single Active EVPN mode. It is connected to CE. It acts as standby router during the tests.

CE is acting as bridge configured with multiple vlans. The same vlans are configured on MHPE1, MHPE2, SHPE3. Traffic generator is connected to CE. The traffic generator acts as sender or receiver of traffic.

Depending up on the test scenarios the traffic generators will be used to generate uni directional or bi directional flows.

The above configuration will be serving as the base configuration for all test cases.

The X below is used as variable to denote scale factor of the testing parameters. It must be in the multiples of 100.

3. Test Cases for EVPN Benchmarking

3.1. Data Plane MAC Learning

Objective:

Measure the time taken to learn the Data Plane MAC in DUT.

Topology : Topology 1

Procedure:

The data plane MAC learning can be measured using the parameters defined in RFC 2889 section 5.8.

Confirm the DUT is up and running with EVPN.
Traffic generator connected to CE must send frames with X different source and destination MAC address for one vlan, the same vlan must be present in all the devices except RR.

Send X unicast frames from CE to MHPE1(DUT) for one EVPN instance working in SA mode.

The DUT will learn these X MAC in data plane.

Measurement:

Measure the time taken to learn X MAC locally in DUT EVPN MAC table. The data plane measurement is taken by considering DUT as black box. The range of MAC are known from traffic generator, the same must be learned in DUT, the time taken to learn X MAC is measured. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp. Note at present there are no standard models to measure the counters.

The test is repeated for N times and the values are collected. The MAC learning rate is calculated by averaging the values obtained from N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1,T2...Tn.

MAC learning rate = (T1+T2+..Tn)/N

3.2. Control Plane MAC Learning

Objective:

Measure the time taken to learn the control plane MAC.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with EVPN.

Traffic generator connected to SHPE3 must send frames with X different source and destination MAC address for one vlan, the same vlan must be present in all the devices except RR.

Ensure the frames must be destined to one EVPN instance.

The DUT will learn these X MAC in control plane.
Measurement:

Measure the time taken by the DUT to learn the X MAC in the data plane. The test is repeated for N times and the values are collected. The remote MAC learning rate is calculated by averaging the values obtained from N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1, T2...Tn. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp. Note at present there are no standard models to measure the counters.

MAC learning rate = (T1+T2+..Tn)/N

3.3. MAC Flush-Local Link Failure and Relearning

Objective:

Measure the time taken to flush the Data Plane MAC and the time taken to relearn the same amount of MAC.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with EVPN.

Send X frames with X different source and destination MAC addresses to DUT from CE using traffic generator for one vlan.

Ensure the DUT learns all X MAC addresses in data plane.

Fail the DUT-CE link and measure the time taken to flush these X MAC from the EVPN MAC table.

Bring up the link which was made Down (the link between DUT and CE). Measure time taken by the DUT to relearn these X MAC.

The DUT and MHPE2 are running SA mode.

Measurement:

Measure the time taken for flushing these X MAC addresses. Measure the time taken to relearn these X MAC in DUT. The test is repeated for N times and the values are collected. The flush and the relearning time is calculated by averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The
time measured for each sample is denoted by T1, T2, ..., Tn. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp.

\[ \text{Flush rate} = \frac{(T_1 + T_2 + \cdots + T_n)}{N} \]

\[ \text{Relearning rate} = \frac{(T_1 + T_2 + \cdots + T_n)}{N} \]

3.4. MAC Flush-Remote Link Failure and Relearning.

Objective:

Measure the time taken to flush the Control plane MAC learned in DUT during remote link failure and the time taken to relearn.

Topology: Topology 1

Procedure:

Confirm the DUT is up and running with EVPN.

Send X frames with X different source and destination MAC addresses to DUT from SHPE3 using traffic generator for one vlan.

Bring down the link between SHPE3 and traffic generator.

SHPE3 will withdraw the routes from DUT due to link failure.

Measure the time taken to flush the DUT EVPN MAC table. The DUT and MHPE2 are running SA mode.

Bring up the link which was made Down (the link between SHPE3 and traffic generator).

Measure time taken by the DUT to relearn these X MAC from control plane.

Measurement:

Measure the time taken to flush X remote MAC from EVPN MAC table of the DUT. Measure the time taken to relearn these X MAC in DUT. The test is repeated for N times and the values are collected. The flush rate is calculated by averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1, T2, ..., Tn. The measurement is carried
out using external server which polls the DUT using automated scripts
which measures the MAC counter value with timestamp.

Flush rate = (T1+T2+..Tn)/N
Relearning rate = (T1+T2+..Tn)/N

3.5. MAC Aging

Objective:
To measure the MAC Aging time.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with EVPN.

Send X frames with X different source and destination MAC addresses
to DUT from CE using traffic generator for one vlan.

Ensure these X MAC addresses are learned in DUT.

Then stop the traffic.

Ensure the DUT and other devices in the test are using the default
timers for aging.

Measure the time taken to flush X MAC from DUT EVPN MAC table due to
aging.

The DUT and MHPE2 are running SA mode.

Measurement :

Measure the time taken to flush X MAC addresses due to aging. The
test is repeated for N times and the values are collected. The aging
is calculated by averaging the values obtained by N samples. N is an
arbitrary number to get a sufficient sample. The time measured for
each sample is denoted by T1,T2...Tn. The measurement is carried out
using external server which polls the DUT using automated scripts
which measures the MAC counter value with timestamp.

Aging time for X MAC in sec = (T1+T2+..Tn)/N
3.6. Remote MAC Aging

Objective:
Measure the control plane learned MAC Aging time.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with EVPN.

Send X frames with X different source and destination MAC addresses to DUT from SHPE3 using traffic generator for one vlan.

Ensure these X MAC addresses are learned in DUT via control plane.

Then stop the traffic.

Ensure the DUT and other devices in the test are using the default timers for aging.

Measure the time taken to flush X MAC from DUT EVPN MAC table due to aging.

The DUT and MHPE2 are running SA mode.

Measurement:

Measure the time taken to flush X remote MAC learned in DUT EVPN MAC table due to aging. The test is repeated for N times and the values are collected. The aging is calculated by averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1, T2...Tn. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp.

Aging time for X MAC in sec = (T1+T2+..Tn)/N

3.7. Control and Data plane MAC Learning

Objective:

To record the time taken to learn both local and remote MAC.

Topology : Topology 1
Procedure:

Confirm the DUT is up and running with EVPN.

Send X frames with X different source and destination MAC addresses to DUT from SHPE3 using traffic generator for one vlan.

Send X frames with different source and destination MAC addresses from traffic generator connected to CE for one vlan.

The source and destination addresses of flows must be complimentary to have unicast flows.

Measure the time taken by the DUT to learn 2X in EVPN MAC table.

DUT and MHPE2 are running in SA mode.

Measurement:

Measure the time taken to learn 2X MAC addresses in DUT EVPN MAC table. The test is repeated for N times and the values are collected. The MAC learning time is calculated by averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1, T2...Tn. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp.

\[
\text{MAC learning rate} = \frac{(T1+T2+..Tn)}{N}
\]


Objective:

Measure traffic loss during routing engine fail over.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with EVPN.

Send X frames from CE to DUT from traffic generator with X different source and destination MAC addresses.
Send X frames from traffic generator to SHPE3 with X different source and destination MAC addresses, so that 2X MAC address will be learned in the DUT.

There is a bi directional traffic flow with X pps in each direction.

Ensure the DUT learn 2X MAC.

Then do a routing engine fail-over.

Measurement:

The expectation of the test is 0 traffic loss with no change in the DF role. DUT should not withdraw any routes. But in cases where the DUT is not property synchronised between master and standby, due to that packet loss are observed. In that scenario the packet loss is measured. The test is repeated for N times and the values are collected. The packet loss is calculated by averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1, T2...Tn. The measurement is carried out using external server which polls the DUT using automated scripts to ensure the DUT learned 2X MAC. The packet drop is measured using traffic generator.

Packet loss in sec with 2X MAC addresses = (T1+T2+..Tn)/N

3.9. ARP/ND Scale

Measure the DUT scaling limit of ARP/ND.

Objective:

Measure the ARP/ND scale of the DUT.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with EVPN.

Send X arp/neighbour discovery(ND) from the traffic generator to DUT with different sender ip/ipv6, MAC addresses to the target IRB address configured in EVPN instance.

The EVPN instance learns the MAC+ip and MAC+ipv6 addresses from these request and advertise as type 2 MAC+ip/MAC+ipv6 route to remote provide edge routers which have same EVPN configurations.
The value of X must be increased at a incremental value of 5% of X, till the limit is reached. The limit is where the DUT cant learn any more type 2 MAC+ip/MAC+ipv6. The test must be separately conducted for arp and ND.

Measurement:

Measure the scale limit of type 2 MAC+ip/MAC+ipv6 route which DUT can learn. The test is repeated for N times and the values are collected. The scale limit is calculated by averaging the values obtained by N samples for both MAC+ip and MAC+ipv6. N is an arbitrary number to get a sufficient sample. The scale value obtained by each sample be v1,v2..vn for MAC+ipv4 and s1,s2..sn for MAC+ipv6. The measurement is carried out using external server which polls the DUT using automated scripts to find the scale limit of MAC+ipv4/MAC+ipv6.

Scale limit for MAC+ip = (v1+v2+..vn)/N
Scale limit for MAC+ipv6 = (s1+s2+..sn)/N

3.10. Scaling of Services

Objective:

Measure the scale of EVPN instances that a DUT can hold.

Topology: Topology 1

Procedure:

Confirm the DUT is up and running with EVPN.

The DUT, MHPE2 and SHPE3 are scaled to N EVI.

Ensure routes received from MHPE2 and SHPE3 for N EVI in the DUT.

Then increment the scale of N by 5% of N till the limit is reached.

The limit is where the DUT cant learn any EVPN routes from its peers.

Measurement:

There should not be any loss of route types 1, 2, 3 and 4 in DUT. DUT must relearn all type 1, 2, 3 and 4 from remote routers. The DUT must be subjected to various values of N to find the optimal scale limit. The scope of the test is find out the maximum EVPN instance
that a DUT can hold. The measurement is carried out using external server which polls the DUT using automated scripts to find the scale limit of EVPN instances.

3.11. Scale Convergence

Objective:

Measure the convergence time of DUT when the DUT is scaled with EVPN instance along with traffic.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with EVPN.

Scale N EVIs in DUT, SHPE3 and MHPE2.

Send F frames to DUT from CE using traffic generator with X different source and destination MAC addresses for N EVI’s.

Send F frames from traffic generator to SHPE3 with X different source and destination MAC addresses.

There will be 2X number of MAC addresses will be learned in DUT EVPN MAC table.

There is a bi directional traffic flow with F pps in each direction.

Then clear the BGP neighbours in the DUT.

Once the BGP session is in established state in DUT.

Measure the time taken to learn 2X MAC address in DUT MAC table.

Measurement :

The DUT must learn 2X MAC addresses. Measure the time taken to learn 2X MAC in DUT. The test is repeated for N times and the values are collected. The convergence time is calculated by averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by $T_1, T_2, \ldots, T_n$. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp.
3.12. SOAK Test.

Objective:

This test is carried out to measure the stability of the DUT in a scaled environment with traffic over a period of time "T'". In each interval "t1" the DUT CPU usage, memory usage are measured. The DUT is checked for any crashes during this time period.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with EVPN.

Scale N EVI’s in DUT, SHPE3 and MHPE2. Send F frames to DUT from CE using traffic generator with different X source and destination MAC addresses for N EVI’s.

Send F frames from traffic generator to SHPE3 with X different source and destination MAC addresses.

There will be 2X number of MAC addresses will be learned in DUT EVPN MAC table.

There is a bi directional traffic flow with F pps in each direction.

The DUT must run with traffic for 24 hours.

Every hour check for memory leak in EVPN process, CPU usage and crashes in DUT.

Measurement :

Take the hourly reading of CPU, process memory. There should not be any memory leak, crashes, CPU spikes. The CPU spike is determined as the sudden increase of CPU usage to 100 percent compared to the average usage. The average value vary from device to device. Memory leak is determined as the increase of memory usage with respect to time. The expectation is under steady state the memory usage for EVPN process should not increase with respect to time. The measurement is carried out using external server which polls the DUT using automated scripts which captures the CPU usage and process memory.
4. Test Cases for PBB-EVPN Benchmarking

4.1. Data Plane Local MAC Learning

Objective:
Measure the time taken to learn the Data Plane MAC in DUT.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with PBB-EVPN.

Traffic generator connected to CE must send frames with X different source and destination MAC address for one vlan, the same vlan must be present in all the devices except RR.

Send X unicast frames from CE to MHPE1(DUT) for one PBB-EVPN instance working in SA mode.

The DUT will learn these X MAC in data plane.

Measurement :

Measure the time taken to learn X MAC locally in DUT PBB-EVPN MAC table. The data plane measurement is taken by considering DUT as black box. The range of MAC are known from traffic generator, the same must be learned in DUT, the time taken to learn X MAC is measured.The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp.

The test is repeated for N times and the values are collected. The MAC learning rate is calculated by averaging the values obtained from N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1,T2...Tn.

MAC learning rate = (T1+T2+...Tn)/N

4.2. Data Plane Remote MAC Learning

Objective:
To Record the time taken to learn the remote MAC.

Topology : Topology 1
Procedure:

Confirm the DUT is up and running with PBB-EVPN.

Traffic generator connected to SHPE3 must send frames with X different source and destination MAC address for one vlan, the same vlan must be present in all the devices except RR.

Ensure the frames must be destined to one PBB-EVPN instance.

The DUT will learn these X MAC in data plane.

Measurement:

Measure the time taken by the DUT to learn the X MAC in the data plane. The test is repeated for N times and the values are collected. The remote MAC learning rate is calculated by averaging the values obtained from N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1, T2...Tn. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp.

MAC learning rate = (T1+T2+..Tn)/N

4.3. MAC Flush-Local Link Failure

Objective:

Measure the time taken to flush the locally learned MAC and the time taken to relearn the same amount of MAC.

Topology: Topology 1

Procedure:

Confirm the DUT is up and running with PBB-EVPN.

Send X frames with X different source and destination MAC addresses to DUT from CE using traffic generator for one vlan.

Ensure the DUT learns all X MAC addresses in data plane.

Fail the DUT-CE link and measure the time taken to flush these X MAC from the PBB-EVPN MAC table.
Bring up the link which was made Down (the link between DUT and CE). Measure time taken by the DUT to relearn these X MAC.

The DUT and MHPE2 are running SA mode.

Measurement:

Measure the time taken for flushing these X MAC addresses. Measure the time taken to relearn these X MAC in DUT. The test is repeated for N times and the values are collected. The flush and the relearning time is calculated by averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1, T2... Tn. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp.

Flush rate = (T1+T2+...Tn)/N

Relearning rate = (T1+T2+...Tn)/N

4.4. MAC Flush-Remote Link Failure

Objective:

Measure the time taken to flush the remote MAC learned in DUT due to remote link failure and relearning it.

Topology: Topology 1

Procedure:

confirm the DUT is up and running with PBB-EVPN.

Send X frames with X different source and destination MAC addresses to DUT from SHPE3 using traffic generator for one vlan.

Bring down the link between SHPE3 and traffic generator.

Measure the time taken to flush the DUT PBB-EVPN MAC table. The DUT and MHPE2 are running SA mode.

Bring up the link which was made Down (the link between SHPE3 and traffic generator).

Measure time taken by the DUT to relearn these X MAC.
Measurement:

Measure the time taken to flush X remote MAC from PBB-EVPN MAC table of the DUT. Measure the time taken to relearn these X MAC in DUT. The test is repeated for N times and the values are collected. The flush rate is calculated by averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1,T2...Tn. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp.

Flush rate = \(\frac{T1+T2+..+Tn}{N}\)

Relearning rate = \(\frac{T1+T2+..+Tn}{N}\)

4.5. MAC Aging

Objective:

Measure the MAC Aging time.

Topology: Topology 1

Procedure:

Confirm the DUT is up and running with PBB-EVPN.

Send X frames with X different source and destination MAC addresses to DUT from CE using traffic generator for one vlan.

Ensure these X MAC addresses are learned in DUT.

Then stop the traffic.

Ensure the DUT and other devices in the test are using the default timers for aging.

Measure the time taken to flush X MAC from DUT PBB-EVPN MAC table due to aging.

The DUT and MHPE2 are running SA mode.

Measurement:

Measure the time taken to flush X MAC addresses due to aging. The test is repeated for N times and the values are collected. The aging
is calculated averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1, T2...Tn. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp.

Aging time for X MAC in sec = (T1+T2+..Tn)/N


Objective:
Measure the remote MAC Aging time.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with PBB-EVPN.

Send X frames with X different source and destination MAC addresses to DUT from SHPE3 using traffic generator for one vlan.

Ensure these X MAC addresses are learned in DUT.

Then stop the traffic.

Ensure the DUT and other devices in the test are using the default timers for aging.

Measure the time taken to flush X MAC from DUT PBB-EVPN MAC table due to aging.

The DUT and MHPE2 are running SA mode.

Measurement:

Measure the time taken to flush X remote MAC learned in DUT EVPN MAC table due to aging. The test is repeated for N times and the values are collected. The aging is calculated by averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1, T2...Tn. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp.

Aging time for X MAC in sec = (T1+T2+..Tn)/N
4.7. Local and Remote MAC Learning

Objective:

Measure the time taken to learn both local and remote MAC.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with PBB-EVPN.

Send X frames with X different source and destination MAC addresses to DUT from SHPE3 using traffic generator for one vlan.

Send X frames with different source and destination MAC addresses from traffic generator connected to CE for one vlan.

The source and destination addresses of flows must be complimentary to have unicast flows.

Measure the time taken by the DUT to learn 2X in PBB-EVPN MAC table.

DUT and MHPE2 are running in SA mode.

Measurement:

Measure the time taken to learn 2X MAC addresses in DUT PBB-EVPN MAC table. The test is repeated for N times and the values are collected. The MAC learning time is calculated by averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T1,T2,...Tn. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp.

MAC learning rate = (T1+T2+..Tn)/N

4.8. High Availability

Objective:

Measure traffic loss during routing engine failover.

Topology : Topology 1

Procedure:
Confirm the DUT is up and running with PBB-EVPN.

Send X frames from CE to DUT from traffic generator with X different source and destination MAC addresses.

Send X frames from traffic generator to SHPE3 with X different source and destination MAC addresses, so that 2X MAC address will be learned in the DUT.

There is a bi-directional traffic flow with X pps in each direction.

Ensure the DUT learn 2X MAC.

Then do a routing engine fail-over.

Measurement:

The expectation of the test is 0 traffic loss with no change in the DF role. DUT should not withdraw any routes. But in cases where the DUT is not properly synchronized between master and standby, due to that packet loss are observed. In that scenario the packet loss is measured. The test is repeated for N times and the values are collected. The packet loss is calculated by averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by $T_1, T_2, \ldots, T_n$. The measurement is carried out using an external server which polls the DUT using automated scripts to ensure the DUT learned 2X MAC. The packet drop is measured using a traffic generator.

Packet loss in sec with 2X MAC addresses = $\frac{(T_1+T_2+\ldots+T_n)}{N}$

4.9. Scale

Objective:

Measure the scale limit of DUT for PBB-EVPN.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with PBB-EVPN.

The DUT, MHPE2 and SHPE3 are scaled to N PBB-EVI.

Ensure routes received from MHPE2 and SHPE3 for N PBB-EVI in the DUT.
Then increment the scale of N by 5% of N till the limit is reached.

The limit is where the DUT cant learn any EVPN routes from its peers.

Measurement:

There should not be any loss of route types 2, 3 and 4 in DUT. DUT must relearn all type 2, 3 and 4 from remote routers. The DUT must be subjected to various values of N to find the optimal scale limit. The scope of the test is find out the maximum EVPN instance that a DUT can hold. The measurement is carried out using external server which polls the DUT using automated scripts to find the scale limit of PBB-EVPN instances.

4.10. Scale Convergence

Objective:

To measure the convergence time of DUT when the DUT is scaled with EVPN instance along with traffic.

Topology: Topology 1

Procedure:

Confirm the DUT is up and running with PBB-EVPN.

Scale N PBB-EVIs in DUT, SHPE3 and MHPE2.

Send F frames to DUT from CE using traffic generator with X different source and destination MAC addresses for N PBB-EVI’s.

Send F frames from traffic generator to SHPE3 with X different source and destination MAC addresses.

There will be 2X number of MAC addresses will be learned in DUT PBB-EVPN MAC table.

There is a bi directional traffic flow with F pps in each direction.

Then clear the BGP neighbours in the DUT.

Once the BGP session is in established state in DUT.

Measure the time taken to learn 2X MAC address in DUT MAC table.

Measurement:
The DUT must learn 2X MAC addresses. Measure the time taken to learn 2X MAC in DUT. The test is repeated for N times and the values are collected. The convergence time is calculated by averaging the values obtained by N samples. N is an arbitrary number to get a sufficient sample. The time measured for each sample is denoted by T₁, T₂...Tₙ. The measurement is carried out using external server which polls the DUT using automated scripts which measures the MAC counter value with timestamp.

\[
\text{Time taken to learn 2X MAC in DUT} = \frac{(T₁+T₂+..Tₙ)}{N}
\]

4.11. Soak Test

Objective:

To measure the stability of the DUT in a scaled environment with traffic.

Topology : Topology 1

Procedure:

Confirm the DUT is up and running with PBB-EVPN.

Scale N PBB-EVI’s in DUT, SHPE3 and MHPE2. Send F frames to DUT from CE using traffic generator with different X source and destination MAC addresses for N EVI’s.

Send F frames from traffic generator to SHPE3 with X different source and destination MAC addresses.

There will be 2X number of MAC addresses will be learned in DUT PBB-EVPN MAC table.

There is a bi directional traffic flow with F pps in each direction.

The DUT must run with traffic for 24 hours.

Every hour check for memory leak in PBB-EVPN process, CPU usage and crashes in DUT.

Measurement:

Take the hourly reading of CPU, process memory. There should not be any memory leak, crashes, CPU spikes. The CPU spike is determined as the sudden increase of CPU usage to 100 percent compared to the average usage. The average value vary from device to device. Memory
leak is determined as the increase of memory usage with respect to time. The expectation is under steady state the memory usage for PBB-EVPN process should not increase with respect to time. The measurement is carried out using external server which polls the DUT using automated scripts which captures the CPU usage and process memory.

5. Acknowledgments

We would like to thank Fioccola Giuseppe of Telecom Italia reviewing our draft and commenting it. We would like to thank Sarah Banks for guiding and mentoring us. We take the opportunity to thank Al for reviewing our draft and gave us valuable comments.

6. IANA Considerations

This memo includes no request to IANA.

7. Security Considerations

The benchmarking tests described in this document are limited to the performance characterisation of controllers in a lab environment with isolated networks. The benchmarking network topology will be an independent test setup and MUST NOT be connected to devices that may forward the test traffic into a production network or misroute traffic to the test management network. Further, benchmarking is performed on a "black-box" basis, relying solely on measurements observable external to the controller. Security features mentioned in the RFC 7432 will affect the test results. Special capabilities SHOULD NOT exist in the controller specifically for benchmarking purposes. Any implications for network security arising from the controller SHOULD be identical in the lab and in production networks.

8. References

8.1. Normative References


8.2. Informative References


Appendix A. Appendix

Authors’ Addresses

Sudhin Jacob (editor)
Independent
Bangalore
India

Phone: +91 8061212543
Email: sudhinjacob@rediffmail.com

Kishore Tiruveedhula
Juniper Networks
10 Technology Park Dr
Westford, MA  01886
USA

Phone: +1 9785890861
Email: kishoret@juniper.net
Multiple Loss Ratio Search for Packet Throughput (MLRsearch)
draft-ietf-bmwg-mlrsearch-00

Abstract

This document proposes changes to [RFC2544], specifically to packet throughput search methodology, by defining a new search algorithm referred to as Multiple Loss Ratio search (MLRsearch for short). Instead of relying on binary search with pre-set starting offered load, it proposes a novel approach discovering the starting point in the initial phase, and then searching for packet throughput based on defined packet loss ratio (PLR) input criteria and defined final trial duration time. One of the key design principles behind MLRsearch is minimizing the total test duration and searching for multiple packet throughput rates (each with a corresponding PLR) concurrently, instead of doing it sequentially.

The main motivation behind MLRsearch is the new set of challenges and requirements posed by NFV (Network Function Virtualization), specifically software based implementations of NFV data planes. Using [RFC2544] in the experience of the authors yields often not repetitive and not replicable end results due to a large number of factors that are out of scope for this draft. MLRsearch aims to address this challenge in a simple way of getting the same result sooner, so more repetitions can be done to describe the replicability.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at https://datatracker.ietf.org/drafts/current/.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."
1. Terminology

- Frame size: size of an Ethernet Layer-2 frame on the wire, including any VLAN tags (dot1q, dot1ad) and Ethernet FCS, but excluding Ethernet preamble and inter-frame gap. Measured in bytes.

- Packet size: same as frame size, both terms used interchangeably.
Device Under Test (DUT): In software networking, "device" denotes a specific piece of software tasked with packet processing. Such device is surrounded with other software components (such as operating system kernel). It is not possible to run devices without also running the other components, and hardware resources are shared between both. For purposes of testing, the whole set of hardware and software components is called "system under test" (SUT). As SUT is the part of the whole test setup performance of which can be measured by [RFC2544] methods, this document uses SUT instead of [RFC2544] DUT. Device under test (DUT) can be re-introduced when analysing test results using whitebox techniques, but this document sticks to blackbox testing.

System Under Test (SUT): System under test (SUT) is a part of the whole test setup whose performance is to be benchmarked. The complete test setup contains other parts, whose performance is either already established, or not affecting the benchmarking result.

Bi-directional throughput tests: involve packets/frames flowing in both transmit and receive directions over every tested interface of SUT/DUT. Packet flow metrics are measured per direction, and can be reported as aggregate for both directions and/or separately for each measured direction. In most cases bi-directional tests use the same (symmetric) load in both directions.

Uni-directional throughput tests: involve packets/frames flowing in only one direction, i.e. either transmit or receive direction, over every tested interface of SUT/DUT. Packet flow metrics are measured and are reported for measured direction.

Packet Loss Ratio (PLR): ratio of packets received relative to packets transmitted over the test trial duration, calculated using formula: PLR = ( pkts_transmitted - pkts_received ) / pkts_transmitted. For bi-directional throughput tests aggregate PLR is calculated based on the aggregate number of packets transmitted and received.

Packet Throughput Rate: maximum packet offered load DUT/SUT forwards within the specified Packet Loss Ratio (PLR). In many cases the rate depends on the frame size processed by DUT/SUT. Hence packet throughput rate MUST be quoted with specific frame size as received by DUT/SUT during the measurement. For bi-directional tests, packet throughput rate should be reported as aggregate for both directions. Measured in packets-per-second (pps) or frames-per-second (fps), equivalent metrics.
Internet-Draft

Multiple Loss Ratio Search for Packet Throughput

February 2021

o Bandwidth Throughput Rate: a secondary metric calculated from packet throughput rate using formula: \( bw\_rate = \text{pkt}\_rate \times (\text{frame}\_size + \text{L1}\_\text{overhead}) \times 8 \), where \( \text{L1}\_\text{overhead} \) for Ethernet includes preamble (8 bytes) and inter-frame gap (12 bytes). For bi-directional tests, bandwidth throughput rate should be reported as aggregate for both directions. Expressed in bits-per-second (bps).

o Non Drop Rate (NDR): maximum packet/bandwidth throughput rate sustained by DUT/SUT at PLR equal zero (zero packet loss) specific to tested frame size(s). MUST be quoted with specific packet size as received by DUT/SUT during the measurement. Packet NDR measured in packets-per-second (or fps), bandwidth NDR expressed in bits-per-second (bps).

o Partial Drop Rate (PDR): maximum packet/bandwidth throughput rate sustained by DUT/SUT at PLR greater than zero (non-zero packet loss) specific to tested frame size(s). MUST be quoted with specific packet size as received by DUT/SUT during the measurement. Packet PDR measured in packets-per-second (or fps), bandwidth PDR expressed in bits-per-second (bps).

o Maximum Receive Rate (MRR): packet/bandwidth rate regardless of PLR sustained by DUT/SUT under specified Maximum Transmit Rate (MTR) packet load offered by traffic generator. MUST be quoted with both specific packet size and MTR as received by DUT/SUT during the measurement. Packet MRR measured in packets-per-second (or fps), bandwidth MRR expressed in bits-per-second (bps).

o Trial: a single measurement step. See [RFC2544] section 23.

o Trial duration: amount of time over which packets are transmitted in a single measurement step.

2. MLRsearch Background

Multiple Loss Ratio search (MLRsearch) is a packet throughput search algorithm suitable for deterministic systems (as opposed to probabilistic systems). MLRsearch discovers multiple packet throughput rates in a single search, with each rate associated with a distinct Packet Loss Ratio (PLR) criteria.

For cases when multiple rates need to be found, this property makes MLRsearch more efficient in terms of time execution, compared to traditional throughput search algorithms that discover a single packet rate per defined search criteria (e.g. a binary search specified by [RFC2544]). MLRsearch reduces execution time even further by relying on shorter trial durations of intermediate steps,
with only the final measurements conducted at the specified final trial duration. This results in the shorter overall search execution time when compared to a traditional binary search, while guaranteeing the same results for deterministic systems.

In practice two rates with distinct PLRs are commonly used for packet throughput measurements of NFV systems: Non Drop Rate (NDR) with PLR=0 and Partial Drop Rate (PDR) with PLR>0. The rest of this document describes MLRsearch for NDR and PDR. If needed, MLRsearch can be adapted to discover more throughput rates with different pre-defined PLRs.

Similarly to other throughput search approaches like binary search, MLRsearch is effective for SUTs/DUTs with PLR curve that is continuously flat or increasing with growing offered load. It may not be as effective for SUTs/DUTs with abnormal PLR curves.

MLRsearch relies on traffic generator to qualify the received packet stream as error-free, and invalidate the results if any disqualifying errors are present e.g. out-of-sequence frames.

MLRsearch can be applied to both uni-directional and bi-directional throughput tests.

For bi-directional tests, MLRsearch rates and ratios are aggregates of both directions, based on the following assumptions:

- Traffic transmitted by traffic generator and received by SUT/DUT has the same packet rate in each direction, in other words the offered load is symmetric.
- SUT/DUT packet processing capacity is the same in both directions, resulting in the same packet loss under load.

3. MLRsearch Overview

The main properties of MLRsearch:

- MLRsearch is a duration aware multi-phase multi-rate search algorithm:
  * Initial Phase determines promising starting interval for the search.
  * Intermediate Phases progress towards defined final search criteria.
* Final Phase executes measurements according to the final search criteria.

* Final search criteria are defined by following inputs:
  + PLRs associated with NDR and PDR.
  + Final trial duration.
  + Measurement resolution.

o Initial Phase:
  * Measure MRR over initial trial duration.
  * Measured MRR is used as an input to the first intermediate phase.

o Multiple Intermediate Phases:
  * Trial duration:
    + Start with initial trial duration in the first intermediate phase.
    + Converge geometrically towards the final trial duration.
  * Track two values for NDR and two for PDR:
    + The values are called lower_bound and upper_bound.
    + Each value comes from a specific trial measurement:
      - Most recent for that transmit rate.
      - As such the value is associated with that measurement’s duration and loss.
    + A bound can be valid or invalid:
      - Valid lower_bound must conform with PLR search criteria.
      - Valid upper_bound must not conform with PLR search criteria.
      - Example of invalid NDR lower_bound is if it has been measured with non-zero loss.
Invalid bounds are not real boundaries for the searched value:
  - They are needed to track interval widths.
- Valid bounds are real boundaries for the searched value.
- Each non-initial phase ends with all bounds valid.
- Bound can become invalid if it re-measured at a longer trial duration in a subsequent phase.

* Search:
  + Start with a large (lower_bound, upper_bound) interval width, that determines measurement resolution.
  + Geometrically converge towards the width goal of the phase.
  + Each phase halves the previous width goal.
    - First measurement of the next phase will be internal search which always gives a valid bound and brings the width to the new goal.
    - Only one bound then needs to be re-measured with new duration.

* Use of internal and external searches:
  + External search:
    - Measures at transmit rates outside the (lower_bound, upper_bound) interval.
    - Activated when a bound is invalid, to search for a new valid bound by multiplying (for example doubling) the interval width.
    - It is a variant of "exponential search".
  + Internal search:
    - A "binary search" that measures at transmit rates within the (lower_bound, upper_bound) valid interval, halving the interval width.

  o Final Phase:
The main benefits of MLRsearch vs. binary search include:

- In general MLRsearch is likely to execute more trials overall, but likely less trials at a set final trial duration.
- In well behaving cases, e.g. when results do not depend on trial duration, it greatly reduces (>50%) the overall duration compared to a single PDR (or NDR) binary search over duration, while finding multiple drop rates.
- In all cases MLRsearch yields the same or similar results to binary search.
- Note: both binary search and MLRsearch are susceptible to reporting non-repeatable results across multiple runs for very bad behaving cases.

Caveats:

- Worst case MLRsearch can take longer than a binary search e.g. in case of drastic changes in behaviour for trials at varying durations.

4.  Sample Implementation

Following is a brief description of a sample MLRsearch implementation, which is a simplified version of the existing implementation.

4.1.  Input Parameters

1.  *maximum_transmit_rate* - Maximum Transmit Rate (MTR) of packets to be used by external traffic generator implementing MLRsearch, limited by the actual Ethernet link(s) rate, NIC model or traffic generator capabilities.

2.  *minimum_transmit_rate* - minimum packet transmit rate to be used for measurements. MLRsearch fails if lower transmit rate needs to be used to meet search criteria.

3.  *final_trial_duration* - required trial duration for final rate measurements.
4. *initial_trial_duration* - trial duration for initial MLR search phase.

5. *final_relative_width* - required measurement resolution expressed as (lower_bound, upper_bound) interval width relative to upper_bound.


7. *number_of_intermediate_phases* - number of phases between the initial phase and the final phase. Impacts the overall MLR search duration. Less phases are required for well behaving cases, more phases may be needed to reduce the overall search duration for worse behaving cases.

4.2. Initial Phase

1. First trial measures at configured maximum transmit rate (MTR) and discovers maximum receive rate (MRR).
   * IN: trial_duration = initial_trial_duration.
   * IN: offered_transmit_rate = maximum_transmit_rate.
   * DO: single trial.
   * OUT: measured loss ratio.
   * OUT: MRR = measured receive rate. If loss ratio is zero, MRR is set below MTR so that interval width is equal to the width goal of the first intermediate phase.

2. Second trial measures at MRR and discovers MRR2.
   * IN: trial_duration = initial_trial_duration.
   * IN: offered_transmit_rate = MRR.
   * DO: single trial.
   * OUT: measured loss ratio.
   * OUT: MRR2 = measured receive rate. If loss ratio is zero, MRR2 is set above MRR so that interval width is equal to the width goal of the first intermediate phase. MRR2 could end up being equal to MTR (for example if both measurements so far
had zero loss), which was already measured, step 3 is skipped in that case.

3. Third trial measures at MRR2.
   * IN: trial_duration = initial_trial_duration.
   * IN: offered_transmit_rate = MRR2.
   * DO: single trial.
   * OUT: measured loss ratio.

4.3. Non-Initial Phases

1. Main loop:
   1. IN: trial_duration for the current phase. Set to initial_trial_duration for the first intermediate phase; to final_trial_duration for the final phase; or to the element of interpolating geometric sequence for other intermediate phases. For example with two intermediate phases, trial_duration of the second intermediate phase is the geometric average of initial_trial_duration and final_trial_duration.
   2. IN: relative_width_goal for the current phase. Set to final_relative_width for the final phase; doubled for each preceding phase. For example with two intermediate phases, the first intermediate phase uses quadruple of final_relative_width and the second intermediate phase uses double of final_relative_width.
   3. IN: ndr_interval, pdr_interval from the previous main loop iteration or the previous phase. If the previous phase is the initial phase, both intervals are formed by a (correctly ordered) pair of MRR2 and MRR. Note that the initial phase is likely to create intervals with invalid bounds.
   4. DO: According to the procedure described in point 2., either exit the phase (by jumping to 1.7.), or calculate new transmit rate to measure with.
   5. DO: Perform the trial measurement at the new transmit rate and trial_duration, compute its loss ratio.
   6. DO: Update the bounds of both intervals, based on the new measurement. The actual update rules are numerous, as NDR
external search can affect PDR interval and vice versa, but the result agrees with rules of both internal and external search. For example, any new measurement below an invalid lower_bound becomes the new lower_bound, while the old measurement (previously acting as the invalid lower_bound) becomes a new and valid upper_bound. Go to next iteration (1.3.), taking the updated intervals as new input.

7. OUT: current ndr_interval and pdr_interval. In the final phase this is also considered to be the result of the whole search. For other phases, the next phase loop is started with the current results as an input.

2. New transmit rate (or exit) calculation (for point 1.4.):

1. If there is an invalid bound then prepare for external search:

   + IF the most recent measurement at NDR lower_bound transmit rate had the loss higher than zero, then the new transmit rate is NDR lower_bound decreased by two NDR interval widths.

   + Else, IF the most recent measurement at PDR lower_bound transmit rate had the loss higher than PLR, then the new transmit rate is PDR lower_bound decreased by two PDR interval widths.

   + Else, IF the most recent measurement at NDR upper_bound transmit rate had no loss, then the new transmit rate is NDR upper_bound increased by two NDR interval widths.

   + Else, IF the most recent measurement at PDR upper_bound transmit rate had the loss lower or equal to PLR, then the new transmit rate is PDR upper_bound increased by two PDR interval widths.

2. Else, if interval width is higher than the current phase goal:

   + IF NDR interval does not meet the current phase width goal, prepare for internal search. The new transmit rate is a in the middle of NDR lower_bound and NDR upper_bound.

   + IF PDR interval does not meet the current phase width goal, prepare for internal search. The new transmit rate is a in the middle of PDR lower_bound and PDR upper_bound.
3. Else, if some bound has still only been measured at a lower
duration, prepare to re-measure at the current duration (and
the same transmit rate). The order of priorities is:

+ NDR lower_bound,
+ PDR lower_bound,
+ NDR upper_bound,
+ PDR upper_bound.

4. Else, do not prepare any new rate, to exit the phase. This
ensures that at the end of each non-initial phase all
intervals are valid, narrow enough, and measured at current
phase trial duration.

5. FD.io CSIT Implementation

The only known working implementation of MLRsearch is in the open-
source code running in Linux Foundation FD.io CSIT project
[FDio-CSIT-MLRsearch] as part of a Continuous Integration /
Continuous Development (CI/CD) framework.

MLRsearch is also available as a Python package in [PyPI-MLRsearch].

5.1. Additional details

This document so far has been describing a simplified version of
MLRsearch algorithm. The full algorithm as implemented in CSIT
contains additional logic, which makes some of the details (but not
general ideas) above incorrect. Here is a short description of the
additional logic as a list of principles, explaining their main
differences from (or additions to) the simplified description, but
without detailing their mutual interaction.

1. Logarithmic transmit rate.

* In order to better fit the relative width goal, the interval
doubling and halving is done differently.
* For example, the middle of 2 and 8 is 4, not 5.

2. Optimistic maximum rate.

* The increased rate is never higher than the maximum rate.
* Upper bound at that rate is always considered valid.
3. Pessimistic minimum rate.
   * The decreased rate is never lower than the minimum rate.
   * If a lower bound at that rate is invalid, a phase stops refining the interval further (until it gets re-measured).

   * Measurements above the current upper bound never update a valid upper bound, even if drop ratio is low.
   * Measurements below the current lower bound always update any lower bound if drop ratio is high.

5. Ensure sufficient interval width.
   * Narrow intervals make external search take more time to find a valid bound.
   * If the new transmit increased or decreased rate would result in width less than the current goal, increase/decrease more.
   * This can happen if the measurement for the other interval makes the current interval too narrow.
   * Similarly, take care the measurements in the initial phase create wide enough interval.

6. Timeout for bad cases.
   * The worst case for MLRsearch is when each phase converges to intervals way different than the results of the previous phase.
   * Rather than suffer total search time several times larger than pure binary search, the implemented tests fail themselves when the search takes too long (given by argument _timeout_).

7. Pessimistic external search.
   * Valid bound becoming invalid on re-measurement with higher duration is frequently a sign of SUT behaving in non-deterministic way (from blackbox point of view). If the final width interval goal is too narrow compared to width of rate region where SUT is non-deterministic, it is quite likely that there will be multiple invalid bounds before the external search finds a valid one.
* In this case, external search can be sped up by increasing interval width more rapidly. As only powers of two ensure the subsequent internal search will not result in needlessly narrow interval, a parameter _doublings_ is introduced to control the pessimism of external search. For example three doublings result in interval width being multiplied by eight in each external search iteration.

5.1.1. FD.io CSIT Input Parameters

1. *maximum_transmit_rate* - Typical values: 2 * 14.88 Mpps for 64B 10GE link rate, 2 * 18.75 Mpps for 64B 40GE NIC (specific model).
2. *minimum_transmit_rate* - Value: 2 * 10 kpps (traffic generator limitation).
3. *final_trial_duration* - Value: 30 seconds.
4. *initial_trial_duration* - Value: 1 second.
5. *final_relative_width* - Value: 0.005 (0.5%).
6. *packet_loss_ratio* - Value: 0.005 (0.5%).
7. *number_of_intermediate_phases* - Value: 2. The value has been chosen based on limited experimentation to date. More experimentation needed to arrive to clearer guidelines.
8. *timeout* - Limit for the overall search duration (for one search). If MLRsearch oversteps this limit, it immediately declares the test failed, to avoid wasting even more time on a misbehaving SUT. Value: 600 (seconds).
9. *doublings* - Number of doublings when computing new interval width in external search. Value: 2 (interval width is quadrupled). Value of 1 is best for well-behaved SUTs, but value of 2 has been found to decrease overall search time for worse-behaved SUT configurations, contributing more to the overall set of different SUT configurations tested.

5.2. Example MLRsearch Run

The following table shows data from a real test run in CSIT (using the default input values as above). The first column is the phase, the second is the trial measurement performed (aggregate bidirectional offered load in megapackets per second, and trial duration in seconds). Each of last four columns show one bound as updated after the measurement (duration truncated to save space).
Loss ratio is not shown, but invalid bounds are marked with a plus sign.

<table>
<thead>
<tr>
<th>Phase</th>
<th>Trial</th>
<th>NDR lower</th>
<th>NDR upper</th>
<th>PDR lower</th>
<th>PDR upper</th>
</tr>
</thead>
<tbody>
<tr>
<td>init.</td>
<td>37.50</td>
<td>N/A</td>
<td>37.50 1.</td>
<td>N/A</td>
<td>37.50 1.</td>
</tr>
<tr>
<td>init.</td>
<td>10.55</td>
<td>+10.55 1.</td>
<td>37.50 1.</td>
<td>+10.55 1.</td>
<td>37.50 1.</td>
</tr>
<tr>
<td>init.</td>
<td>9.437</td>
<td>+9.437 1.</td>
<td>10.55 1.</td>
<td>+9.437 1.</td>
<td>10.55 1.</td>
</tr>
<tr>
<td>int 1</td>
<td>6.053</td>
<td>6.053 1.</td>
<td>9.437 1.</td>
<td>6.053 1.</td>
<td>9.437 1.</td>
</tr>
<tr>
<td>int 1</td>
<td>7.558</td>
<td>7.558 1.</td>
<td>9.437 1.</td>
<td>7.558 1.</td>
<td>9.437 1.</td>
</tr>
<tr>
<td>int 1</td>
<td>8.446</td>
<td>8.446 1.</td>
<td>9.437 1.</td>
<td>8.446 1.</td>
<td>9.437 1.</td>
</tr>
<tr>
<td>int 1</td>
<td>8.928</td>
<td>8.928 1.</td>
<td>9.437 1.</td>
<td>8.928 1.</td>
<td>9.437 1.</td>
</tr>
</tbody>
</table>
6. IANA Considerations

No requests of IANA.

7. Security Considerations

Benchmarking activities as described in this memo are limited to technology characterization of a DUT/SUT using controlled stimuli in a laboratory environment, with dedicated address space and the constraints specified in the sections above.

The benchmarking network topology will be an independent test setup and MUST NOT be connected to devices that may forward the test traffic into a production network or misroute traffic to the test management network.
Further, benchmarking is performed on a "black-box" basis, relying solely on measurements observable external to the DUT/SUT.

Special capabilities SHOULD NOT exist in the DUT/SUT specifically for benchmarking purposes. Any implications for network security arising from the DUT/SUT SHOULD be identical in the lab and in production networks.

8. Acknowledgements

Many thanks to Alec Hothan of OPNFV NFVbench project for thorough review and numerous useful comments and suggestions.

9. References

9.1. Normative References


9.2. Informative References


[PyPI-MLRsearch] "MLRsearch 0.3.0, Python Package Index", February 2020, <https://pypi.org/project/MLRsearch/0.3.0/>.

Authors’ Addresses

Maciek Konstantynowicz (editor)
Cisco Systems

Email: mkonstan@cisco.com
Vratko Polak (editor)
Cisco Systems

Email: vrpolak@cisco.com
This document proposes changes to [RFC2544], specifically to packet throughput search methodology, by defining a new search algorithm referred to as Multiple Loss Ratio search (MLRsearch for short). Instead of relying on binary search with pre-set starting offered load, it proposes a novel approach discovering the starting point in the initial phase, and then searching for packet throughput based on defined packet loss ratio (PLR) input criteria and defined final trial duration time. One of the key design principles behind MLRsearch is minimizing the total test duration and searching for multiple packet throughput rates (each with a corresponding PLR) concurrently, instead of doing it sequentially.

The main motivation behind MLRsearch is the new set of challenges and requirements posed by NFV (Network Function Virtualization), specifically software based implementations of NFV data planes. Using [RFC2544] in the experience of the authors yields often not repetitive and not replicable end results due to a large number of factors that are out of scope for this draft. MLRsearch aims to address this challenge in a simple way of getting the same result sooner, so more repetitions can be done to describe the replicability.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at https://datatracker.ietf.org/drafts/current/.
1. Terminology

TODO: Update after most other sections are updated.

* TODO: The current text uses Throughput for the zero loss ratio load. Is the capital T needed/useful?

* DUT and SUT: see the definitions in https://gerrit.fd.io/r/c/csit/+/35545

* Traffic Generator (TG) and Traffic Analyzer (TA): see https://datatracker.ietf.org/doc/html/rfc6894#section-4 TODO: Maybe there is an earlier RFC?

* Overall search time: the time it takes to find all required loads within their precision goals, starting from zero trials measured at given DUT configuration and traffic profile.

* TODO: traffic profile?

* Intended load: https://datatracker.ietf.org/doc/html/rfc2285#section-3.5.1

* Offered load: https://datatracker.ietf.org/doc/html/rfc2285#section-3.5.2

* Maximum offered load (MOL): see https://datatracker.ietf.org/doc/html/rfc2285#section-3.5.3

* Forwarding rate at maximum offered load (FRMOL) https://datatracker.ietf.org/doc/html/rfc2285#section-3.6.2

* Trial Loss Count: the number of frames transmitted minus the number of frames received. Negative count is possible, e.g. when SUT duplicates some frames.
* Trial Loss Ratio: ratio of frames received relative to frames transmitted over the trial duration. For bi-directional throughput tests, the aggregate ratio is calculated, based on the aggregate number of frames transmitted and received. If the trial loss count is negative, its absolute value MUST be used to keep compliance with RFC2544.

* Safe load: any value, such that trial measurement at this (or lower) intended load is correctly handled by both TG and TA, regardless of SUT behavior. Frequently, it is not known what the safe load is.

* Max load (TODO rename?): Maximal intended load to be used during search. Benchmarking team decides which value is low enough to guarantee values reported by TG and TA are reliable. It has to be a safe load, but it can be lower than a safe load estimate for added safety. See the subsection on unreliable test equipment below. This value MUST NOT be higher than MOL, which itself MUST NOT be higher than Maximum Frame Rate, as per RFC2544 §20

* Min load: Minimal intended load to be used during search. Benchmarking team decides which value is high enough to guarantee the trial measurement results are valid. E.g. considerable overall search time can be saved by declaring SUT faulty if min load trial shows too high loss rate. Zero frames per second is a valid min load value.

* Effective loss ratio: a corrected value of trial loss ratio chosen to avoid difficulties if SUT exhibits decreasing loss ratio with increasing load. It is the maximum of trial loss ratios measured at the same duration on all loads smaller than (and including) the current one.

* Target loss ratio: a loss ratio value acting as an input for the search. The search is finding tight enough lower and upper bounds in intended load, so that the measurement at the lower bound has smaller or equal trial loss ratio, and upper bound has strictly larger trial loss ratio. For the tightest upper bound, the effective loss ratio is the same as trial loss ratio at that upper bound load. For the tightest lower bound, the effective loss ratio can be higher than the trial loss ratio at that lower bound, but still not larger than the target loss ratio.

* TODO: Search algorithm.

* TODO: Precision goal.
* TODO: Define a "benchmarking group".
* TODO: Upper and lower bound.
* TODO: Valid and invalid bound?
* TODO: Interval and interval width?

TODO: Mention NIC/PCI bandwidth/pps limits can be lower than bandwidth of medium.

2. Intentions of this document

The intention of this document is to provide recommendations for: * optimizing search for multiple target loss ratios at once, * speeding up the overall search time, * improve search results repeatability and comparability.

No part of RFC2544 is intended to be obsoleted by this document.

3. RFC2544

3.1. Throughput search

It is useful to restate the key requirements of RFC2544 using the new terminology (see section Terminology).

The following sections of RFC2544 are of interest for this document.

* https://datatracker.ietf.org/doc/html/rfc2544#section-20 Mentions the max load SHOULD not be larger than the theoretical maximum rate for the frame size on the media.

* https://datatracker.ietf.org/doc/html/rfc2544#section-23 Lists the actions to be done for each trial measurement, it also mentions loss rate as an example of trial measurement results. This document uses loss count instead, as that is the quantity that is easier for the current test equipment to measure, e.g. it is not affected by the real traffic duration. TODO: Time uncertainty again.

* https://datatracker.ietf.org/doc/html/rfc2544#section-24 Mentions "full length trials" leading to the Throughput found, as opposed to shorter trial durations, allowed in an attempt to "minimize the length of search procedure". This document talks about "final trial duration" and aims to "optimize overall search time".

No additional constraints are placed on the load selection, and there is no mention of an exit condition, e.g. when there is enough trial measurements to proclaim the largest load with zero trial loss (and final trial duration) to be the Throughput found.

4. Problems

4.1. Repeatability and Comparability

RFC2544 does not suggest to repeat Throughput search, and from just one Throughput value, it cannot be determined how repeatable that value is (how likely it is for a repeated Throughput search to end up with a value less than the precision goal away from the first value).

Depending on SUT behavior, different benchmark groups can report significantly different Throughput values, even when using identical SUT and test equipment, just because of minor differences in their search algorithm (e.g. different max load value).

While repeatability can be addressed by repeating the search several times, the differences in the comparability scenario may be systematic, e.g. seeming like a bias in one or both benchmark groups.

MLRsearch algorithm does not really help with the repeatability problem. This document *RECOMMENDS* to repeat a selection of "important" tests ten times, so users can ascertain the repeatability of the results.

TODO: How to report? Average and standard deviation?

Following MLRsearch algorithm leaves less freedom for the benchmark groups to encounter the comparability problem, although more research is needed to determine the effect of MLRsearch’s tweakable parameters.

4.2. Non-Zero Target Loss Ratios

[https://datatracker.ietf.org/doc/html/rfc1242#section-3.17 defines Throughput as: The maximum rate at which none of the offered frames are dropped by the device.](https://datatracker.ietf.org/doc/html/rfc1242#section-3.17)
and then it says: Since even the loss of one frame in a data stream can cause significant delays while waiting for the higher level protocols to time out, it is useful to know the actual maximum data rate that the device can support.

New "software DUTs" (traffic forwarding programs running on commercial-off-the-shelf compute server hardware) frequently exhibit quite low repeatability of Throughput results per above definition.

This is due to, in general, throughput rates of software DUTs (programs) being sensitive to server resource allocation by OS during runtime, as well as any interrupts or blocking of software threads involved in packet processing.

To deal with this, this document recommends discovery of multiple throughput rates of interest for software DUTs that run on general purpose COTS servers (with x86, AArch64 Instruction Set Architectures): * throughput rate with target of zero packet loss ratio. * at least one throughput rate with target of non-zero packet loss ratio.

In our experience, the higher the target loss ratio is, the better is the repeatability of the corresponding load found.

TODO: Define a good name for a load corresponding to a specific non-zero target loss ration, while keeping Throughput for the load corresponding to zero target loss ratio.

This document RECOMMENDS the benchmark groups to search for corresponding loads to at least one non-zero target loss ratio. This document does not suggest any particular non-zero target loss ratio value to search the corresponding load for.

5. Solution ideas

This document gives several independent ideas on how to lower the (average) overall search time, while remaining unconditionally compliant with RFC2544 (and adding some of extensions).

This document also specifies one particular way to combine all the ideas into a single search algorithm class (single logic with few tweakable parameters).

Little to no research has been done into the question of which combination of ideas achieves the best compromise with respect to overall search time, high repeatability and high comparability.
TODO: How important it is to discuss particular implementation choices, especially when motivated by non-deterministic SUT behavior?

5.1. Short duration trials

https://datatracker.ietf.org/doc/html/rfc2544#section-24 already mentions the possibility of using shorter duration for trials that are not part of "final determination".

Obviously, the upper and lower bound from a smaller duration trial can be used as the initial upper and lower bound for the final determination.

MLRsearch makes it clear a re-measurement is always needed (new trial measurement with the same load but longer duration). It also specifies what to do if the longer trial is no longer a valid bound (TODO define?), e.g. start an external search. Additionally one halving can be saved during the shorter duration search.

5.2. FRMOL as reasonable start

TODO expand: Overall search ends with "final determination" search, preceded by "shorter duration search" preceded by "bound initialization", where the bounds can be considerably different from min and max load.

For SUTs with high repeatability, the FRMOL is usually a good approximation of Throughput. But for less repeatable SUTs, forwarding rate (TODO define) is frequently a bad approximation to Throughput, therefore halving and other robust-to-worst-case approaches have to be used. Still, forwarding rate at FRMOL load can be a good initial bound.

5.3. Non-zero loss ratios

See the "Popularity of non-zero target loss ratios" section above.

TODO: Define "trial measurement result classification criteria", or keep reusing long phrases without definitions?

A search for a load corresponding to a non-zero target loss rate is very similar to a search for Throughput, just the criterion when to increase or decrease the intended load for the next trial measurement uses the comparison of trial loss ratio to the target loss ratio (instead of comparing loss count to zero) Any search algorithm that works for Throughput can be easily used also for non-zero target loss rates, perhaps with small modifications in places where the measured forwarding rate is used.
Note that it is possible to search for multiple loss ratio goals if needed.

5.4. Concurrent ratio search

A single trial measurement result can act as an upper bound for a lower target loss ratio, and as a lower bound for a higher target loss ratio at the same time. This is an example of how it can be advantageous to search for all loss ratio goals "at once", or at least "reuse" trial measurement result done so far.

Even when a search algorithm is fully deterministic in load selection while focusing on a single loss ratio and trial duration, the choice of iteration order between target loss ratios and trial durations can affect the obtained results in subtle ways. MLRsearch offers one particular ordering.

5.5. Load selection heuristics and shortcuts

Aside of the two heuristics already mentioned (FRMOL based initial bounds and saving one halving when increasing trial duration), there are other tricks that can save some overall search time at the cost of keeping the difference between final lower and upper bound intentionally large (but still within the precision goal).


The impact on overall duration is probably small, and the effect on result distribution maybe even smaller. TODO: Is the two-liner above useful at all?

6. Non-compliance with RFC2544

It is possible to achieve even faster search times by abandoning some requirements and suggestions of RFC2544, mainly by reducing the wait times at start and end of trial.

Such results are therefore no longer compliant with RFC2544 (or at least not unconditionally), but they may still be useful for internal usage, or for comparing results of different DUTs achieved with an identical non-compliant algorithm.

TODO: Refer to the subsection with CSIT customizations.
7. Additional Requirements

RFC2544 can be understood as having a number of implicit requirements. They are made explicit in this section (as requirements for this document, not for RFC2544).

Recommendations on how to properly address the implicit requirements are out of scope of this document.

7.1. TODO: Search Stop Criteria

TODO: Mention the timeout parameter?

7.2. Reliability of Test Equipment

Both TG and TA MUST be able to handle correctly every intended load used during the search.

On TG side, the difference between Intended Load and Offered Load MUST be small.

TODO: How small? Difference of one packet may not be measurable due to time uncertainties.

TODO expand: time uncertainty.

To ensure that, max load (see Terminology) has to be set to low enough value. Benchmark groups MAY list the max load value used, especially if the Throughput value is equal (or close) to the max load.

Solutions (even problem formulations) for the following open problems are outside of the scope of this document: * Detecting when the test equipment operates above its safe load. * Finding a large but safe load value. * Correcting any result affected by max load value not being a safe load.

7.2.1. Very late frames

RFC2544 requires quite conservative time delays see https://datatracker.ietf.org/doc/html/rfc2544#section-23 to prevent frames buffered in one trial measurement to be counted as received in a subsequent trial measurement.
However, for some SUTs it may still be possible to buffer enough frames, so they are still sending them (perhaps in bursts) when the next trial measurement starts. Sometimes, this can be detected as a negative trial loss count, e.g. TA receiving more frames than TG has sent during this trial measurement. Frame duplication is another way of causing the negative trial loss count.

https://datatracker.ietf.org/doc/html/rfc2544#section-10 recommends to use sequence numbers in frame payloads, but generating and verifying them requires test equipment resources, which may be not plenty enough to support at high loads. (Using low enough max load would work, but frequently that would be smaller than SUT’s actual Throughput.)

RFC2544 does not offer any solution to the negative loss problem, except implicitly treating negative trial loss counts the same way as positive trial loss counts.

This document also does not offer any practical solution.

Instead, this document SUGGESTS the search algorithm to take any precaution necessary to avoid very late frames.

This document also REQUIRES any detected duplicate frames to be counted as additional lost frames. This document also REQUIRES, any negative trial loss ratio to be treated as positive trial loss ratio of the same absolute value.

!!! Nothing below is up-to-date with draft v02. !!!

8. MLRssearch Background

TODO: Old section, probably obsoleted by preceding section(s).

Multiple Loss Ratio search (MLRssearch) is a packet throughput search algorithm suitable for deterministic systems (as opposed to probabilistic systems). MLRssearch discovers multiple packet throughput rates in a single search, each rate is associated with a distinct Packet Loss Ratio (PLR) criterion.
For cases when multiple rates need to be found, this property makes MLRsearch more efficient in terms of time execution, compared to traditional throughput search algorithms that discover a single packet rate per defined search criteria (e.g. a binary search specified by [RFC2544]). MLRsearch reduces execution time even further by relying on shorter trial durations of intermediate steps, with only the final measurements conducted at the specified final trial duration. This results in the shorter overall search execution time when compared to a traditional binary search, while guaranteeing the same results for deterministic systems.

In practice, two rates with distinct PLRs are commonly used for packet throughput measurements of NFV systems: Non Drop Rate (NDR) with PLR=0 and Partial Drop Rate (PDR) with PLR>0. The rest of this document describes MLRsearch with NDR and PDR pair as an example.

Similarly to other throughput search approaches like binary search, MLRsearch is effective for SUTs/DUTs with PLR curve that is non-decreasing with growing offered load. It may not be as effective for SUTs/DUTs with abnormal PLR curves, although it will always converge to some value.

MLRsearch relies on traffic generator to qualify the received packet stream as error-free, and invalidate the results if any disqualifying errors are present e.g. out-of-sequence frames.

MLRsearch can be applied to both uni-directional and bi-directional throughput tests.

For bi-directional tests, MLRsearch rates and ratios are aggregates of both directions, based on the following assumptions:

* Traffic transmitted by traffic generator and received by SUT/DUT has the same packet rate in each direction, in other words the offered load is symmetric.

* SUT/DUT packet processing capacity is the same in both directions, resulting in the same packet loss under load.

MLRsearch can be applied even without those assumptions, but in that case the aggregate loss ratio is less useful as a metric.

MLRsearch can be used for network transactions consisting of more than just one packet, or anything else that has intended load as input and loss ratio as output (duration as input is optional). This text uses mostly packet-centric language.
9. MLRsearch Overview

The main properties of MLRsearch:

* MLRsearch is a duration aware multi-phase multi-rate search algorithm:
  - Initial Phase determines promising starting interval for the search.
  - Intermediate Phases progress towards defined final search criteria.
  - Final Phase executes measurements according to the final search criteria.
  - Final search criteria are defined by following inputs:
    - Target PLRs (e.g. 0.0 and 0.005 when searching for NDR and PDR).
    - Final trial duration.
    - Measurement resolution.

* Initial Phase:
  - Measure MRR over initial trial duration.
  - Measured MRR is used as an input to the first intermediate phase.

* Multiple Intermediate Phases:
  - Trial duration:
    - Start with initial trial duration in the first intermediate phase.
    - Converge geometrically towards the final trial duration.
  - Track all previous trial measurement results:
    - Duration, offered load and loss ratio are tracked.
    - Effective loss ratios are tracked.
While in practice, real loss ratios can decrease with increasing load, effective loss ratios never decrease. This is achieved by sorting results by load, and using the effective loss ratio of the previous load if the current loss ratio is smaller than that.

- The algorithm queries the results to find best lower and upper bounds.
  + Effective loss ratios are always used.

- The phase ends if all target loss ratios have tight enough bounds.

- Search:
  - Iterate over target loss ratios in increasing order.
  - If both upper and lower bound are in measurement results for this duration, apply bisect until the bounds are tight enough, and continue with next loss ratio.
  - If a bound is missing for this duration, but there exists a bound from the previous duration (compatible with the other bound at this duration), re-measure at the current duration.
  - If a bound in one direction (upper or lower) is missing for this duration, and the previous duration does not have a compatible bound, compute the current "interval size" from the second tightest bound in the other direction (lower or upper respectively) for the current duration, and choose next offered load for external search.

  - The logic guarantees that a measurement is never repeated with both duration and offered load being the same.

  - The logic guarantees that measurements for higher target loss ratio iterations (still within the same phase duration) do not affect validity and tightness of bounds for previous target loss ratio iterations (at the same duration).

- Use of internal and external searches:
  - External search:
    + It is a variant of "exponential search".
+ The "interval size" is multiplied by a configurable constant (powers of two work well with the subsequent internal search).

o Internal search:

+ A variant of binary search that measures at offered load between the previously found bounds.

+ The interval does not need to be split into exact halves, if other split can get to the target width goal faster.

  * The idea is to avoid returning interval narrower than the current width goal. See sample implementation details, below.

* Final Phase:

  - Executed with the final test trial duration, and the final width goal that determines resolution of the overall search.

* Intermediate Phases together with the Final Phase are called Non-Initial Phases.

* The returned bounds stay within prescribed min_rate and max_rate.

  - When returning min_rate or max_rate, the returned bounds may be invalid.

    o E.g. upper bound at max_rate may come from a measurement with loss ratio still not higher than the target loss ratio.

The main benefits of MLRsearch vs. binary search include:

* In general, MLRsearch is likely to execute more trials overall, but likely less trials at a set final trial duration.

* In well behaving cases, e.g. when results do not depend on trial duration, it greatly reduces (>50%) the overall duration compared to a single PDR (or NDR) binary search over duration, while finding multiple drop rates.

* In all cases MLRsearch yields the same or similar results to binary search.

* Note: both binary search and MLRsearch are susceptible to reporting non-repeatable results across multiple runs for very bad behaving cases.
Caveats:

* Worst case MLRsearch can take longer than a binary search, e.g. in case of drastic changes in behaviour for trials at varying durations.

  - Re-measurement at higher duration can trigger a long external search. That never happens in binary search, which uses the final duration from the start.

10. Sample Implementation

Following is a brief description of a sample MLRsearch implementation, which is a simplified version of the existing implementation.

10.1. Input Parameters

1. *max_rate* - Maximum Transmit Rate (MTR) of packets to be used by external traffic generator implementing MLRsearch, limited by the actual Ethernet link(s) rate, NIC model or traffic generator capabilities.

2. *min_rate* - minimum packet transmit rate to be used for measurements. MLRsearch fails if lower transmit rate needs to be used to meet search criteria.

3. *final_trial_duration* - required trial duration for final rate measurements.

4. *initial_trial_duration* - trial duration for initial MLRsearch phase.

5. *final_relative_width* - required measurement resolution expressed as (lower_bound, upper_bound) interval width relative to upper_bound.

6. *packet_loss_ratios* - list of maximum acceptable PLR search criteria.

7. *number_of_intermediate_phases* - number of phases between the initial phase and the final phase. Impacts the overall MLRsearch duration. Less phases are required for well behaving cases, more phases may be needed to reduce the overall search duration for worse behaving cases.
10.2. Initial Phase

1. First trial measures at configured maximum transmit rate (MTR) and discovers maximum receive rate (MRR).
   * IN: trial_duration = initial_trial_duration.
   * IN: offered_transmit_rate = maximum_transmit_rate.
   * DO: single trial.
   * OUT: measured loss ratio.
   * OUT: MRR = measured receive rate. Received rate is computed as intended load multiplied by pass ratio (which is one minus loss ratio). This is useful when loss ratio is computed from a different metric than intended load. For example, intended load can be in transactions (multiple packets each), but loss ratio is computed on level of packets, not transactions.
     * Example: If MTR is 10 transactions per second, and each transaction has 10 packets, and receive rate is 90 packets per second, then loss rate is 10\%, and MRR is computed to be 9 transactions per second.

   If MRR is too close to MTR, MRR is set below MTR so that interval width is equal to the width goal of the first intermediate phase. If MRR is less than min_rate, min_rate is used.

2. Second trial measures at MRR and discovers MRR2.
   * IN: trial_duration = initial_trial_duration.
   * IN: offered_transmit_rate = MRR.
   * DO: single trial.
   * OUT: measured loss ratio.
   * OUT: MRR2 = measured receive rate. If MRR2 is less than min_rate, min_rate is used. If loss ratio is less or equal to the smallest target loss ratio, MRR2 is set to a value above MRR, so that interval width is equal to the width goal of the first intermediate phase. MRR2 could end up being equal to MTR (for example if both measurements so far had zero loss), which was already measured, step 3 is skipped in that case.

3. Third trial measures at MRR2.
* IN: trial_duration = initial_trial_duration.
* IN: offered_transmit_rate = MRR2.
* DO: single trial.
* OUT: measured loss ratio.
* OUT: MRR3 = measured receive rate. If MRR3 is less than min_rate, min_rate is used. If step 3 is not skipped, the first trial measurement is forgotten. This is done because in practice (if MRR2 is above MRR), external search from MRR and MRR2 is likely to lead to a faster intermediate phase than a bisect between MRR2 and MTR.

10.3. Non-Initial Phases

1. Main phase loop:
   1. IN: trial_duration for the current phase. Set to initial_trial_duration for the first intermediate phase; to final_trial_duration for the final phase; or to the element of interpolating geometric sequence for other intermediate phases. For example with two intermediate phases, trial_duration of the second intermediate phase is the geometric average of initial_trial_duration and final_trial_duration.
   2. IN: relative_width_goal for the current phase. Set to final_relative_width for the final phase; doubled for each preceding phase. For example with two intermediate phases, the first intermediate phase uses quadruple of final_relative_width and the second intermediate phase uses double of final_relative_width.
   3. IN: Measurement results from the previous phase (previous duration).
   4. Internal target ratio loop:
      1. IN: Target loss ratio for this iteration of ratio loop.
      2. IN: Measurement results from all previous ratio loop iterations of current phase (current duration).
      3. DO: According to the procedure described in point 2:
         1. either exit the phase (by jumping to 1.5),
2. or exit loop iteration (by continuing with next target loss ratio, jumping to 1.4.1),
3. or calculate new transmit rate to measure with.

4. DO: Perform the trial measurement at the new transmit rate and current trial duration, compute its loss ratio.
5. DO: Add the result and go to next iteration (1.4.1), including the added trial result in 1.4.2.

5. OUT: Measurement results from this phase.
6. OUT: In the final phase, bounds for each target loss ratio are extracted and returned.
   1. If a valid bound does not exist, use min_rate or max_rate.

2. New transmit rate (or exit) calculation (for point 1.4.3):
   1. If the previous duration has the best upper and lower bound, select the middle point as the new transmit rate.
      1. See 2.5.3. below for the exact splitting logic.
      2. This can be a no-op if interval is narrow enough already, in that case continue with 2.2.
   3. Discussion, assuming the middle point is selected and measured:
      1. Regardless of loss rate measured, the result becomes either best upper or best lower bound at current duration.
      2. So this condition is satisfied at most once per iteration.
      3. This also explains why previous phase has double width goal:
         1. We avoid one more bisection at previous phase.
         2. At most one bound (per iteration) is re-measured with current duration.
3. Each re-measurement can trigger an external search.

4. Such surprising external searches are the main hurdle in achieving low overall search durations.

5. Even without 1.1, there is at most one external search per phase and target loss ratio.

6. But without 1.1 there can be two re-measurements, each coming with a risk of triggering external search.

2. If the previous duration has one bound best, select its transmit rate. In deterministic case this is the last measurement needed this iteration.

3. If only upper bound exists in current duration results:
   1. This can only happen for the smallest target loss ratio.
   2. If the upper bound was measured at min_rate, exit the whole phase early (not investigating other target loss ratios).
   3. Select new transmit rate using external search:
      1. For computing previous interval size, use:
         1. second tightest bound at current duration,
         2. or tightest bound of previous duration, if compatible and giving a more narrow interval,
         3. or target interval width if none of the above is available.
         4. In any case increase to target interval width if smaller.
      2. Quadruple the interval width.
      3. Use min_rate if the new transmit rate is lower.
   4. If only lower bound exists in current duration results:
      1. If the lower bound was measured at max_rate, exit this iteration (continue with next lowest target loss ratio).
2. Select new transmit rate using external search:
   1. For computing previous interval size, use:
      1. second tightest bound at current duration,
      2. or tightest bound of previous duration, if compatible and giving a more narrow interval,
      3. or target interval width if none of the above is available.
      4. In any case increase to target interval width if smaller.
   2. Quadruple the interval width.
   3. Use max_rate if the new transmit rate is higher.
   5. The only remaining option is both bounds in current duration results.
      1. This can happen in two ways, depending on how the lower bound was chosen.
         1. It could have been selected for the current loss ratio, e.g. in re-measurement (2.2) or in initial bisect (2.1).
         2. It could have been found as an upper bound for the previous smaller target loss ratio, in which case it might be too low.
         3. The algorithm does not track which one is the case, as the decision logic works well regardless.
      2. Compute "extending down" candidate transmit rate exactly as in 2.3.
      3. Compute "bisecting" candidate transmit rate:
         1. Compute the current interval width from the two bounds.
         2. Express the width as a (float) multiple of the target width goal for this phase.
3. If the multiple is not higher than one, it means the width goal is met. Exit this iteration and continue with next higher target loss ratio.

4. If the multiple is two or less, use half of that for new width if the lower subinterval.

5. Round the multiple up to nearest even integer.

6. Use half of that for new width if the lower subinterval.

7. Example: If lower bound is 2.0 and upper bound is 5.0, and width goal is 1.0, the new candidate transmit rate will be 4.0. This can save a measurement when 4.0 has small loss. Selecting the average (3.5) would never save a measurement, giving more narrow bounds instead.

4. If either candidate computation want to exit the iteration, do as bisecting candidate computation says.

5. The remaining case is both candidates wanting to measure at some rate. Use the higher rate. This prefers external search down narrow enough interval, competing with perfectly sized lower bisect subinterval.

11. FD.io CSIT Implementation

The only known working implementation of MLRsearch is in the open-source code running in Linux Foundation FD.io CSIT project [FDio-CSIT-MLRsearch] as part of a Continuous Integration / Continuous Development (CI/CD) framework.

MLRsearch is also available as a Python package in [PyPI-MLRsearch].

11.1. Additional details

This document so far has been describing a simplified version of MLRsearch algorithm. The full algorithm as implemented in CSIT contains additional logic, which makes some of the details (but not general ideas) above incorrect. Here is a short description of the additional logic as a list of principles, explaining their main differences from (or additions to) the simplified description, but without detailing their mutual interaction.

1. Logarithmic transmit rate.
* In order to better fit the relative width goal, the interval
doubling and halving is done differently.

* For example, the middle of 2 and 8 is 4, not 5.

2. Timeout for bad cases.

* The worst case for MLRsearch is when each phase converges to
intervals way different than the results of the previous
phase.

* Rather than suffer total search time several times larger than
pure binary search, the implemented tests fail themselves when
the search takes too long (given by argument _timeout_).

3. Intended count.

* The number of packets to send during the trial should be equal
to the intended load multiplied by the duration.

  - Also multiplied by a coefficient, if loss ratio is
calculated from a different metric.

    o Example: If a successful transaction uses 10 packets,
    load is given in transactions per second, but loss ratio
    is calculated from packets, so the coefficient to get
    intended count of packets is 10.

* But in practice that does not work.

  - It could result in a fractional number of packets,
  - so it has to be rounded in a way traffic generator chooses,
  - which may depend on the number of traffic flows and traffic
generator worker threads.

4. Attempted count. As the real number of intended packets is not
known exactly, the computation uses the number of packets traffic
generator reports as sent. Unless overridden by the next point.

5. Duration stretching.

* In some cases, traffic generator may get overloaded, causing
it to take significantly longer (than duration) to send all
packets.

* The implementation uses an explicit stop,
- causing lower attempted count in those cases.

* The implementation tolerates some small difference between attempted count and intended count.

- 10 microseconds worth of traffic is sufficient for our tests.

* If the difference is higher, the unsent packets are counted as lost.

- This forces the search to avoid the regions of high duration stretching.

- The final bounds describe the performance of not just SUT, but of the whole system, including the traffic generator.


* In some test (e.g. using TCP flows) Traffic generator reacts to packet loss by retransmission. Usually, such packet loss is already affecting loss ratio. If a test also wants to treat retransmissions due to heavily delayed packets also as a failure, this is once again visible as a mismatch between the intended count and the attempted count.

* The CSIT implementation simply looks at absolute value of the difference, so it offers the same small tolerance before it starts marking a "loss".

7. For result processing, we use lower bounds and ignore upper bounds.

11.1.1. FD.io CSIT Input Parameters

1. *max_rate* - Typical values: 2 * 14.88 Mpps for 64B 10GE link rate, 2 * 18.75 Mpps for 64B 40GE NIC (specific model).

2. *min_rate* - Value: 2 * 9001 pps (we reserve 9000 pps for latency measurements).

3. *final_trial_duration* - Value: 30.0 seconds.

4. *initial_trial_duration* - Value: 1.0 second.

5. *final_relative_width* - Value: 0.005 (0.5%).
6. *packet_loss_ratios* - Value: 0.0, 0.005 (0.0% for NDR, 0.5% for PDR).

7. *number_of_intermediate_phases* - Value: 2. The value has been chosen based on limited experimentation to date. More experimentation needed to arrive to clearer guidelines.

8. *timeout* - Limit for the overall search duration (for one search). If MLRsearch oversteps this limit, it immediately declares the test failed, to avoid wasting even more time on a misbehaving SUT. Value: 600.0 (seconds).

9. *expansion_coefficient* - Width multiplier for external search. Value: 4.0 (interval width is quadroupled). Value of 2.0 is best for well-behaved SUTs, but value of 4.0 has been found to decrease overall search time for worse-behaved SUT configurations, contributing more to the overall set of different SUT configurations tested.

11.2. Example MLRsearch Run

The following list describes a search from a real test run in CSIT (using the default input values as above).

* Initial phase, trial duration 1.0 second.

Measurement 1, intended load 18750000.0 pps (MTR), measured loss ratio 0.7089514628479618 (valid upper bound for both NDR and PDR).

Measurement 2, intended load 5457160.071600716 pps (MRR), measured loss ratio 0.018650817320118702 (new tightest upper bounds).

Measurement 3, intended load 5348832.933500009 pps (slightly less than MRR2 in preparation for first intermediate phase target interval width), measured loss ratio 0.00964383362905351 (new tightest upper bounds).

* First intermediate phase starts, trial duration still 1.0 seconds.

Measurement 4, intended load 4936605.579021453 pps (no lower bound, performing external search downwards, for NDR), measured loss ratio 0.0 (valid lower bound for both NDR and PDR).

Measurement 5, intended load 5138587.208637197 pps (bisecting for NDR), measured loss ratio 0.0 (new tightest lower bounds).

Measurement 6, intended load 5242656.244044665 pps (bisecting), measured loss ratio 0.013523745379347257 (new tightest upper bounds).
* Both intervals are narrow enough.

* Second intermediate phase starts, trial duration 5.477225575051661 seconds.

Measurement 7, intended load 5190360.904111567 pps (initial bisect for NDR), measured loss ratio 0.0023533920869969953 (NDR upper bound, PDR lower bound).

Measurement 8, intended load 5138587.208637197 pps (re-measuring NDR lower bound), measured loss ratio 1.208022912800403e-06 (new tightest NDR upper bound).

* The two intervals have separate bounds from now on.

Measurement 9, intended load 4936605.381062318 pps (external NDR search down), measured loss ratio 0.0 (new valid NDR lower bound).

Measurement 10, intended load 5036583.888432355 pps (NDR bisect), measured loss ratio 0.0 (new tightest NDR lower bound).

Measurement 11, intended load 5087329.903232804 pps (NDR bisect), measured loss ratio 0.0 (new tightest NDR lower bound).

* NDR interval is narrow enough, PDR interval not ready yet.

Measurement 12, intended load 5242656.244044665 pps (re-measuring PDR upper bound), measured loss ratio 0.0101174866190136 (still valid PDR upper bound).

* Also PDR interval is narrow enough, with valid bounds for this duration.

* Final phase starts, trial duration 30.0 seconds.

Measurement 13, intended load 5112894.3238511775 pps (initial bisect for NDR), measured loss ratio 0.0 (new tightest NDR lower bound).

Measurement 14, intended load 5138587.208637197 (re-measuring NDR upper bound), measured loss ratio 2.030389804256833e-06 (still valid PDR upper bound).

* NDR interval is narrow enough, PDR interval not yet.

Measurement 15, intended load 5216443.04126728 pps (initial bisect for PDR), measured loss ratio 0.005620871287975237 (new tightest PDR upper bound).
Measurement 16, intended load 5190360.904111567 (re-measuring PDR lower bound), measured loss ratio 0.0027629971184465604 (still valid PDR lower bound).

* PDR interval is also narrow enough.

* Returning bounds:

* NDR_LOWER = 5112894.3238511775 pps; NDR_UPPER = 5138587.208637197 pps;

* PDR_LOWER = 5190360.904111567 pps; PDR_UPPER = 5216443.04126728 pps.

12. IANA Considerations

No requests of IANA.

13. Security Considerations

Benchmarking activities as described in this memo are limited to technology characterization of a DUT/SUT using controlled stimuli in a laboratory environment, with dedicated address space and the constraints specified in the sections above.

The benchmarking network topology will be an independent test setup and MUST NOT be connected to devices that may forward the test traffic into a production network or misroute traffic to the test management network.

Further, benchmarking is performed on a "black-box" basis, relying solely on measurements observable external to the DUT/SUT.

Special capabilities SHOULD NOT exist in the DUT/SUT specifically for benchmarking purposes. Any implications for network security arising from the DUT/SUT SHOULD be identical in the lab and in production networks.

14. Acknowledgements

Many thanks to Alec Hothan of OPNFV NFVbench project for thorough review and numerous useful comments and suggestions.

15. References

15.1. Normative References
15.2. Informative References

[FDio-CSIT-MLRsearch]  "FD.io CSIT Test Methodology - MLRsearch", November 2021,

[PyPI-MLRsearch]  "MLRsearch 0.4.0, Python Package Index", April 2021,
<https://pypi.org/project/MLRsearch/0.4.0/>.

Authors’ Addresses

Maciek Konstantynowicz (editor)
Cisco Systems
Email: mkonstan@cisco.com

Vratko Polak
Cisco Systems
Email: vrpolak@cisco.com
Abstract

This document provides benchmarking terminology and methodology for next-generation network security devices including next-generation firewalls (NGFW), next-generation intrusion prevention systems (NGIPS), and unified threat management (UTM) implementations. The main areas covered in this document are test terminology, test configuration parameters, and benchmarking methodology for NGFW and NGIPS. This document aims to improve the applicability, reproducibility, and transparency of benchmarks and to align the test methodology with today’s increasingly complex layer 7 security centric network application use cases. As a result, this document makes [RFC3511] obsolete.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at https://datatracker.ietf.org/drafts/current/.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on 5 July 2022.

Copyright Notice

Copyright (c) 2022 IETF Trust and the persons identified as the document authors. All rights reserved.
This document is subject to BCP 78 and the IETF Trust’s Legal Provisions Relating to IETF Documents (https://trustee.ietf.org/license-info) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Revised BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Revised BSD License.

Table of Contents

1. Introduction ........................................... 4
2. Requirements .......................................... 4
3. Scope .................................................. 4
4. Test Setup ............................................ 4
4.1. Testbed Configuration ............................... 5
4.2. DUT/SUT Configuration ............................... 6
4.2.1. Security Effectiveness Configuration .......... 12
4.3. Test Equipment Configuration ....................... 12
4.3.1. Client Configuration ............................. 12
4.3.2. Backend Server Configuration ................... 15
4.3.3. Traffic Flow Definition ......................... 17
4.3.4. Traffic Load Profile ............................ 17
5. Testbed Considerations ............................... 18
6. Reporting ............................................. 19
6.1. Introduction ....................................... 19
6.2. Detailed Test Results ............................. 21
6.3. Benchmarks and Key Performance Indicators ....... 21
7. Benchmarking Tests .................................... 23
7.1. Throughput Performance with Application Traffic Mix 23
7.1.1. Objective ......................................... 23
7.1.2. Test Setup ....................................... 23
7.1.3. Test Parameters ................................... 23
7.1.4. Test Procedures and Expected Results .......... 25
7.2. TCP/HTTP Connections Per Second .................. 26
7.2.1. Objective ......................................... 26
7.2.2. Test Setup ....................................... 27
7.2.3. Test Parameters ................................... 27
7.2.4. Test Procedures and Expected Results .......... 28
7.3. HTTP Throughput .................................... 30
7.3.1. Objective ......................................... 30
7.3.2. Test Setup ....................................... 30
7.3.3. Test Parameters ................................... 30
7.3.4. Test Procedures and Expected Results .......... 32
7.4. HTTP Transaction Latency .......................... 33
7.4.1. Objective ......................................... 33
7.4.2. Test Setup ....................................... 33
7.4.3. Test Parameters ................................... 34
7.4.4. Test Procedures and Expected Results .................................. 35
7.5. Concurrent TCP/HTTP Connection Capacity ................................. 36
  7.5.1. Objective .......................................................... 36
  7.5.2. Test Setup ......................................................... 36
  7.5.3. Test Parameters ................................................... 37
  7.5.4. Test Procedures and Expected Results ................................ 38
7.6. TCP/HTTPS Connections per Second ........................................ 39
  7.6.1. Objective .......................................................... 40
  7.6.2. Test Setup ......................................................... 40
  7.6.3. Test Parameters ................................................... 40
  7.6.4. Test Procedures and Expected Results ................................ 42
7.7. HTTPS Throughput .......................................................... 43
  7.7.1. Objective .......................................................... 43
  7.7.2. Test Setup ......................................................... 43
  7.7.3. Test Parameters ................................................... 43
  7.7.4. Test Procedures and Expected Results ................................ 45
7.8. HTTPS Transaction Latency ................................................... 46
  7.8.1. Objective .......................................................... 46
  7.8.2. Test Setup ......................................................... 46
  7.8.3. Test Parameters ................................................... 46
  7.8.4. Test Procedures and Expected Results ................................ 48
7.9. Concurrent TCP/HTTPS Connection Capacity ................................. 49
  7.9.1. Objective .......................................................... 49
  7.9.2. Test Setup ......................................................... 49
  7.9.3. Test Parameters ................................................... 49
  7.9.4. Test Procedures and Expected Results ................................ 51
8. IANA Considerations .............................................................. 52
9. Security Considerations .......................................................... 53
10. Contributors ........................................................................... 53
11. Acknowledgements ..................................................................... 53
12. References ............................................................................. 53
  12.1. Normative References .................................................... 53
  12.2. Informative References ................................................... 53
Appendix A. Test Methodology — Security Effectiveness
  Evaluation ............................................................................. 54
  A.1. Test Objective ................................................................ 55
  A.2. Testbed Setup ............................................................ 55
  A.3. Test Parameters ............................................................ 55
    A.3.1. DUT/SUT Configuration Parameters .............................. 55
    A.3.2. Test Equipment Configuration Parameters .................... 55
  A.4. Test Results Validation Criteria ....................................... 56
  A.5. Measurement .................................................................. 56
  A.6. Test Procedures and Expected Results ................................... 57
    A.6.1. Step 1: Background Traffic ....................................... 57
    A.6.2. Step 2: CVE Emulation ............................................. 58
Appendix B. DUT/SUT Classification ................................................. 58
Authors’ Addresses ....................................................................... 58
1. Introduction

18 years have passed since IETF recommended test methodology and terminology for firewalls initially ([RFC3511]). The requirements for network security element performance and effectiveness have increased tremendously since then. In the eighteen years since [RFC3511] was published, recommending test methodology and terminology for firewalls, requirements and expectations for network security elements has increased tremendously. Security function implementations have evolved to more advanced areas and have diversified into intrusion detection and prevention, threat management, analysis of encrypted traffic, etc. In an industry of growing importance, well-defined, and reproducible key performance indicators (KPIs) are increasingly needed to enable fair and reasonable comparison of network security functions. All these reasons have led to the creation of a new next-generation network security device benchmarking document, which makes [RFC3511] obsolete.

2. Requirements

The keywords "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119], [RFC8174] when, and only when, they appear in all capitals, as shown here.

3. Scope

This document provides testing terminology and testing methodology for modern and next-generation network security devices that are configured in Active ("Inline", see Figure 1 and Figure 2) mode. It covers the validation of security effectiveness configurations of network security devices, followed by performance benchmark testing. This document focuses on advanced, realistic, and reproducible testing methods. Additionally, it describes testbed environments, test tool requirements, and test result formats.

4. Test Setup

Test setup defined in this document applies to all benchmarking tests described in Section 7. The test setup MUST be contained within an Isolated Test Environment (see Section 3 of [RFC6815]).
4.1. Testbed Configuration

Testbed configuration MUST ensure that any performance implications that are discovered during the benchmark testing aren’t due to the inherent physical network limitations such as the number of physical links and forwarding performance capabilities (throughput and latency) of the network devices in the testbed. For this reason, this document recommends avoiding external devices such as switches and routers in the testbed wherever possible.

In some deployment scenarios, the network security devices (Device Under Test/System Under Test) are connected to routers and switches, which will reduce the number of entries in MAC or ARP tables of the Device Under Test/System Under Test (DUT/SUT). If MAC or ARP tables have many entries, this may impact the actual DUT/SUT performance due to MAC and ARP/ND (Neighbor Discovery) table lookup processes. This document also recommends using test equipment with the capability of emulating layer 3 routing functionality instead of adding external routers in the testbed.

The testbed setup Option 1 (Figure 1) is the RECOMMENDED testbed setup for the benchmarking test.

```plaintext
+-----------------------+                   +-----------------------+
| +-------------------+ |   +-----------+   | +-------------------+ |
| | Emulated Router(s)| |   |           |   | | Emulated Router(s)| |
| |    (Optional)     | +----- DUT/SUT +-----+    (Optional)     | |
| +-------------------+ |   |           |   | +-------------------+ |
| +-------------------+ |   +-----------+   | +-------------------+ |
| |     Clients       | |                   | |      Servers      | |
| +-------------------+ |                   | +-------------------+ |
|                       |                   |                       |
|   Test Equipment      |                   |   Test Equipment      |
+-----------------------+                   +-----------------------+
```

**Figure 1: Testbed Setup - Option 1**

If the test equipment used is not capable of emulating layer 3 routing functionality or if the number of used ports is mismatched between test equipment and the DUT/SUT (need for test equipment port aggregation), the test setup can be configured as shown in Figure 2.
4.2. DUT/SUT Configuration

A unique DUT/SUT configuration MUST be used for all benchmarking tests described in Section 7. Since each DUT/SUT will have its own unique configuration, users SHOULD configure their device with the same parameters and security features that would be used in the actual deployment of the device or a typical deployment in order to achieve maximum network security coverage. The DUT/SUT MUST be configured in "Inline" mode so that the traffic is actively inspected by the DUT/SUT. Also "Fail-Open" behavior MUST be disabled on the DUT/SUT.

Table 1 and Table 2 below describe the RECOMMENDED and OPTIONAL sets of network security feature list for NGFW and NGIPS respectively. The selected security features SHOULD be consistently enabled on the DUT/SUT for all benchmarking tests described in Section 7.

To improve repeatability, a summary of the DUT/SUT configuration including a description of all enabled DUT/SUT features MUST be published with the benchmarking results.
<table>
<thead>
<tr>
<th>DUT/SUT (NGFW) Features</th>
<th>RECOMMENDED</th>
<th>OPTIONAL</th>
</tr>
</thead>
<tbody>
<tr>
<td>SSL Inspection</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>IDS/IPS</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>Anti-Spyware</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>Anti-Virus</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>Anti-Botnet</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>Web Filtering</td>
<td></td>
<td>x</td>
</tr>
<tr>
<td>Data Loss Protection (DLP)</td>
<td></td>
<td>x</td>
</tr>
<tr>
<td>DDoS</td>
<td></td>
<td>x</td>
</tr>
<tr>
<td>Certificate Validation</td>
<td></td>
<td>x</td>
</tr>
<tr>
<td>Logging and Reporting</td>
<td></td>
<td>x</td>
</tr>
<tr>
<td>Application Identification</td>
<td></td>
<td>x</td>
</tr>
</tbody>
</table>

Table 1: NGFW Security Features
### DUT/SUT (NGIPS) Features

<table>
<thead>
<tr>
<th>Feature</th>
<th>RECOMMENDED</th>
<th>OPTIONAL</th>
</tr>
</thead>
<tbody>
<tr>
<td>SSL Inspection</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>Anti-Malware</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>Anti-Spyware</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>Anti-Botnet</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>Logging and Reporting</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>Application Identification</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>Deep Packet Inspection</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>Anti-Evasion</td>
<td>x</td>
<td></td>
</tr>
</tbody>
</table>

**Table 2: NGIPS Security Features**

The following table provides a brief description of the security features.

<table>
<thead>
<tr>
<th>Feature</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>SSL Inspection</td>
<td>DUT/SUT intercepts and decrypts inbound HTTPS traffic between servers and clients. Once the content inspection has been completed, DUT/SUT encrypts the HTTPS traffic with ciphers and keys used by the clients and servers.</td>
</tr>
<tr>
<td>IDS/IPS</td>
<td>DUT/SUT detects and blocks exploits targeting known and unknown vulnerabilities across the monitored network.</td>
</tr>
<tr>
<td>Anti-Malware</td>
<td>DUT/SUT detects and prevents the transmission of malicious executable code and any associated communications across the monitored network. This includes data exfiltration as well as command and control channels.</td>
</tr>
<tr>
<td>Anti-Spyware</td>
<td>Anti-Spyware is a subcategory of Anti Malware. Spyware transmits information without the user’s knowledge or permission. DUT/SUT</td>
</tr>
</tbody>
</table>
detects and block initial infection or transmission of data.

**Anti-Botnet**
DUT/SUT detects traffic to or from botnets.

**Anti-Evasion**
DUT/SUT detects and mitigates attacks that have been obfuscated in some manner.

**Web Filtering**
DUT/SUT detects and blocks malicious website including defined classifications of website across the monitored network.

**DLP**
DUT/SUT detects and prevents data breaches and data exfiltration, or it detects and blocks the transmission of sensitive data across the monitored network.

**Certificate Validation**
DUT/SUT validates certificates used in encrypted communications across the monitored network.

**Logging and Reporting**
DUT/SUT logs and reports all traffic at the flow level across the monitored network.

**Application Identification**
DUT/SUT detects known applications as defined within the traffic mix selected across the monitored network.

| **Table 3: Security Feature Description** |

Below is a summary of the DUT/SUT configuration:

* DUT/SUT MUST be configured in "inline" mode.

* "Fail-Open" behavior MUST be disabled.

* All RECOMMENDED security features are enabled.

* Logging SHOULD be enabled. DUT/SUT SHOULD log all traffic at the flow level - Logging to an external device is permissible.

* Geographical location filtering, and Application Identification and Control SHOULD be configured to trigger based on a site or application from the defined traffic mix.
In addition, a realistic number of access control rules (ACL) SHOULD be configured on the DUT/SUT where ACLs are configurable and reasonable based on the deployment scenario. This document determines the number of access policy rules for four different classes of DUT/SUT: Extra Small (XS), Small (S), Medium (M), and Large (L). A sample DUT/SUT classification is described in Appendix B.

The Access Control Rules (ACL) defined in Figure 3 MUST be configured from top to bottom in the correct order as shown in the table. This is due to ACL types listed in specificity decreasing order, with "block" first, followed by "allow", representing a typical ACL based security policy. The ACL entries SHOULD be configured with routable IP subnets by the DUT/SUT. (Note: There will be differences between how security vendors implement ACL decision making.) The configured ACL MUST NOT block the security and measurement traffic used for the benchmarking tests.
<table>
<thead>
<tr>
<th>Rules Type</th>
<th>Match Criteria</th>
<th>Description</th>
<th>Action</th>
<th>XS</th>
<th>S</th>
<th>M</th>
<th>L</th>
</tr>
</thead>
<tbody>
<tr>
<td>Application</td>
<td>Application</td>
<td>Any application not included in the measurement traffic</td>
<td>block</td>
<td>5</td>
<td>10</td>
<td>20</td>
<td>50</td>
</tr>
<tr>
<td>Transport</td>
<td>SRC IP and TCP/UDP DST ports</td>
<td>Any SRC IP subnet used and any DST ports not used in the measurement traffic</td>
<td>block</td>
<td>25</td>
<td>50</td>
<td>100</td>
<td>250</td>
</tr>
<tr>
<td>IP</td>
<td>SRC/DST IP</td>
<td>Any SRC/DST IP subnet not used in the measurement traffic</td>
<td>block</td>
<td>25</td>
<td>50</td>
<td>100</td>
<td>250</td>
</tr>
<tr>
<td>Application</td>
<td>Application</td>
<td>Half of the applications included in the measurement traffic (see the note below)</td>
<td>allow</td>
<td>10</td>
<td>10</td>
<td>10</td>
<td>10</td>
</tr>
<tr>
<td>Transport</td>
<td>SRC IP and TCP/UDP DST ports</td>
<td>Half of the SRC IPs used and any DST ports used in the measurement traffic (one rule per subnet)</td>
<td>allow</td>
<td>&gt;1</td>
<td>&gt;1</td>
<td>&gt;1</td>
<td>&gt;1</td>
</tr>
<tr>
<td>IP</td>
<td>SRC IP</td>
<td>The rest of the SRC IP subnet range used in the measurement traffic (one rule per subnet)</td>
<td>allow</td>
<td>&gt;1</td>
<td>&gt;1</td>
<td>&gt;1</td>
<td>&gt;1</td>
</tr>
</tbody>
</table>

Figure 3: DUT/SUT Access List
Note: If half of the applications included in the measurement traffic is less than 10, the missing number of ACL entries (dummy rules) can be configured for any application traffic not included in the measurement traffic.

4.2.1. Security Effectiveness Configuration

The Security features (defined in Table 1 and Table 2) of the DUT/SUT MUST be configured effectively to detect, prevent, and report the defined security vulnerability sets. This section defines the selection of the security vulnerability sets from Common vulnerabilities and Exposures (CVE) list for the testing. The vulnerability set SHOULD reflect a minimum of 500 CVEs from no older than 10 calendar years to the current year. These CVEs SHOULD be selected with a focus on in-use software commonly found in business applications, with a Common vulnerability Scoring System (CVSS) Severity of High (7-10).

This document is primarily focused on performance benchmarking. However, it is RECOMMENDED to validate the security features configuration of the DUT/SUT by evaluating the security effectiveness as a prerequisite for performance benchmarking tests defined in the section 7. In case the benchmarking tests are performed without evaluating security effectiveness, the test report MUST explain the implications of this. The methodology for evaluating security effectiveness is defined in Appendix A.

4.3. Test Equipment Configuration

In general, test equipment allows configuring parameters in different protocol layers. These parameters thereby influence the traffic flows which will be offered and impact performance measurements.

This section specifies common test equipment configuration parameters applicable for all benchmarking tests defined in Section 7. Any benchmarking test specific parameters are described under the test setup section of each benchmarking test individually.

4.3.1. Client Configuration

This section specifies which parameters SHOULD be considered while configuring clients using test equipment. Also, this section specifies the RECOMMENDED values for certain parameters. The values are the defaults used in most of the client operating systems currently.
4.3.1.1. TCP Stack Attributes

The TCP stack SHOULD use a congestion control algorithm at client and server endpoints. The IPv4 and IPv6 Maximum Segment Size (MSS) SHOULD be set to 1460 bytes and 1440 bytes respectively and a TX and RX initial receive windows of 64 KByte. Client initial congestion window SHOULD NOT exceed 10 times the MSS. Delayed ACKs are permitted and the maximum client delayed ACK SHOULD NOT exceed 10 times the MSS before a forced ACK. Up to three retries SHOULD be allowed before a timeout event is declared. All traffic MUST set the TCP PSH flag to high. The source port range SHOULD be in the range of 1024 - 65535. Internal timeout SHOULD be dynamically scalable per RFC 793. The client SHOULD initiate and close TCP connections. The TCP connection MUST be initiated via a TCP three-way handshake (SYN, SYN/ACK, ACK), and it MUST be closed via either a TCP three-way close (FIN, FIN/ACK, ACK), or a TCP four-way close (FIN, ACK, FIN, ACK).

4.3.1.2. Client IP Address Space

The sum of the client IP space SHOULD contain the following attributes.

* The IP blocks SHOULD consist of multiple unique, discontinuous static address blocks.

* A default gateway is permitted.

* The DSCP (differentiated services code point) marking is set to DF (Default Forwarding) ‘000000’ on IPv4 Type of Service (ToS) field and IPv6 traffic class field.

The following equation can be used to define the total number of client IP addresses that will be configured on the test equipment.

Desired total number of client IP = Target throughput [Mbit/s] / Average throughput per IP address [Mbit/s]

As shown in the example list below, the value for "Average throughput per IP address" can be varied depending on the deployment and use case scenario.

(Option 1) DUT/SUT deployment scenario 1 : 6-7 Mbit/s per IP (e.g. 1,400-1,700 IPs per 10Gbit/s throughput)

(Option 2) DUT/SUT deployment scenario 2 : 0.1-0.2 Mbit/s per IP (e.g. 50,000-100,000 IPs per 10Gbit/s throughput)
Based on deployment and use case scenario, client IP addresses SHOULD be distributed between IPv4 and IPv6. The following options MAY be considered for a selection of traffic mix ratio.

(Option 1) 100 % IPv4, no IPv6
(Option 2) 80 % IPv4, 20% IPv6
(Option 3) 50 % IPv4, 50% IPv6
(Option 4) 20 % IPv4, 80% IPv6
(Option 5) no IPv4, 100% IPv6

Note: The IANA has assigned IP address range for the testing purpose as described in Section 8. If the test scenario requires more IP addresses or subnets than the IANA assigned, this document recommends using non routable Private IPv4 address ranges or Unique Local Address (ULA) IPv6 address ranges for the testing.

4.3.1.3. Emulated Web Browser Attributes

The client emulated web browser (emulated browser) contains attributes that will materially affect how traffic is loaded. The objective is to emulate modern, typical browser attributes to improve realism of the result set.

For HTTP traffic emulation, the emulated browser MUST negotiate HTTP version 1.1 or higher. Depending on test scenarios and chosen HTTP version, the emulated browser MAY open multiple TCP connections per Server endpoint IP at any time depending on how many sequential transactions need to be processed. For HTTP/2 or HTTP/3, the emulated browser MAY open multiple concurrent streams per connection (multiplexing). HTTP/3 emulated browser uses QUIC ([RFC9000]) as transport protocol. HTTP settings such as number of connection per server IP, number of requests per connection, and number of streams per connection MUST be documented. This document refers to [RFC8446] for HTTP/2. The emulated browser SHOULD advertise a User-Agent header. The emulated browser SHOULD enforce content length validation. Depending on test scenarios and selected HTTP version, HTTP header compression MAY be set to enable or disable. This setting (compression enabled or disabled) MUST be documented in the report.

For encrypted traffic, the following attributes SHALL define the negotiated encryption parameters. The test clients MUST use TLS version 1.2 or higher. TLS record size MAY be optimized for the HTTPS response object size up to a record size of 16 KByte. If
Server Name Indication (SNI) is required in the traffic mix profile, the client endpoint MUST send TLS extension Server Name Indication (SNI) information when opening a security tunnel. Each client connection MUST perform a full handshake with server certificate and MUST NOT use session reuse or resumption.

The following TLS 1.2 supported ciphers and keys are RECOMMENDED to use for HTTPS based benchmarking tests defined in Section 7.

1. ECDHE-ECDSA-AES128-GCM-SHA256 with Prime256v1 (Signature Hash Algorithm: ecdsa_secp256r1_sha256 and Supported group: secp256r1)
2. ECDHE-RSA-AES128-GCM-SHA256 with RSA 2048 (Signature Hash Algorithm: rsa_pkcs1_sha256 and Supported group: secp256r1)
3. ECDHE-ECDSA-AES256-GCM-SHA384 with Secp521 (Signature Hash Algorithm: ecdsa_secp384r1_sha384 and Supported group: secp521r1)
4. ECDHE-RSA-AES256-GCM-SHA384 with RSA 4096 (Signature Hash Algorithm: rsa_pkcs1_sha384 and Supported group: secp256r1)

Note: The above ciphers and keys were those commonly used enterprise grade encryption cipher suites for TLS 1.2. It is recognized that these will evolve over time. Individual certification bodies SHOULD use ciphers and keys that reflect evolving use cases. These choices MUST be documented in the resulting test reports with detailed information on the ciphers and keys used along with reasons for the choices.

[RFC8446] defines the following cipher suites for use with TLS 1.3.

1. TLS_AES_128_GCM_SHA256
2. TLS_AES_256_GCM_SHA384
3. TLS_CHACHA20_POLY1305_SHA256
4. TLS_AES_128_CCM_SHA256
5. TLS_AES_128_CCM_8_SHA256

4.3.2. Backend Server Configuration

This section specifies which parameters should be considered while configuring emulated backend servers using test equipment.
4.3.2.1. TCP Stack Attributes

The TCP stack on the server side SHOULD be configured similar to the client side configuration described in Section 4.3.1.1. In addition, server initial congestion window MUST NOT exceed 10 times the MSS. Delayed ACKs are permitted and the maximum server delayed ACK MUST NOT exceed 10 times the MSS before a forced ACK.

4.3.2.2. Server Endpoint IP Addressing

The sum of the server IP space SHOULD contain the following attributes.

* The server IP blocks SHOULD consist of unique, discontinuous static address blocks with one IP per server Fully Qualified Domain Name (FQDN) endpoint per test port.

* A default gateway is permitted. The DSCP (differentiated services code point) marking is set to DF (Default Forwarding) ‘000000’ on IPv4 Type of Service (ToS) field and IPv6 traffic class field.

* The server IP addresses SHOULD be distributed between IPv4 and IPv6 with a ratio identical to the clients distribution ratio.

Note: The IANA has assigned IP address range for the testing purpose as described in Section 8. If the test scenario requires more IP addresses or subnets than the IANA assigned, this document recommends using non routable Private IPv4 address ranges or Unique Local Address (ULA) IPv6 address ranges for the testing.

4.3.2.3. HTTP / HTTPS Server Pool Endpoint Attributes

The server pool for HTTP SHOULD listen on TCP port 80 and emulate the same HTTP version (HTTP 1.1 or HTTP/2 or HTTP/3) and settings chosen by the client (emulated web browser). The Server MUST advertise server type in the Server response header [RFC7230]. For HTTPS server, TLS 1.2 or higher MUST be used with a maximum record size of 16 KByte and MUST NOT use ticket resumption or session ID reuse. The server SHOULD listen on TCP port 443 for HTTP version 1.1 and 2. For HTTP/3 (HTTP over QUIC) the server SHOULD listen on UDP 443. The server SHALL serve a certificate to the client. The HTTPS server MUST check host SNI information with the FQDN if SNI is in use. Cipher suite and key size on the server side MUST be configured similar to the client side configuration described in Section 4.3.1.3.
4.3.3. Traffic Flow Definition

This section describes the traffic pattern between client and server endpoints. At the beginning of the test, the server endpoint initializes and will be ready to accept connection states including initialization of the TCP stack as well as bound HTTP and HTTPS servers. When a client endpoint is needed, it will initialize and be given attributes such as a MAC and IP address. The behavior of the client is to sweep through the given server IP space, generating a recognizable service by the DUT. Sequential and pseudorandom sweep methods are acceptable. The method used MUST be stated in the final report. Thus, a balanced mesh between client endpoints and server endpoints will be generated in a client IP and port to server IP and port combination. Each client endpoint performs the same actions as other endpoints, with the difference being the source IP of the client endpoint and the target server IP pool. The client MUST use the server IP address or FQDN in the host header [RFC7230].

4.3.3.1. Description of Intra-Client Behavior

Client endpoints are independent of other clients that are concurrently executing. When a client endpoint initiates traffic, this section describes how the client steps through different services. Once the test is initialized, the client endpoints randomly hold (perform no operation) for a few milliseconds for better randomization of the start of client traffic. Each client will either open a new TCP connection or connect to a TCP persistence stack still open to that specific server. At any point that the traffic profile may require encryption, a TLS encryption tunnel will form presenting the URL or IP address request to the server. If using SNI, the server MUST then perform an SNI name check with the proposed FQDN compared to the domain embedded in the certificate. Only when correct, will the server process the HTTPS response object. The initial response object to the server is based on benchmarking tests described in Section 7. Multiple additional sub-URLs (response objects on the service page) MAY be requested simultaneously. This MAY be to the same server IP as the initial URL. Each sub-object will also use a canonical FQDN and URL path, as observed in the traffic mix used.

4.3.4. Traffic Load Profile

The loading of traffic is described in this section. The loading of a traffic load profile has five phases: Init, ramp up, sustain, ramp down, and collection.
1. **Init phase**: Testbed devices including the client and server endpoints should negotiate layer 2-3 connectivity such as MAC learning and ARP. Only after successful MAC learning or ARP/ND resolution SHALL the test iteration move to the next phase. No measurements are made in this phase. The minimum RECOMMENDED time for Init phase is 5 seconds. During this phase, the emulated clients SHOULD NOT initiate any sessions with the DUT/SUT, in contrast, the emulated servers should be ready to accept requests from DUT/SUT or from emulated clients.

2. **Ramp up phase**: The test equipment SHOULD start to generate the test traffic. It SHOULD use a set of the approximate number of unique client IP addresses to generate traffic. The traffic SHOULD ramp up from zero to desired target objective. The target objective is defined for each benchmarking test. The duration for the ramp up phase MUST be configured long enough that the test equipment does not overwhelm the DUT/SUTs stated performance metrics defined in Section 6.3 namely, TCP Connections Per Second, Inspected Throughput, Concurrent TCP Connections, and Application Transactions Per Second. No measurements are made in this phase.

3. **Sustain phase**: Starts when all required clients are active and operating at their desired load condition. In the sustain phase, the test equipment SHOULD continue generating traffic to constant target value for a constant number of active clients. The minimum RECOMMENDED time duration for sustain phase is 300 seconds. This is the phase where measurements occur. The test equipment SHOULD measure and record statistics continuously. The sampling interval for collecting the raw results and calculating the statistics SHOULD be less than 2 seconds.

4. **Ramp down phase**: No new connections are established, and no measurements are made. The time duration for ramp up and ramp down phase SHOULD be the same.

5. **Collection phase**: The last phase is administrative and will occur when the test equipment merges and collates the report data.

5. **Testbed Considerations**

This section describes steps for a reference test (pre-test) that control the test environment including test equipment, focusing on physical and virtualized environments and as well as test equipments. Below are the RECOMMENDED steps for the reference test.
1. Perform the reference test either by configuring the DUT/SUT in the most trivial setup (fast forwarding) or without presence of the DUT/SUT.

2. Generate traffic from traffic generator. Choose a traffic profile used for HTTP or HTTPS throughput performance test with smallest object size.

3. Ensure that any ancillary switching or routing functions added in the test equipment does not limit the performance by introducing network metrics such as packet loss and latency. This is specifically important for virtualized components (e.g., vSwitches, vRouters).

4. Verify that the generated traffic (performance) of the test equipment matches and reasonably exceeds the expected maximum performance of the DUT/SUT.

5. Record the network performance metrics packet loss latency introduced by the test environment (without DUT/SUT).

6. Assert that the testbed characteristics are stable during the entire test session. Several factors might influence stability specifically, for virtualized testbeds. For example, additional workloads in a virtualized system, load balancing, and movement of virtual machines during the test, or simple issues such as additional heat created by high workloads leading to an emergency CPU performance reduction.

The reference test SHOULD be performed before the benchmarking tests (described in section 7) start.

6. Reporting

This section describes how the benchmarking test report should be formatted and presented. It is RECOMMENDED to include two main sections in the report, namely the introduction and the detailed test results sections.

6.1. Introduction

The following attributes SHOULD be present in the introduction section of the test report.

1. The time and date of the execution of the tests

2. Summary of testbed software and hardware details
a. DUT/SUT hardware/virtual configuration
   * This section SHOULD clearly identify the make and model of the DUT/SUT
   * The port interfaces, including speed and link information
   * If the DUT/SUT is a Virtual Network Function (VNF), host (server) hardware and software details, interface acceleration type such as DPDK and SR-IOV, used CPU cores, used RAM, resource sharing (e.g. Pinning details and NUMA Node) configuration details, hypervisor version, virtual switch version
   * details of any additional hardware relevant to the DUT/SUT such as controllers

b. DUT/SUT software
   * Operating system name
   * Version
   * Specific configuration details (if any)

c. DUT/SUT enabled features
   * Configured DUT/SUT features (see Table 1 and Table 2)
   * Attributes of the above-mentioned features
   * Any additional relevant information about the features

d. Test equipment hardware and software
   * Test equipment vendor name
   * Hardware details including model number, interface type
   * Test equipment firmware and test application software version

e. Key test parameters
   * Used cipher suites and keys
   * IPv4 and IPv6 traffic distribution
3. Results Summary / Executive Summary

   a. Results SHOULD resemble a pyramid in how it is reported, with the introduction section documenting the summary of results in a prominent, easy to read block.

6.2. Detailed Test Results

   In the result section of the test report, the following attributes SHOULD be present for each benchmarking test.

   a. KPIs MUST be documented separately for each benchmarking test. The format of the KPI metrics SHOULD be presented as described in Section 6.3.

   b. The next level of details SHOULD be graphs showing each of these metrics over the duration (sustain phase) of the test. This allows the user to see the measured performance stability changes over time.

6.3. Benchmarks and Key Performance Indicators

   This section lists key performance indicators (KPIs) for overall benchmarking tests. All KPIs MUST be measured during the sustain phase of the traffic load profile described in Section 4.3.4. All KPIs MUST be measured from the result output of test equipment.

   * Concurrent TCP Connections
The aggregate number of simultaneous connections between hosts across the DUT/SUT, or between hosts and the DUT/SUT (defined in [RFC2647]).

* **TCP Connections Per Second**

The average number of successfully established TCP connections per second between hosts across the DUT/SUT, or between hosts and the DUT/SUT. The TCP connection MUST be initiated via a TCP three-way handshake (SYN, SYN/ACK, ACK). Then the TCP session data is sent. The TCP session MUST be closed via either a TCP three-way close (FIN, FIN/ACK, ACK), or a TCP four-way close (FIN, ACK, FIN, ACK), and MUST NOT by RST.

* **Application Transactions Per Second**

The average number of successfully completed transactions per second. For a particular transaction to be considered successful, all data MUST have been transferred in its entirety. In case of HTTP(S) transactions, it MUST have a valid status code (200 OK), and the appropriate FIN, FIN/ACK sequence MUST have been completed.

* **TLS Handshake Rate**

The average number of successfully established TLS connections per second between hosts across the DUT/SUT, or between hosts and the DUT/SUT.

* **Inspected Throughput**

The number of bits per second of examined and allowed traffic a network security device is able to transmit to the correct destination interface(s) in response to a specified offered load. The throughput benchmarking tests defined in Section 7 SHOULD measure the average Layer 2 throughput value when the DUT/SUT is "inspecting" traffic. This document recommends presenting the inspected throughput value in Gbit/s rounded to two places of precision with a more specific Kbit/s in parenthesis.

* **Time to First Byte (TTFB)**

TTFB is the elapsed time between the start of sending the TCP SYN packet from the client and the client receiving the first packet of application data from the server or DUT/SUT. The benchmarking tests HTTP Transaction Latency (Section 7.4) and HTTPS Transaction Latency (Section 7.8) measure the minimum, average and maximum TTFB. The value SHOULD be expressed in milliseconds.
* URL Response time / Time to Last Byte (TTLB)

URL Response time / TTLB is the elapsed time between the start of sending the TCP SYN packet from the client and the client receiving the last packet of application data from the server or DUT/SUT. The benchmarking tests HTTP Transaction Latency (Section 7.4) and HTTPS Transaction Latency (Section 7.8) measure the minimum, average and maximum TTLB. The value SHOULD be expressed in millisecond.

7. Benchmarking Tests

7.1. Throughput Performance with Application Traffic Mix

7.1.1. Objective

Using a relevant application traffic mix, determine the sustainable inspected throughput supported by the DUT/SUT.

Based on the test customer’s specific use case, testers can choose the relevant application traffic mix for this test. The details about the traffic mix MUST be documented in the report. At least the following traffic mix details MUST be documented and reported together with the test results:

- Name of applications and layer 7 protocols
- Percentage of emulated traffic for each application and layer 7 protocol
- Percentage of encrypted traffic and used cipher suites and keys (The RECOMMENDED ciphers and keys are defined in Section 4.3.1.3.)
- Used object sizes for each application and layer 7 protocols

7.1.2. Test Setup

Testbed setup MUST be configured as defined in Section 4. Any benchmarking test specific testbed configuration changes MUST be documented.

7.1.3. Test Parameters

In this section, the benchmarking test specific parameters SHOULD be defined.
7.1.3.1. DUT/SUT Configuration Parameters

DUT/SUT parameters MUST conform to the requirements defined in Section 4.2. Any configuration changes for this specific benchmarking test MUST be documented. In case the DUT/SUT is configured without SSL inspection, the test report MUST explain the implications of this to the relevant application traffic mix encrypted traffic.

7.1.3.2. Test Equipment Configuration Parameters

Test equipment configuration parameters MUST conform to the requirements defined in Section 4.3. The following parameters MUST be documented for this benchmarking test:

   - Client IP address range defined in Section 4.3.1.2
   - Server IP address range defined in Section 4.3.2.2
   - Traffic distribution ratio between IPv4 and IPv6 defined in Section 4.3.1.2
   - Target inspected throughput: Aggregated line rate of interface(s) used in the DUT/SUT or the value defined based on requirement for a specific deployment scenario
   - Initial throughput: 10% of the "Target inspected throughput" Note: Initial throughput is not a KPI to report. This value is configured on the traffic generator and used to perform Step 1: "Test Initialization and Qualification" described under the Section 7.1.4.

   One of the ciphers and keys defined in Section 4.3.1.3 are RECOMMENDED to use for this benchmarking test.

7.1.3.3. Traffic Profile

Traffic profile: This test MUST be run with a relevant application traffic mix profile.

7.1.3.4. Test Results Validation Criteria

The following criteria are the test results validation criteria. The test results validation criteria MUST be monitored during the whole sustain phase of the traffic load profile.
a. Number of failed application transactions (receiving any HTTP response code other than 200 OK) MUST be less than 0.001% (1 out of 100,000 transactions) of total attempted transactions.

b. Number of Terminated TCP connections due to unexpected TCP RST sent by DUT/SUT MUST be less than 0.001% (1 out of 100,000 connections) of total initiated TCP connections.

7.1.3.5. Measurement

Following KPI metrics MUST be reported for this benchmarking test:

Mandatory KPIs (benchmarks): Inspected Throughput, TTFB (minimum, average, and maximum), TTLB (minimum, average, and maximum) and Application Transactions Per Second

Note: TTLB MUST be reported along with the object size used in the traffic profile.

Optional KPIs: TCP Connections Per Second and TLS Handshake Rate

7.1.4. Test Procedures and Expected Results

The test procedures are designed to measure the inspected throughput performance of the DUT/SUT at the sustaining period of traffic load profile. The test procedure consists of three major steps: Step 1 ensures the DUT/SUT is able to reach the performance value (initial throughput) and meets the test results validation criteria when it was very minimally utilized. Step 2 determines the DUT/SUT is able to reach the target performance value within the test results validation criteria. Step 3 determines the maximum achievable performance value within the test results validation criteria.

This test procedure MAY be repeated multiple times with different IP types: IPv4 only, IPv6 only, and IPv4 and IPv6 mixed traffic distribution.

7.1.4.1. Step 1: Test Initialization and Qualification

Verify the link status of all connected physical interfaces. All interfaces are expected to be in "UP" status.
Configure traffic load profile of the test equipment to generate test traffic at the "Initial throughput" rate as described in Section 7.1.3.2. The test equipment SHOULD follow the traffic load profile definition as described in Section 4.3.4. The DUT/SUT SHOULD reach the "Initial throughput" during the sustain phase. Measure all KPI as defined in Section 7.1.3.5. The measured KPIs during the sustain phase MUST meet all the test results validation criteria defined in Section 7.1.3.4.

If the KPI metrics do not meet the test results validation criteria, the test procedure MUST NOT be continued to step 2.

7.1.4.2. Step 2: Test Run with Target Objective

Configure test equipment to generate traffic at the "Target inspected throughput" rate defined in Section 7.1.3.2. The test equipment SHOULD follow the traffic load profile definition as described in Section 4.3.4. The test equipment SHOULD start to measure and record all specified KPIs. Continue the test until all traffic profile phases are completed.

Within the test results validation criteria, the DUT/SUT is expected to reach the desired value of the target objective ("Target inspected throughput") in the sustain phase. Follow step 3, if the measured value does not meet the target value or does not fulfill the test results validation criteria.

7.1.4.3. Step 3: Test Iteration

Determine the achievable average inspected throughput within the test results validation criteria. Final test iteration MUST be performed for the test duration defined in Section 4.3.4.

7.2. TCP/HTTP Connections Per Second

7.2.1. Objective

Using HTTP traffic, determine the sustainable TCP connection establishment rate supported by the DUT/SUT under different throughput load conditions.

To measure connections per second, test iterations MUST use different fixed HTTP response object sizes (the different load conditions) defined in Section 7.2.3.2.
7.2.2. Test Setup

Testbed setup SHOULD be configured as defined in Section 4. Any specific testbed configuration changes (number of interfaces and interface type, etc.) MUST be documented.

7.2.3. Test Parameters

In this section, benchmarking test specific parameters SHOULD be defined.

7.2.3.1. DUT/SUT Configuration Parameters

DUT/SUT parameters MUST conform to the requirements defined in Section 4.2. Any configuration changes for this specific benchmarking test MUST be documented.

7.2.3.2. Test Equipment Configuration Parameters

Test equipment configuration parameters MUST conform to the requirements defined in Section 4.3. The following parameters MUST be documented for this benchmarking test:

- Client IP address range defined in Section 4.3.1.2
- Server IP address range defined in Section 4.3.2.2
- Traffic distribution ratio between IPv4 and IPv6 defined in Section 4.3.1.2
- Target connections per second: Initial value from product datasheet or the value defined based on requirement for a specific deployment scenario
- Initial connections per second: 10% of "Target connections per second" (Note: Initial connections per second is not a KPI to report. This value is configured on the traffic generator and used to perform the Step1: "Test Initialization and Qualification" described under the Section 7.2.4.

The client SHOULD negotiate HTTP and close the connection with FIN immediately after completion of one transaction. In each test iteration, client MUST send GET request requesting a fixed HTTP response object size.

The RECOMMENDED response object sizes are 1, 2, 4, 16, and 64 KByte.
7.2.3.3. Test Results Validation Criteria

The following criteria are the test results validation criteria. The test results validation criteria MUST be monitored during the whole sustain phase of the traffic load profile.

a. Number of failed application transactions (receiving any HTTP response code other than 200 OK) MUST be less than 0.001% (1 out of 100,000 transactions) of total attempted transactions.

b. Number of terminated TCP connections due to unexpected TCP RST sent by DUT/SUT MUST be less than 0.001% (1 out of 100,000 connections) of total initiated TCP connections.

c. During the sustain phase, traffic SHOULD be forwarded at a constant rate (considered as a constant rate if any deviation of traffic forwarding rate is less than 5%).

d. Concurrent TCP connections MUST be constant during steady state and any deviation of concurrent TCP connections SHOULD be less than 10%. This confirms the DUT opens and closes TCP connections at approximately the same rate.

7.2.3.4. Measurement

TCP Connections Per Second MUST be reported for each test iteration (for each object size).

7.2.4. Test Procedures and Expected Results

The test procedure is designed to measure the TCP connections per second rate of the DUT/SUT at the sustaining period of the traffic load profile. The test procedure consists of three major steps: Step 1 ensures the DUT/SUT is able to reach the performance value (Initial connections per second) and meets the test results validation criteria when it was very minimally utilized. Step 2 determines the DUT/SUT is able to reach the target performance value within the test results validation criteria. Step 3 determines the maximum achievable performance value within the test results validation criteria.

This test procedure MAY be repeated multiple times with different IP types: IPv4 only, IPv6 only, and IPv4 and IPv6 mixed traffic distribution.
7.2.4.1. Step 1: Test Initialization and Qualification

Verify the link status of all connected physical interfaces. All interfaces are expected to be in "UP" status.

Configure the traffic load profile of the test equipment to establish "Initial connections per second" as defined in Section 7.2.3.2. The traffic load profile SHOULD be defined as described in Section 4.3.4.

The DUT/SUT SHOULD reach the "Initial connections per second" before the sustain phase. The measured KPIs during the sustain phase MUST meet all the test results validation criteria defined in Section 7.2.3.3.

If the KPI metrics do not meet the test results validation criteria, the test procedure MUST NOT continue to "Step 2".

7.2.4.2. Step 2: Test Run with Target Objective

Configure test equipment to establish the target objective ("Target connections per second") defined in Section 7.2.3.2. The test equipment SHOULD follow the traffic load profile definition as described in Section 4.3.4.

During the ramp up and sustain phase of each test iteration, other KPIs such as inspected throughput, concurrent TCP connections and application transactions per second MUST NOT reach the maximum value the DUT/SUT can support. The test results for specific test iterations SHOULD NOT be reported, if the above-mentioned KPI (especially inspected throughput) reaches the maximum value. (Example: If the test iteration with 64 KByte of HTTP response object size reached the maximum inspected throughput limitation of the DUT/SUT, the test iteration MAY be interrupted and the result for 64 KByte SHOULD NOT be reported.)

The test equipment SHOULD start to measure and record all specified KPIs. Continue the test until all traffic profile phases are completed.

Within the test results validation criteria, the DUT/SUT is expected to reach the desired value of the target objective ("Target connections per second") in the sustain phase. Follow step 3, if the measured value does not meet the target value or does not fulfill the test results validation criteria.
7.2.4.3. Step 3: Test Iteration

Determine the achievable TCP connections per second within the test results validation criteria.

7.3. HTTP Throughput

7.3.1. Objective

Determine the sustainable inspected throughput of the DUT/SUT for HTTP transactions varying the HTTP response object size.

7.3.2. Test Setup

Testbed setup SHOULD be configured as defined in Section 4. Any specific testbed configuration changes (number of interfaces and interface type, etc.) MUST be documented.

7.3.3. Test Parameters

In this section, benchmarking test specific parameters SHOULD be defined.

7.3.3.1. DUT/SUT Configuration Parameters

DUT/SUT parameters MUST conform to the requirements defined in Section 4.2. Any configuration changes for this specific benchmarking test MUST be documented.

7.3.3.2. Test Equipment Configuration Parameters

Test equipment configuration parameters MUST conform to the requirements defined in Section 4.3. The following parameters MUST be documented for this benchmarking test:

Client IP address range defined in Section 4.3.1.2

Server IP address range defined in Section 4.3.2.2

Traffic distribution ratio between IPv4 and IPv6 defined in Section 4.3.1.2

Target inspected throughput: Aggregated line rate of interface(s) used in the DUT/SUT or the value defined based on requirement for a specific deployment scenario
Initial throughput: 10% of "Target inspected throughput" Note: Initial throughput is not a KPI to report. This value is configured on the traffic generator and used to perform Step 1: "Test Initialization and Qualification" described under Section 7.3.4.

Number of HTTP response object requests (transactions) per connection: 10

RECOMMENDED HTTP response object size: 1, 16, 64, 256 KByte, and mixed objects defined in Table 4.

<table>
<thead>
<tr>
<th>Object size (KByte)</th>
<th>Number of requests/ Weight</th>
</tr>
</thead>
<tbody>
<tr>
<td>0.2</td>
<td>1</td>
</tr>
<tr>
<td>6</td>
<td>1</td>
</tr>
<tr>
<td>8</td>
<td>1</td>
</tr>
<tr>
<td>9</td>
<td>1</td>
</tr>
<tr>
<td>10</td>
<td>1</td>
</tr>
<tr>
<td>25</td>
<td>1</td>
</tr>
<tr>
<td>26</td>
<td>1</td>
</tr>
<tr>
<td>35</td>
<td>1</td>
</tr>
<tr>
<td>59</td>
<td>1</td>
</tr>
<tr>
<td>347</td>
<td>1</td>
</tr>
</tbody>
</table>

Table 4: Mixed Objects

7.3.3.3. Test Results Validation Criteria

The following criteria are the test results validation criteria. The test results validation criteria MUST be monitored during the whole sustain phase of the traffic load profile.

a. Number of failed application transactions (receiving any HTTP response code other than 200 OK) MUST be less than 0.001% (1 out of 100,000 transactions) of attempt transactions.
b. Traffic SHOULD be forwarded at a constant rate (considered as a constant rate if any deviation of traffic forwarding rate is less than 5%).

c. Concurrent TCP connections MUST be constant during steady state and any deviation of concurrent TCP connections SHOULD be less than 10%. This confirms the DUT opens and closes TCP connections at approximately the same rate.

7.3.3.4. Measurement

Inspected Throughput and HTTP Transactions per Second MUST be reported for each object size.

7.3.4. Test Procedures and Expected Results

The test procedure is designed to measure HTTP throughput of the DUT/SUT. The test procedure consists of three major steps: Step 1 ensures the DUT/SUT is able to reach the performance value (Initial throughput) and meets the test results validation criteria when it was very minimal utilized. Step 2 determines the DUT/SUT is able to reach the target performance value within the test results validation criteria. Step 3 determines the maximum achievable performance value within the test results validation criteria.

This test procedure MAY be repeated multiple times with different IPv4 and IPv6 traffic distribution and HTTP response object sizes.

7.3.4.1. Step 1: Test Initialization and Qualification

Verify the link status of all connected physical interfaces. All interfaces are expected to be in "UP" status.

Configure traffic load profile of the test equipment to establish "Initial inspected throughput" as defined in Section 7.3.3.2.

The traffic load profile SHOULD be defined as described in Section 4.3.4. The DUT/SUT SHOULD reach the "Initial inspected throughput" during the sustain phase. Measure all KPI as defined in Section 7.3.3.4.

The measured KPIs during the sustain phase MUST meet the test results validation criteria "a" defined in Section 7.3.3.3. The test results validation criteria "b" and "c" are OPTIONAL for step 1.

If the KPI metrics do not meet the test results validation criteria, the test procedure MUST NOT be continued to "Step 2".
7.3.4.2. Step 2: Test Run with Target Objective

Configure test equipment to establish the target objective ("Target inspected throughput") defined in Section 7.3.3.2. The test equipment SHOULD start to measure and record all specified KPIs. Continue the test until all traffic profile phases are completed.

Within the test results validation criteria, the DUT/SUT is expected to reach the desired value of the target objective in the sustain phase. Follow step 3, if the measured value does not meet the target value or does not fulfill the test results validation criteria.

7.3.4.3. Step 3: Test Iteration

Determine the achievable inspected throughput within the test results validation criteria and measure the KPI metric Transactions per Second. Final test iteration MUST be performed for the test duration defined in Section 4.3.4.

7.4. HTTP Transaction Latency

7.4.1. Objective

Using HTTP traffic, determine the HTTP transaction latency when DUT is running with sustainable HTTP transactions per second supported by the DUT/SUT under different HTTP response object sizes.

Test iterations MUST be performed with different HTTP response object sizes in two different scenarios. One with a single transaction and the other with multiple transactions within a single TCP connection. For consistency both the single and multiple transaction test MUST be configured with the same HTTP version.

Scenario 1: The client MUST negotiate HTTP and close the connection with FIN immediately after completion of a single transaction (GET and RESPONSE).

Scenario 2: The client MUST negotiate HTTP and close the connection FIN immediately after completion of 10 transactions (GET and RESPONSE) within a single TCP connection.

7.4.2. Test Setup

Testbed setup SHOULD be configured as defined in Section 4. Any specific testbed configuration changes (number of interfaces and interface type, etc.) MUST be documented.
7.4.3. Test Parameters

In this section, benchmarking test specific parameters SHOULD be defined.

7.4.3.1. DUT/SUT Configuration Parameters

DUT/SUT parameters MUST conform to the requirements defined in Section 4.2. Any configuration changes for this specific benchmarking test MUST be documented.

7.4.3.2. Test Equipment Configuration Parameters

Test equipment configuration parameters MUST conform to the requirements defined in Section 4.3. The following parameters MUST be documented for this benchmarking test:

- Client IP address range defined in Section 4.3.1.2
- Server IP address range defined in Section 4.3.2.2
- Traffic distribution ratio between IPv4 and IPv6 defined in Section 4.3.1.2

Target objective for scenario 1: 50% of the connections per second measured in benchmarking test TCP/HTTP Connections Per Second (Section 7.2)

Target objective for scenario 2: 50% of the inspected throughput measured in benchmarking test HTTP Throughput (Section 7.3)

Initial objective for scenario 1: 10% of "Target objective for scenario 1"

Initial objective for scenario 2: 10% of "Target objective for scenario 2"

Note: The Initial objectives are not a KPI to report. These values are configured on the traffic generator and used to perform the Step1: "Test Initialization and Qualification" described under the Section 7.4.4.

HTTP transaction per TCP connection: Test scenario 1 with single transaction and test scenario 2 with 10 transactions.
HTTP with GET request requesting a single object. The RECOMMENDED object sizes are 1, 16, and 64 KByte. For each test iteration, client MUST request a single HTTP response object size.

7.4.3.3. Test Results Validation Criteria

The following criteria are the test results validation criteria. The Test results validation criteria MUST be monitored during the whole sustain phase of the traffic load profile.

a. Number of failed application transactions (receiving any HTTP response code other than 200 OK) MUST be less than 0.001% (1 out of 100,000 transactions) of attempt transactions.

b. Number of terminated TCP connections due to unexpected TCP RST sent by DUT/SUT MUST be less than 0.001% (1 out of 100,000 connections) of total initiated TCP connections.

c. During the sustain phase, traffic SHOULD be forwarded at a constant rate (considered as a constant rate if any deviation of traffic forwarding rate is less than 5%).

d. Concurrent TCP connections MUST be constant during steady state and any deviation of concurrent TCP connections SHOULD be less than 10%. This confirms the DUT opens and closes TCP connections at approximately the same rate.

e. After ramp up the DUT MUST achieve the "Target objective" defined in Section 7.4.3.2 and remain in that state for the entire test duration (sustain phase).

7.4.3.4. Measurement

TTFB (minimum, average, and maximum) and TTLB (minimum, average and maximum) MUST be reported for each object size.

7.4.4. Test Procedures and Expected Results

The test procedure is designed to measure TTFB or TTLB when the DUT/SUT is operating close to 50% of its maximum achievable connections per second or inspected throughput. The test procedure consists of two major steps: Step 1 ensures the DUT/SUT is able to reach the initial performance values and meets the test results validation criteria when it was very minimally utilized. Step 2 measures the latency values within the test results validation criteria.
This test procedure MAY be repeated multiple times with different IP types (IPv4 only, IPv6 only and IPv4 and IPv6 mixed traffic distribution), HTTP response object sizes and single and multiple transactions per connection scenarios.

7.4.4.1. Step 1: Test Initialization and Qualification

Verify the link status of all connected physical interfaces. All interfaces are expected to be in "UP" status.

Configure traffic load profile of the test equipment to establish "Initial objective" as defined in Section 7.4.3.2. The traffic load profile SHOULD be defined as described in Section 4.3.4.

The DUT/SUT SHOULD reach the "Initial objective" before the sustain phase. The measured KPIs during the sustain phase MUST meet all the test results validation criteria defined in Section 7.4.3.3.

If the KPI metrics do not meet the test results validation criteria, the test procedure MUST NOT be continued to "Step 2".

7.4.4.2. Step 2: Test Run with Target Objective

Configure test equipment to establish "Target objective" defined in Section 7.4.3.2. The test equipment SHOULD follow the traffic load profile definition as described in Section 4.3.4.

The test equipment SHOULD start to measure and record all specified KPIs. Continue the test until all traffic profile phases are completed.

Within the test results validation criteria, the DUT/SUT MUST reach the desired value of the target objective in the sustain phase.

Measure the minimum, average, and maximum values of TTFB and TTLB.

7.5. Concurrent TCP/HTTP Connection Capacity

7.5.1. Objective

Determine the number of concurrent TCP connections that the DUT/SUT sustains when using HTTP traffic.

7.5.2. Test Setup

Testbed setup SHOULD be configured as defined in Section 4. Any specific testbed configuration changes (number of interfaces and interface type, etc.) MUST be documented.
7.5.3. Test Parameters

In this section, benchmarking test specific parameters SHOULD be defined.

7.5.3.1. DUT/SUT Configuration Parameters

DUT/SUT parameters MUST conform to the requirements defined in Section 4.2. Any configuration changes for this specific benchmarking test MUST be documented.

7.5.3.2. Test Equipment Configuration Parameters

Test equipment configuration parameters MUST conform to the requirements defined in Section 4.3. The following parameters MUST be noted for this benchmarking test:

- Client IP address range defined in Section 4.3.1.2
- Server IP address range defined in Section 4.3.2.2
- Traffic distribution ratio between IPv4 and IPv6 defined in Section 4.3.1.2
- Target concurrent connection: Initial value from product datasheet or the value defined based on requirement for a specific deployment scenario.

  Initial concurrent connection: 10% of "Target concurrent connection" Note: Initial concurrent connection is not a KPI to report. This value is configured on the traffic generator and used to perform the Step1: "Test Initialization and Qualification" described under the Section 7.5.4.

- Maximum connections per second during ramp up phase: 50% of maximum connections per second measured in benchmarking test TCP/HTTP Connections per second (Section 7.2)
- Ramp up time (in traffic load profile for "Target concurrent connection"): "Target concurrent connection" / "Maximum connections per second during ramp up phase"
- Ramp up time (in traffic load profile for "Initial concurrent connection"): "Initial concurrent connection" / "Maximum connections per second during ramp up phase"

The client MUST negotiate HTTP and each client MAY open multiple concurrent TCP connections per server endpoint IP.
Each client sends 10 GET requests requesting 1 KByte HTTP response object in the same TCP connection (10 transactions/TCP connection) and the delay (think time) between each transaction MUST be X seconds.

\[ X = \frac{\text{"Ramp up time" + "steady state time"}}{10} \]

The established connections SHOULD remain open until the ramp down phase of the test. During the ramp down phase, all connections SHOULD be successfully closed with FIN.

### 7.5.3.3. Test Results Validation Criteria

The following criteria are the test results validation criteria. The Test results validation criteria MUST be monitored during the whole sustain phase of the traffic load profile.

a. Number of failed application transactions (receiving any HTTP response code other than 200 OK) MUST be less than 0.001% (1 out of 100,000 transaction) of total attempted transactions.

b. Number of terminated TCP connections due to unexpected TCP RST sent by DUT/SUT MUST be less than 0.001% (1 out of 100,000 connections) of total initiated TCP connections.

c. During the sustain phase, traffic SHOULD be forwarded at a constant rate (considered as a constant rate if any deviation of traffic forwarding rate is less than 5%).

### 7.5.3.4. Measurement

Average Concurrent TCP Connections MUST be reported for this benchmarking test.

### 7.5.4. Test Procedures and Expected Results

The test procedure is designed to measure the concurrent TCP connection capacity of the DUT/SUT at the sustaining period of traffic load profile. The test procedure consists of three major steps: Step 1 ensures the DUT/SUT is able to reach the performance value (Initial concurrent connection) and meets the test results validation criteria when it was very minimally utilized. Step 2 determines the DUT/SUT is able to reach the target performance value within the test results validation criteria. Step 3 determines the maximum achievable performance value within the test results validation criteria.
This test procedure MAY be repeated multiple times with different IPv4 and IPv6 traffic distribution.

7.5.4.1. Step 1: Test Initialization and Qualification

Verify the link status of all connected physical interfaces. All interfaces are expected to be in "UP" status.

Configure test equipment to establish "Initial concurrent TCP connections" defined in Section 7.5.3.2. Except ramp up time, the traffic load profile SHOULD be defined as described in Section 4.3.4.

During the sustain phase, the DUT/SUT SHOULD reach the "Initial concurrent TCP connections". The measured KPIs during the sustain phase MUST meet all the test results validation criteria defined in Section 7.5.3.3.

If the KPI metrics do not meet the test results validation criteria, the test procedure MUST NOT be continued to "Step 2".

7.5.4.2. Step 2: Test Run with Target Objective

Configure test equipment to establish the target objective ("Target concurrent TCP connections"). The test equipment SHOULD follow the traffic load profile definition (except ramp up time) as described in Section 4.3.4.

During the ramp up and sustain phase, the other KPIs such as inspected throughput, TCP connections per second, and application transactions per second MUST NOT reach the maximum value the DUT/SUT can support.

The test equipment SHOULD start to measure and record KPIs defined in Section 7.5.3.4. Continue the test until all traffic profile phases are completed.

Within the test results validation criteria, the DUT/SUT is expected to reach the desired value of the target objective in the sustain phase. Follow step 3, if the measured value does not meet the target value or does not fulfill the test results validation criteria.

7.5.4.3. Step 3: Test Iteration

Determine the achievable concurrent TCP connections capacity within the test results validation criteria.

7.6. TCP/HTTPS Connections per Second
7.6.1. Objective

Using HTTPS traffic, determine the sustainable SSL/TLS session establishment rate supported by the DUT/SUT under different throughput load conditions.

Test iterations MUST include common cipher suites and key strengths as well as forward looking stronger keys. Specific test iterations MUST include ciphers and keys defined in Section 7.6.3.2.

For each cipher suite and key strengths, test iterations MUST use a single HTTPS response object size defined in Section 7.6.3.2 to measure connections per second performance under a variety of DUT/SUT security inspection load conditions.

7.6.2. Test Setup

Testbed setup SHOULD be configured as defined in Section 4. Any specific testbed configuration changes (number of interfaces and interface type, etc.) MUST be documented.

7.6.3. Test Parameters

In this section, benchmarking test specific parameters SHOULD be defined.

7.6.3.1. DUT/SUT Configuration Parameters

DUT/SUT parameters MUST conform to the requirements defined in Section 4.2. Any configuration changes for this specific benchmarking test MUST be documented.

7.6.3.2. Test Equipment Configuration Parameters

Test equipment configuration parameters MUST conform to the requirements defined in Section 4.3. The following parameters MUST be documented for this benchmarking test:

Client IP address range defined in Section 4.3.1.2

Server IP address range defined in Section 4.3.2.2

Traffic distribution ratio between IPv4 and IPv6 defined in Section 4.3.1.2

Target connections per second: Initial value from product datasheet or the value defined based on requirement for a specific deployment scenario.
Initial connections per second: 10% of "Target connections per second" Note: Initial connections per second is not a KPI to report. This value is configured on the traffic generator and used to perform the Step1: "Test Initialization and Qualification" described under the Section 7.6.4.

RECOMMENDED ciphers and keys defined in Section 4.3.1.3

The client MUST negotiate HTTPS and close the connection with FIN immediately after completion of one transaction. In each test iteration, client MUST send GET request requesting a fixed HTTPS response object size. The RECOMMENDED object sizes are 1, 2, 4, 16, and 64 KByte.

7.6.3.3. Test Results Validation Criteria

The following criteria are the test results validation criteria. The test results validation criteria MUST be monitored during the whole test duration.

a. Number of failed application transactions (receiving any HTTP response code other than 200 OK) MUST be less than 0.001% (1 out of 100,000 transactions) of attempt transactions.

b. Number of terminated TCP connections due to unexpected TCP RST sent by DUT/SUT MUST be less than 0.001% (1 out of 100,000 connections) of total initiated TCP connections.

c. During the sustain phase, traffic SHOULD be forwarded at a constant rate (considered as a constant rate if any deviation of traffic forwarding rate is less than 5%).

d. Concurrent TCP connections MUST be constant during steady state and any deviation of concurrent TCP connections SHOULD be less than 10%. This confirms the DUT opens and closes TCP connections at approximately the same rate.

7.6.3.4. Measurement

TCP connections per second MUST be reported for each test iteration (for each object size).

The KPI metric TLS Handshake Rate can be measured in the test using 1 KByte object size.
7.6.4. Test Procedures and Expected Results

The test procedure is designed to measure the TCP connections per second rate of the DUT/SUT at the sustaining period of traffic load profile. The test procedure consists of three major steps: Step 1 ensures the DUT/SUT is able to reach the performance value (Initial connections per second) and meets the test results validation criteria when it was very minimally utilized. Step 2 determines the DUT/SUT is able to reach the target performance value within the test results validation criteria. Step 3 determines the maximum achievable performance value within the test results validation criteria.

This test procedure MAY be repeated multiple times with different IPv4 and IPv6 traffic distribution.

7.6.4.1. Step 1: Test Initialization and Qualification

Verify the link status of all connected physical interfaces. All interfaces are expected to be in "UP" status.

Configure traffic load profile of the test equipment to establish "Initial connections per second" as defined in Section 7.6.3.2. The traffic load profile SHOULD be defined as described in Section 4.3.4.

The DUT/SUT SHOULD reach the "Initial connections per second" before the sustain phase. The measured KPIs during the sustain phase MUST meet all the test results validation criteria defined in Section 7.6.3.3.

If the KPI metrics do not meet the test results validation criteria, the test procedure MUST NOT be continued to "Step 2".

7.6.4.2. Step 2: Test Run with Target Objective

Configure test equipment to establish "Target connections per second" defined in Section 7.6.3.2. The test equipment SHOULD follow the traffic load profile definition as described in Section 4.3.4.

During the ramp up and sustain phase, other KPIs such as inspected throughput, concurrent TCP connections, and application transactions per second MUST NOT reach the maximum value the DUT/SUT can support. The test results for specific test iteration SHOULD NOT be reported, if the above mentioned KPI (especially inspected throughput) reaches the maximum value. (Example: If the test iteration with 64 KByte of HTTPS response object size reached the maximum inspected throughput limitation of the DUT, the test iteration MAY be interrupted and the result for 64 KByte SHOULD NOT be reported).
The test equipment SHOULD start to measure and record all specified KPIs. Continue the test until all traffic profile phases are completed.

Within the test results validation criteria, the DUT/SUT is expected to reach the desired value of the target objective ("Target connections per second") in the sustain phase. Follow step 3, if the measured value does not meet the target value or does not fulfill the test results validation criteria.

7.6.4.3. Step 3: Test Iteration

Determine the achievable connections per second within the test results validation criteria.

7.7. HTTPS Throughput

7.7.1. Objective

Determine the sustainable inspected throughput of the DUT/SUT for HTTPS transactions varying the HTTPS response object size.

Test iterations MUST include common cipher suites and key strengths as well as forward looking stronger keys. Specific test iterations MUST include the ciphers and keys defined in Section 7.7.3.2.

7.7.2. Test Setup

Testbed setup SHOULD be configured as defined in Section 4. Any specific testbed configuration changes (number of interfaces and interface type, etc.) MUST be documented.

7.7.3. Test Parameters

In this section, benchmarking test specific parameters SHOULD be defined.

7.7.3.1. DUT/SUT Configuration Parameters

DUT/SUT parameters MUST conform to the requirements defined in Section 4.2. Any configuration changes for this specific benchmarking test MUST be documented.

7.7.3.2. Test Equipment Configuration Parameters

Test equipment configuration parameters MUST conform to the requirements defined in Section 4.3. The following parameters MUST be documented for this benchmarking test:
Client IP address range defined in Section 4.3.1.2

Server IP address range defined in Section 4.3.2.2

Traffic distribution ratio between IPv4 and IPv6 defined in
Section 4.3.1.2

Target inspected throughput: Aggregated line rate of interface(s) used in the DUT/SUT or the value defined based on requirement for a specific deployment scenario.

Initial throughput: 10% of "Target inspected throughput" Note: Initial throughput is not a KPI to report. This value is configured on the traffic generator and used to perform the Step1: "Test Initialization and Qualification" described under the Section 7.7.4.

Number of HTTPS response object requests (transactions) per connection: 10

RECOMMENDED ciphers and keys defined in Section 4.3.1.3

RECOMMENDED HTTPS response object size: 1, 16, 64, 256 KByte, and mixed objects defined in Table 4 under Section 7.3.3.2.

7.7.3.3. Test Results Validation Criteria

The following criteria are the test results validation criteria. The test results validation criteria MUST be monitored during the whole sustain phase of the traffic load profile.

a. Number of failed Application transactions (receiving any HTTP response code other than 200 OK) MUST be less than 0.001% (1 out of 100,000 transactions) of attempt transactions.

b. Traffic SHOULD be forwarded at a constant rate (considered as a constant rate if any deviation of traffic forwarding rate is less than 5%).

c. Concurrent TCP connections MUST be constant during steady state and any deviation of concurrent TCP connections SHOULD be less than 10%. This confirms the DUT opens and closes TCP connections at approximately the same rate.

7.7.3.4. Measurement

Inspected Throughput and HTTP Transactions per Second MUST be reported for each object size.
7.7.4. Test Procedures and Expected Results

The test procedure consists of three major steps: Step 1 ensures the DUT/SUT is able to reach the performance value (Initial throughput) and meets the test results validation criteria when it was very minimally utilized. Step 2 determines the DUT/SUT is able to reach the target performance value within the test results validation criteria. Step 3 determines the maximum achievable performance value within the test results validation criteria.

This test procedure MAY be repeated multiple times with different IPv4 and IPv6 traffic distribution and HTTPS response object sizes.

7.7.4.1. Step 1: Test Initialization and Qualification

Verify the link status of all connected physical interfaces. All interfaces are expected to be in "UP" status.

Configure traffic load profile of the test equipment to establish "Initial throughput" as defined in Section 7.7.3.2.

The traffic load profile SHOULD be defined as described in Section 4.3.4. The DUT/SUT SHOULD reach the "Initial throughput" during the sustain phase. Measure all KPI as defined in Section 7.7.3.4.

The measured KPIs during the sustain phase MUST meet the test results validation criteria "a" defined in Section 7.7.3.3. The test results validation criteria "b" and "c" are OPTIONAL for step 1.

If the KPI metrics do not meet the test results validation criteria, the test procedure MUST NOT be continued to "Step 2".

7.7.4.2. Step 2: Test Run with Target Objective

Configure test equipment to establish the target objective ("Target inspected throughput") defined in Section 7.7.3.2. The test equipment SHOULD start to measure and record all specified KPIs. Continue the test until all traffic profile phases are completed.

Within the test results validation criteria, the DUT/SUT is expected to reach the desired value of the target objective in the sustain phase. Follow step 3, if the measured value does not meet the target value or does not fulfill the test results validation criteria.
7.7.4.3. Step 3: Test Iteration

Determine the achievable average inspected throughput within the test results validation criteria. Final test iteration MUST be performed for the test duration defined in Section 4.3.4.

7.8. HTTPS Transaction Latency

7.8.1. Objective

Using HTTPS traffic, determine the HTTPS transaction latency when DUT/SUT is running with sustainable HTTPS transactions per second supported by the DUT/SUT under different HTTPS response object size.

Scenario 1: The client MUST negotiate HTTPS and close the connection with FIN immediately after completion of a single transaction (GET and RESPONSE).

Scenario 2: The client MUST negotiate HTTPS and close the connection with FIN immediately after completion of 10 transactions (GET and RESPONSE) within a single TCP connection.

7.8.2. Test Setup

Testbed setup SHOULD be configured as defined in Section 4. Any specific testbed configuration changes (number of interfaces and interface type, etc.) MUST be documented.

7.8.3. Test Parameters

In this section, benchmarking test specific parameters SHOULD be defined.

7.8.3.1. DUT/SUT Configuration Parameters

DUT/SUT parameters MUST conform to the requirements defined in Section 4.2. Any configuration changes for this specific benchmarking test MUST be documented.

7.8.3.2. Test Equipment Configuration Parameters

Test equipment configuration parameters MUST conform to the requirements defined in Section 4.3. The following parameters MUST be documented for this benchmarking test:

Client IP address range defined in Section 4.3.1.2
Server IP address range defined in Section 4.3.2.2
Traffic distribution ratio between IPv4 and IPv6 defined in Section 4.3.1.2

RECOMMENDED cipher suites and key sizes defined in Section 4.3.1.3

Target objective for scenario 1: 50% of the connections per second measured in benchmarking test TCP/HTTPS Connections per second (Section 7.6)

Target objective for scenario 2: 50% of the inspected throughput measured in benchmarking test HTTPS Throughput (Section 7.7)

Initial objective for scenario 1: 10% of "Target objective for scenario 1"

Initial objective for scenario 2: 10% of "Target objective for scenario 2"

Note: The Initial objectives are not a KPI to report. These values are configured on the traffic generator and used to perform the Step 1: "Test Initialization and Qualification" described under the Section 7.8.4.

HTTPS transaction per TCP connection: Test scenario 1 with single transaction and scenario 2 with 10 transactions

HTTPS with GET request requesting a single object. The RECOMMENDED object sizes are 1, 16, and 64 KByte. For each test iteration, client MUST request a single HTTPS response object size.

7.8.3.3. Test Results Validation Criteria

The following criteria are the test results validation criteria. The Test results validation criteria MUST be monitored during the whole sustain phase of the traffic load profile.

a. Number of failed application transactions (receiving any HTTP response code other than 200 OK) MUST be less than 0.001% (1 out of 100,000 transactions) of attempt transactions.

b. Number of terminated TCP connections due to unexpected TCP RST sent by DUT/SUT MUST be less than 0.001% (1 out of 100,000 connections) of total initiated TCP connections.

c. During the sustain phase, traffic SHOULD be forwarded at a constant rate (considered as a constant rate if any deviation of traffic forwarding rate is less than 5%).
d. Concurrent TCP connections MUST be constant during steady state and any deviation of concurrent TCP connections SHOULD be less than 10%. This confirms the DUT opens and closes TCP connections at approximately the same rate.

e. After ramp up the DUT/SUT MUST achieve the "Target objective" defined in the parameter Section 7.8.3.2 and remain in that state for the entire test duration (sustain phase).

7.8.3.4. Measurement

TTFB (minimum, average, and maximum) and TTLB (minimum, average and maximum) MUST be reported for each object size.

7.8.4. Test Procedures and Expected Results

The test procedure is designed to measure TTFB or TTLB when the DUT/SUT is operating close to 50% of its maximum achievable connections per second or inspected throughput. The test procedure consists of two major steps: Step 1 ensures the DUT/SUT is able to reach the initial performance values and meets the test results validation criteria when it was very minimally utilized. Step 2 measures the latency values within the test results validation criteria.

This test procedure MAY be repeated multiple times with different IP types (IPv4 only, IPv6 only and IPv4 and IPv6 mixed traffic distribution), HTTPS response object sizes and single, and multiple transactions per connection scenarios.

7.8.4.1. Step 1: Test Initialization and Qualification

Verify the link status of all connected physical interfaces. All interfaces are expected to be in "UP" status.

Configure traffic load profile of the test equipment to establish "Initial objective" as defined in the Section 7.8.3.2. The traffic load profile SHOULD be defined as described in Section 4.3.4.

The DUT/SUT SHOULD reach the "Initial objective" before the sustain phase. The measured KPIs during the sustain phase MUST meet all the test results validation criteria defined in Section 7.8.3.3.

If the KPI metrics do not meet the test results validation criteria, the test procedure MUST NOT be continued to "Step 2".
7.8.4.2. Step 2: Test Run with Target Objective

Configure test equipment to establish "Target objective" defined in Section 7.8.3.2. The test equipment SHOULD follow the traffic load profile definition as described in Section 4.3.4.

The test equipment SHOULD start to measure and record all specified KPIs. Continue the test until all traffic profile phases are completed.

Within the test results validation criteria, the DUT/SUT MUST reach the desired value of the target objective in the sustain phase.

Measure the minimum, average, and maximum values of TTFB and TTLB.

7.9. Concurrent TCP/HTTPS Connection Capacity

7.9.1. Objective

Determine the number of concurrent TCP connections the DUT/SUT sustains when using HTTPS traffic.

7.9.2. Test Setup

Testbed setup SHOULD be configured as defined in Section 4. Any specific testbed configuration changes (number of interfaces and interface type, etc.) MUST be documented.

7.9.3. Test Parameters

In this section, benchmarking test specific parameters SHOULD be defined.

7.9.3.1. DUT/SUT Configuration Parameters

DUT/SUT parameters MUST conform to the requirements defined in Section 4.2. Any configuration changes for this specific benchmarking test MUST be documented.

7.9.3.2. Test Equipment Configuration Parameters

Test equipment configuration parameters MUST conform to the requirements defined in Section 4.3. The following parameters MUST be documented for this benchmarking test:

- Client IP address range defined in Section 4.3.1.2
- Server IP address range defined in Section 4.3.2.2
Traffic distribution ratio between IPv4 and IPv6 defined in Section 4.3.1.2

RECOMMENDED cipher suites and key sizes defined in Section 4.3.1.3

Target concurrent connections: Initial value from product datasheet or the value defined based on requirement for a specific deployment scenario.

Initial concurrent connections: 10% of "Target concurrent connections" Note: Initial concurrent connection is not a KPI to report. This value is configured on the traffic generator and used to perform the Step1: "Test Initialization and Qualification" described under the Section 7.9.4.

Connections per second during ramp up phase: 50% of maximum connections per second measured in benchmarking test TCP/HTTPS

Connections per second (Section 7.6)

Ramp up time (in traffic load profile for "Target concurrent connections"): "Target concurrent connections" / "Maximum connections per second during ramp up phase"

Ramp up time (in traffic load profile for "Initial concurrent connections"): "Initial concurrent connections" / "Maximum connections per second during ramp up phase"

The client MUST perform HTTPS transaction with persistence and each client can open multiple concurrent TCP connections per server endpoint IP.

Each client sends 10 GET requests requesting 1 KByte HTTPS response objects in the same TCP connections (10 transactions/TCP connection) and the delay (think time) between each transaction MUST be X seconds.

X = ("Ramp up time" + "steady state time") /10

The established connections SHOULD remain open until the ramp down phase of the test. During the ramp down phase, all connections SHOULD be successfully closed with FIN.

7.9.3.3. Test Results Validation Criteria

The following criteria are the test results validation criteria. The Test results validation criteria MUST be monitored during the whole sustain phase of the traffic load profile.
a. Number of failed application transactions (receiving any HTTP response code other than 200 OK) MUST be less than 0.001% (1 out of 100,000 transactions) of total attempted transactions.

b. Number of terminated TCP connections due to unexpected TCP RST sent by DUT/SUT MUST be less than 0.001% (1 out of 100,000 connections) of total initiated TCP connections.

c. During the sustain phase, traffic SHOULD be forwarded at a constant rate (considered as a constant rate if any deviation of traffic forwarding rate is less than 5%).

7.9.3.4. Measurement

Average Concurrent TCP Connections MUST be reported for this benchmarking test.

7.9.4. Test Procedures and Expected Results

The test procedure is designed to measure the concurrent TCP connection capacity of the DUT/SUT at the sustaining period of traffic load profile. The test procedure consists of three major steps: Step 1 ensures the DUT/SUT is able to reach the performance value (Initial concurrent connection) and meets the test results validation criteria when it was very minimally utilized. Step 2 determines the DUT/SUT is able to reach the target performance value within the test results validation criteria. Step 3 determines the maximum achievable performance value within the test results validation criteria.

This test procedure MAY be repeated multiple times with different IPv4 and IPv6 traffic distribution.

7.9.4.1. Step 1: Test Initialization and Qualification

Verify the link status of all connected physical interfaces. All interfaces are expected to be in "UP" status.

Configure test equipment to establish "Initial concurrent TCP connections" defined in Section 7.9.3.2. Except ramp up time, the traffic load profile SHOULD be defined as described in Section 4.3.4.

During the sustain phase, the DUT/SUT SHOULD reach the "Initial concurrent TCP connections". The measured KPIs during the sustain phase MUST meet the test results validation criteria "a" and "b" defined in Section 7.9.3.3.
If the KPI metrics do not meet the test results validation criteria, the test procedure MUST NOT be continued to "Step 2".

7.9.4.2.  Step 2: Test Run with Target Objective

Configure test equipment to establish the target objective ("Target concurrent TCP connections"). The test equipment SHOULD follow the traffic load profile definition (except ramp up time) as described in Section 4.3.4.

During the ramp up and sustain phase, the other KPIs such as inspected throughput, TCP connections per second, and application transactions per second MUST NOT reach to the maximum value that the DUT/SUT can support.

The test equipment SHOULD start to measure and record KPIs defined in Section 7.9.3.4. Continue the test until all traffic profile phases are completed.

Within the test results validation criteria, the DUT/SUT is expected to reach the desired value of the target objective in the sustain phase. Follow step 3, if the measured value does not meet the target value or does not fulfill the test results validation criteria.

7.9.4.3.  Step 3: Test Iteration

Determine the achievable concurrent TCP connections within the test results validation criteria.

8.  IANA Considerations

This document makes no specific request of IANA.

The IANA has assigned IPv4 and IPv6 address blocks in [RFC6890] that have been registered for special purposes. The IPv6 address block 2001:2::/48 has been allocated for the purpose of IPv6 Benchmarking [RFC5180] and the IPv4 address block 198.18.0.0/15 has been allocated for the purpose of IPv4 Benchmarking [RFC2544]. This assignment was made to minimize the chance of conflict in case a testing device were to be accidentally connected to part of the Internet.
9. Security Considerations

The primary goal of this document is to provide benchmarking terminology and methodology for next-generation network security devices for use in a laboratory isolated test environment. However, readers should be aware that there is some overlap between performance and security issues. Specifically, the optimal configuration for network security device performance may not be the most secure, and vice-versa. The cipher suites recommended in this document are for test purpose only. The cipher suite recommendation for a real deployment is outside the scope of this document.

10. Contributors

The following individuals contributed significantly to the creation of this document:

Alex Samonte, Amritam Putatunda, Aria Eslamboichizadeh, Chao Guo, Chris Brown, Cory Ford, David DeSanto, Jurrie Van Den Breekel, Michelle Rhines, Mike Jack, Ryan Liles, Samaresh Nair, Stephen Goudreault, Tim Carlin, and Tim Otto.

11. Acknowledgements

The authors wish to acknowledge the members of NetSecOPEN for their participation in the creation of this document. Additionally, the following members need to be acknowledged:


12. References

12.1. Normative References


12.2. Informative References
Appendix A.  Test Methodology - Security Effectiveness Evaluation
A.1. Test Objective

This test methodology verifies the DUT/SUT is able to detect, prevent, and report the vulnerabilities.

In this test, background test traffic will be generated to utilize the DUT/SUT. In parallel, the CVEs will be sent to the DUT/SUT as encrypted and as well as clear text payload formats using a traffic generator. The selection of the CVEs is described in Section 4.2.1.

The following KPIs are measured in this test:
* Number of blocked CVEs
* Number of bypassed (nonblocked) CVEs
* Background traffic performance (verify if the background traffic is impacted while sending CVE toward DUT/SUT)
* Accuracy of DUT/SUT statistics in term of vulnerabilities reporting

A.2. Testbed Setup

The same testbed MUST be used for security effectiveness test and as well as for benchmarking test cases defined in Section 7.

A.3. Test Parameters

In this section, the benchmarking test specific parameters SHOULD be defined.

A.3.1. DUT/SUT Configuration Parameters

DUT/SUT configuration parameters MUST conform to the requirements defined in Section 4.2. The same DUT configuration MUST be used for Security effectiveness test and as well as for benchmarking test cases defined in Section 7. The DUT/SUT MUST be configured in inline mode and all detected attack traffic MUST be dropped and the session SHOULD be reset.

A.3.2. Test Equipment Configuration Parameters

Test equipment configuration parameters MUST conform to the requirements defined in Section 4.3. The same client and server IP ranges MUST be configured as used in the benchmarking test cases. In addition, the following parameters MUST be documented for this benchmarking test:
* Background Traffic: 45% of maximum HTTP throughput and 45% of Maximum HTTPS throughput supported by the DUT/SUT (measured with object size 64 KByte in the benchmarking tests "HTTP(S) Throughput" defined in Section 7.3 and Section 7.7).

* RECOMMENDED CVE traffic transmission Rate: 10 CVEs per second

* It is RECOMMENDED to generate each CVE multiple times (sequentially) at 10 CVEs per second

* Ciphers and keys for the encrypted CVE traffic MUST use the same cipher configured for HTTPS traffic related benchmarking tests (Section 7.6 - Section 7.9)

A.4. Test Results Validation Criteria

The following criteria are the test results validation criteria. The test results validation criteria MUST be monitored during the whole test duration.

a. Number of failed application transaction in the background traffic MUST be less than 0.01% of attempted transactions.

b. Number of terminated TCP connections of the background traffic (due to unexpected TCP RST sent by DUT/SUT) MUST be less than 0.01% of total initiated TCP connections in the background traffic.

c. During the sustain phase, traffic SHOULD be forwarded at a constant rate (considered as a constant rate if any deviation of traffic forwarding rate is less than 5%).

d. False positive MUST NOT occur in the background traffic.

A.5. Measurement

Following KPI metrics MUST be reported for this test scenario:

Mandatory KPIs:

* Blocked CVEs: It SHOULD be represented in the following ways:
  - Number of blocked CVEs out of total CVEs
  - Percentage of blocked CVEs

* Unblocked CVEs: It SHOULD be represented in the following ways:
- Number of unblocked CVEs out of total CVEs
- Percentage of unblocked CVEs

* Background traffic behavior: It SHOULD be represented one of the followings ways:

- No impact: Considered as "no impact'" if any deviation of traffic forwarding rate is less than or equal to 5 % (constant rate)
- Minor impact: Considered as "minor impact" if any deviation of traffic forwarding rate is greater than 5% and less than or equal to10% (i.e. small spikes)
- Heavily impacted: Considered as "Heavily impacted" if any deviation of traffic forwarding rate is greater than 10% (i.e. large spikes) or reduced the background HTTP(S) throughput greater than 10%

* DUT/SUT reporting accuracy: DUT/SUT MUST report all detected vulnerabilities.

Optional KPIs:

* List of unblocked CVEs

A.6. Test Procedures and Expected Results

The test procedure is designed to measure the security effectiveness of the DUT/SUT at the sustaining period of the traffic load profile. The test procedure consists of two major steps. This test procedure MAY be repeated multiple times with different IPv4 and IPv6 traffic distribution.

A.6.1. Step 1: Background Traffic

Generate background traffic at the transmission rate defined in Appendix A.3.2.

The DUT/SUT MUST reach the target objective (HTTP(S) throughput) in sustain phase. The measured KPIs during the sustain phase MUST meet all the test results validation criteria defined in Appendix A.4.

If the KPI metrics do not meet the acceptance criteria, the test procedure MUST NOT be continued to "Step 2".

A.6.2. Step 2: CVE Emulation

While generating background traffic (in sustain phase), send the CVE traffic as defined in the parameter section.

The test equipment SHOULD start to measure and record all specified KPIs. Continue the test until all CVEs are sent.

The measured KPIs MUST meet all the test results validation criteria defined in Appendix A.4.

In addition, the DUT/SUT SHOULD report the vulnerabilities correctly.

Appendix B. DUT/SUT Classification

This document aims to classify the DUT/SUT in four different categories based on its maximum supported firewall throughput performance number defined in the vendor datasheet. This classification MAY help user to determine specific configuration scale (e.g., number of ACL entries), traffic profiles, and attack traffic profiles, scaling those proportionally to DUT/SUT sizing category.

The four different categories are Extra Small (XS), Small (S), Medium (M), and Large (L). The RECOMMENDED throughput values for the following categories are:

Extra Small (XS) - Supported throughput less than or equal to 1Gbit/s

Small (S) - Supported throughput greater than 1Gbit/s and less than or equal to 5Gbit/s

Medium (M) - Supported throughput greater than 5Gbit/s and less than or equal to 10Gbit/s

Large (L) - Supported throughput greater than 10Gbit/s

Authors’ Addresses

Balamuhunthan Balarajah
Berlin
Germany

Email: bm.balarajah@gmail.com
Problems and Requirements of Evaluation Methodology for Integrated Space and Terrestrial Networks
draft-lai-bmwg-istn-methodology-01

Abstract

With the rapid evolution of the aerospace industry, many "NewSpace" upstarts are actively deploying their mega-constellations in low earth orbits (LEO) and building integrated space and terrestrial networks (ISTN), promising to provide pervasive, low-latency, and high-throughput Internet service globally. Due to the high manufacturing, launching, and updating cost of LEO mega-constellations, it is expected that ISTNs can be well designed and evaluated before the launch of satellites. However, the progress of designing, assessing, and understanding new network functionalities and protocols for futuristic ISTNs faces a substantial obstacle: lack of standardized evaluation methodology with acceptable realism (e.g. can involve the unique dynamic behaviors of ISTNs), flexibility, and cost. This memo first reviews the unique characteristics of LEO mega-constellations. Further, it analyzes the limitation of existing evaluation and analysis methodologies under ISTN environments. Finally, it outlines the key requirements of future evaluation methodology tailored for ISTNs.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at https://datatracker.ietf.org/drafts/current/.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."
Table of Contents

1. Introduction .................................................. 3
2. Notation and Terminology ..................................... 4
3. Quick Primer for Integrated Space and Terrestrial Networks . 5
   3.1. Mega-constellation ......................................... 5
   3.2. Topological Dynamics ....................................... 6
   3.3. Limited Resources .......................................... 7
   3.4. Long Manufacturing and Deployment Duration ............... 8
4. Problem Statement: We Need the Right Evaluation Methodology .................................................. 8
   4.1. Live networks and platforms .............................. 9
   4.2. Network Simulators ......................................... 10
   4.3. Network Emulators ......................................... 11
   4.4. Summary .................................................... 12
5. Requirements: New Evaluation Methodology Tailored for ISTNs .................................................. 12
   5.1. Realism ..................................................... 13
   5.2. Flexibility ................................................ 13
   5.3. Low-cost and Easy-to-use .................................. 13
   5.4. Cross-domain Dataset Support ................................ 13
6. Conclusion .................................................... 14
7. Acknowledgements ............................................... 14
8. IANA Considerations ............................................ 14
9. Security Considerations ........................................ 14
10. References ..................................................... 14
    10.1. Normative References ..................................... 14
    10.2. Informative References ................................... 15
Authors’ Addresses ................................................ 17
1. Introduction

Integrated Space and Terrestrial Networks (ISTN), combining diverse spacecrafts and ground infrastructures, are extending the frontier of today’s terrestrial network, promising to provide low-latency, high-bandwidth Internet access with broader coverage globally.

Low earth orbit (LEO) satellites are the key building block for constructing ISTNs. Recently, we have witnessed a renaissance in the space industry, stimulating an exponential increase in constructing mega-constellations. As compared with their predecessor, cutting-edge satellites can be equipped with high-resolution sensors, space-grade multi-core processors, high-data-rate communication links, and multifunctional space software.

While ISTNs hold great promise, to completely unleash the network potential of emerging ISTN, it still needs to address a series of new technical issues. The unique characteristics of LEO satellites (e.g., high-dynamics), not only impose new challenges at various layers of the ISTN networking stack but also open the door to many new technical problems. With many unexplored problems facing the "NewSpace" industry, it is thus foreseen that in the near future, there will be a surge of new efforts (e.g. topology, addressing, routing, transport, etc.) to rethink and reshape the networking stack in ISTNs. In addition, the cost/timeline of manufacturing, launching, operating, and updating satellite constellations is typically much higher/longer than that in traditional terrestrial networks. Therefore, it is expected that new network functionalities and protocols can be well evaluated before they are launched and deployed in realistic satellite constellations.

However, the network community lacks the proper analysis tools and evaluation methodologies that can mimic the unique dynamic behavior to analyze many of the ISTN challenges that have been highlighted by prior works. At high level, existing evaluation methodologies in the network community can typically be grouped into three major categories: live networks or platforms, simulation, and emulation. However, the feasibility and flexibility of live satellite networks are technically and economically limited. The abstraction level of network simulation could be too high to capture low-level system effects. Existing network emulators fail to characterize the high dynamicity of LEO satellites and thus cannot accomplish an environment with acceptable fidelity. The community hence needs a reasonable and standardized evaluation methodology to build proper experimental environments which can mimic the behavior of ISTNs, supporting the community to deeply understand the problems, and to evaluate new functionalities and protocols (e.g. for topology, addressing, routing, transport, etc.) for ISTNs, before the mega-

constellation is completely deployed. In this memo, we first review the unique characteristics of emerging LEO mega-constellations and the key challenges of integrating satellites and terrestrial Internet. Further, we analyze the limitation of existing network analysis tools and evaluation methodologies in ISTNs. Finally, we outline key requirements of evaluation methodologies tailored for ISTNs.

2. Notation and Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

This document uses the following acronyms and terminologies:

Mega-constellation: A group of satellites working as a system.
LEO: Low Earth Orbit with an altitude no more than 2000 km.
MEO: Medium Earth Orbit with an altitude from 2000 km to 35786 km.
GEO: Geostationary Earth Orbit with an altitude of 35786 km.
NGSO: Non-Geostationary Orbit
LSN: LEO Satellite Networks
ISTN: Integrated Satellite and Terrestrial Network
ISL: Inter-satellite Links
EO: Earth Observation
GS: Ground Station
AS: Autonomous System
EOS: Earth Observation Satellite
BGP: Border Gateway Protocol [RFC4271]
OSPF: Open Shortest Path First [RFC2328]
VM: Virtual Machine
3. Quick Primer for Integrated Space and Terrestrial Networks

Emerging mega-constellations with inter-satellite links (ISLs) can build a satellite network in outer space, and further be integrated with terrestrial ground infrastructures to construct an integrated space and terrestrial network (ISTN).

3.1. Mega-constellation

A constellation is a group of satellites working as a system to give a coverage of the earth surface, among which satellites are positioned in fixed orbital planes with regular trajectories. LEO and MEO satellites often belong to a constellation, because a single satellite only covers a small area with high angular velocity. Thus, continuous coverage over an area could be maintained by the relay within a constellation, as compared with GEO satellites that only provides a permanent coverage over a target area. Walker Delta constellation is the most common formation for constellations. It is defined as a bunch of circular orbits with a fixed inclination, satellite number, number of equally spaced planes and the relative spacing between satellites in adjacent planes. The famous Ballard rosette constellation is another name of Walker Delta constellation, where it uses a different notation. Near-polar Walker Star is one of this kind, initially used by Iridium [Iridium]. Constellations with a higher inclination give the polar regions more chances to get accessed. The well-known emerging commercial constellations are Starlink [Starlink-Fcc], Kuiper [Kuiper-Fcc] and Telesat [Telesat-Fcc], as shown in Table 1 below. And all of them contain more than one shell.
<table>
<thead>
<tr>
<th>Name and Shell</th>
<th>Altitude (km)</th>
<th>Inclination (degree)</th>
<th># of orbits</th>
<th># of satellites per orbit</th>
</tr>
</thead>
<tbody>
<tr>
<td>Starlink S1</td>
<td>550</td>
<td>53</td>
<td>72</td>
<td>22</td>
</tr>
<tr>
<td>Starlink S2</td>
<td>1110</td>
<td>53.8</td>
<td>32</td>
<td>50</td>
</tr>
<tr>
<td>Starlink S3</td>
<td>1130</td>
<td>74</td>
<td>8</td>
<td>50</td>
</tr>
<tr>
<td>Starlink S4</td>
<td>1275</td>
<td>81</td>
<td>5</td>
<td>75</td>
</tr>
<tr>
<td>Starlink S5</td>
<td>1325</td>
<td>70</td>
<td>6</td>
<td>75</td>
</tr>
<tr>
<td>Kuiper K1</td>
<td>630</td>
<td>51.9</td>
<td>34</td>
<td>34</td>
</tr>
<tr>
<td>Kuiper K2</td>
<td>610</td>
<td>42</td>
<td>36</td>
<td>36</td>
</tr>
<tr>
<td>Kuiper K3</td>
<td>590</td>
<td>33</td>
<td>28</td>
<td>28</td>
</tr>
<tr>
<td>Telesat T1</td>
<td>1015</td>
<td>98.98</td>
<td>27</td>
<td>13</td>
</tr>
<tr>
<td>Telesat T2</td>
<td>1325</td>
<td>50.88</td>
<td>40</td>
<td>33</td>
</tr>
</tbody>
</table>

Table 1: Mega-constellation information.

3.2. Topological Dynamics

Unlike geostationary satellite networks or terrestrial core infrastructure that keep a stable topology, LEO satellite networks suffer from high topological dynamics, since LEO satellites move fast, causing short-lived coverage for fixed terrestrial users. For example, considering the first shell of Starlink Phase-I, a fixed user sees each satellite for only up to 3 minutes in one pass, after which the satellite moves away from the user’s perspective. Table 2 shows the medium space-ground link churn intervals [link-churn-interval] between existing GS and constellations. If
each GS only uses one antenna to connect the satellite with the shortest distance, the medium interval is no more than one minute. This kind of high dynamic motion incurs frequent link changes between LEO satellites and GS or users, thus causing frequent topology changes. Moreover, inter-satellites visibility may also change if LEO satellites move in different directions or in different shells, resulting in connectivity change of ISLs.

Such high LEO dynamics can impose significant challenges in the networking stack of ISTNs. The high dynamics make the logical network and mega-constellations and physical ISTN inconsistent. One big challenge is how to overcome the routing oscillation properly in the high dynamic ISTN environment. Frequent satellite-GS link changes make the inter-connectivity of space and ground segments in ISTNs unstable. Thus, the routing have to be re-calculated every time the link changes. In addition, the topological dynamics also result in RTT fluctuations in end-to-end paths, involving new challenges for congestion control in ISTNs, as a RTT variation observed by end-host might not indicate congestions.

<table>
<thead>
<tr>
<th>Name</th>
<th>Interval (s)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Starlink</td>
<td>3.0901</td>
</tr>
<tr>
<td>Kuiper</td>
<td>5.0562</td>
</tr>
<tr>
<td>OneWeb</td>
<td>10.6824</td>
</tr>
<tr>
<td>Telesat</td>
<td>45.5696</td>
</tr>
</tbody>
</table>

Table 2: Space-ground link churn interval.

3.3. Limited Resources

Space resources (e.g. CPU, energy) on satellites are limited, as compared with terrestrial network. Since resource-constrained satellites such as nanosatellites are only able to carry certain sensing or transferring missions, energy-consuming or complex tasks may not be achievable in these satellites. Such complicated tasks include on-board target identification and instant and continuous disaster monitoring.

For example, the CPU frequency of current spaceborne processors (e.g. RAD5545 [RAD5545], RAD750 [RAD750]) is only up to 466MHz per core. More recently, some low energy-consuming commodity processors are
used in space to complete certain remote sensing missions under a limited CPU capacity. [raspberry-pi] With a constrained computation ability and limited storage and energy, satellite functions and lifetime are greatly repressed.

3.4. Long Manufacturing and Deployment Duration

Different from terrestrial network infrastructures, the timeline of manufacturing and deploying satellite networks could be much longer due to the high cost and complex process during the development and launch period. Satellites, as well as the orbit and spectrum they used, have to be regulated, and launches have to be carefully scheduled (e.g. to avoid the impact of poor weather conditions). In addition, the maintenance and update cost of a satellite network is also typically much higher than that in a terrestrial network.

For example, a review of 24 Air Force and Navy space vehicle (SV) development programs found that on average it took about 7.5 years from contract start to launch a government satellite. [Development-Timeline] Commercial satellite programs typically take 2 to 3 years from contract start to launch. [Production-Cycles] SpaceX’s Starlink constellation plan to launch about 42,000 satellites to construct a mega-constellation in outer space. On 15 October 2019, the United States Federal Communications Commission (FCC) submitted filings to the International Telecommunication Union (ITU) on SpaceX’s behalf to arrange spectrum for 30,000 additional Starlink satellites to supplement the 12,000 Starlink satellites already approved by the FCC. As of the date of April 2022, SpaceX has launched about 2,100 Starlink satellites, which is about 5% of the ultimate constellation plan consisting of 42,000 satellites. Foreseeably, it may take many years to complete the entire constellation deployment. Even the first phase of Starlink which consists of about 4400 satellites is not expected to be completed until 2024.

4. Problem Statement: We Need the Right Evaluation Methodology

The unique characteristics of LEO mega-constellations involve new challenges on various layers of the networking stack of ISTNs. On one hand, it is foreseen that in the near future, there will be a surge of new network functionalities and protocols designed or optimized for ISTNs. On the other hand, because the cost/timeline of manufacturing, launching, operating, and updating satellite constellations is typically much higher/longer than that in traditional terrestrial networks, it is expected that those new network functionalities and protocols tailored for ISTNs should be well evaluated before they are launched and deployed in realistic satellite constellations.
Existing methodologies for testing, assessing, and understanding a network functionality or protocol can typically be classified into three categories: (1) live networks; (2) network simulators; and (3) network emulators. The subsections discuss these three categories of network evaluation methodologies, along with their using deficiencies and possible remedies respectively.

4.1. Live networks and platforms

Representative platforms such as Emulab [Emulab] and Sparta [Sparta] are successful pioneers that build a large-scale experimental network environment. These test environments are originally designed to provide special and exclusive test services for affiliated universities, scientific research institutions or Internet business companies. And for the resource competition, each independent experiment needs to completely monopolize a part of the test bed, so the researcher cannot deploy the experiment until being allocated with enough nodes. PlanetLab [PlanetLab] is truly global ground testbed prototype. Started from 2003, it consists of 1353 nodes at 717 sites spanning 48 countries. Together the nodes form a global network system to support new design of network services.

The live platforms described above were initially proposed for terrestrial networks and they are developed and repaired at the same time. The key limitation of them in an ISTN environment is that they are designed for terrestrial network experiments, and do not incorporate the realistic characteristic of LEO mega-constellations to support experiments and evaluations in ISTNs.

We may search for help from live satellites, but still there is only limited help. It seems that with the help of live ISTN, researchers are capable to assess, verify and evaluate their ideas and thoughts. Live ISTN can give a real constellation-consistency and stack-consistency testing environment. However, current satellites only provide users a bent-pipe service, which is purely relaying the transmission messages, such as the current deployment of Starlink [Starlink]. The construction is far from a comprehensive ISTN, so the research scope is limited. Even if there is a live ISTN, it lacks flexibility, owning to the inconvenient control over satellites. Besides, the access to satellites is also limited.

Therefore, live networks or platforms for terrestrial networks can give us a large-scale experimental environment but they lack the support for ISTN characteristics. On the other hand, live ISTN is able to guarantee a real space environment, but it is not that affordable and flexible.
4.2. Network Simulators

Simulators are tools that enable researchers to reproduce their testing experiments by simulating a real-world process or system over time. Simulators work by using discrete event simulation to calculate the interactive states among all the network entities, ranging from switches, routers, nodes, access points, links and so on. While working fast and efficiently, the fidelity is only brought by the state variable changes at discrete points.

Such tools like Systems Tool Kit (STK) [Systems-Tool-Kit] and General Mission Analysis Tool (GMAT) [General-Mission-Analysis-Tool] are good for orbit analysis. STK is a powerful tool to help researchers to model the behavior of mission entities in aerospace, telecommunications and so forth. It also provides visualization and analysis functions. GMAT is a similar tool for space trajectory optimization and mission modeling. Nevertheless, these tools do not support networking simulations such as topology and protocol simulations. ns-3 [ns-3] goes a step further with support for Internet simulation, but on the contrary, it was not designed for ISTN and lacks the support for high-dynamics of ISTN. StarPerf [StarPerf] is a simulator that helps researchers to study network performance under a range of constellation conditions. But still, it lacks the ability to support interactive network traffic simulation and system codes in the systems.

Overall, while flexible and low-cost, the realism of simulators is not content enough, because they are difficult to describe the low-level characteristics. In other words, simulators are being too object-oriented to involve additional overhead in the actual execution of programs. Besides, when accessing the network performance, a number of recent emerging algorithms for congestion control, reliable transmission or even protocols are not supported, for example ns-3 [ns-3] only supports basic congestion control like Reno [RFC6582] and so forth, so the need to work with some new algorithms cannot be satisfied and the research to discover new mechanisms, such as new routing algorithms and re-transmission schemes, is extensively prohibited. Another problem of simulators, such as ns-3 [ns-3], is that it difficult to trace or understand the previous codes, without appropriate documentations. Simulators usually face the additional compatibility problem, which means they are not portable with other systems, or they do not support kernel codes. Since there are multiple simulators developed by different group of users, sometimes users are required to be familiar with the writing language, scripting style and modelling technique.
4.3. Network Emulators

Emulators are another kind of paradigm for network evaluation over a virtual network. The difference between a simulator and an emulator is that emulators leverage VM or containers to keep the realism which is close to actual performances. Therefore, in emulators, virtual nodes, virtual network links, virtual models of traffic, and protocols are all applied. Emulators are capable to run real kernel and application code. Thus, emulators not only support diverse topology design, but also protocol emulation in a synthetic network environment. They emulate the network behavior in a more real way. Mininet [Mininet] is commonly regarded as the most illustrious emulator for networking with its strong ability to support experiments with Software-Defined Networking (SDN) [Software-defined-networking] systems. EstiNet [EstiNet] is another emulator that supports evaluating and testing the performances of software-defined networks. Based on containers, they can emulate real TCP/IP protocol stack in the Linux kernel. However, existing emulation tools lack the ability to construct the dynamic links and orbits in ISTN like simulators. Thus, more problems could happen in higher-level protocols such as routing protocols (e.g. OSPF and BGP). Besides, since emulators run containers or virtual machines which occupy more software overhead, as compared with simulators, it will be hard to emulate the large-scale mega-constellations.

To conclude, emulators are relatively good methodologies for network experiments, but emulators still have limitations when using them for ISTN research. While keeping a moderate realism by using VM or containers for entity emulation and flexibility, emulators still lack the supports for ISTN characteristics, such as frequent link changes, satellite network topology uncertainty, and so on. More specifically, current emulators only support fixed network topology emulation. It is not flexible to emulate the time-varying link packet loss, bandwidth, and other traits. A possible way is to frequently replace the link with a new one from time to time sequentially for the entire ISTN. However, it is far from the real situation. Besides, VM or containers are able to deploy a range of network nodes in a physical server, but the actual CPU, memory and other resources should not be shared in reality for each satellite. In addition, it is still difficult to emulate thousands or ten thousand of satellites for ISTN even with VM or containers, subject to hardware limitations. For flexibility, some emulators do not support a good network animator tool. Especially in ISTN emulation, GUI is important for users to observe and analyze orbit trajectories and real time satellite positions.
4.4. Summary

In this section, we explain the necessity of an evaluation methodology specifically for ISTNs. Then we demonstrate the problems with existing methodologies related to ISTNs. The performance comparison result is shown in Table 3. Above all, ISTNs should be designed first and then launched. Live satellites enable good realism but they lack flexibility and require very high cost as well as a very long deployment period. Other testing tools such as simulators and emulators are either functional for merely aerospace analysis or simply terrestrial networks. None of the existing methodologies guarantees a practical and user-friendly methodology while keeping the evaluation environment realism with low costs.

<table>
<thead>
<tr>
<th>Platform/Tool</th>
<th>Realism</th>
<th>Flexibility</th>
<th>Cost</th>
<th>Cross-domain Dataset Support</th>
</tr>
</thead>
<tbody>
<tr>
<td>Live satellite network</td>
<td>Y</td>
<td>N</td>
<td>High</td>
<td>Y</td>
</tr>
<tr>
<td>ns-3 [ns-3]</td>
<td>N</td>
<td>Y</td>
<td>Low</td>
<td>N</td>
</tr>
<tr>
<td>Hypatia [Hypatia]</td>
<td>N</td>
<td>Y</td>
<td>Low</td>
<td>N</td>
</tr>
<tr>
<td>StarPerf [StarPerf]</td>
<td>N</td>
<td>Y</td>
<td>Low</td>
<td>N</td>
</tr>
<tr>
<td>Mininet [Mininet]</td>
<td>N</td>
<td>Y</td>
<td>Low</td>
<td>N</td>
</tr>
</tbody>
</table>

Table 3: Existing platforms/tools for network analysis and evaluation. (Y for yes/N for no)

5. Requirements: New Evaluation Methodology Tailored for ISTNs

A proper evaluation methodology tailored for ISTNs is expected to help developers, researchers, engineers to explore various design-space of the networking stack of ISTNs in a technically and economically feasible manner. Based on the comparative analysis results in the prior section, we sum up the following requirements for the new evaluation methodology in ISTNs.
5.1. Realism

The first requirement is realism. Realism represents the testing authenticity and fidelity. The evaluation methodology is expected to keep the actual characteristics of mega-constellations. In other words, the orbit-level information including the latitude, longitude, and height of each satellite in any given time and the same information for GS and elevation angles of antennas of each GS. Note that the constellation information also determines the visibility, links and even topology of ISTNs. Since the mega-constellations are unstable, how the temporal satellite locations, visibility, link propagation delays and so on should also be considered carefully. In addition, it requires the network nodes to communicate and negotiate their messages following the actual protocol process. For example, when doing a test for OSPF in an ISTN, we would like the nodes to send Hello packets, Link-State-Request (LSR) packets, Link-State-Update (LSU) packets and so on. A real network stack is preferred to provide researchers an opportunity to see the performance of different protocols in ISTNs.

5.2. Flexibility

Another requirement is flexibility and feasibility. The testing methodology should be technically easy to use and easy to learn. Without extra modifications or process, the methodology should help researchers learn and use it without much effort and can evaluate their ideas as they wish, which means it should support flexible, controllable environments for researchers.

5.3. Low-cost and Easy-to-use

Meanwhile, the evaluation methodology is expected to be low-cost. A well-acceptable methodology should be economically feasible for users to create an evaluation environment. Researchers do not want to conduct their tests all in live ISTN, which is over-cumbersome and unaffordable, let alone launching their own spacecraft. Even if there are a number of orbiting satellites, whether users can easily gain access to satellites is also a problem.

5.4. Cross-domain Dataset Support

The evaluation methodology is expected to be driven by realistic datasets from multi-dimensions to support its realism. Multi-dimension refers to multi-disciplinary research on ISTNs. Since a standard ISTN evaluation methodology not only contains high-level benchmarks from topology, routing to transmission, but also considers the low-level traits such as wireless link conditions, weather conditions and Earth rotations. To be more concrete, the former one
requires knowledge in networks while the latter one relies more on aerospace. Hence, to build a high-fidelity methodology, we need community efforts both from networks and aerospace. On the other hand, an authentic dataset is an indispensable element for data driven testing methodology. Actual data is the first step to obtain a realistic emulation. with characteristics of a real ISTN. Thus, the dataset is a collection of messages for testing, in which geographical mega-constellation information (orbit number, satellite number, height), orbital information (orbit inclination angle and link strategies), weather information as well as ground station information (positions, antenna angle and so forth) are involved.

6. Conclusion

To conclude, the emergence of mega-constellations brings us new opportunities for the development of ISTN that extends the Internet to the space era. Combined with terrestrial networks, ISTN is expected to supply pervasive, low-latency and high-speed services to users globally, which greatly enhances the current Internet. At the same time, the unique characteristics (e.g. high-dynamics) of ISTN impose challenges in topology, routing, transportation, application, and security. However, we simply believe addressing the challenges also gives us open opportunities for future research by our community-driven effort. To accelerate the research speed and to help make testing more feasible, new methodologies that satisfy user requirements should be proposed. To this extent, this draft reviews the limitation of existing network analysis tools in ISTNs, considering the unique characteristics of emerging LSNs and the key challenges. This draft further analyzes the limitation of existing evaluation methodologies in ISTN environments. Finally, this draft outlines key requirements of evaluation methodologies tailored for future ISTNs.

7. Acknowledgements

8. IANA Considerations

This memo includes no request to IANA.

9. Security Considerations

This entire draft discusses security considerations from different perspectives of ISTN in Section 3.

10. References

10.1. Normative References
10.2. Informative References

[Development-Timeline]


[General-Mission-Analysis-Tool]


[Kuiper-Fcc]

[link-churn-interval]


[Telesat-Fcc]

Authors’ Addresses

Zeqi Lai
Tsinghua University
30 ShuangQing Ave
Beijing
100089
China
Email: zeqilai@tsinghua.edu.cn

Hewu Li
Tsinghua University
30 ShuangQing Ave
Beijing
100084
China
Email: lihewu@cernet.edu.cn

Yangtao Deng
Tsinghua University
30 ShuangQing Ave
Beijing
100084
China
Email: dengyt21@mails.tsinghua.edu.cn

Qian Wu
Tsinghua University
30 ShuangQing Ave
Beijing
100084
China
Email: wuqian@cernet.edu.cn

Jun Liu
Tsinghua University
30 ShuangQing Ave
Beijing
100084
China
Benchmarking Methodology Working Group                         G. Lencse
Internet-Draft                               Szechenyi Istvan University
Intended status: Informational                                  K. Shima
Expires: 5 September 2022                       IIJ Innovation Institute
4 March 2022

Benchmarking Methodology for Stateful NATxy Gateways using RFC 4814
Pseudorandom Port Numbers
draft-lencse-bmwg-benchmarking-stateful-03

Abstract

RFC 2544 has defined a benchmarking methodology for network interconnect devices. RFC 5180 addressed IPv6 specificities and it also provided a technology update, but excluded IPv6 transition technologies. RFC 8219 addressed IPv6 transition technologies, including stateful NAT64. However, none of them discussed how to apply RFC 4814 pseudorandom port numbers to any stateful NAT (NAT44, NAT64, NAT66) technologies. We discuss why using pseudorandom port numbers with stateful NAT gateways is a hard problem and recommend a solution.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at https://datatracker.ietf.org/drafts/current/.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on 5 September 2022.

Copyright Notice

Copyright (c) 2022 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust’s Legal Provisions Relating to IETF Documents (https://trustee.ietf.org/license-info) in effect on the date of publication of this document.
1. Introduction

[RFC2544] has defined a comprehensive benchmarking methodology for network interconnect devices, which is still in use. It was mainly IP version independent, but it used IPv4 in its examples. [RFC5180] addressed IPv6 specificities and also added technology updates, but declared IPv6 transition technologies out of its scope. [RFC8219] addressed the IPv6 transition technologies, including stateful NAT64. It has reused several benchmarking procedures from [RFC2544] (e.g. throughput, frame loss rate), it has redefined the latency measurement, and added further ones, e.g. the PDV (packet delay
variation) measurement.

However, none of them discussed, how to apply [RFC4814] pseudorandom port numbers, when benchmarking stateful NATxy (NAT44, NAT64, NAT66) gateways. We are not aware of any other RFCs that address this question.

First, we discuss why using pseudorandom port numbers with stateful NATxy gateways is a hard problem.

Then we recommend a solution.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

2. Pseudorandom Port Numbers and Stateful Translation

In its appendix, [RFC2544] has defined a frame format for test frames including specific source and destination port numbers. [RFC4814] recommends to use pseudorandom and uniformly distributed values for both source and destination port numbers. However, stateful NATxy (NAT44, NAT64, NAT66) solutions use the port numbers to identify connections. The usage of pseudorandom port numbers causes different problems depending on the direction.

* As for the private to public direction, pseudorandom source and destination port numbers could be used, however, this approach would be a denial of service attack against the stateful NATxy gateway, because it would exhaust its connection tracking table. To that end, let us see some calculations using the recommendations of RFC 4814:

  - The recommended source port range is: 1024-65535, thus its size is: 64512.
  - The recommended destination port range is: 1-49151, thus its size is: 49151.
  - The number of source and destination port number combinations is: 3,170,829,312.
We note that section 12 of [RFC2544] also requires testing with 256 destination networks, which further increases the number of connection tracking table entries.

* As for the public to private direction, the stateful DUT (Device Under Test) would drop any packets that do not belong to an existing connection, therefore, the direct usage of pseudorandom port numbers from the above-mentioned ranges is not feasible.

3. Test Setup and Terminology

Our methodology works with any IP version. We use IPv4 in the Test Setup shown in Figure 1 to facilitate its easy understanding based on the well-known stateful NAT44 (also called NAPT: Network Address and Port Translation) solution.

![Test Setup Diagram]

Figure 1: Test Setup for benchmarking stateful NATxy gateways

As for transport layer protocol, [RFC2544] recommended testing with UDP, and it was kept also in [RFC8219]. For the general recommendation, we also keep UDP, thus the port numbers in the following text are to be understood as UDP port numbers. We discuss the limitation of this approach in Section 6.

We define the most important elements of our proposed benchmarking system as follows.

* Connection tracking table: The stateful NATxy gateway uses a connection tracking table to be able to perform the stateful translation in the public to private direction. Its size, policy and content are unknown for the Tester.
* Four tuple: The four numbers that identify a connection are source IP address, source port number, destination IP address, destination port number.

* State table: The Responder of the Tester extracts the four tuple from each received test frame and stores it in its state table. Recommendation is given for writing and reading order of the state table in Section 4.8.

* Initiator: The port of the Tester that may initiate a connection through the stateful DUT in the private to public direction. Theoretically, it can use any source and destination port numbers from the ranges recommended by [RFC4814]: if the used four tuple does not belong to an existing connection, the DUT will register a new connection into its connection tracking table.

* Responder: The port of the Tester that may not initiate a connection through the stateful DUT in the public to private direction. It may send only frames that belong to an existing connection. To that end, it uses four tuples that have been previously extracted from the received test frames and stored in its state table.

* Preliminary test phase: Test frames are sent only by the Initiator to the Responder through the DUT to fill both the connection tracking table of the DUT and the state table of the Responder. This is a newly introduced operation phase for stateful NATxy benchmarking. The necessity of this phase is explained in Section 4.2.

* Real test phase: The actual test (e.g. throughput, latency, etc.) is performed in this phase after the completion of the preliminary test phase. Test frames are sent as required (e.g. bidirectional test or unidirectional test in any of the two directions).

4. Recommended Benchmarking Method

4.1. Restricted Port Number Ranges

The Initiator SHOULD use restricted ranges for source and destination port numbers to avoid the denial of service attack like event against the connection tracking table of the DUT described in Section 2. The size of the source port number range SHOULD be larger (e.g. in the order of a few times ten thousand), whereas the size of the destination port number range SHOULD be smaller (may vary from a few to several hundreds or thousands as needed). The rationale is that source and destination port numbers that can be observed in the Internet traffic are not symmetrical. Whereas source port numbers

Lencse & Shima Expires 5 September 2022 [Page 5]
may be random, there are a few very popular destination port numbers (e.g. 443, 80, etc., see [IIR2020]) and others hardly occur. And we have found that their role is also asymmetric in the Linux kernel routing hash function [LEN2020].

The product of the sizes of the two ranges can be used as a parameter. The performance of the stateful NATxy gateway MAY be examined as a function of this parameter.

4.2. Preliminary Test Phase

The preliminary phase serves two purposes:

1. The connection tracking table of the DUT is filled. It is important, because its maximum connection establishment rate may be lower than its maximum frame forwarding rate (that is throughput).

2. The state table of the Responder is filled with valid four tuples. It is a precondition for the Responder to be able to transmit frames that belong to connections exist in the connection tracking table of the DUT.

Whereas the above two things are always necessary before the real test phase, the preliminary phase can be used without the real test phase. It is done so, when the maximum connection establishment rate is measured (as described in Section 4.5).

A preliminary test phase MUST be performed before all tests performed in the real test phase. In this phase, the following things happen:

1. The Initiator sends test frames to the Responder through the DUT at a specific frame rate.

2. The DUT performs the stateful translation of the test frames and it also stores the new combinations in its connection tracking table.

3. The Responder receives the translated test frames and updates its state table with the received four tuples. The responder transmits no test frames during the preliminary phase.

When the preliminary test phase is performed in preparation to the real test phase, the applied frame rate and the duration of the preliminary phase SHOULD be carefully selected so that:

* The applied frame rate be safely lower than the maximum connection establishment rate.
* Enough four tuples be stored in the state table of the Responder so that it can generate frames with the proper distribution of the four tuples.

Please refer to Section 4.4 for further conditions regarding timeout and port number combinations.

4.3. Consideration of the Cases of Stateful Operation

We consider the most important Events that may happen during the operation of a stateful NATxy gateway, and the Actions of the gateway as follows.

1. EVENT: A packet not belonging to an existing connection arrives in the private to public direction. ACTION: A new connection is registered into the connection tracking table and the packet is translated and forwarded.

2. EVENT: A packet not belonging to an existing connection arrives in the public to private direction. ACTION: The packet is discarded.

3. EVENT: A packet belonging to an existing connection arrives (in any direction). ACTION: The packet is translated and forwarded and the timeout counter of the corresponding connection tracking table entry is reset.

4. EVENT: A connection tracking table entry times out. ACTION: The entry is deleted from the connection tracking table.

Due to "black box" testing, the Tester is not able to directly examine (or delete) the entries of the connection tracking table. But the entries can be and MUST be controlled by setting an appropriate timeout value and carefully selecting the port numbers of the packets (as described in Section 4.4) to be able to produce meaningful and repeatable measurement results.

We aim to support the measurement of the following performance characteristics of a stateful NATxy gateway:

1. maximum connection establishment rate

2. all "classic" performance metrics like throughput, frame loss rate, latency, etc.

3. connection tear down rate.
4.4. Control of the Connection Tracking Table Entries

It is necessary to control the connection tracking table entries of the DUT in order to achieve clear conditions for the measurements. We can simply achieve the following two extreme situations:

1. All frames create a new entry in the connection tracking table of the DUT and no old entries are deleted during the test. This is required for measuring the maximum connection establishment rate.

2. No new entries are created in the connection tracking table of the DUT and no old ones are deleted during the test. This is ideal for the real test phase measurements, like throughput, latency, etc.

From this point we use the following three assumptions:

1. A single source address destination address pair is used for all tests. We make this assumption for simplicity. Of course, we are aware that [RFC2544] requires testing also with 256 different destination networks.

2. The connection tracking table of the stateful NATxy is large enough to store all connections defined by the different source port number destination port number combinations.

3. Each experiment is started with an empty connection tracking table. (It can be ensured by deleting its content before the experiment.)

The first extreme situation can be achieved by

* using all different source port number destination port number combinations in the preliminary phase and

* setting the UDP timeout of the NATxy gateway to a value higher than the length of the preliminary phase.

The second extreme situation can be achieved by

* using all different source port number destination port number combinations in the preliminary phase and

* enumerating all the possible source port number destination port number combinations in the preliminary phase and
* setting the UDP timeout of the NATxy gateway to a value higher than the length of the preliminary phase plus the gap between the two phases plus the length of the real test phase.

[RFC4814] REQUIRES pseudorandom port numbers, which we believe is a good approximation of the distribution of the source port numbers a NATxy gateway on the Internet may face with.

We note that pseudorandom all different source port number destination port number combinations may be computing efficiently generated by preparing a random permutation of the previously enumerated all possible source port number destination port number combinations using Dustenfeld’s random shuffle algorithm [DUST1964]. This method also satisfies the criterion for the second case that all possible source port number destination port number combinations must be enumerated during the preliminary phase.

Important warning: in normal (non-NAT) router testing, the port number selection algorithm, whether it is pseudo-random or enumerated in increasing (or decreasing) order does not affect final results. However, our experience with iptables shows that if the connection tracking table is filled using port number enumeration in increasing order, then the maximum connection establishment rate of iptables degrades significantly compared to its performance using pseudorandom port numbers [LEN2021].

The enumeration of the source port number destination port number combinations in increasing or decreasing order (or in any other specific order) MAY be used as an additional measurement.

4.5. Measurement of the Maximum Connection Establishment Rate

The maximum connection establishment rate is an important characteristic of the stateful NATxy gateway and its determination is necessary for the safe execution of the preliminary test phase (without frame loss) before the real test phase.

The measurement procedure of the maximum connection establishment rate is very similar to the throughput measurement procedure defined in [RFC2544].
Procedure: The Initiator sends a specific number of test frames using all different source port number destination port number combinations at a specific rate through the DUT. The Responder counts the frames that are successfully translated by the DUT. If the count of offered frames is equal to the count of received frames, the rate of the offered stream is raised and the test is rerun. If fewer frames are received than were transmitted, the rate of the offered stream is reduced and the test is rerun.

The maximum connection establishment rate is the fastest rate at which the count of test frames successfully translated by the DUT is equal to the number of test frames sent to it by the Initiator.

Notes:

1. In practice, we RECOMMEND the usage of binary search.

2. As for the successful translation, the Responder MAY (or SHOULD?) check that the source IP address is different than the original source IP address set by the Initiator.

4.6. Real Test Phase

As for the traffic direction, there are three possible cases during the real test phase:

* bidirectional traffic: The Initiator sends test frames to the Responder and the Responder sends test frames to the Initiator.

* unidirectional traffic from the Initiator to the Responder: The Initiator sends test frames to the Responder but the Responder does not send test frames to the Initiator.

* unidirectional traffic from the Responder to the Initiator: The Responder sends test frames to the Initiator but the Initiator does not send test frames to the Responder.

If the Initiator sends test frames, then it uses pseudorandom source port numbers and destination port numbers from the restricted port number ranges. The responder receives the test frames, updates its state table and processes the test frames as required by the given measurement procedure (e.g. only counts them for throughput test, handles timestamps for latency or PDV tests, etc.).
If the Responder sends test frames, then it uses the four tuples from its state table. The reading order of the state table may follow different policies (discussed in Section 4.8). The Initiator receives the test frames, and processes them as required by the given measurement procedure.

As for the actual measurement procedures, we RECOMMEND to use the updated ones from Section 7 of [RFC8219].

4.7. Measurement of the Connection Tear Down Rate

Connection tear down can cause significant load for the NATxy gateway. The connection tear down performance can be measured as follows:

1. Load a certain number of connections (N) into the connection tracking table of the DUT (in the same way as done to measure the maximum connection establishment rate).

2. Record TimestampA.

3. Delete the content of the connection tracking table of the DUT.

4. Record TimestampB.

The connection tear down rate can be computed as:

\[
\text{connection tear down rate} = \frac{N}{(\text{TimestampB} - \text{TimestampA})}
\]

The connection tear down rate SHOULD be measured for various (important) values of N.

We assume that the content of the connection tracking table may be deleted by an out-of-band control mechanism specific to the given NATxy gateway implementation. (E.g. by removing the appropriate kernel module under Linux.)

We are aware that the performance of removing the entire content of the connection tracking table at one time may be different from removing all the entries one by one.

4.8. Writing and Reading Order of the State Table

As for writing policy of the state table of the Responder, we RECOMMEND round robin, because it ensures that its entries are automatically kept fresh and consistent with that of the connection tracking table of the DUT.
The Responder can read its state table in various orders, for example:

* pseudorandom
* round robin

We RECOMMEND pseudorandom to follow the spirit of [RFC4814]. Round robin may be used as a computationally cheaper alternative.

5. Implementation and Experience

The "stateful" branch of siitperf [SIITPERF] is an implementation of this concept. It is documented in a paper currently under second review [LEN2022].

Our experience with this methodology using siitperf for measuring the scalability of the iptables stateful NAT44 and Jool stateful NAT64 implementations is described in [I-D.lencse-v6ops-transition-scalability].

6. Limitations of using UDP as Transport Layer Protocol

Stateful NATxy solutions handle TCP and UDP differently, e.g. iptables uses 30s timeout for UDP and 60s timeout for TCP. Thus benchmarking results produced using UDP do not necessarily characterize the performance of a NATxy gateway well enough, when they are used for forwarding Internet traffic. As for the given example, timeout values of the DUT may be adjusted, but it requires extra consideration.

Other differences in handling UDP or TCP are also possible. Thus we recommend that further investigations are to be performed in this field.

As a mitigation of this problem, we recommend that testing with protocols using TCP (like HTTP and HTTPS) can be performed as described in [I-D.ietf-bmwg-ngfw-performance]. This approach also solves the potential problem of protocol helpers may be present in the stateful DUT.

7. Acknowledgements

The authors would like to thank Al Morton, Sarah Banks, Edwin Cordeiro, Lukasz Bromirski and Sandor Repas for their comments.
8. IANA Considerations

This document does not make any request to IANA.

9. Security Considerations

We have no further security considerations beyond that of [RFC8219]. Perhaps they should be cited here so that they be applied not only for the benchmarking of IPv6 transition technologies, but also for the benchmarking of stateful NATxy gateways.

10. References

10.1. Normative References


10.2. Informative References


Appendix A. Change Log

A.1. 00

Initial version.

A.2. 01

Updates based on the comments received on the BMWG mailing list and minor corrections.

A.3. 02

Section 4.4 was completely re-written. As a consequence, the occurrences of the now undefined "mostly different" source port number destination port number combinations were deleted from Section 4.5, too.

A.4. 03

Added Section 4.3 about the consideration of the cases of stateful operation.

Consistency checking. Removal of some parts obsolated by the previous re-writing of Section 4.4.

Added Section 4.7 about the method for measuring connection tear down rate.

Updates for Section 5 about the implementation and experience.

Authors’ Addresses
Abstract

Several IPv6 transition technologies have been developed to provide customers with IPv4-as-a-Service (IPv4aaS) for ISPs with an IPv6-only access and/or core network. All these technologies have their advantages and disadvantages, and depending on existing topology, skills, strategy and other preferences, one of these technologies may be the most appropriate solution for a network operator.

This document examines the scalability of the five most prominent IPv4aaS technologies (464XLAT, Dual Stack Lite, Lightweight 4over6, MAP-E, MAP-T) considering two aspects: (1) how their performance scales up with the number of CPU cores, (2) how their performance degrades, when the number of concurrent sessions is increased until hardware limit is reached.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at https://datatracker.ietf.org/drafts/current/.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on 8 September 2022.

Copyright Notice

Copyright (c) 2022 IETF Trust and the persons identified as the document authors. All rights reserved.
This document is subject to BCP 78 and the IETF Trust’s Legal Provisions Relating to IETF Documents (https://trustee.ietf.org/license-info) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Revised BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Revised BSD License.

Table of Contents

1.  Introduction .................................................. 2
  1.1. Requirements Language .................................... 3
2.  Scalability of iptables ....................................... 3
  2.1. Measurement Method ........................................ 3
  2.2. Performance scale up against the number of CPU cores .. 4
  2.3. Performance degradation caused by the number of sessions ..................................................... 6
  2.4. Connection tear down rate .................................. 8
3.  Scalability of Jool ............................................ 9
  3.1. Measurement Method ........................................ 10
  3.2. Performance scale up against the number of CPU cores .. 10
  3.3. Performance degradation caused by the number of sessions ..................................................... 11
  3.4. Connection tear down rate .................................. 12
4.  Acknowledgements ............................................... 13
5.  IANA Considerations ........................................... 13
6.  Security Considerations ....................................... 13
7.  References .................................................... 13
  7.1. Normative References ...................................... 13
  7.2. Informative References .................................... 14
Appendix A. Change Log ........................................ 15
  A.1. 00 ................................................................ 15
  A.2. 01 ................................................................ 15
Author’s Address .................................................. 15

1.  Introduction

IETF has standardized several IPv6 transition technologies [LEN2019] and occupied a neutral position trusting the selection of the most appropriate ones to the market. [I-D.ietf-v6ops-transition-comparison] provides a comprehensive comparative analysis of the five most prominent IPv4aaS technologies to assist operators with this problem. This document adds one more detail: measurement data regarding the scalability of the examined IPv4aaS technologies.
Currently, this document contains only the scalability measurements of the iptables stateful NAT44 implementation. It serves as a sample to test if the disclosed results are (1) useful and (2) sufficient for the network operators.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

2. Scalability of iptables

2.1. Measurement Method

[RFC8219] has defined a benchmarking methodology for IPv6 transition technologies. [I-D.lencse-bmwg-benchmarking-stateful] has amended it by addressing how to benchmark stateful NATxy gateways using pseudorandom port numbers recommended by [RFC4814]. It has defined a measurement procedure for maximum connection establishment rate and reused the classic measurement procedures like throughput, latency, frame loss rate, etc. from [RFC8219]. We used two of them: maximum connection establishment rate and throughput to characterize the performance of the examined system.

The scalability of iptables is examined in two aspects:

* How its performance scales up with the number of CPU cores?
* How its performance degrades, when the number of concurrent sessions is increased?
The test setup in Figure 1 was followed. The two devices, the Tester and the DUT (Device Under Test), were both Dell PowerEdge R430 servers having two 2.1GHz Intel Xeon E5-2683 v4 CPUs, 384GB 2400MHz DDR4 RAM and Intel 10G dual port X540 network adapters. The NICs of the servers were interconnected by direct cables, and the CPU clock frequency was set to fixed 2.1 GHz on both servers. They had Debian 9.13 Linux operating system with 4.9.0-16-amd64 kernel. The measurements were performed by siitperf [LEN2021] using the "stateful" branch (latest commit Aug. 16, 2021). The DPDK version was 16.11.11-1+deb9u2. The version of iptables was 1.6.0.

The ratio of number of connections in the connection tracking table and the value of the hashsize parameter of iptables significantly influences its performance. Although the default setting is hashsize=nf_conntrack_max/8, we have usually set hashsize=nf_conntrack_max to increase the performance of iptables, which was crucial, when high number of connections were used, because then the execution time of the tests was dominated by the preliminary phase, when several hundreds of millions connections had to be established. (In some cases, we had to use different settings due to memory limitations. The tables presenting the results always contain these parameters.)

The size of the port number pool is an important parameter of the benchmarking method for stateful NATxy gateways, thus it is also given for all tests.

2.2. Performance scale up against the number of CPU cores

To examine how the performance of iptables scales up with the number of CPU cores, the number of active CPU cores was set to 1, 2, 4, 8, 16 using the "maxcpus=" kernel parameter.

The number of connections was always 4,000,000 using 4,000 different source port numbers and 1,000 different destination port numbers. Both the connection tracking table size and the hash table size was set to $2^{23}$.

The error of the binary search was chosen to be lower than 0.1% of the expected results. The experiments were executed 10 times.

Besides the connection establishment rate and the throughput of iptables, also the throughput of the IPv4 packet forwarding of the Linux kernel was measured to provide a basis for comparison.
The results are presented in Figure 2. The unit for the maximum connection establishment rate is 1,000 connections per second. The unit for throughput is 1,000 packets per second (measured with bidirectional traffic, and the number of all packets per second is displayed).

<table>
<thead>
<tr>
<th>num. CPU cores</th>
<th>1</th>
<th>2</th>
<th>4</th>
<th>8</th>
<th>16</th>
</tr>
</thead>
<tbody>
<tr>
<td>src ports</td>
<td>4,000</td>
<td>4,000</td>
<td>4,000</td>
<td>4,000</td>
<td>4,000</td>
</tr>
<tr>
<td>dst ports</td>
<td>1,000</td>
<td>1,000</td>
<td>1,000</td>
<td>1,000</td>
<td>1,000</td>
</tr>
<tr>
<td>num. conn.</td>
<td>4,000,000</td>
<td>4,000,000</td>
<td>4,000,000</td>
<td>4,000,000</td>
<td>4,000,000</td>
</tr>
<tr>
<td>conntrack t. s.</td>
<td>2^23</td>
<td>2^23</td>
<td>2^23</td>
<td>2^23</td>
<td>2^23</td>
</tr>
<tr>
<td>hash table size</td>
<td>2^23</td>
<td>2^23</td>
<td>2^23</td>
<td>2^23</td>
<td>2^23</td>
</tr>
<tr>
<td>c.t.s/num.conn.</td>
<td>2.097</td>
<td>2.097</td>
<td>2.097</td>
<td>2.097</td>
<td>2.097</td>
</tr>
<tr>
<td>num. experiments</td>
<td>10</td>
<td>10</td>
<td>10</td>
<td>10</td>
<td>10</td>
</tr>
<tr>
<td>error</td>
<td>100</td>
<td>100</td>
<td>100</td>
<td>1,000</td>
<td>1,000</td>
</tr>
<tr>
<td>cps median</td>
<td>223.5</td>
<td>371.1</td>
<td>708.7</td>
<td>1,341</td>
<td>2,383</td>
</tr>
<tr>
<td>cps min</td>
<td>221.6</td>
<td>367.7</td>
<td>701.7</td>
<td>1,325</td>
<td>2,304</td>
</tr>
<tr>
<td>cps max</td>
<td>226.7</td>
<td>375.9</td>
<td>723.6</td>
<td>1,376</td>
<td>2,417</td>
</tr>
<tr>
<td>cps rel. scale up</td>
<td>1</td>
<td>0.830</td>
<td>0.793</td>
<td>0.750</td>
<td>0.666</td>
</tr>
<tr>
<td>throughput median</td>
<td>414.9</td>
<td>742.3</td>
<td>1,379</td>
<td>2,336</td>
<td>4,557</td>
</tr>
<tr>
<td>throughput min</td>
<td>413.9</td>
<td>740.6</td>
<td>1,373</td>
<td>2,311</td>
<td>4,436</td>
</tr>
<tr>
<td>throughput max</td>
<td>416.1</td>
<td>746.9</td>
<td>1,395</td>
<td>2,361</td>
<td>4,627</td>
</tr>
<tr>
<td>tp. rel. scale up</td>
<td>1</td>
<td>0.895</td>
<td>0.831</td>
<td>0.704</td>
<td>0.686</td>
</tr>
</tbody>
</table>

IPv4 packet forwarding (using the same port number ranges)

| error     | 200 | 500 | 1,000 | 1,000 |
| throughput median | 910.9| 1,523| 3,016| 5,920 |
| throughput min    | 874.8| 1,485| 2,951| 5,811 |
| throughput max    | 914.3| 1,534| 3,037| 5,940 |
| tp. rel. scale up | 1   | 0.836| 0.828| 0.812|
| throughput ratio (%) | 45.5| 48.8| 45.7| 39.5|

Figure 2: Scale up of iptables against the number of CPU cores
(Please refer to the next figure for the explanation of the abbreviations.)
Whereas the throughput of IPv4 packet forwarding scaled up from 0.91Mpps to 11.56Mpps showing a relative scale up of 0.793, the throughput of iptables scaled up from 414.9kpps to 4,557kpps showing a relative scale up of 0.686 (and the relative scale up of the maximum connection establishment rate is only 0.666). On the one hand, this is the price of the stateful operation. On the other hand, this result is quite good compared to the scale-up results of NSD (a high performance authoritative DNS server) presented in Table 9 of [LEN2020], which is only 0.52. (1,454,661/177,432=8.2-fold performance using 16 cores.) And DNS is not a stateful technology.

2.3. Performance degradation caused by the number of sessions

To examine how the performance of iptables degrades with the number connections in the connection tracking table, the number of connections was increased fourfold by doubling the size of both the source port number range and the destination port number range. Both the connection tracking table size and the hash table size was also increased four fold. However, we reached the limits of the hardware at 400,000,000 connections: we could not set the size of the hash table to $2^{29}$ but only to $2^{28}$. The same value was used at 800,000,000 connections too, when the number of connections was only
doubled, because 1.6 billion connections would not fit into the memory.

The error of the binary search was chosen to be lower than 0.1% of the expected results. The experiments were executed 10 times (except for the very long lasting measurements with 800,000,000 connections).

The results are presented in Figure 4. The unit for the maximum connection establishment rate is 1,000,000 connections per second. The unit for throughput is 1,000,000 packets per second (measured with bidirectional traffic, and the number of all packets per second is displayed).

<table>
<thead>
<tr>
<th>num. conn.</th>
<th>1.56M</th>
<th>6.25M</th>
<th>25M</th>
<th>100M</th>
<th>400M</th>
<th>800M</th>
</tr>
</thead>
<tbody>
<tr>
<td>src ports</td>
<td>2,500</td>
<td>5,000</td>
<td>10,000</td>
<td>20,000</td>
<td>40,000</td>
<td>40,000</td>
</tr>
<tr>
<td>dst ports</td>
<td>625</td>
<td>1,250</td>
<td>2,500</td>
<td>5,000</td>
<td>10,000</td>
<td>20,000</td>
</tr>
<tr>
<td>conntrack t. s.</td>
<td>2^21</td>
<td>2^23</td>
<td>2^25</td>
<td>2^27</td>
<td>2^29</td>
<td>2^30</td>
</tr>
<tr>
<td>hash table size</td>
<td>2^21</td>
<td>2^23</td>
<td>2^25</td>
<td>2^27</td>
<td>2^28</td>
<td>2^28</td>
</tr>
<tr>
<td>num. exp.</td>
<td>10</td>
<td>10</td>
<td>10</td>
<td>10</td>
<td>10</td>
<td>5</td>
</tr>
<tr>
<td>error</td>
<td>1,000</td>
<td>1,000</td>
<td>1,000</td>
<td>1,000</td>
<td>1,000</td>
<td>1,000</td>
</tr>
<tr>
<td>n.c./h.t.s.</td>
<td>0.745</td>
<td>0.745</td>
<td>0.745</td>
<td>0.745</td>
<td>1.490</td>
<td>2.980</td>
</tr>
<tr>
<td>cps median</td>
<td>2.406</td>
<td>2.279</td>
<td>2.278</td>
<td>2.237</td>
<td>2.013</td>
<td>1.405</td>
</tr>
<tr>
<td>cps min</td>
<td>2.358</td>
<td>2.226</td>
<td>2.226</td>
<td>2.124</td>
<td>1.983</td>
<td>1.390</td>
</tr>
<tr>
<td>cps max</td>
<td>2.505</td>
<td>2.315</td>
<td>2.317</td>
<td>2.290</td>
<td>2.050</td>
<td>1.440</td>
</tr>
<tr>
<td>throughput med.</td>
<td>5.326</td>
<td>4.369</td>
<td>4.510</td>
<td>4.516</td>
<td>4.244</td>
<td>3.689</td>
</tr>
<tr>
<td>throughput max</td>
<td>5.533</td>
<td>4.408</td>
<td>4.572</td>
<td>4.537</td>
<td>4.342</td>
<td>3.709</td>
</tr>
</tbody>
</table>

Figure 4: Performance of iptables against the number of sessions

The performance of iptables shows degradation at 6.25M connections compared to 1.56M connections very likely due to the exhaustion of the L3 cache of the CPU of the DUT. Then the performance of iptables is fairly constant up to 100M connections. A small performance decrease can be observed at 400M connections due to the lower hash table size. A more significant performance decrease can be observed at 800M connections. It is caused by two factors:

* on average, about 3 connections were hashed to the same place
* non NUMA local memory was also used.

We note that the CPU has 2 NUMA nodes, cores 0, 2, ... 14 belong to NUMA node 0, and cores 1, 3, ... 15 belong to NUMA node 1. The maximum memory consumption with 400,000,000 connections was below 150GB, thus it could be stored in NUMA local memory.
Therefore, we have pointed out important limitations of the stateful NAT44 technology:

* there is a performance decrease, when approaching hardware limits

* there is a hardware limit, beyond which the system cannot handle the connections at all (e.g. 1600M connections would not fit into the memory).

Therefore, we can conclude that, on the one hand, a well tailored hashing may guarantee an excellent scale-up of stateful NAT44 regarding the number of connections in a wide range, however, on the other hand, stateful operation has its limits resulting both in performance decrease, when approaching hardware limits and also in inability to handle more sessions, when reaching the memory limits.

2.4. Connection tear down rate

[I-D.lencse-bmwg-benchmarking-stateful] has defined connection tear down rate measurement as an aggregate measurement, that is, N number of connections are loaded into the connection tracking table of the DUT and then the entire content of the connection tracking table is deleted, and its deletion time is measured (T). Finally, the connection tear down rate is computed as: N/T.)

We have observed that the deletion of an empty connection tracking table of iptables may take a significant amount of time depending on its size. Therefore, we made our measurements more accurate by subtracting the deletion time of the empty connection tracking table from that of the filled one, thus we got the time spent with the deleting of the connections.

The same setup and parameters were used as in Section 2.3 and the experiments were executed 10 times (except for the long lasting measurements with 800,000,000 connections).

The results are presented in Figure 5.
The connection tear down performance of iptables shows significant degradation at 6.25M connections compared to 1.56M connections very likely due to the exhaustion of the L3 cache of the CPU of the DUT. Then it shows only a minor degradation up to 100M connections. A small performance increase can be observed at 400M connections due to the relatively lower hash table size. A more visible performance decrease can be observed at 800M connections. It is likely caused by keeping the hash table size constant and doubling the number of connections. The same thing that caused performance degradation of the maximum connection establishment rate and throughput, made now the deletion of the connections faster and thus caused an increase of the connection tear down rate.

We note that according to the recommended settings of iptables, 8 connections are hashed to each place of the hash table on average, but we wilfully used much smaller number (0.745 whenever it was possible) to increase the maximum connection establishment rate and thus to speed up experimenting. However, finally this choice significantly slowed down our experiments due to the very low connection tear down rate.

3. Scalability of Jool
3.1. Measurement Method

The same methodology was used as in Section 2, but now the test setup in Figure 6 was followed. The same Tester and DUT devices were used as before, but the operating system of the DUT was updated to Debian 10.11 with 4.19.0-18-amd64 kernel to meet the requirement of the jool-tools package. The version of Jool was 4.1.6. (The most mature version of Jool at the date of starting the measurements, Release Date: 2021-12-10.)

Unlike with iptables, we did not find any way to tune the hashsize or any other parameters of Jool.

3.2. Performance scale up against the number of CPU cores

The number of connections was always 1,000,000 using 2,000 different source port numbers and 500 different destination port numbers.

The error of the binary search was chosen to be lower than 0.1% of the expected results. The experiments were executed 10 times.

The results are presented in Figure 7. The unit for the maximum connection establishment rate is 1,000 connections per second. The unit for throughput is 1,000 packets per second (measured with bidirectional traffic, and the number of all packets per second is displayed).
Both the maximum connection establishment rate and the throughput scaled up poorly with the number of active CPU cores. The increase of the performance was very low above 4 CPU cores.

3.3. Performance degradation caused by the number of sessions

To examine how the performance of Jool degrades with the number of connections, the number of connections was increased fourfold by doubling the size of both the source port number range and the destination port number range. We did not reach the limits of the hardware regarding the number of connections, because unlike iptables, Jool worked also with 1.6 billion connections.

The error of the binary search was chosen to be lower than 0.1% of the expected results and the experiments were executed 10 times (except for the very long lasting measurements with 800 million and 1.6 billion connections to save execution time).

The results are presented in Figure 8. The unit for the maximum connection establishment rate is 1,000 connections per second. The unit for throughput is 1,000 packets per second (measured with bidirectional traffic, and the number of all packets per second is displayed).
The performance of Jool shows degradation at the entire range of the number of connections. We did not analyze the root cause of the degradation yet. And we are not aware of the implementation of its connection tracking table. We also plan to check the memory consumption of Jool, what is definitely lower that that of iptables.

### 3.4. Connection tear down rate

Basically, the same measurement method was used as in Section 2.4, however having no parameter of Jool to tune, only a single measurement series was performed to determine the deletion time of the empty connection tracking table. The median, minimum and maximum values of the 10 measurements were 0.46s, 0.42s and 0.50s respectively.

The same setup and parameters were used as in Section 2.3 and the experiments were executed 10 times (except for the long lasting measurements with 800,000,000 connections).

The results are presented in Figure 9. The unit for the connection tear down rate is 1,000,000 connections per second.

<table>
<thead>
<tr>
<th>Num. conn.</th>
<th>1.56M</th>
<th>6.35M</th>
<th>25M</th>
<th>100M</th>
<th>400M</th>
<th>1600M</th>
</tr>
</thead>
<tbody>
<tr>
<td>src ports</td>
<td>2,500</td>
<td>5,000</td>
<td>10,000</td>
<td>20,000</td>
<td>40,000</td>
<td>40,000</td>
</tr>
<tr>
<td>dst ports</td>
<td>625</td>
<td>1,250</td>
<td>2,500</td>
<td>5,000</td>
<td>10,000</td>
<td>40,000</td>
</tr>
<tr>
<td>num. exp.</td>
<td>10</td>
<td>10</td>
<td>10</td>
<td>10</td>
<td>5</td>
<td>5</td>
</tr>
<tr>
<td>full contr. del med</td>
<td>0.87</td>
<td>2.05</td>
<td>7.84</td>
<td>36.38</td>
<td>126.09</td>
<td>474.68</td>
</tr>
<tr>
<td>full contr. del min</td>
<td>0.80</td>
<td>2.02</td>
<td>7.80</td>
<td>36.27</td>
<td>125.84</td>
<td>473.20</td>
</tr>
<tr>
<td>full contr. del max</td>
<td>0.91</td>
<td>2.09</td>
<td>7.94</td>
<td>36.80</td>
<td>127.54</td>
<td>481.38</td>
</tr>
<tr>
<td>empty contr. del med</td>
<td>0.46</td>
<td>0.46</td>
<td>0.46</td>
<td>0.46</td>
<td>0.46</td>
<td>0.46</td>
</tr>
<tr>
<td>conn. deletion time</td>
<td>0.41</td>
<td>1.59</td>
<td>7.38</td>
<td>35.92</td>
<td>125.63</td>
<td>474.22</td>
</tr>
<tr>
<td>conn. t. d. r. (M)</td>
<td>3.811</td>
<td>3.931</td>
<td>3.388</td>
<td>2.784</td>
<td>3.184</td>
<td>3.374</td>
</tr>
</tbody>
</table>
The connection tear down performance of Jool is excellent at any number of connections. It is about an order of magnitude higher than its connection establishment rate and than the connection tear down rate of iptables. (A slight degradation can be observed at 100M connections.)

4. Acknowledgements

The measurements were carried out by remotely using the resources of NICT StarBED, 2-12 Asahidai, Nomi-City, Ishikawa 923-1211, Japan. The author would like to thank Shuuhei Takimoto for the possibility to use StarBED, as well as to Satoru Gonno and Makoto Yoshida for their help and advice in StarBED usage related issues.

The author would like to thank Ole Troan for his comments on the v6ops mailing list, while the scalability measurements of iptables were intended to be a part of [I-D.ietf-v6ops-transition-comparison].

5. IANA Considerations

This document does not make any request to IANA.

6. Security Considerations

TBD.

7. References

7.1. Normative References


7.2. Informative References

[I-D.ietf-v6ops-transition-comparison]

[I-D.lencse-bmwg-benchmarking-stateful]


Appendix A. Change Log

A.1. 00

Initial version: scale up of iptables.

A.2. 01

Added the scale up of Jool.

Author’s Address

Gabor Lencse
Szechenyi Istvan University
Gyor
Egyetem ter 1.
H-9026
Hungary
Email: lencse@sze.hu
A YANG Data Model for Network Interconnect Tester Management
draft-vassilev-bmwg-network-interconnect-tester-07

Abstract

This document introduces new YANG model for use in network interconnect testing containing modules of traffic generator and traffic analyzer.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at https://datatracker.ietf.org/drafts/current/.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on 28 April 2022.

Copyright Notice

Copyright (c) 2021 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust’s Legal Provisions Relating to IETF Documents (https://trustee.ietf.org/license-info) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Simplified BSD License.
1. Introduction

There is a need for standard mechanism to allow the specification and implementation of the transactions part of network tests. The mechanism should allow the control and monitoring of the data plane traffic in a transactional way. This document defines two YANG modules for test traffic generator and analyzer.

The YANG modules in this document conform to the Network Management Datastore Architecture (NMDA) defined in RFC 8342.

1.1. Terminology

1.1.1. Definitions and Acronyms

DUT: Device Under Test

TA: Traffic Analyzer

TG: Traffic Generator
1.1.2. Tree Diagram

For a reference to the annotations used in tree diagrams included in this document, please see YANG Tree Diagrams [RFC8340].

1.2. Problem Statement

Network interconnect tests require active network elements part of the tested network that generate test traffic and network elements that analyze the test traffic at one or more points of its path. A network interconnect tester is a device that can either generate test traffic, analyze test traffic or both. Here is a figure borrowed from [RFC2544] representing the horseshoe test setup topology consisting of a single tester and a single DUT connected in a network interconnect loop.

```
+------------+
|            |
+------------|  tester    |<-------------+
|            |            |              |
|            +------------+              |
|                                        |
|            +------------+              |
|            |            |              |
+----------->|    DUT     |--------------+
|            |            |              |
|            +------------+
```

This document attempts to address the problem of defining YANG model of a network interconnect tester that can be used for development of vendor independent network interconnect tests and utilize the advantages of transactional management using standard protocols like NETCONF.

1.3. Objectives

This section describes some of the design objectives for the model. It should:

* provide means to specify the generated traffic as streams of cyclic sequence of bursts with configurable frame size, frame data, interframe gap and interburst gap.

* have a mandatory single stream mode and optional multi stream mode.
* provide means for configuration of traffic streams with static frame data where frames with identical frame data are sent during the lifetime of the stream.

* provide means for configuration of traffic streams with dynamic frame data where frames contain fields with dynamic data like generation time and sequence number.

* allow third parties to augment the base module with alternative dynamic fields of frame data extensions.

* provide means for realtime synchronization and orchestration of the generated streams.

* provide counters for received test traffic frames and octets.

* provide latency statistic in the case of test traffic with dynamic frame data that includes timestamp.

* provide sequence number errors in the case of test traffic with dynamic frame data that includes sequence number.

1.4. Solution

The proposed model splits the design into 2 modules - 1) Traffic Generator module (TG), 2) Traffic Analyzer module (TA). The modules are implemented as augmentations of the ietf-interfaces [RFC8343] module adding configuration and state data that models the functionality of a network interconnect tester. The TA and TG modules concept is illustrated with the following diagram of a tester with two interfaces (named e0 and e1) connected in a loop with single DUT:
2. Using the network interconnect tester model

Basic example of how the model can be used in transactional network test program to control the testers part of a network and report counter statistics and timing measurement data is presented in Appendix A. All example cases present the configuration and state data from a single test trial. The search algorithm logic that operates to control the trial configuration is outside the scope of this document. One of the examples demonstrates the use of the [RFC2544] defined testframe packet.

3. Traffic Generator Module Tree Diagram

```yang
module: ietf-traffic-generator
augment /if:interfaces/if:interface:
  +--rw traffic-generator (egress-direction)?
  |  +--rw (type)?
  |     |  +--:(single-stream)
  |     |     |  +--rw testframe-type? identityref
  |     |     |  +--rw frame-size uint32
  |     |     |  +--rw frame-data? string
  |     |     |  +--rw interframe-gap uint32
  |     |     |  +--rw interburst-gap? uint32
  |     |     |  +--rw frames-per-burst? uint32
  |     |     |  +--rw src-mac-address? yang:mac-address {ethernet}?
  |     |     |  +--rw dst-mac-address? yang:mac-address {ethernet}?
  |     |     |  +--rw ether-type? uint16 {ethernet}?
  |     |  +--:(multi-stream)
  |     |     +--rw streams
  |     |        +--rw stream* [id]
  |     |        |  +--rw id uint32
  |     |        |  +--rw testframe-type? identityref
  |     |        |  +--rw frame-size uint32
  |     |        |  +--rw frame-data? string
  |     |        |  +--rw interframe-gap uint32
  |     |        |  +--rw interburst-gap? uint32
  |     |        |  +--rw frames-per-burst? uint32
  |     |        |  +--rw frames-per-stream uint32
  |     |        |  +--rw interstream-gap uint32
  |     |        |  +--rw src-mac-address? yang:mac-address {ethernet}?
  |     |        |  |  +--rw dst-mac-address? yang:mac-address {ethernet}?
  |     |        |  |  +--rw ether-type? uint16 {ethernet}?
  |     |        |  |  +--rw realtime-epoch?
  |     |        |  |     |  +--rw date-and-time {realtime-epoch}?
  |     |        |  |  +--rw total-frames? uint64
  |     |  +--rw traffic-generator-ingress {ingress-direction)}?
```
4. Traffic Analyzer Module Tree Diagram

module: ietf-traffic-analyzer

  augment /if:interfaces/if:interface:
    +--rw traffic-analyzer! {ingress-direction}? 
      +--rw filter! {filter}? 
        +--rw type identityref
        +--rw ether-type? uint16
      +--rw capture {capture}? 
      +--rw start-trigger 
        +--rw (start-trigger)? 
          +--:(frame-index) 
            | +--rw frame-index? uint64 
            +--:(testframe-index) 
              +--rw testframe-index? uint64
5. Traffic Generator Module YANG

<CODE BEGINS> file "ietf-traffic-generator@2021-10-25.yang"

define module ietf-traffic-generator {
    yang-version 1.1;
    namespace "urn:ietf:params:xml:ns:yang:ietf-traffic-generator";
    prefix nttg;

    import ietf-interfaces {
        prefix if;
        reference "RFC 8343: A YANG Data Model For Interface Management";
    }
    import ietf-yang-types {
        prefix yang;
        reference "RFC 6991: Common YANG Data Types";
    }
    import iana-if-type {
        prefix ianaift;
        reference "RFC 7224: IANA Interface Type YANG Module";
    }

    organization "IETF Benchmarking Methodology Working Group";
    contact "WG Web: <http://tools.ietf.org/wg/bmwg/>
             WG List: <mailto:bmwg@ietf.org>
             Editor: Vladimir Vassilev
             <mailto:vladimir@lightside-instruments.com>";
description
"This module contains a collection of YANG definitions for
description and management of network interconnect testers.

Copyright (c) 2021 IETF Trust and the persons identified as
authors of the code. All rights reserved.

Redistribution and use in source and binary forms, with or
without modification, is permitted pursuant to, and subject
to the license terms contained in, the Simplified BSD License
set forth in Section 4.c of the IETF Trust’s Legal Provisions
Relating to IETF Documents
(http://trustee.ietf.org/license-info).
This version of this YANG module is part of RFC XXXX; see
the RFC itself for full legal notices.";

revision 2021-10-25 {
  description
    "Initial revision.";
  reference
    "RFC XXXX: A YANG Data Model for
     Network Interconnect Tester Management";
}

feature egress-direction {
  description
    "The device can generate traffic in the egress direction.";
}

feature ingress-direction {
  description
    "The device can generate traffic in the ingress direction.";
}

feature multi-stream {
  description
    "The device can generate multi-stream traffic.";
}

feature ethernet {
  description
    "The device can generate ethernet traffic.";
}

feature realtime-epoch {
  description
    "The device can generate traffic precisely
     at configured realtime epoch.";
}
identity testframe-type {
    description
        "Base identity for all testframe types.";
}

identity static {
    base testframe-type;
    description
        "Identity for static testframe.
        The frame data and size are constant.";
}

identity dynamic {
    base testframe-type;
    description
        "Identity to be used as base for dynamic
        testframe type identities defined
        in external modules.

        When used itself it identifies dynamic testframe
        where the last 18 octets of the payload contain
        incrementing sequence number field (8 octets)
        followed by timestamp field in the
        IEEE 1588-2008 format (10 octets). If frame data is defined
        for the last 18 octets of the payload it will be ignored
        and overwritten with dynamic data according to this
        specification.";
}

grouping common-data {
    description
        "Common configuration data.";
    leaf realtime-epoch {
        if-feature "realtime-epoch";
        type yang:date-and-time;
        description
            "If this leaf is present the stream generation will start
            at the specified realtime epoch.";
    }
    leaf total-frames {
        type uint64;
        description
            "If this leaf is present the traffic generation will stop
            after the specified number of frames are generated.";
    }
}
grouping burst-data {

description "Generated traffic burst parameters.";

leaf testframe-type {

type identityref {
   base nttg:testframe-type;
}

default "nttg:static";

description "In case of dynamic testframes this leaf specifies the dynamic testframe identity.";
}

leaf frame-size {

type uint32;

mandatory true;

description "Size of the frames generated. For example for ethernet interfaces the following definition applies:

Ethernet frame-size in octets includes:
* Destination Address (6 octets),
* Source Address (6 octets),
* Frame Type (2 octets),
* Data (min 46 octets or 42 octets + 4 octets 802.1Q tag),
* CRC Checksum (4 octets).

Ethernet frame-size does not include:
* Preamble (dependent on MAC configuration by default 7 octets),
* Start of frame delimiter (1 octet)

Minimum standard ethernet frame-size is 64 bytes but generators might support smaller sizes for validation.";
}

leaf frame-data {

type string {
   pattern '([0-9A-F]{2})*';
}

must 'string-length(.)<=(../frame-size*2)';

description "The raw frame data specified as hexadecimal string. The specified data can be shorter than the ../frame-size value specifying only the header or the header and the payload with or without the 4 byte CRC Checksum in the case of a Ethernet frame.";
}

leaf interframe-gap {
type uint32;
mandatory true;
description
"Length of the idle period between generated frames. For example for ethernet interfaces the following definition applies:

Ethernet interframe-gap between transmission of frames known as the interframe gap (IFG). A brief recovery time between frames allows devices to prepare for reception of the next frame. The minimum interframe gap is 96 bit times (12 octet times) (the time it takes to transmit 96 bits (12 octets) of raw data on the medium). However the preamble (7 octets) and start of frame delimiter (1 octet) are considered a constant gap that should be included in the interframe-gap. Thus the minimum value for standard ethernet transmission should be considered 20 octets.";
}
leaf interburst-gap {
    type uint32;
description
"Similar to the interframe-gap but takes place between any two bursts of the stream.";
}
leaf frames-per-burst {
    type uint32;
description
"Number of frames contained in a burst";
}
}

grouping multi-stream-data {
description
"Multi stream traffic generation parameters.";
container streams {
description
"Non-presence container holding the configured stream list.";
list stream {
    key "id";
description
"Each stream repeats a burst until frames-per-stream count is reached followed by interstream-gap delay.";
leaf id {
    type uint32;
description
"Number specifying the order of the stream.";
}
uses burst-data;
leaf frames-per-stream {
  type uint32;
  mandatory true;
  description
  "The count of frames to be generated before
   generation of the next stream is started.";
}
leaf interstream-gap {
  type uint32;
  mandatory true;
  description
  "Idle period after the last frame of the last burst.";
}


grouping ethernet-data {
  description
  "Ethernet frame data specific parameters.";
  reference
  "IEEE 802-2014 Clause 9.2";
  leaf src-mac-address {
    type yang:mac-address;
    description
    "Source Address field of the generated Ethernet packet.";
  }
  leaf dst-mac-address {
    type yang:mac-address;
    description
    "Destination Address field of the generated Ethernet packet.";
  }
  leaf ether-type {
    type uint16;
    description
    "Length/Type field of the generated Ethernet packet.";
  }
}

augment "/if:interfaces/if:interface" {
  description
  "Traffic generator augmentations of ietf-interfaces.";
  container traffic-generator {
    if-feature "egress-direction";
    description
    "Traffic generator for egress direction.";
    choice type {

description
  "Choice of the type of the data model of the generator. 
    Single or multi stream.";
case single-stream {
  uses burst-data;
}
case multi-stream {
  uses multi-stream-data;
}
}
uses common-data;
}
container traffic-generator-ingress {
  if-feature "ingress-direction";
  description
    "Traffic generator for ingress direction.";
  choice type {
    description
      "Choice of the type of the data model of the generator. 
        Single or multi stream.";
    case single-stream {
      uses burst-data;
    }
    case multi-stream {
      uses multi-stream-data;
    }
    }
  uses common-data;
}
}

augment "/if:interfaces/if:interface/nttg:traffic-generator/"
  + "nttg:type/nttg:single-stream" {
    when "derived-from-or-self(../if:type, 'ianaift:ethernetCsmacd')" {
      description
        "Ethernet interface type.";
    }
  }
if-feature "ethernet";
description
  "Ethernet specific augmentation for egress 
  single stream generator type.";
uses ethernet-data;
}

augment "/if:interfaces/if:interface/nttg:traffic-generator/"
  + "nttg:type/nttg:multi-stream/nttg:streams/nttg:stream" {
    when "derived-from-or-self(../../if:type,"
    + "'ianaift:ethernetCsmacd')" {
      
    }

Vassilev                  Expires 28 April 2022                
[Page 14]
description
   "Ethernet interface type.";
}
if-feature "ethernet";
description
   "Ethernet specific augmentation for egress
   multi stream generator type.";
uses ethernet-data;
}

augment "/if:interfaces/if:interface/nttg:traffic-generator-ingress/
   + "nttg:type/nttg:single-stream" {
   when "derived-from-or-self(../if:type, 'ianaift:ethernetCsmacd')" {
      description
         "Ethernet interface type.";
   }
   if-feature "ethernet";
description
   "Ethernet specific augmentation for ingress
   single stream generator type.";
uses ethernet-data;
}

augment "/if:interfaces/if:interface/nttg:traffic-generator-ingress/
   + "nttg:type/nttg:multi-stream/nttg:streams/nttg:stream"
   + "'ianaift:ethernetCsmacd'" {
      description
         "Ethernet interface type.";
   }
   if-feature "ethernet";
description
   "Ethernet specific augmentation for ingress
   multi stream generator type.";
uses ethernet-data;
}

<CODE ENDS>

6. Traffic Analyzer Module YANG

<CODE BEGINS> file "ietf-traffic-analyzer@2021-10-25.yang"
module ietf-traffic-analyzer {
  yang-version 1.1;
  namespace "urn:ietf:params:xml:ns:yang:ietf-traffic-analyzer";
  prefix ntta;

  import ietf-interfaces {
    prefix if;
    reference
      "RFC 8343: A YANG Data Model For Interface Management";
  }
  import ietf-yang-types {
    prefix yang;
    reference
      "RFC 6991: Common YANG Data Types";
  }

  organization
    "IETF Benchmarking Methodology Working Group";
  contact
    "WG Web:  <http://tools.ietf.org/wg/bmwg/>
     WG List:  <mailto:bmwg@ietf.org>
    Editor:   Vladimir Vassilev
              <mailto:vladimir@lightside-instruments.com>"
  description
    "This module contains a collection of YANG definitions for
description and management of network interconnect testers.

  Copyright (c) 2021 IETF Trust and the persons identified as
  authors of the code.  All rights reserved.

  Redistribution and use in source and binary forms, with or
  without modification, is permitted pursuant to, and subject
  to the license terms contained in, the Simplified BSD License
  set forth in Section 4.c of the IETF Trust's Legal Provisions
  Relating to IETF Documents
  (http://trustee.ietf.org/license-info).

  This version of this YANG module is part of RFC XXXX; see
  the RFC itself for full legal notices.";

  revision 2021-10-25 {
    description
      "Initial revision.";
    reference
      "RFC XXXX: A YANG Data Model for
      Network Interconnect Tester Management";
  }
}
feature egress-direction {
    description
    "The device can analyze traffic from the egress direction.";
}

feature ingress-direction {
    description
    "The device can generate traffic from the ingress direction.";
}

feature filter {
    description
    "This feature indicates that the device implements filter that can specify a subset of packets to be analyzed when filtering is enabled.";
}

feature idle-octets-counter {
    description
    "This feature indicates that the device implements idle-octets counter that accumulates the time the link is not utilized. The minimum required idle gaps are not counted as idle octets.";
}

feature capture {
    description
    "This feature indicates that the device implements packet capture functionality.";
}

identity filter {
    description
    "Base filter identity.";
}

identity ethernet {
    base ntta:filter;
    description
    "Ethernet packet fields filter.";
}

grouping statistics-data {
    description
    "Analyzer statistics.";
    leaf pkts {
        type yang:counter64;
        description
        "Total number of packets analyzed.";
    }
}
leaf octets {
    type yang:counter64;
    description
    "This counter is identical with the in-octets/out-octets
    counters defined in RFC8343 except that it counts the
    octets since the analyzer was created.";
}
leaf idle-octets {
    if-feature "idle-octets-counter";
    type yang:counter64;
    description
    "Total accumulated period with no frame transmission
    taking place measured in octets at the current link
    speed. Octets not counted in ../octets but not idle are
    for example layer 1 framing octets - for Ethernet links
    7+1 preamble octets per packet.";
}
leaf errors {
    type yang:counter64;
    description
    "Count of packets with errors.
    Not counted in the pkts or captured.
    For example packets with CRC error.";
}
container testframe-stats {
    description
    "Statistics for testframe packets containing
    either sequence number, payload checksum,
    timestamp or any combination of these features.";
    leaf testframe-pkts {
        type yang:counter64;
        description
        "Total count of detected testframe packets.";
    }
    leaf sequence-errors {
        type yang:counter64;
        description
        "Total count of testframe packets with
        unexpected sequence number. After each sequence
        error the expected next sequence number is
        updated.";
    }
    leaf payload-errors {
        type yang:counter64;
        description
        "Total count of testframe packets with
        payload errors.";
    }
container latency {
  description  
    "Latency statistics.";
  leaf samples {
    type uint64;
    description  
      "Total count of packets used for estimating 
      the latency statistics. Ideally 
      samples=../testframe-stats.";
  }
  leaf min {
    type uint64;
    units "nanoseconds";
    description  
      "Minimum measured latency.";
  }
  leaf max {
    type uint64;
    units "nanoseconds";
    description  
      "Maximum measured latency.";
  }
  leaf average {
    type uint64;
    units "nanoseconds";
    description  
      "The sum of all sampled latencies divided 
      by the number of samples.";
  }
  leaf latest {
    type uint64;
    units "nanoseconds";
    description  
      "Latency of the latest sample.";
  }
}

grouping capture-config-data {
  description  
    "Grouping with a capture configuration container.";
  container capture {
    if-feature "capture";
    description  
      "Contains capture parameters.";
  }
}
container start-trigger {
    description
    "Configures when the capture start is triggered.";
    choice start-trigger {
        description
        "If none of the cases in this choice are configured the capture process starts from the first frame received.";
        case frame-index {
            description
            "Start capturing frames at the specified frame index.";
            leaf frame-index {
                type uint64;
                description
                "First captured frame index.";
            }
        }
        case testframe-index {
            description
            "Start capturing frames at the specified testframe index.";
            leaf testframe-index {
                type uint64;
                description
                "Starts capture as specified testframe index.";
            }
        }
    }
}

container stop-trigger {
    description
    "Configures when the capture is stopped.";
    choice stop-trigger {
        description
        "If none of the cases in this choice are configured the captured frames are always the last frames received for as many frames the implementation can buffer.";
        case when-full {
            description
            "Stops capturing when the implementation can not store more frames.";
            leaf when-full {
                type empty;
                description
                "When present in configuration capture stops when the capture buffer is full.";
            }
        }
    }
}
grouping capture-data {
    description
        "Grouping with statistics and data
        of one or more captured frames.";
    container capture {
        if-feature "capture";
        description
            "Statistics and data of
            one or more captured frames.";
        list frame {
            key "sequence-number";
            description
                "Statistics and data of a captured frame.";
            leaf sequence-number {
                type uint64;
                description
                    "Incremental counter of frames captured.";
            }
            leaf timestamp {
                type yang:date-and-time;
                description
                    "Timestamp of the moment the frame was captured.";
            }
            leaf length {
                type uint32;
                description
                    "Frame length. Ideally the data captured will be
                    of the same length but can be shorter
                    depending on implementation limitations.";
            }
            leaf preceding-interframe-gap {
                type uint32;
                units "nanoseconds";
                description
                    "Measured delay between the reception of the previous
                    frame was completed and the reception of the current
                    frame was started.";
            }
            leaf data {
                type string {
                    pattern '([0-9A-F]{2})*';
                }
                description
                    "Raw data of the captured frame.";
            }
        }
    }
}
grouping filter-data {
  description
    "Grouping with a filter container specifying the filtering
    rules for processing only a specific subset of the frames.";
  container filter {
    if-feature "filter";
    presence "When present packets are
        filtered before analyzed according
        to the filter type";
    description
      "Contains the filtering rules for processing only
        a specific subset of the frames.";
    leaf type {
      type identityref {
        base ntta:filter;
      }
      mandatory true;
      description
        "Type of the applied filter. External modules can
        define alternative filter type identities.";
    }
  }
}

augment "/if:interfaces/if:interface" {
  description
    "Traffic analyzer augmentations of ietf-interfaces.";
  container traffic-analyzer {
    if-feature "ingress-direction";
    presence "Enables the traffic analyzer for ingress traffic.";
    description
      "Traffic analyzer for ingress direction.";
    uses filter-data;
    uses capture-config-data;
    container state {
      config false;
      description
        "State data.";
      uses statistics-data;
      uses capture-data;
    }
  }
}
container traffic-analyzer-egress {
  if-feature "egress-direction";
  presence "Enables the traffic analyzer for egress traffic.";
  description "Traffic analyzer for egress direction.";
  uses filter-data;
  uses capture-config-data;
  container state {
    config false;
    description "State data.”;
    uses statistics-data;
    uses capture-data;
  }
}

augment "/if:interfaces/if:interface/ntta:traffic-analyzer/"
  + "ntta:filter" {
    when "derived-from-or-self(ntta:type, 'ntta:ethernet')";
    description "Ethernet frame specific filter type.";
    leaf ether-type {
      type uint16;
      description "The Ethernet Type (or Length) value defined by IEEE 802.";
      reference "IEEE 802-2014 Clause 9.2";
    }
}

<CODE ENDS>

7. IANA Considerations

This document registers two URIs and two YANG modules.

7.1. URI Registration

This document registers two URIs in the IETF XML registry [RFC3688]. Following the format in RFC 3688, the following registration is requested to be made:

Registrant Contact: The IESG.

XML: N/A, the requested URI is an XML namespace.

7.2. YANG Module Name Registration

This document registers two YANG module in the YANG Module Names registry YANG [RFC6020].

name: ietf-traffic-generator
prefix: nttg
reference: RFC XXXX

name: ietf-traffic-analyzer
prefix: ntta
reference: RFC XXXX

8. Security Considerations

The YANG modules defined in this document are designed to be accessed via the NETCONF protocol RFC 6241 [RFC6241]. The lowest NETCONF layer is the secure transport layer and the mandatory to implement secure transport is SSH RFC 6242 [RFC6242]. The NETCONF access control model RFC 6536 [RFC6536] provides the means to restrict access for particular NETCONF users to a pre-configured subset of all available NETCONF protocol operations and content.

There are a number of data nodes defined in this YANG module which are writable/creatable/deletable (i.e. config true, which is the default). These data nodes may be considered sensitive or vulnerable in some network environments. Write operations (e.g. edit-config) to these data nodes without proper protection can have a negative effect on network operations. These are the subtrees and data nodes and their sensitivity/vulnerability:

8.1. ietf-traffic-generator.yang

The ietf-traffic-generator YANG module controls a stateless traffic generator which is intended to be used for testing and verification purposes but can be used for malicious purposes like generating network traffic part of a Denial-of-Service (DoS) attack. This should be taken into consideration when granting write access to the following container and descendant data nodes:

* /if:interfaces/if:interface/nttg:traffic-generator
8.2. ietf-traffic-analyzer.yang

The ietf-traffic-analyzer YANG module controls a traffic analyzer which is designed for use in testing and verification but can be used for reading information contained in packets sent and received on any of the interfaces on systems that implement the capture feature. This should be taken into consideration when granting read access to the following container and descendant data nodes:

* /if:interfaces/if:interface/ntta:traffic-analyzer/ntta:capture

9. References

9.1. Normative References


Appendix A. Examples

The following topology will be used for the examples in this section:

```
+-------------+          +------------+         +------------+
|             | e0    e0 |            | e1   e0 |            |
| tester0   TG|>-------->|    dut0    |>------->|TA  tester1 |
|             |          |            |         |            |
+-------------+          +------------+         +------------+
```

A.1. Basic Test Program

This pseudo code program orchestrates a network test and shows how the model can be used:

```python
# Connect to network
net=connect("topology.xml")

# Configure DUTs and enable traffic-analyzers
net.node("dut0").edit( \
  "create /interfaces/interface[name='e0'] -- type=ethernetCsmacd")
net.node("dut0").edit(
```
"create /interfaces/interface[name='e1'] -- type=ethernetCsmacd"
net.node("dut0").edit(
    "create /flows/flow[id='t0'] -- match/in-port=e0 
    "actions/action[order='0']/output-action/out-port=e1")

net.node("tester1").edit(
    "create /interfaces/interface[name='e0']/traffic-analyzer")
net.commit()

#Get network state - before
before=net.get()

# Start traffic
net.node("tester0").edit(
    "create /interfaces/interface[name='e0']/traffic-generator -- "
    "frame-size=64 interframe-gap=20")
net.commit()

time.sleep(60)

# Stop traffic
net.node("tester1").edit("delete /interfaces/interface[name='e0']/"
    "traffic-generator")
net.commit()

#Get network state - after
after=net.get()

#Report
sent_pkts=delta("tester0",before,after,
    "/interfaces/interface[name='e0']/statistics/out-unicast-pkts")

received_pkts=delta("tester1",before,after,
    "/interfaces/interface[name='e0']/statistics/in-unicast-pkts")

latency_max=absolute(after,
    "/interfaces/interface[name='e0']/traffic-analyzer/state/"
    "testframe-stats/latency/max")

#Cleanup
net.node("tester1").edit(
    "delete /interfaces/interface/traffic-analyzer")
net.node("dut0").edit("delete /flows")
net.node("dut0").edit("delete /interfaces")
net.commit()
A.2. Generating RFC2544 Testframes

In sec. C.2.6.4 Test Frames a detailed format is specified. The frame-data leaf allows full control over the generated frames payload.

...  
net.node("tester1").edit(
  "merge /interfaces/interface[name='e0']/"
  "traffic-generator -- frame-data="
  "6CA96F0000026CA96F00000108004500"
  "002ED4A500000A115816C0000201C000"
  "0202C0200007001A000010203040506"
  "0708090A0B0C0D0E0F101112")
...
Benchmarking Methodology for MPLS Segment Routing

draft-vfv-bmwg-srmls-bench-meth-01

Abstract

This document defines a methodology for benchmarking Segment Routing (SR) performance for Segment Routing over MPLS (SR-MPLS). It builds upon [RFC2544], [RFC5695] and [RFC8402].

Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at https://datatracker.ietf.org/drafts/current/.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on September 8, 2022.

Copyright Notice

Copyright (c) 2022 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust’s Legal Provisions Relating to IETF Documents
1. Introduction

Segment Routing (SR), defined in [RFC8402], leverages the source routing paradigm. The headend node steers a packet through an SR Policy [I-D.ietf-spring-segment-routing-policy], instantiated as an ordered list of segments. A segment, referred to by its Segment Identifier (SID), can have a semantic local to an SR node or global within an SR domain. SR supports per-flow explicit routing while maintaining per-flow state only at the ingress nodes to the SR domain.

However, there is no standard method defined to compare and contrast the foundational SR packet forwarding capabilities of network
devices. This document aims to extend the efforts of [RFC1242] and [RFC2544] to SR network.

The SR architecture can be instantiated on two data-plane: SR over MPLS (SR-MPLS) and SR over IPv6 (SRv6). This document is limited to SR-MPLS.

SR can be directly applied to the Multiprotocol Label Switching (MPLS) architecture with no change to the forwarding plane [RFC8660]. A segment is encoded as an MPLS label. An SR Policy is instantiated as a stack of labels.

For Segment Routing, PUSH, NEXT, and CONTINUE are operations applied by the forwarding plane.

PUSH consists of the insertion of a segment at the top of the segment list. In SR-MPLS, the top of the segment list is the outer label of the label stack.

NEXT consists of the inspection of the next segment. The active segment is completed and the next segment becomes active. In SR-MPLS, NEXT is implemented as a POP of the top label.

CONTINUE happens when the active segment is not completed; hence, it remains active. In SR-MPLS, the CONTINUE operation is implemented as a SWAP of the top label.

[RFC5695] describes a methodology specific to the benchmarking of MPLS forwarding devices, by considering the most common MPLS packet forwarding scenarios and corresponding performance measurements.

The purpose of this document is to describe a methodology specific to the benchmarking of Segment Routing. The methodology described is a complement for [RFC5695].

2. SR-MPLS Forwarding

In MPLS, a Prefix-SID is allocated in the form of an MPLS label. For SR-MPLS, Segment Routing does not require any change to the MPLS forwarding plane. An SR Policy is instantiated through the MPLS Label Stack: the Segment IDs (SIDs) of a Segment List are inserted as MPLS Labels. The classical forwarding functions available for MPLS networks allow implementing the SR operations.

The operations applied by the SR-MPLS forwarding plane are PUSH, NEXT, and CONTINUE.
The PUSH operation corresponds to the Label Push function, according to the MPLS label pushing rules specified in [RFC3032]. It consists of pushing one or more MPLS labels on top of an incoming packet then sending it out of a particular physical interface or virtual interface towards a particular next hop.

The NEXT operation corresponds to the Label Pop function, that consists of removing the topmost label. The action before and/or after the popping depends on the instruction associated with the active SID on the received packet prior to the popping. It is equivalent to Penultimate Hop Popping (PHP).

The CONTINUE operation corresponds to the Label Swap function, according to the MPLS label-swapping rules in [RFC3031]. It consists of associating an incoming label with an outgoing interface and outgoing label and forwarding the packet on the outgoing interface. It is equivalent to Ultimate Hop Popping (UHP).

The encapsulation of an IP packet into an SR-MPLS packet is performed at the edge of an SR-MPLS domain, reusing the MPLS Forwarding Equivalent Class (FEC) concept. A Forwarding Equivalent Class (FEC) can be associated with an SR Policy ([RFC8660]). When pushing labels onto a packet’s label stack, the Time-to-Live (TTL) field and the Traffic Class (TC) field of each label stack entry must also be set.

All SR nodes in the SR domain use an IGP signaling extension to advertise their own prefix SIDs. After receiving advertised prefix SIDs, each SR node calculates the prefix SIDs to the advertisers. The prefix SID advertisement can be an absolute value advertisement or an index value advertisement. In this regard, the mapping of Segments to MPLS Labels (SIDs) is an important process in the SR-MPLS data plane. Each router can advertise its own available label space to be used for Global Segments called Segment Routing Global Block (SRGB) and an identical range of labels (SRGB) should be used in all routers in order to simplify services and operations. In the SR domain Global Segments can be identified by an index, which has to be re-mapped into a label, or by an absolute value. This is relevant for the nodes that perform the NEXT operation to the segments, because the label for the next segments needs to be crafted accordingly.

[I-D.ietf-spring-segment-routing-policy] specifies the concepts of SR Policy and steering into an SR Policy. The header of a packet steered in an SR Policy is augmented with the ordered list of segments associated with that SR Policy. SR Policy state is instantiated only on the headend node, that steers a flow into an SR Policy. Indeed intermediate and endpoint nodes do not require any state to be maintained. SR Policies can be instantiated on the
headend dynamically and on demand basis. Moreover, signaling can be used in the case of a controller based deployment. For all these reasons, SR Policies scale better than traditional TE mechanisms.

3. Test Methodology

3.1. Test Setup

The Device Under Test (DUT) is connected to the test ports on the test tool according to [RFC2544].

The recommended topology for SR-MPLS Forwarding Benchmarking should be the same as MPLS and it is described in [RFC5695] for both single-port and multi-port scenarios. Indeed, the number of ports is a parameter that MUST be reported.

3.2. IGP and BGP Support

It is RECOMMENDED that all of the ports on the DUT and test tool support a Segment Routing extensions for dynamic Interior Gateway Protocol (IGP) for routing such as IS-IS [RFC8667] and OSPF [RFC8665] as well as Border Gateway Protocol (BGP) [RFC8669].

As specified in [RFC8402], in the context of an IGP-based distributed control plane, two topological segments are defined: the IGP-Adjacency segment and the IGP-Prefix segment; while, in the context of a BGP-based distributed control plane, two topological segments are defined: the BGP peer segment and the BGP Prefix segment.

The distribution method that is used (e.g. OSPF, IS-IS, BGP) MUST be reported.

3.3. Frame Formats and Sizes

The tests for SR-MPLS will use the Frame characteristics as described in [RFC5695].

Note that [RFC5695] requires exactly a single entry in the MPLS label stack in an MPLS packet. In other words, the depth of the label stack is set to one.

To ensure successful delivery of Layer 2 frames carrying SR-MPLS packets and realistic benchmarking, it is RECOMMENDED to set the media MTU value to the effective maximum frame payload size (payload of 1500 octets for Ethernet).

The number of entries in the label stack MUST be reported. In addition, it MUST be chosen taking into account this condition.
4. Reporting Format

There are new parameters that MUST be replaced or added to the parameters specified in [RFC5695]:

- SR-MPLS Forwarding Operations (PUSH/ NEXT/ CONTINUE).
- Number of Segments considered in the MPLS Label Stack.
- Global SIDs or Local SID forwarding behavior.
- SR Policy headend or endpoint behavior.

5. SR-MPLS Forwarding Benchmarking Tests

This document recommends the same benchmarking tests described in [RFC2544] and [RFC5695] while observing the DUT setup and the traffic setup considerations specific for SR-MPLS as described above. It may require additional benchmarking steps.

5.1. Throughput

This section contains the description of the tests that are related to the characterization of a DUT’s SR-MPLS traffic forwarding throughput.

The list of segments for SR-MPLS is represented as a stack of MPLS labels. There are three distinct operations to be tested: PUSH, NEXT and CONTINUE. These correspond to the three forwarding operations of an MPLS packet: PUSH (or LSP Ingress), SWAP, or POP (or LSP Egress).

5.1.1. Throughput for SR-MPLS PUSH

Objective: To obtain the DUT’s Throughput during PUSH forwarding operation. It is similar to label Push or LSP Ingress forwarding operation, as per [RFC5695]. Non-reserved MPLS label values MUST be used.

Procedure: Same as [RFC5695].

Reporting Format: Same as [RFC5695] but adding the additional parameters specified in Section 4.

5.1.2. Throughput for SR-MPLS NEXT

Objective: To obtain the DUT’s Throughput during NEXT forwarding operation. It is equivalent to MPLS Label Pop or Penultimate Hop
Popping (PHP), as per [RFC5695]. Non-reserved MPLS label values MUST be used.

Procedure: Same as [RFC5695].

Reporting Format: Same as [RFC5695] but adding the additional parameters specified in Section 4.

5.1.3. Throughput for SR-MPLS CONTINUE

Objective: To obtain the DUT’s Throughput during CONTINUE forwarding operation. It is equivalent to MPLS Label Swap or Ultimate Hop Popping (UHP), as per [RFC5695]. Non-reserved MPLS label values MUST be used.

Procedure: Same as [RFC5695].

Reporting Format: Same as [RFC5695] but adding the additional parameters specified in Section 4.

5.2. Latency

Objective: To determine the latency as defined in [RFC5695] for each of the SR-MPLS forwarding operations.

Procedure: Same as [RFC5695].

Reporting Format: Same as [RFC5695] but adding the additional parameters specified in Section 4.

5.3. Frame Loss

Objective: To determine the frame-loss rate (as defined in [RFC5695]) for each of the SR-MPLS forwarding operations of a DUT throughout the entire range of input data rates and frame sizes.

Procedure: Same as [RFC5695].

Reporting Format: Same as [RFC5695] but adding the additional parameters specified in Section 4.

5.4. System Recovery

Objective: To characterize the speed at which a DUT recovers from an overload condition for each of the SR-MPLS forwarding operations.

Procedure: Same as [RFC5695].
Reporting Format: Same as [RFC5695], but adding the additional parameters specified in Section 4.

5.5. Reset

Objective: To characterize the speed at which a DUT recovers from a device or software reset for each of the SR-MPLS forwarding operations.

Procedure: Same as [RFC5695].

Reporting Format: Same as [RFC5695] but adding the additional parameters specified in Section 4.

6. SR Policy: protection performance

[RFC6414] provides common terminology and metrics for benchmarking the performance of protection mechanisms. [RFC6894] provides detailed test cases with different topologies and scenarios that should be considered to effectively benchmark MPLS-FRR protection mechanisms and failover times on the data plane. The same approach can be considered also for Segment Routing protection mechanisms.

An SR Policy can be used for Traffic Engineering (TE), Operations, Administration, and Maintenance (OAM), or Fast Reroute (FRR) reasons. Protection allows that, in the event the interface associated with the Adj-SID is down, the packet can still be forwarded via an alternate path. The use of protection is clearly a policy-based decision that determines, for example, that a PUSH operation is done to forward a packet over a backup path calculated using TI-LFA. There are 2 different protection mechanisms for SR-TE: Segment protection specified in [I-D.ietf-spring-segment-protection-sr-te-paths] and Path protection introduced in [I-D.ietf-spring-segment-routing-policy].

7. Security Considerations

Benchmarking methodologies are limited to technology characterization in a laboratory environment, with dedicated address space and constraints. Special capabilities SHOULD NOT exist in the DUT/SUT specifically for benchmarking purposes. Any implications for network security arising from the DUT/SUT SHOULD be identical in the lab and in production networks. The benchmarking network topology is an independent test setup and MUST NOT be connected to devices that may forward the test traffic into a production network or misroute traffic to the test management network.
There are no specific security considerations within the scope of this document.

8. IANA Considerations

This document has no IANA actions.

9. Acknowledgements

TBD

10. References

10.1. Normative References


10.2. Informative References

[I-D.ietf-spring-segment-protection-sr-te-paths]
Hegde, S., Bowers, C., Litkowski, S., Xu, X., and F. Xu,
"Segment Protection for SR-TE Paths", draft-ietf-spring-
segment-protection-sr-te-paths-02 (work in progress),
January 2022.

[I-D.ietf-spring-segment-routing-policy]
Filsfils, C., Talaulikar, K., Voyer, D., Bogdanov, A., and
P. Mattes, "Segment Routing Policy Architecture",
draft-ietf-spring-segment-routing-policy-19 (work in progress),
March 2022.

Label Switching Architecture", RFC 3031,
DOI 10.17487/RFC3031, January 2001,

[RFC3032] Rosen, E., Tappan, D., Fedorkow, G., Rekhter, Y.,
Farinacci, D., Li, T., and A. Conta, "MPLS Label Stack
Encoding", RFC 3032, DOI 10.17487/RFC3032, January 2001,

H., Shakir, R., Henderickx, W., and J. Tantsura, "OSPF
Extensions for Segment Routing", RFC 8665,
DOI 10.17487/RFC8665, December 2019,

Bashandy, A., Gredler, H., and B. Decraene, "IS-IS
Extensions for Segment Routing", RFC 8667,
DOI 10.17487/RFC8667, December 2019,

Authors’ Addresses

Giuseppe Fioccola
Huawei Technologies
Riesstrasse, 25
Munich  80992
Germany

Email: giuseppe.fioccola@huawei.com

Eduard Vasilenko
Huawei Technologies
17/4 Krylatskaya str.
Moscow  121614
Russia

Email: vasilenko.eduard@huawei.com

Paolo Volpato
Huawei Technologies
Via Lorenteggio, 240
Milan  20147
Italy

Email: paolo.volpato@huawei.com

Luis Miguel Contreras Murillo
Telefonica
Spain

Email: luismiguel.contrerasmurillo@telefonica.com
Benchmarking Methodology for IPv6 Segment Routing
draft-vfv-bmwg-srv6-bench-meth-01

Abstract

This document defines a methodology for benchmarking Segment Routing (SR) performance for Segment Routing over IPv6 (SRv6). It builds upon [RFC2544], [RFC5180] and [RFC8402].

Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at https://datatracker.ietf.org/drafts/current/.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on September 8, 2022.

Copyright Notice

Copyright (c) 2022 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust’s Legal Provisions Relating to IETF Documents
1. Introduction

Segment Routing (SR), defined in [RFC8402], leverages the source routing paradigm. The headend node steers a packet through an SR Policy [I-D.ietf-spring-segment-routing-policy], instantiated as an ordered list of segments. A segment, referred to by its Segment Identifier (SID), can have a semantic local to an SR node or global within an SR domain. SR supports per-flow explicit routing while maintaining per-flow state only at the ingress nodes to the SR domain.
However, there is no standard method defined to compare and contrast the foundational SR packet forwarding capabilities of network devices. This document aims to extend the efforts of [RFC1242] and [RFC2544] to SR network.

The SR architecture can be instantiated on two data-plane: SR over MPLS (SR-MPLS) and SR over IPv6 (SRv6). This document is limited to SRv6.

SR can be applied to the IPv6 architecture with a new type of routing header called the SR Header (SRH) [RFC8754]. An instruction is associated with a segment and encoded as an IPv6 address. An SRv6 segment is also called an SRv6 SID. An SR Policy is instantiated as an ordered list of SRv6 SIDs in the routing header. The active segment is indicated by the Destination Address (DA) of the packet.

For Segment Routing, PUSH, NEXT, and CONTINUE are operations applied by the forwarding plane.

PUSH consists of the insertion of a segment at the top of the segment list. In SRv6, the top of the segment list is represented by the first segment in the SRH.

NEXT consists of the inspection of the next segment. The active segment is completed and the next segment becomes active. In SRv6, NEXT is implemented as the copy of the next segment from the SRH to the destination address of the IPv6 header.

CONTINUE happens when the active segment is not completed; hence, it remains active. In SRv6, the CONTINUE operation is the plain IPv6 forwarding action of a regular IPv6 packet according to its destination address.

[RFC5180] provides benchmarking methodology recommendations that address IPv6-specific aspects, such as evaluating the forwarding performance of traffic containing extension headers.

The purpose of this document is to describe a methodology specific to the benchmarking of Segment Routing. The methodology described is a complement for [RFC5180].

2. SRv6 Forwarding

In IPv6, a Prefix-SID is allocated in the form of an IPv6 address. For the IPv6 data plane, a new type of IPv6 Routing Extension Header, called Segment Routing Header (SRH) has been defined. The SRH contains the Segment List as an ordered list of IPv6 addresses: each address in the list is a SID. A dedicated field, referred to as
Segments Left, is used to maintain the pointer to the active SID of the Segment List.

There are three different categories of nodes that may be involved in segment routing networks.

The SR source node is the headend node and steers a packet into an SR Policy. It can be a host originating an IPv6 packet or an SR domain ingress router encapsulating a received packet into an outer IPv6 packet and inserts the SRH in the outer IPv6 header. It sets the first SID of the SR Policy as IPv6 Destination Address of the packet.

The SR transit node forwards packets destined to a remote segment as a normal IPv6 packet on the basis of the IPv6 destination address, because the IPv6 destination address does not locally match with a segment. Indeed, according to [RFC8200] the only node allowed to inspect the Routing Extension Header (and therefore the SRH) is the node corresponding to the destination address of the packet.

The SR segment endpoint node receives packets whose IPv6 destination address is locally configured as a segment. It creates Forwarding Information Base (FIB) entries for its local SIDs. For each SR packet, it inspects the SRH and replaces the IPv6 destination address with the new active segment.

The operations applied by the SRv6 packet processing are different at the SR source, Transit and SR segment endpoint nodes.

The processing of the SR source node corresponds to the sequence of the insertion of the SRH, composed of SIDs stored in reverse order, and setting of the IPv6 Destination Address as first SID of the SR Policy. It can be performed by encapsulating a packet into an outer IPv6 packet with an SRH. Another possibility is to perform the insertion of an SRH as a new header between the IPv6 header and the Next Header (e.g. the Transport Layer Header, TCP or UDP). This option only applies to IPv6 packets and it is especially suited in case the source host is acting as headend node.

The processing of the SR segment endpoint node corresponds to the detection of the new active segment, which is the next segment in the Segment List and the related modification of the IPv6 destination address of the outer IPv6 header. Then packets are forwarded on the basis of the IPv6 forwarding table.

The processing of the SR transit node corresponds to normal forwarding of the packets containing the SR header. In SRv6 the transit nodes do not need to be SRv6 aware, as every IPv6 router can act as an SRv6 transit node since any IPv6 node will maintain a plain
IPv6 FIB entry for any prefix, no matter if the prefix represents a segment or not.

[I-D.ietf-spring-segment-routing-policy] specifies the concepts of SR Policy and steering into an SR Policy. The header of a packet steered in an SR Policy is augmented with the ordered list of segments associated with that SR Policy. SR Policy state is instantiated only on the headend node, that steers a flow into an SR Policy. Indeed intermediate and endpoint nodes do not require any state to be maintained. SR Policies can be instantiated on the headend dynamically and on demand basis. Moreover, signaling can be used in the case of a controller based deployment. For all these reasons, SR Policies scale better than traditional TE mechanisms.

In addition to the basic SRv6 packet processing, the SRv6 Network Programming model [RFC8986] describes a set of functions that can be associated to segments and executed in a given SRv6 node.

Examples of such functions are described in [RFC8986], but, in practice, any behavior and function can be associated to a local SID in a node, in order to apply any special processing on the packet. Obviously, the definition of a standardized set of segment routing functions facilitates the deployment of SR domains with interoperable equipment from multiple vendors.

According to [RFC8986], 128 bit SID can be logically split into three fields and interpreted as LOCATOR:FUNCTION:ARGS (in short LOC:FUNCT:ARG) where LOC includes the L most significant bits, FUNCT the following F bits and ARG the remaining A bits, where 128=L+F+A. The LOC corresponds to an IPv6 prefix (for example with a length of 48, 56 or 64 bits) that can be distributed by the routing protocols and provides the reachability of a node that hosts a number of functions. All the different functions residing in a node have a different FUNCT code, so that their SIDs will be different. The ARG bits are used to provide information (arguments) to a function. From the routing point of view, the solution is scalable, as a single prefix is distributed for a node, which implements a potentially large number of functions and related arguments.

3. Test Methodology

3.1. Test Setup

The Device Under Test (DUT) is connected to the test ports on the test tool according to [RFC2544].

The test topology recommended for the SRv6 performance evaluation are the same as IPv6 and are described in [RFC5180] and [RFC2544], in
both single-port and multi-port scenarios. Single-port testing measures per-interface forwarding performance, while multi-port testing measures the scalability of forwarding performance across the entire platform.

3.2. IGP and BGP Support

It is RECOMMENDED that all of the ports on the DUT and test tool support a Segment Routing extensions for dynamic Interior Gateway Protocol (IGP) for routing such as IS-IS [I-D.ietf-lsr-isis-srv6-extensions] and OSPF [I-D.ietf-lsr-ospfv3-srv6-extensions] as well as Border Gateway Protocol (BGP) [I-D.ietf-bess-srv6-services].

As specified in [RFC8402], in the context of an IGP-based distributed control plane, two topological segments are defined: the IGP-Adjacency segment and the IGP-Prefix segment; while, in the context of a BGP-based distributed control plane, two topological segments are defined: the BGP peering segment and the BGP-Prefix segment.

The distribution method that is used (e.g. OSPF, IS-IS, BGP) MUST be reported.

3.3. Frame Formats and Sizes

The tests for SRv6 will use the Frame characteristics as described in [RFC5180].

As specified in [RFC5180], for Ethernet, the following frame sizes SHOULD be used for benchmarking over this media type: 64, 128, 256, 512, 1024, 1280, and 1518 bytes. Note that the recommended 1518-byte frame size represents the maximum size of an untagged Ethernet frame. A frame size commonly used in operational environments is 1522 bytes, the max length for a VLAN-tagged frame.

3.4. Protocol Addresses

IANA reserved an IPv6 address block for use with IPv6 benchmark testing (see [RFC5180]). IPv6 source and destination addresses for the test streams SHOULD belong to the IPv6 range assigned by IANA.

3.5. Traffic with SRH

The extension header chain recommended in [RFC5180] for testing is: Routing header (24-32 bytes), Destination options header (8 bytes), Fragment header (8 bytes). This was considered a real-life extension-header chain but it does not fit well for SRv6.
The length of the SRH is \((n \times 16 + 8)\) bytes, where \(n\) is the number of segments. So, for most of the SRv6 application the recommendation of [RFC5180] is not enough. In addition, it is worth mentioning that the length of SRv6 packets is increased in Topology Independent Loop-Free Alternate (TI-LFA) Fast Reroute (FRR), binding SID, and microloop avoidance scenarios.

For SRv6, the extension header chain characteristics and length that are used MUST be reported and the DUT MUST traverse the chain of extension headers, so the impact on performance can be observed.

4. Reporting Format

There are new parameters that MUST be added to the parameters specified in [RFC5180] and [RFC2544]:

- SRv6 types of nodes.
- Number of Segments considered in the SRH.
- Extension header chain (including SRH) characteristics and length.
- Global SIDs or Local SID forwarding behavior.
- SR Headend or Endpoint Behaviors eventually associated with a SID, as specified in [RFC8986].

For the sake of completeness, the following Figure 1 reports all the SR Headend or Endpoint Behaviors, as defined in [RFC8986]. But, in most cases, it may not be necessary to test all the services and it is possible to select a subset.

<table>
<thead>
<tr>
<th>H.Encaps</th>
<th>SR Headend with Encapsulation in an SR Policy</th>
</tr>
</thead>
<tbody>
<tr>
<td>H.Encaps.Red</td>
<td>H.Encaps with Reduced Encapsulation</td>
</tr>
<tr>
<td>H.Encaps.L2</td>
<td>H.Encaps Applied to Received L2 Frames</td>
</tr>
<tr>
<td>End</td>
<td>Endpoint</td>
</tr>
<tr>
<td>End.X</td>
<td>Endpoint with L3 cross-connect</td>
</tr>
<tr>
<td>End.T</td>
<td>Endpoint with specific IPv6 table lookup</td>
</tr>
<tr>
<td>End.DX6</td>
<td>Endpoint with decapsulation and IPv6 cross-</td>
</tr>
<tr>
<td>End.DX4</td>
<td>Endpoint with decapsulation and IPv4 cross-connect</td>
</tr>
<tr>
<td>------------------</td>
<td>-----------------------------------------------</td>
</tr>
<tr>
<td>End.DT6</td>
<td>Endpoint with decapsulation and specific IPv6 table lookup</td>
</tr>
<tr>
<td>End.DT4</td>
<td>Endpoint with decapsulation and specific IPv4 table lookup</td>
</tr>
<tr>
<td>End.DT46</td>
<td>Endpoint with decapsulation and specific IP table lookup</td>
</tr>
<tr>
<td>End.DX2</td>
<td>Endpoint with decapsulation and L2 cross-connect</td>
</tr>
<tr>
<td>End.DX2V</td>
<td>Endpoint with decapsulation and VLAN L2 table lookup</td>
</tr>
<tr>
<td>End.DT2U</td>
<td>Endpoint with decapsulation and unicast MAC L2 table lookup</td>
</tr>
<tr>
<td>End.DT2M</td>
<td>Endpoint with decapsulation and L2 table flooding</td>
</tr>
<tr>
<td>End.B6.Encaps</td>
<td>Endpoint bound to an SRv6 Policy with encapsulation</td>
</tr>
<tr>
<td>End.BM</td>
<td>Endpoint bound to an SR-MPLS Policy</td>
</tr>
</tbody>
</table>

Figure 1: SR Policy Headend and Endpoint Behaviors

5. SRv6 Forwarding Benchmarking Tests

This document recommends the same benchmarking tests described in [RFC2544] and [RFC5180] while observing the DUT setup and the traffic setup considerations described above. Indeed, the specificities of SRv6, for example the SRH processing, require additional benchmarking steps.
5.1. Throughput

This section contains the description of the tests that are related to the characterization of a DUT’s SRv6 traffic forwarding throughput.

The list of segments for SRv6 is represented as a list of IPv6 addresses, included in the SRH. There are three distinct types of nodes that are involved in segment routing networks.

5.1.1. Throughput of a Source Node

Objective: To obtain the DUT’s Throughput during the packet processing of a Source Node. It is when the Source SR node, which corresponds to the headend node, encapsulates a received packet into an outer IPv6 packet and inserts the SR Header (SRH) as a Routing Extension Header in the outer IPv6 header. The Segment List in the SRH is composed of SIDs and the Source SR node sets the first SID of the SR Policy as IPv6 Destination Address of the packet.

Procedure: Same as [RFC5180].

Reporting Format: Same as [RFC5180] but adding the additional parameters specified in Section 4.

5.1.2. Throughput of a Segment Endpoint Node

Objective: To obtain the DUT’s Throughput during the packet processing of a Segment Endpoint Node. It is when the SR Segment Endpoint node receives packets whose IPv6 destination address is locally configured as a segment. The SR Segment Endpoint node inspects the SR header: it detects the new active segment, i.e. the next segment in the Segment List, modifies the IPv6 destination address of the outer IPv6 header and forwards the packet on the basis of the IPv6 forwarding table.

Procedure: Same as [RFC5180].

Reporting Format: Same as [RFC5180] but adding the additional parameters specified in Section 4.

5.1.3. Throughput of a Transit Node

Objective: To obtain the DUT’s Throughput during the packet processing of a Transit Node. It is when a Transit node forwards the packet containing the SR header as a normal IPv6 packet because the IPv6 destination address does not locally match with a segment.
Procedure: Same as [RFC5180].

Reporting Format: Same as [RFC5180] but adding the additional parameters specified in Section 4.

5.2. Latency

Objective: To determine the latency as defined in [RFC5180] for each of the SRv6 forwarding operations.

Procedure: Same as [RFC5180].

Reporting Format: Same as [RFC5180] but adding the additional parameters specified in Section 4.

5.3. Frame Loss

Objective: To determine the frame-loss rate (as defined in [RFC5180]) for each of the SRv6 forwarding operations of a DUT throughout the entire range of input data rates and frame sizes.

Procedure: Same as [RFC5180].

Reporting Format: Same as [RFC5180] but adding the additional parameters specified in Section 4.

5.4. System Recovery

Objective: To characterize the speed at which a DUT recovers from an overload condition for each of the SRv6 forwarding operations.

Procedure: Same as [RFC5180].

Reporting Format: Same as [RFC5180] but adding the additional parameters specified in Section 4.

5.5. Reset

Objective: To characterize the speed at which a DUT recovers from a device or software reset for each of the SRv6 forwarding operations.

Procedure: Same as [RFC5180].

Reporting Format: Same as [RFC5180] but adding the additional parameters specified in Section 4.
6. SR Policy: protection performance

[RFC6414] provides common terminology and metrics for benchmarking the performance of protection mechanisms.

An SR Policy can be used for Traffic Engineering (TE), Operations, Administration, and Maintenance (OAM), or Fast Reroute (FRR) reasons. Protection allows that, in the event the interface associated with the Adj-SID is down, the packet can still be forwarded via an alternate path. The use of protection is clearly a policy-based decision that determines, for example, that the packet processing by the source node is done to forward a packet over a backup path calculated using TI-LFA. There are 2 different protection mechanisms for SR-TE: Segment protection specified in [I-D.ietf-spring-segment-protection-sr-te-paths] and Path protection introduced in [I-D.ietf-spring-segment-routing-policy].

7. Security Considerations

Benchmarking methodologies are limited to technology characterization in a laboratory environment, with dedicated address space and constraints. Special capabilities SHOULD NOT exist in the DUT/SUT specifically for benchmarking purposes. Any implications for network security arising from the DUT/SUT SHOULD be identical in the lab and in production networks. The benchmarking network topology is an independent test setup and MUST NOT be connected to devices that may forward the test traffic into a production network or misroute traffic to the test management network.

There are no specific security considerations within the scope of this document.

8. IANA Considerations

This document has no IANA actions.

9. Acknowledgements

TBD

10. References

10.1. Normative References


10.2. Informative References

[I-D.ietf-isis-srv6-extensions]

[I-D.ietf-ospfv3-srv6-extensions]

[I-D.ietf-spring-segment-protection-sr-te-paths]

[I-D.ietf-spring-segment-routing-policy]


Authors’ Addresses

Giuseppe Fioccola
Huawei Technologies
Riesstrasse, 25
Munich  80992
Germany

Email: giuseppe.fioccola@huawei.com

Eduard Vasilenko
Huawei Technologies
17/4 Krylatskaya str.
Moscow  121614
Russia

Email: vasilenko.eduard@huawei.com
Paolo Volpato
Huawei Technologies
Via Lorenteggio, 240
Milan  20147
Italy
Email: paolo.volpato@huawei.com

Luis Miguel Contreras Murillo
Telefonica
Spain
Email: luismiguel.contrerasmurillo@telefonica.com