Rate-Limited Issuance

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- Rate-limiting is a common part of fraud prevention and anonymous access
- It also often relies on tracking cookies or client IP addresses
 - A common way to implement this is with "token buckets"

Token buckets Status quo rate-limiting



Token buckets Status quo rate-limiting

Token Bucket 1. Identity request bucket 2. Increment count by number of tokens Token Replenish Mapping 1250123010

Resource Request



	Token Count
339	N (N + T)

Yes, service request!



No, don't service request



Token buckets Status quo rate-limiting



Resource Request



Token Bucket

2. Decrement count associated with bucket 3. Process if non-zero, otherwise discard

	Token Count
339	N (N - 1) > 0?

Yes, service request!

No, don't service request



Why Privacy Pass?

- - Proxies
 - VPNs
 - Shared IPs on public networks
- A basic Privacy Pass token isn't always enough
 - Attests to the fact that a device or user passed some check
 - metered paywall)
 - Degenerates to blocking access

Existing rate-limiting schemes break down when clients have more privacy (shared rate-limiting buckets)

Does not prevent that legitimate device or user doing too many actions (click farm, or abuse of



Private Token variants

Basic Token (OPRF, RSA Blind Signature)

Attests to user/device legitimacy

Replaces captcha for improving confidence in user



Rate-Limited Token (RSA Blind Signature)

Attests to user/device legitimacy + access rate below threshold

Adds mitigations against a device in a click farm

Allows metered paywall access

Rate-limited tokens *extend* the basic issuance protocol with new properties:

- Attester maintains counters for client + anonymized origin 1.
 - Attesters learn stable mapping between per-client secret and per-origin secret, without learning only per-origin information
- 2. Issuer provides a rate limit to enforce when issuing tokens
 - Issuers learn origin associated with a token challenge, encrypted with HPKE
- Attesters fail requests if the per-origin rate limit is exceeded 3.

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- Attesters learn stable mapping between per-client secret and per-origin secret, without learning only per-origin information
- 2. Issuer provides a rate limit

- Issuers learn origin associated with a token challenge, encrypted with HPKE
- 3. Attesters fail requests if the per-origin rate limit is exceeded

This is the main challenge for the protocol

Attester state

The "token buckets" used for rate limiting now are "private token buckets" maintained on the attester

- e.g., F(client secret, origin secret)
- Attester uses mapping as index into data structure tracking per-client state



A stable mapping is a deterministic function between per-client and per-origin information,

The mapping is used to enforce rate limits based on individual clients for individual origins

Count
Ν



Stable Mappings and Rate Limits



Attester

Compute stable mapping, decrement count, compare against origin limit, accept or reject response accordingly







Stable Mappings and Rate Limits



Attester req Compute stable mapping, decrement count, resp, L_{origin} compare against origin limit, accept or reject response accordingly $1234 \neq F(\text{client}, \text{origin})$ Count . . . **№** -> N-1 . . . $N-1 < L_{\text{origin}}$





Current Status

mapping

Requires split deployment model for meaningful privacy

Builds on signature schemes with key blinding for private computing the stable

- Several interoperable implementations exist and security analysis is underway

Is the WG interested in adopting this draft?

Backup

Rate-limited tokens extend the basic issuance protocol with new properties:

1. Attester maintains counters for client + anonymized origin

- without learning only per-origin information
- 2. Issuer provides a rate limit
- 3. Attesters fail requests if the per-origin rate limit is exceeded

Attesters learn stable mapping between per-client secret and per-origin secret,

Can we use an OPRF to compute this?...

Issuers learn origin associated with a token challenge, encrypted with HPKE

An OPRF protocol computes F(k, x) for per-origin k and per-client x



Clients can encrypt the origin identifier under the Issuer's public key



An Attester can relay the encrypted origin name and complete the OPRF



... Attester can perform a dictionary attack to learn F(k, x)



Rate-limited tokens extend the basic issuance protocol with new properties:

1. Attester maintains counters for client + anonymized origin

- without learning only per-origin information
- 2. Issuer provides a rate limit

Issuers learn origin assoc

An OPRF alone isn't sufficient because of dictionary attacks. Computation of the mapping requires proof of ownership for the per-client secret.

3. Attesters fail requests if the per-origin rate limit is exceeded

Attesters learn stable mapping between per-client secret and per-origin secret,

HPKE

Signature Scheme with Key Blinding

BlindPublicKey and UnblindPublicKey: Given public key and secret blind, produce blinded public key

BlindKeySign: Sign message with secret key and secret blind

Verify(BlindPublicKey(pkS, skB), msg, BlindKeySign(skS, skB, msg)) = true

Draft specification: https://datatracker.ietf.org/doc/draft-dew-cfrg-signature-key-blinding/

- Extend digital signature schemes with two functionalities for signing requests

UnblindPublicKey(BlindPublicKey(pkS, skB), skB) = pkS

Detour: Signature Scheme with Key Blinding

Use signature public key blinding to compute an OPRF*

and let skR be a random client-generated blind per request

F(skC, skO) = Hash(pkC, Blind(pkC, skO))

Close -- but not identical -- to the OPRF construction in draft-irtf-cfrg-voprf

* Security analysis is underway and results will be published in before IETF 114

- Let skC (skO) be the per-Client (per-Origin) secret with public key pkC (pkO),

 - = Hash(pkC, Unblind(Blind(Blind(pkC, skR), skO), skR))





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Drop request



