

draft-piriaux-tcpls

TCPLS: Modern Transport Services with TCP and TLS

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Content

- Introduction
- Using TLS for transport protocol extensibility
- Opportunities for the transport stack
- The TCPLS protocol
- Conclusion

The design of MPTCP

- In 2009, the mptcp WG formed with an initial design involving TCP Options
- In 2013, v0 shipped and enabled
 - Bandwidth aggregation of several TCP subflows
 - Failover in case of network failure
 - Backwards compatibility with TCP
- MPTCP has several issues
 - Address exchange is not secure, improved in MPTCP v1
 - TCP is prone to middlebox interference
 - Can be difficult to implement
 - 7-year journey from specification to mainline Linux

The design of QUIC

- In 2016, the quic WG formed to design an UDP-based transport protocol
- In 2021, QUIC v1 shipped and enabled:
 - Stream multiplexing
 - Connection migration, failover
- TLS secures most of the QUIC header and all QUIC payloads
- QUIC can be implemented in user-space and shipped with applications

Using TLS for transport protocol extensibility

- TLS is the most used protocol atop TCP
- TLS version 1.3 used encryption to extend the protocol
 - Encrypted TLS records and Encrypted Extensions allows securely exchanging control and application data
- TCP support in the network and in operating systems remains wider
- **Given the ubiquity of TLS, can we provide new transport services with TCP and TLS ?**

Opportunities for the transport stack

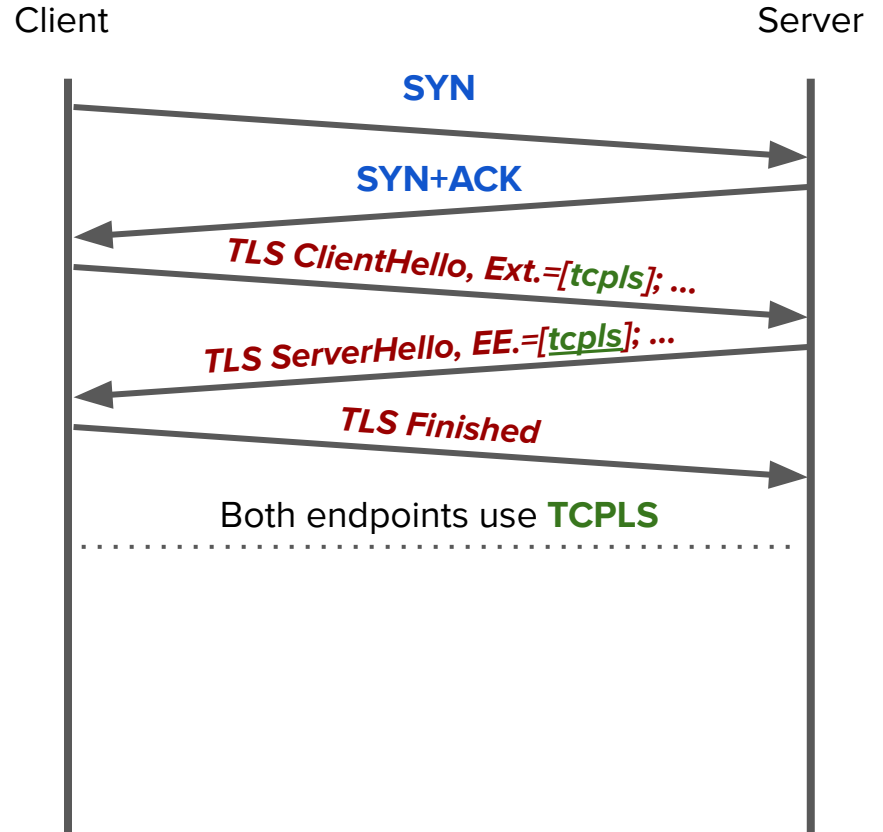
- Build an encrypted transport protocol
 - Stream multiplexing
 - App-chosen HoL blocking resilience
 - Connection Migration
 - Based on app triggers and network conditions
 - Multipath
 - Scheduling at the TLS record level
- More efficient than the HTTP/2+TLS+MPTCP stack
 - Built on a strict layering assumption
- Clean slate for other transport extensions

The TCPLS protocol

- Session establishment
- Exchanging application and control data
- Adding TCP connections
- Record acknowledgements
- Modern Transport Services
 - Stream multiplexing
 - Failover
 - Bandwidth aggregation

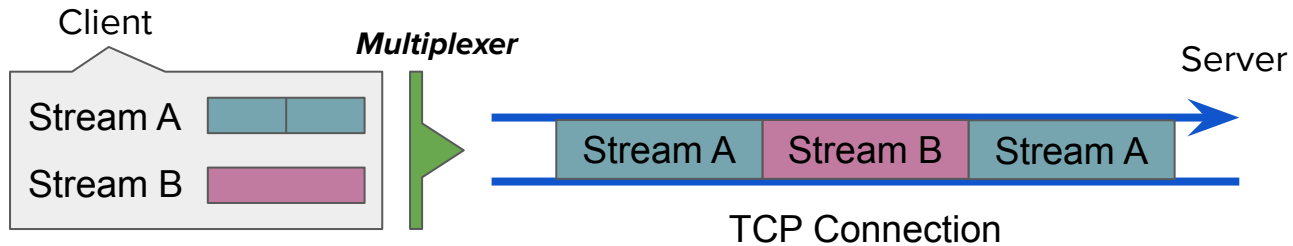
Session establishment

- TCPLS does not modify the TCP and TLS handshake
- *tcpls* is a TLS Extension indicating the support of TCPLS
- Compatible with TCP TFO and TLS 0-RTT Handshake



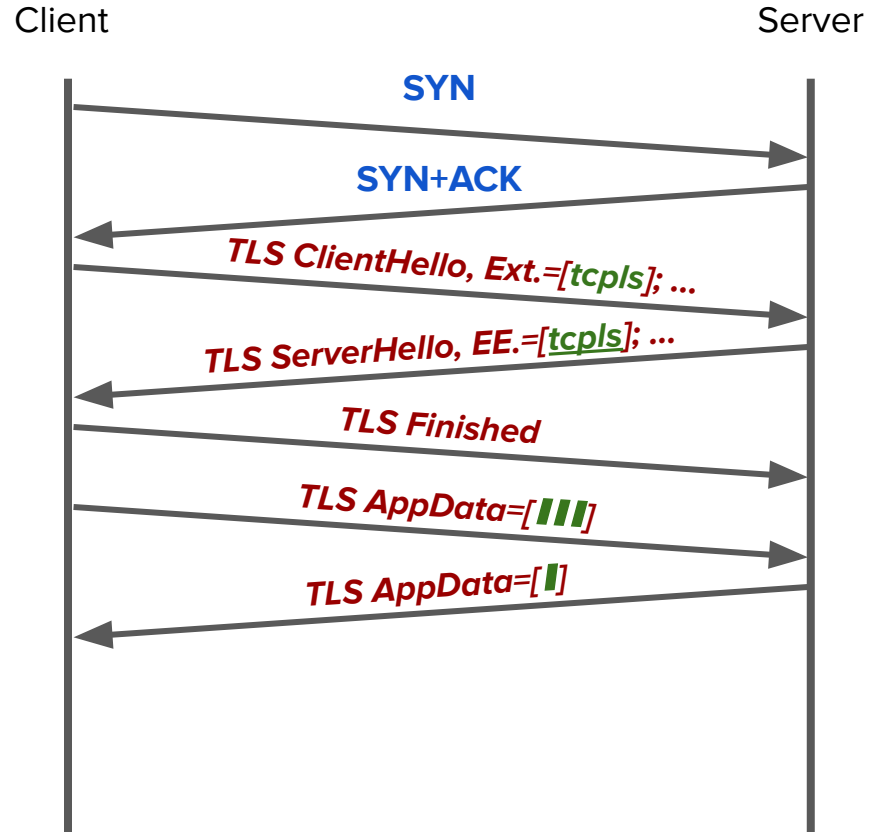
Stream multiplexing

- Streams provide concurrent bytestreams to applications
- TCPLS manages the streams and multiplexes them



Exchanging data

- Application and control data can then be sent in TLS encrypted records using **TCPLS frames**
- **Frames** compose **TLS records**



Example: A TLS record containing a TCPLS Stream frame



TLS Ciphertext header




TLS Encrypted record

- TCPLS
- TLS Ciphertext
- TLS Cleartext

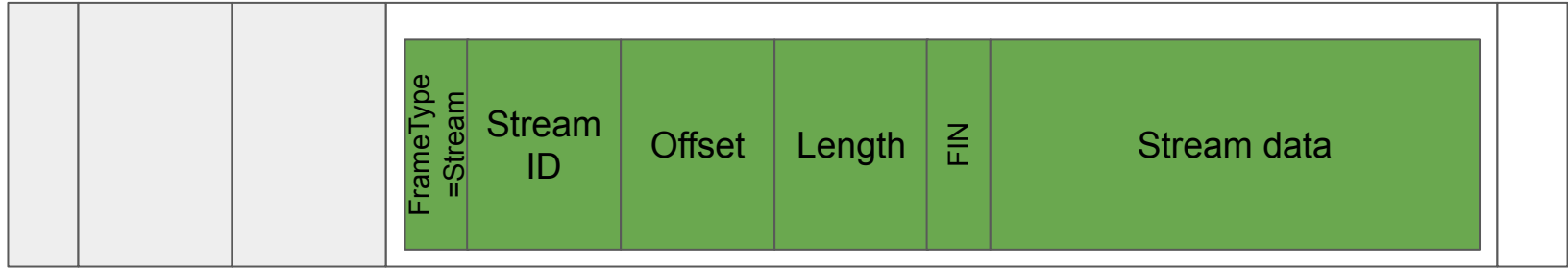
Example: A TLS record containing a Stream frame






TLS Application Data record

-  TCPLS
-  TLS Ciphertext
-  TLS Cleartext

Example: A TLS record containing a Stream frame

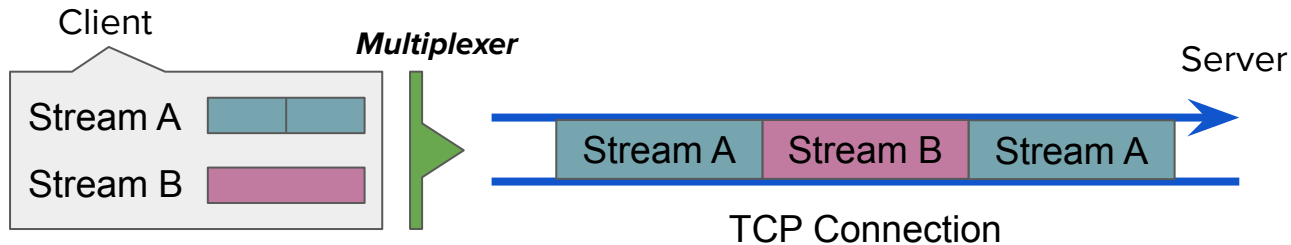


Stream frame

-  TCPLS
-  TLS Ciphertext
-  TLS Cleartext

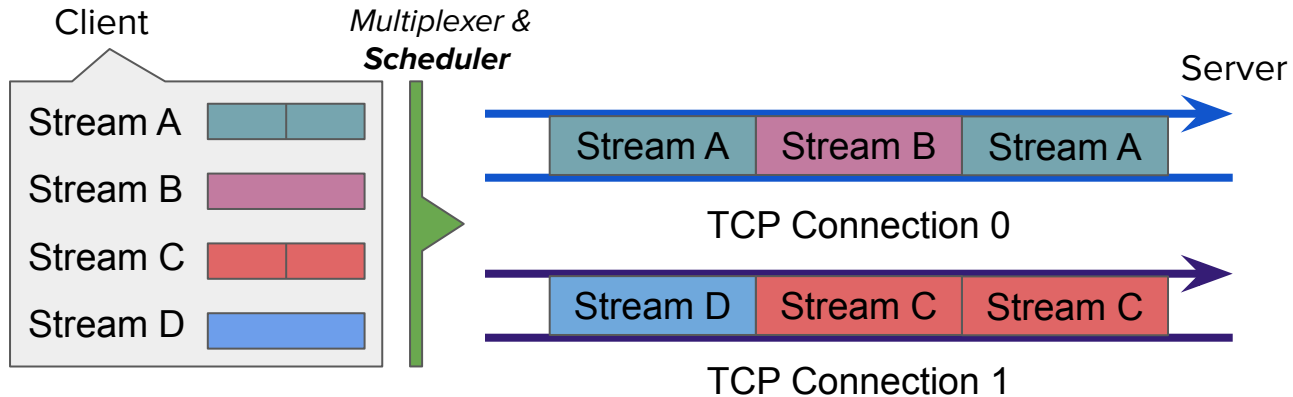
Stream multiplexing

- Streams provide concurrent bytestreams to applications
- TCPLS manages the streams and multiplexes them
- Streams multiplexed on a single connection are subject to HoL blocking



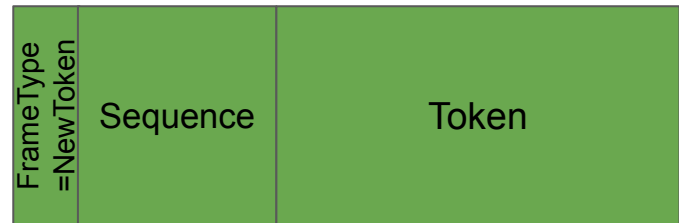
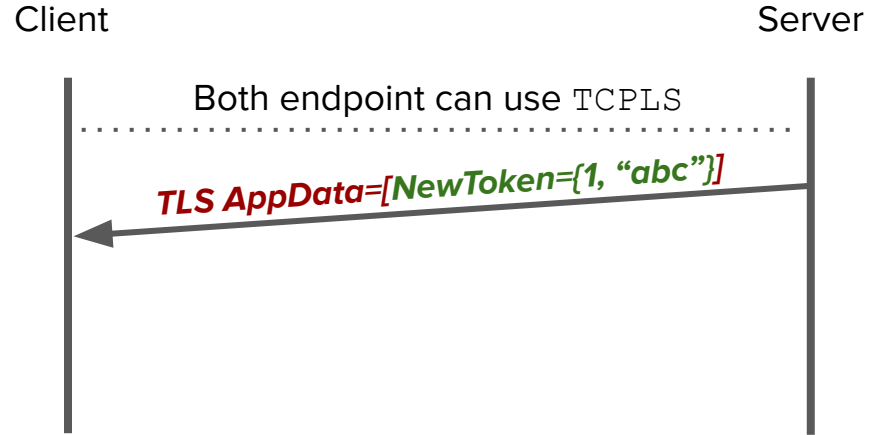
Stream multiplexing

- TCPLS manages TCP connections and schedules the TLS records
- By mapping streams to connections, the app choose the streams it wants to protect, and the ones that are bound together



Adding TCP connections

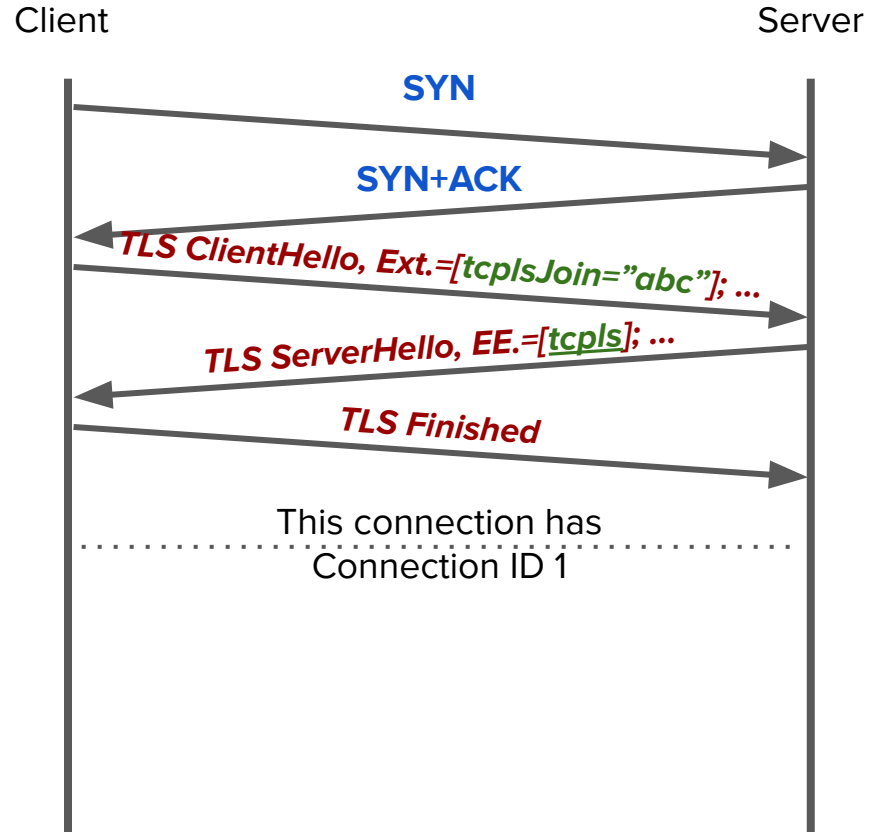
- Server gives tokens to the client
- Each token can be used by the client to open and join an additional TCP connection
- Server can limit the connections by limiting the tokens
- The Sequence number of the Token becomes the Connection ID



New Token frame

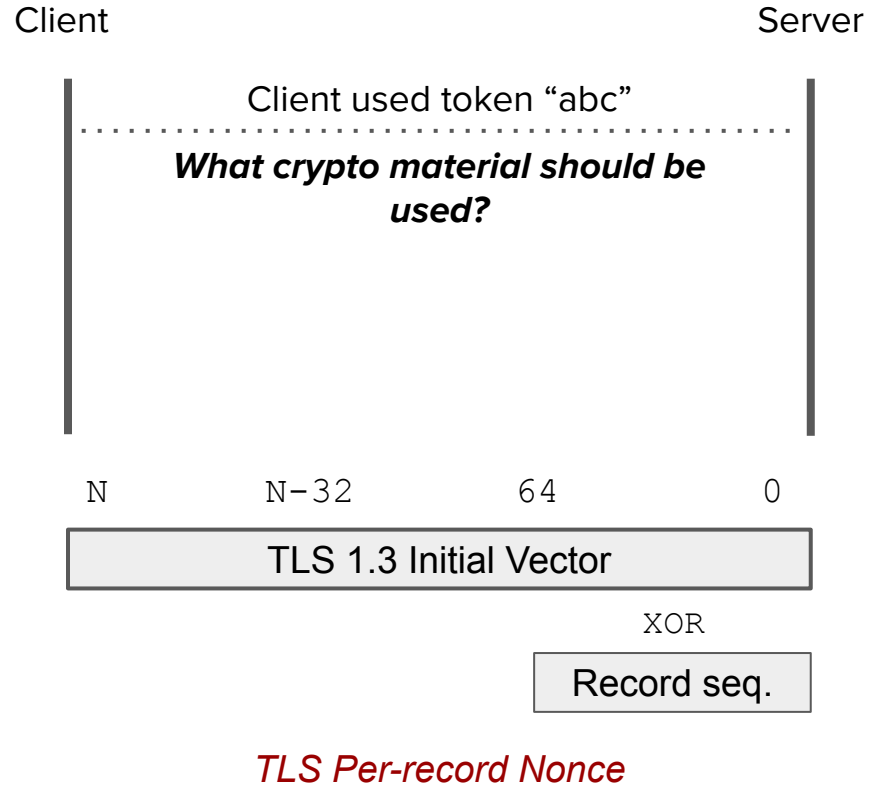
Adding TCP connections

- The client put the token in the *tcplsJoin* TLS Extension
- The server validates the token and joins the TCP connection to the session



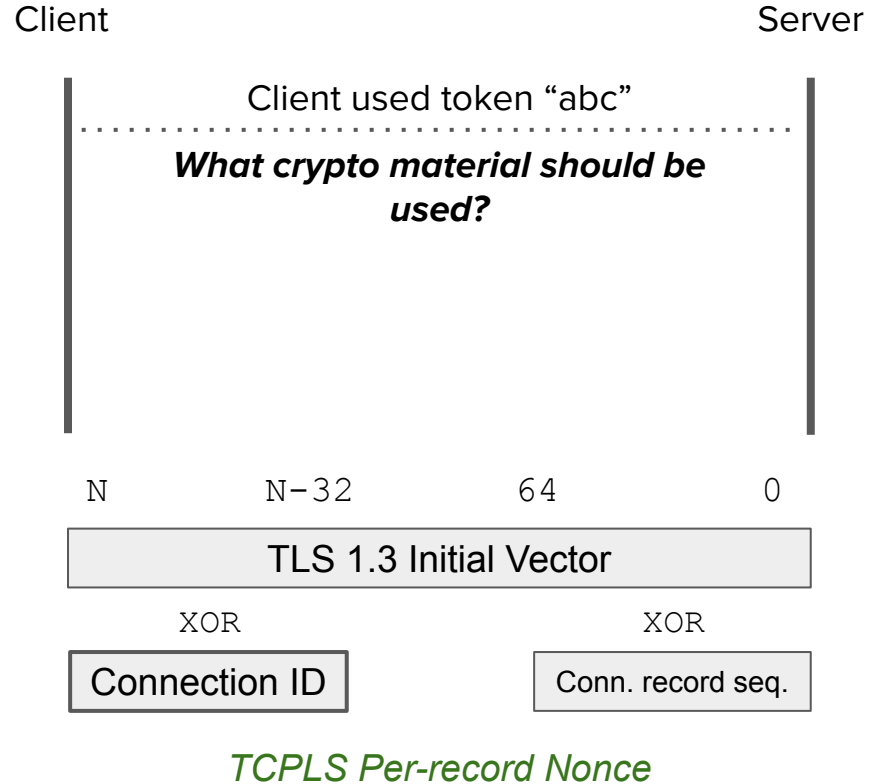
Adding TCP connections

- Each TLS record is encrypted with a unique nonce
- Record sequence is kept implicit
- The record sequence cannot be shared among TCP connections
- **We do not want to do a full TLS handshake, which is costly**



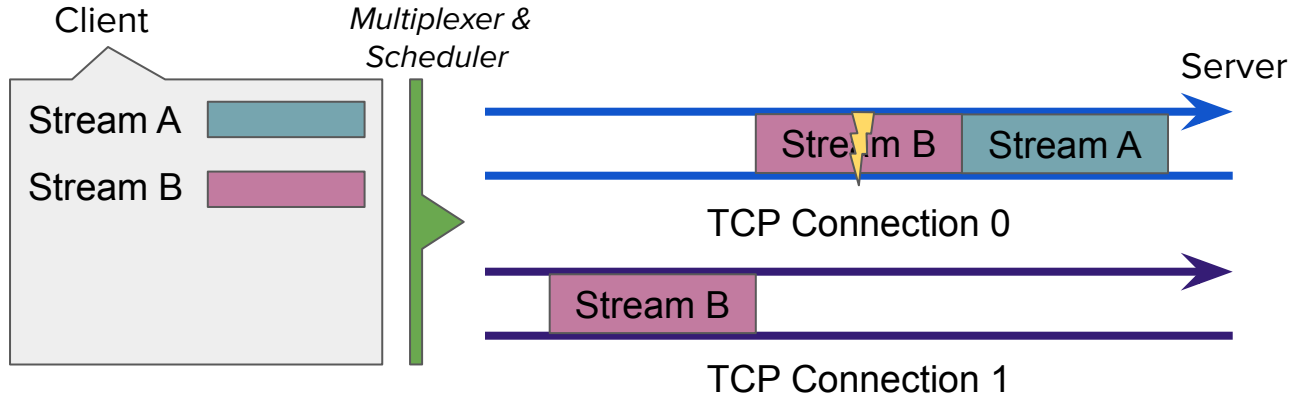
Adding TCP connections

- Each TLS record is encrypted with a unique nonce
- Record sequence is kept implicit
- The record sequence cannot be shared among TCP connections
- **We XOR the Connection ID to the nonce and add a per-connection record sequence**

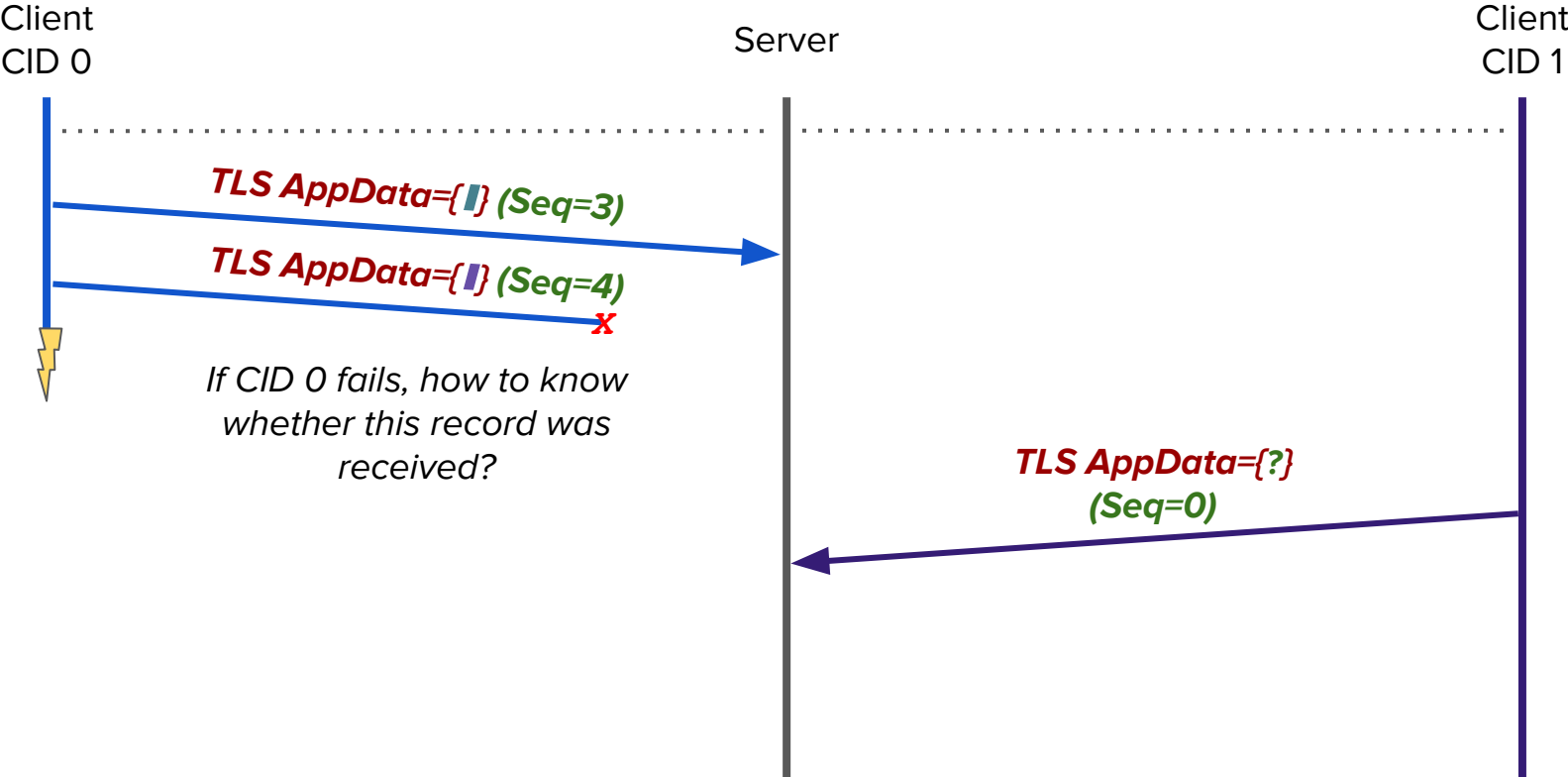


Failover

- Endpoints can reinject frames from lost records onto other TCP connections
- They know which records have been received

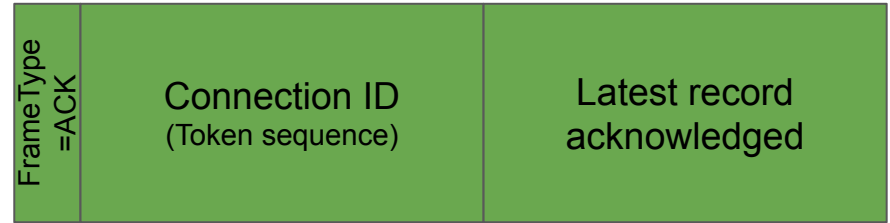


Record acknowledgements



Record acknowledgements

- Each record is identified by:
 - its TLS record sequence number
 - the Connection ID (=Token sequence number) it was sent on
- ACK frame indicates the sequence number of the latest record received over a connection
- As TCP delivers data in sequence, only cumulative ACKs are needed



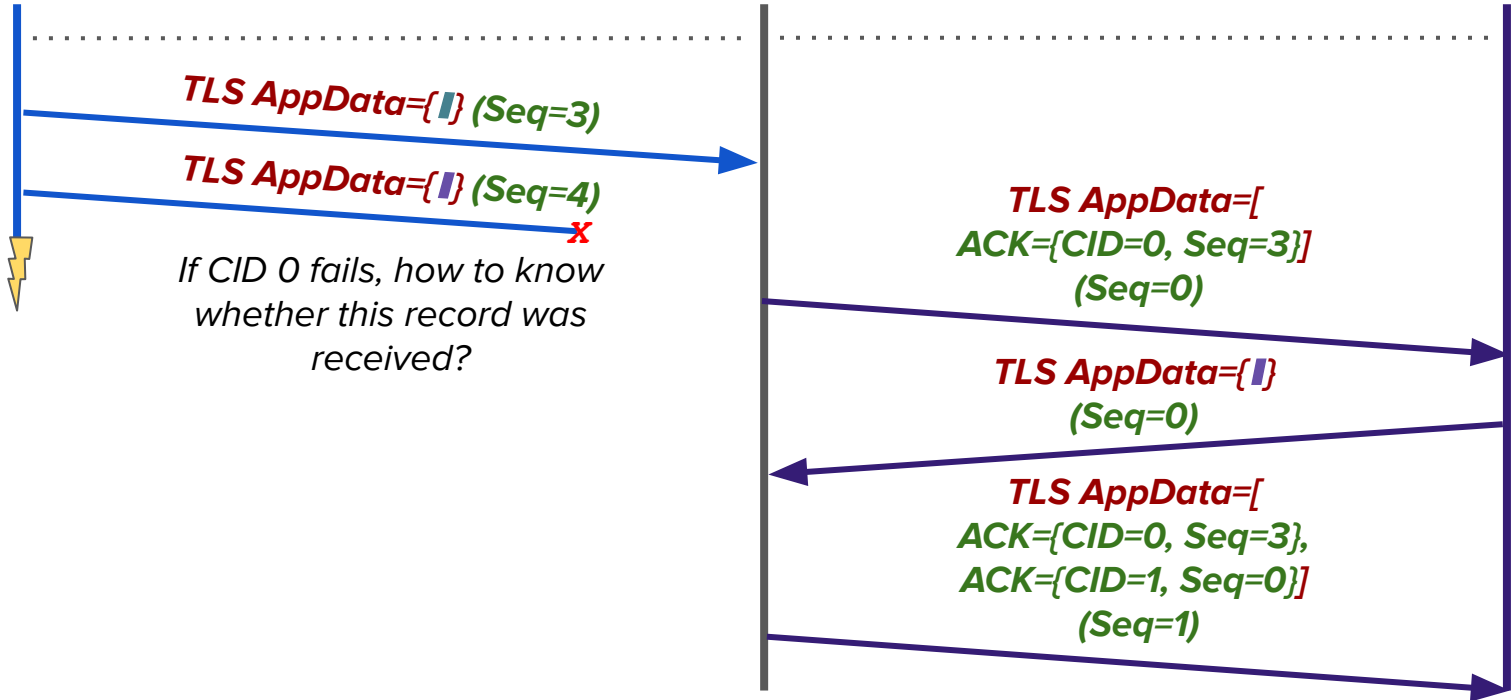
ACK frame

Record acknowledgements

Client
CID 0

Server

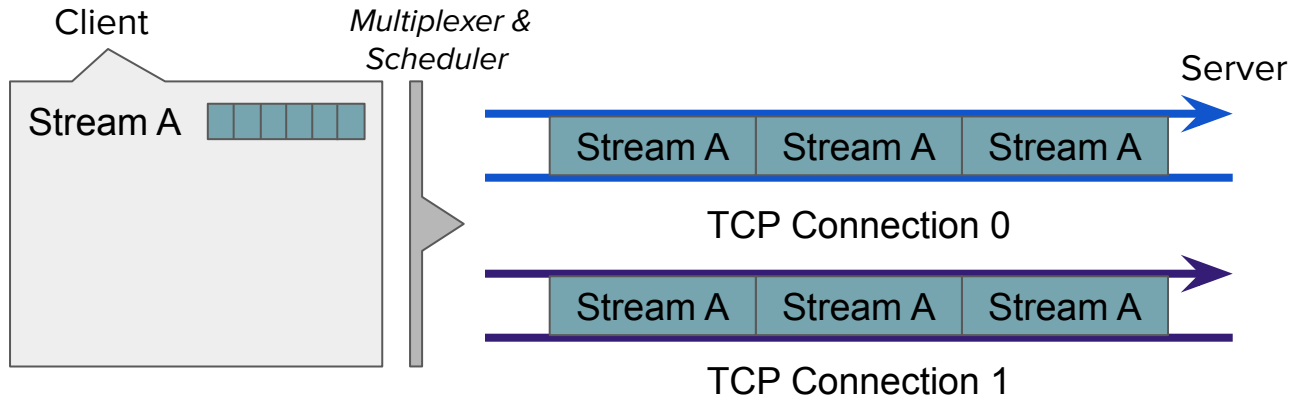
Server
CID 1



If CID 0 fails, how to know whether this record was received?

Bandwidth aggregation

- Endpoints can send Stream frames of a given stream on several TCP connections, benefiting from bandwidth aggregation



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- It describes the protocol presented here
- We welcome feedback and comments on the draft
 - For both the protocol and the use-cases
- We will continue working on improving the protocol
- Some parts will be discussed in future versions
 - Congestion control
 - Flow control
- Followed a preliminary version of the TCPLS protocol presented at CoNEXT'21 [1]

Prototype

- We implemented *draft-piroux-tcpls-01* on top of `picotls`, a TLS 1.3 implementation in C
- We modified 50 lines of `picotls` for the required TCPLS interface
- The prototype implements stream multiplexing, failover and multipath
- It consists of 2.5k lines of C
- We will release the prototype under an open-source license

Conclusion

- TCPLS is a secure, user-space, transport protocol bringing
 - Stream multiplexing
 - Connection migration, Failover
 - Multipath
- TCPLS leverages in-kernel high performance TCP implementations
- We implemented a prototype in 2.5k lines of C
 - We will publish the code
- We are interested pursuing this work within the IETF
 - Should we start with a dedicated mailing list ?

Backup – HTTP/2

- HTTP/2+TLS+MPTCP is built on strict layering assumption
- TCPLS offers more control to the application over the TCP connections of the session

