

# Alternative Approach for Mixing Preshared Keys in IKEv2 for Post-quantum Security

`draft-smyslov-ipsecme-ikev2-qr-alt`

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IETF 116

# PPK for IKEv2

Defined in [RFC 8784](#):

Initiator

Responder

**IKE\_SA\_INIT**

HDR, SAI1, KEi, Ni, N(USE\_PPK)

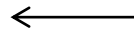
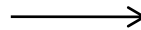


**IKE\_SA\_INIT**

HDR, SAr1, KEr, Nr, N(USE\_PPK)

**IKE\_AUTH**

HDR, SK{IDi, AUTH, SAI2, TSi, TSr,  
N(PPK\_IDENTITY) [, N(NO\_PPK\_AUTH) ] }



**IKE\_AUTH**

HDR, SK{IDr, AUTH, SAr2, TSi, TSr,  
N(PPK\_IDENTITY) }

# The Problem

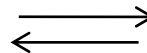
- Initial IKE SA is not protected with PPK (WG decision)
  - it was assumed that no sensitive information was transferred over initial IKE SA, and one could immediately rekey it to get protection
- G-IKEv2 ([draft-ietf-ipsecme-g-ikev2](#)) uses initial IKE SA to immediately transfer session keys from Group Controller/Key Server (GCKS) to Group Member (GM)
  - these keys **are not protected** with PPK

GM

GCKS

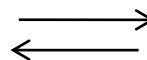
**IKE\_SA\_INIT**

HDR, SAi1, KEi, Ni, N(USE\_PPK)



**GSA\_AUTH**

HDR, SK{ IDi, AUTH, IDg, SAg,  
N(PPK\_IDENTITY) [, N(NO\_PPK\_AUTH) ] }



**IKE\_SA\_INIT**

HDR, SAr1, KEr, Nr, N(USE\_PPK)

**GSA\_AUTH**

HDR, SK{ IDr, AUTH, N(PPK\_IDENTITY) ,  
GSA, **KD** }

# Current Use of PPK with G-IKEv2

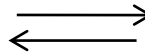
Currently G-IKEv2 draft suggests the following sequence of exchanges to get the protection with PPK:

GM

GCKS

## **IKE\_SA\_INIT**

HDR, SAi1, KEi, Ni, N(USE\_PPK)

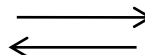


## **IKE\_SA\_INIT**

HDR, SAR1, KEr, Nr, N(USE\_PPK)

## **GSA\_AUTH**

HDR, SK{IDi, AUTH, IDg, SAg,  
N(PPK\_IDENTITY) [, N(NO\_PPK\_AUTH) ] }

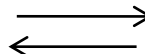


## **GSA\_AUTH**

HDR, SK{IDr, AUTH, N(PPK\_IDENTITY),  
N(REKEY\_IS\_NEEDED) }

## **CREATE\_CHILD\_SA**

HDR, SK{SAi, KEi, Ni}

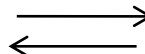


## **CREATE\_CHILD\_SA**

HDR, SK{SAr, KEr, Nr}

## **INFORMATIONAL**

HDR, SK{D}

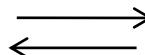


## **INFORMATIONAL**

HDR, SK{ }

## **GSA\_REGISTRATION**

HDR, SK{IDg, SAg}



## **GSA\_REGISTRATION**

HDR, SK{GSA, KD }

# Alternative Approach

Proposed in [draft-smyslov-ipsecme-ikev2-qr-alt](#):

GM

GCKS

## **IKE\_SA\_INIT**

HDR, SAi1, KEi, Ni, N(USE\_PPK),  
N(INTERMEDIATE\_EXCHANGE\_SUPPORTED)



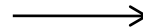
## **IKE\_SA\_INIT**

HDR, SAr1, KEr, Nr, N(USE\_PPK),  
N(INTERMEDIATE\_EXCHANGE\_SUPPORTED)



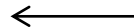
## **IKE\_INTERMEDIATE**

HDR, SK{...N(PPK\_IDENTITY)  
[, N(PPK\_IDENTITY)...]}



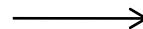
## **IKE\_INTERMEDIATE**

HDR, SK{...N(PPK\_IDENTITY)}



## **GSA\_AUTH**

HDR, SK{IDi, AUTH, IDg, SAg}



## **GSA\_AUTH**

HDR, SK{IDr, AUTH, GSA, KD}



# Fallback to RFC 8784

- If the responder doesn't support this extension, then it doesn't respond with any PPK\_IDENTITY in IKE\_INTERMEDIATE
  - the initiator MAY fallback to RFC 8784 in this case
  - the same situation happens if the responder isn't configured with any of the proposed PPK\_IDs
    - no need to fallback to RFC 8784 in this case, but allowed in the draft for simplicity
- It is possible to modify draft to distinguish between these two cases and disallow fallback if extension is supported, but no PPK found

# Double PPK

- Do we need to support using both RFC 8784 and this draft's approaches for a single SA?
  - Currently is not supported in the draft
  - It seems that this is too complex with no benefits
    - Should be explicitly prohibited in the draft?

# Session Keys Calculation

- RFC 8784:

```
SKEYSEED = prf(Ni | Nr, g^ir)
{SK_d' | SK_ai | SK_ar | SK_ei | SK_er | SK_pi' | SK_pr'} =
    prf+ (SKEYSEED, Ni | Nr | SPIi | SPIr)

SK_d = prf+ (PPK, SK_d')
SK_pi = prf+ (PPK, SK_pi')
SK_pr = prf+ (PPK, SK_pr')
```

- This proposal

```
SKEYSEED' = prf+ (PPK, SK_d)
{SK_d | SK_ai | SK_ar | SK_ei | SK_er | SK_pi | SK_pr} =
    prf+ (SKEYSEED', Ni | Nr | SPIi | SPIr )
```



# Mismatched PPK

- If PPKs with the same PPK\_ID are different, then we run into the problem that the responder cannot decrypt and authenticate IKE\_AUTH messages and will drop them
  - MUST be fixed in the next version of the draft
    - need to have key confirmation payload in the IKE\_INTERMEDIATE exchange, perhaps  $\text{prf}(\text{PPK}, N_i \parallel N_r \parallel \text{SPI}_i \parallel \text{SPI}_r)$
    - Who should send it – initiator or responder? Seems like more appropriate for initiator

# Comparison

- For G-IKEv2:
  - 3 exchanges instead of 5
  - 1 DH shared key computation instead of 2
  - 1 computation of AUTH in case of optional PPK instead of 2
  - initiator can propose several PPK\_IDs
- Can also be used in IKEv2:
  - 3 exchanges instead of 2
    - but PPK\_ID can be piggybacked if IKE\_INTERMEDIATE is also used for other purposes
  - 1 computation of AUTH instead of 2 if PPK is optional
  - initiator can propose several PPK\_IDs

# Coexistence

- The proposed approach is **not intended to replace** the existing one, both can co-exist:
  - for G-IKEv2 the proposed approach can be a primary one (or the only one?)
  - for IKEv2 the proposed approach can be an alternative one (e.g. if IKE identities need to be protected)

# Implementations

- At least 2 implementations of -06 draft exists:
  - ELVIS-PLUS
  - libreswan
- Successfully interoperated during hackathon

# Thanks

- Comments? Questions?
- More details in the draft
- WG adoption?