





P4 Tofino Implementation Experiences with Advanced Stateless Multicast Source Routing

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Motivation

- Concept of Advanced Stateless Multicast Source Routing
 - BitString-Encoded Explicit Trees (BEET)
 - Segment-Encoded Explicit Trees (SEET)
 - Combination thereof
- Conclusion



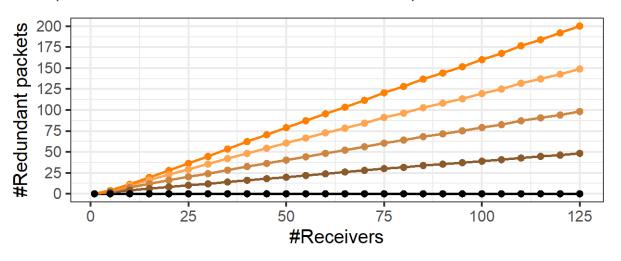


- ► BIER bitstring typically 256 bits
 - Not large enough to cover large networks
- ► Partition BIER domain into sets ≤ 256 nodes
 - Use set identifier (SIs) to indicate bitstrings for that set
 - Helps to scale BIER to large domains

▶ But ...

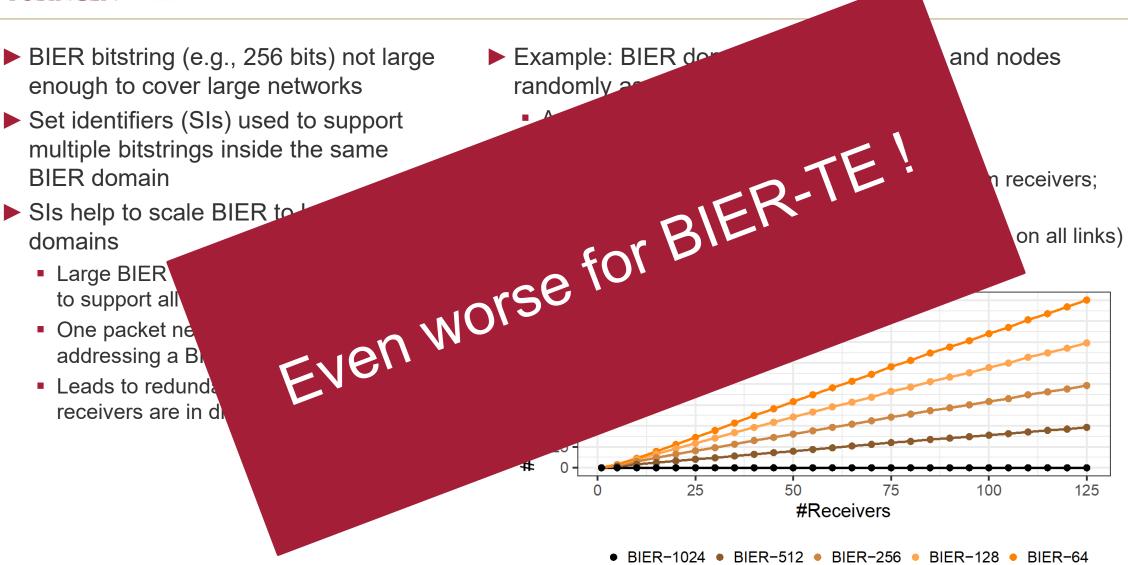
- Large BIER domains require many sets to support all BFERs
- One packet sent per set w/ a receiver
- Leads to redundant packet copies if receivers are in different sets

- Experiment
 - BIER domain with 1024 nodes, average node degree: 4
 - BIER-X: bitstring length of x bit
 - BFERs randomly assigned to SIs
 - Random source sends BIER packet to *n* random receivers; averaged over 50 runs
 - #Redundant packets = (#Packets in BIER domain on all links)
 (#Packets in IPMC domain on all links)



• BIER-1024 • BIER-512 • BIER-256 • BIER-128 • BIER-64



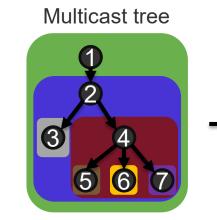


Motivation

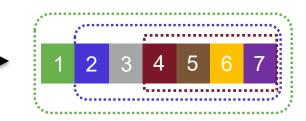


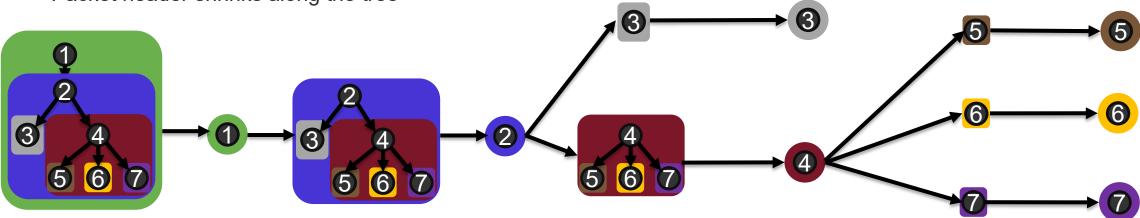
Idea

- Convert distribution tree into a list and encode it in the header
 - Instead of using a flat bitstring
- No need for sets or SIs
- Send multiple packets only if header size does not suffice
- Forwarding principle for replication nodes
 - Partition tree in header into subtrees
 - Send packet copies encapsulated in headers with single subtree to next hops
 - Packet header shrinks along the tree



Sequential order

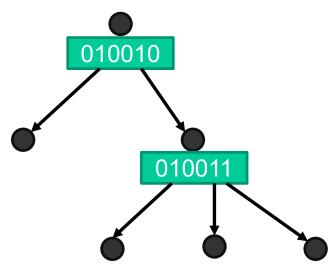






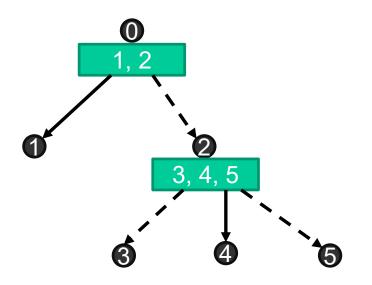
Bitstring-Encoded Explicit Trees (BEET)

- CGM2 / RBS
 - Use bitstring to address next hops
 - Already presented in BIER-WG
 - draft-eckert-bier-cgm2-rbs-01
 - draft-eckert-bier-rbs-00



Segment-Encoded Explicit Trees (SEET)

- Use domain-wide IDs to address next hops
 - New concept, included in
 - draft-eckert-pim-rts-forwarding

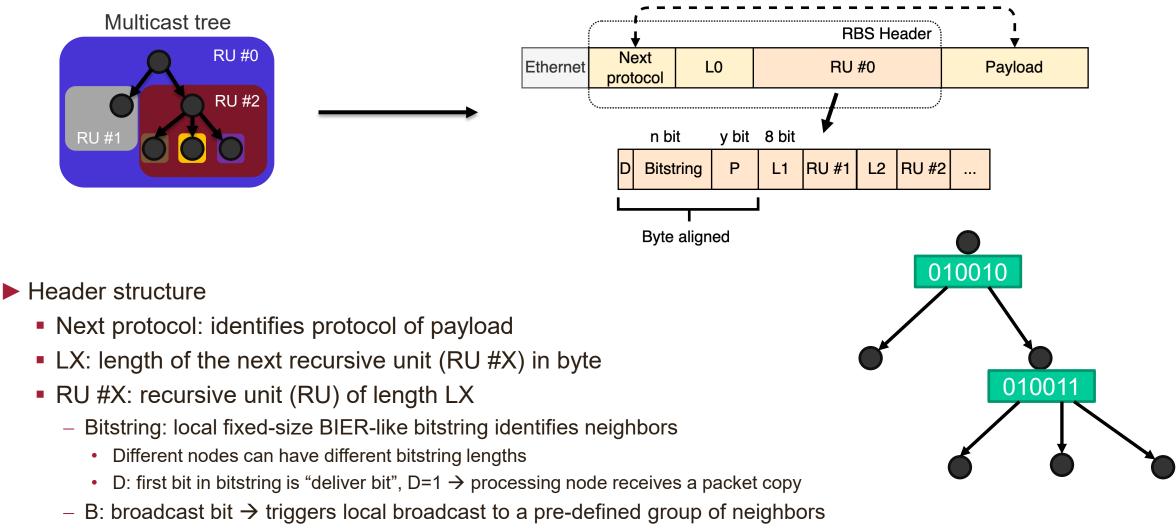


Proof of concept

- Prototype of "RBS Light" and SEET on P4/Intel Tofino[™] (100 Gb/s)
- Whatever is implemented with P4 on Intel Tofino should be simple enough for other platforms



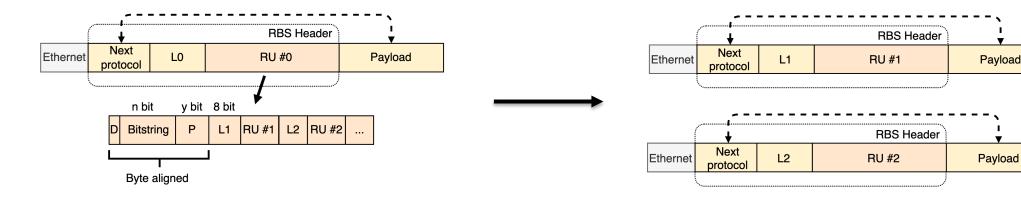
Recursive BitString Structure (RBS) (1)



- P: padding \rightarrow bitstring + B + P need to be byte-aligned







- Forwarding operation
 - For each activated bit in Bitstring
 - Copy LX and the next LX byte (RU #X) into a new packet
 - Forward packet copy to the neighbor that is identified by that bit

- Optimization
 - Explicit leaf nodes
 - Enough to send them a non-encapsulated packet copy

Payload

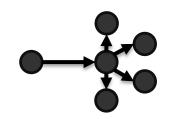
- Do not require an own $RU \Rightarrow$ saves header space
- Not implemented



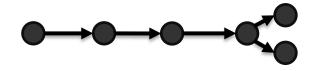
Recursive BitString Structure (RBS) (3)

Advantage

Efficiently encodes packet replication to many neighbors

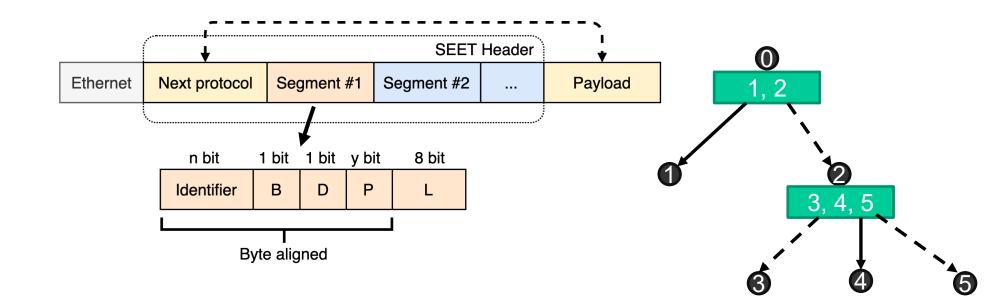


- Disadvantage
 - Less efficient if multicast tree has long "line" paths



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- ► Use segment IDs (SIDs) instead of "local bitstring" to address neighbors
- SIDs can have domain-wide meaning
 - Allows addressing of remote nodes several hops away
- ► Also implemented on P4/Intel TofinoTM



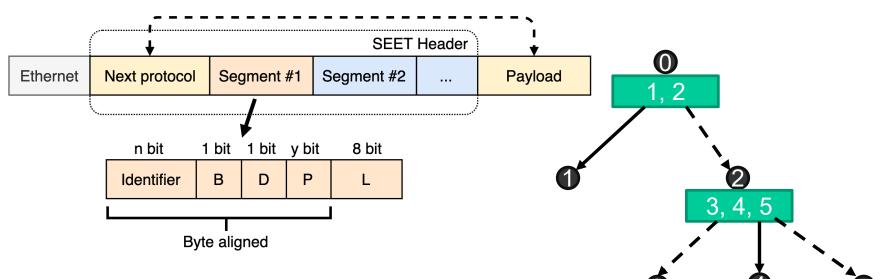


 Use segment IDs (SIDs) instead of "local bitstring" to address neighbors

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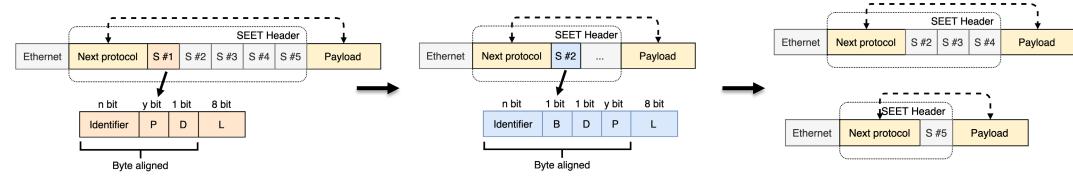
- SIDs can have domain-wide meaning
- Allow addressing of remote nodes several hops away
- Header structure
 - Next protocol: identifies protocol of payload
 - Multiple segments
 - First segment relates to next hop



- Segment structure
 - Identifier: SID
 - Instructions for the denoted node
 - B: broadcast bit → triggers local broadcast to a pre-defined group of neighbors
 - D: deliver bit \rightarrow denoted node receives a packet copy
 - P: padding \rightarrow Identifier + B + D + P needs to be byte-aligned
 - L: number of following header bytes for the denoted node



Segment-Encoded Explicit Trees (SEET) (2)

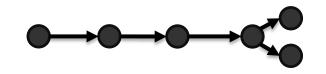


- Forwarding operation
 - Check SID in first segment identifies processing node
 - Check D-bit
 - Yes: deliver a copy to processing node
 - Remove first segment
 - Repeat until original packet empty
 - Copy next segment and the next L byte into a new packet
 - Forward new packet according to SID in first segment
 - Remove the segment and the next L byte from original packet



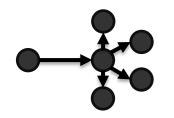
Advantage

• Efficiently encodes multicast trees with long "line" paths



Disadvantage

Less efficient for replication to many neighbors

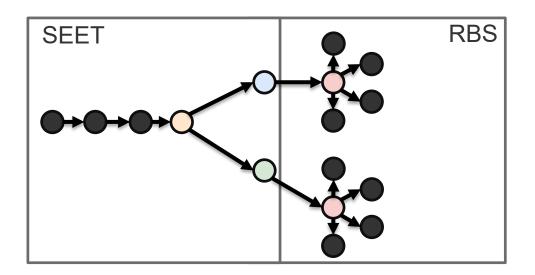


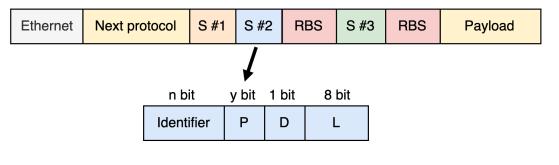


- SEET efficient for sparse multicast trees with long lines
- BEET efficient for multicast trees with many leaves at some penultimate node

Idea

- Use SEET to reach penultimate nodes
- Use encapsulated bitstring at penultimate hops to efficiently replicate to many leaf nodes



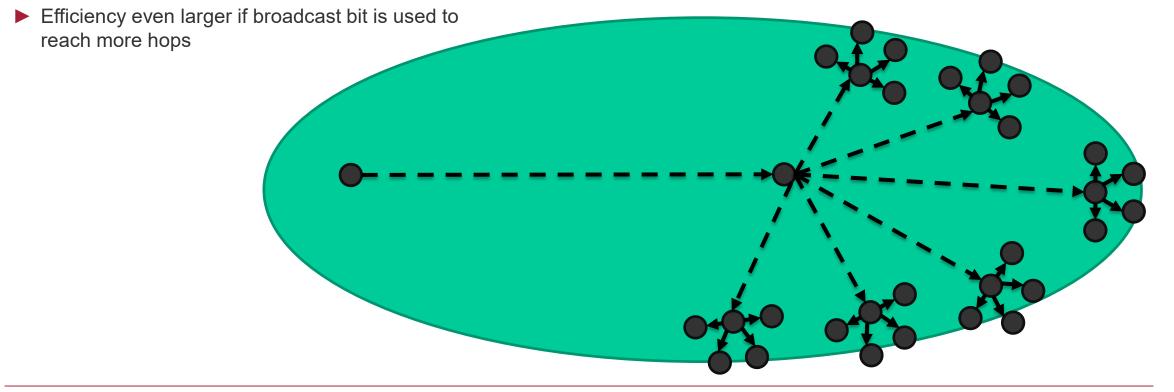


Encode in S #2 that RBS header follows



- "Tunnel 8-bit bitstrings" to penultimate hops over two hops
 - SID (2 bytes)
 - Length field (1 byte)
 - Bitstring (1 byte)

- Header size 256 bits (32 bytes)
 - \Rightarrow 256/8/4-1=7 penultimate hops can be addressed
 - \Rightarrow 7*8=56 receivers can be addressed
- Header size 1024 bits (128 bytes)
 - \Rightarrow 1024/8/4-1=31 penultimate hops can be addressed
 - \Rightarrow 31*8=243 receivers can be addressed





- ► BEET / SEET allow encoding of generic multicast trees
 - Scales better than BIER(-TE) in large domains with sparse multicast trees
- ▶ Implemented on P4/Intel Tofino[™] with 100 Gb/s
- Advantages of SEET and BEET can be combined
 - Tunneling of bitstrings over SEET to address leaf nodes of the tree
 - Not yet implemented
- Technical report with documentation to come