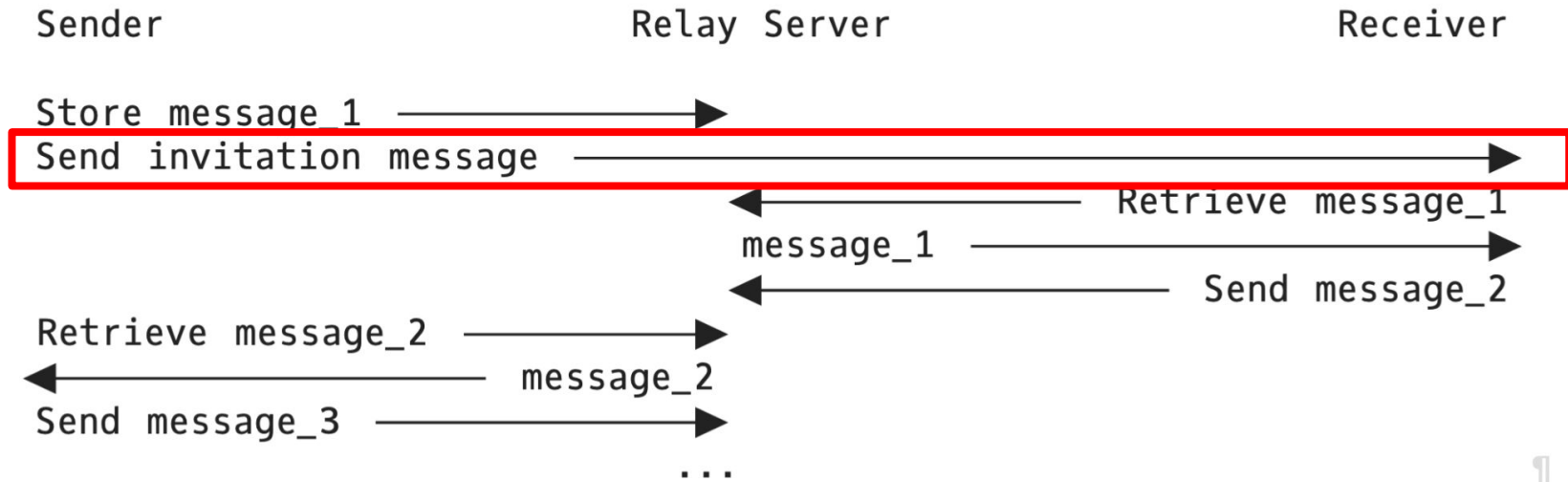


# TIGRESS Introduction Channel Security

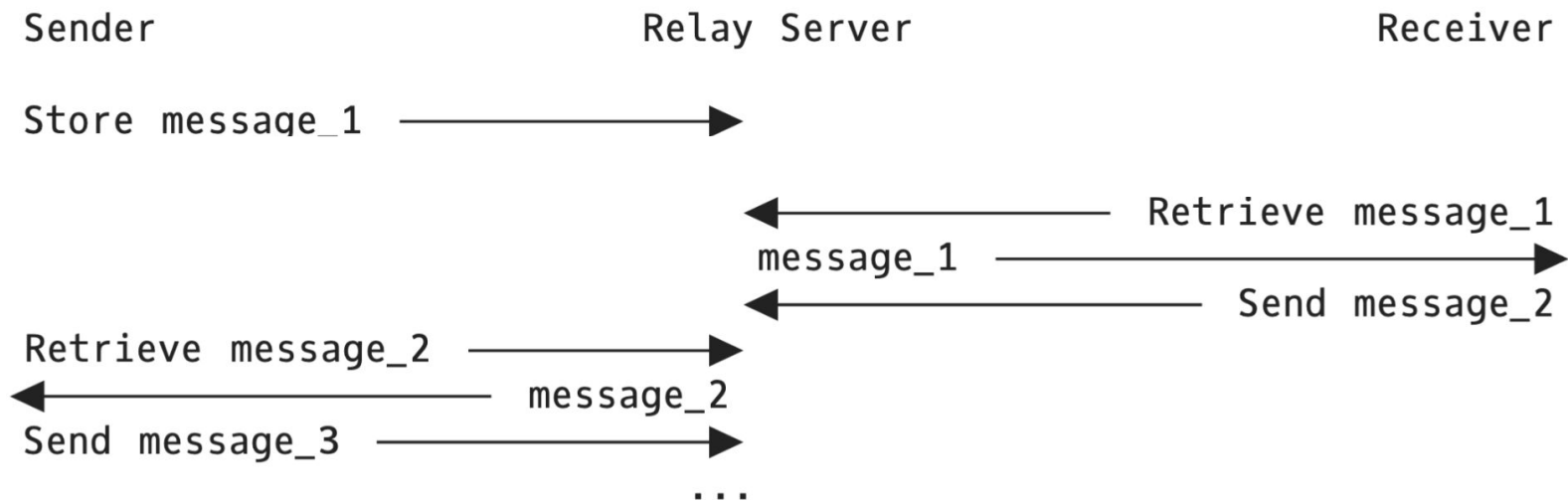
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IETF 118

# Introduction Channel



# Relay Channel



# Secret Invitations

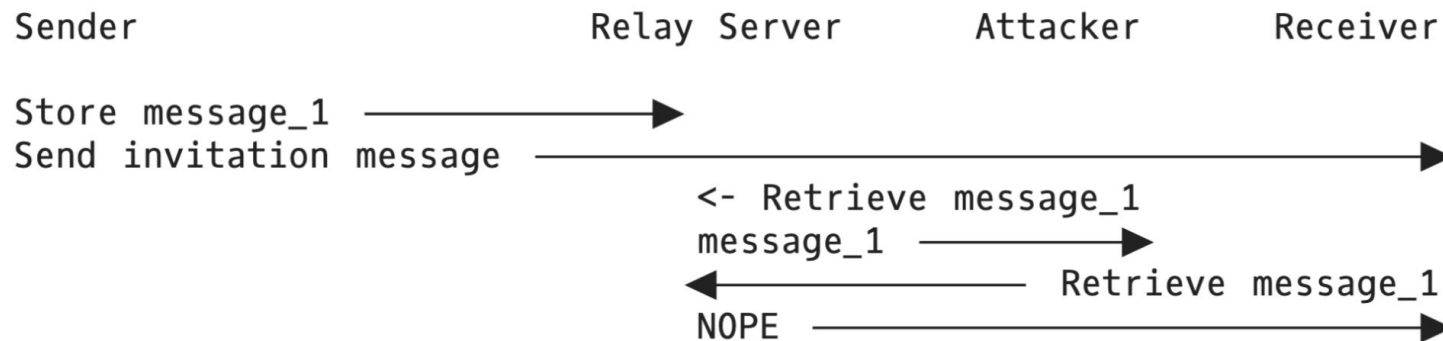
- Invitations identify Relay Channel
  - And are high entropy
- Correct invitation is needed to access the appropriate Relay Channel
- As long as Invitation remains secret, only the Receiver can access the Relay Channel

# At-most-once semantics

- Both protocols are designed to provide at-most-once semantics
- The first agents to initiate a connection to the Relay Channel can access it
- Subsequent agents get denied access to the channel

# Public Invitations

- If the attacker learns the invitation, then there is a race with the receiver
- The winner of the race gets the credential
- The attacker is often going to win this race

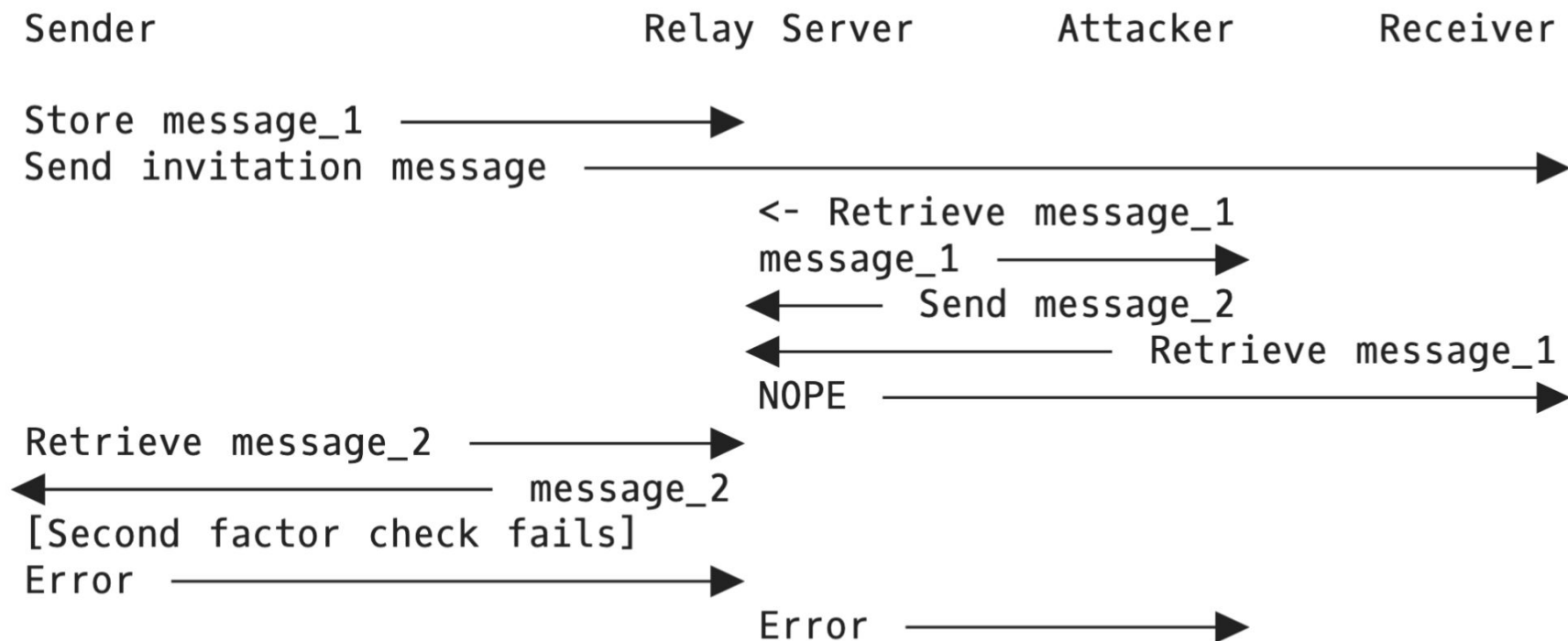


# Second Factor

- What if you have a second factor?
  - Delivered over a secure channel
- Examples
  - A cryptographic key
  - A PIN
  - An identity check\*
- This controls credential issuance **but not the channel**
- Result: the attacker can access the channel but not get the credential

\* Strictly speaking that's not delivered

# Denial of Service





# Properties of Introduction Channel

- Lots of different channels
  - e-mail, SMS, iMessage, WhatsApp, NFC
- Widely varying security properties
- In practice we tend to treat these all as “secure” for other applications...
  - Password reset
  - 2FA
  - WebPKI certificate issuance

# Options

1. Assume that the Introduction Channel is secure and move forward to protocol selection.
2. Assume that the Introduction Channel is insecure and ask what the properties of the “second factor” are and what needs to be done to bind it to the Relay Channel. This should happen prior to protocol selection.

Questions?