

# Verifying Security Protocols End-to-End with Owl

**Joshua Gancher**  
**Carnegie Mellon University**



# Owl

new language for building security protocols

- TLS, WireGuard, Signal, ...
- Supports **verified implementations**

# Owl

new language for building security protocols

- TLS, WireGuard, Signal, ...
- Supports **verified implementations**

Our vision:

**formally verified**, drop-in  
replacements of protocol implementations

# Owl

new language for building security protocols

- TLS, WireGuard, Signal, ...
- Supports **verified implementations**

Our vision:

**formally verified**, drop-in  
replacements of protocol implementations



# Owl

new language for building security protocols

- TLS, WireGuard, Signal, ...
- Supports **verified implementations**

Our vision:

**formally verified**, drop-in  
replacements of protocol implementations

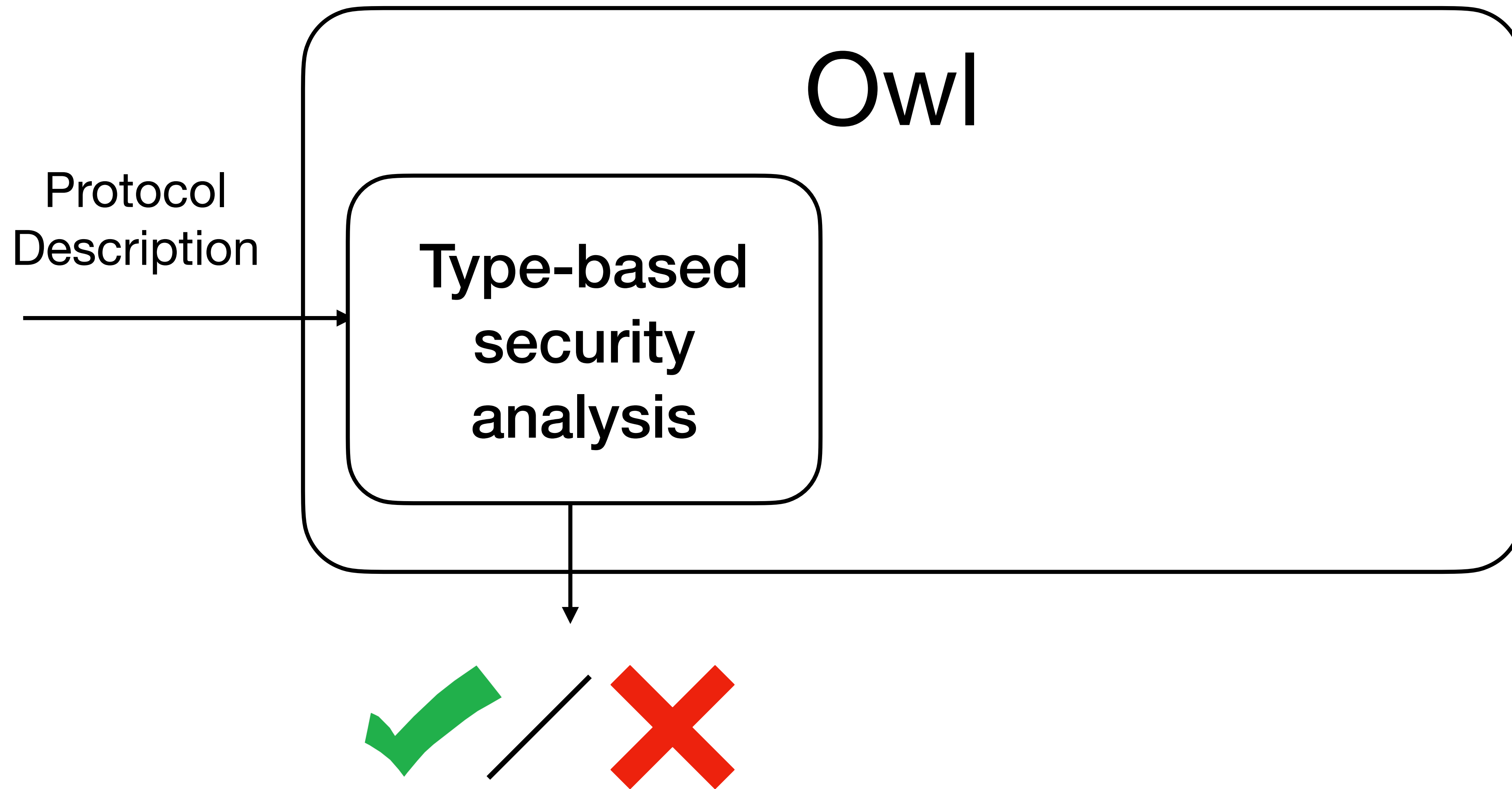


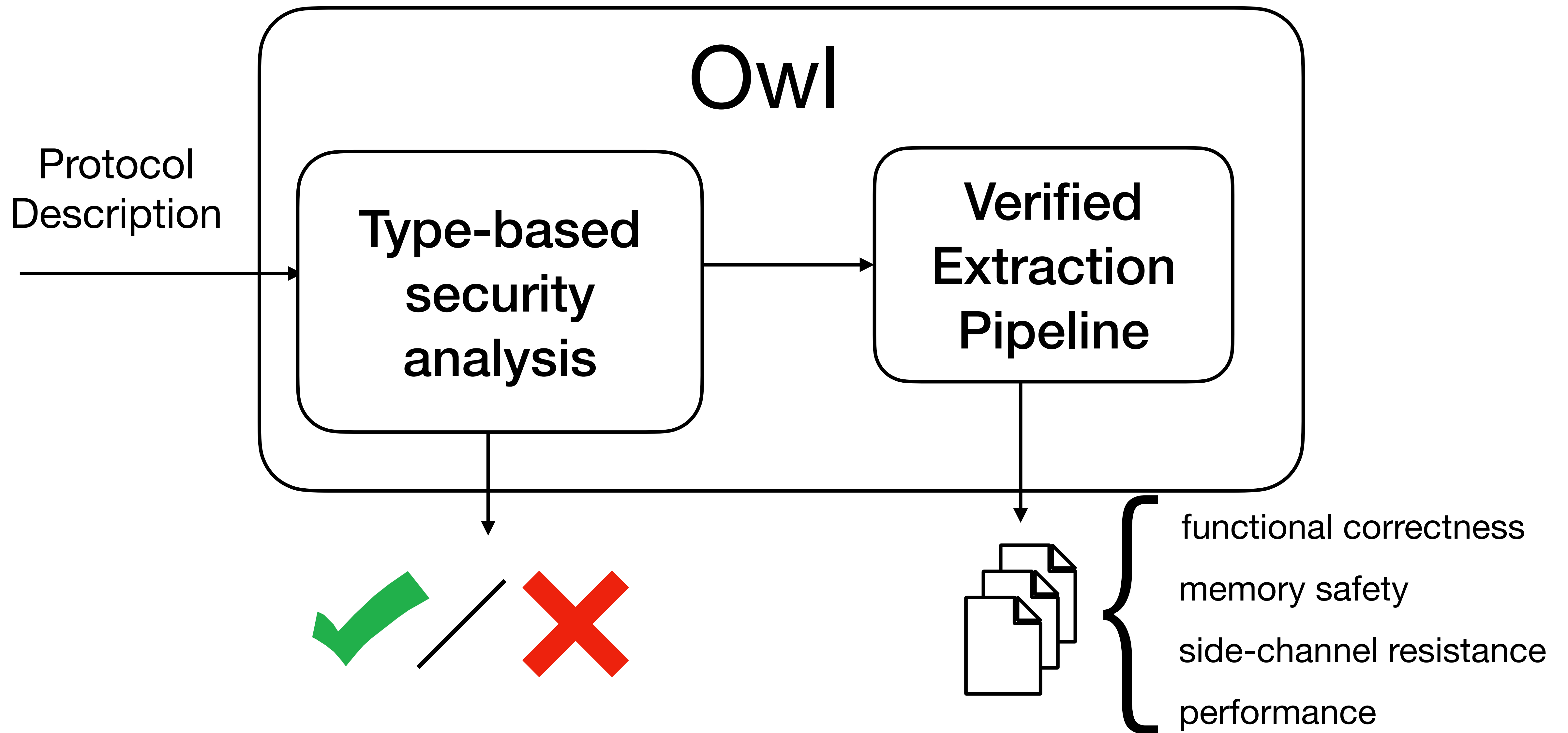
{ functional correctness  
memory safety  
secure designs

Protocol  
Description

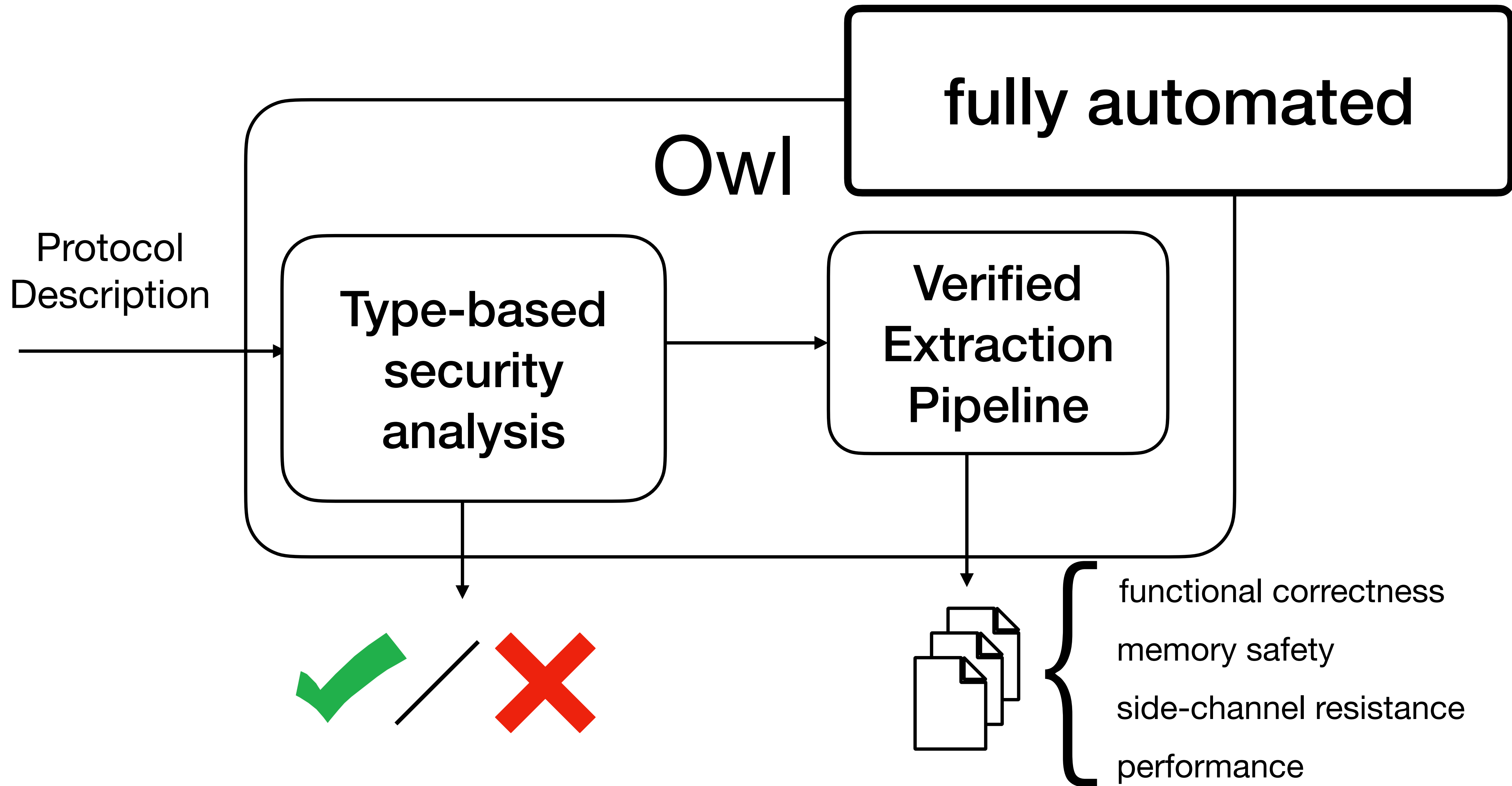


Owl









# OWL: Compositional Verification of Security Protocols via an Information-Flow Type System

Joshua Gancher

Sydney Gibson

Pratap Singh

Samvid Dharanikota

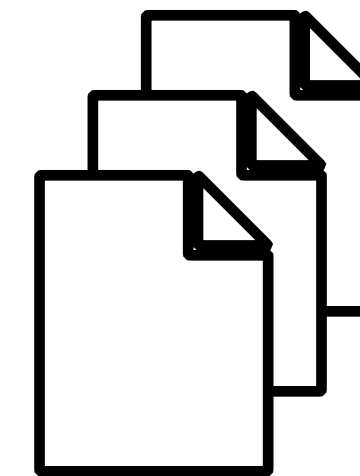
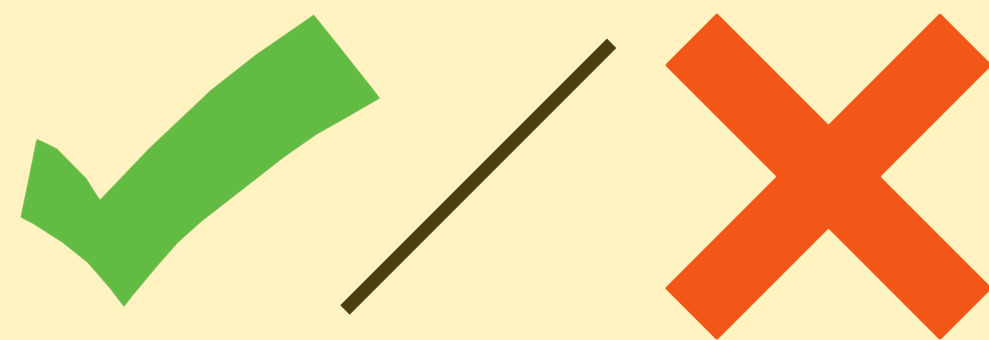
Bryan Parno

Carnegie Mellon University

Protocol  
Description

**Type-based  
security  
analysis**

**Verified  
Extraction  
Pipeline**



functional correctness  
memory safety  
side-channel resistance  
performance

# **Symbolic Security**

# **Computational Security**

# Symbolic Security

Crypto modeled by  
**abstract terms**,  
equations on functions

# Computational Security

# Symbolic Security

Crypto modeled by  
**abstract terms**,  
equations on functions

<b>functions:</b> $\text{enc}/2, \text{dec}/2$ <b>equations:</b> $\text{dec}(\text{enc}(m, k), k) = m$
---

# Computational Security

# Symbolic Security

---

Crypto modeled by  
**abstract terms**,  
equations on functions

<b>functions:</b> $\text{enc}/2, \text{dec}/2$ <b>equations:</b> $\text{dec}(\text{enc}(m, k), k) = m$
---

attacker can only use  
specified functions,  
equations

# Computational Security

---

# Symbolic Security

---

Crypto modeled by  
**abstract terms**,  
equations on functions

<b>functions:</b> $\text{enc}/2, \text{dec}/2$ <b>equations:</b> $\text{dec}(\text{enc}(m, k), k) = m$
---

attacker can only use  
specified functions,  
equations

# Computational Security

---

Crypto is given by algorithms on  
bytes

Cryptography specified by  
**security properties:**  
secrecy of messages,  
unforgeability of ciphertexts,  
...

# Symbolic Security

---

Crypto modeled by  
**abstract terms**,  
equations on functions

<b>functions:</b> $\text{enc}/2, \text{dec}/2$ <b>equations:</b> $\text{dec}(\text{enc}(m, k), k) = m$
---

attacker can only use  
specified functions,  
equations

# Computational Security

---

Crypto is given by algorithms on  
bytes

Cryptography specified by  
**security properties:**  
secrecy of messages,  
unforgeability of ciphertexts,  
...

Strong attacker model;  
closer to implementations



# Types for Security Protocols

# Types for Security Protocols

$K : \text{enckey for } T$   
 $m : T$   $\Longrightarrow$   $\text{enc}(K, m) : \text{public}$

# Types for Security Protocols

$K : \text{enckey for } T$   
 $m : T \quad \Longrightarrow \quad \text{enc}(K, m) : \text{public}$

security via type checking:  $\vdash P \Longrightarrow P \text{ secure}$

# Types for Security Protocols

$K : \text{enckey for } T$   
 $m : T$   $\Longrightarrow$   $\text{enc}(K, m) : \text{public}$

security via type checking:  $\vdash P \Longrightarrow P \text{ secure}$   
 $\downarrow$   
**one time proof effort**

# Types for Security Protocols

$$\begin{array}{l} K : \text{enckey for } T \\ m : T \end{array} \Longrightarrow \text{enc}(K, m) : \text{public}$$

security via type checking:  $\vdash P \Longrightarrow P \text{ secure}$   
 $\downarrow$   
**one time proof effort**

developer only  
**uses type system!**

# Types for Security Protocols

$K$  : enckey for  $T$   
 $m$  :  $T$   $\Longrightarrow$   $\text{enc}(K, m)$  : public

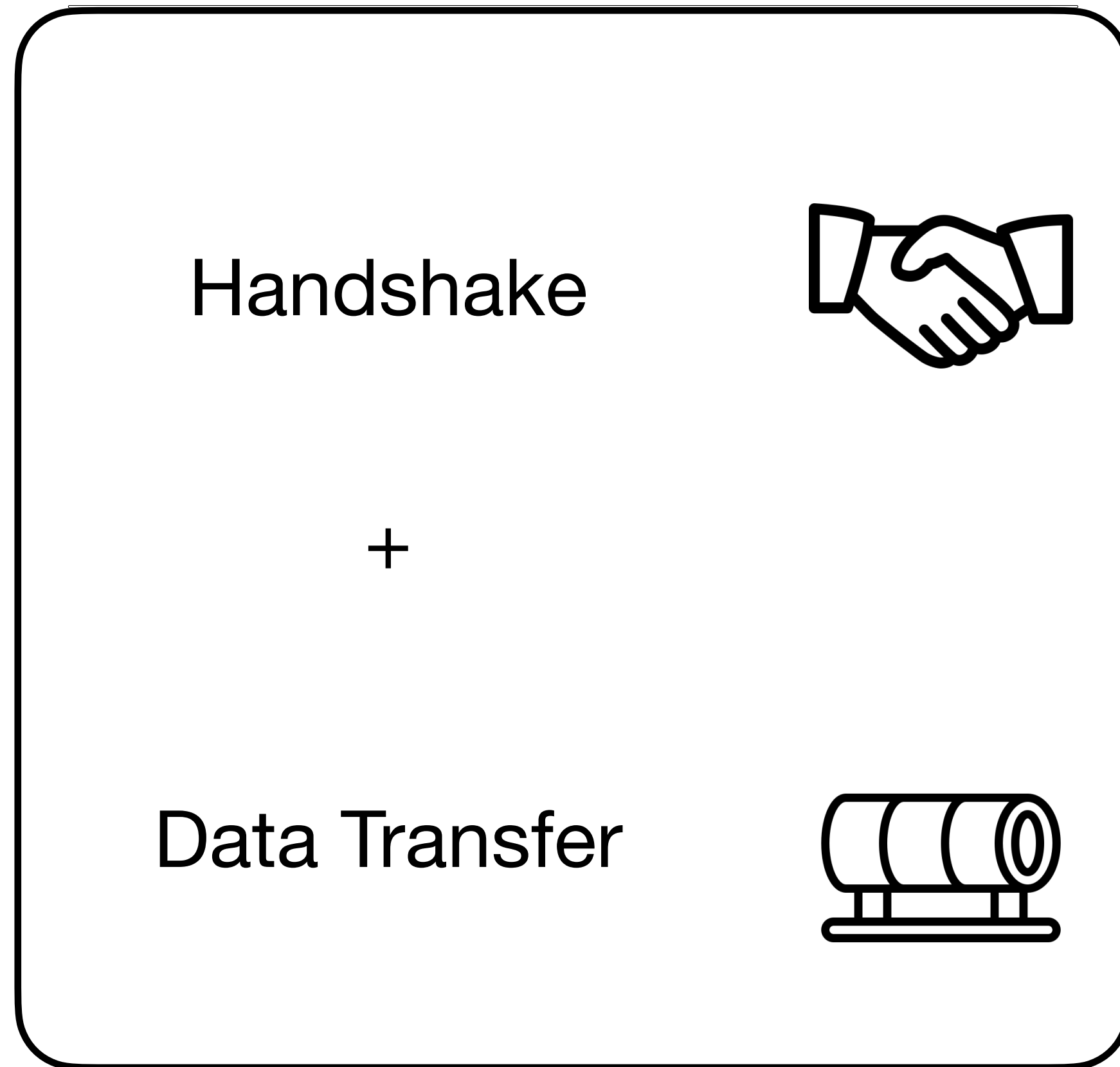
security via type checking:  $\vdash P \Longrightarrow P$  secure  
 $\downarrow$   
one time proof effort

developer only  
**uses type system!**

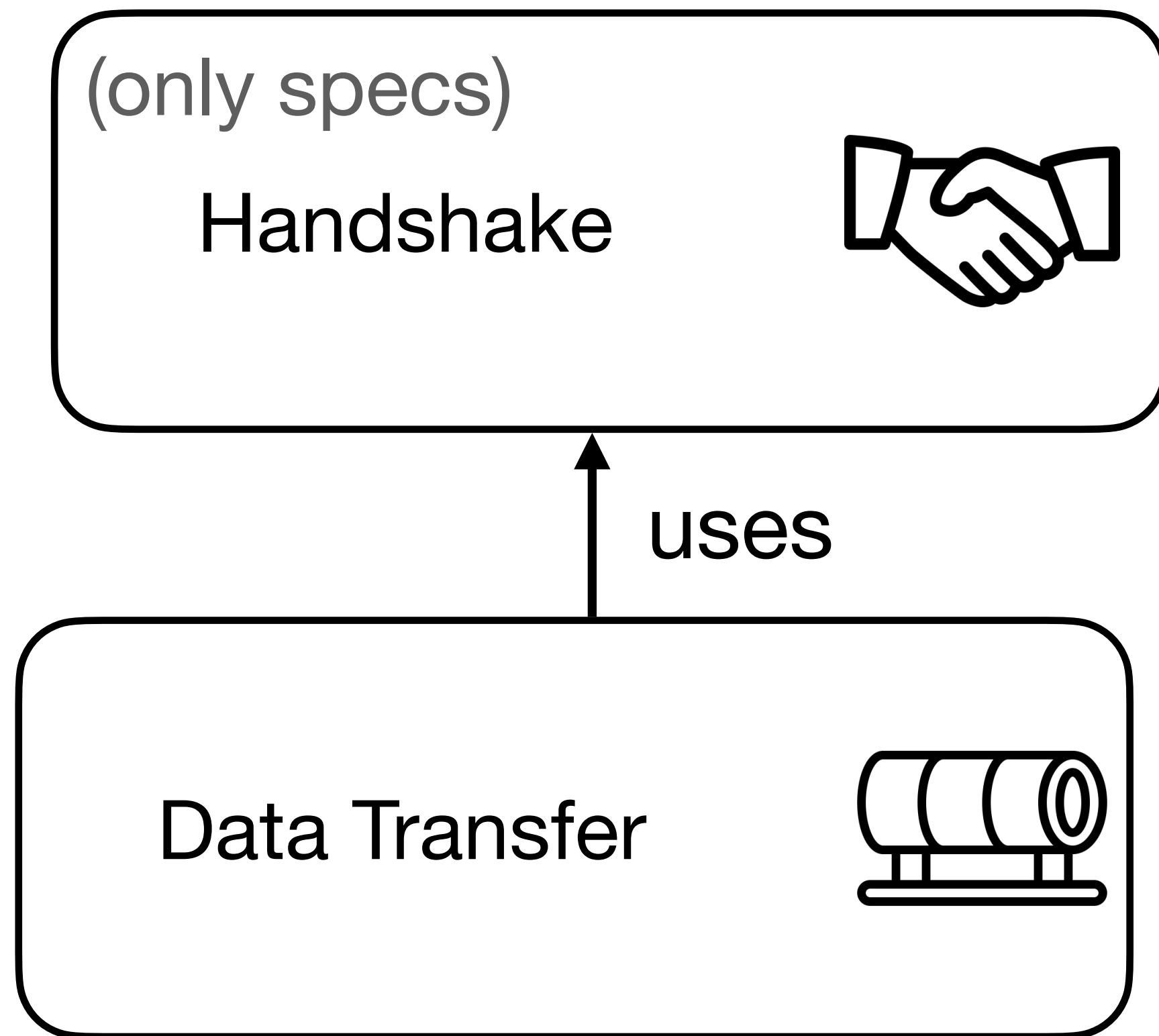


automatic,  
modular proof effort

# Protocol-Level Modularity

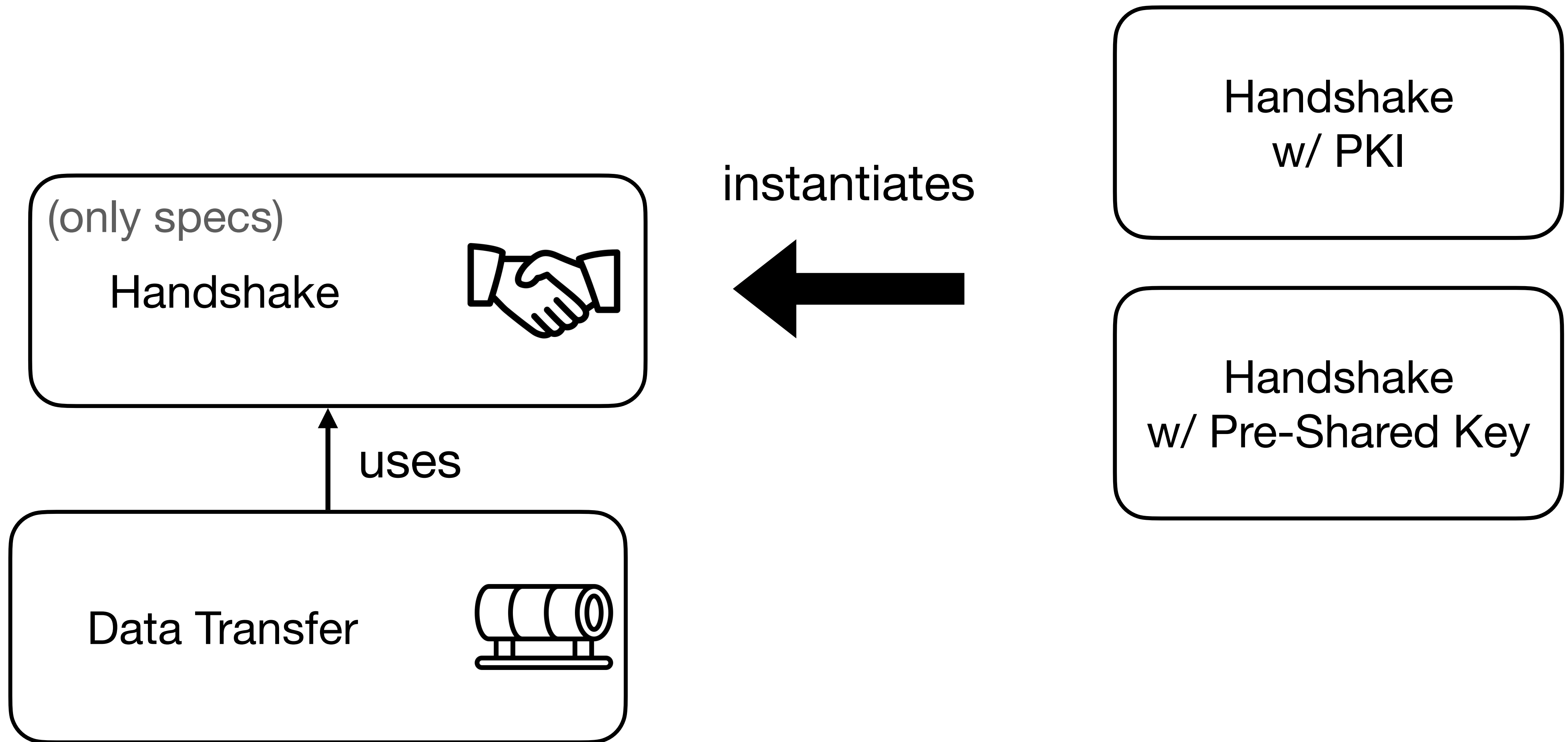


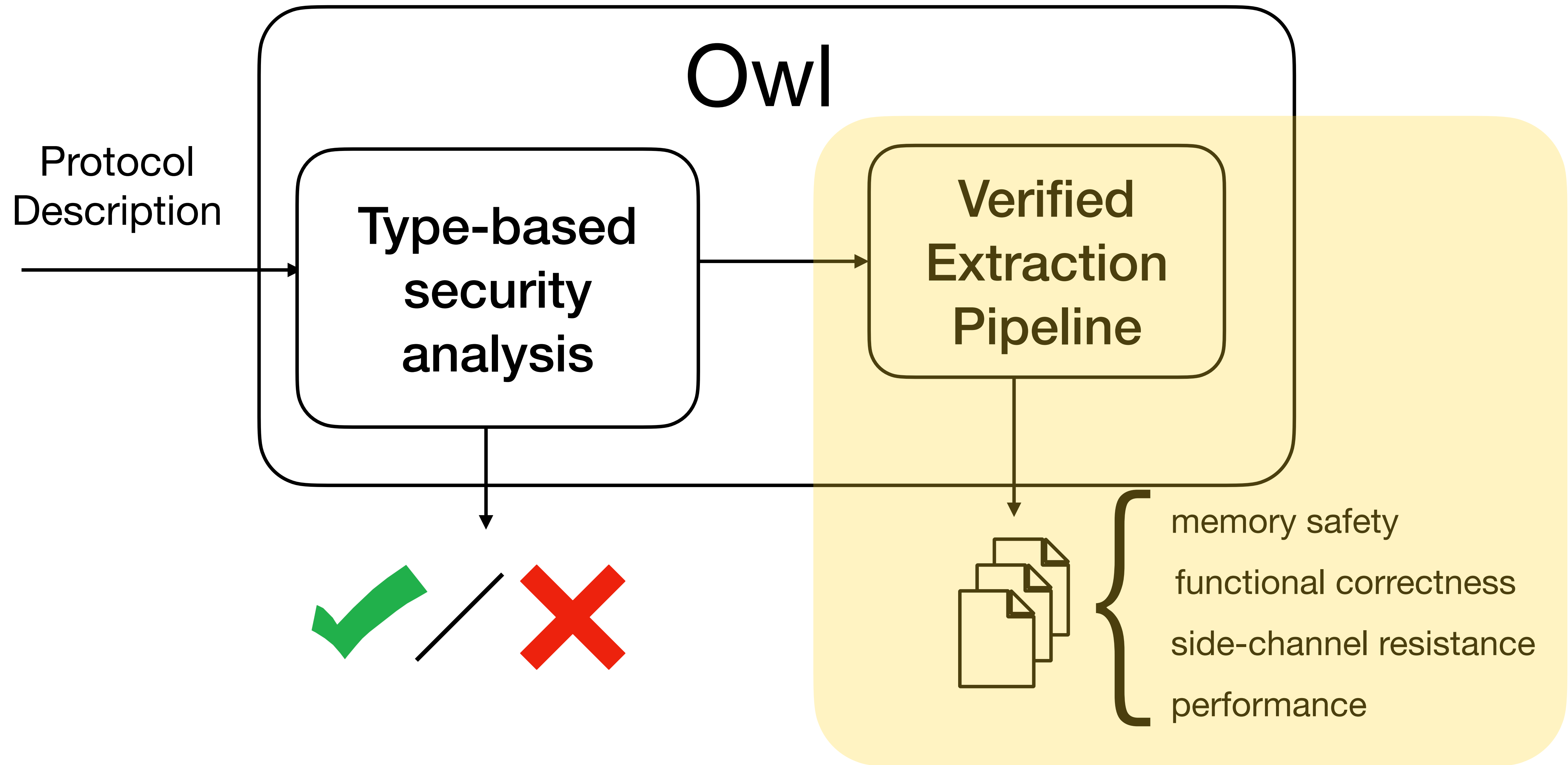
# Protocol-Level Modularity

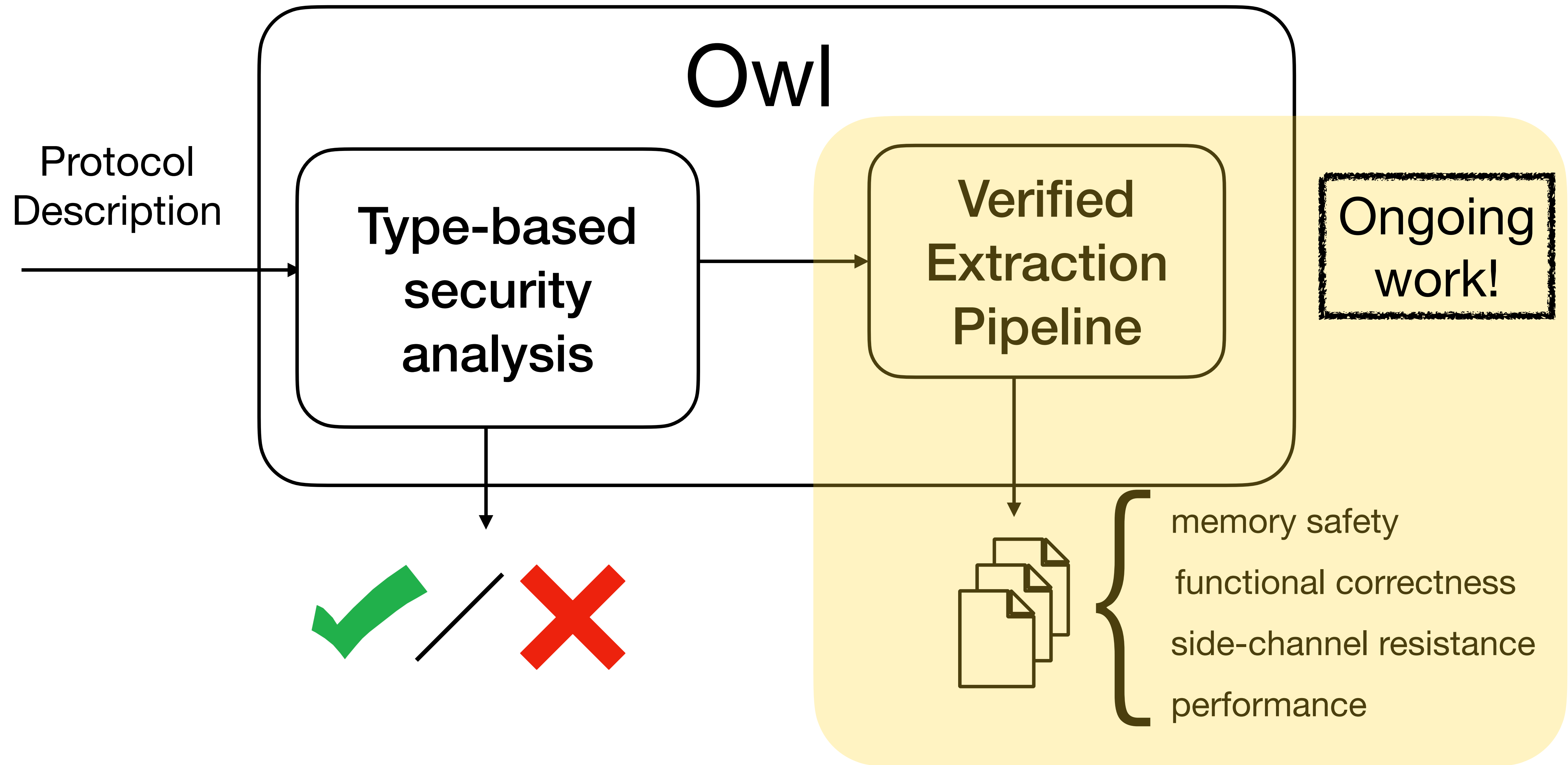




# Protocol-Level Modularity

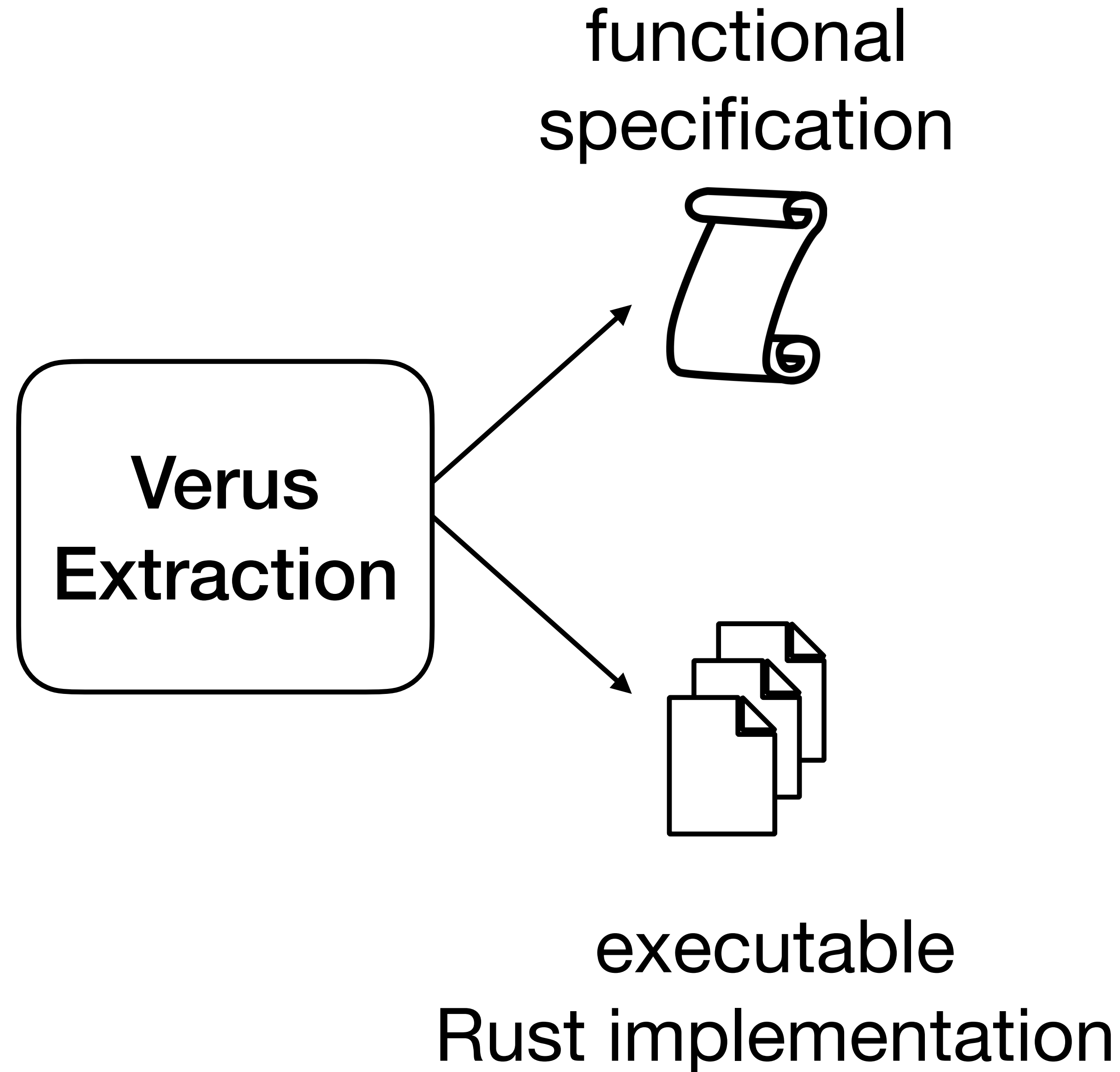




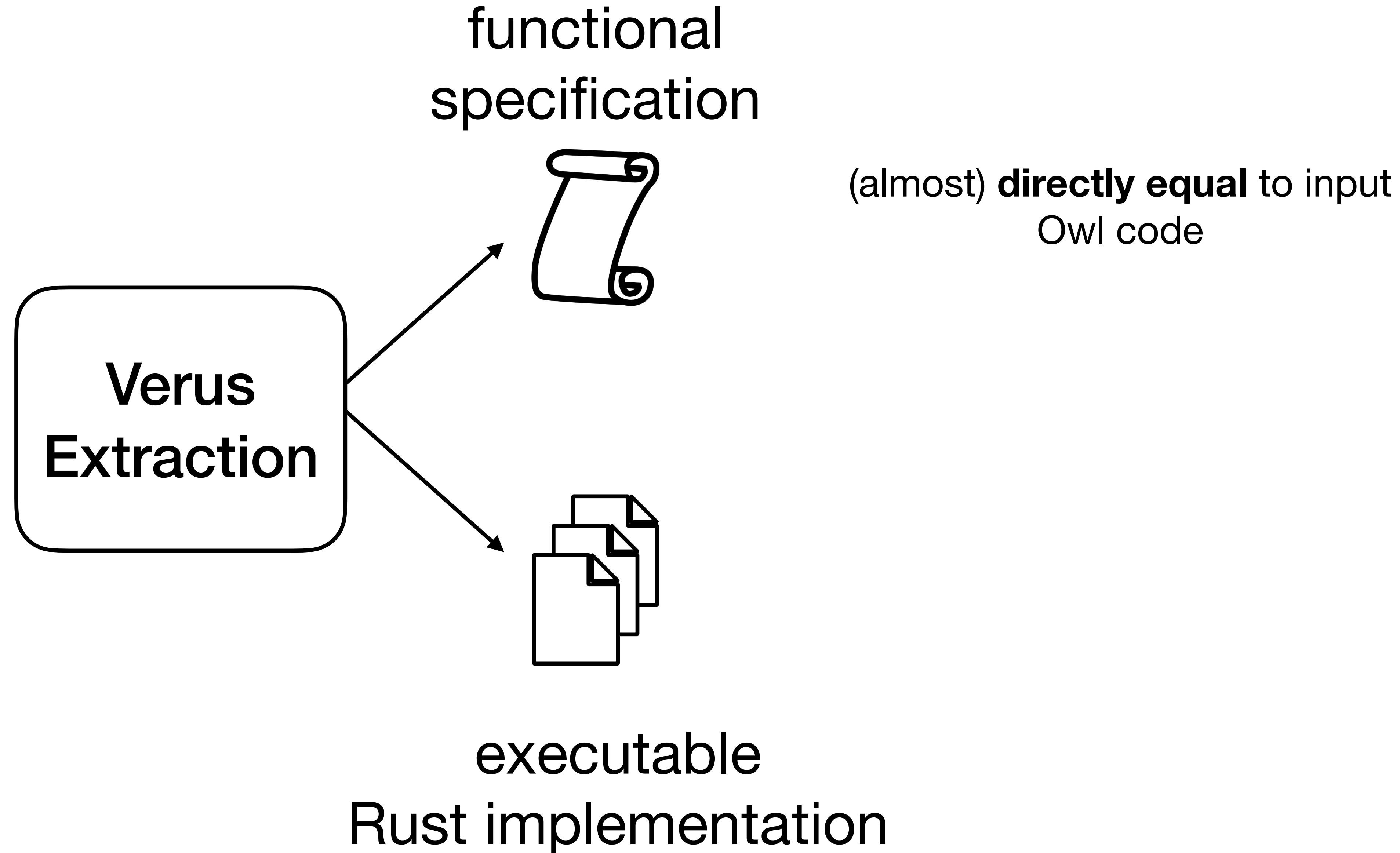


# Verus: Verifying Rust Programs using Linear Ghost Types

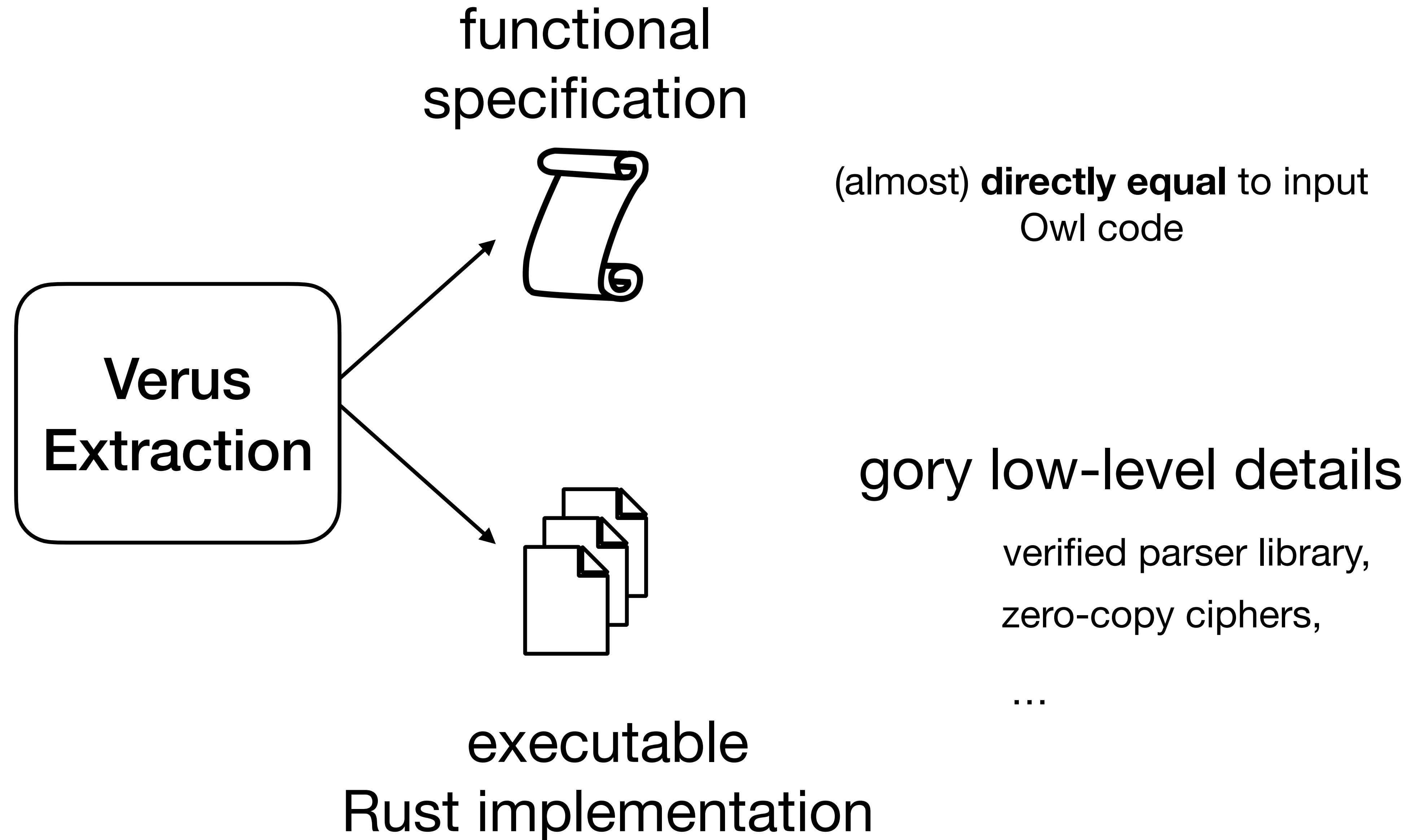
# Verus: Verifying Rust Programs using Linear Ghost Types



# Verus: Verifying Rust Programs using Linear Ghost Types

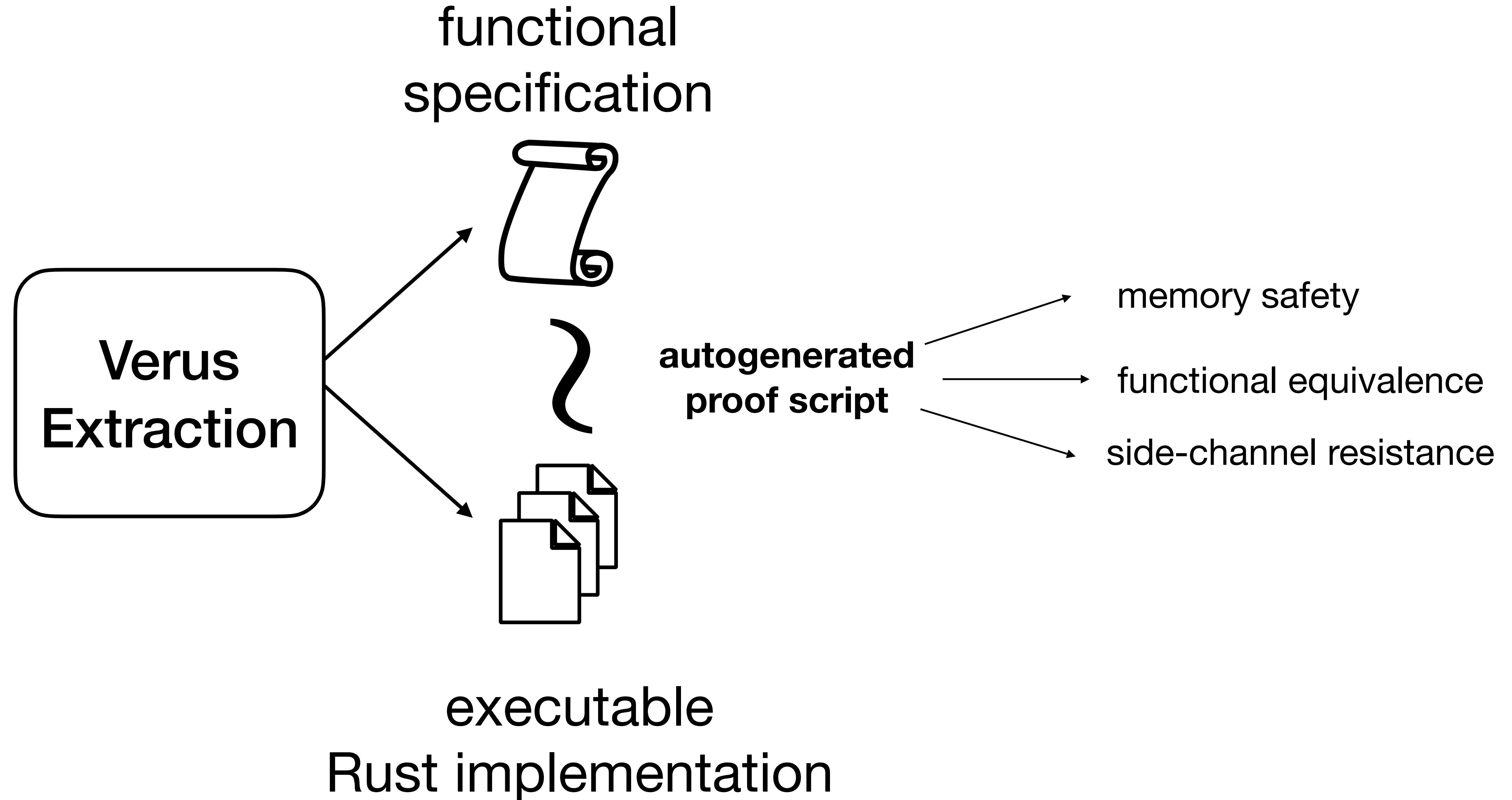


# Verus: Verifying Rust Programs using Linear Ghost Types





# Verus: Verifying Rust Programs using Linear Ghost Types





# Ongoing Work: a verified VPN



# Ongoing Work: a verified VPN



widely used:

inside Linux kernel

# Ongoing Work: a verified VPN



widely used:

inside Linux kernel

very lean:

implementable in 4K LoC

# Ongoing Work: a verified VPN



widely used:

inside Linux kernel

very lean:

implementable in 4K LoC

**Goal: verified, drop-in replacement**



# Owl: End-to-End Verification of Security Protocols via a Secure Type System

New tool for **modular, automated** proofs of security protocols

- Novel use of **type systems** for constructing secure cryptographic protocols
  - Security is proved once-and-for-all;
  - Protocols checked **via type checker**
- Ongoing work: **verified extraction** and drop-in implementation of WireGuard

[owl-lang.org](https://owl-lang.org)

[jgancher@andrew.cmu.edu](mailto:jgancher@andrew.cmu.edu)