

RFC 4895bis: SCTP Authentication

draft-ietf-tsvwg-rfc4895-bis-02

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Scope

- Incorporate relevant changes from draft-nagesh-sctp-auth-4895bis-00
- Address two security issues reported by Ericsson:
 - Use direction dependent keys to mitigate reflection attacks.
 - Don't use different HMAC algorithms with the same keys.
- Generalize HMAC to MAC.
- Add more algorithms, potentially retire HMAC-SHA-1.
- Add socket API considerations for improved control of the SCTP AUTH usage.

Status

- draft-tuexen-tsvwg-rfc4895-bis-00
Submit RFC 4895 as an ID.
- draft-tuexen-tsvwg-rfc4895-bis-01
Update to xmlv3.
- draft-tuexen-tsvwg-rfc4895-bis-02
Wordsmithing and updating references.
- draft-tuexen-tsvwg-rfc4895-bis-03
Minor editorial change.
- draft-tuexen-tsvwg-rfc4895-bis-04
Add socket API related updates required for DTLS/SCTP.
- draft-tuexen-tsvwg-rfc4895-bis-05
Remove ekr from list of authors, improve socket API.
- draft-tuexen-tsvwg-rfc4895-bis-06
Update Acknowledgements.
- draft-tuexen-tsvwg-rfc4895-bis-00
Same as above.
- draft-tuexen-tsvwg-rfc4895-bis-01
Incorporate draft-nagesh-sctp-auth-4895bis-00, editorial changes, update IANA section.
- draft-tuexen-tsvwg-rfc4895-bis-02
Introduce directional keys.

SCTP AUTH Handshake

```
----- INIT[RANDOM; CHUNKS; HMAC-ALGO] ----->  
<----- INIT-ACK[RANDOM; CHUNKS; HMAC-ALGO] -----  
----- COOKIE-ECHO ----->  
<----- COOKIE-ACK -----
```

Key Management (1)

- Ensure that the RANDOM parameter values used by both endpoints are different.
 - In client/server scenarios this is easy to accomplish.
 - In case of INIT collision the sending of ABORTs are required, if both sides chose the same 32 byte random number. However, this is very rare.
- Local and remote key vector
 - Local key vector based on parameters sent.
 - Remote key vector based on parameters received.
 - key vector = RANDOM || CHUNKS || HMAC ALGO
 - Local key vector and remote key vector are different.

Key Management (2)

- send context = local key vector || remote key vector
- receive context = remote key vector || local key vector
- send key = KDF(shared_key, send_context, key len)
- receive key = KDF(shared_key, receive_context, key len)
- Using KDF as described in RFC 5926, Section 3.1.
- shared_key is controlled by the API and might be empty!
- Provided in the document based on HMAC-SHA-256.

Backwards Compatibility

- Must be enabled by the upper layer.
- Non directional keys are only used if no directional keys are supported by the peer.
- send key and receive key are the concatenation of
 - $\min(\text{local key vector}, \text{remote key vector})$
 - shared key
 - $\max(\text{local key vector}, \text{remote key vector})$

Computing the AUTH chunk

- When sending, compute $\text{MAC}(\text{send_key}, \text{chunks})$ and put it into the value of the AUTH chunk.
- When receiving, compute $\text{MAC}(\text{receive_key}, \text{chunks})$ and verify that this is the chunk value of the AUTH chunk.

Next Steps

- Double checking: out of scope is improving the replay protection for chunks with sequence numbers:
 - DATA and SACK/NR-SACK
 - I-DATA and SACK/NR-SACK
 - ASCONF and ASCONF ACK
 - RE CONFIG
- Use a formula-based description instead of a text based one.
- Generalize HMAC to MAC.
- Add more algorithms. Which ones?
- Address comments already received and all upcoming comments.