

FrodoKEM

A simple and conservative KEM from generic lattices

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FrodoKEM: Introduction

- ❑ FrodoKEM is a quantum-safe IND-CCA2 secure Key Encapsulation Mechanism (KEM) based on the hardness of the plain Learning With Errors (LWE) problem
 - It uses generic, algebraically unstructured lattices
- ❑ It was a **Round 3 alternate** in the NIST PQC standardization process
 - Dropping FrodoKEM was primarily motivated by performance: “In terms of security, Frodo’s conservative design choices are laudable.” (NIST Round 3 Status Report)
- ❑ FrodoKEM has been recommended for use by several European countries (Germany [BSI24], France, Sweden)
- ❑ Ongoing standardization by ISO (started on April’23)
 - To be included as Amendment 2 of ISO/IEC 18033-2 (together with ML-KEM and Classic McEliece)
 - Currently approved for Draft Amendment (DAM), Sept/Oct’24

Design principles

A simple, conservative yet practical design

- ❑ Plain LWE: generic, algebraically unstructured lattices
 - Minimizes potential attack surface: no algebraic ring structure
- ❑ Cautious parameterization: ‘medium-sized’ errors conforming to a worst-case/average-case reduction
 - (Non-tight) reduction from worst-case bounded distance decoding with discrete Gaussian samples (BDDwDGS) problem to decision LWE
 - Narrower errors \Rightarrow smaller parameters, better efficiency
- ❑ Concrete parameters chosen according to ‘core-SVP’ methodology
 - Lower-bound the first-order exponential time and space of SVP

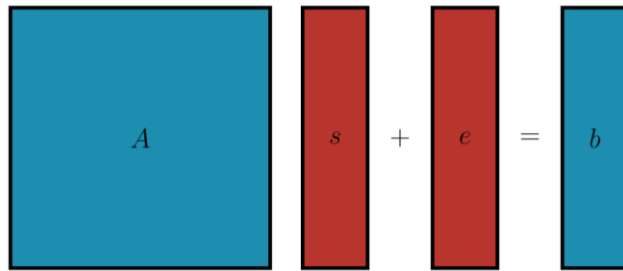
IND-CCA KEM transformation

- ❑ Begin with an **IND-CPA PKE** scheme called FrodoPKE
- ❑ Transform it to IND-CCA KEM using a variant of Fujisaki-Okamoto (FO) transform
 - Hofheinz–Hovelmanns–Kiltz (HHK) give explicit variant **IND-CPA PKE** → **IND-CCA KEM**
 - HHK transform is secure in **both the classical and quantum ROM models**
 - Jiang et al. [JZC+18] eliminates need of some extra hashing
- ❑ Similar to ML-KEM, FrodoKEM uses a slight variant based on Jiang et al. (some additional values in hash computations to reduce risk of multi-target attacks)

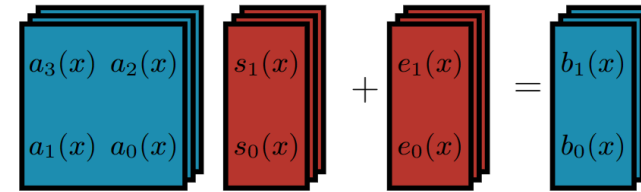
Implementation aspects (part I)

FrodoKEM has a simple design and implementation

- ❑ Matrix-vector products over \mathbb{Z}_q^n with a power-of-2 modulus q
 - Not use of NTT, no prime modulus



Matrix multiplication based (FrodoKEM)



Polynomial multiplication based (ML-KEM)

Implementation aspects (part I)

FrodoKEM has a simple design and implementation

- ❑ Matrix-vector products over \mathbb{Z}_q^n with a power-of-2 modulus q
 - Not use of NTT
- ❑ Straightforward error sampling: approximation to rounded Gaussian
 - E.g., using inversion sampling:
 - Table T_χ stores $(s + 1)$ integers related to discrete cumulative distribution function
 - Given a random value r , determine smallest index i such that $r \leq T_\chi[i]$
 - Output $(-1)^b i$ for a random bit b

Implementation aspects (part II)

FrodoKEM has a simple design and implementation

- ❑ x64 implementation consists of ~350 lines of C code (+ existing symmetric primitives)
- ❑ Facilitates compiler's work to produce optimized code
 - E.g., x64 implementation does not require hand-assembly code
- ❑ Easier to secure against side-channel attacks
 - E.g., negligible overhead of first-order masking for arithmetic operations

FrodoKEM: Variants and parameters

- ❑ Two (2) options for generating the public matrix A
 - Uses either AES128 or SHAKE128
- ❑ Two variants depending on the reuse of public keys
 - Standard and ephemeral FrodoKEM
- ❑ Twelve (12) parameter sets in total:

- (e)FrodoKEM-640-XXX: targets security level 1 (\geq AES-128)
- (e)FrodoKEM-976-XXX: targets security level 3 (\geq AES-192)
- (e)FrodoKEM-1344-XXX: targets security level 5 (\geq AES-256)

Dimension $n \in \{640, 976, 1344\}$, $XXX \in \{\text{AES}, \text{SHAKE}\}$

Performance profile

- ❑ On an x64 AMD Ryzen 9 3900XT @3.8GHz:
 - One full FrodoKEM execution (at level 1) is completed in **0.79 msec.**
 - Encaps + Decaps runs in **0.55 msec.**
- ❑ Public keys and ciphertexts are roughly 9KB, 15KB and 21KB for levels 1, 3 and 5 (resp.)
- Speed and bandwidth comparison with [ML-KEM](#): FrodoKEM roughly **one order of magnitude slower and larger**
- In comparison with [Classic McEliece](#): FrodoKEM public keys are **more than an order of magnitude smaller**, Encaps + Decaps is **an order of magnitude slower**, full FrodoKEM run is **more than an order of magnitude faster**

LWE security profile

- Quantum security estimates based on core-SVP hardness (in log2)

Level	ML-KEM	FrodoKEM
1	107	137
3	165	195
5	232	255

- Classical security estimates in terms of gates (in log2) and memory (bits in log2)

Level	ML-KEM		FrodoKEM	
	gates	mem	gates	mem
1	152	94	175	110
3	215	139	240	156
5	287	190	305	202

Summary

- ❑ FrodoKEM is practical for many applications
 - Takes more running time and bandwidth than ML-KEM
 - Brings attractive trade-offs in comparison with Classic McEliece
- ❑ It offers several attractive features:
 - Conservative LWE security, and parameters with larger (classical and quantum) security margins against known attacks
 - Enables simple, compact and easier-to-protect implementations
- **FrodoKEM** is a great alternative for security-sensitive applications

Proposal for Internet-Draft

- ❑ We have a draft ready that is in good shape and is aligned with the upcoming ISO standard amendment
 - It can be improved further with the help of the CFRG community
- ❑ Editors:
 - Joppe Bos (NXP)
 - Stephan Ehlen (BSI)
 - Patrick Longa (Microsoft)
 - Douglas Stebila (University of Waterloo)

Links

- ❑ Website (specification and updates):

<https://frodokem.org/>

- ❑ Official implementation:

<https://github.com/microsoft/PQCrypto-LWEKE>

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