

Insights into BBRv3's Performance and Behavior by Experimental Evaluation

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Experimental Evaluation – Our Contributions

- Real hardware (no emulation, no virtualization)
- Bottleneck data rates 100 Mbit/s – 10 Gbit/s
- Traced BBR internal variables → root cause analysis

1) Self-Induced Queuing Delay

This talk

2) Delay-Jitter

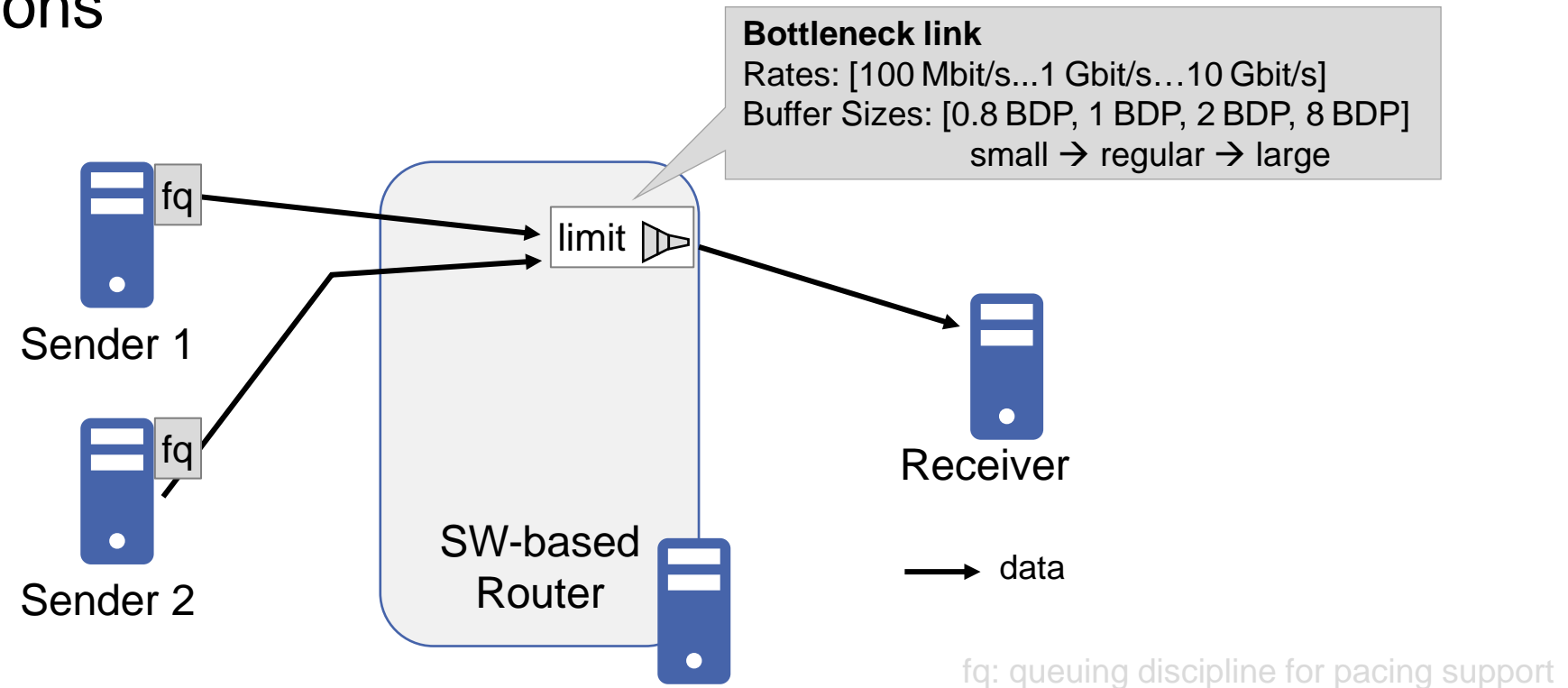
See paper

3) Short-flows using Real-World Traffic

This talk

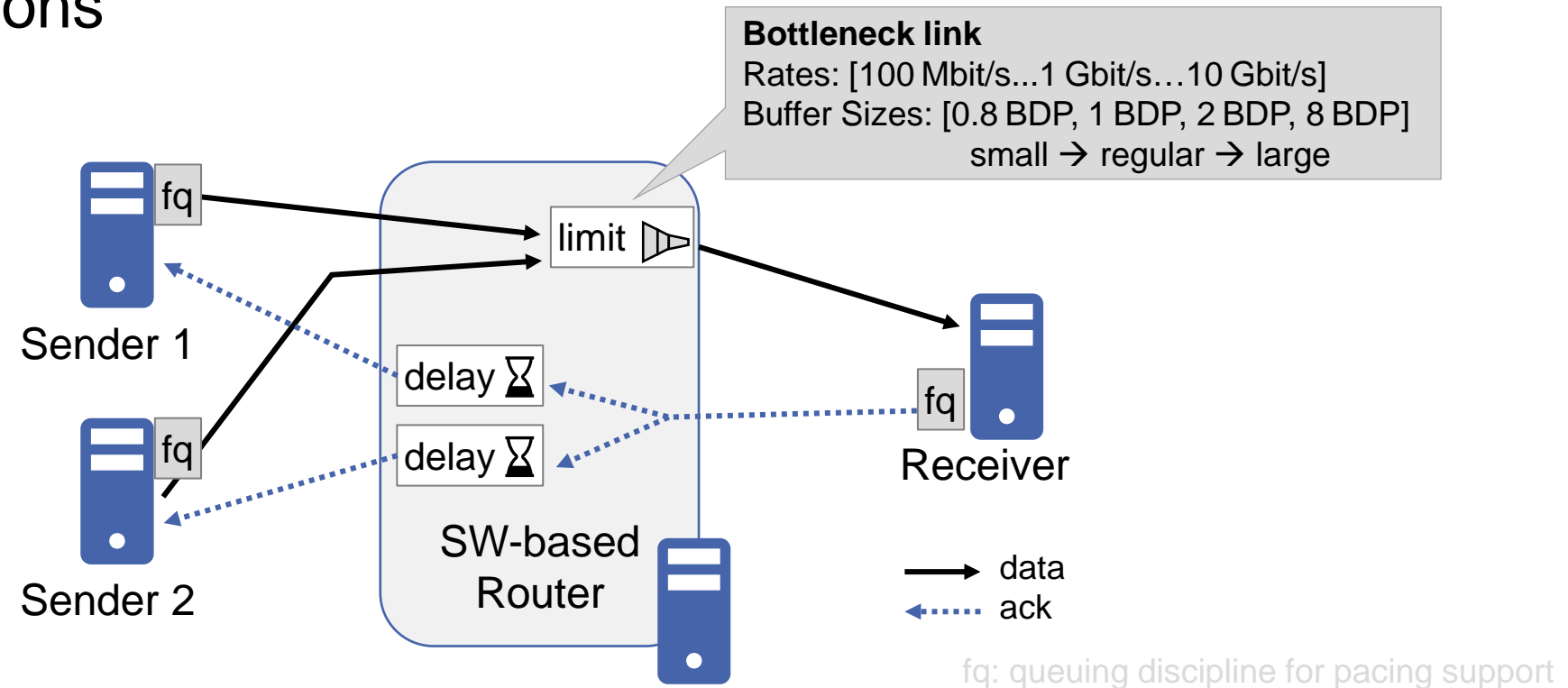
Testbed Setup

- Linux-based testbed with 100 Mbit/s – 10 Gbit/s
- Using Google's BBRv3 implementation
- 60s runs, 30 repetitions



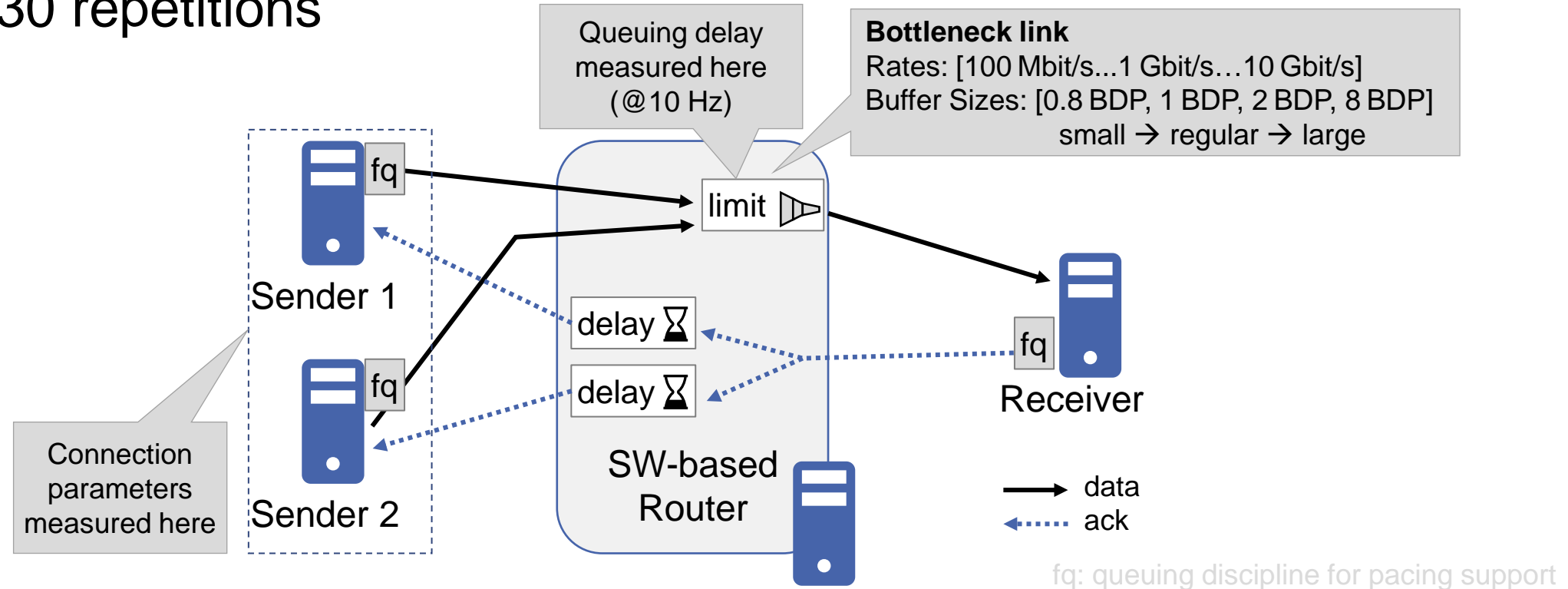
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Experimental Results

1) Self-Induced Queuing Delay

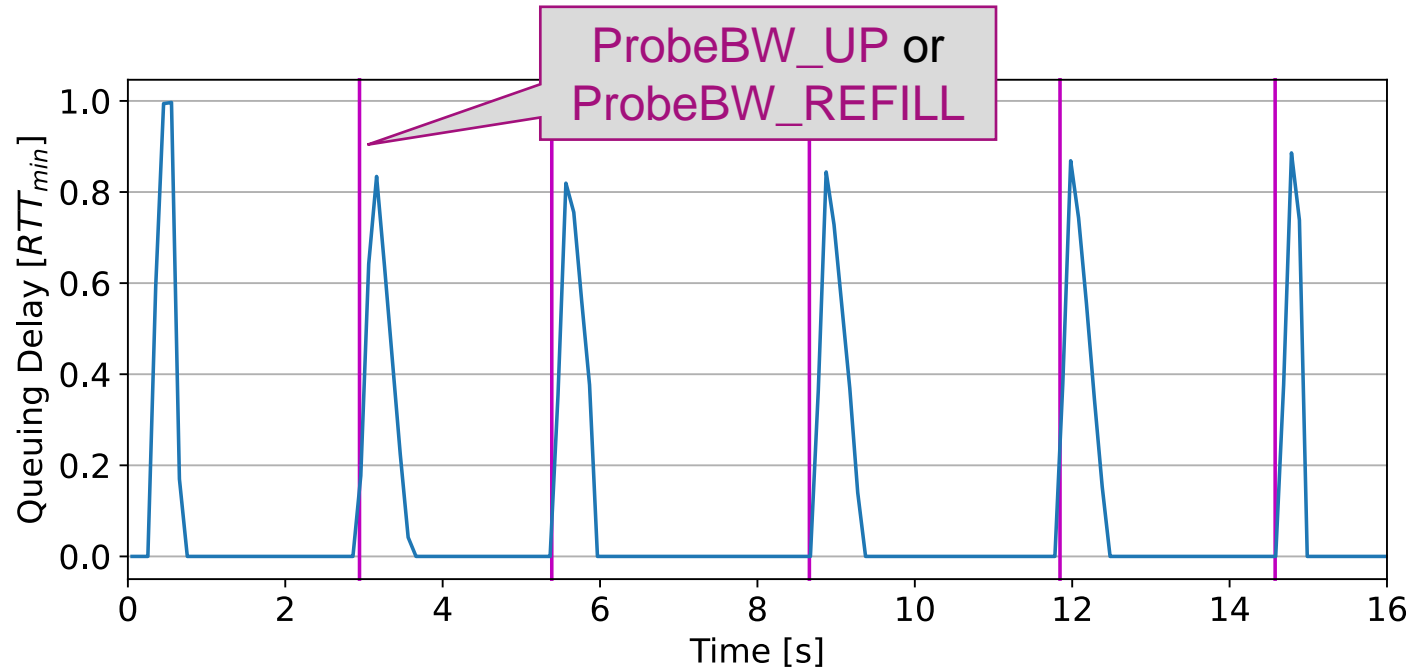
NOW

(traffic **not** limited
by application)

2) Delay-Jitter

3) Short-flows using Real-World Traffic

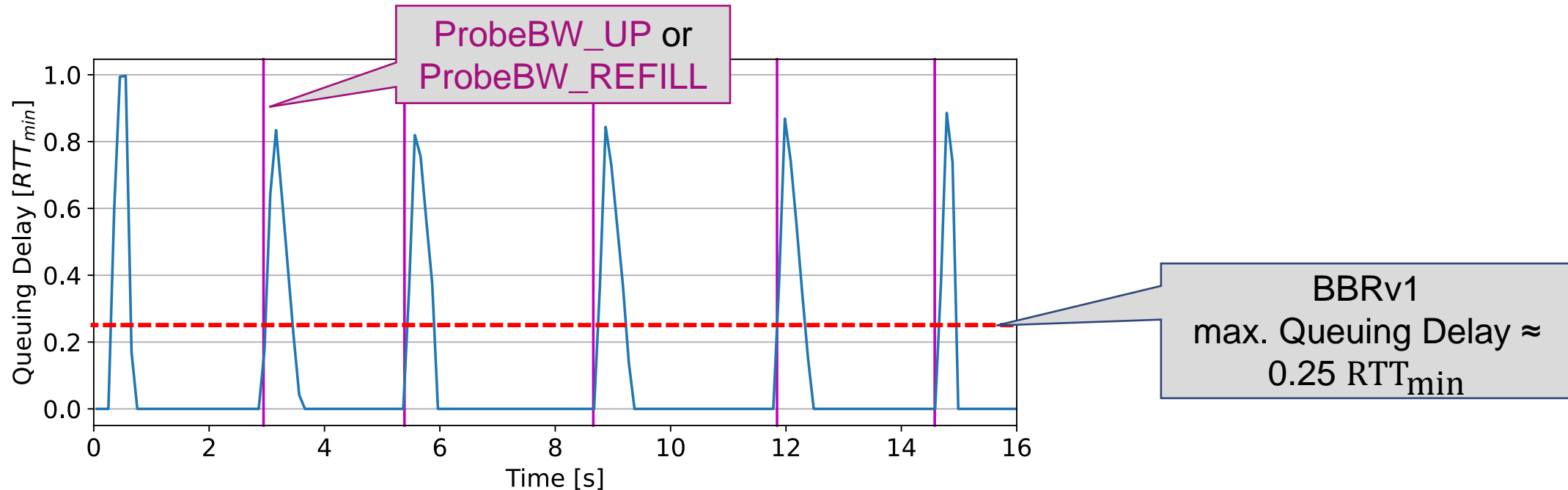
Self-induced Queuing Delay – Single Flow



100 Mbit/s, $RTT_{min} = 50$ ms, buffer size = 8 BDP

- **1st Cause:** ProbeBW_UP Phase → estimated bandwidth · 1.25, for 3 RTTs
- Self-induced queuing delay prolongs round-trips

Self-induced Queuing Delay – Single Flow

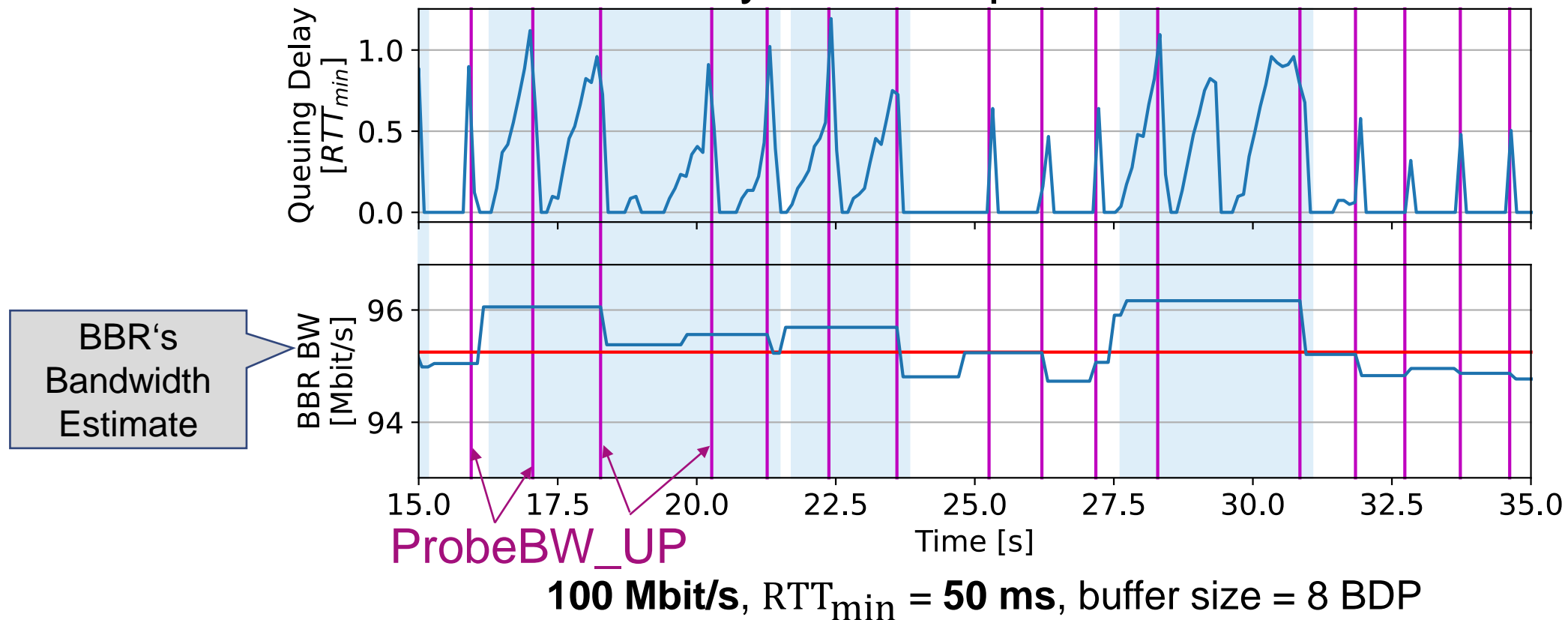


100 Mbit/s, $RTT_{min} = 50$ ms, buffer size = 8 BDP

- **1st Cause:** ProbeBW_UP Phase \rightarrow estimated bandwidth $\cdot 1.25$, for 3 RTTs
- Self-induced queuing delay prolongs round-trips
- BBRv3 creates higher peaks – not present to that extent with BBRv1

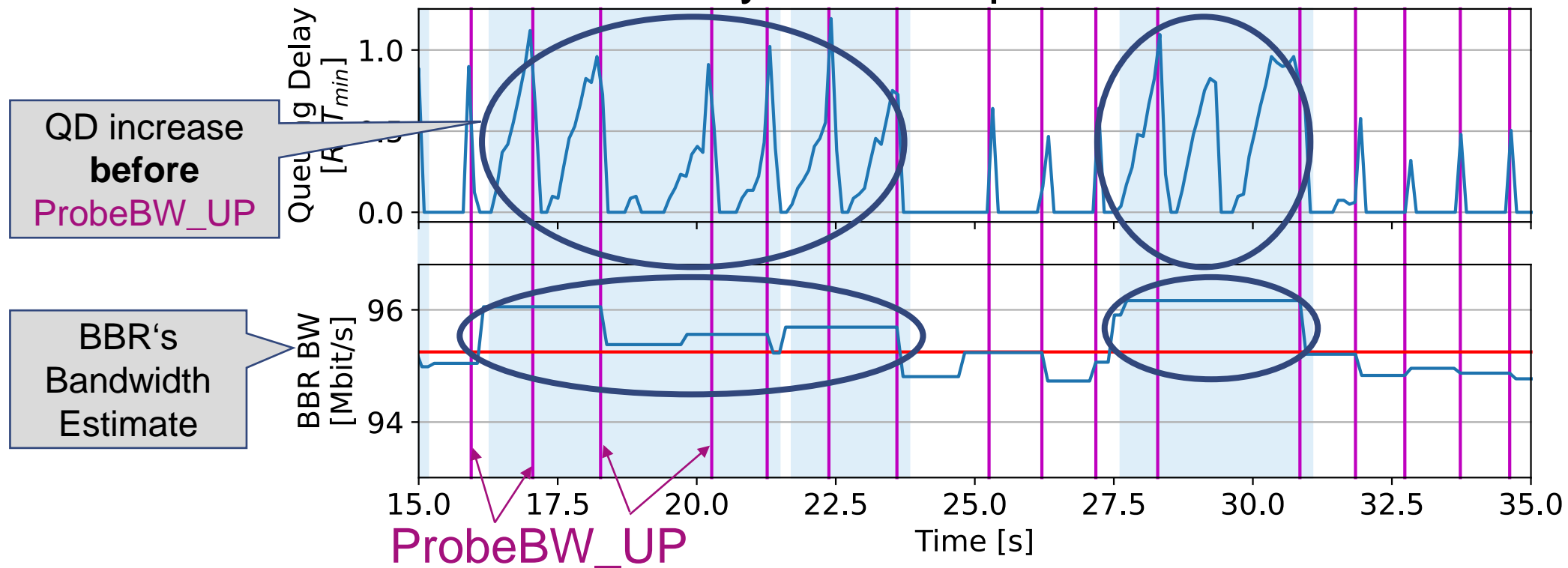
Self-induced Queuing Delay – Single Flow

- 2nd Cause: Bandwidth overestimation due to natural jitter
- Measurement noise in delivery rate samples



Self-induced Queuing Delay – Single Flow

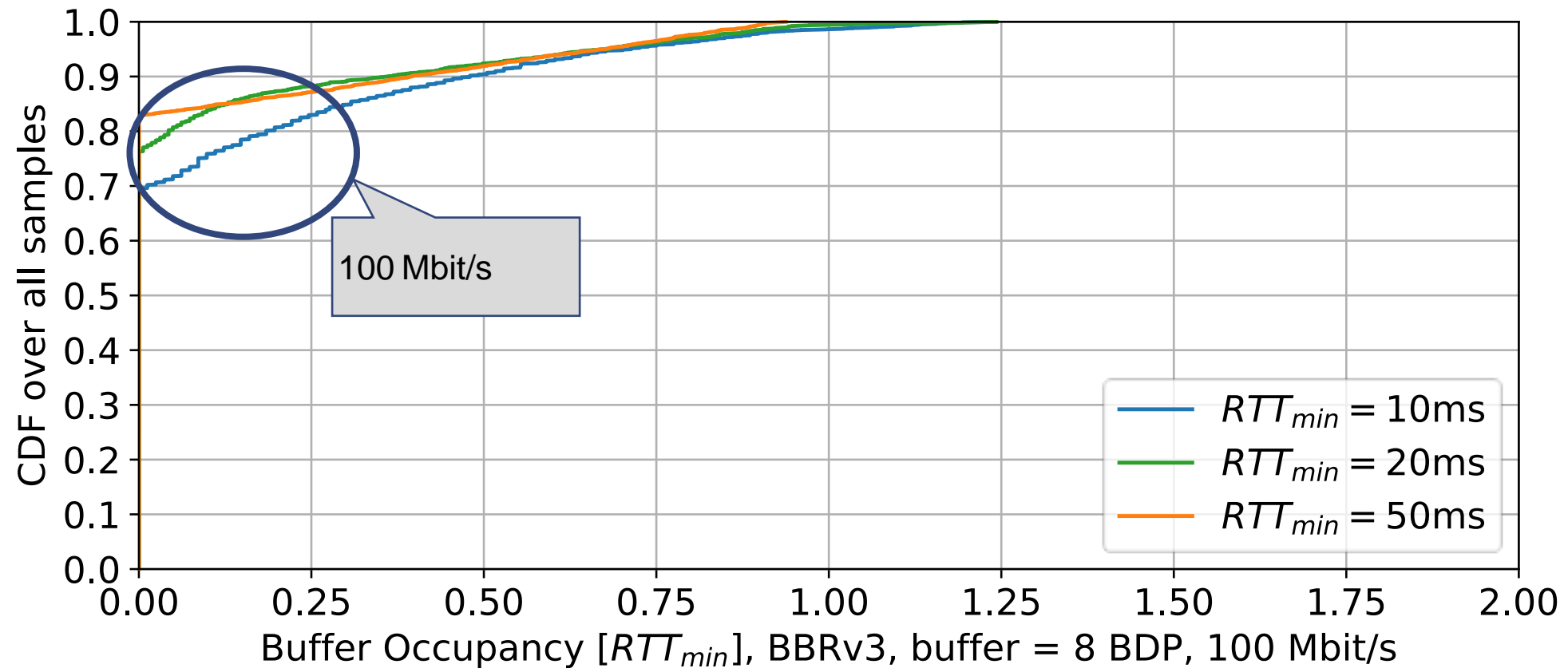
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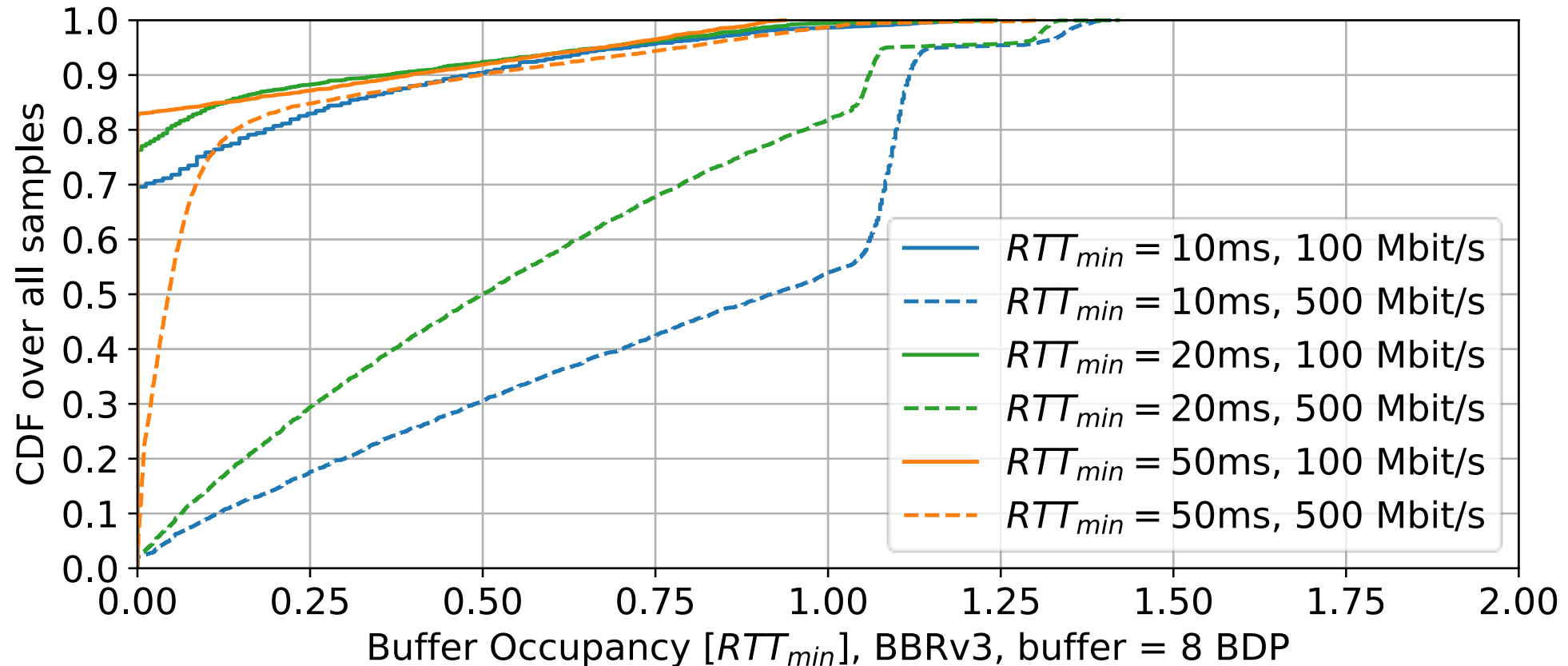
Self-induced Queuing Delay – Single Flow

- It gets worse at higher speeds and smaller RTT_{min} values...



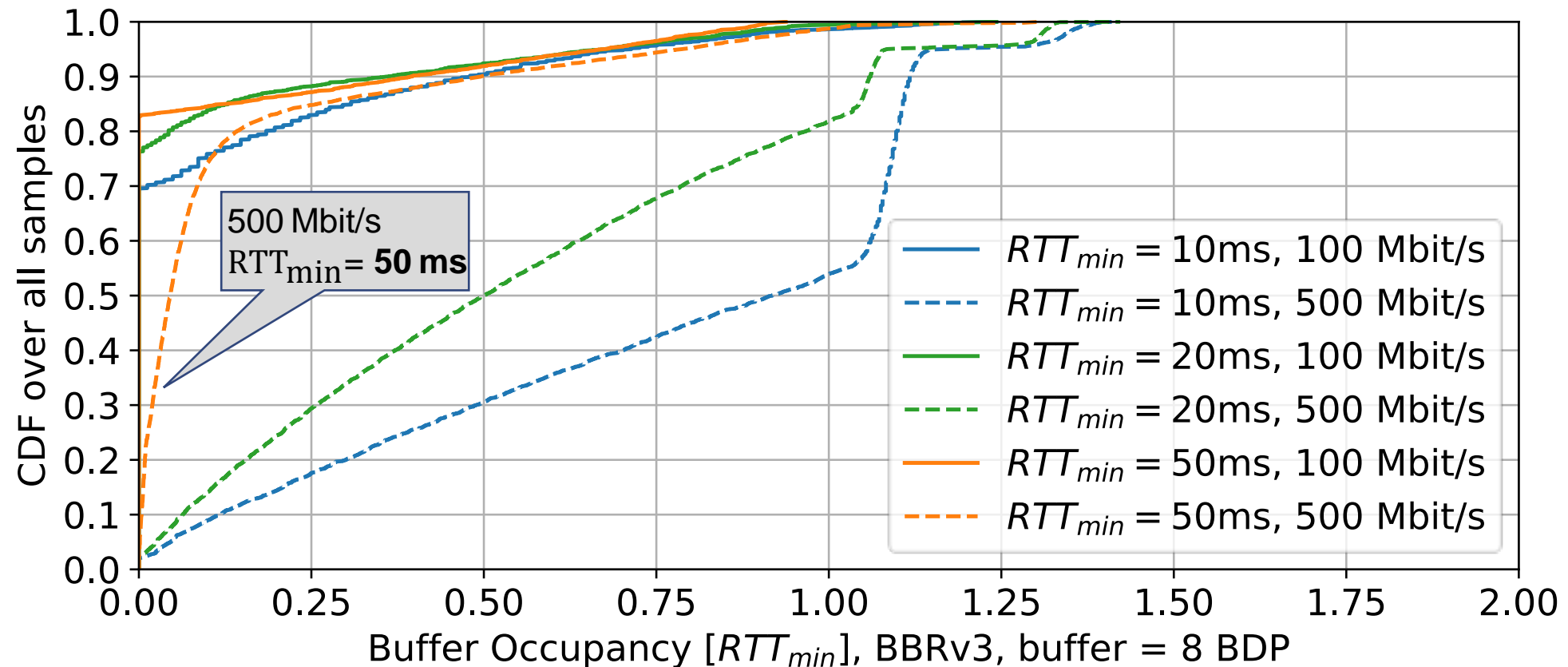
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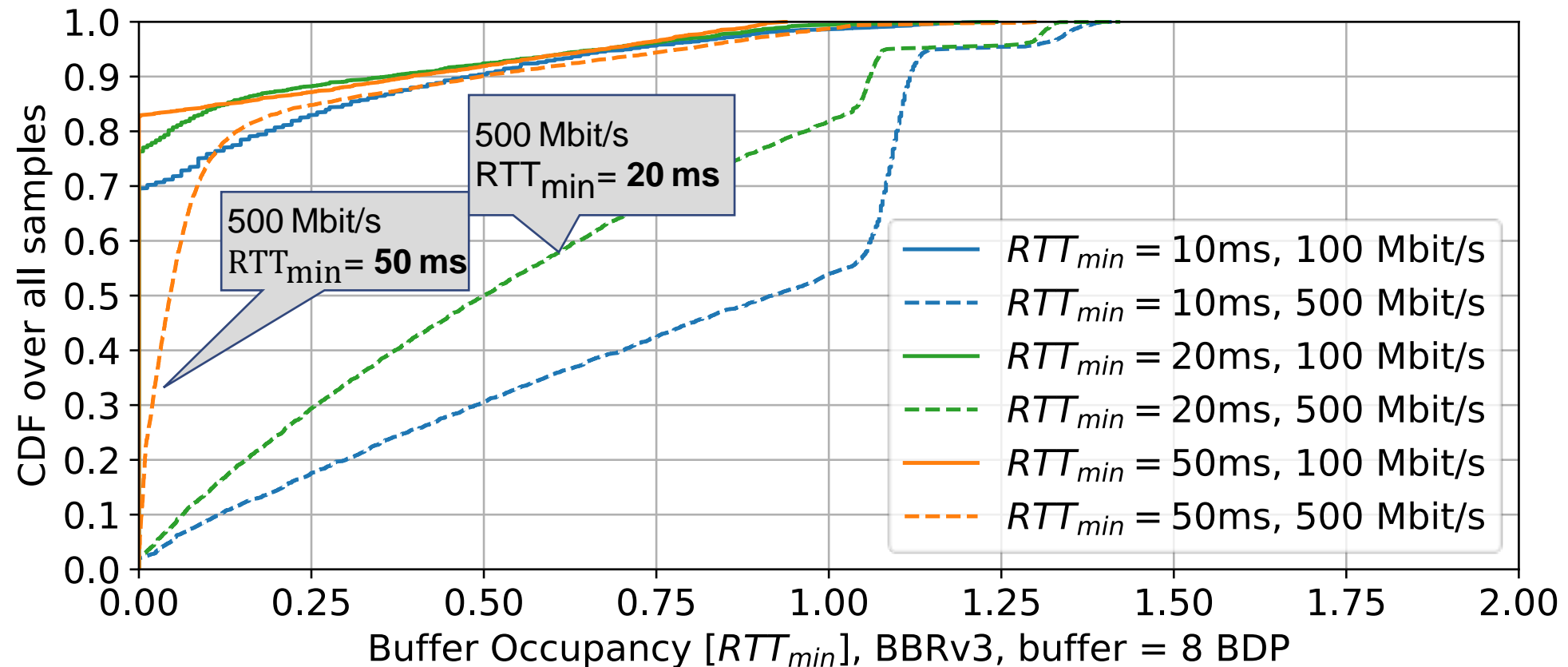
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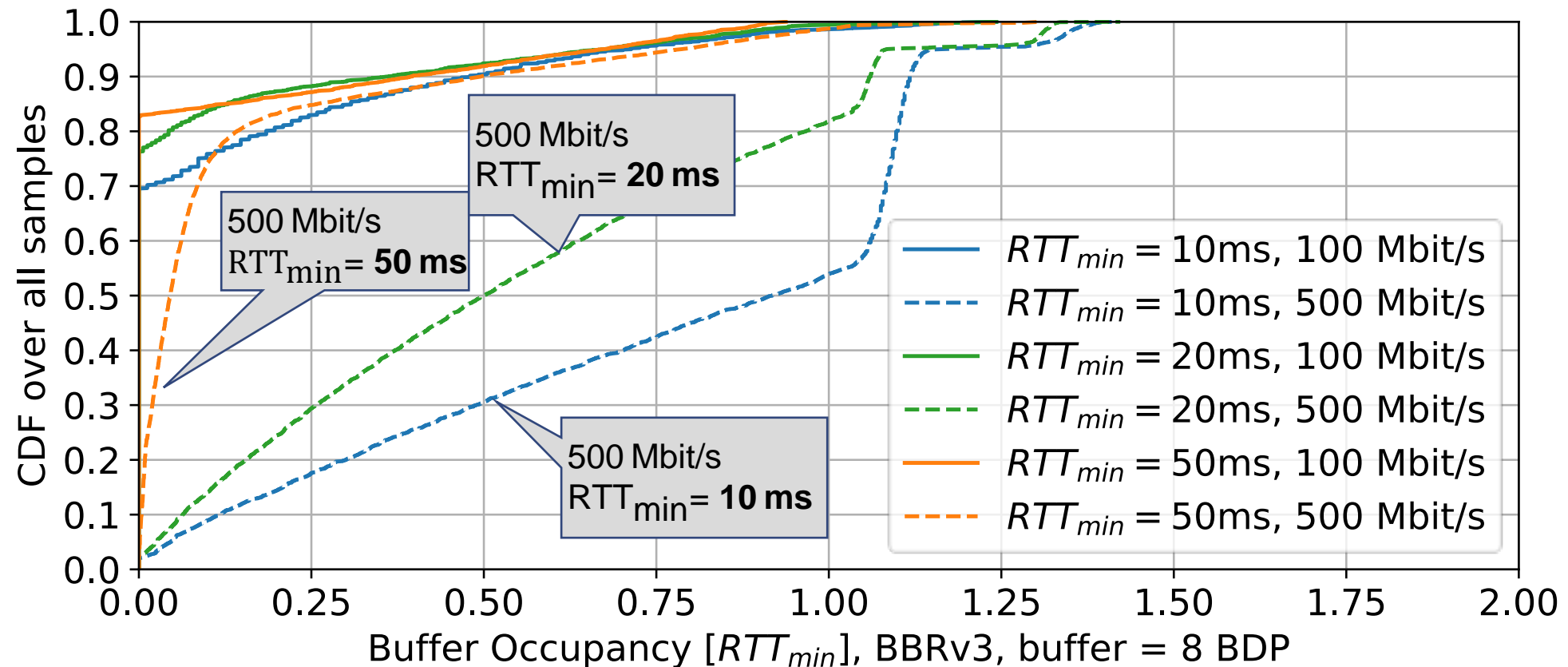
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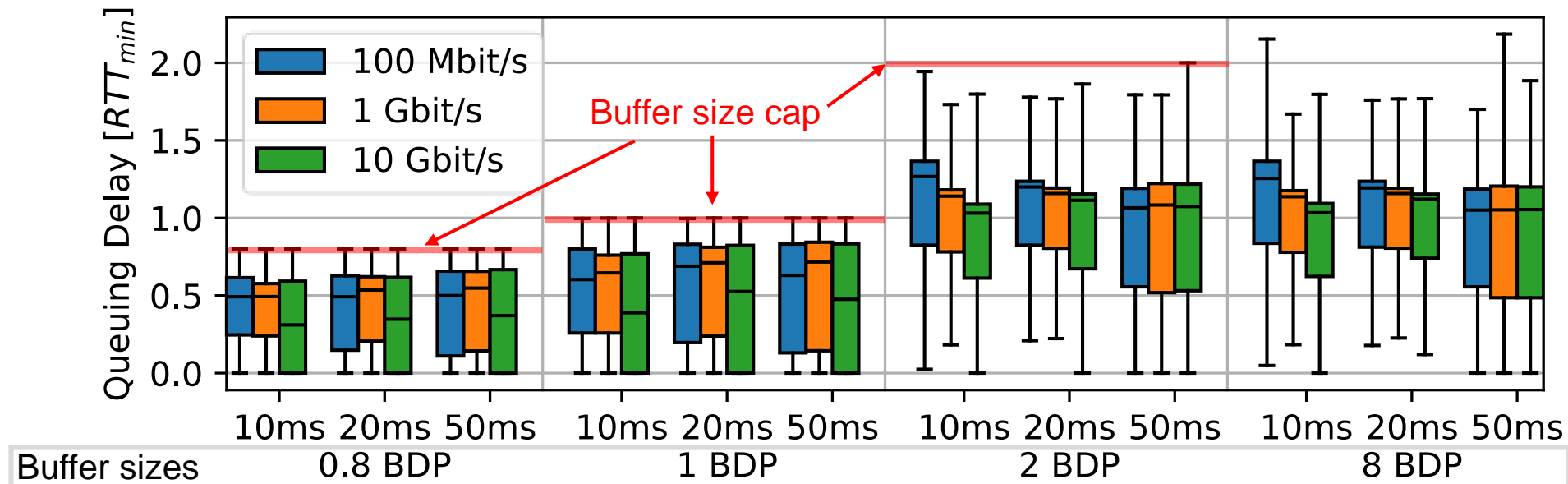


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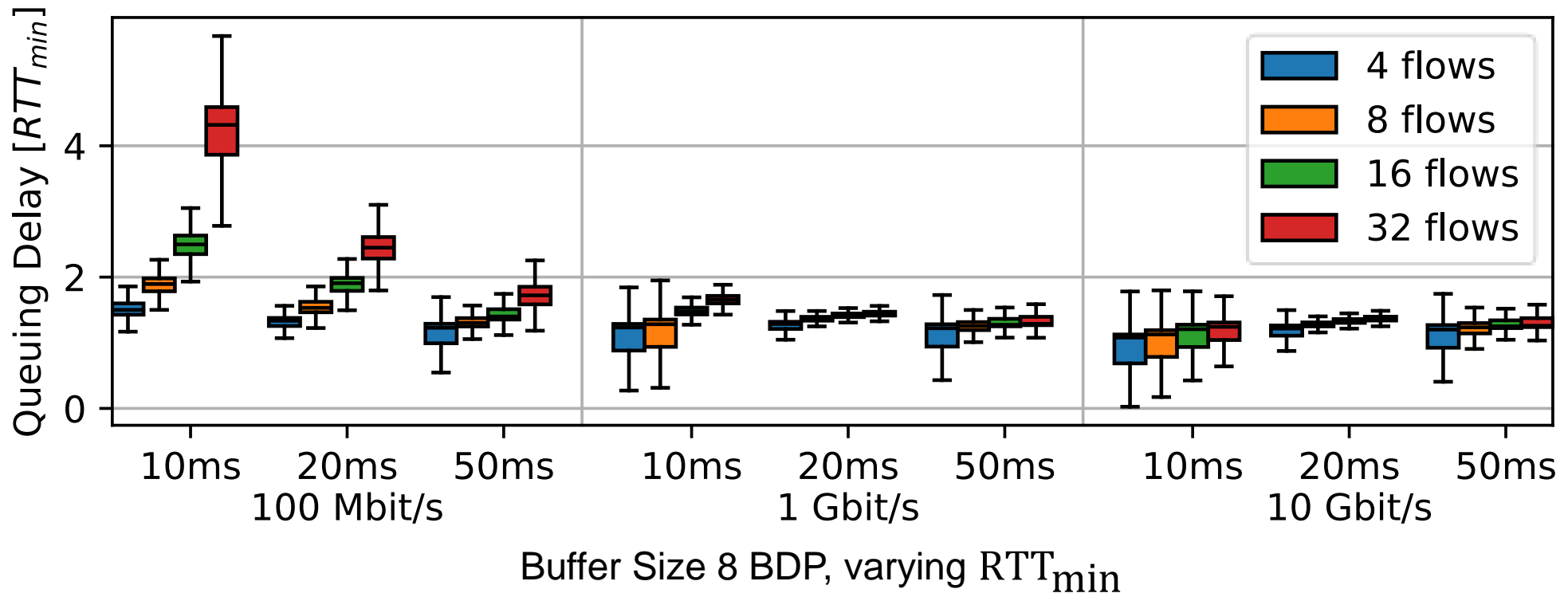
Self-induced Queuing Delay – Two Flows



- Queuing Delay around 1 RTT_{min} for >50% of the time for larger buffers
- **Cause:** overestimation of available bandwidth in ProbeBW_UP (present since BBRv1)

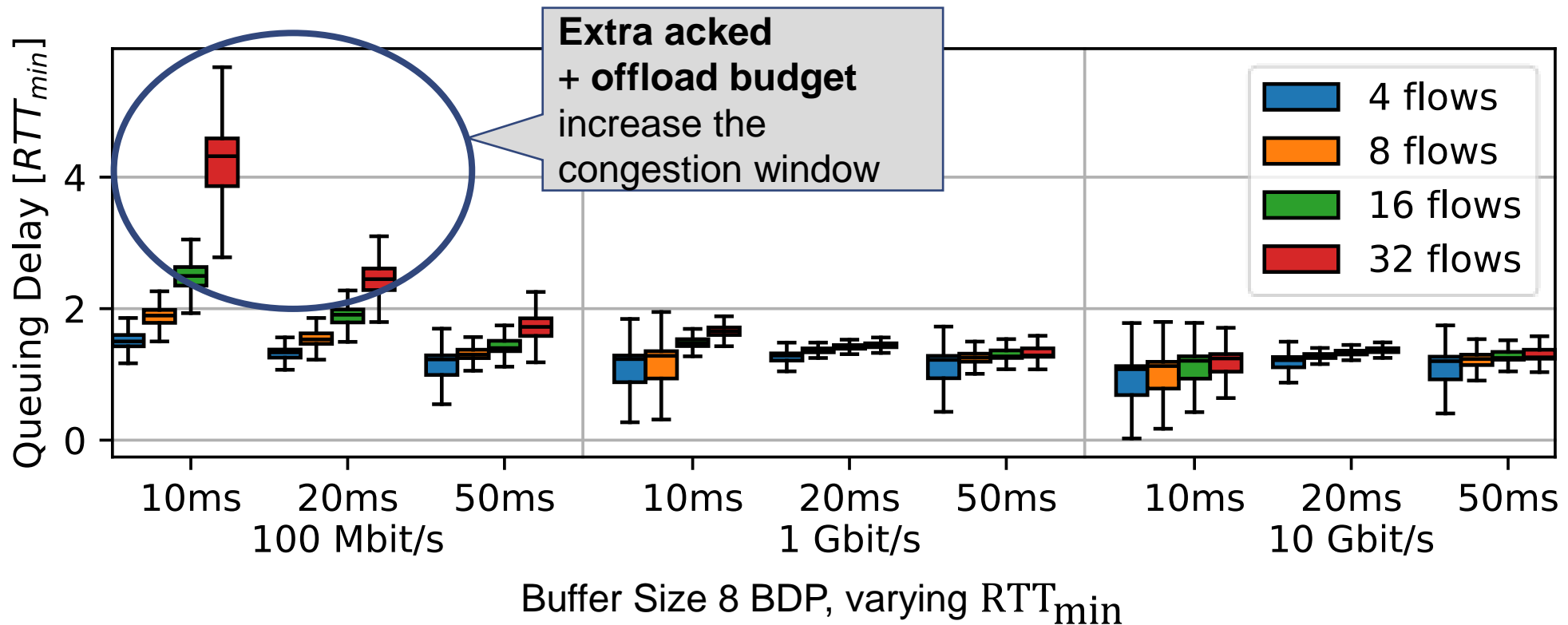
Self-induced Queuing Delay – Multiple Flows

- Gets a bit worse with increasing number of flows



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Experimental Results

1) Self-Induced Queuing Delay

2) Delay-Jitter

See paper

3) Short-flows using Real-World Traffic

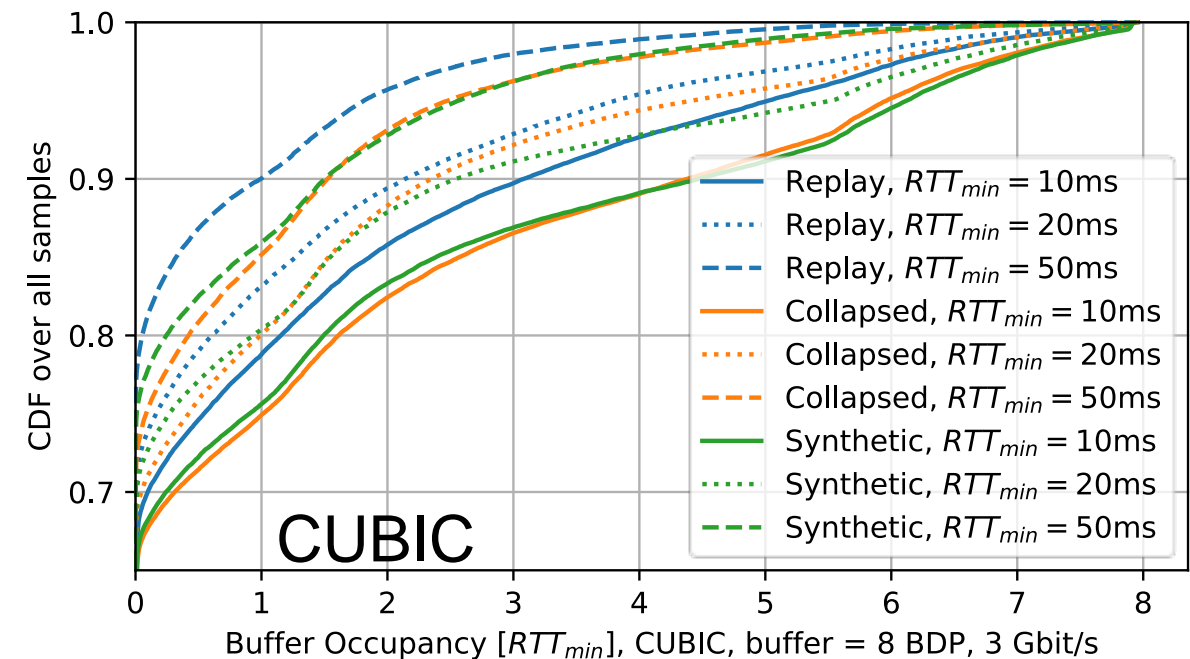
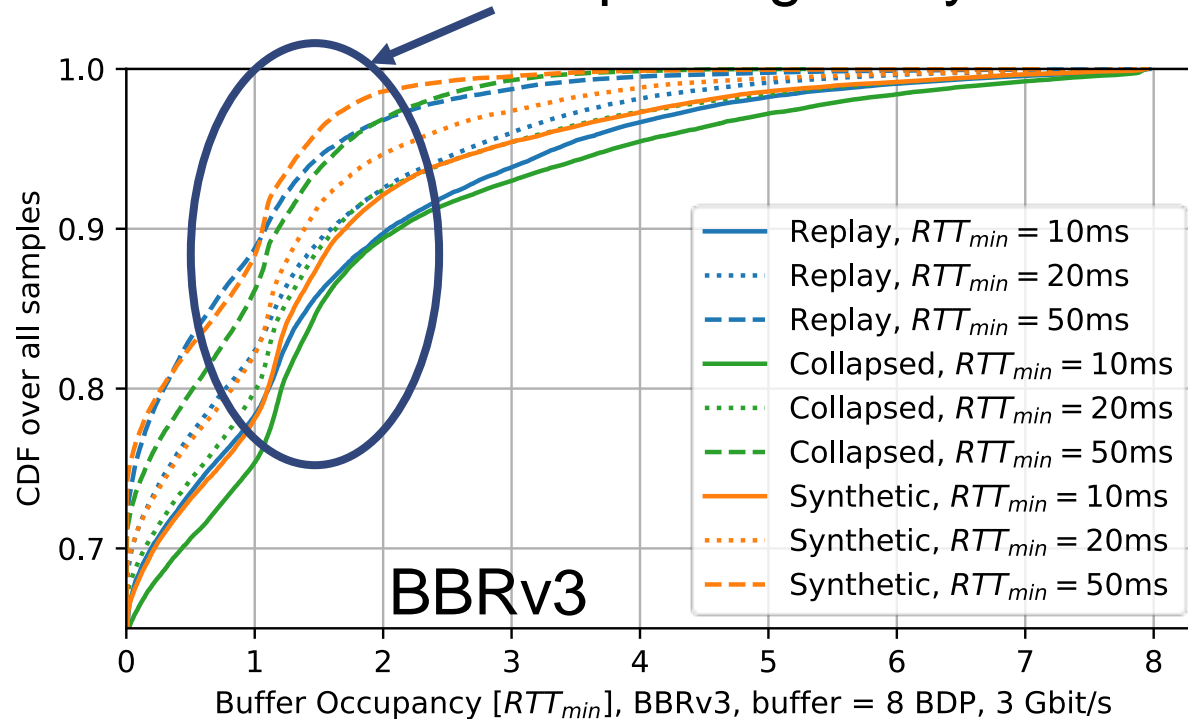
NOW

(traffic often limited
by application)

Continuous application-unlimited flows → short application-limited flows?

Short Flows – Real World Trace

- 15min MAWI trace (2020/06/10, 10 Gbit/s link) → Replay at 3 Gbit/s
- BBR avoids to fill deep buffers completely
- BBR often limits queuing delay to $2 RTT_{min}$



Conclusions

- We also confirmed other known fairness issues
 - RTT unfairness (large RTT_{min} flows prevail), unfairness to CUBIC
 - Long convergence to fairness

1) Self-Induced Queuing Delay

- BBR creates noticeable self-inflicted queuing delay
 - around 1 RTT_{min} for a single BBR flow
 - $>1 RTT_{min}$ more than 50% of the time for multiple flows

2) Delay-Jitter

- see paper: BBR is not as jitter sensitive as predicted by other research

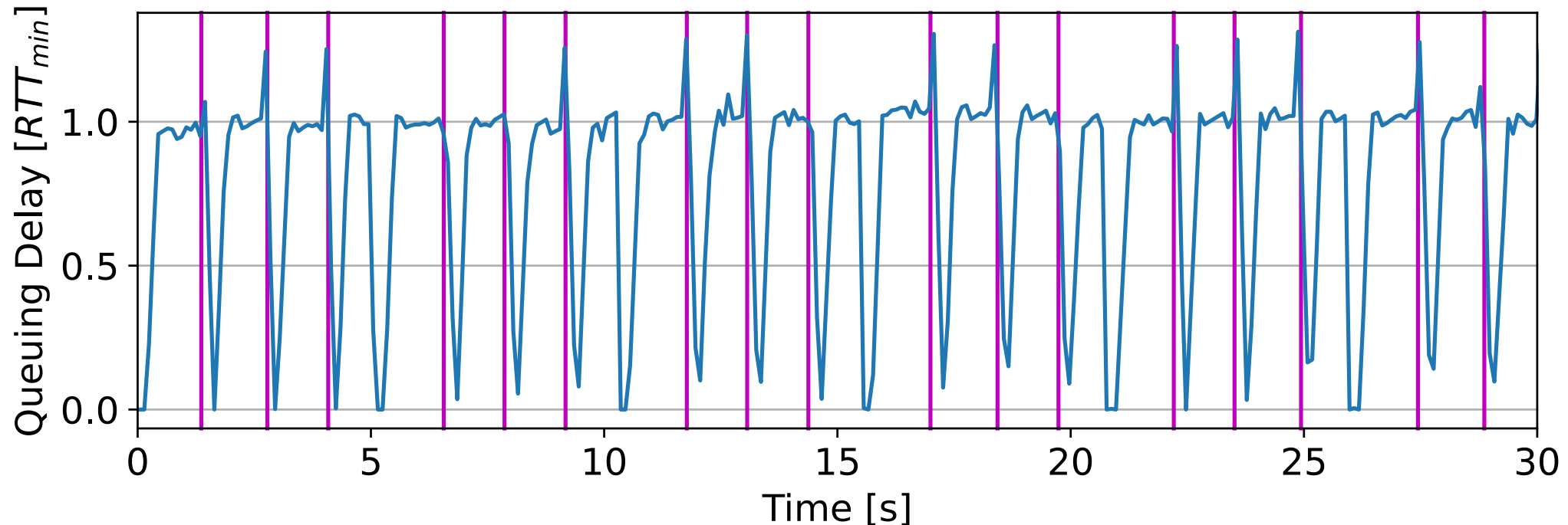
3) Short-flows using Real-World Traffic

- BBR shows slightly improved performance over CUBIC for short-lived flows
- BBR avoids completely filling deep buffers

Backup Slides

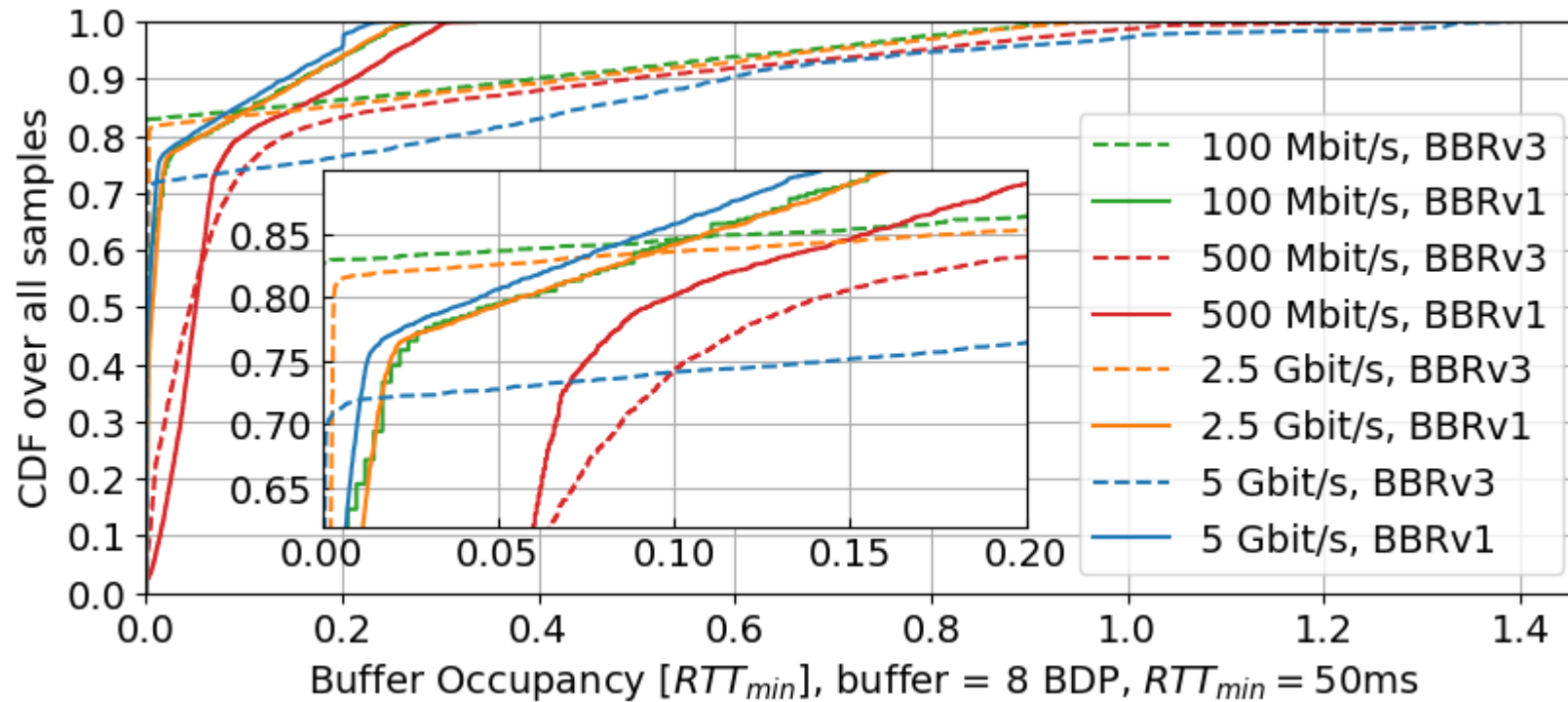
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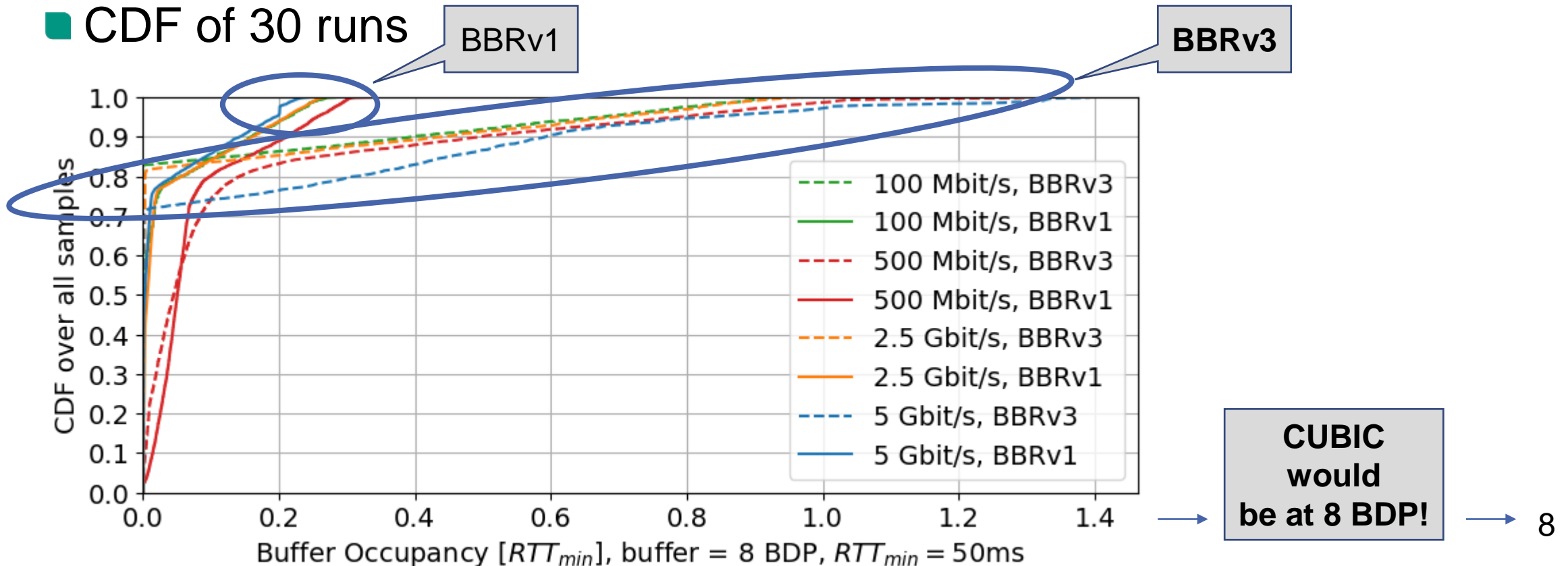
1 Gbit/s, $RTT_{min} = 10$ ms, buffer size = 8 BDP

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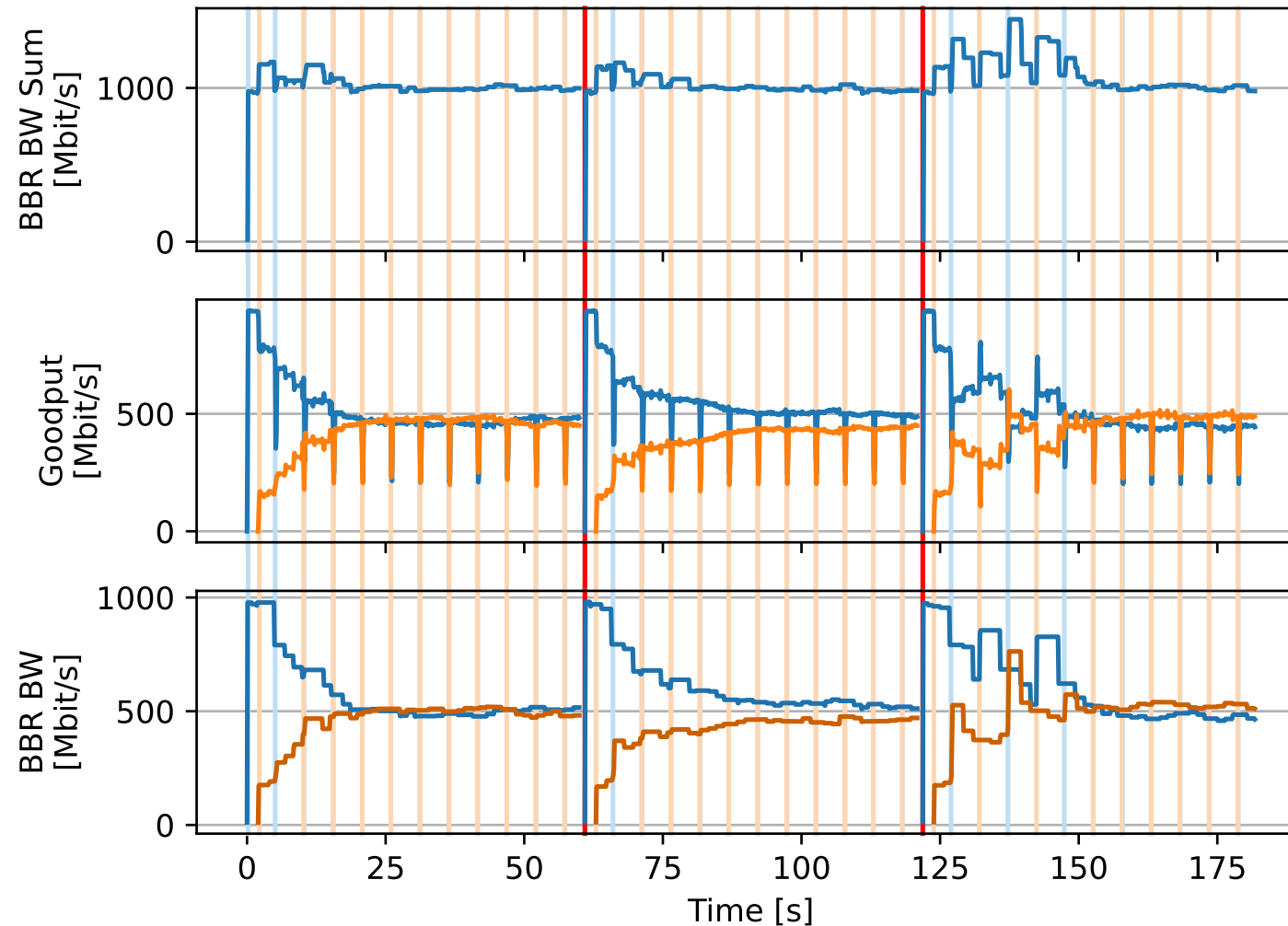
Self-induced Queuing Delay – Single Flow

■ CDF of 30 runs



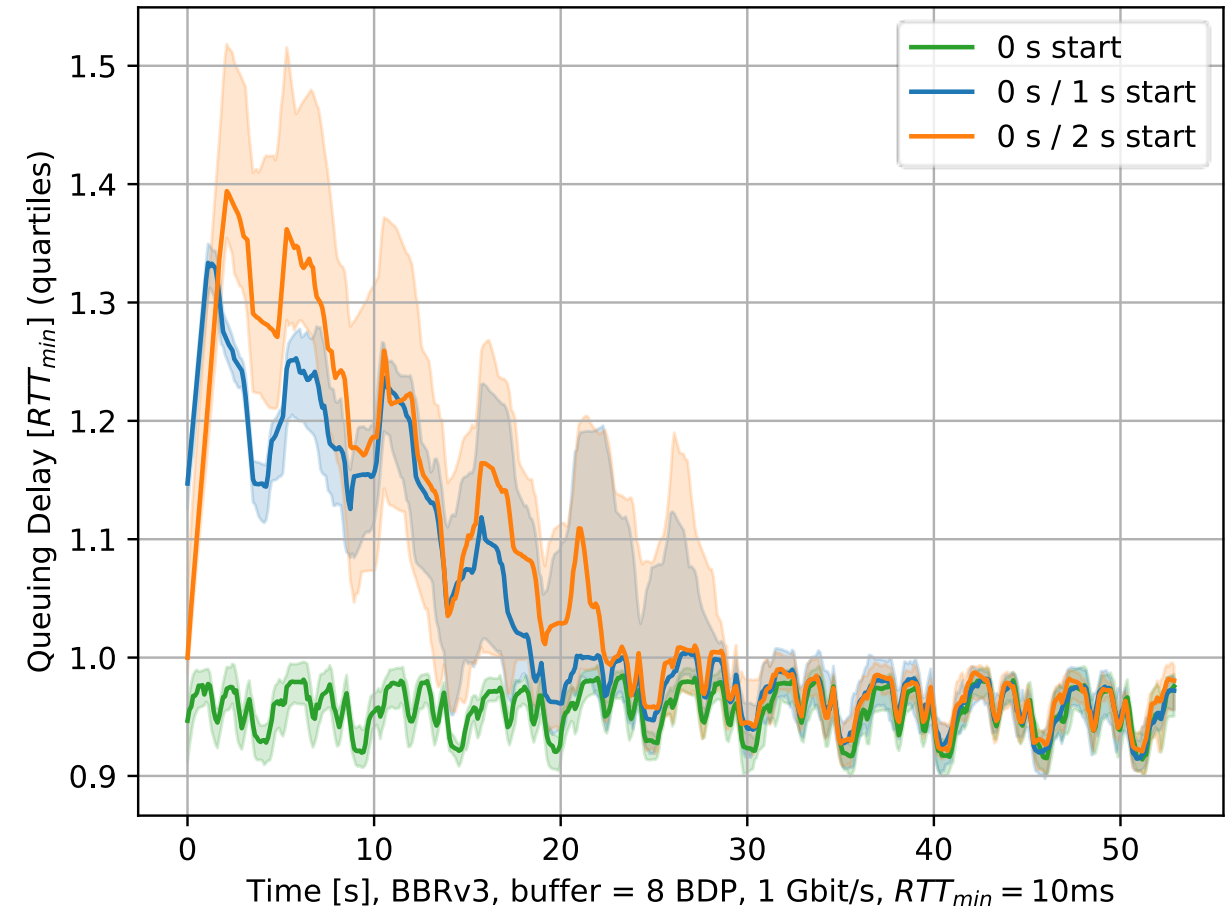
■ BBRv3 creates higher peaks in queuing delay

Bandwidth Overestimation

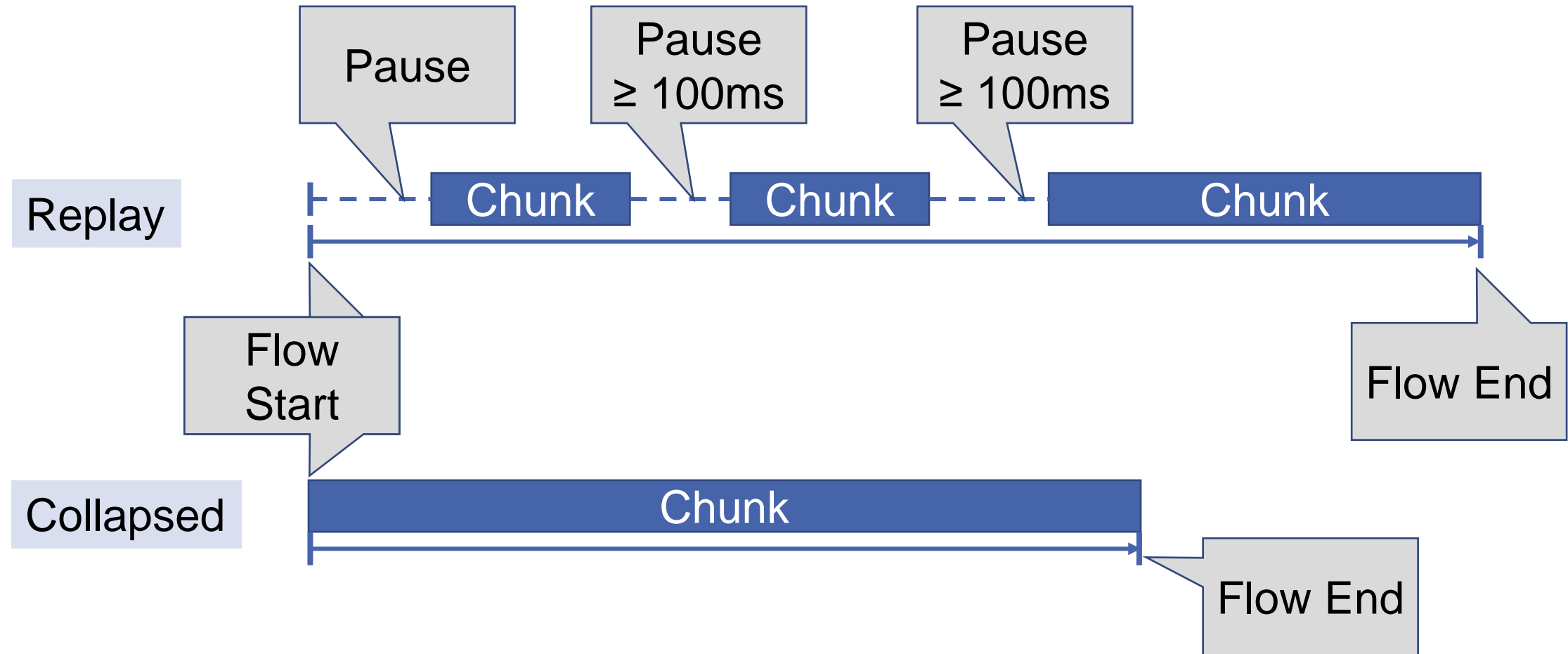


Latecomer Disadvantage

- Long convergence time for flows that start later

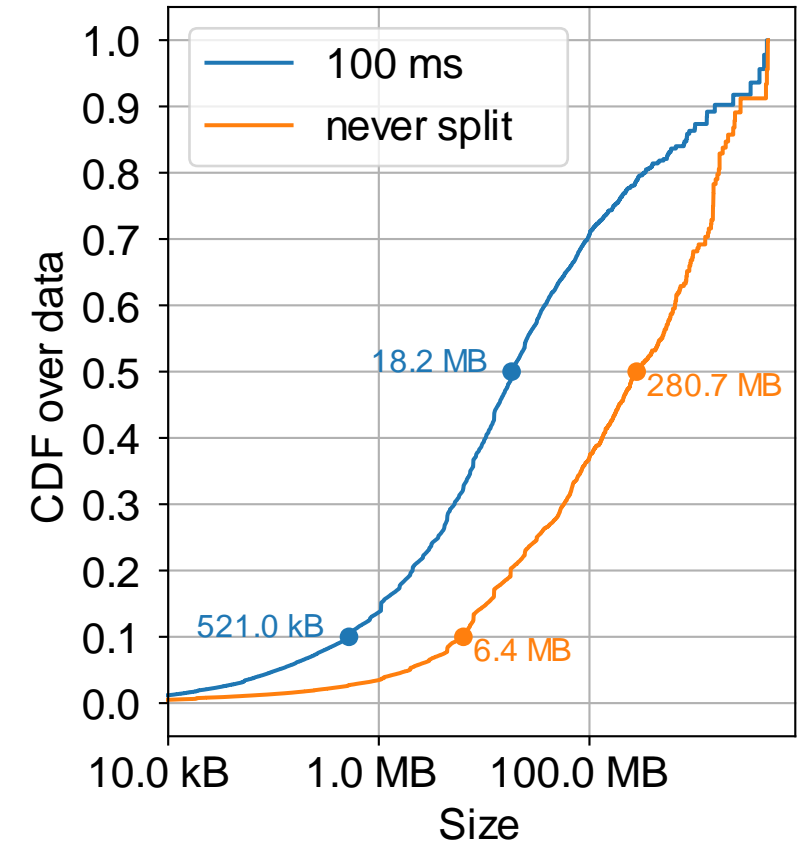
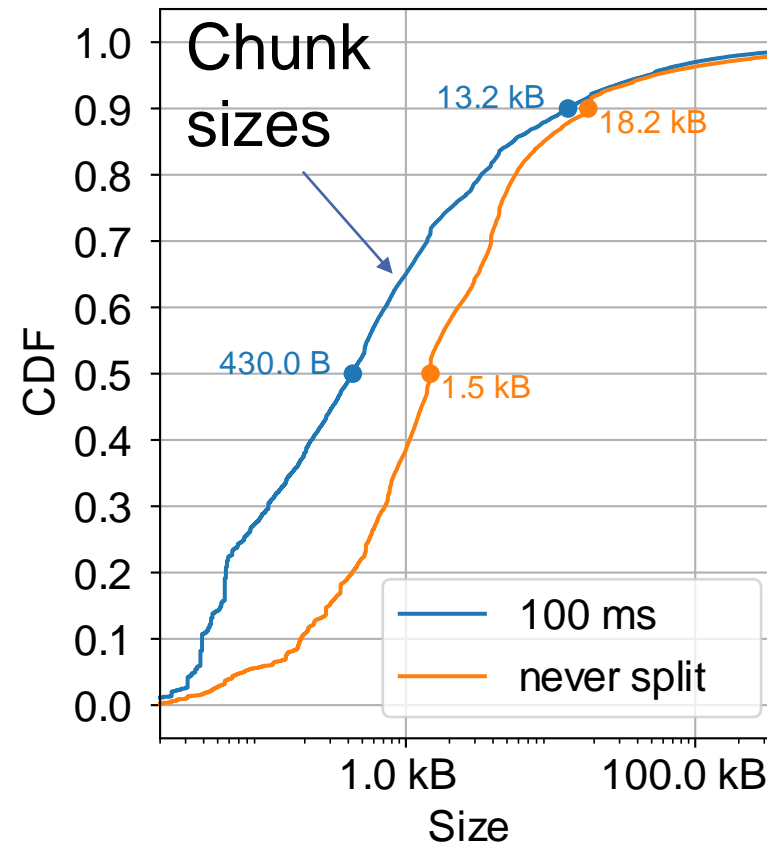


Flow-based Traffic Models



Flow Properties of MAWI Trace

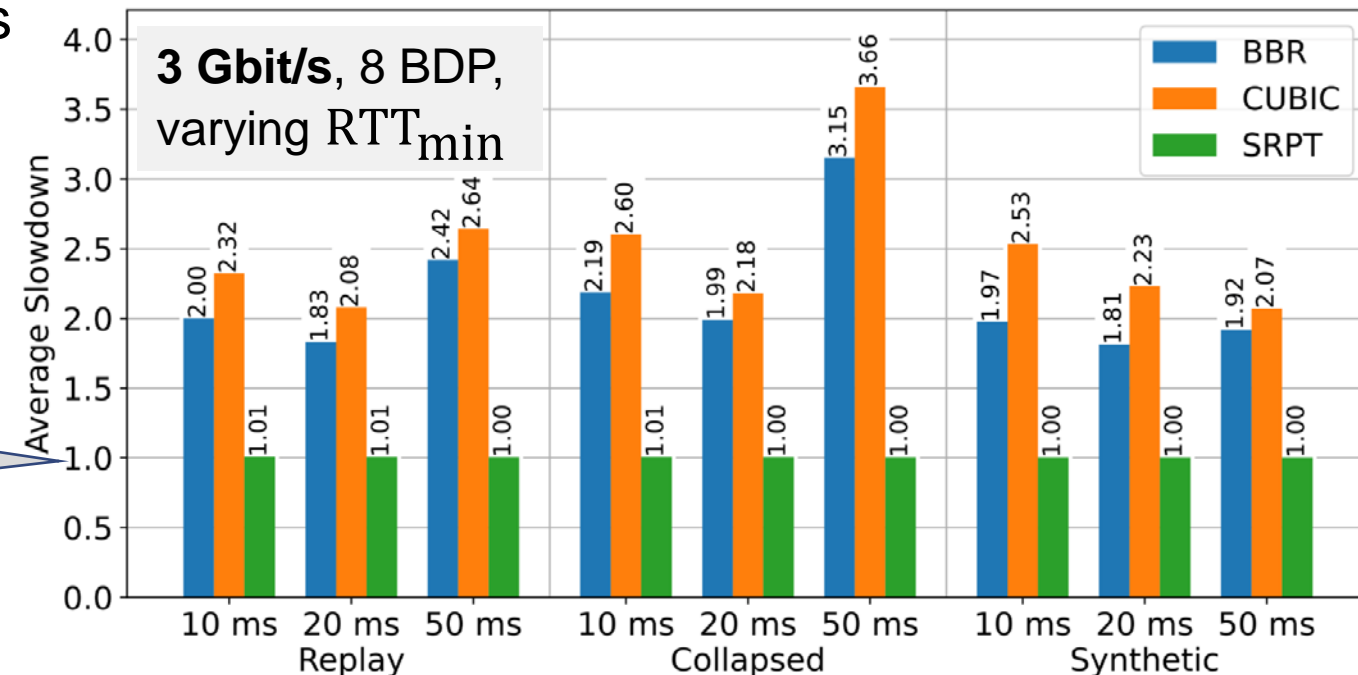
- 90% of flows smaller than 18.2kByte
- If you select a byte randomly, 90% probability that it belongs to a flow >6.4MByte



Short Flows – Real-World Trace

- Used 15min MAWI trace (2020/06/10, 10 Gbit/s link, 890 000 TCP flows)
 - with different replay models at 3 Gbit/s
 - **Replay** considers pauses >100ms
 - **Collapsed** (no pauses)
 - **Synthetic** (independent use of flow inter-arrival time and flow volume distributions)
- **Slowdown: Normalized Flow Completion Time**
 - how well are short flows supported?
- SRPT knows flow duration a priori
 - near optimum scheduling

optimal



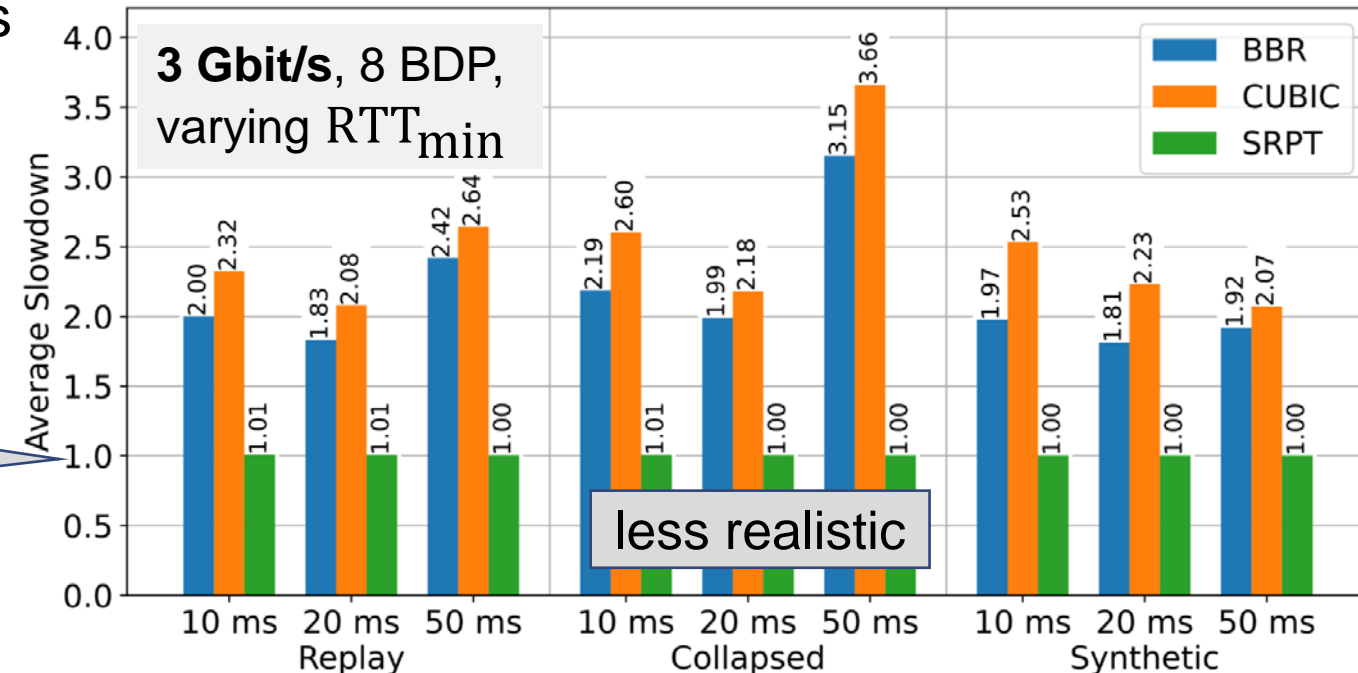
→ BBR performs slightly but consistently better than Cubic

SRPT: Shortest Remaining Processing Time

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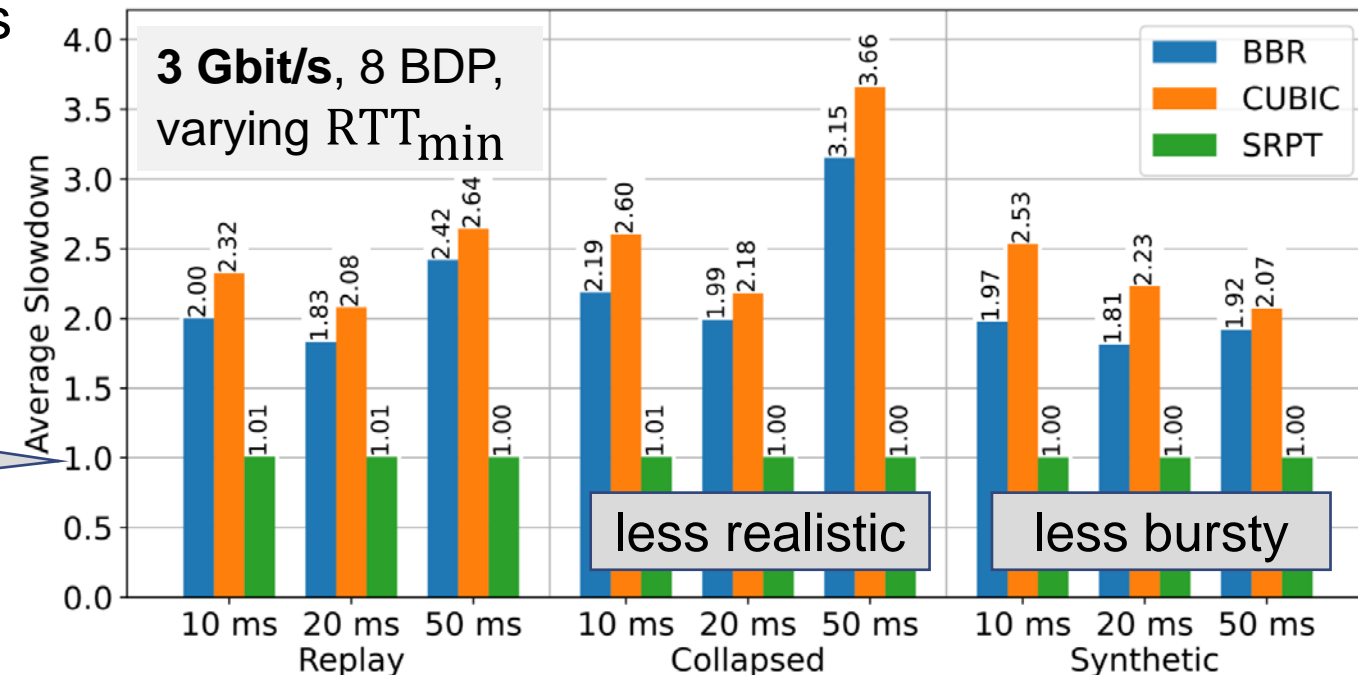


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