

# Post-Quantum EDHOC

## Initiator and Responder using signature and/or KEM

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IETF 125 Shenzhen

draft-papon-lake-pq-edhoc-00

- 1 Context and objectives
- 2 Post-Quantum EDHOC when Initiator knows Responder
- 3 Post-Quantum EDHOC - Authentication with Signature and/or KEM
- 4 Security considerations
- 5 Bytes analysis

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## Context

- Post-quantum secure extension for EDHOC
- Use of post-quantum signatures (same workflow as classical signatures)
- Use of a KEM (different mechanism than Diffie–Hellman)
- Maintain the Identity Protection security property

## Objectives

- Propose a variant of [draft-pocero-authkem-ikr-edhoc-01] where the Initiator authenticates using a PQ-signature
- Propose a post-quantum alternative to EDHOC methods 1–3 (trade-off between number of messages, message size, and computational cost)
- Discussions on security properties (identity protection)

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# Initiator authenticates with a PQ-signature

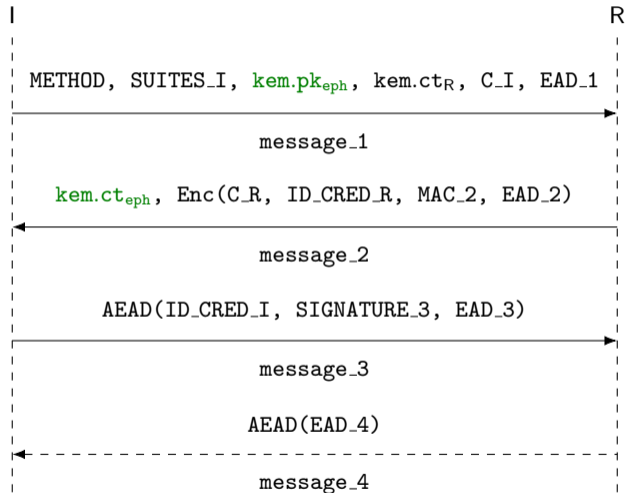


Figure: PQ-EDHOC-IKR, I Signs, R KEM message flow

# Initiator authenticates with a PQ-signature

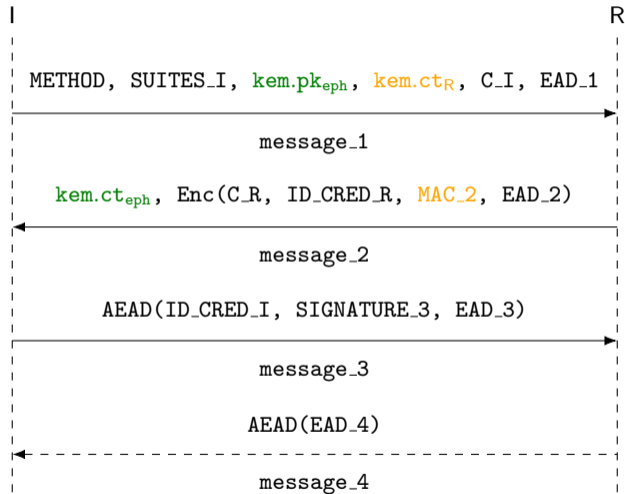


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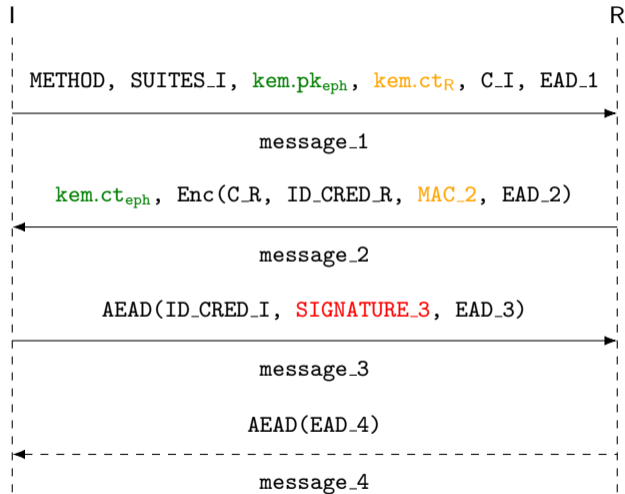


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# Case 1: Initiator Signs, Responder KEM & Signs

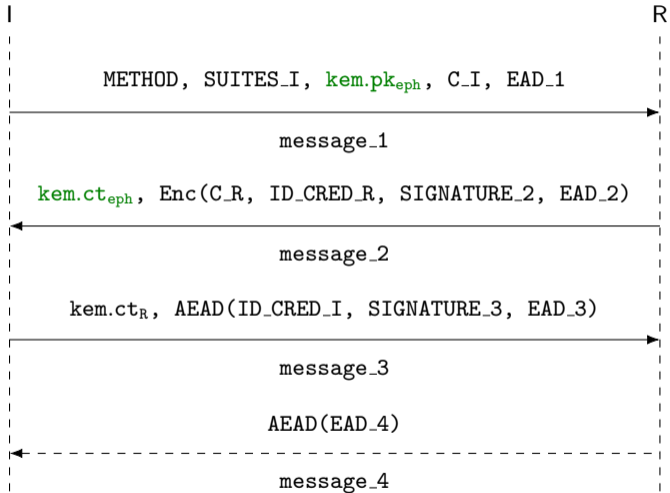


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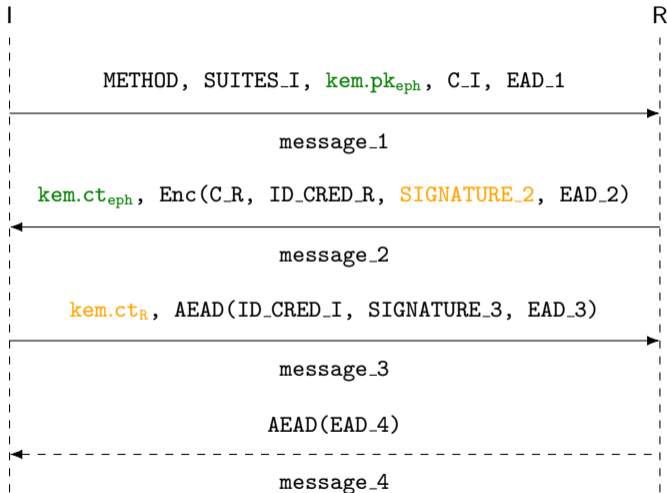


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# Case 1: Initiator Signs, Responder KEM & Signs



Figure: PQ-EDHOC, I Signs, R KEM & Signs message flow

## Case 2: Initiator KEM & Signs, Responder KEM

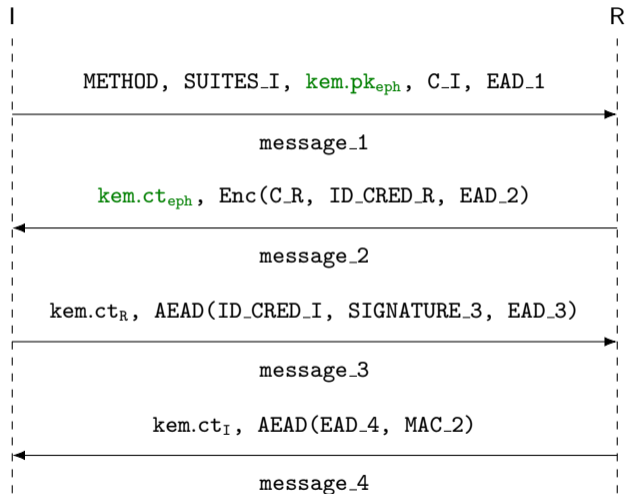


Figure: PQ-EDHOC, I KEM & Signs, R KEM message flow

## Case 2: Initiator KEM & Signs, Responder KEM

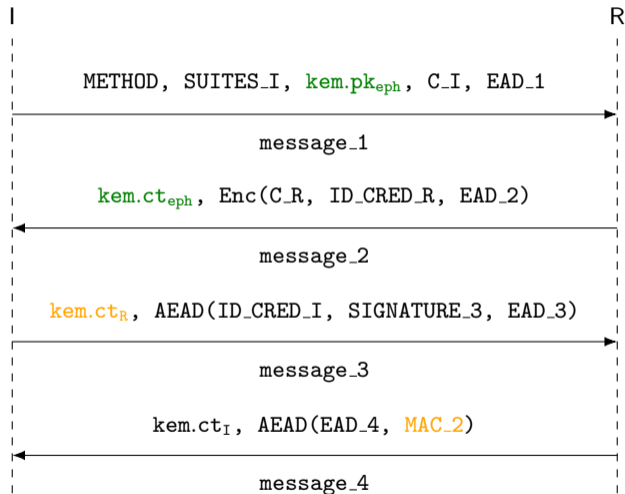


Figure: PQ-EDHOC, I KEM & Signs, R KEM message flow

## Case 2: Initiator KEM & Signs, Responder KEM

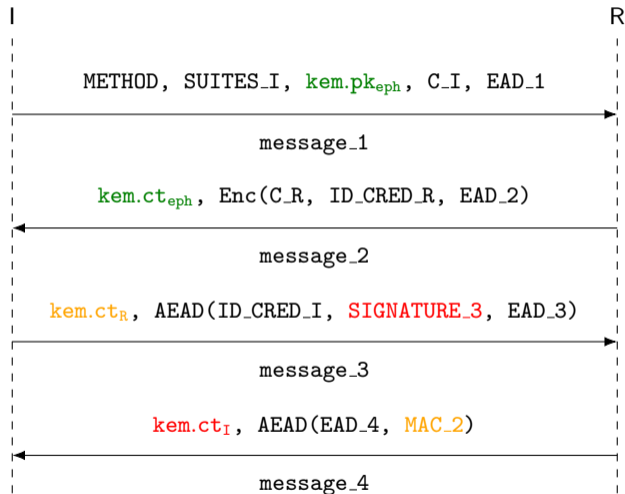


Figure: PQ-EDHOC, I KEM & Signs, R KEM message flow

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- Global security: at the end of each handshake, both endpoints:
  - securely compute the session key  $PRK_{out}$
  - securely authenticate their partner
- Forward secrecy
- Identity protection:
  - each protocol verifies the Identity Protection security property
  - we suggest a more global approach for this property, based on PPAKE (Privacy-Preserving Authenticated Key Exchange)
- Downgrade protection, Transcript hash binding, EAD security

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# Bytes comparison

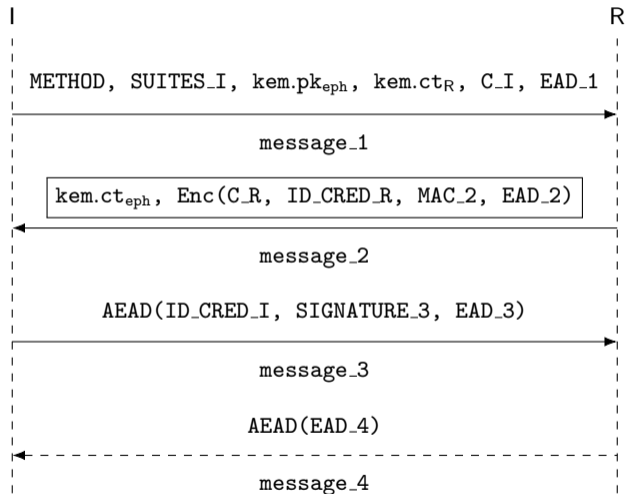


Figure: PQ-EDHOC-IKR, I Signs, R KEM message flow

# Bytes comparison

$$\begin{array}{rcccccc} \text{message\_2} & = & \text{kem.ct}_{\text{eph}} & + & \text{Enc}(\text{C\_R}, & \text{ID\_CRED\_R} & \text{MAC\_2}, & \text{EAD\_2}) \\ & & \downarrow & & \downarrow & \downarrow & \downarrow & \downarrow \\ & & 768 \text{ bytes} & & 1 \text{ byte} & 4 \text{ bytes} & 8 \text{ bytes} & 0 \text{ byte} \\ \\ \text{Total} & = & 781 \text{ bytes} & & & & & \end{array}$$

- Connection Identifier, Method, Suites  $\rightarrow$  1 byte each
- Credentials (kid)  $\rightarrow$  4 or 14 bytes depending on authentication method
- AEAD: AES-CCM-16-64-128  $\rightarrow$  tag = 8 bytes
- MAC  $\rightarrow$  8 bytes
- KEM: ML-KEM-512  $\rightarrow$  pk = 800 bytes ; ciphertext = 768 bytes (FIPS 203)
- DSA: ML-DSA-44  $\rightarrow$  signature = 2420 bytes ; pk = 1312 bytes (FIPS 204)

# Bytes comparison

	Protocol 1	Protocol 2	Protocol 3	Protocol 4	Protocol 5
message_1	1571	803	803	803	803
message_2	781	3203	3203	783	3203
message_3	2442	3210	2442	3210	3202
message_4	0	0	776	784	776
<b>Total</b>	<b>4794</b>	<b>7216</b>	<b>7224</b>	<b>5580</b>	<b>7984</b>

Figure: Bytes analysis of the 5 protocols.

**Thank you!**  
Questions?