
SPINDLE-II Project Overview

Disruption-Tolerant Networking R&D at BBN

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BBN SPINDLE-II Project Team

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Outline

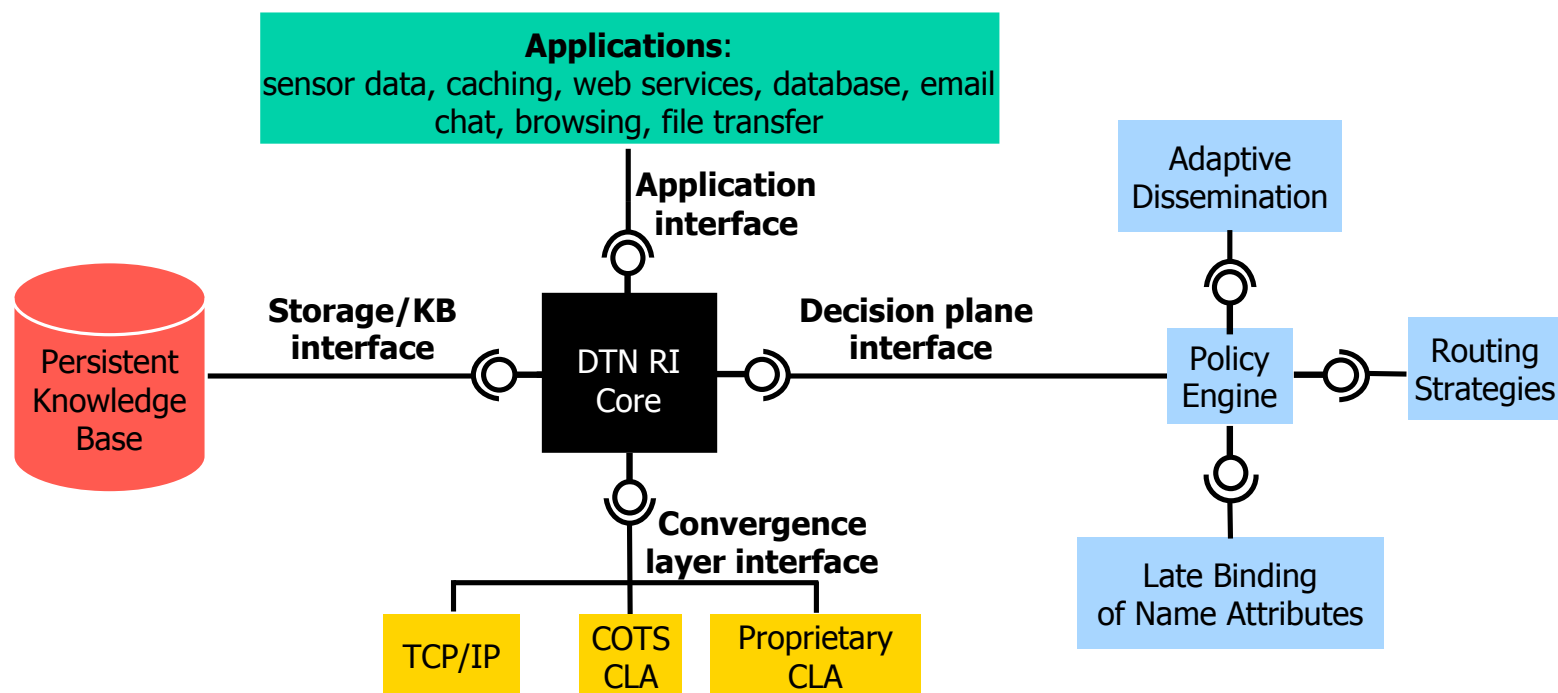
- BBN's SPINDLE-II project
- A declarative knowledge-based approach to DTN
- Late binding of intentional names for endpoints
- Disruption-tolerant access to content
- Routing, network state dissemination, and policy support

BBN's SPINDLE-II project

- System Research and Development “thread”
 - Led by Christopher Small.
 - Working with DTN RI (DTN2) developer community.
 - Completing missing features.
 - Adding to testing infrastructure, doing scalability testing.
 - Designing and implementing plug-in framework.
 - Documenting internal and external interfaces.
- Technology Innovation “thread”
 - Led by Prithwish Basu.
 - Balance of talk is about SPINDLE-II research.

SPINDLE-II

DTN RI architecture + plug-ins



- DTN RI core platform: open source
- Plug-ins: open source or proprietary
 - convergence layers, storage/KB, routing, naming

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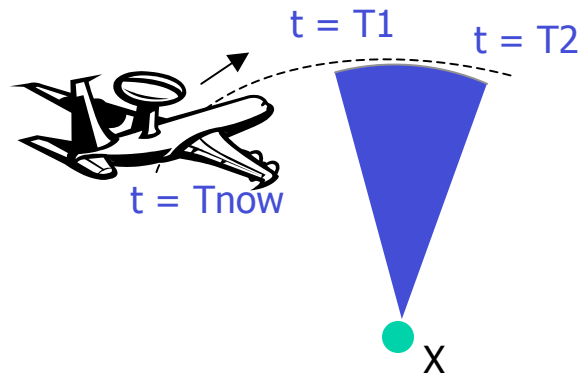
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A declarative knowledge-based DTN

- Knowledge bases offer flexible, extensible framework
 - Use inference to derive new facts from simpler facts.
 - Offer uniform query interface for explicit and derived facts.
 - Perform decision making and search of dynamic facts/rules.
 - Support rapid prototyping, deployment of new capabilities.
- Integrated declarative routing, naming, and policy
 - Approach extends naturally to content-based networking.

Specifying rich DTN information

Example: inference rules for link formation



Ground node X infers a predictedLink for plane, given its trajectory information. Prediction is useful as plane is in range too short a time for traditional network topology discovery.

- Rule to predicted link:

`predictedLink` :: spindleLink.

S : `predictedLink` [from→X, to→plane, upAt→T1, downAt→T2] :-
plane [trajectory→Trj1], X [trajectory→Trj2], walltime (Tnow),
trajectory_crossing (Tnow, Trj1, Trj2, [T1, T2]).

- KB contents can be shared, for informed routing decisions
 - forwarding, routing, and other decisions based on queries into the KB

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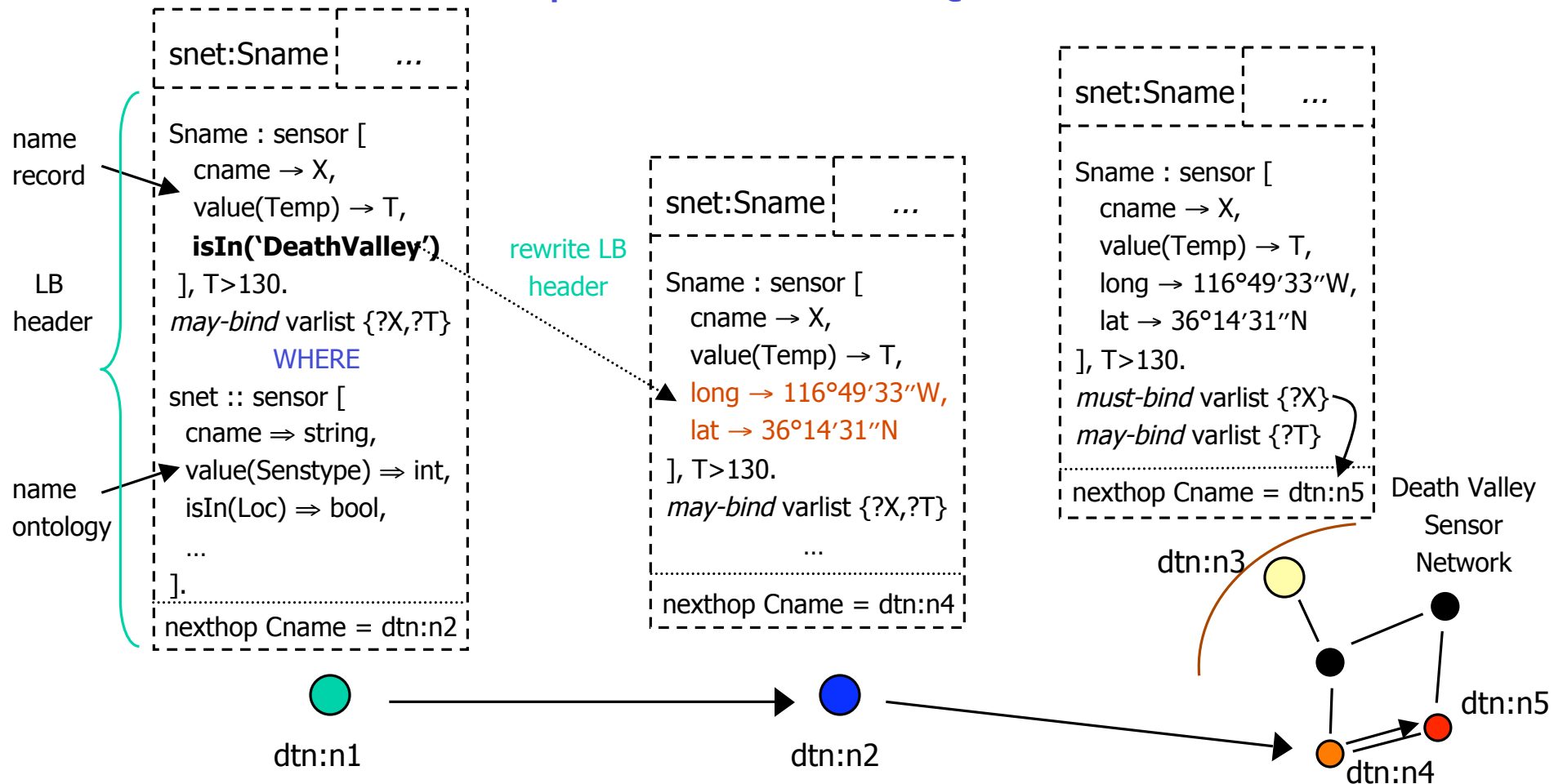
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Late binding of intensional names

- *Intensional* naming is via attributes, not identity:
 - Canonical name: Christopher.Small@bbn.com
 - Intentional name: techLead@{org=BBN,proj=SPINDLE-II,thread=dev}.
- Disruption-tolerant network often needs late name resolution
 - Final destination not known at the originator.
 - Destination learned progressively within the network.
- Mapping of names from a rich namespace to routable endpoints
 - Extensible ontology to express multi-attributed intentional names.
 - Protocols to maintain name KBs and progressive resolution within network.
- Defer binding of next hop to convergence layer address and parameters
 - Capability for several communication modes (e.g., WiFi, GPRS, ...) may exist.
 - Parameters (address, protocol, data rate) may be negotiated on contact.
 - Can route with only coarse-grained information about future adjacencies.

Progressive resolution of attributes

Endpoint Names As Queries



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Disruption-tolerant access to content

- Boy Scout troop 47 is hiking through woods, led by Scout Master Chuck, who has maps of planned route on his PDA.
- Scout Master Chuck is lousy guide, and gets troop lost in the woods. He requests new maps from Google.
- Internet scenario:
 - Chuck is off the net, TCP connection request fails, no maps.
 - Scout troop stays lost until discovered days later by Ranger Rick.
- DTN scenario:
 - Request sent to local nodes (PDAs, iPods, etc.), and stay there.
 - Scout troop stays lost until discovered days later by Ranger Rick.
 - Request forwarded to Google after scouts are returned safely home.

Disruption-tolerant access to content

- What Scout Master Chuck didn't know is that Junior Scout Billy has maps of the area in the browser cache of his Sony PlayStation Personal.
- With disruption tolerant access to content, Chuck's request could have been satisfied by Billy's PSP.
- Disruption-tolerant scenario
 - Chuck's request for new maps is received by local PDAs, iPods, etc. Each checks to see if request can be satisfied by cached content.
 - Billy's PSP determines it can satisfy the request, and replies with the cached map.
 - Chuck successfully leads troop out of the woods.

Challenges for content-based DTN

- Need distributed indexing, querying and retrieval.
 - Conventional distributed database models don't work for DTN.
- Indexing will need to include reachable caches.
- Data may be encrypted; source must be authenticated.
- Users should be told of data freshness, time to retrieve, and have control over how cache misses are handled.

Note: we understand that this is a large problem space, and don't expect to solve it completely.

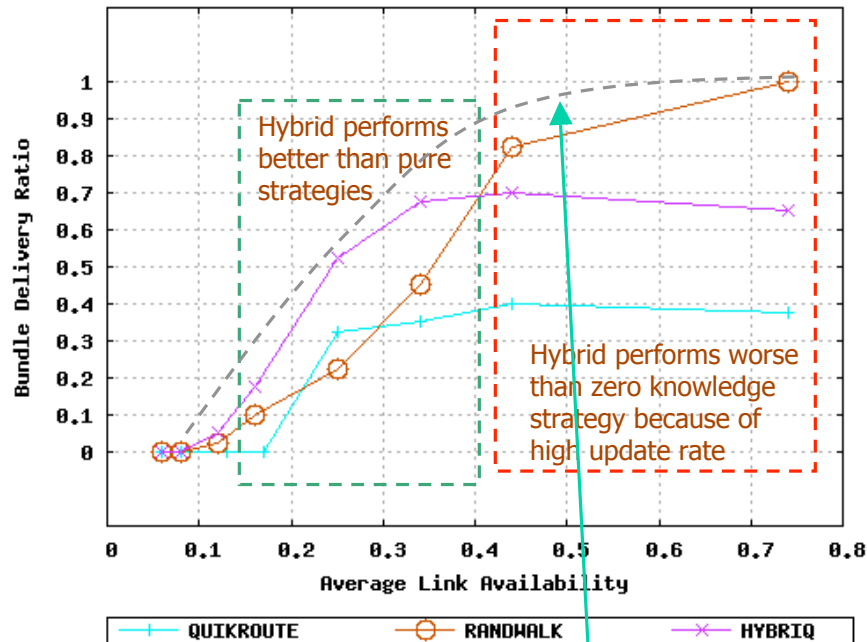
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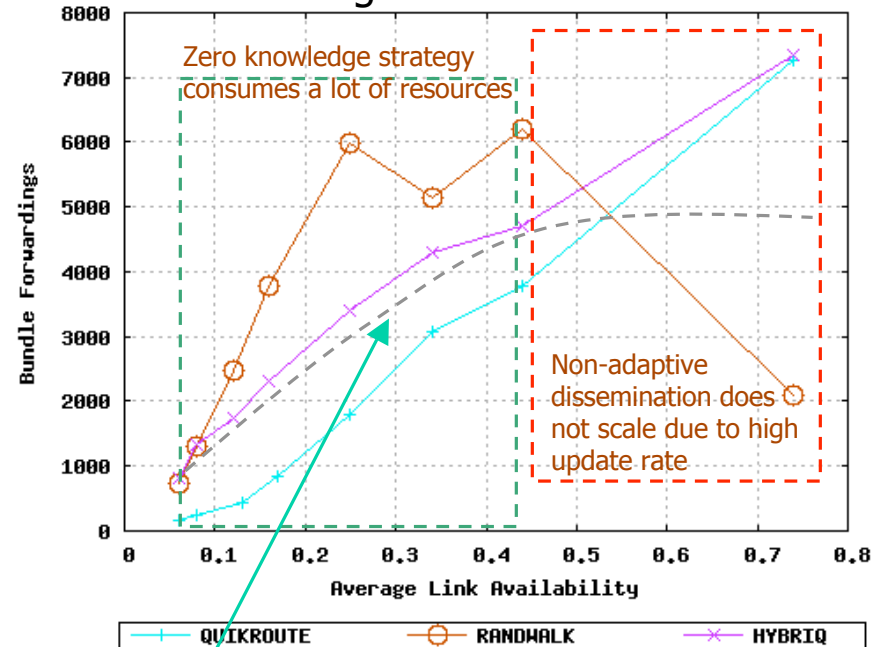
Adaptive routing and network state dissemination

Offered load: 40 bundles of 2800 bytes; Time:1200s*; Source-Destination path: 6-7 hops
 In case of dissemination-based strategies, LS updates are sent every 30s

Delivery Ratio



Bundle Forwarding



Adaptive network state dissemination will help us track these curves

Adaptive routing and network state dissemination

- Adaptive hybrid routing strategies
 - Make use of routing information when available.
 - Explore locally when path is unknown or highly uncertain.
 - Dynamic choice of strategies based on policy.
- Adaptive state dissemination service shared by routing strategies
 - Exchange subset of topology information with neighbors.
 - Algorithm is not quite link state or distance vector.
 - Uses succinct space-time representation of topology.
 - Control rate, scope, and content of updates.
 - Tradeoff control overhead versus opportunistic data transfer.
 - Service can be used to distribute policy/naming/content metadata.

Policy-based resource utilization

- Declare and control resource use in DTN nodes
 - Link formation/use based on costs and delivery requirements.
 - Storage management.
 - Choice of security services.
- Use a declarative language to express and process policies
 - Node primitives controlled by policy are explicitly declared.
 - Deductive database rules check for policy consistency and conformance of usage to policy.
 - Constraint solver searches for communication opportunities that are authorized by policy
- Protocol for policy dissemination and consistent use in a DTN.

Summary

- Engineering of DTN Reference Implementation (DTN2)
 - Feature development
 - Testing
 - Documentation
 - Plug-in interfaces and plug-ins.
- Research
 - Intentional naming and late binding
 - Content-based networking
 - Routing and adaptive dissemination

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