

# Formal Models for Path-Vector Protocols

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# Goals

- Are there areas for collaboration?
  - What have we been doing?
  - What have others been doing?
- Feedback on previous/current work
  - In which ways this could be refined to be more useful?
- Questions for future work
  - Are there aspects of routing policies that would be good to look at next?



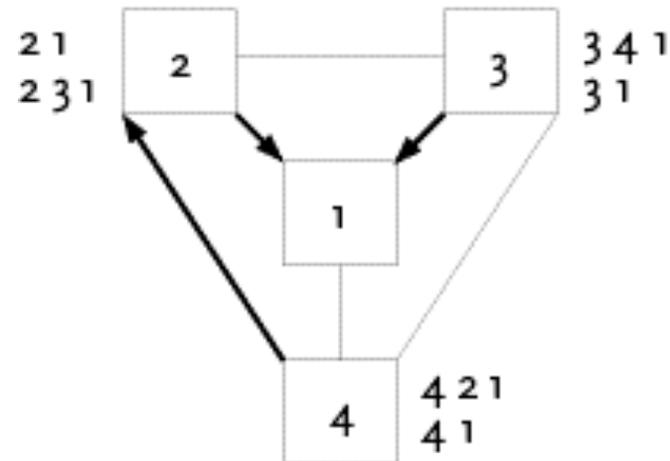
# The Design Space of Path-Vector Protocols [GJR '03]

- **Robustness:** Predictable routing tree, even after link/node failure
  - Primary concern
- **Expressiveness:** What routing policies are permitted?
  - Use the *Stable Paths Problem* as semantic domain
- **Autonomy:** What degree of independence do operators have in local-policy configuration?
  - One example: *next-hop policies*, which can contradict shortest-paths routing
- **Global Constraint:** What assumptions about the network are needed?
- **Protocol details:**
  - **Policy Opaqueness:** Can local route settings be kept private?
  - **Protocol Transparency:** How directly does the protocol apply local policy to route data?

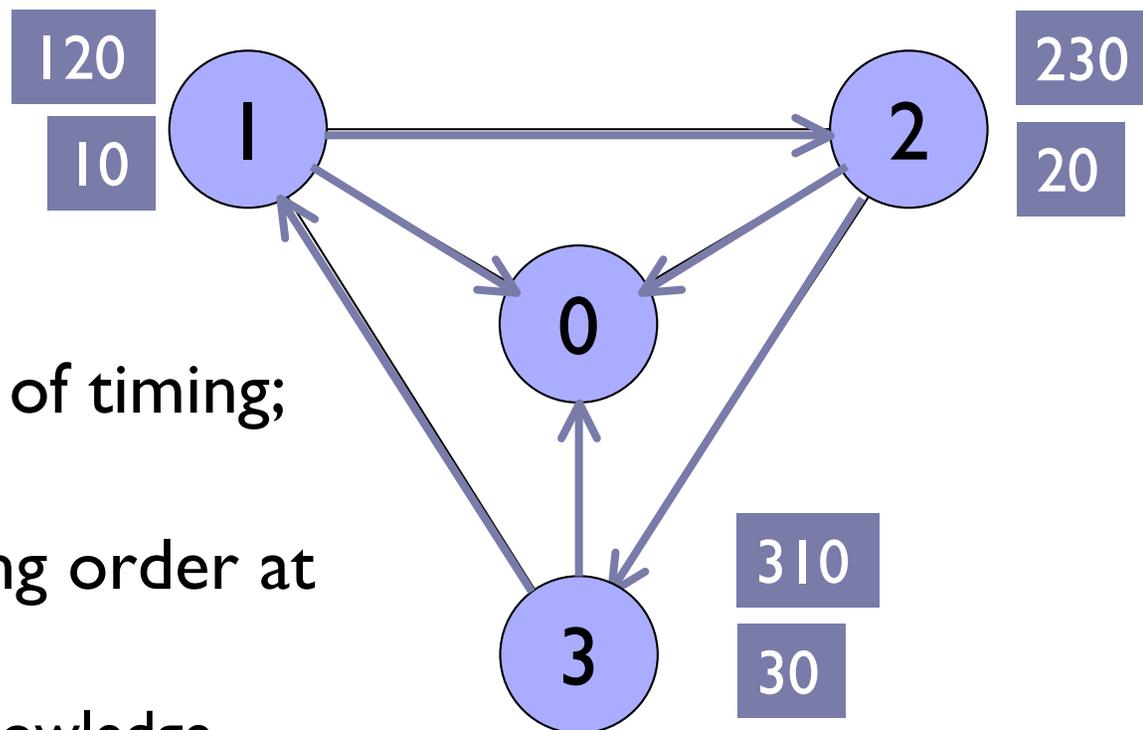
# Formally Modeling Policy Semantics

- The Stable Paths Problem (SPP) models the underlying theoretical problem that eBGP is trying to solve [Griffin-Shepherd-Wilfong '02]
- SPP solvability is NP-complete; solvability  $\not\Rightarrow$  convergence.

An SPP instance is a graph in which each node represents one AS and has a policy in the form of a linear preference ordering on paths.

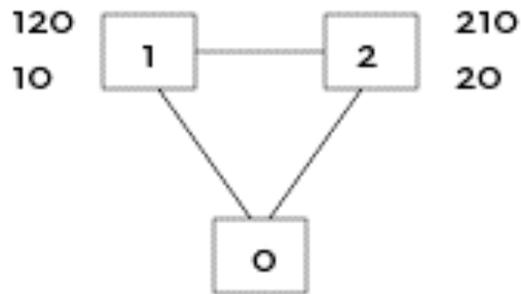


# Bad Gadget [GSW'02]

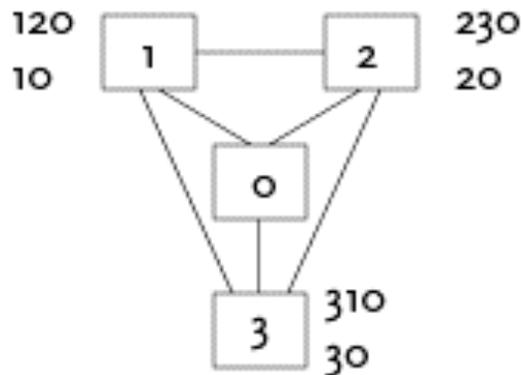


- Cycles regardless of timing;  
no stable solution
- Can fix by changing order at  
any node
  - Fix uses global knowledge
- Note: Each node prefers rim  
to spoke

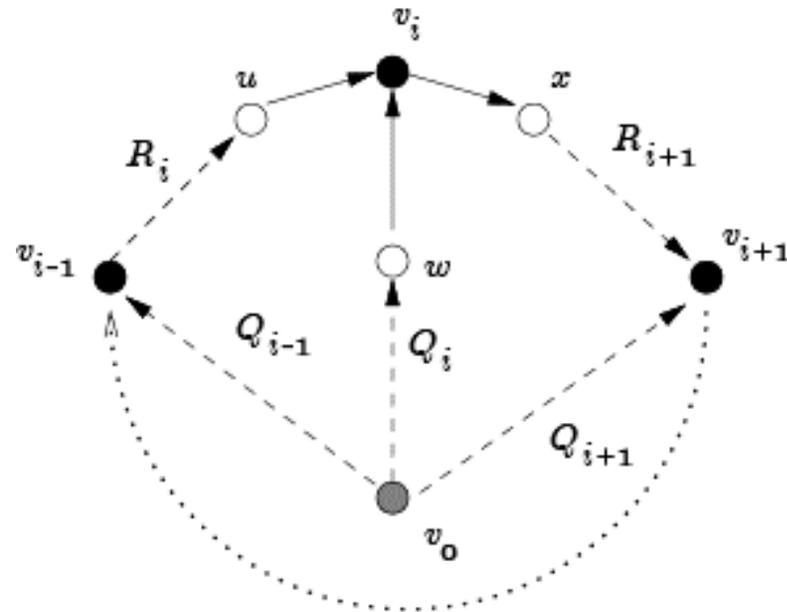
# SPP Results [GSW '02]



DISAGREE (multiple solutions)



BAD GADGET (no solution)



Dispute Wheel

No dispute wheel implies robust convergence.

# Path-Vector Policy Systems

[GJR '03]

Formal model of path-vector routing:

$(PV, PL, K)$

Path-Vector System:

The underlying message-exchange system for route information. What is exchanged and how?

Global Constraint:

What assumptions about the network must be true to achieve robustness?

Policy Language:

How can policies be described?  
PL acts as a local constraint on the expressiveness of policies.

Question:

What role do these components play in achieving protocol design goals?



## Path-Vector Algebras [Sobrinho '03]

- A *path-vector algebra* defines:
  - Signatures (path data objects)
  - Labels (combines import and export policies)
    - Apply label to signature to obtain new signature (the path data after export/import)
  - Weight function on signatures (rank)
  - Operation to apply labels to signatures
  - Rank criteria for tie-breaking
- These abstract away some protocol-level details



# Robustness Condition

[GJR '03, S '03]

**Theorem:** A protocol in which a path's (global) rank always **increases** as it is extended (by export/import) is dispute-wheel-free (and thus robust).

(Assume that we prefer the path with smallest rank, as with cost.)

- Increasing systems generalize cost functions
  - Cost now assigned to (path, edge) pairs



# Trade-Offs in Implementation

## [GJR '03]

**Theorem.** A transparent, robust PVPS that supports next-hop policies and is at least as expressive as shortest paths **must have a non-trivial global constraint.**

**Corollary.** A globally unconstrained, robust PVPS that is expressive enough to capture all increasing configurations either **does not support next-hop policies** or is **not transparent**, or both.



# Hierarchical BGP (HBGP)

- Partition neighbors into customers, providers, and peers
- Local constraints on policies
  - Scoping: Share route data from customers with everyone, share data from everyone with customers, do not share other data
  - Relative-preference: Prefer peer routes to provider routes, customer routes to both peer and provider
- No customer/provider cycles
- HBGP is robust [Gao-Rexford '01]
  - Are constraints violated often? Why?



# Extending HBGP [JR' 04]

- Use the PVPS framework to generalize the HBGP constraints of [GR' 01, GGR' 01].
- A class-based PVPS is described by:
  - A set of classes (types of neighbor assignments, e.g., customer/provider/peer) and consistency relationships between them
  - Scoping rules
  - Relative-preference rules
- These systems are transparent and support next-hop policies enough to require a nontrivial global constraint.



## Class-Based Robustness [JR '04]

From the class description alone, we can construct a global constraint involving a check on pairs of class assignments.

- Prevent cycles that could form dispute wheel rims by checking two cases
- Networks obeying this constraint are robust.
- Networks violating this constraint allow nodes to write policies that induce routing anomalies.



# Dispute Rings [FJB '05]

- Dispute rings specialize dispute wheels
  - No node appears more than once
- Safety under filtering generalizes robustness
  - Allow arbitrary filtering, not just all paths through a certain node or edge

**Theorem: Dispute-ring-freeness is necessary for safety under filtering**

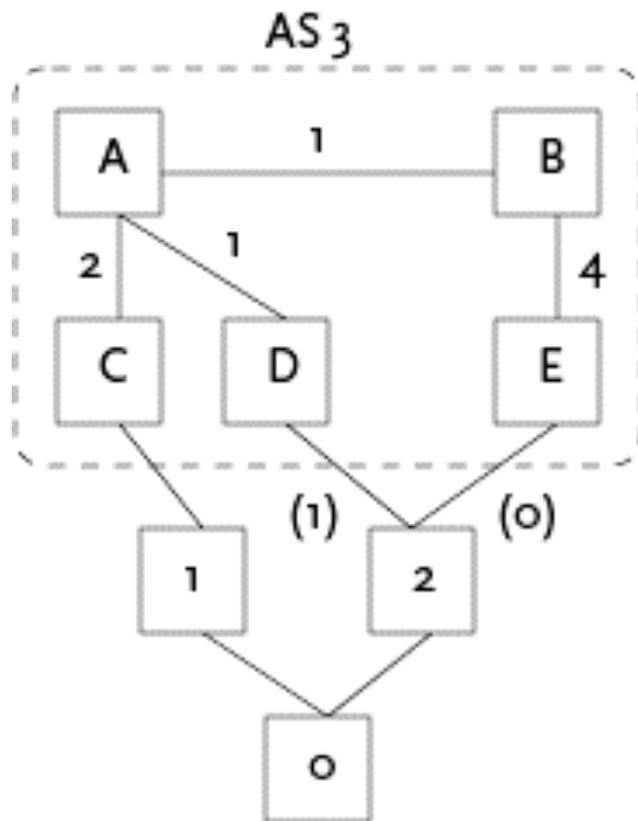
- Still open: Is there a necessary condition for robustness?



# How to Model MEDs?

- Have a selection function choose one route (according to local policy) from a set of routes
  - No longer ranking paths linearly
- A singleton-valued selection function  $f$  satisfies Independent Route Ranking if, for  $T$  containing  $S$ ,
$$f(T) = P_2 \text{ implies } f(S) = P_2 \text{ or } P_2 \text{ is in } T \setminus S$$
  - Learning new routes shouldn't cause the selection of a new, previously known route
  - Potentially violated by use of MED attribute
- Second condition for set-valued selection functions
  - Source of future work; focus on singleton-valued here

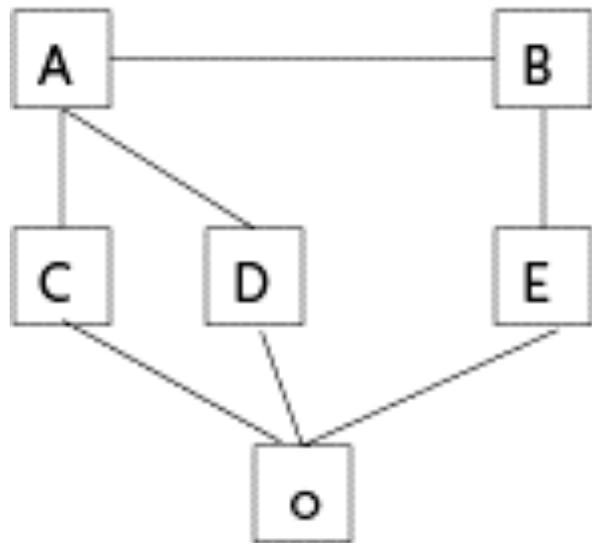
# Generalized SPP [GW '02]



MED-EVIL [GW'02] (no solution)

- BGP selection:
  - lowest MED value from paths to the same AS; then
  - shortest IGP distance.
- IGP distances are shown near intra-domain links.
- MED values are shown in parentheses near inter-domain links.
- This example oscillates.

# Violate IRR Using MEDs



MED-EVIL [GW'02]  
(condensed)

$$\begin{aligned} \sigma_A(AC0, AD0) &= AD0 \\ \sigma_A(AD0, ABE0) &= ABE0 \\ \sigma_A(AC0, ABE0) &= AC0 \\ \sigma_A(AC0, AD0, ABE0) &= AC0 \end{aligned}$$

$$\begin{aligned} \sigma_B(BE0) &= BE0 \\ \sigma_B(BAD0, BE0) &= BE0 \\ \sigma_B(BAC0, BE0) &= BAC0 \end{aligned}$$

# Generalized Path Relations [JR'06]

**Subpath**  $P_1 \ominus P_2$  iff  $P_1 = v \cdots v_0$ ,  $P_2 = u \cdots v_0$ , and  $uP_1 = P_2$

**Linear Selection**  $P_1 \otimes P_2$  iff  $P_1 = v \cdots v_0$ ,  $P_2 = u \cdots v_0$ , and  
 $\exists S : \sigma_u(\{uP_1, P_2\} \cup S) = uP_1$

**Nonlinear Selection 1**  $P_1 \odot_1 P_2$  iff  $P_1 = v \cdots v_0$ ,  $P_2 = u \cdots v_0$ ,  
and there exists a set of routes  $S \not\supseteq uP_1$  such that  
 $\sigma_u(\{P_2\} \cup S) \neq P_2$  and  $\sigma_u(\{uP_1, P_2\} \cup S) = P_2$

**Nonlinear Selection 2**  $P_1 \odot_2 P_2$  iff  $P_1 = v \cdots v_0$ ,  $P_2 = u \cdots v_0$ ,  
and there exists a set of routes  $S \not\supseteq uP_1$  such that  
 $\sigma_u(S) = P_2$  and  $\sigma_u(\{uP_1\} \cup S) \notin \{uP_1, P_2\}$



# Generalized Dispute Wheels [JR'06]

- Extend original notion of dispute wheel to include new relations between paths
  - Gives sufficient condition for robustness in generalized SPP
- Not considered in initial generalization of SPP
  - GW'02 translated limited class of GSPPs to SPP, applied SPP convergence conditions
  - Here, generalized dispute wheels apply directly to GSPP



# Summary of Previous Work

- PVPS framework for study of path-vector protocols
  - Conditions needed for robustness
  - Tradeoffs involved in implementing these conditions
- Concrete and reasonable guidelines for class-based systems
- Extended framework to allow nonlinear selection
  - Start to model interactions between internal and external routing



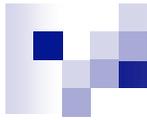
# Questions (I)

- Feedback on our work
  - How to refine it to make it more useful?
    - Can people work within class-based routing?
    - Should we focus on next-hop (plus tweaks)?
    - What are pressing (non-implementation) questions?
  - Related areas for joint work?
    - In RRG?



## Questions (II)

- What sort of policies should we look at?
  - What sort of policies are typically written?
  - Are there policies you'd like to write but can't?
  - How do you want to be able to reconfigure policies?
  - What anomalies do you see? (Often?)
  - What are typical iBGP policies?



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